Memory Access Management with LLVM and OpenCL

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- 2 Low Level Virtual Machine: Compiler Infrastructure
- 3 Open Computing Language: Open Standard For Parallel Programming
- 4 Memory Management Done By LLVM On OpenCL
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Context

- Computing devices opt for parallel programming
- Shared memory management problem
- Tool capable to give enough information concerning the memory to the developer to avoid these problems
- Existing tools:
 - Implemented for LLVM: AddressSanitizer, MemorySanitizer, ThreadSanitizer
 - Used on simulated environments: Symbolic OpenCL, OCLgrind

Problem

- Compile time: syntax, typechecking and compiler errors
- Runtime: only division by zero, deferencing a null pointer or running out of memory errors

Main idea:

- Go from an OpenCL program
- Construct LLVM tools and passes to get as much memory information as possible during compile time
- Provide them to the developer

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Definition and purpose

- Framework with many tools, implemented in C++
- Open Source and providing low level code representation in Single Static Assignment
- Compilation with LLVM: clang -c -emit-llvm const.c -o const.bc

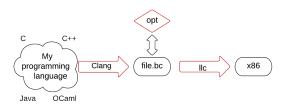


Figure: Global view of LLVM use

LLVM Pass

- Collection of libraries ready-to-use allowing code analyses and optimizations. Each analyse or optimization is called a pass.
- Example of pass: Memory allocation, common subexpression elimination ...

Remark

A pass can be run on a complete code or only on a portion. This leads to the distinction of various types of passes. In our case, we will focus on ModulePass (ran on the whole code) and FunctionPass (ran in functions)

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Definition

- Framework dedicated to writting C programs executed across multiple heterogeneous platforms
- OpenCL programs are compiled at runtime
- Platforms are composed of many compute devices
- Compute device execute programs: kernels
- One kernel can be executed in different compute devices
- One kernel can have multiple memory areas: buffers
- 1 compute device = many processing elements

Hierarchy

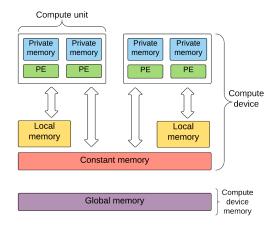


Figure: Memory hierarchy in OpenCL

From Runtime to Compile Time

- Information about kernel buffers needed (size, access right...)
- Known by OpenCL during runtime: communication
- Link between runtime and compile time

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Organization

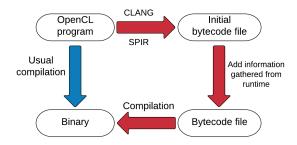


Figure: Steps of our work

Information of kernels needed. Two options:

- Force the user to insert information
- LLVM pass changing the function signatures at compile time

Adding arguments

Function signature before LLVM pass :

```
define cc76 void @addition(<2 x float> %alpha, float addrspace(1)* nocapture %x, float addrspace(1)* nocapture %y) nounwind {/*Body of the function*/}
```

Function signature after LLVM pass :

```
define cc76 void @addition(<2 x float> %alpha, float addrspace(1)* nocapture %x, float addrspace(1)* nocapture %y, int %size_x, int %size_y) nounwind {/*Body of the function*/}
```

```
!opencl.kernels = !{!0} /*Metadata before LLVM pass*/
!0 = metadata !{void (<2 x float>, float addrspace(1)
   *, float addrspace(1)*)* @addition, metadata !1,
   metadata !2, metadata !3, metadata !4, metadata !5}
```

```
!opencl.kernels = !{!0} /*Metadata before LLVM pass*/
!0 = metadata !{void (<2 x float>, float addrspace(1)
   *, float addrspace(1)*)* @addition, metadata !1,
   metadata !2, metadata !3, metadata !4, metadata !5}
!1 = metadata !{metadata !"kernel_arg_addr_space", i32
   0, i32 1, i32 1}
```

```
!opencl.kernels = !{!0} /*Metadata before LLVM pass*/
!0 = metadata !{void (<2 x float>, float addrspace(1)
 *, float addrspace(1)*)* @addition, metadata !1,
 metadata !2, metadata !3, metadata !4, metadata !5}
!1 = metadata !{metadata !"kernel_arg_addr_space", i32
    0, i32 1, i32 1}
!2 = metadata !{metadata !"kernel_arg_access_qual",
    metadata !"none", metadata !"none", metadata !"none
"}
```

```
!opencl.kernels = !{!0} /*Metadata before LLVM pass*/
!0 = metadata !{void (<2 x float>, float addrspace(1)
    *, float addrspace(1)*)* @addition, metadata !1,
    metadata !2, metadata !3, metadata !4, metadata !5}
!1 = metadata !{metadata !"kernel_arg_addr_space", i32
    0, i32 1, i32 1}
!2 = metadata !{metadata !"kernel_arg_access_qual",
    metadata !"none", metadata !"none
    "}
!3 = metadata !{metadata !"kernel_arg_type", metadata
    !"float2", metadata !"float*"}
```

```
!opencl.kernels = !{!0} /*Metadata before LLVM pass*/
 !0 = metadata !{void (<2 x float>, float addrspace(1)
 *, float addrspace(1)*)* @addition, metadata !1,
     metadata !2. metadata !3. metadata !4. metadata !5}
0, i32 1, i32 1}
 | !2 = metadata !{metadata !"kernel_arg_access_qual",
| metadata !"none", metadata !"none", metadata !"none
 !3 = metadata !{metadata !"kernel_arg_type", metadata
     !"float2", metadata !"float*", metadata !"float*"}
o||!4 = metadata !{metadata !"kernel_arg_type_qual"
     metadata !"", metadata !"const", metadata !""}
```

```
!opencl.kernels = !{!0} /*Metadata before LLVM pass*/
|!O = metadata !{void (<2 x float>, float addrspace(1)
    *, float addrspace(1)*)* @addition, metadata !1,
    metadata !2, metadata !3, metadata !4, metadata !5}
!1 = metadata !{metadata !"kernel_arg_addr_space", i32
    0. i32 1. i32 1}
| !2 = metadata !{metadata !"kernel_arg_access_qual",
| metadata !"none", metadata !"none", metadata !"none
!3 = metadata !{metadata !"kernel_arg_type", metadata
    !"float2", metadata !"float*", metadata !"float*"}
|!4 = metadata !{metadata !"kernel_arg_type_qual"
    metadata !"", metadata !"const", metadata !"")
!5 = metadata !{metadata !"kernel_arg_base_type",
metadata !"float2", metadata !"float*", metadata !"
    float*"}
```

Metada before LLVM pass:

```
!3 = metadata !{metadata !"kernel_arg_type",
    metadata !"float2", metadata !"float*",
    metadata !"float*"}
```

Metadata after LLVM pass:

```
!3 = metadata !{metadata !"kernel_arg_type",
   metadata !"float2", metadata !"float*",
   metadata !"float*", metadata !"int",
   metadata !"int"}
```

Aliasing

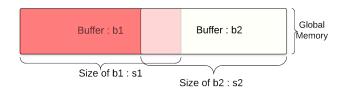


Figure: Example of aliasing

Out Of Bound

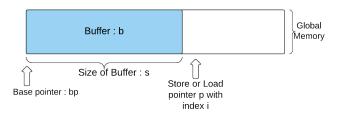


Figure: Example of out-of-bound memory access

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Future work:

- Function call using buffers as arguments inside kernel functions
- Access flag testing
- Concurrent memory access

Application:

- Created to overcome memory management problem
- Next big step: Detecting Direct Memory Access races

End

Thank you for your attention!