

CS202 - Algorithm Analysis

Tree Algorithms - Module 2

Aravind Mohan

Allegheny College

March 21, 2023

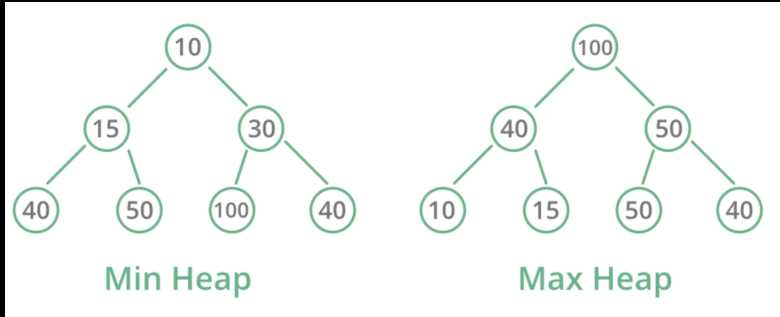


Sedgewick 2.4 Heap Sort

What is a Binary Heap?

- Each node has atmost two children.
- Complete binary tree or atmost complete binary tree qualified as binary heap.
- Node with no children is also qualified as heap.
- Left skewed or right skewed tree is not a heap.
- There are two types of heap, namely:
 Max heap and Min heap.

What is a Binary Heap?



Binary Heap Properties

- Binary Heap has two main properties:
 - 1 Order property
 - 2 Shape property

Binary Heap Properties

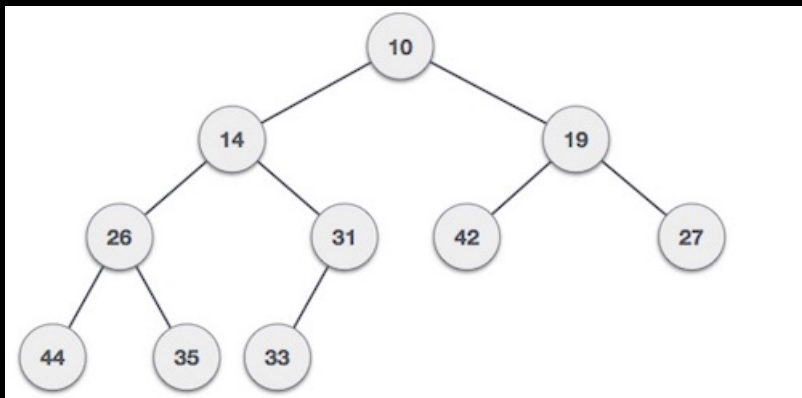
- **Order** property: The value in node n is \geq the values in its children, for every node n (MAX heap).
- How about MIN heap?

Binary Heap Properties

Shape Property:

- All leaves are either at depth d or $d - 1$ for some d
- All of the leaves at depth $d - 1$ are to the right of the leaves at depth d
- And the following:
 - 1 There is at most 1 node with just 1 child v .
 - 2 v is the left child of its parent.
 - 3 v is the rightmost leaf at depth d .

Binary Heap Example



Heap Sort

- **Phase 1:** convert the array into an n -element heap
- **Phase 2:** repeatedly remove maximum element from the heap, and place that element in its proper position in the array
 - swap element at 0th position with element at $(n - 1)$ th position and then “reheapify” considering only the first $n - 1$ elements
 - repeat this process until heap size is reduced to 1 (minimum element remains, at 0th position)

Heap Sort: Phase 1 - build the heap

```
for  $i = 1$  to  $n - 1$  do  
    insert element  $s[i]$  into the heap consisting  
    of the elements  $s[0] \dots s[i - 1]$  [heapify]
```

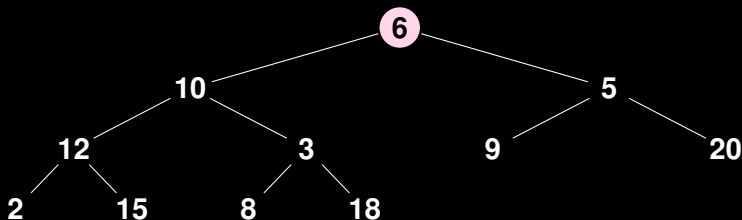
Once the heap is built, $s[0]$ will contain the maximum element

Heap Sort Example

6	10	5	12	3	9	20	2	15	8	18
---	----	---	----	---	---	----	---	----	---	----

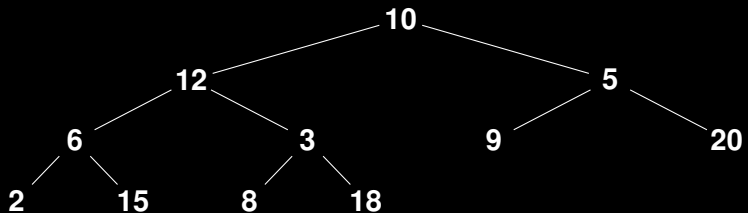
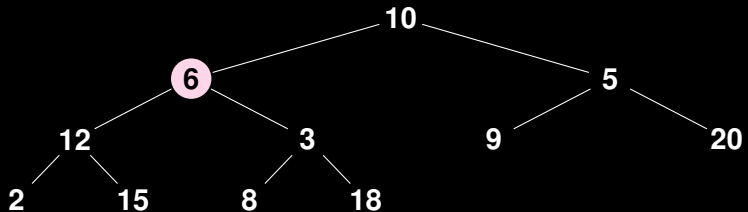
Step 1: Initial Heapify

Fix-1



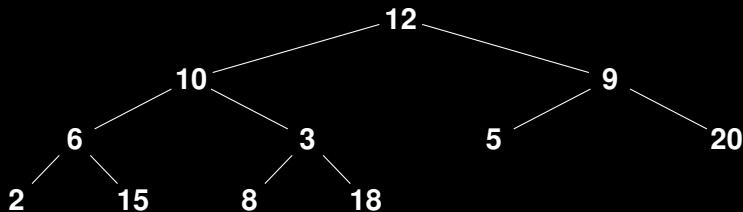
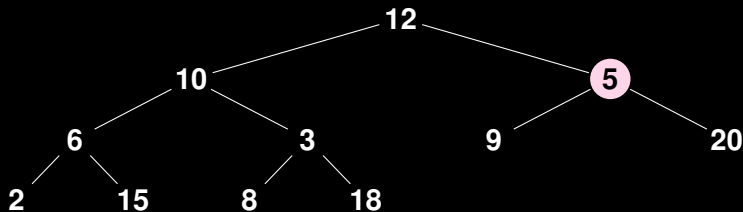
Heap Sort Example

Fix-2

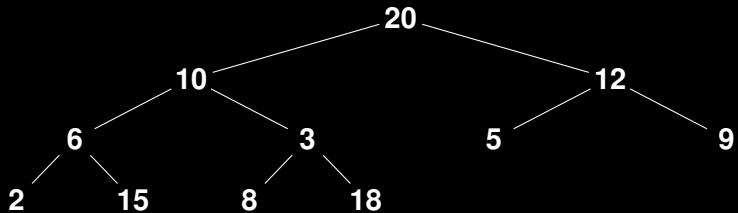
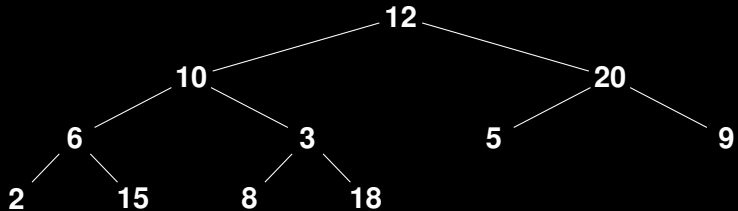


Heap Sort Example

Fix-3

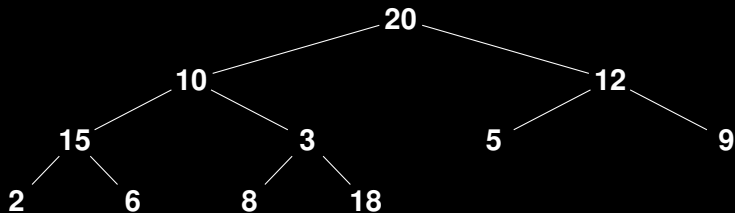
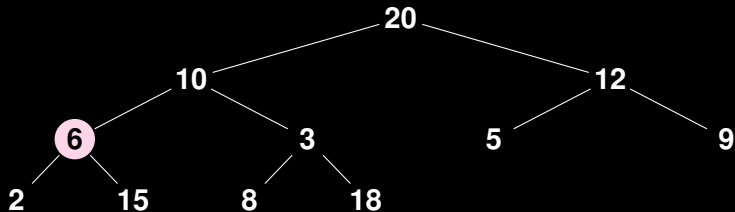


Heap Sort Example

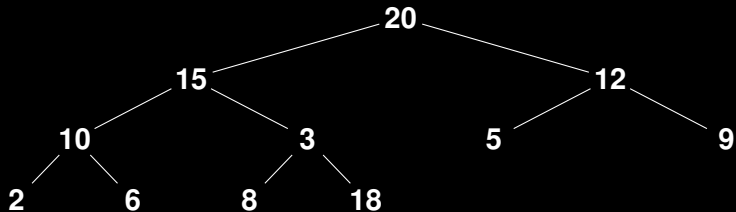


Heap Sort Example

Fix-4

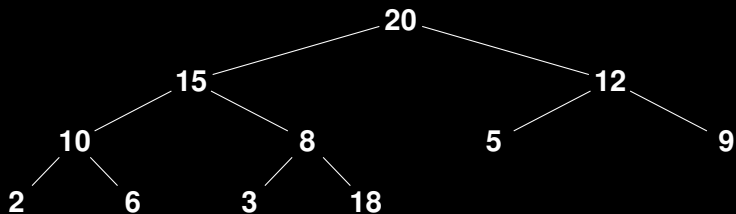
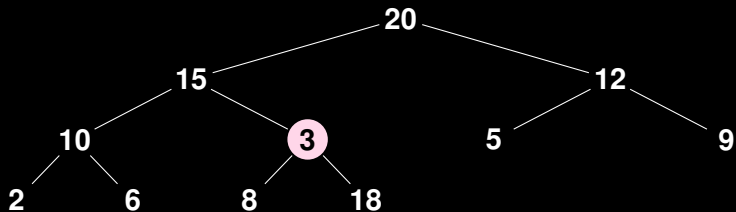


Heap Sort Example

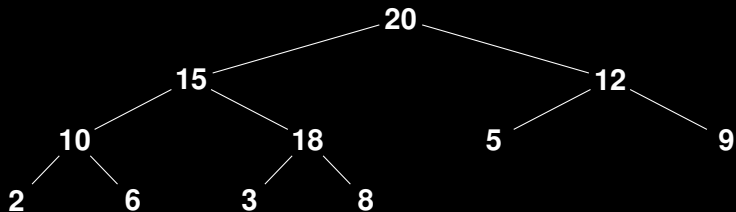


Heap Sort Example

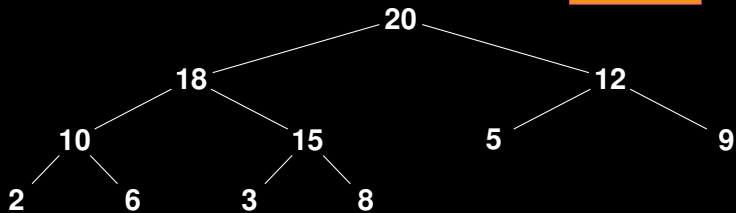
Fix-5



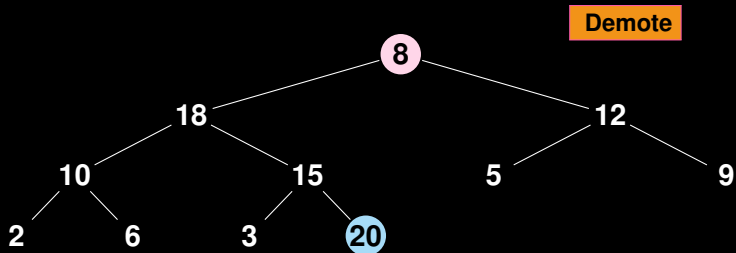
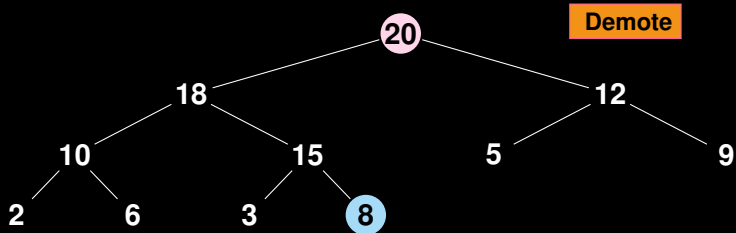
Heap Sort Example



Heapified

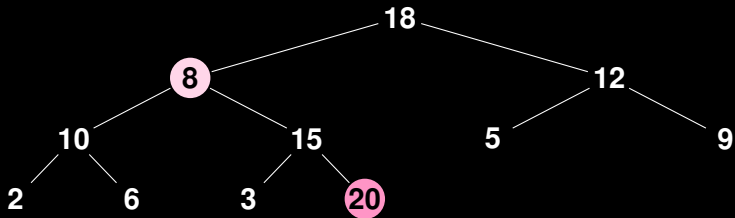
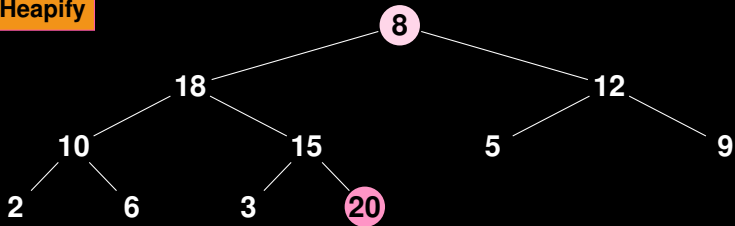


Heap Sort Example

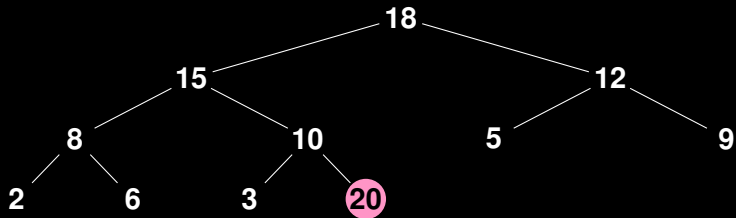
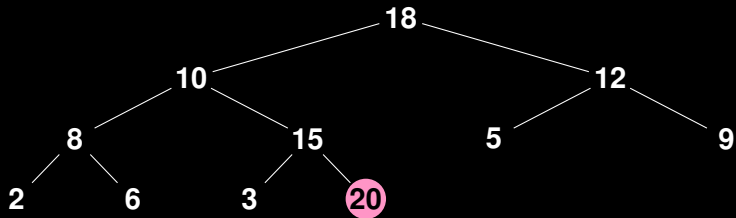


Heap Sort Example

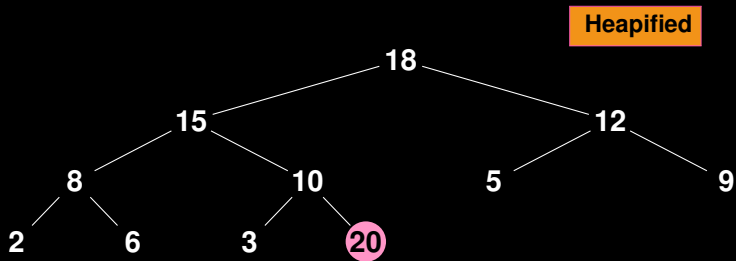
Step 2: Heapify



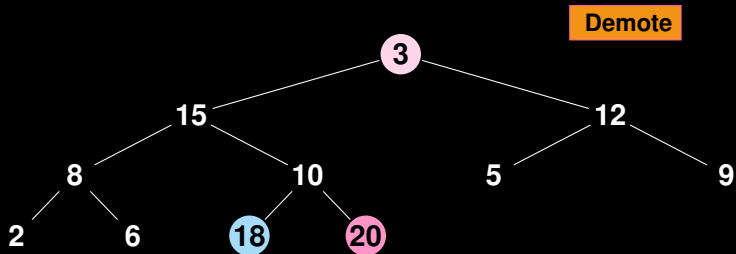
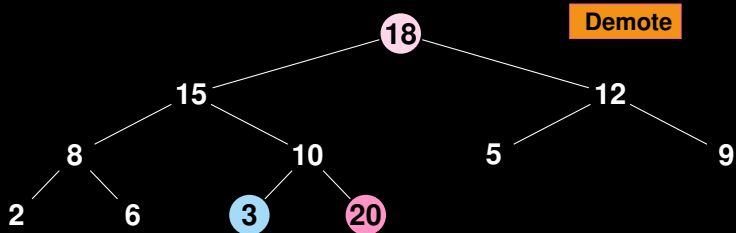
Heap Sort Example



Heap Sort Example

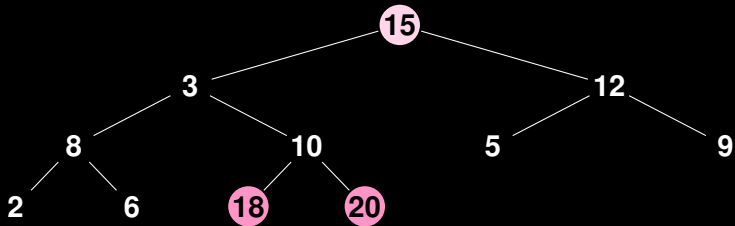
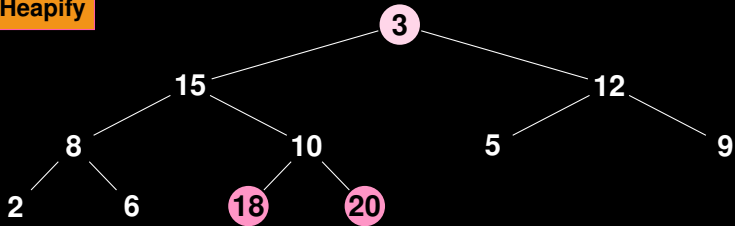


Heap Sort Example

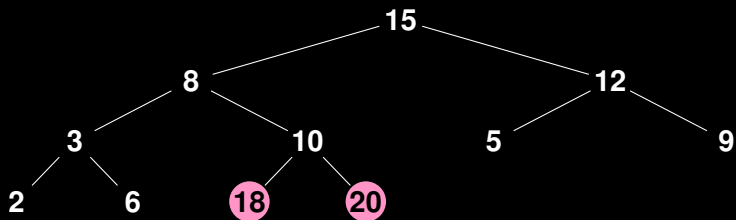
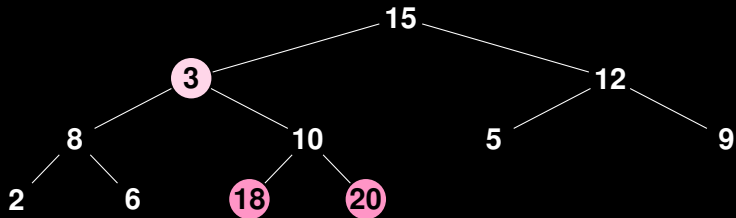


Heap Sort Example

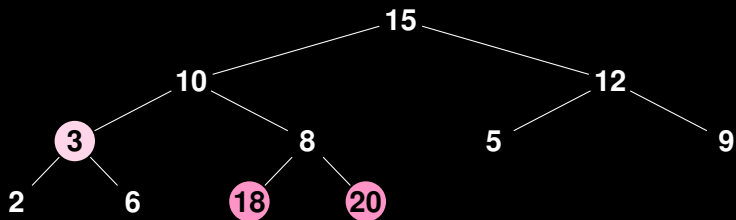
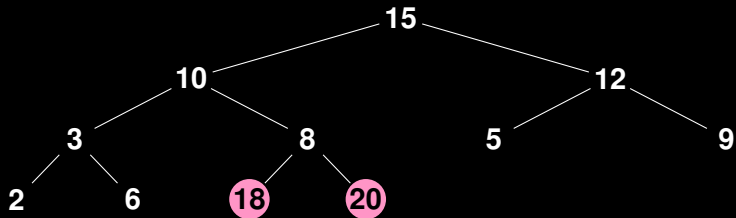
Step 3: Heapify



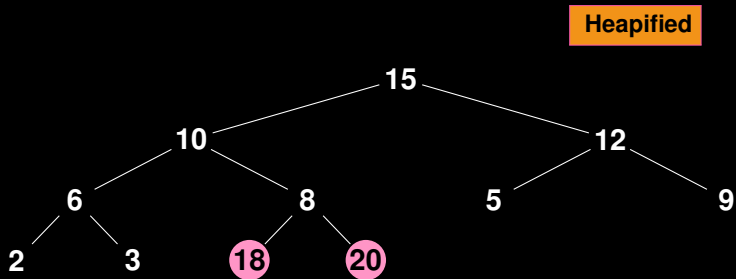
Heap Sort Example



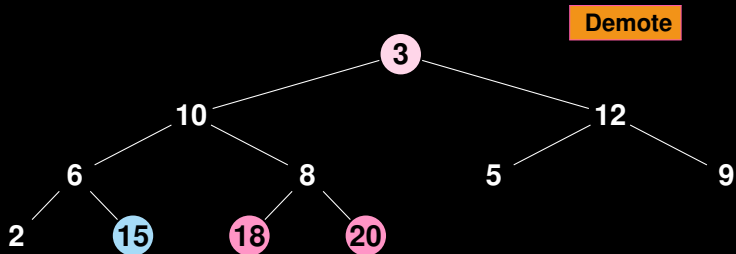
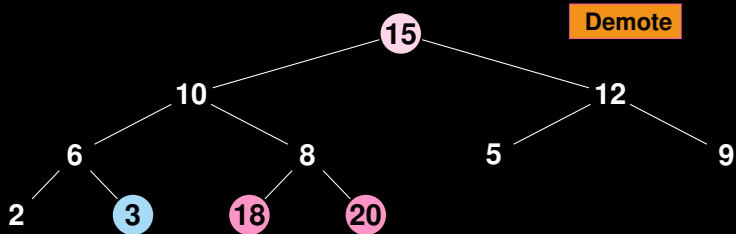
Heap Sort Example



Heap Sort Example

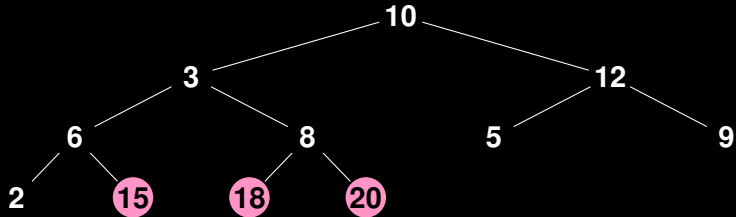
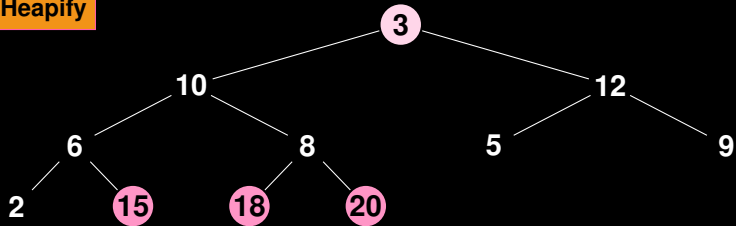


Heap Sort Example

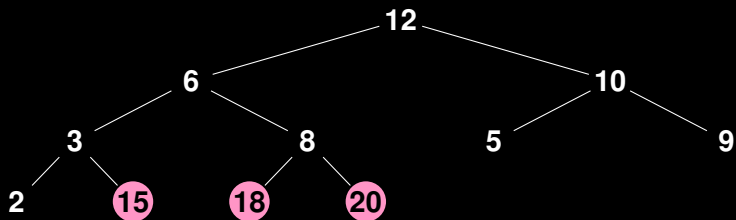
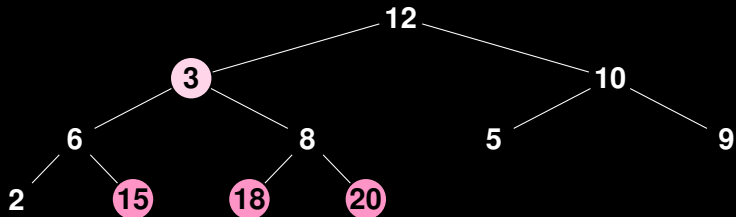


Heap Sort Example

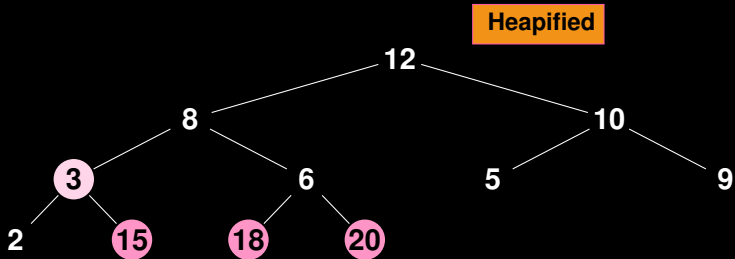
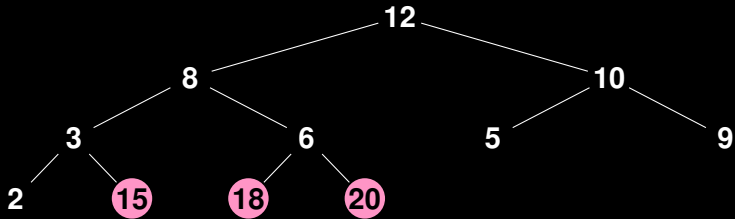
Step 4: Heapify



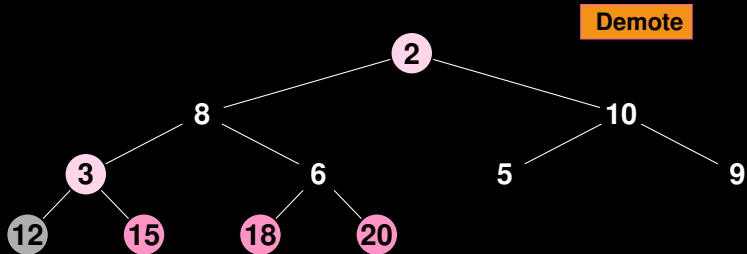
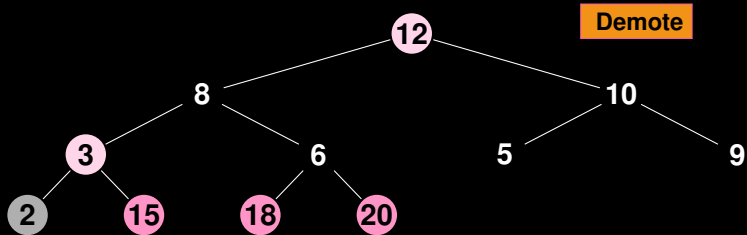
Heap Sort Example



Heap Sort Example



Heap Sort Example



Heap Sort

Complete other steps.

Repeat steps till all nodes becomes **Red**
At the end of all steps, given array is **Sorted**.

Heap Sort Complexity

for $i \leftarrow 1$ to $n-1$ do

insert element $s[i]$ into the
heap consisting of the
elements $s[0] \dots s[i-1]$

$O(n \log n)$

$O(\log n)$ operations

for $i \leftarrow n-1$ down to 1 do

swap $s[0]$ and $s[i]$
"demote" $s[0]$ to its proper place
in the heap consisting of the
elements $s[0] \dots s[i-1]$

$O(n \log n)$

Heap Sort

Note that heap sort is just a more clever version of selection sort since a maximum is repeatedly selected and placed in its proper position.

Selection Sort in next Lab.

Heap Sort

Best, Average, and Worst case is $O(n \log n)$

Unstable, and Inefficient for complex data.

Sorting Algorithms - Comparison

Algorithm	Time	Notes
selection-sort	$O(n^2)$	<ul style="list-style-type: none">slowin-placefor small data sets ($< 1K$)
insertion-sort	$O(n^2)$	<ul style="list-style-type: none">slowin-placefor small data sets ($< 1K$)
heap-sort	$O(n \log n)$	<ul style="list-style-type: none">fastin-placefor large data sets ($1K - 1M$)
merge-sort	$O(n \log n)$	<ul style="list-style-type: none">fastsequential data accessfor huge data sets ($> 1M$)

Sedgewick 2.4 Heap Sort

Questions?

Please ask if there are any Questions!