

Franck Petit INRIA/LIP6-CNRS/UPMC

 A team of k "weak" robots evolving into a ring of n nodes

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- A team of k"weak" robots evolving into a ring of n nodes
 - Autonomous : No central authority

- A team of k "weak" robots evolving into a ring of n nodes
 - Autonomous
 - Anonymous : Undistinguishable

- A team of k"weak" robots evolving into a ring of n nodes
 - Autonomous
 - Anonymous
 - Oblivious : No mean to know the past

- A team of k "weak" robots evolving into a ring of n nodes
 - Autonomous
 - Anonymous
 - Oblivious
 - Disoriented : No mean to agree on a common direction or orientation

 A team of k "weak" robots evolving into a ring of n nodes

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- A team of k "weak" robots evolving into a ring of n nodes
 - Atomicity

: In every configuration, each robot is located at exactly one node

- A team of k "weak" robots evolving into a ring of n nodes
 - Atomicity
 - Multiplicity
- : In every configuration, each node contains zero, one, or more than one robot

(every robot is able to detect it)

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 - SSM

: In every configuration, k' robots are activated $(0 < k' \le k)$

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 - SSM : In every configuration, k' robots are activated (0 < k' ≤ k)
 - The k' activated robots execute the cycle:
 - Look : Instantaneous snapshot with multiplicity detection

- A team of k "weak" robots evolving into a ring of n nodes
 - SSM

- : In every configuration, k' robots are activated $(0 < k' \le k)$
- The k' activated robots execute the cycle:
 - 1. Look
 - Compute: Based on this observation, decides to either stay idle or move to one of the neighboring nodes

- A team of k "weak" robots evolving into a ring of n nodes
 - SSM : In every configuration, k' robots are activated (0 < k' ≤ k)
 - The k' activated robots execute the cycle:
 - 1. Look
 - 2. Compute
 - 3. Move : Move toward its destination

Starting from a configuration where no two robots are located at the same node:

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Exploration:

Each node must be visited by at least one robot

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Termination:

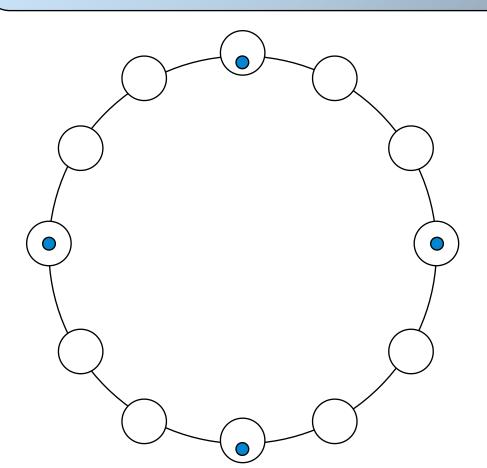
Eventually, every robot stays idle

Starting from a configuration where no two robots are located at the same node:

- Exploration:
 - Each node must be visited by at least one robot
- Termination:
 - Eventually, every robot stays idle
- Performance: Number of robots (k<n)

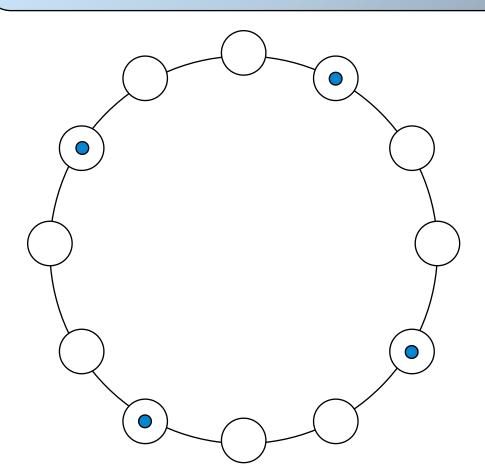
Lower Bound (1/2)

Deterministic Exploration impossible if $k \mid n$



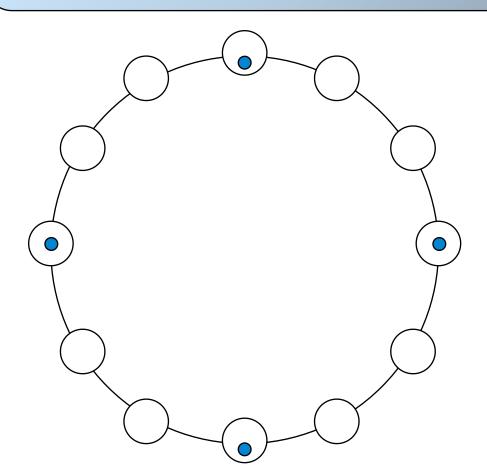
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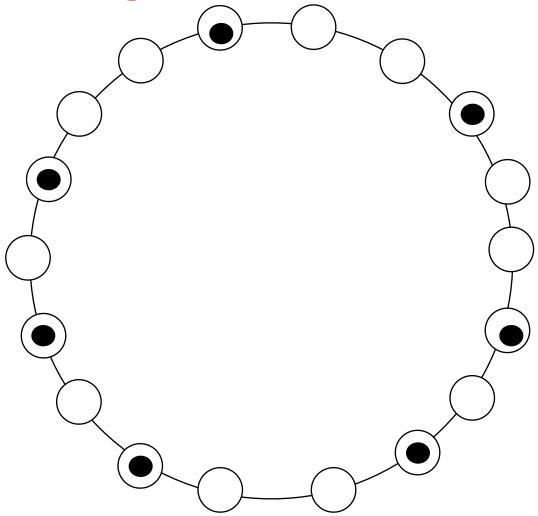


Lower Bound (1/2)

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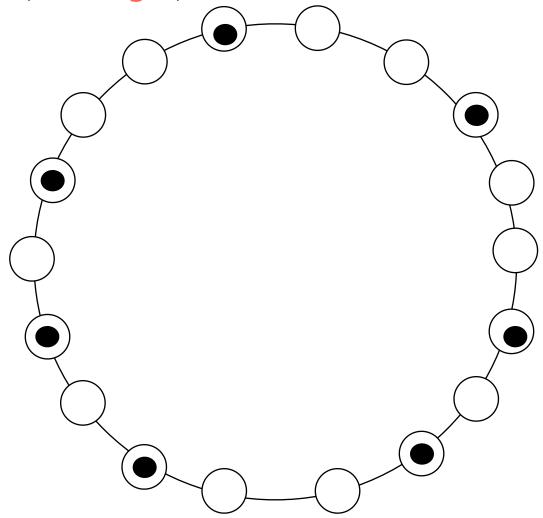
 $k \nmid n$, $k > \log n$, and n > 16



 $^{\circ}$ $k \nmid n$, $k > \log n$, and n > 16Interdistance (d): Minimum distance taken over all pairs of distinct robots.

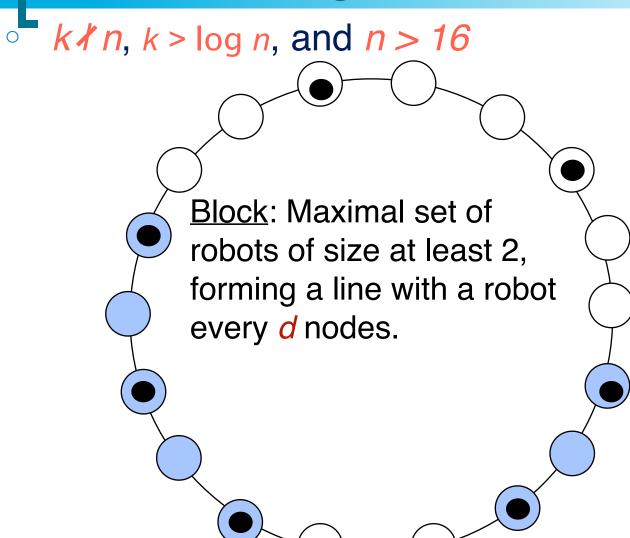
 $^{\circ}$ $k \nmid n$, $k > \log n$, and n > 16Interdistance (d): Minimum distance taken over all pairs of distinct debots. d = 2d = 2.

 $k \nmid n$, $k > \log n$, and n > 16





Block: Maximal set of robots of size at least 2, forming a line with a robot every *d* nodes.



 $k \nmid n, k > \log n, \text{ and } n > 16$

- $^{\circ}$ $k \nmid n$, $k > \log n$, and n > 16
- Setup Phase:

Goal: Transform the (arbitrary) initial configuration into a configuration of interdistance 1 where there is a single block or two blocks of the same size.

Method: Decrease the number of blocks whenever possible. Otherwise, decrease the interdistance.

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Tower Phase:

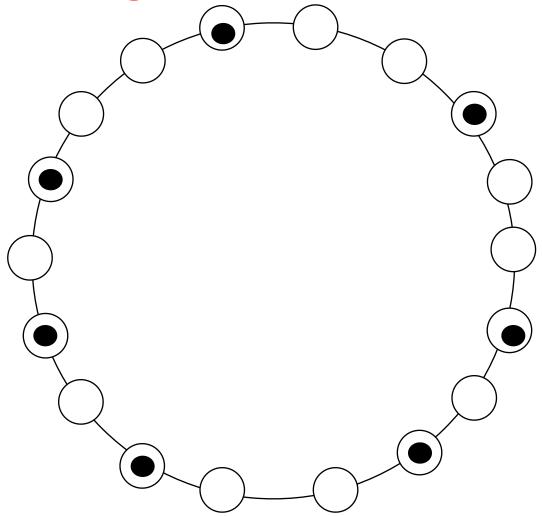
Goal: Create one or two multiplicities inside each block.

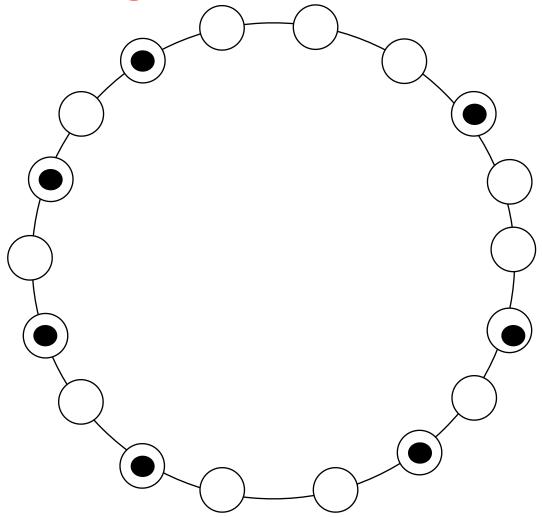
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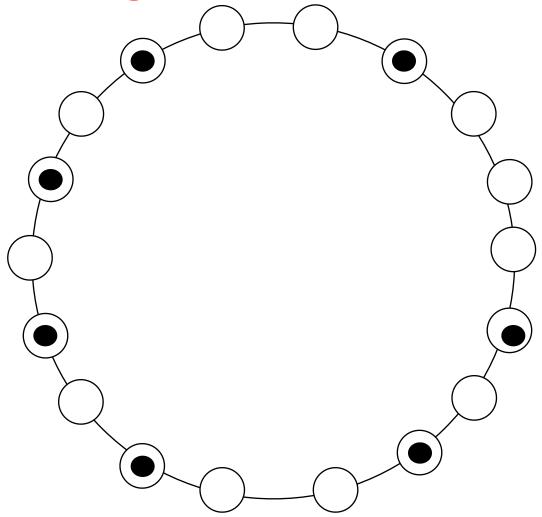
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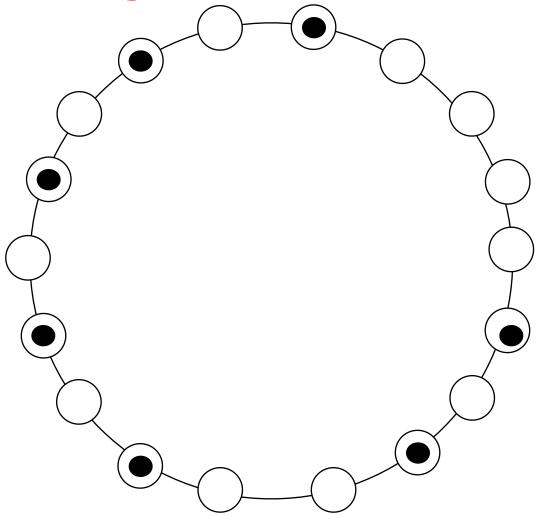
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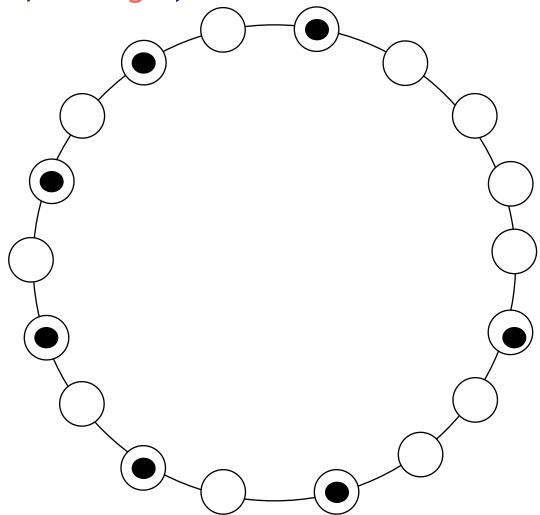
- Tower Phase:
 - Goal: Create one or two multiplicities inside each block.
- Exploration Phase:
 - Goal: Perform exploration until reaching an identified final configuration.

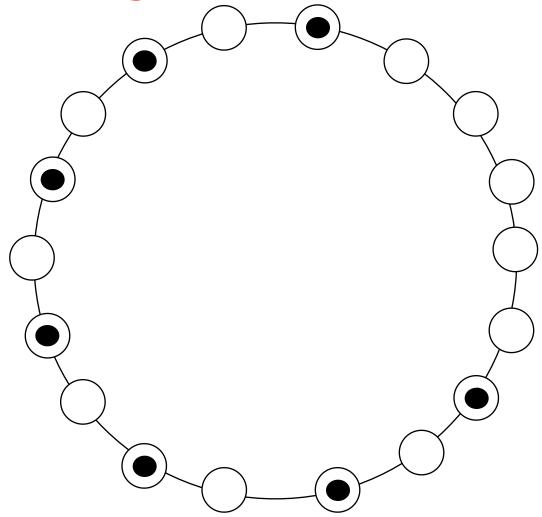


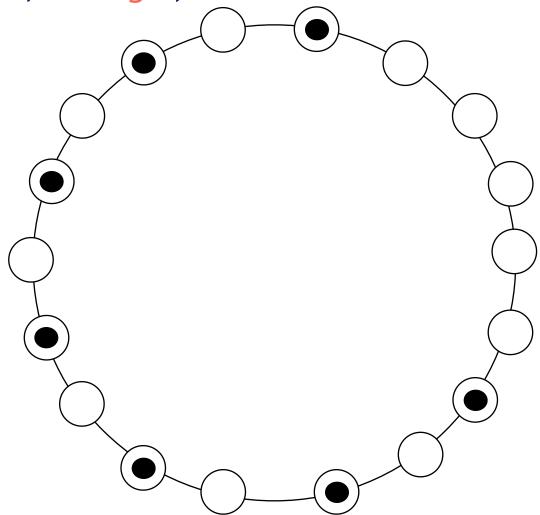


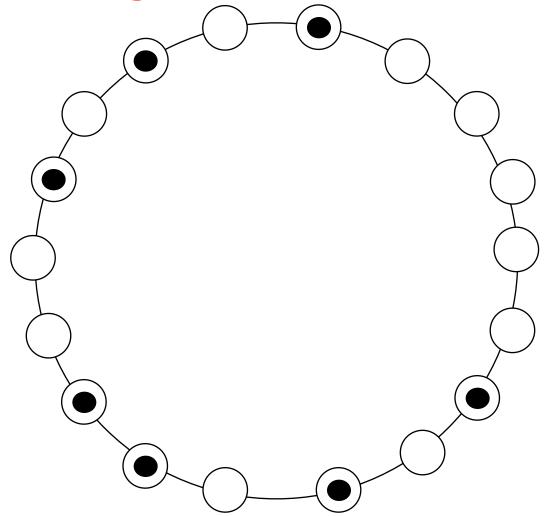


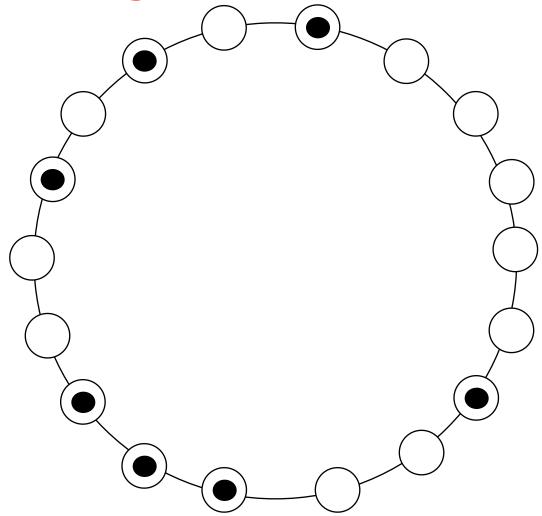


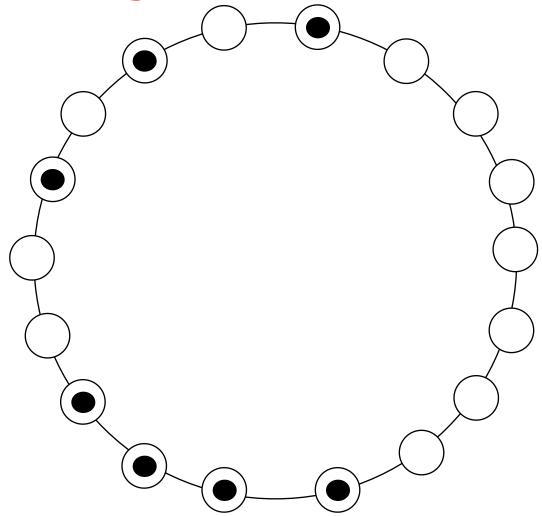


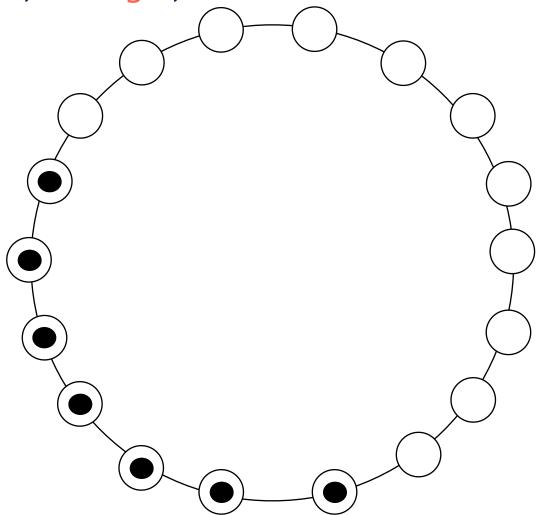


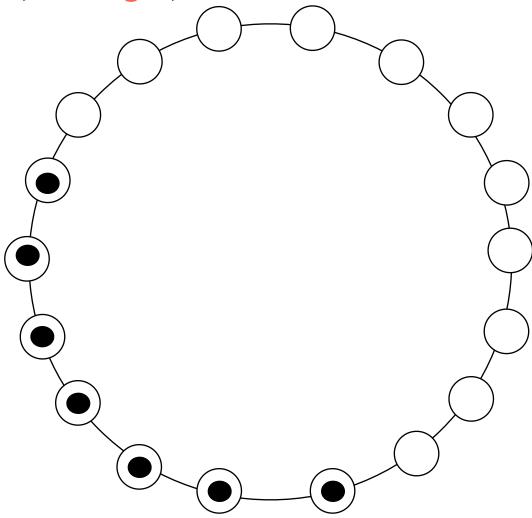


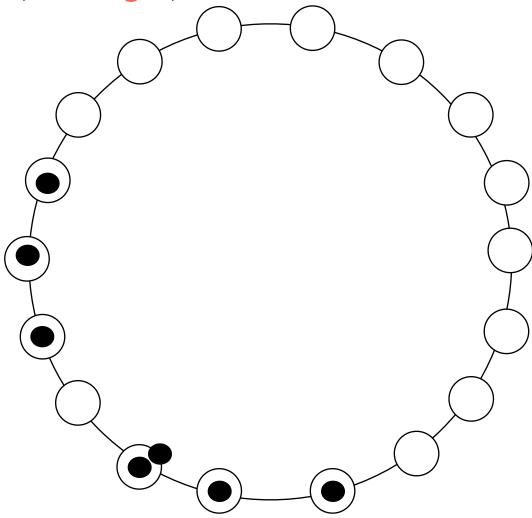


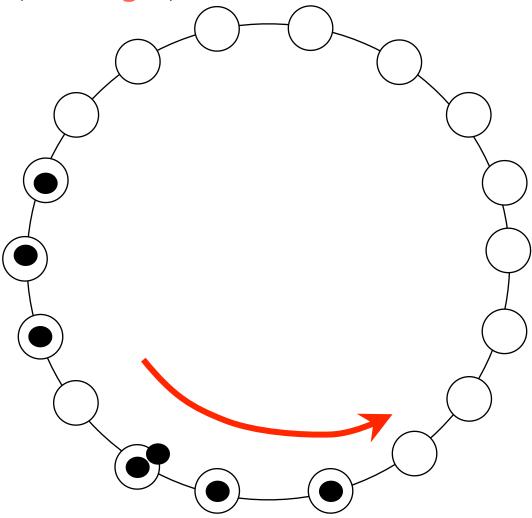


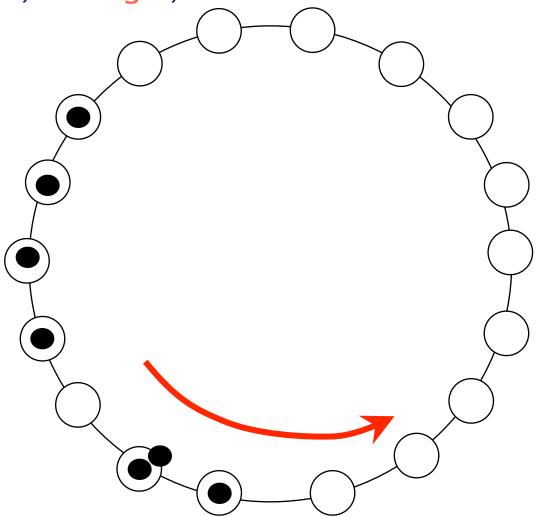












Minimal Number of

Robots?

Theorem.

4 probabilistic robots are necessary and sufficient, provided that n > 4

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1. Exploration impossible with less than 4 robots

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- 1. Exploration impossible with less than 4 robots
- 2. An algorithm working with 4 probabilistic robots (n > 4)

Theorem.

4 probabilistic robots are necessary and sufficient, provided that n > 4

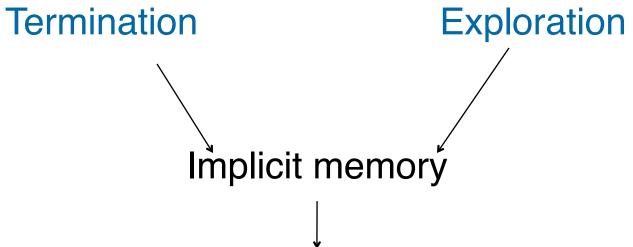
The theorem holds even if k divides n.

- Exploration impossible with less than 4 robots
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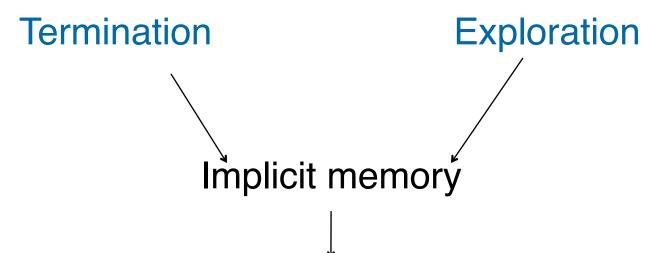
Termination

Exploration





At least one configuration that cannot be an initial configuration



At least one configuration that cannot be an initial configuration

Remark.

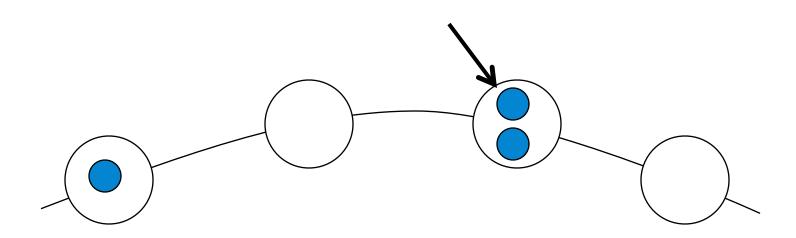
If n > k, any terminal configuration of any protocol contains at least one *tower*.

Definition.

A node with at least two robots.

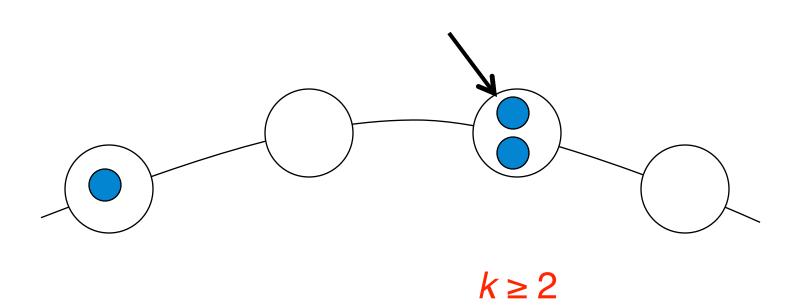
Definition.

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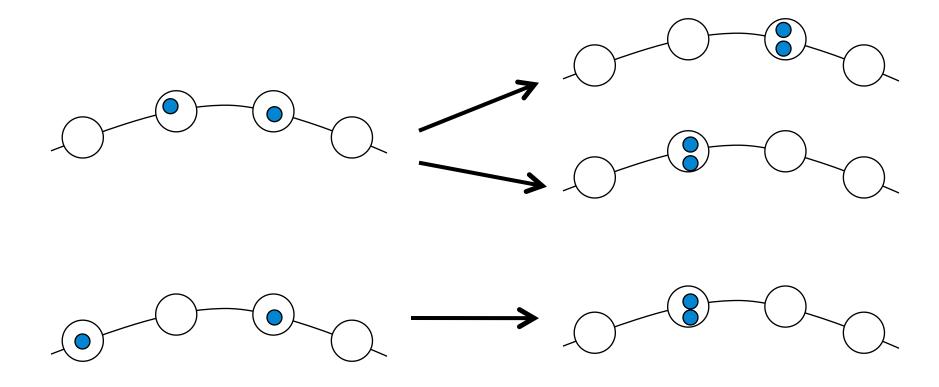
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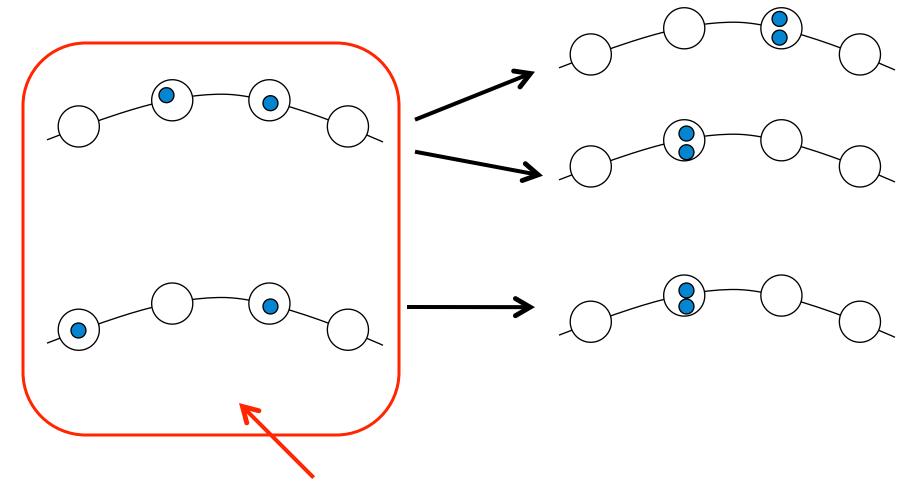


Tower Building

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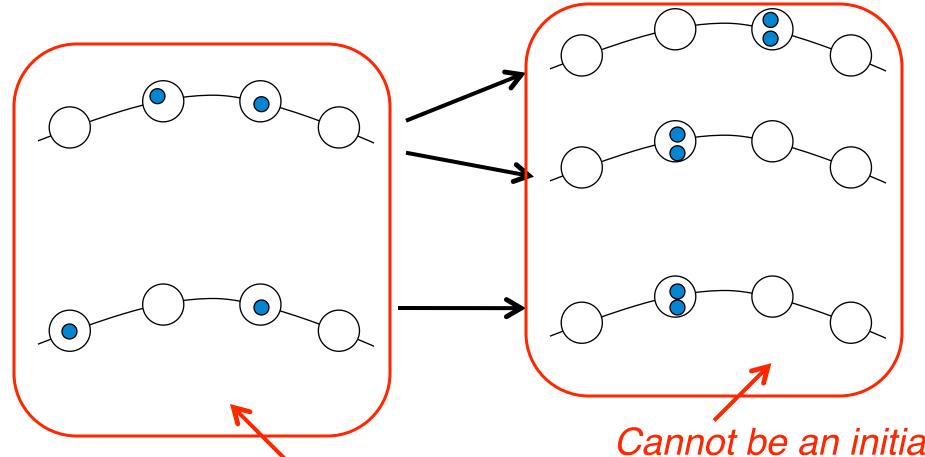


Tower Building



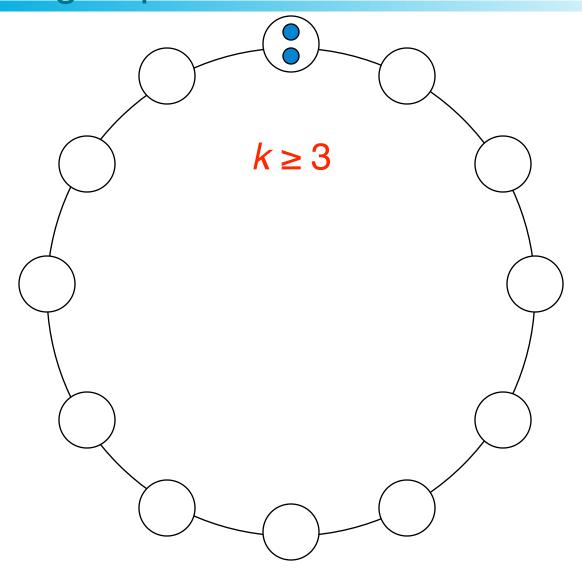
Can be an initial configuration

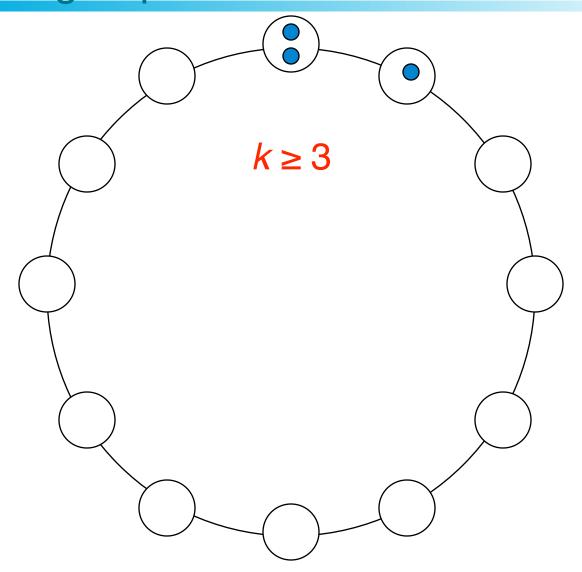
Tower Building

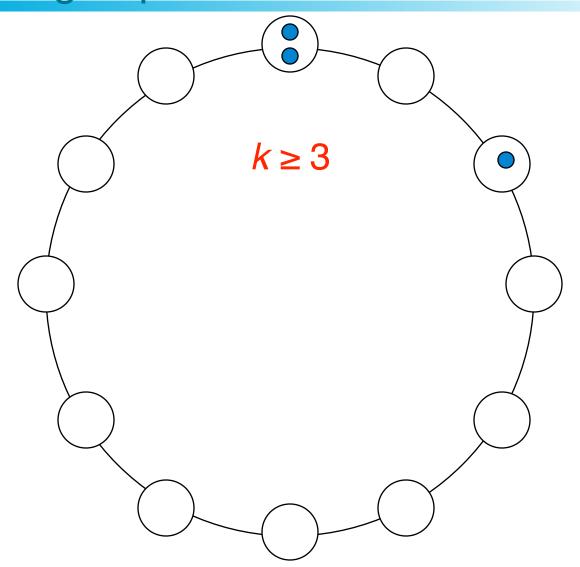


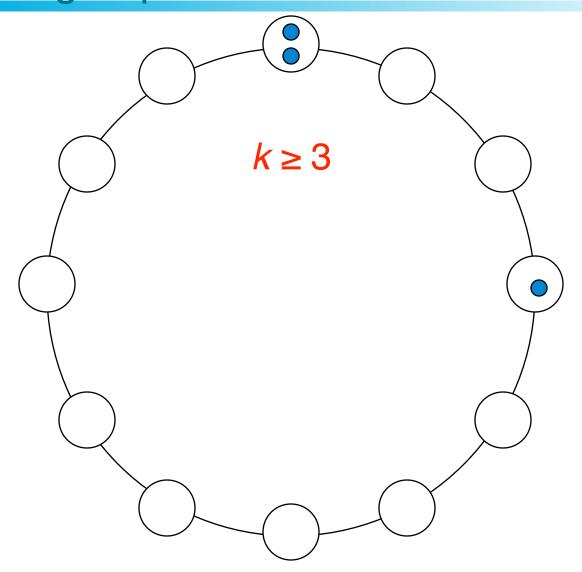
Can be an initial configuration

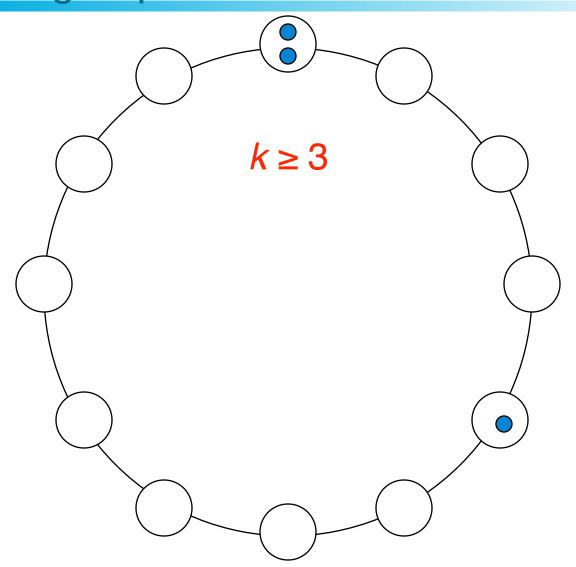
configuration

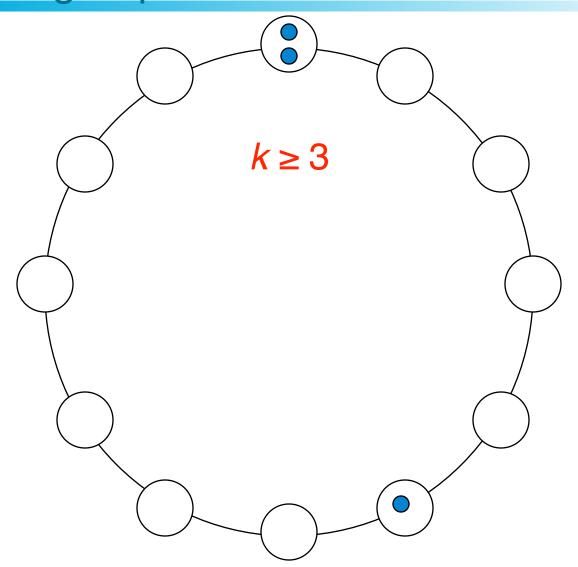


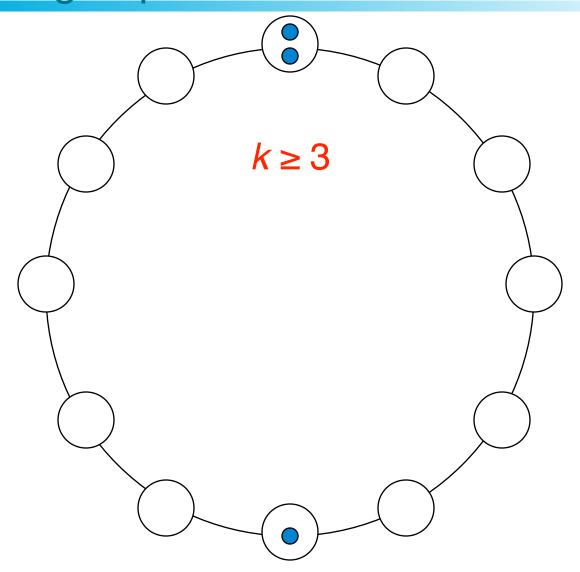


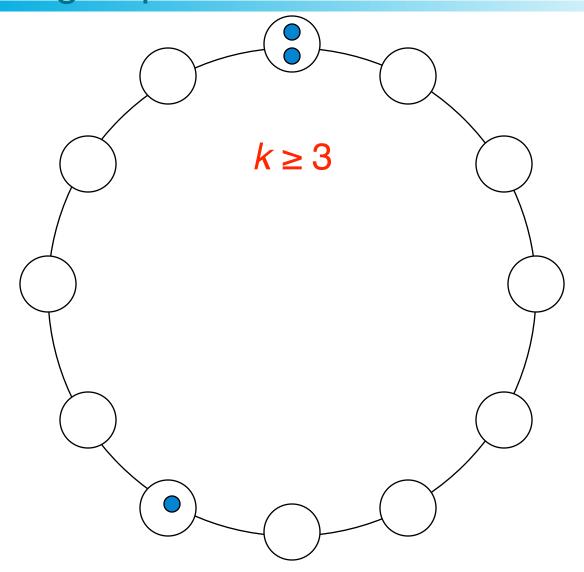


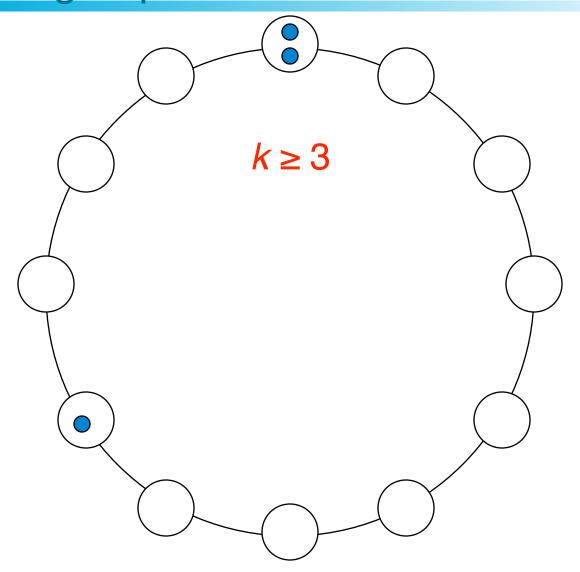


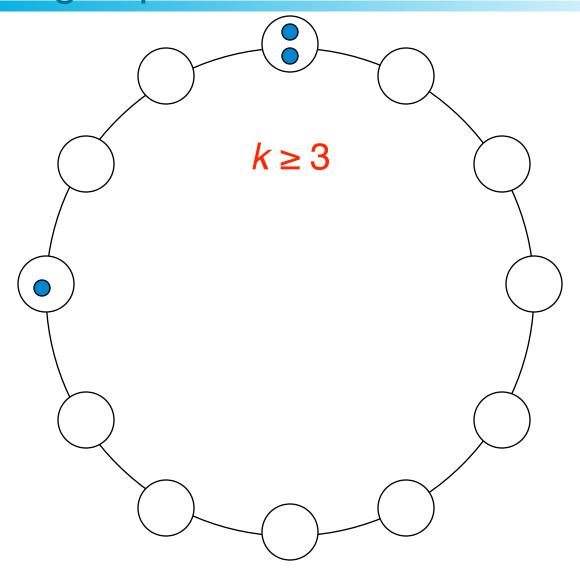


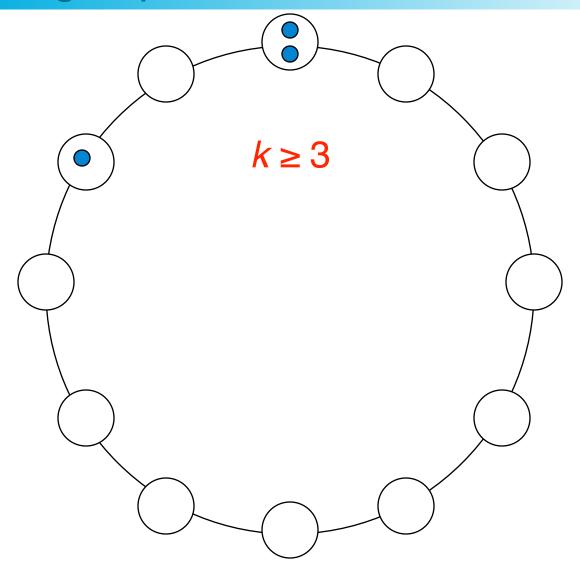


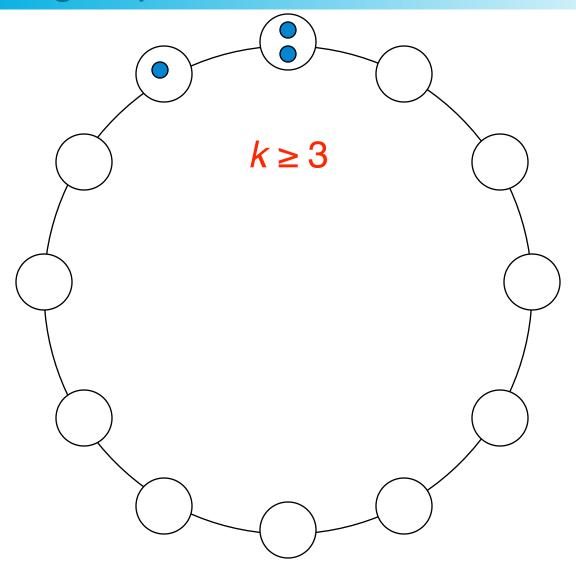


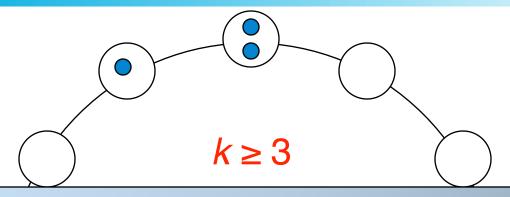






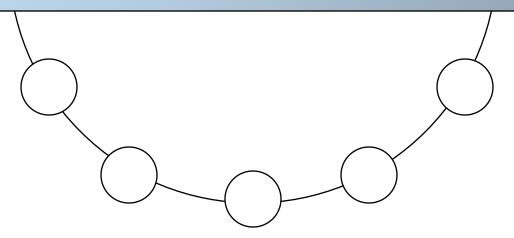


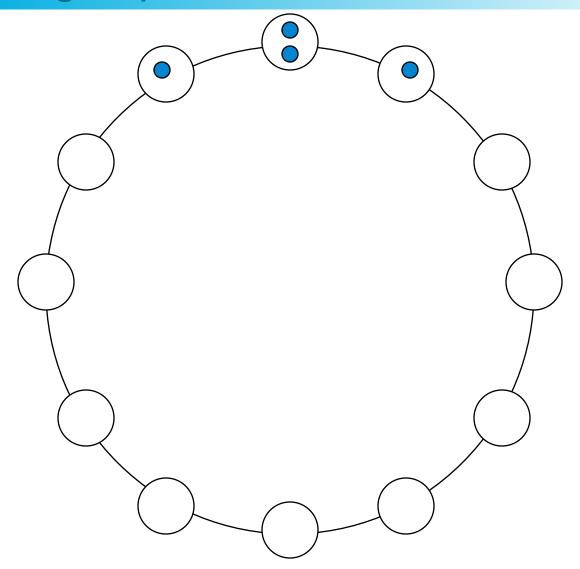


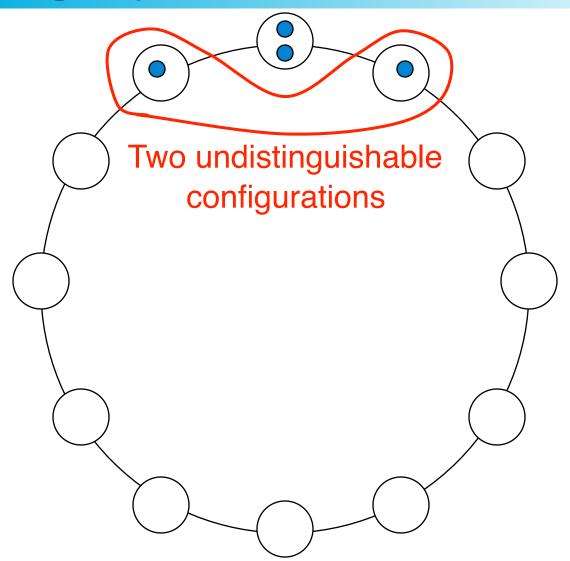


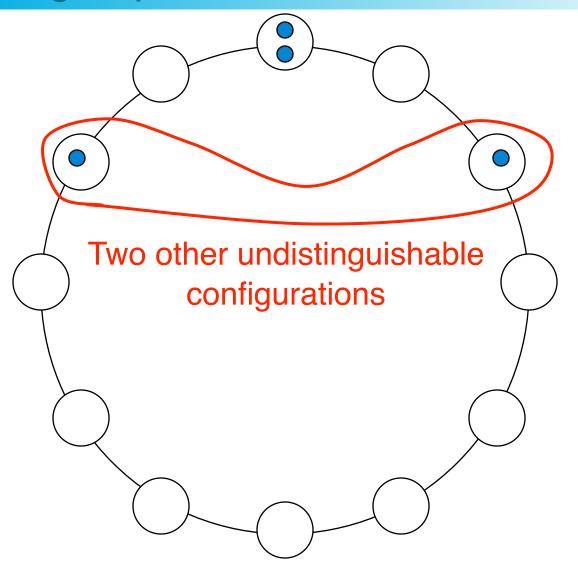
Lemma.

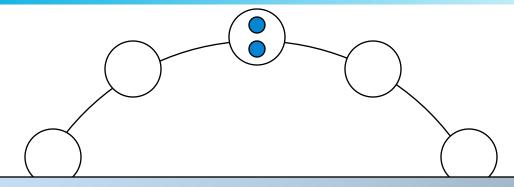
Every execution must contain a suffix of at least n-k+1 configurations containing a tower of less than k robots and any two of them are distinguishable.





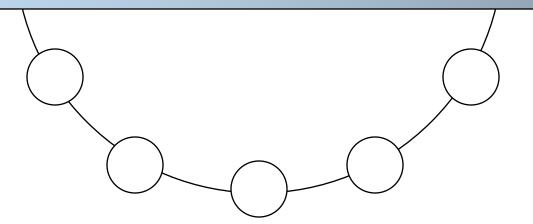






<u>Lemma.</u>

With 3 robots and a fixed tower of 2 robots, the maximum number of distinguishable configurations is equal to $\frac{n}{2}$



Theorem.

For every n > 4, there exists no exploration protocol (even probabilistic) of an n-size ring with 3 robots.

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Proof:

$$\left\lfloor \frac{n}{2} \right\rfloor \ge n - k + 1 \Rightarrow n \le 4$$

Contribution

Theorem.

4 probabilistic robots are necessary and sufficient, provided that n > 4

- The theorem holds even if k divides n.
- 1. Exploration impossible with less than 4 robots
- 2. Give an algorithm working with 4 probabilistic robots

Contribution

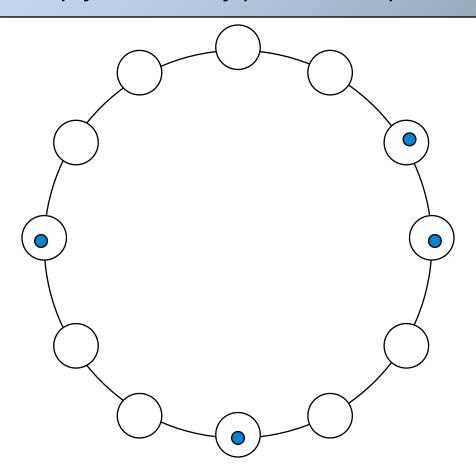
Theorem.

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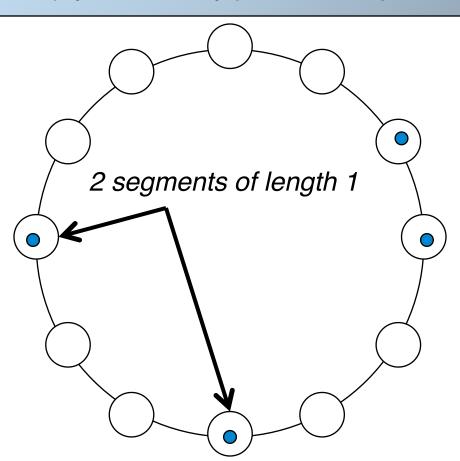
- The theorem holds even if *k* divides *n*.
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Segment.

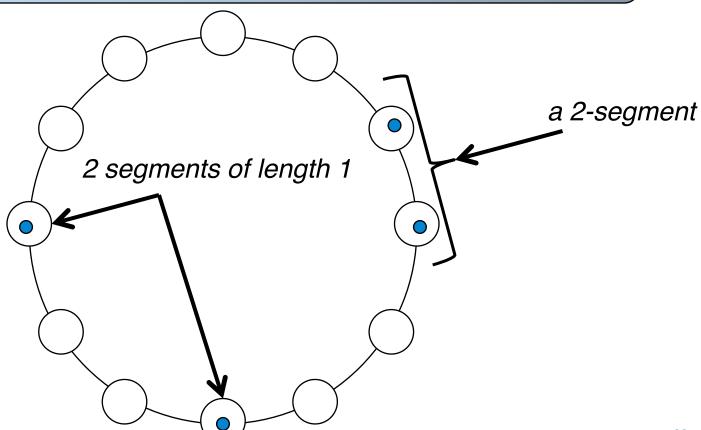
Segment.

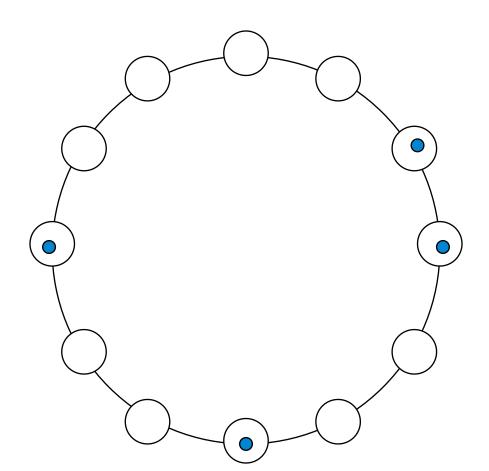


Segment.

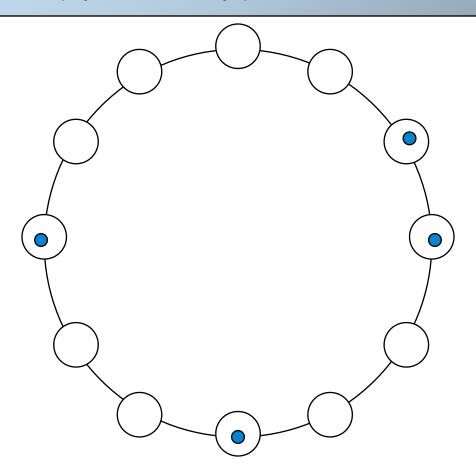


Segment.

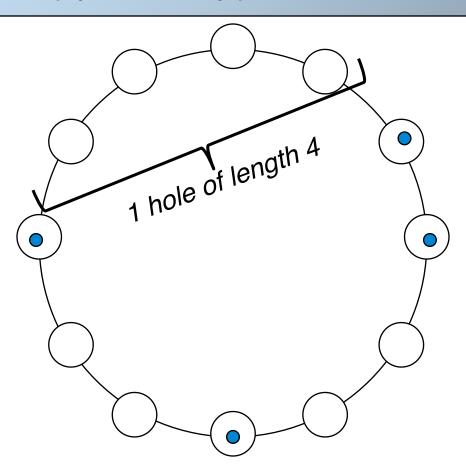




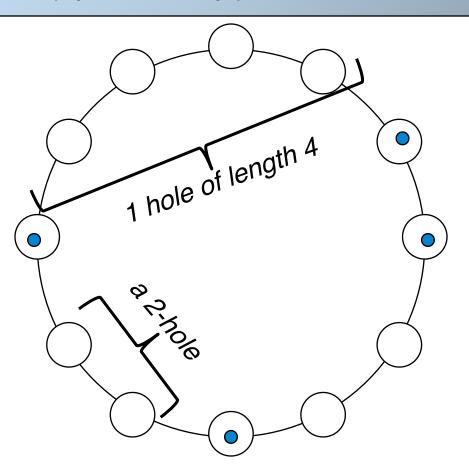
Hole.

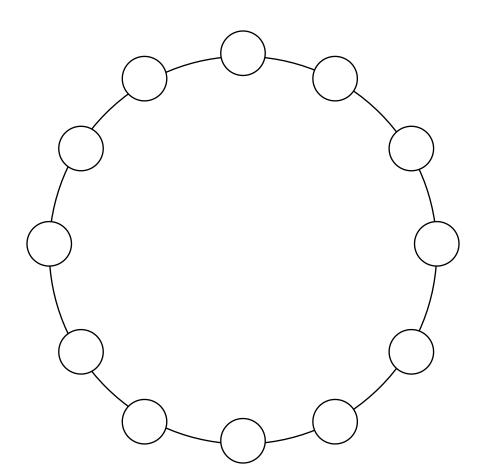


Hole.

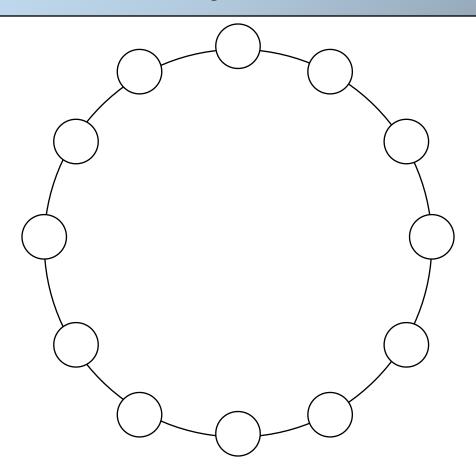


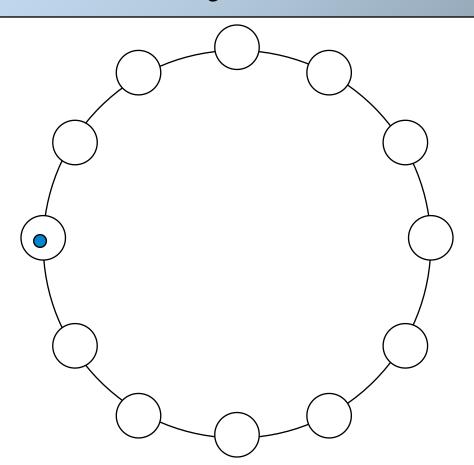
Hole.

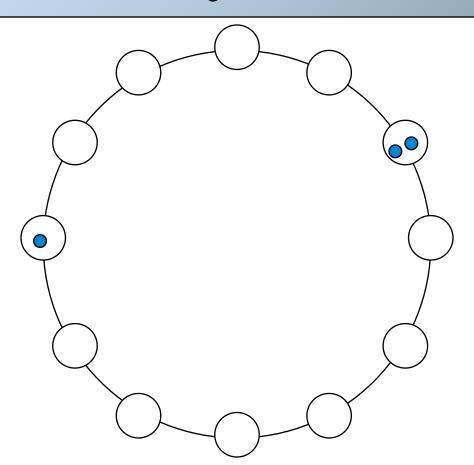


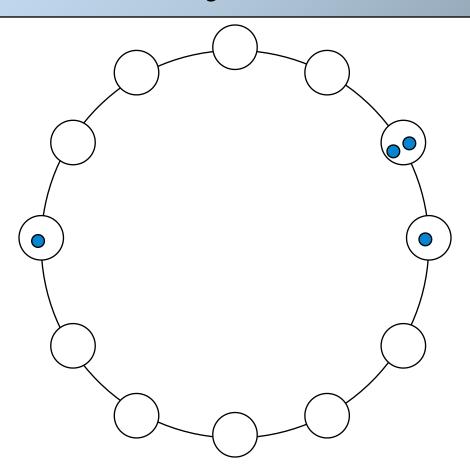


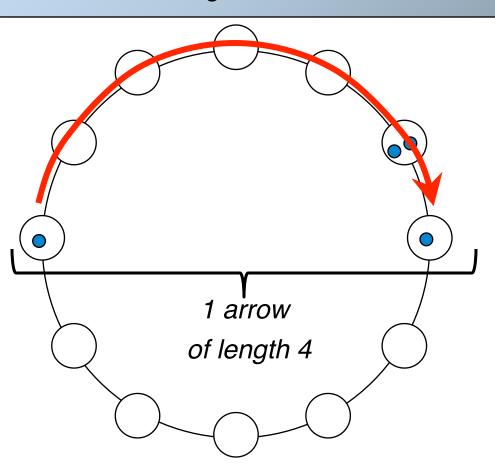
Arrow. A 1-segment, followed by a non-empty elementary path of free nodes, a tower, and a 1-segment.

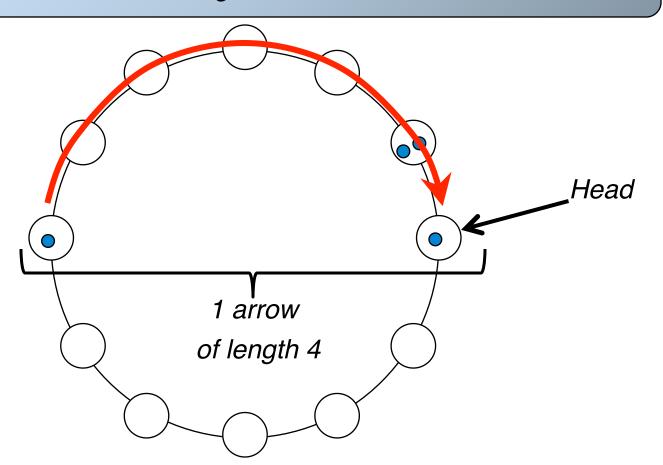


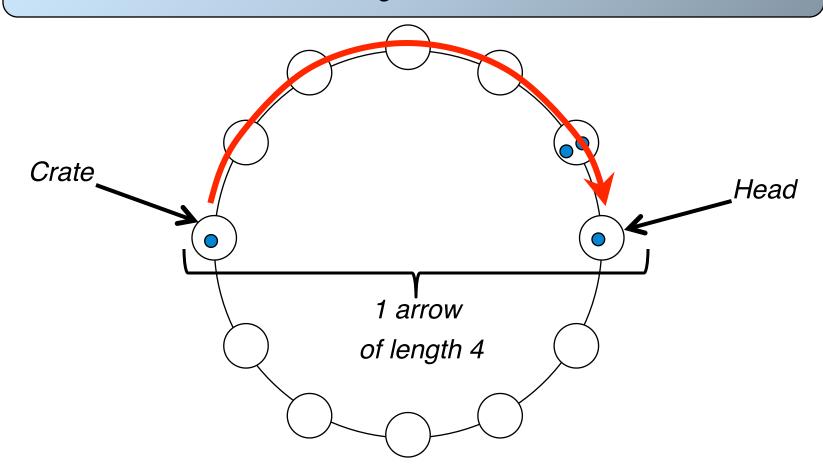


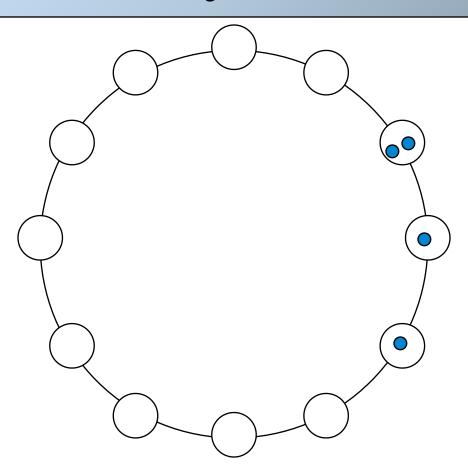


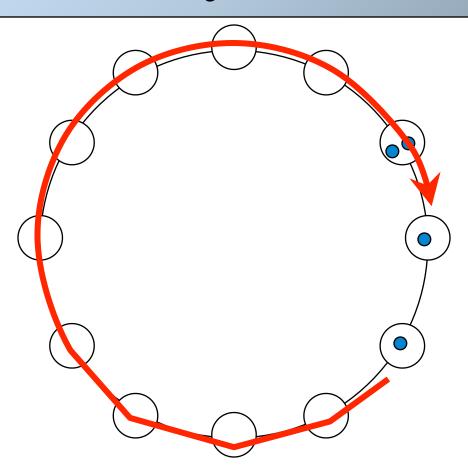


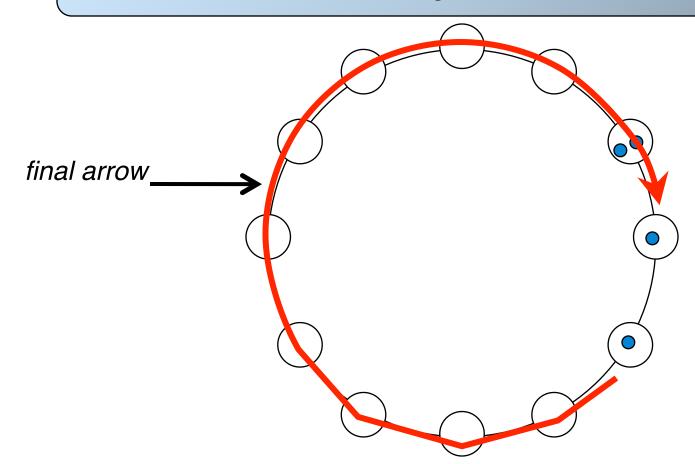


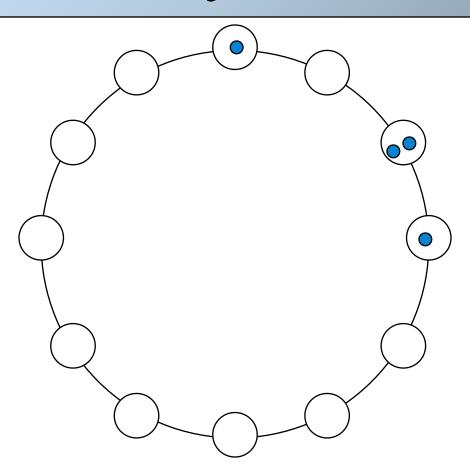


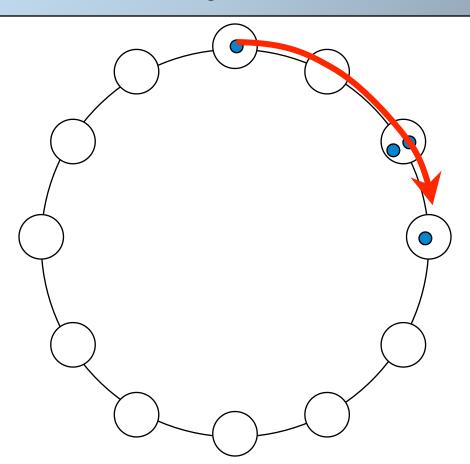


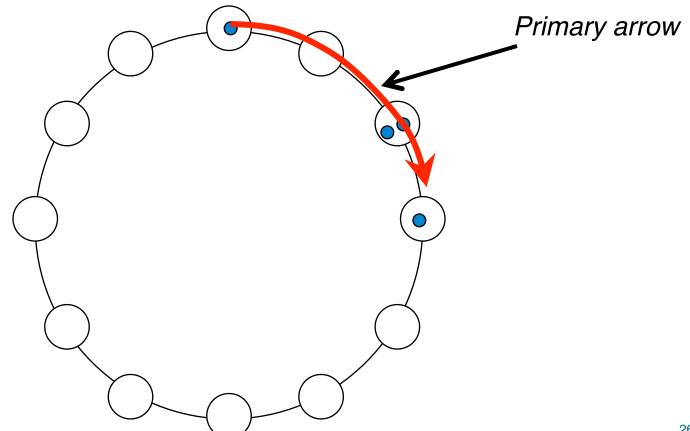












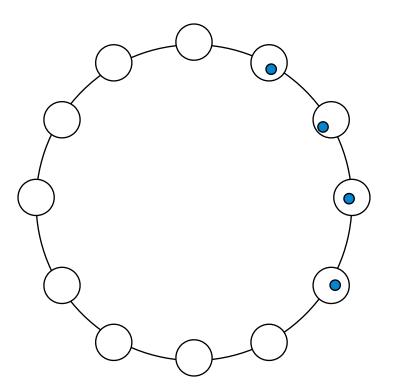
Initially, there is no tower

Converge toward a 4-segment

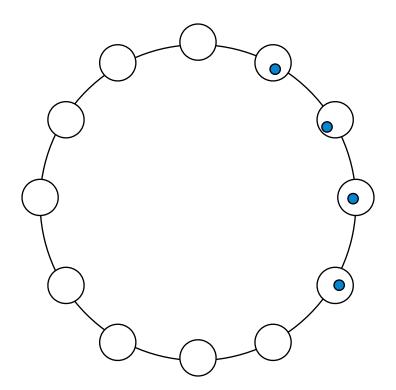
- 1. Converge toward a 4-segment
- 2. Build a tower

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- 2. Build a tower
- 3. Visit the ring and terminate

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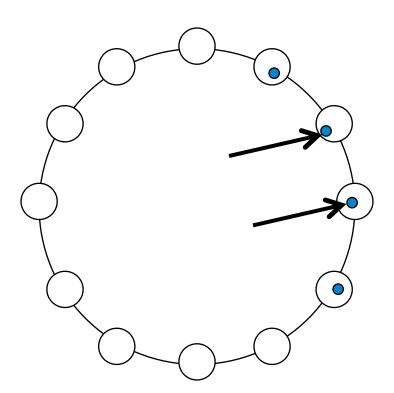
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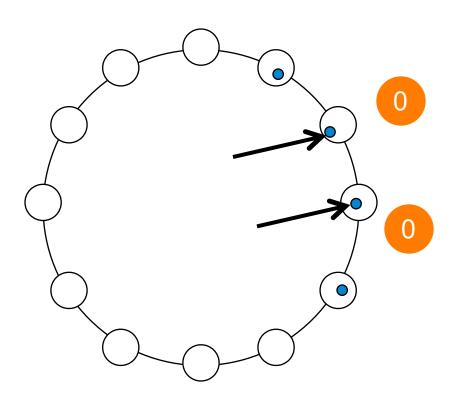
If I am an internal node, then I try to move on the other internal node.



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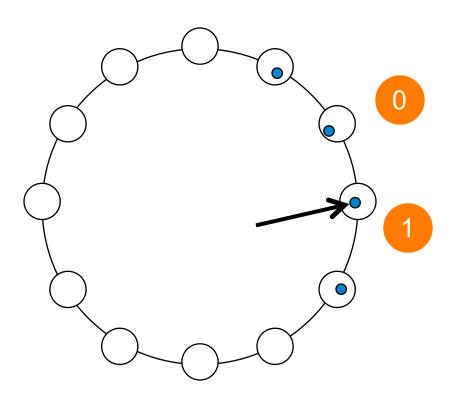
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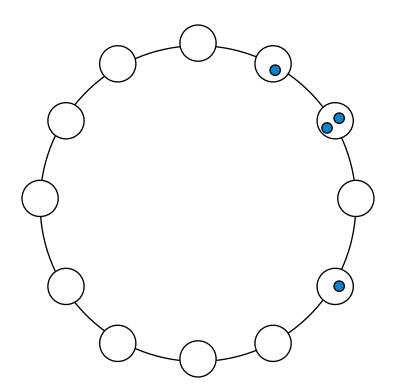
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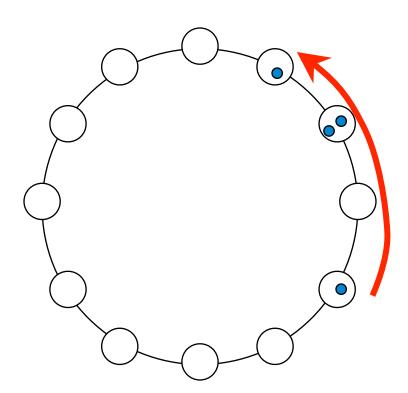
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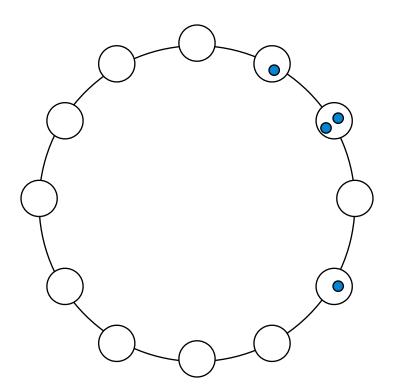
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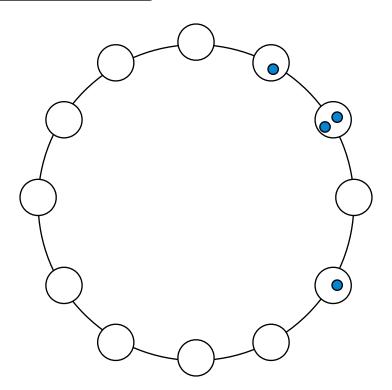
- Converge toward a 4-segment
- 2. Build a tower
- → Primary arrow
- 3. Visit the ring and terminate



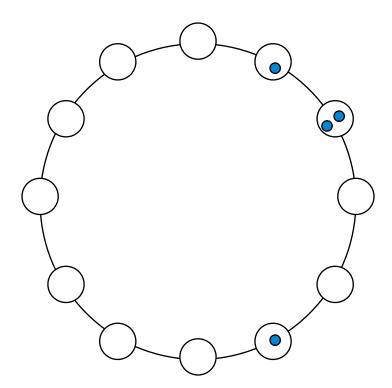
- Converge toward a 4-segment
- 2. Build a tower
 - Primary arrow
- 3. Visit the ring and terminate



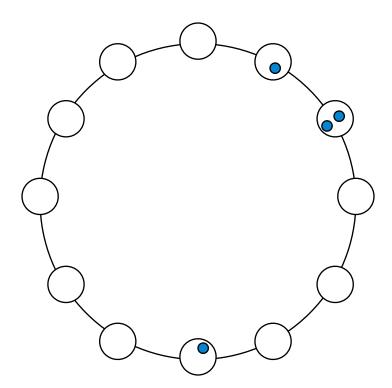
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- Visit the ring and terminate



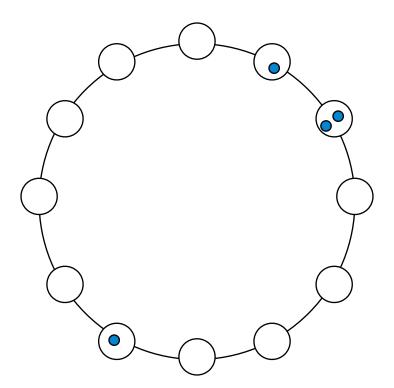
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate



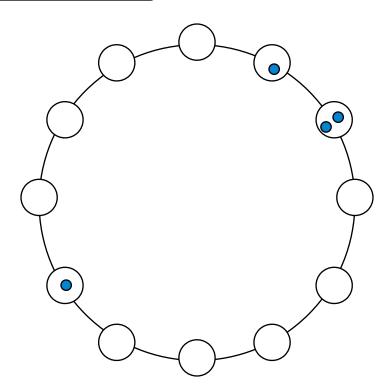
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate



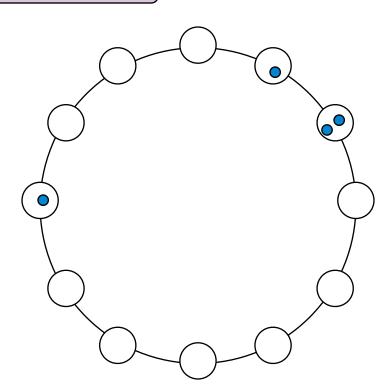
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate



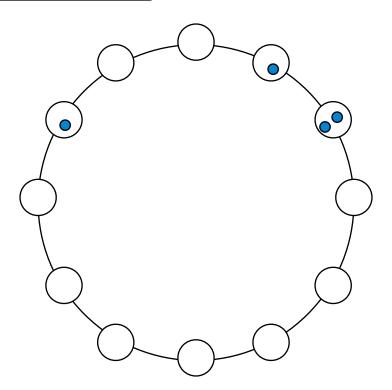
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate



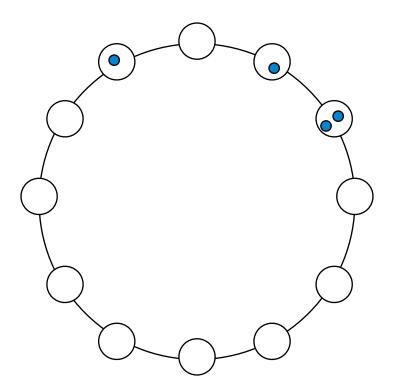
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- Visit the ring and terminate



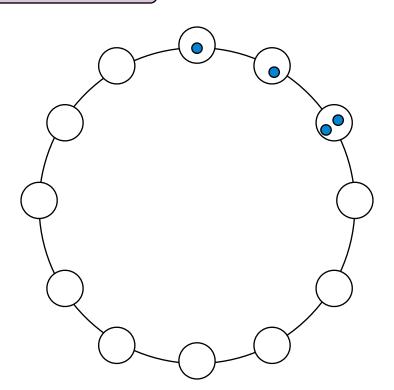
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate



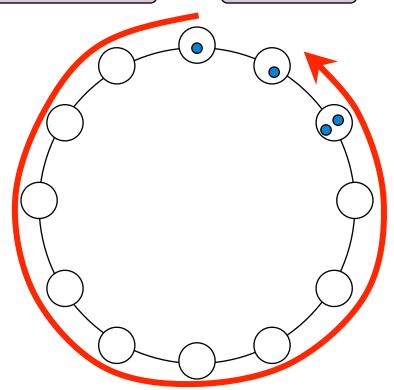
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate



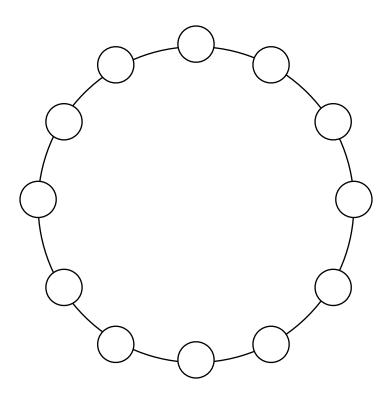
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate



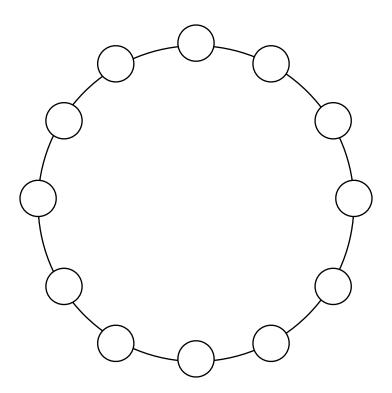
- Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate → Final arrow



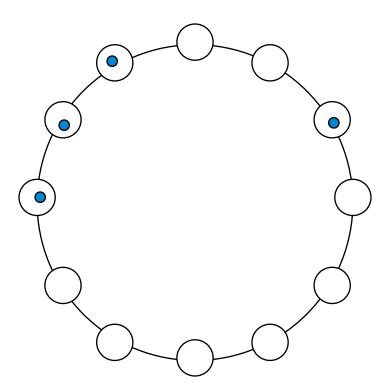
- Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow



- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow



- 1. Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- Visit the ring and terminate → Final arrow
- a) 3-segment

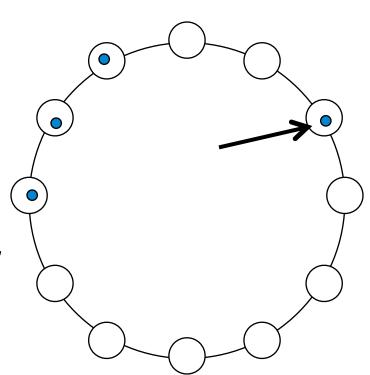


Initially, there is no tower

- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- Visit the ring and terminate → Final arrow

a) 3-segment

If I am the isolated node, then I move through a shortest hole.

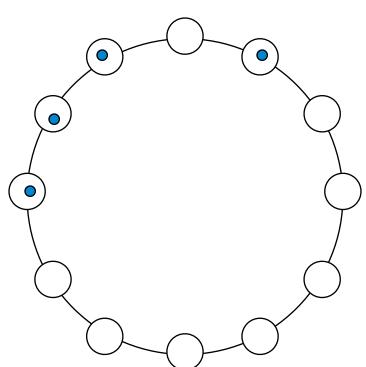


Initially, there is no tower

- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow

a) 3-segment

If I am the isolated node, then I move through a shortest hole.

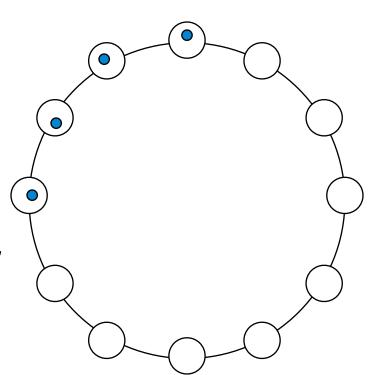


Initially, there is no tower

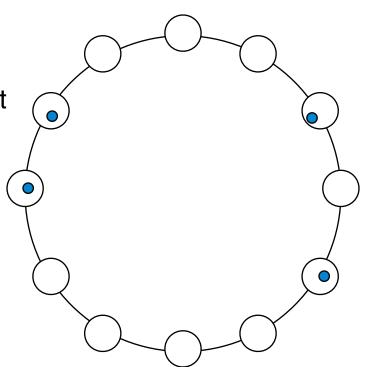
- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- Visit the ring and terminate → Final arrow

a) 3-segment

If I am the isolated node, then I move through a shortest hole.



- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- Visit the ring and terminate → Final arrow
- a) 3-segment
- b) a unique 2-segment

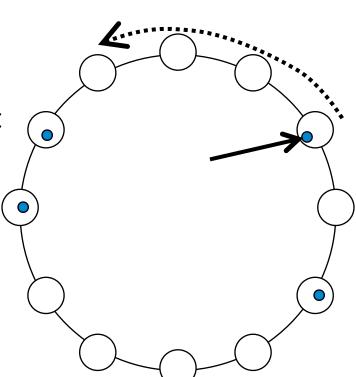


Initially, there is no tower

- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- Visit the ring and terminate → Final arrow

- a) 3-segment
- b) a unique 2-segment

If I am at the closest distance from the 2segment, then I move toward the closest extremity.

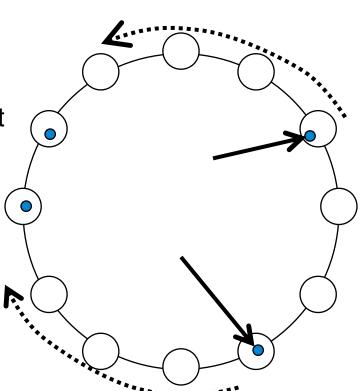


Initially, there is no tower

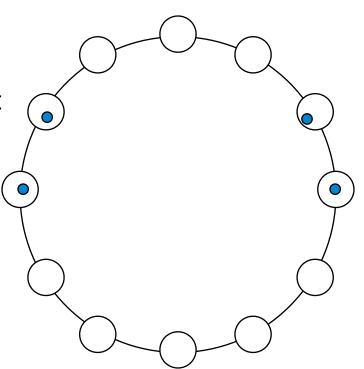
- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow

- a) 3-segment
- b) a unique 2-segment

If I am at the closest distance from the 2segment, then I move toward the closest extremity.



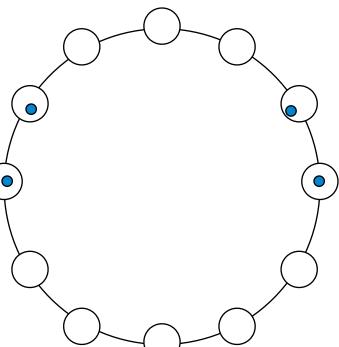
- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow
- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments



Initially, there is no tower

- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow
- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments

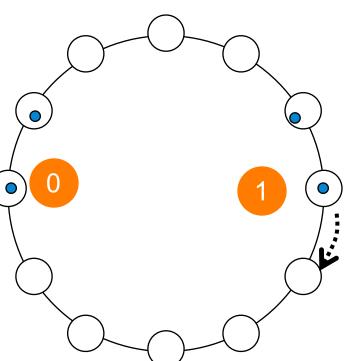
If I am a neighbor of the longest hole, then I try to move toward the other 2-segment.



Initially, there is no tower

- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow
- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments

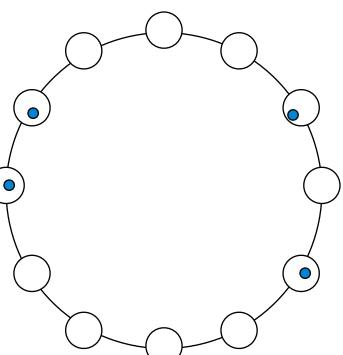
If I am a neighbor of the longest hole, then I try to move toward the other 2-segment.



Initially, there is no tower

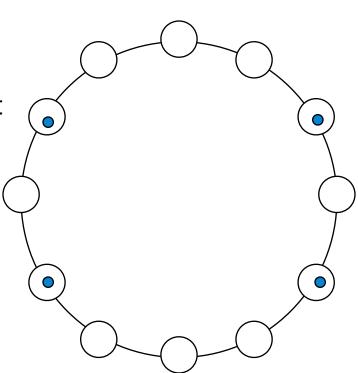
- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow
- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments

If I am a neighbor of the longest hole, then I try to move toward the other 2-segment.

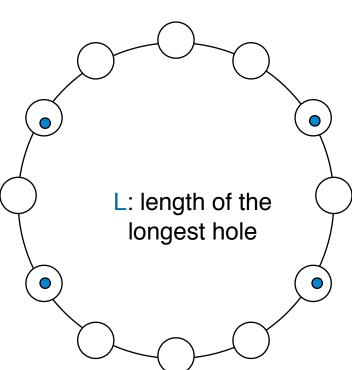


- 1. Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- Visit the ring and terminate
 Final arrow

- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments
- d) four isolated nodes



- 1. Converge toward a 4-segment
- 2. Build a tower
- Primary arrow
- 3. Visit the ring and terminate → Final arrow
- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments
- d) four isolated nodes

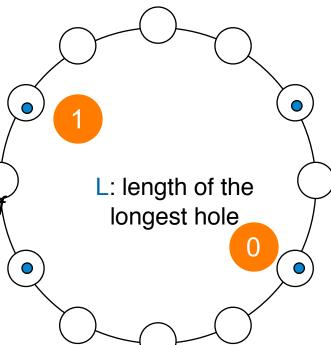


Initially, there is no tower

- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow

- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments
- d) four isolated nodes

If 4 robots are neighbors of an L-hole, then I try to move through my longest neighboring hole.

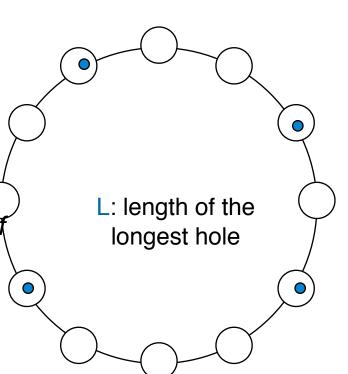


Initially, there is no tower

- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- 3. Visit the ring and terminate → Final arrow

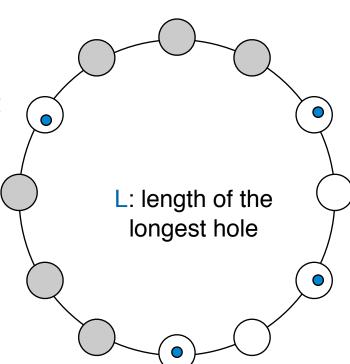
- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments
- d) four isolated nodes

If 4 robots are neighbors of an L-hole, then I try to move through my longest neighboring hole.



- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- Visit the ring and terminate → Final arrow

- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments
- d) four isolated nodes

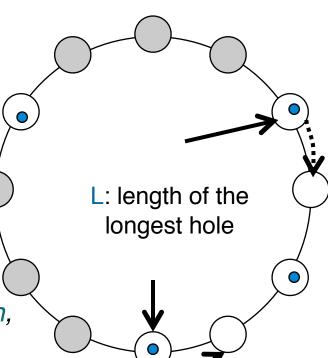


Initially, there is no tower

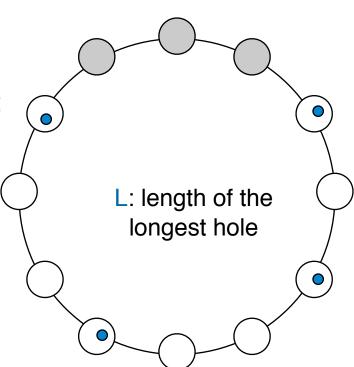
- 1. Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- Visit the ring and terminate → Final arrow

- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments
- d) four isolated nodes

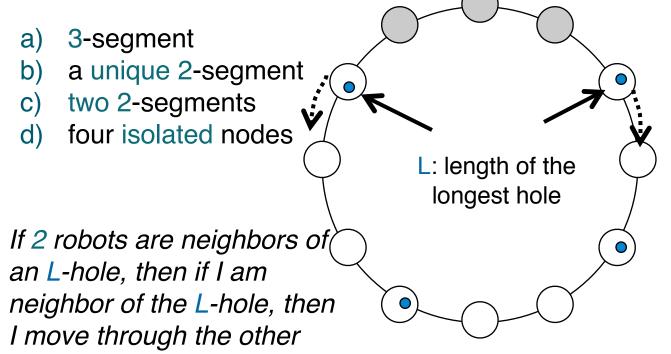
If 3 robots are neighbors of an L-hole, then if I am one of this 3 robots and a neighbor of a smaller hole h, then I move through h.



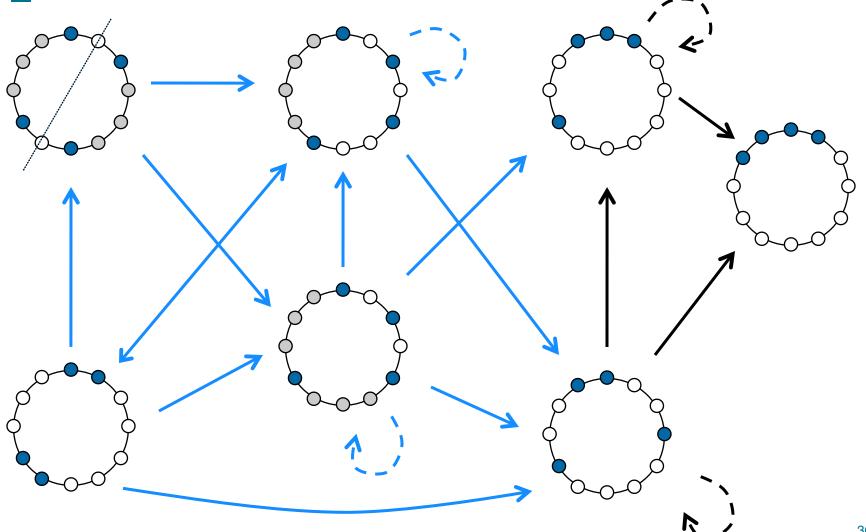
- 1. Converge toward a 4-segment
- 2. Build a tower
 - Primary arrow
- 3. Visit the ring and terminate → Final arrow
- a) 3-segment
- b) a unique 2-segment
- c) two 2-segments
- d) four isolated nodes



- Converge toward a 4-segment
- 2. Build a tower → Primary arrow
- Visit the ring and terminate → Final arrow



Phase 1, Summary



<u>Lemma</u>.

No tower is created during Phase 1 in a n-ring with n > 8.

Lemma.

No tower is created during Phase 1 in a n-ring with n > 8.

Proof Bas:

With n > 8 and 4 robots, there always exists a

hole of length greater than 1.

<u>Lemma</u>.

No tower is created during Phase 1 in a n-ring with n > 8.

Lemma.

No tower is created during Phase 1 in a n-ring with n > 8.

<u>Lemma.</u>

Starting from any initial configuration, the system reaches in finite expected time a configuration containing a 4-segment.

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No tower is created during Phase 1 in a n-ring with n > 8.

<u>Lemma</u>.

Starting from any initial configuration, the system reaches in finite expected time a configuration containing a 4-segment.

Theorem.

The algorithm (Phases 1 to 3) is a probabilistic exploration protocol for 4 robots in a ring of n > 8 nodes.

✓ Ring

n: Number of nodes

n: Number of nodes

- ✓ Ring [Flocchini et al., OPODIS 2007]
 - Deterministic exploration impossible if k divides n (except if k = n)
 - Asynchronous deterministic algorithm with k > 16
 - Deterministic or probabilistic exploration impossible if k < 4
 - Probabilistic algorithm impossible in asynchronous settings
 - Optimal Semi-synchronous Probabilistic Algorithm
 - Deterministic exploration impossible if k < 5 and n even
 - Optimal asynchronous deterministic algorithm, k = 5 and n even
 - Optimal semi-synchronous deterministic algorithm, k = 4 and n odd

n: Number of nodes

- ✓ Ring [Flocchini et al., OPODIS 2007] [Devismes et al., SIROCCO 2009]
 - Deterministic exploration impossible if k divides n (except if k = n)
 - Asynchronous deterministic algorithm with k > 16
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n: Number of nodes

- ✓ Ring [Flocchini et al., OPODIS 2007] [Devismes et al., SIROCCO 2009] [Lamani et al., SIROCCO 2010]
 - Deterministic exploration impossible if k divides n (except if k = n)
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 - Optimal Semi-synchronous Probabilistic Algorithm
 - Deterministic exploration impossible if k < 5 and n even
 - Optimal asynchronous deterministic algorithm, k = 5 and n even
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n: Number of nodes

k: Number of agents

✓ Ring [Flocchini et al., OPODIS 2007] [Lamani et al., SIROCCO 2010]

[Devismes et al., SIROCCO 2009]

n: Number of nodes

- ✓ Ring [Flocchini et al., OPODIS 2007] [Devismes et al., SIROCCO 2009] [Lamani et al., SIROCCO 2010]
- ✓ Tree [Flocchini et al., SIROCCO 2008]
 - Asynchronous deterministic algorithm for trees with maximum degree equal to 3: $k \in \Theta$ (log $n/\log \log n$)
 - Arbitrary tree: $k \in \Theta (\log n)$
- ✓ Chain [Flocchini et al., IPL 2011]
 - Characterization of k: k = 3, k > 4, or k = 4 and n odd

n: Number of nodes

- ✓ Ring [Flocchini et al., OPODIS 2007] [Devismes et al., SIROCCO 2009] [Lamani et al., SIROCCO 2010]
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 - Arbitrary tree: $k \in \Theta (\log n)$
- ✓ Chain [Flocchini et al., IPL 2011]
 - Characterization of k: k = 3, k > 4, or k = 4 and n odd
- ✓ Grid [Devismes et al., SSS 2012]
 - Deterministic or probabilistic exploration impossible if k < 2
 - Optimal Semi-synchronous Deterministic Algorithm, k = 3