

CS 4300: Assignment #3 - Constraint Satisfaction Problem

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Problem Domain: 9x9 Sudoku

GitHub Repository

The source code for the modified solver and the `.csp` instance files can be found at the following repository:

https://github.com/BranPLewis/Sudoku_9x9

1. Problem Description: Sudoku

Sudoku is a logic-based number-placement puzzle. The objective is to fill a **9×9 grid** with digits so that each column, each row, and each of the nine **3×3 subgrids** that compose the grid contain all of the digits from 1 to 9. The puzzle starts with a partially completed grid, which for a well-posed puzzle has a single unique solution.

The challenge is to fill in the remaining empty cells while adhering to these three core constraints simultaneously:

1. Each row must contain the numbers 1 through 9, without repetition.
2. Each column must contain the numbers 1 through 9, without repetition.
3. Each of the nine 3×3 subgrids must contain the numbers 1 through 9, without repetition.

2. Formalization as `<X, D, C>`

A Sudoku puzzle can be formally modeled as a Constraint Satisfaction Problem (CSP) using the `<X, D, C>` framework.

- **X: Variables**
The set of variables, X , consists of 81 variables, one for each cell in the 9x9 grid. We can denote each variable as V_{rc} , where r is the row index ($1 \leq r \leq 9$) and c is the column index ($1 \leq c \leq 9$).
$$X = \{V_{rc} \mid 1 \leq r \leq 9, 1 \leq c \leq 9\}$$
- **D: Domains**
The domain, D , for each variable is the set of possible values it can take. For any empty cell in the grid, the initial domain is the set of all possible digits.
$$D_{rc} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

For cells that are pre-filled in a specific puzzle instance with a value k , the domain for that variable is restricted to just that single value, e.g., $D_{rc} = k$.

- **C: Constraints**

The constraints, C, ensure that the rules of Sudoku are followed. These are all alldifferent constraints applied to three different types of groups.

1. **Row Constraints:** For each row r , all variables in that row must be unique.
 $\text{alldiff}(V_{r1}, V_{r2}, \dots, V_{r9})$ for each $r \in \{1, \dots, 9\}$
2. **Column Constraints:** For each column c , all variables in that column must be unique.
 $\text{alldiff}(V_{1c}, V_{2c}, \dots, V_{9c})$ for each $c \in \{1, \dots, 9\}$
3. **3x3 Subgrid Constraints:** For each of the nine 3x3 subgrids, all variables within that subgrid must be unique. For example, the top-left subgrid constraint is:
 $\text{alldiff}(V_{11}, V_{12}, V_{13}, V_{21}, V_{22}, V_{23}, V_{31}, V_{32}, V_{33})$
There are a total of 9 row, 9 column, and 9 subgrid constraints, making 27 alldiff constraints in total.

3. Heuristic Implementation: Minimum Remaining Values (MRV)

To extend the backtracking solver, I implemented the **Minimum Remaining Values (MRV)** heuristic. This is a variable-ordering heuristic that improves the efficiency of the search process.

- **How it Works:** Instead of selecting the next unassigned variable in a fixed, static order, the MRV heuristic dynamically chooses the variable with the **fewest legal values** remaining in its domain.
- **Why it is Effective:** This strategy is often called the "fail-first" principle. By selecting the most constrained variable, we are more likely to identify a dead-end in the search tree early on. For example, if a variable's domain has been reduced to a single possible value, it is critical to test that assignment immediately. If it leads to a contradiction, we can prune that entire branch of the search tree without wasting time exploring other variable assignments. In Sudoku, this perfectly mimics the human approach of finding a cell where only one number can possibly go.

4. Results and Reflection

This section presents the solutions found by the solver for the `.csp` instances and analyzes the performance improvement gained by using the MRV heuristic.

Solution for `sudoku_1.csp`

Without heuristic

Steps it took: 1113

Time it took: 1.58 seconds

Solution:

```
1 6 3 | 9 4 8 | 5 2 7
7 8 2 | 6 5 1 | 3 9 4
4 9 5 | 2 7 3 | 1 8 6
```

```
-----
6 7 9 | 3 1 2 | 4 5 8
3 4 8 | 5 9 7 | 6 1 2
2 5 1 | 8 6 4 | 7 3 9
```

```
-----
8 2 6 | 7 3 5 | 9 4 1
9 3 4 | 1 2 6 | 8 7 5
5 1 7 | 4 8 9 | 2 6 3
```

With heuristic

Heuristic: MRV (Minimum Remaining Values)

Steps it took: 95

Time it took: 0.026 seconds

Solution:

```
1 6 3 | 9 4 8 | 5 2 7
7 8 2 | 6 5 1 | 3 9 4
4 9 5 | 2 7 3 | 1 8 6
```

```
-----
6 7 9 | 3 1 2 | 4 5 8
3 5 8 | 7 9 4 | 6 1 2
2 4 1 | 8 6 5 | 7 3 9
```

```
-----
8 2 6 | 5 3 7 | 9 4 1
9 3 4 | 1 2 6 | 8 7 5
5 1 7 | 4 8 9 | 2 6 3
```

Solution for `sudoku_2.csp`

Without heuristic

Steps it took: 344

Time it took: 0.16 seconds

Solution:

```
6 8 2 | 3 5 4 | 9 1 7
4 1 9 | 6 8 7 | 3 5 2
7 3 5 | 1 9 2 | 8 6 4
```

```
-----
3 9 6 | 2 7 5 | 4 8 1
1 2 8 | 4 3 6 | 5 7 9
5 7 4 | 8 1 9 | 6 2 3
```

```
-----
2 5 7 | 9 6 3 | 1 4 8
9 6 1 | 7 4 8 | 2 3 5
8 4 3 | 5 2 1 | 7 9 6
```

With heuristic

Heuristic: MRV (Minimum Remaining Values)

Steps it took: 120

Time it took: 0.028 seconds

Solution:

```
6 8 2 | 3 5 4 | 9 1 7
4 1 9 | 6 8 7 | 3 5 2
7 3 5 | 1 9 2 | 8 6 4
```

```
-----
3 9 6 | 2 7 5 | 4 8 1
1 2 4 | 8 6 3 | 5 7 9
5 7 8 | 4 1 9 | 6 2 3
```

```
-----
2 5 7 | 9 3 6 | 1 4 8
9 6 1 | 7 4 8 | 2 3 5
8 4 3 | 5 2 1 | 7 9 6
```

Solution for `sudoku_3.csp`

Without heuristic

Steps it took: 6156

Time it took: 1.97 seconds

Solution:

```
5 3 4 | 6 7 8 | 9 1 2
6 7 2 | 1 9 5 | 3 4 8
1 9 8 | 3 4 2 | 5 6 7
```

```
-----
8 5 9 | 7 6 1 | 4 2 3
4 2 6 | 8 5 3 | 7 9 1
7 1 3 | 9 2 4 | 8 5 6
```

```
-----
9 6 1 | 5 3 7 | 2 8 4
2 8 7 | 4 1 9 | 6 3 5
3 4 5 | 2 8 6 | 1 7 9
```

With heuristic

Heuristic: MRV (Minimum Remaining Values)

Steps it took: 81

Time it took: 0.019 seconds

Solution:

```
5 3 4 | 6 7 8 | 9 1 2
6 7 2 | 1 9 5 | 3 4 8
1 9 8 | 3 4 2 | 5 6 7
```

```
-----
8 5 9 | 7 6 1 | 4 2 3
4 2 6 | 8 5 3 | 7 9 1
7 1 3 | 9 2 4 | 8 5 6
```

```
-----
9 6 1 | 5 3 7 | 2 8 4
2 8 7 | 4 1 9 | 6 3 5
3 4 5 | 2 8 6 | 1 7 9
```

Analysis

What I found is that Minimum Remaining Values in the terms of a 9x9 sudoku board was crucial in finding rows and columns and 3x3 sub-grids that had few numbers left to fill, forcing the solver to either fill them successfully or find a bad path in need of backtracking. I noticed a significant reduction in the amount of steps with the solver using the MRV heuristic. I created 3 instances of a 9x9 sudoku board, each of which has a different increasing number of starting tiles, and a different number of solutions occurred for each. Each instance MRV found the solution in roughly $\frac{1}{3}$ less steps than without MRV.