

# ***CS202 - Algorithm Analysis***

## **Merge Sort**

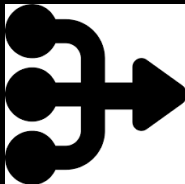
Aravind Mohan

Allegheny College

February 21, 2023



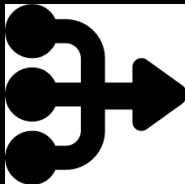
# Merge Sort Algorithm



## Strategy:

- **Divide:** if  $S$  has at least two elements, remove all the elements from  $S$  and put them into two sequences  $S_1$  and  $S_2$ , each containing about half of the elements of  $S$ . (i.e.,  $S_1$  contains the first floor  $(n/2)$  elements and  $S_2$  contains the remaining floor  $(n/2)$  elements.
- **Conquer:** Sort sequences  $S_1$  and  $S_2$  using Merge Sort.
- **Combine:** Put back the elements into  $S$  by merging the sorted sequences  $S_1$  and  $S_2$  into one sorted sequence.

# Merge Sort Algorithm



## Characteristics:

- sort out of "place", i.e., does require an additional array
- uses divide and conquer principle
- worst case running time is  $O(n \times \log(n))$

# Merge Sort Algorithm

## Merge Procedure (linear)

**Algorithm** - Merge( $A, p, m, r$ )

**Input:** an  $n$ -element un-sorted array  $A$  of integer values, a lower bound  $p$  of the array  $A$ , and a pivot  $r$  in the array  $A$ .

**Output:** an  $n$ -element sorted array  $A$  of integer values.

$n_1 \leftarrow m - p$

$n_2 \leftarrow r - m$

Initialize Array  $L$  of size  $n_1 + 1$

Initialize Array  $R$  of size  $n_2 + 1$

**for**  $i = 0$  to  $n_1$  **do**

$L[i] \leftarrow A[p+i]$

**end for**

**for**  $j = 0$  to  $n_2$  **do**

$R[j] \leftarrow A[m+j]$

**end for**

$L[n_1 + 1] \leftarrow \infty$

$R[n_2 + 1] \leftarrow \infty$

CONTINUE ON NEXT PAGE ...

# Merge Sort Algorithm

## Merge Procedure (linear)

```
Initialize  $i, j \leftarrow 0$   
for  $k = p$  to  $r$  do  
  if  $L[i] \leq R[j]$  then  
     $A[k] \leftarrow L[i]$   
     $i \leftarrow i+1$   
  else  
     $A[k] \leftarrow R[j]$   
     $j \leftarrow j+1$   
  end if  
end for
```

# Merge Sort Algorithm

## MergeSort Procedure (logarithmic)

**Algorithm** - MergeSort( $A, p, r$ )

**Input:** an  $n$ -element un-sorted array  $A$  of integer values, a lower bound  $p$  of the array  $A$ , and a pivot  $r$  in the array  $A$ .

**Output:** an  $n$ -element sorted array  $A$  of integer values.

**if**  $p < r$  **then**

$m \leftarrow \text{Floor } (p + r)/2$

MergeSort( $A, p, m$ )

MergeSort( $A, m+1, r$ )

Merge( $A, p, m, r$ )

**end if**

# Merge Example

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
| 1 | 5 | 7 | 8 | 2 | 4 | 6 | 9 |
|---|---|---|---|---|---|---|---|

i

|   |   |   |   |          |
|---|---|---|---|----------|
| 1 | 5 | 7 | 8 | $\infty$ |
|---|---|---|---|----------|

j

|   |   |   |   |          |
|---|---|---|---|----------|
| 2 | 4 | 6 | 9 | $\infty$ |
|---|---|---|---|----------|

i

|   |   |   |   |          |
|---|---|---|---|----------|
| 1 | 5 | 7 | 8 | $\infty$ |
|---|---|---|---|----------|

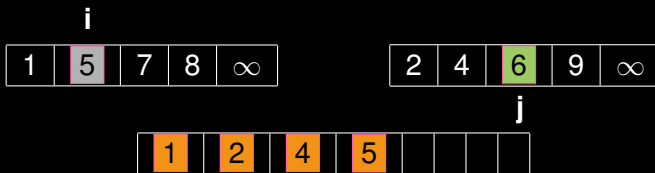
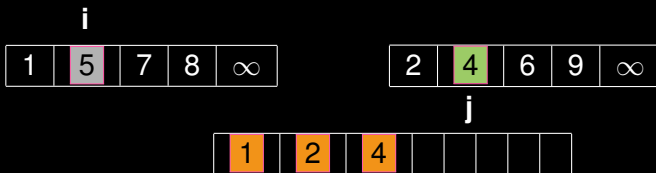
j

|   |   |   |   |          |
|---|---|---|---|----------|
| 2 | 4 | 6 | 9 | $\infty$ |
|---|---|---|---|----------|

|   |  |  |  |  |  |  |  |
|---|--|--|--|--|--|--|--|
| 1 |  |  |  |  |  |  |  |
|---|--|--|--|--|--|--|--|

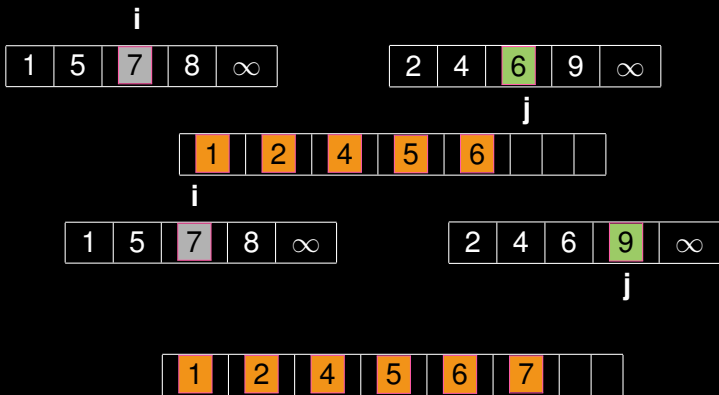
|   |   |  |  |  |  |  |  |
|---|---|--|--|--|--|--|--|
| 1 | 2 |  |  |  |  |  |  |
|---|---|--|--|--|--|--|--|

# Merge Example

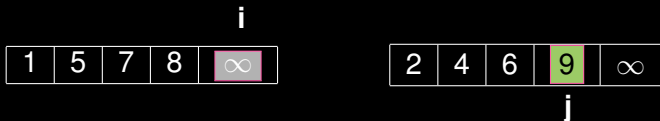




# Merge Example

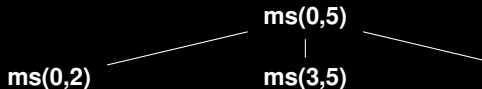


# Merge Example



# Merge Sort Example

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 9 | 6 | 5 | 0 | 8 | 2 |
|---|---|---|---|---|---|

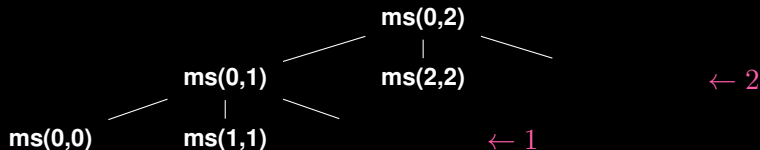


← 5

Let us complete the tree?

# Merge Sort Example

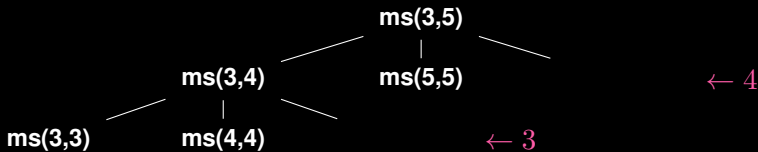
|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 9 | 6 | 5 | 0 | 8 | 2 |
|---|---|---|---|---|---|



Let us complete the tree?

# Merge Sort Example

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 9 | 6 | 5 | 0 | 8 | 2 |
|---|---|---|---|---|---|



**OK done**

# Merge Sort Example

Merge(0,0,1)

|   |   |  |  |  |  |
|---|---|--|--|--|--|
| 9 | 6 |  |  |  |  |
|---|---|--|--|--|--|

|   |          |
|---|----------|
| 9 | $\infty$ |
|---|----------|

|   |          |
|---|----------|
| 6 | $\infty$ |
|---|----------|

|   |   |  |  |  |  |
|---|---|--|--|--|--|
| 6 | 9 |  |  |  |  |
|---|---|--|--|--|--|

# Merge Sort Example

Merge(0,1,2)

|   |   |   |  |  |  |
|---|---|---|--|--|--|
| 6 | 9 | 5 |  |  |  |
|---|---|---|--|--|--|

|   |   |          |
|---|---|----------|
| 6 | 9 | $\infty$ |
|---|---|----------|

|   |          |
|---|----------|
| 5 | $\infty$ |
|---|----------|

|   |   |   |  |  |  |
|---|---|---|--|--|--|
| 5 | 6 | 9 |  |  |  |
|---|---|---|--|--|--|

# Merge Sort Example

**Merge(3,3,4)**

|   |   |   |   |   |  |
|---|---|---|---|---|--|
| 5 | 6 | 9 | 0 | 8 |  |
|---|---|---|---|---|--|

|   |          |
|---|----------|
| 0 | $\infty$ |
|---|----------|

|   |          |
|---|----------|
| 8 | $\infty$ |
|---|----------|

|   |   |   |   |   |  |
|---|---|---|---|---|--|
| 5 | 6 | 9 | 0 | 8 |  |
|---|---|---|---|---|--|



# Merge Sort Example

**Merge(3,4,5)**

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 5 | 6 | 9 | 0 | 8 | 2 |
|---|---|---|---|---|---|

|   |   |          |
|---|---|----------|
| 0 | 8 | $\infty$ |
|---|---|----------|

|   |          |
|---|----------|
| 2 | $\infty$ |
|---|----------|

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 5 | 6 | 9 | 0 | 2 | 8 |
|---|---|---|---|---|---|

# Merge Sort Example

**Merge(0,2,5)**

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 5 | 6 | 9 | 0 | 2 | 8 |
|---|---|---|---|---|---|

|   |   |   |          |
|---|---|---|----------|
| 5 | 6 | 9 | $\infty$ |
|---|---|---|----------|

|   |   |   |          |
|---|---|---|----------|
| 0 | 2 | 8 | $\infty$ |
|---|---|---|----------|

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 0 | 2 | 5 | 6 | 8 | 9 |
|---|---|---|---|---|---|

**Done Sorting!**

# Merge Sort Algorithm - An analysis

- Space complexity:  $O(n)$
- Time complexity:  $O(n \times \log n(n))$  for Worst, Average, and Best case

## **Sedgewick 2.2 Merge Sort**

# Questions?

**Please ask if there are any Questions!**