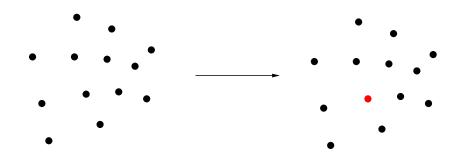
Leader Election in Swarms of Deterministic Robots

Franck Petit

LiP6, UPMC Paris 6

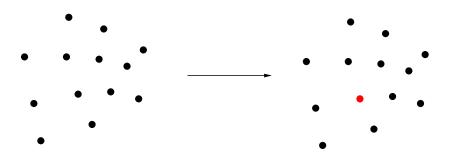
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Question

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SoD and Chirality

[Flocchini et al., 1999]

A solution exists for any n.

SoD and No Chirality

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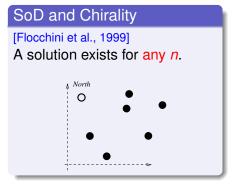
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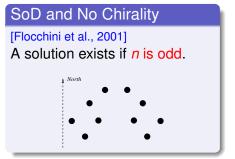
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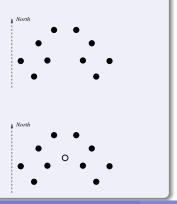




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No SoD

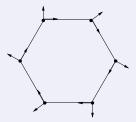
[Prencipe 2002]

Impossible in general.

No SoD

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Impossible in general.



In such a configuration, it is not possible to break the symmetry.

Leader Election With No Sense of Direction

Question

Assuming no sense of direction (with or without chirality), what are the geometric conditions to be able to deterministically agree on a single robot?

To answer to this question, we need tools from the theory on Combinatoric on Words, specifically Lyndon Words.

Leader Election With No Sense of Direction

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Assuming no sense of direction (with or without chirality), what are the geometric conditions to be able to deterministically agree on a single robot?

To answer to this question, we need tools from the theory on Combinatoric on Words, specifically Lyndon Words.

Definition (Word)

Let $A = \{a_0, a_1, \dots, a_n\}$ be an alphabet. A word is a (possibly empty) sequence of letters in A.

$$A = \{a, b, c, d\}$$

$$abcc \quad a \quad \epsilon \quad dddddddd \equiv d^{8}$$

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Definition (Concatenation)

Let $u=a_1,\ldots,a_i,\ldots,a_k$ and $v=b_1,\ldots,b_j,\ldots,b_\ell$. The concatenation of u and v, denoted uv, is equal to the word $a_1,\ldots,a_i,\ldots,a_k,b_1,\ldots,b_\ell$.

$$u = UP, v = MC, uv = UPMC$$

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$$u = UP$$
, $v = MC$, $uv = UPMC$

Definition (Lexicographic Order)

Let *A* be an alphabet totally ordered by \prec , *i.e.*, $a_0 \prec a_1 \prec \ldots \prec a_n$.

A word $u = a_0 a_1 \dots a_s$ is said to be *lexicographically smaller* than or equal to a word $v = b_0 b_1 \dots b_t$, denoted by $u \leq v$, iff:

- either u is a prefix of v,
- or, $\exists k : \forall i \in [1, ..., k-1], a_i = b_i \text{ and } a_k \prec b_k.$

 $ab \prec abc$ $abc \prec abc$ $\epsilon \prec abc$ $abc \prec de$

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abc ≺ def

Definition (Primitive Word)

A word u is said to be *primitive* iff $u = v^k \Rightarrow k = 1$. Otherwise, u is said to be *periodic*.

Primitive Words

ab dabcı

dcba

Periodic Words

 d^8

bcbc



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A word u is said to be a *rotation* of a word v iff there exists two words x, y such that u = xy and v = yx.

$$u = abcd$$
 and $v = cdab$
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A word *u* is a *Lyndon word* iff *u* is primitive and minimal.

Lyndon Word

 $abc (abc \leq cab \text{ and } abc \leq bca)$

Not a Lyndon Word

 $bca (bca \succ abc)$

Definition (Lyndon Word)

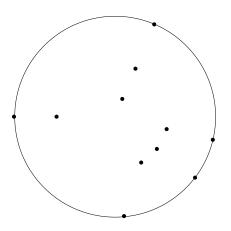
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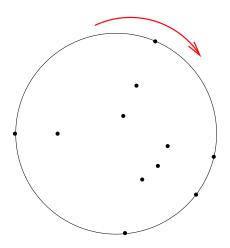
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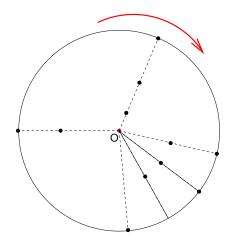
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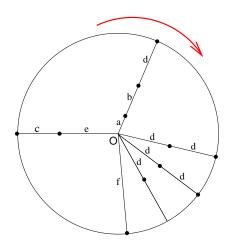
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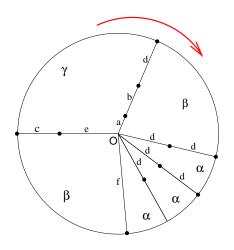
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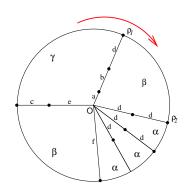












$$W(\rho_1) = (abc, \beta)(d^2, \alpha)^2(d, \alpha)(f, \beta)(ec, \gamma)$$

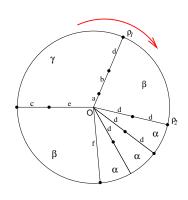
$$W(\rho_2) = (d^2, \alpha)^2(d, \alpha)(f, \beta)(ec, \gamma)(abc, \beta)$$

Lemma

If two distinct radii ρ_1 and ρ_2 exist such that both $W(\rho_1)$ and $W(\rho_2)$ are Lyndon words, then $\forall \rho$, $W(\rho) = (0,0)$.

Lemma (⇒

If there exists a radius ρ such that $W(\rho)$ is a Lyndon word, then the robots are able to deterministically agree on the same leader.



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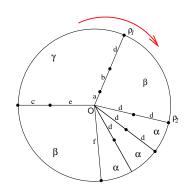
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If there exists a radius ρ such that $W(\rho)$ is a Lyndon word, then the robots are able to deterministically agree on the same leader.

Lemma (⇐)

If there exists no radius ρ such that $W(\rho)$ is a Lyndon word, then the robots are not able to deterministically agree on the same leader.

Property

[Lothaire 1983]

If no rotation of a work *u* is a Lyndon word, then *u* is periodic.

Lemma (⇐)

If there exists no radius ρ such that $W(\rho)$ is a Lyndon word, then the robots are not able to deterministically agree on the same leader.

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If no rotation of a word u is a Lyndon word, then u is periodic.

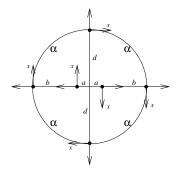
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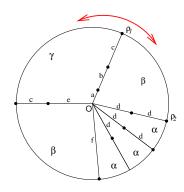
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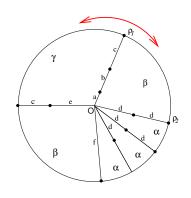
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Theorem

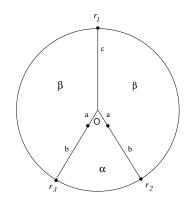
Assuming chirality, a swarm of robots is able to deterministically agree on the same leader if and only if there exists a radius ρ such that $W(\rho)$ is a Lyndon word.

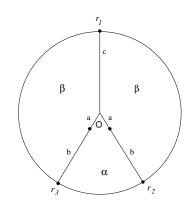




For each ρ , there are 2 ways to compute $W(\rho)$

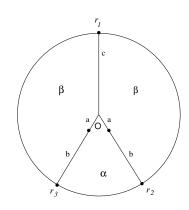
$$W(\rho_1) =$$
 either $(abc, \beta)(d^2, \alpha)^2(d, \alpha)(f, \beta)(ec, \gamma)$ or $(abc, \gamma)(ec, \beta)(f, \alpha)(d, \alpha)(d^2, \alpha)(d^2, \beta)$ depending on either \circ or \circ , respectively.





The word

$$W(\rho_2)^{\circlearrowright} = W(\rho_3)^{\circlearrowleft} = (ab, \alpha)(ab, \beta)(c, \beta)$$
 is a Lyndon word.

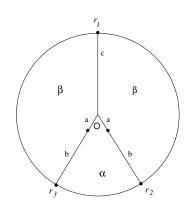


Definition (Type of Symmetry)

A radius ρ_i is of Type (of symmetry) $\mathbf{0}$ if there exists no radius ρ_j such that $W(\rho_i)^{\circlearrowleft} = W(\rho_j)^{\circlearrowright}$. Otherwise, ρ_i is said to be of Type $\mathbf{1}$.

A radius of Type *t* is said to be *t*-symmetric.

 ρ_1 is 0-symmetric. ρ_2 and ρ_3 are 1-symmetric.

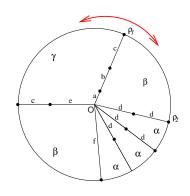


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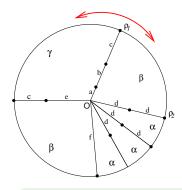
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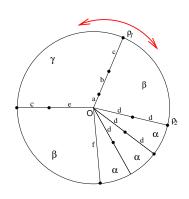
For each radius ρ_i , every robot computes $W(\rho_i)^{\circlearrowleft}$ and $W(\rho_i)^{\circlearrowleft}$ of the form (*type*, *radiusword*, *angle*).



For each radius ρ_i , every robot computes $W(\rho_i)^{\circlearrowright}$ and $W(\rho_i)^{\circlearrowleft}$ of the form (*type*, *radiusword*, *angle*).

$$W(\rho_1)^{\circlearrowleft} = (0, abc, \beta)(0, d^2, \alpha)^2(0, d, \alpha)(0, f, \beta)(0, ec, \gamma)$$

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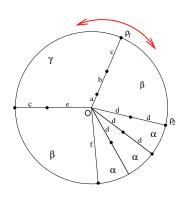


Lemma

If two distinct radii ρ_1 and ρ_2 exist such that both $W(\rho_1)$ and $W(\rho_2)$ are Lyndon words, then $\forall \rho$, $W(\rho) = (0,0,0)$.

Lemma

If there exists a pair of radii $\{\rho_1, \rho_2\}$ so that $W(\rho_i)^{\circlearrowright}$ or $W(\rho_i)^{\circlearrowleft}$ is a Lyndon word ($i \in \{1,2\}$), then the robots are able to deterministically agree on the same leader if and only if ρ_1 and ρ_2 are 0-symmetric.

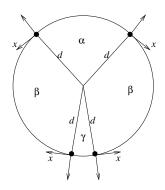


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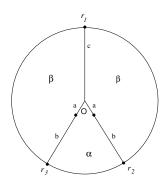
No leader exists.

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The robot on ρ_1 is the leader.

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Theorem

Assuming no chirality, a swarm of robots is able to deterministically agree on the same leader if and only if there exists a radius ρ such that $W(\rho)$ is a 0-symmetric Lyndon word.