

CSC 421 Assignment3

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Question1:

CSC 421
 Assignment 3.
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T1:

clausal form:

$$(p \vee \neg r) \& (\neg r \vee \neg q) \& (\neg q \vee \neg p) \rightarrow ((\neg r \vee q) \& (\neg q \vee p))$$

I: $(p \vee \neg r) \& (\neg r \vee \neg q) \& (\neg q \vee \neg p)$

IV: $(p \vee \neg r) \& [(\neg r \& \neg q) \& (\neg q \& \neg p)] \& [(\neg r \& q) \& (\neg q \& p)]$

D: $(p \vee \neg r) \& [(\neg r \& \neg q) \& (\neg q \vee (\neg p \& \neg r)) \& (\neg r \& \neg q \& \neg p \& \neg r)] \& [(\neg r \& q) \& (\neg q \vee (\neg p \& \neg r)) \& (\neg r \& q \& \neg p \& \neg r)]$

$$= (p \vee \neg r) \& [(\neg r \& \neg q) \& (\neg q \vee (\neg p \& \neg r)) \& (\neg p \vee (\neg r \& \neg q)) \& (\neg r \& \neg q \& \neg p \& \neg r)]$$

$$= (p \vee \neg r) \& [(\neg r \& \neg q) \& (\neg q \vee (\neg p \& \neg r)) \& (\neg p \vee (\neg r \& \neg q)) \& (\neg r \& \neg q \& \neg p \& \neg r)]$$

$$= (p \vee \neg r) \& [(\neg r \& \neg q) \& (\neg q \vee (\neg p \& \neg r)) \& (\neg p \vee (\neg r \& \neg q)) \& (\neg r \& \neg q \& \neg p \& \neg r)]$$

O: $\{p, q, \neg r\}, \{\neg r, \neg q\}, \{\neg q, r, \neg q\}, \{\neg p, r, \neg q\}, \{\neg r, \neg q\}, \{\neg q\}$
 $\{\neg p, \neg q\}, \{\neg r, \neg p\}, \{\neg q, \neg p\}, \{\neg p\}$

resolution step.

1. $\{p, q, \neg r\}$
2. $\{\neg r, \neg q\}$
3. $\{\neg q, r, \neg q\}$
4. $\{\neg p, r, \neg q\}$
5. $\{\neg r, \neg q\}$
6. $\{\neg q\}$
7. $\{\neg p, \neg q\}$
8. $\{\neg r, \neg p\}$
9. $\{\neg q, \neg p\}$
10. $\{\neg p\}$

resolution:

11. $\{p, q\} \vdash 1, 2$
12. $\{\neg q\} \vdash 11, 10$
13. $\{\} \vdash 12, 6$

Question2:

and Johnson
from promise
1500000
E. fasten(y,z)

Q2:

$$\forall x, \forall y (\text{horse}(x) \wedge \text{dog}(y) \Rightarrow \text{faster}(x, y))$$

$$\exists y (\text{Greyhound}(y) \wedge \forall z (\text{Rabbit}(z) \Rightarrow \text{Faster}(y, z)))$$

$$\forall y (\text{Greyhound}(y) \Rightarrow \text{Dog}(y))$$

$$\forall x \forall y \forall z (\text{Faster}(x, y) \wedge \text{Faster}(y, z) \Rightarrow \text{Faster}(x, z))$$

$$\neg \forall x \forall y (\text{Horse}(x) \wedge \text{Rabbit}(y) \Rightarrow \text{Faster}(x, y)) \text{ negated conclusion}$$

do causal form, and do the resolution steps for each statement

(1) $\forall x, \forall y (\text{horse}(x) \wedge \text{dog}(y) \Rightarrow \text{faster}(x, y))$

I: $\forall x, \forall y (\neg \text{horse}(x) \wedge \neg \text{dog}(y)) \vee \text{faster}(x, y)$

N: $\forall x, \forall y ((\neg \text{horse}(x) \vee \neg \text{dog}(y)) \vee \text{faster}(x, y))$

S: $\forall x, \forall y ((\neg \text{horse}(x) \vee \neg \text{dog}(y)) \vee \text{faster}(x, y))$

E: $\forall x, \forall y (((\neg \text{horse}(x) \vee \neg \text{dog}(y)) \vee \text{faster}(x, y))$

A: $((\neg \text{horse}(x) \vee \neg \text{dog}(y)) \vee \text{faster}(x, y))$

D: $\neg \text{horse}(x) \vee \neg \text{dog}(y) \vee \text{faster}(x, y)$

O: $\{ \neg \text{horse}(x), \neg \text{dog}(y), \text{faster}(x, y) \}$

(2) $\exists y (\text{Greyhound}(y) \wedge (\forall z (\text{Rabbit}(z) \Rightarrow \text{faster}(y, z)))$

I: $\exists y ((\text{Greyhound}(y)) \wedge \forall z (\neg \text{Rabbit}(z)) \vee \text{faster}(y, z))$

N: $\exists y ((\text{Greyhound}(y)) \wedge \forall z ((\neg \text{Rabbit}(z)) \vee \text{faster}(y, z)))$

S: $\text{Greyhound}(a) \wedge \forall z ((\neg \text{Rabbit}(z)) \vee \text{faster}(a, z))$

E: $\text{Greyhound}(a) \wedge \forall z ((\neg \text{Rabbit}(z)) \vee \text{faster}(a, z))$

A: $\text{Greyhound}(a) \wedge (\neg \text{Rabbit}(z) \vee \text{faster}(a, z))$

O: $\{ \text{Greyhound}(a) \}, \{ \neg \text{Rabbit}(z), \text{faster}(a, z) \}$

(3) $\forall y (\text{Greyhound}(y) \Rightarrow \text{Dog}(y))$

I: $\forall y ((\text{Greyhound}(y)) \vee \text{Dog}(y))$

A: $\neg \text{Greyhound}(y) \vee \text{Dog}(y)$

O: $\{ \neg \text{Greyhound}(y), \text{Dog}(y) \}$

$$\textcircled{4} \quad \forall x \forall y \forall z (\text{Faster}(x, y) \wedge \text{Faster}(y, z) \Rightarrow \text{Faster}(x, z))$$

$$\text{I: } \forall x \forall y \forall z (\neg(\text{Faster}(x, y) \wedge \text{Faster}(y, z)) \vee \text{Faster}(x, z))$$

$$\text{IV: } \forall x \forall y \forall z ((\neg \text{Faster}(x, y) \vee \neg \text{Faster}(y, z)) \vee \text{Faster}(x, z))$$

$$\text{A: } (\neg \text{Faster}(x, y) \vee \neg \text{Faster}(y, z)) \vee \text{Faster}(x, z)$$

$$\text{D: } \neg \text{Faster}(x, y) \vee \neg \text{Faster}(y, z) \vee \text{Faster}(x, z)$$

$$\text{O: } \left. \begin{array}{l} \neg \text{Faster}(x, y), \neg \text{Faster}(y, z), \text{Faster}(x, z) \end{array} \right\}$$

$$\textcircled{5} \quad \neg \forall x \forall y (\text{Horse}(x) \wedge \text{Rabbit}(y) \Rightarrow \text{Faster}(x, y))$$

$$\text{I: } \neg \forall x \forall y (\neg(\text{Horse}(x) \wedge \text{Rabbit}(y)) \vee \text{Faster}(x, y))$$

$$\text{N: } \exists x \exists y (\neg(\text{Horse}(x) \wedge \text{Rabbit}(y)) \wedge \neg \text{Faster}(x, y))$$

$$\text{E: } (\text{Horse}(c/b) \wedge \text{Rabbit}(c/c)) \wedge \neg \text{Faster}(c/b, c/c)$$

$$\text{D: } (\text{Horse}(c/b) \wedge \text{Rabbit}(c/c) \wedge \neg \text{Faster}(c/b, c/c))$$

$$\text{O: } \left. \begin{array}{l} \text{Horse}(c/b), \text{Rabbit}(c/c), \neg \text{Faster}(c/b, c/c) \end{array} \right\}$$

resolution steps:

$$\textcircled{1} \quad \{ \text{horse}(x), \neg \text{dog}(y), \text{faster}(x, y) \}$$

$$\textcircled{2} \quad \{ \text{Greyhound}(a) \}$$

$$\textcircled{3} \quad \{ \neg \text{Rabbit}(z), \text{faster}(a, z) \}$$

$$\textcircled{4} \quad \{ \neg \text{Greyhound}(y), \text{Dog}(y) \}$$

$$\textcircled{5} \quad \{ \neg \text{Faster}(x, y), \neg \text{faster}(y, z), \text{faster}(x, z) \}$$

$$\textcircled{6} \quad \{ \text{Horse}(b) \}$$

$$\textcircled{7} \quad \{ \text{rabbit}(c) \}$$

$$\textcircled{8} \quad \{ \neg \text{Faster}(b, c) \}$$

$$\textcircled{9} \quad \{ \neg \text{dog}(y), \text{faster}(x, y) \} \quad \textcircled{1} \textcircled{6} \quad c \rightarrow z$$

$$\textcircled{10} \quad \{ \neg \text{greyhound}(y), \text{faster}(x, y) \} \quad \textcircled{4} \textcircled{9} \quad b \rightarrow x$$

$$\textcircled{11} \quad \{ \text{faster}(a, z) \} \quad \textcircled{5} \textcircled{2} \quad a \rightarrow y$$

$$\textcircled{12} \quad \{ \text{faster}(x, y) \} \quad \textcircled{2} \textcircled{10} \quad a \rightarrow y$$

$$\textcircled{13} \quad \{ \neg \text{faster}(y, z), \text{faster}(x, z) \} \quad \textcircled{5} \textcircled{12} \quad b \rightarrow x$$

$$\textcircled{14} \quad \{ \text{faster}(x, z) \} \quad \textcircled{13} \textcircled{11} \quad c \rightarrow z$$

$$\textcircled{15} \quad \{ \} \quad \textcircled{14} \textcircled{8}$$

Question 3:

Assumption:

all x (hummingbird(x) -> bird(x)).
 all x (hummingbird(x) -> richly_colored(x)).
 all x (bird(x) & large(x) -> -live_on_honey(x)).
 all x (bird(x) & -live_on_honey(x) -> -richly_colored(x)).

goal:

all x (hummingbird(x) -> -large(x)).

```
=====
 prooftrans =====
Prover9 (32) version Dec-2007, Dec 2007.
Process 79097 was started by brandonm on Brandons-MacBook-Pro.local,
Mon Nov 18 20:46:03 2019
The command was "/private/var/folders/_q/yybjd8v575d1rrgwsj7r51dr0000gn/T/
AppTranslocation/9DB74EC8-292F-4B76-8F30-33BA37F1B7DD/d/Prover9-Mace4-v05B.app/
Contents/Resources/bin-mac-intel/prover9".
=====
 end of head =====

=====
 end of input =====

=====
 PROOF =====

% ----- Comments from original proof -----
% Proof 1 at 0.00 (+ 0.00) seconds.
% Length of proof is 18.
% Level of proof is 5.
% Maximum clause weight is 0.
% Given clauses 0.

1 (all x (hummingbird(x) -> bird(x))) # label(non_clause). [assumption].
2 (all x (hummingbird(x) -> richly_colored(x))) # label(non_clause). [assumption].
3 (all x (bird(x) & large(x) -> -live_on_honey(x))) # label(non_clause). [assumption].
4 (all x (bird(x) & -live_on_honey(x) -> -richly_colored(x))) # label(non_clause). [assumption].
5 (all x (hummingbird(x) -> -large(x))) # label(non_clause) # label(goal). [goal].
6 hummingbird(c1). [deny(5)].
7 -hummingbird(x) | bird(x). [clausify(1)].
8 -hummingbird(x) | richly_colored(x). [clausify(2)].
9 bird(c1). [resolve(6,a,7,a)].
10 -bird(x) | -large(x) | -live_on_honey(x). [clausify(3)].
11 -bird(x) | live_on_honey(x) | -richly_colored(x). [clausify(4)].
12 -large(c1) | -live_on_honey(c1). [resolve(9,a,10,a)].
13 large(c1). [deny(5)].
14 live_on_honey(c1) | -richly_colored(c1). [resolve(9,a,11,a)].
15 richly_colored(c1). [resolve(6,a,8,a)].
16 live_on_honey(c1). [resolve(14,b,15,a)].
17 -live_on_honey(c1). [resolve(12,a,13,a)].
18 $F. [resolve(16,a,17,a)].
```

```
=====
 end of proof =====
```

Question 4:

Assumption:

```
worth_listening_subjects(gardener).
all x (remember_battle_Waterloo(x) -> old(x)).
all x (worth_listening_subjects(x) -> remember_battle_Waterloo(x)).
```

Goal:

```
old(gardener).
```

===== PROOF =====

% ----- Comments from original proof -----

% Proof 1 at 0.00 (+ 0.00) seconds.

% Length of proof is 10.

% Level of proof is 4.

% Maximum clause weight is 0.

% Given clauses 0.

```
1 (all x (remember_battle_Waterloo(x) -> old(x))) # label(non_clause). [assumption].
2 (all x (worth_listening_subjects(x) -> remember_battle_Waterloo(x))) # label(non_clause).
[assumption].
3 old(gardener) # label(non_clause) # label(goal). [goal].
4 -worth_listening_subjects(x) | remember_battle_Waterloo(x). [clausify(2)].
5 worth_listening_subjects(gardener). [assumption].
6 remember_battle_Waterloo(gardener). [resolve(4,a,5,a)].
7 -remember_battle_Waterloo(x) | old(x). [clausify(1)].
8 old(gardener). [resolve(6,a,7,a)].
9 -old(gardener). [deny(3)].
10 $F. [resolve(8,a,9,a)].
```

===== end of proof =====

Question 5:

Q5: $P_{3,1}$ and $P_{1,3}$ has same probability since it is symmetric.

$$\begin{aligned} P(CP_{13}|b_{12}, b_{21}) &= \alpha [P(b_{12}|P_{13}, P_{22}) \cdot P(b_{21}|P_{31}, P_{13}) + P(P_{13}) \cdot P(P_{22}) \cdot P(P_{31})] \\ &= \alpha [P(b_{12}|P_{13}, P_{22}) \cdot P(b_{21}|P_{31}, P_{13}) \cdot P(P_{13}) \cdot P(P_{22}) \cdot P(P_{31}) + \\ &\quad P(b_{12}|P_{13}, P_{22}) \cdot P(b_{21}|P_{31}, P_{13}) \cdot P(P_{13}) \cdot P(P_{22}) \cdot P(P_{31}) + \\ &\quad P(b_{12}|P_{13}, P_{22}) \cdot P(b_{21}|P_{31}, P_{13}) \cdot P(P_{13}) \cdot P(P_{22}) \cdot P(P_{31})] \\ &= \alpha (1 \times 1 \times 0.01 \times 0.01 \times 0.01 + 1 \times 1 \times 0.99 \times 0.01 + 1 \times 1 \times 0.01 \times 0.01 \times 0.99 + 0) \\ &= 0.000199 \alpha \end{aligned}$$

$$\begin{aligned} P(\neg P_{13}|b_{12}, b_{21}) &= \alpha (1 \times 1 \times 0.99 \times 0.01 \times 0.01 + 0 + 0.99 \times 0.99 \times 0.01 \times 1 \times 1 + 0) \\ &= 0.0099 \alpha \end{aligned}$$

$$\alpha = \frac{1}{0.000199 + 0.0099}$$

$$\begin{aligned} P(P_{13}|b_{12}, b_{21}) &= \frac{0.000199}{\alpha} = 0.019705 = 1.97\% \\ P(\neg P_{13}|b_{12}, b_{21}) &= \frac{0.0099}{\alpha} = 0.10295 = 98.03\% \end{aligned}$$

$$\begin{aligned} P(CP_{22}|b_{12}, b_{21}) &= \alpha (1 \times 1 \times 0.01 \times 0.01 \times 0.01 + 1 \times 1 \times 0.01 \times 0.99 \times 0.01 + 1 \times 1 \times 0.01 \times 0.99 + 1 \times 1 \times 0.01 \times 0.99 \times 0.99) \\ &= 0.01 \alpha \end{aligned}$$

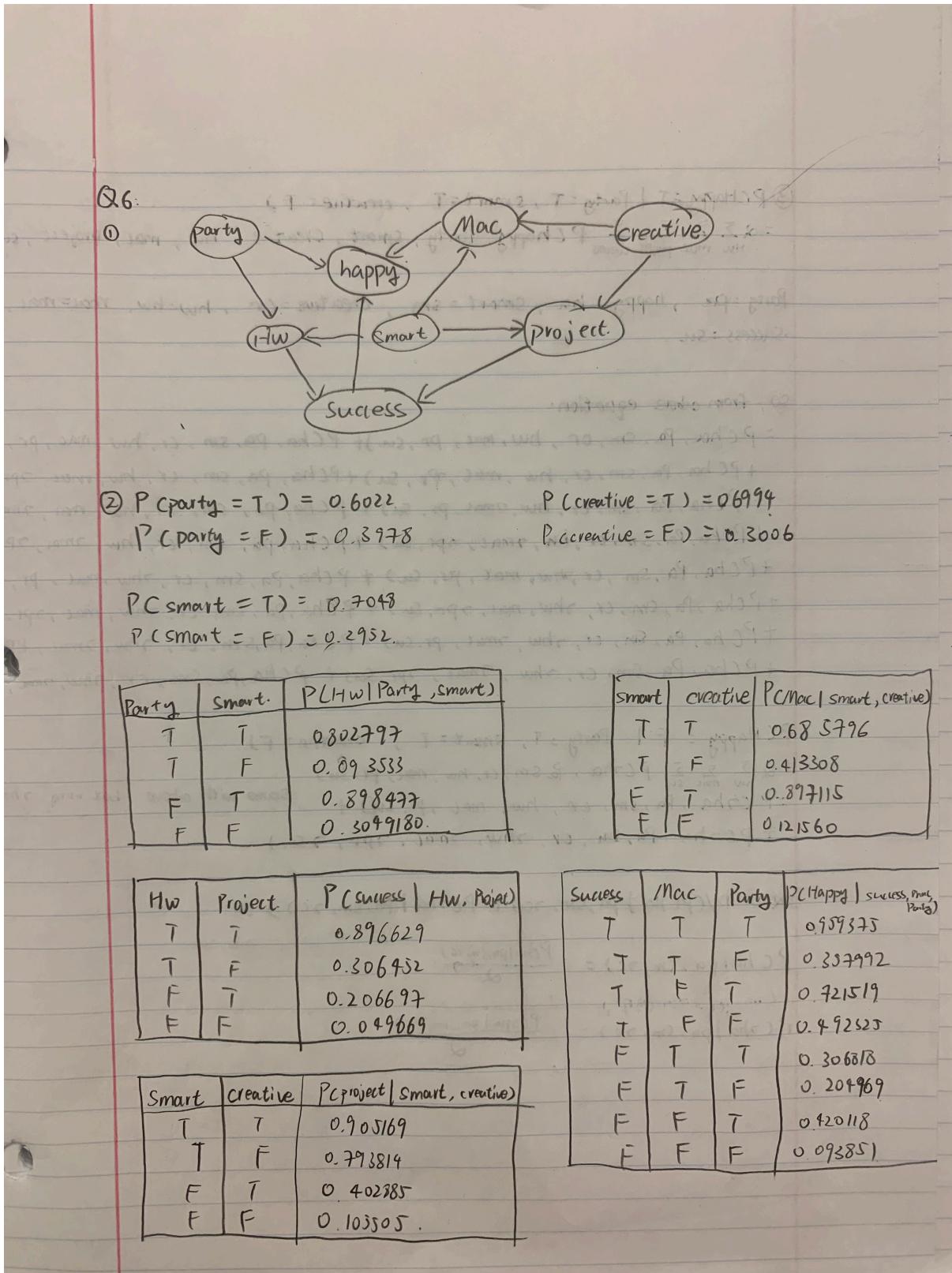
$$\begin{aligned} P(\neg P_{22}|b_{12}, b_{21}) &= \alpha (1 \times 1 \times 0.01 \times 0.99 \times 0.01 + 0 + 0 + 0) \\ &= 0.000099 \alpha \end{aligned}$$

$$\alpha = \frac{1}{0.000099 + 0.01} = \frac{1}{0.010099}$$

$$\begin{aligned} P(CP_{22}|b_{12}, b_{21}) &= \frac{0.01}{\alpha} \approx 0.19 \approx 19.9\% \\ P(\neg P_{22}|b_{12}, b_{21}) &= \frac{0.000099}{\alpha} \approx 0.01\% \end{aligned}$$

Thus, going to $[2, 2]$ is almost certain death $[3, 1]$ and $[1, 3]$ is the probability step that the logical agent goes. But three box are looks have equal opportunity. So $1/3$ chance to die it is same with probability = 0.2,

Question 6:



$$\textcircled{3} \quad P(\text{Happy} = T \mid \text{Party} = T, \text{smart} = T, \text{creative} = F)$$

$$= \alpha \sum_{hw} \sum_{mac} \sum_{proj} \sum_{succ} P(\text{happy}, \text{Party}, \text{smart}, \text{Creative}, hw, mac, \text{proj}, \text{succ})$$

Party = Pa, happy = ha, smart = sm, creative = cr, hw = hw, mac = mac, project = pr, success = su.

So, from above equation:

$$\begin{aligned}
 &= P(\text{ha}, Pa, sm, cr, hw, mac, pr, su) + P(\text{ha}, Pa, sm, cr, hw, mac, pr, \neg su) \\
 &\quad + P(\text{ha}, Pa, sm, cr, hw, mac, \neg pr, su) + P(\text{ha}, Pa, sm, cr, hw, mac, \neg pr, \neg su) \\
 &\quad + P(\text{ha}, Pa, sm, cr, hw, \neg mac, pr, su) + P(\text{ha}, Pa, sm, cr, hw, \neg mac, pr, \neg su) \\
 &\quad + P(\text{ha}, Pa, sm, cr, \neg hw, mac, pr, su) + P(\text{ha}, Pa, sm, cr, \neg hw, mac, pr, \neg su) \\
 &\quad + P(\text{ha}, Pa, sm, cr, \neg hw, mac, \neg pr, su) + P(\text{ha}, Pa, sm, cr, \neg hw, mac, \neg pr, \neg su) \\
 &\quad + P(\text{ha}, Pa, sm, \neg cr, hw, mac, pr, su) + P(\text{ha}, Pa, sm, \neg cr, hw, mac, pr, \neg su) \\
 &\quad + P(\text{ha}, Pa, sm, \neg cr, hw, mac, \neg pr, su) + P(\text{ha}, Pa, sm, \neg cr, hw, mac, \neg pr, \neg su)
 \end{aligned}$$

$$P(\text{Happy} = F \mid \text{Party} = T, \text{smart} = T, \text{creative} = F)$$

$$= \alpha \sum_{hw} \sum_{mac} \sum_{su} P(\neg \text{ha}, Pa, sm, cr, hw, mac, pr, su)$$

= $P(\neg \text{ha}, Pa, sm, cr, hw, mac, pr, su) + \dots$ same with above but using $\neg \text{happy}$.

$$+ P(\neg \text{ha}, Pa, sm, cr, \neg hw, mac, pr, su)$$

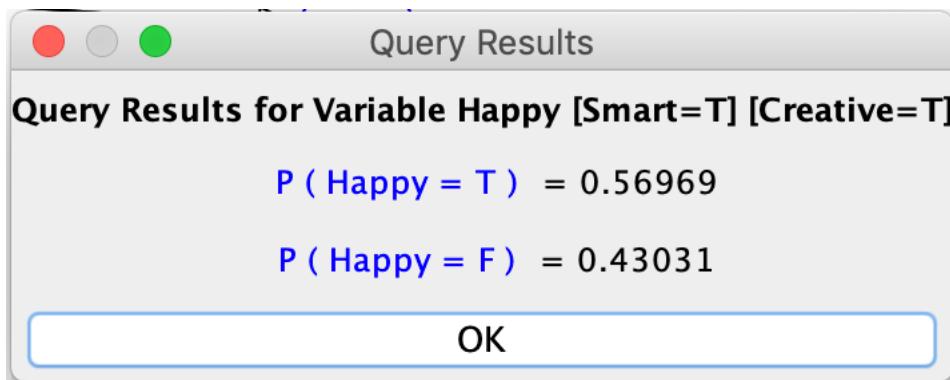
$$\alpha = 1 / (P(\text{ha}) pa, sm, \neg cr) + P(\neg \text{ha}) pa, sm, \neg cr)$$

$$P(\text{ha}) pa, sm, \neg cr = \frac{P(\text{sm} \mid \text{hw}, \text{pa}) P(\text{cr} \mid \text{pa}) P(\text{ha} \mid \text{pa}, sm, \neg cr)}{2}$$

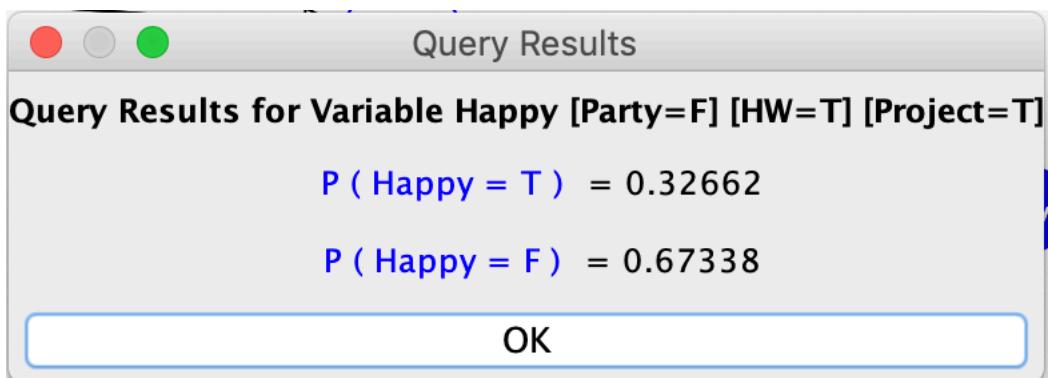
$$P(\neg \text{ha}) pa, sm, \neg cr = \frac{P(\neg \text{sm} \mid \text{hw}, \text{pa}) P(\neg \text{cr} \mid \text{pa}) P(\neg \text{ha} \mid \text{pa}, sm, \neg cr)}{2}$$

$$P(\neg \text{ha}) pa, sm, \neg cr = \frac{P(\neg \text{sm} \mid \text{hw}, \text{pa}) P(\neg \text{cr} \mid \text{pa}) P(\neg \text{ha} \mid \text{pa}, sm, \neg cr)}{2}$$

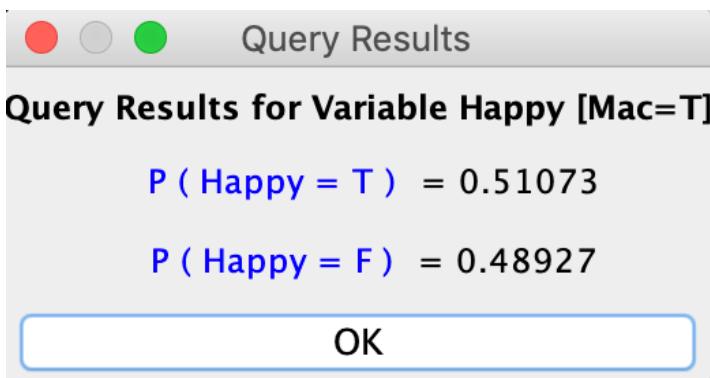
4.



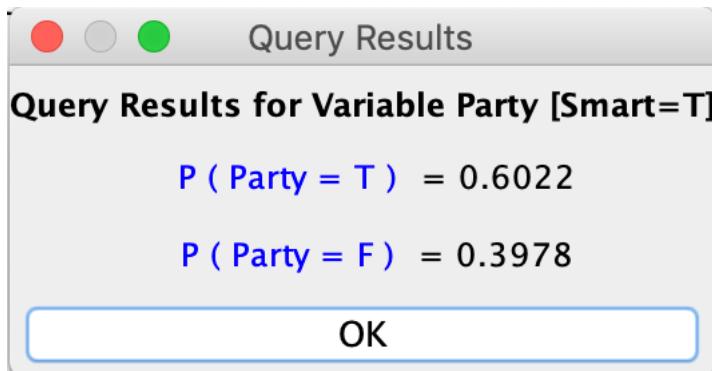
5.



6.



7.



8.

