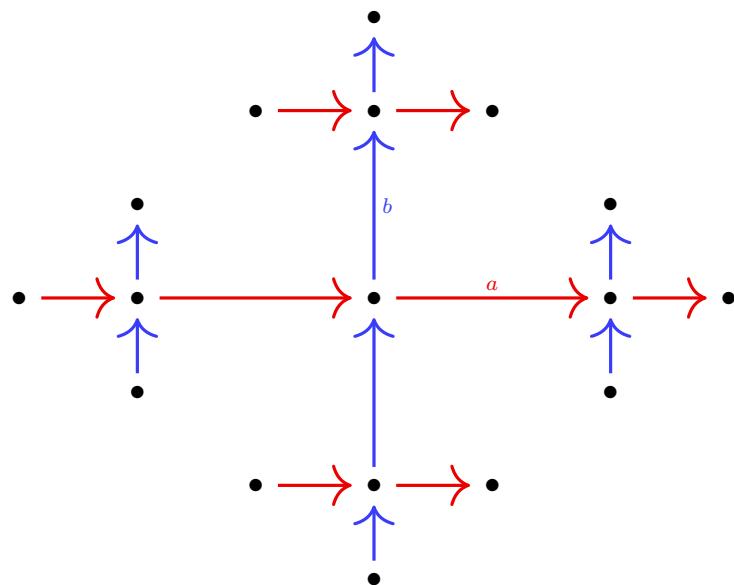


INSERT TITLE HERE

by

Brandon Papandrea



A thesis submitted in conformity with the requirements  
for the completion of MAT492H5S

Department of Mathematical and Computational Sciences  
University of Toronto - Mississauga

Insert Title Here

Brandon Papandrea

Department of Mathematical and Computational Sciences  
University of Toronto - Mississauga  
2026

## Abstract

Lorem ipsum dolor sit amet, consectetur adipiscing elit. Ut purus elit, vestibulum ut, placerat ac, adipiscing vitae, felis. Curabitur dictum gravida mauris. Nam arcu libero, nonummy eget, consectetur id, vulputate a, magna. Donec vehicula augue eu neque. Pellentesque habitant morbi tristique senectus et netus et malesuada fames ac turpis egestas. Mauris ut leo. Cras viverra metus rhoncus sem. Nulla et lectus vestibulum urna fringilla ultrices. Phasellus eu tellus sit amet tortor gravida placerat. Integer sapien est, iaculis in, pretium quis, viverra ac, nunc. Praesent eget sem vel leo ultrices bibendum. Aenean faucibus. Morbi dolor nulla, malesuada eu, pulvinar at, mollis ac, nulla. Curabitur auctor semper nulla. Donec varius orci eget risus. Duis nibh mi, congue eu, accumsan eleifend, sagittis quis, diam. Duis eget orci sit amet orci dignissim rutrum.

Nam dui ligula, fringilla a, euismod sodales, sollicitudin vel, wisi. Morbi auctor lorem non justo. Nam lacus libero, pretium at, lobortis vitae, ultricies et, tellus. Donec aliquet, tortor sed accumsan bibendum, erat ligula aliquet magna, vitae ornare odio metus a mi. Morbi ac orci et nisl hendrerit mollis. Suspendisse ut massa. Cras nec ante. Pellentesque a nulla. Cum sociis natoque penatibus et magnis dis parturient montes, nascetur ridiculus mus. Aliquam tincidunt urna. Nulla ullamcorper vestibulum turpis. Pellentesque cursus luctus mauris.

To people.

## Acknowledgements

I'll write this later

# Contents

<b>1</b>	<b>Introduction</b>	<b>1</b>
<b>2</b>	<b>Free Groups</b>	<b>3</b>
2.1	The Universal Property . . . . .	3
2.2	The Rank of Subgroups of $F_2$ . . . . .	5
<b>3</b>	<b>A Group Acting Freely on a Tree is Free</b>	<b>8</b>
3.1	G-Tilings . . . . .	8
3.2	Proof of Theorem 3.1 . . . . .	11

# Chapter 1

## Introduction

For any set  $S$  and group  $G$ , there is a unique group up to isomorphism satisfying the universal property, called the free group generated by  $S$ . If  $|S| = n$ , then we call this group the free group of rank  $n$ , and denote it as  $F_n$ .

Free groups were first studied as an example of Fuchsian groups, which are discrete subgroups of  $PSL_2(\mathbb{Z})$ , but they did not become relevant on their own until the late 1800s. Recall that we may write any group using its group presentation. von Dyck, who was first to express groups using a presentation, showed that free groups have the simplest possible one: The free group  $F_n$  has presentation

$$\langle a_1, a_2, \dots, a_n \rangle$$

meaning it consists of  $n$  generators with no relations between them; this is often used as an alternate definition of  $F_n$ . It is thus the case that every group is the quotient of a free group. The algebraic properties of free groups were investigated by Nielsen, who also gave them their name, whilst a comprehensive overview of free groups was given by Reidemeister in his 1932 book on combinatorial group theory.

The simple nature of free groups is equal parts useful and deceiving. Consider The Word Problem.

**The Word Problem:** For a group  $G$ , let  $g$  be an element expressed as a product of  $G$ 's generators. Find an algorithm that shows, in finitely many steps, whether or not  $g$  is the identity element.

While it is known that this problem is undecidable for arbitrary groups, it is answerable if  $G$  is a free group. Recall that a free group has no relations between the generators, apart from the obvious one that  $a_i a_i^{-1} = 1$ . Thus, given an arbitrary element of a free group, one can just check if any of these pairs exist, and eliminate them one at a time until none remain; if generators are still present, the element is not the identity, and if no generators are left, the element is the identity. Thus, while The Word Problem is undecidable, it is easily solvable in the case of free groups.

This does not mean that free groups are a trivial class of groups. In fact, understanding their algebraic properties has been an area of study for over a hundred years. To demonstrate this, we consider two questions.

First, suppose we are given the presentation of an arbitrary group  $G$ . Can we easily check whether or not  $G$  is a free group? This question was answered in the negative in the 1950s, where it was shown that such a question is undecidable.

The second question is one that is often asked when a property of a structure is first defined: do the substructures of this structure also have this property? In particular, are the subgroups of a free group also free groups? This was answered in the positive in the 1920s, first in a restricted form by Nielsen, then the full form independently by Dehn and Schreier. However, such proofs were not trivial, and required high level topological or algebraic methods to get to a rigorous proof. Understanding the proofs would require understanding very high level results in group theory, far beyond the level of an undergraduate course in abstract algebra.

Answers to the both of these questions do exist in simpler forms within the field of Combinatorial/Geometric Group Theory, which in part seeks to understand a group's algebraic properties by looking at how it acts on a space. This extra assumption on how a group acts upon a space makes these results much easier to understand at the undergraduate level.

An answer to the first question can be found thanks to the Ping-Pong Lemma, which describes a sufficient condition for a group to be free, so long as it acts in a certain way on a set  $X$ :

**Theorem 1.1** (The Ping-Pong Lemma). *Let  $\{g_1, g_2, \dots, g_n\}$  generate a group  $G$ , which acts on a set  $X$ . If*

(i)  *$X$  contains  $n$  subsets  $X_1, \dots, X_n$  such that  $X_i \cap X_j = \emptyset$ , and*

(ii)  *$g_i^k(X_j) \subset X_i$  for all nonzero powers  $k$  and  $i \neq j$ .*

*Then  $G$  is isomorphic to  $F_n$ .*

More interestingly, an answer to question two is a corollary of a result proven by Serre in 1969, who showed that free groups may be classified by how they act on trees.

**Theorem 1.2** (Serre's Theorem).  *$G$  is a free group if and only if it acts freely on a tree.*

**Corollary 1.3** (Nielsen-Schreier Theorem). *Every subgroup of a free group is itself free.*

In this paper, we present full proofs and explanations of both results. First, we prove the Ping-Pong Lemma both in the case of  $n = 2$ , and in the general case for arbitrary  $n$ . We also apply the Ping-Pong Lemma, and another result, to show that free groups of any rank, including one of infinite rank, lies within the matrix group  $SL_2(\mathbb{Z})$ . Afterwards, we prove both sides of Serre's Theorem. In proving the forwards direction, we define a special graph for our free group to act on called the Cayley Graph, and then show that free groups acts freely on their corresponding Cayley Graphs. To prove the reverse direction, we show that every group acting on a tree induces a special tiling of the tree, before finding a free generating set using this tiling.

# Chapter 2

## Free Groups

### 2.1 The Universal Property

For us to work with free groups, we first need to understand what it means for a group to be “free.”

**Definition 2.1.** Let  $S$  be a set. Then a group  $F_S$  is called the **free group generated by  $S$**  if the following holds: for any function  $f$  from  $S$  to a group  $G$ , there is a unique homomorphism  $\varphi : F_S \rightarrow G$  such that the below diagram commutes (where  $\iota$  denotes inclusion).

$$\begin{array}{ccc} & F_S & \\ \iota \nearrow & \downarrow \varphi & \\ S & \xrightarrow{f} & G \end{array}$$

We call  $S$  the **generating set** of  $F_S$ .

The condition that  $F_S$  must satisfy is called the **universal property**, and plays an important role in developing free structures in category theory. A priori, it is not obvious that such a group exists for every such set. However, not only is there such a group, but this group is unique up to isomorphism.

**Theorem 2.2.** *For any set  $S$ , there is a free group  $F_S$  generated by  $S$ .*

*$F_S$  is also unique up to isomorphism; if  $F_S^1, F_S^2$  are both free groups generated by  $S$ , with inclusion maps given by  $\iota_1 : S \rightarrow F_S^1, \iota_2 : S \rightarrow F_S^2$ , then there are unique isomorphisms  $\hat{\iota}_2 : F_S^1 \rightarrow F_S^2$  and  $\hat{\iota}_1 : F_S^2 \rightarrow F_S^1$  such that*

$$\hat{\iota}_2 \circ \iota_1 = \iota_2, \quad \hat{\iota}_1 \circ \iota_2 = \iota_1, \quad \hat{\iota}_1^{-1} = \hat{\iota}_2$$

*Proof.* We let  $F_S$  be the set of reduced words with letters given by elements of  $S$ , with the operation given by concatenation and reduction. Now, for any  $f : S \rightarrow G$ , where  $G$  is a group, we define  $\varphi$  as follows: we send the empty word  $\emptyset$  to  $1_G$ , the identity element of  $G$ , and for any  $s \in S$ ,

$$\varphi(s) = f(s)$$

where we think of  $s$  as a word consisting of a single letter in  $F_S$ .  $\varphi$  can then be uniquely extended to all other words by the fact that it is a homomorphism.

We now show uniqueness. Setting  $G = F_S^1$  and  $f = \iota_1$  in the diagram from Definition 2.1 gives us

$$\begin{array}{ccc} & F_S^1 & \\ \iota_1 \nearrow & \downarrow \text{id}_{F_S^1} & \\ S & \xrightarrow{\iota_1} & F_S^1 \end{array}$$

Similarly, setting  $G = F_S^2$  and  $f = \iota_2$ , we get

$$\begin{array}{ccc} & F_S^1 & \\ \iota_1 \nearrow & \downarrow \hat{\iota}_2 & \\ S & \xrightarrow[\iota_2]{} & F_S^2 \end{array}$$

from which we get that  $\hat{\iota}_2 \circ \iota_1 = \iota_2$ . Swapping  $F_S^1$  and  $F_S^2$  and their respective maps gives us  $\hat{\iota}_1$  such that  $\hat{\iota}_1 \circ \iota_2 = \iota_1$ .

Notice that  $\hat{\iota}_1 \circ \hat{\iota}_2 : F_S^1 \rightarrow F_S^1$  satisfies

$$\hat{\iota}_1 \circ \hat{\iota}_2 \circ \iota_1 = \hat{\iota}_1 \circ \iota_2 = \iota_1$$

satisfying the same equation that  $\text{id}_{F_S^1}$  does. By uniqueness, we get that  $\hat{\iota}_1 \circ \hat{\iota}_2 = \text{id}_{F_S^1}$ . A similar argument shows that  $\hat{\iota}_2 \circ \hat{\iota}_1 = \text{id}_{F_S^2}$ . This completes the proof.  $\square$

The presence of a set  $S$  seems to complicate our definition, as we must define free groups in terms of this generating set. However, we can greatly reduce the number of sets we must consider by making a key observation. Suppose  $S, T$  are sets such that  $|S| = |T|$ , and let  $F_S, F_T$  be the corresponding free groups. Recall that we construct each free group by sending each element of the generating set to a letter in the free group, and define the homomorphisms  $\varphi$  by how they map these letters. As  $S, T$  are the same size, there is a bijective map  $h : S \rightarrow T$ . More importantly, this map lifts to one on the free groups

$$\psi : F_S \rightarrow F_T, \quad \psi(\iota_S(s)) = \iota_T(h(s))$$

As  $h$  is a bijection on the generating sets, and each free group is uniquely determined by the letters these sets map to,  $\psi$  is an isomorphism of the groups. We have just proven

**Theorem 2.3.** *Every free group generated by sets of the same cardinality are isomorphic.*

Because of this, we only need to consider free groups generated by sets of the form  $\{1, 2, \dots, n\}$ ; the size of this set is called the free group's **rank**. We will thus denote these groups by  $F_n$ , called the **free group of rank  $n$** . For the purposes of this paper, we will restrict our study to free groups of countable rank, meaning we consider  $F_n$  for all  $n \in \mathbb{N}$ , as well as  $F_\infty$ , the free group with generators  $x_i$  for all natural  $i$ .

The construction of free groups makes them naturally important in the study of arbitrary groups. The first to notice this fact was Walther von Dyck, one of the founders of Combinatorial Group Theory, and the first to study groups via generators and relations. It was through this study that he found the usefulness of free groups:

**Theorem 2.4.** *Every group  $G$  is the quotient of a free group.*

*Proof.* Let  $G$  have a presentation given by  $\langle g_1, g_2, \dots, g_n | R \rangle$ . Let  $S = \{1, 2, \dots, n\}$  and let  $F_n$  be the free group generated by  $S$ . Consider the map

$$f : S \rightarrow G, \quad f(i) = g_i$$

Then by the universal property, let  $\varphi : F_n \rightarrow G$  be the unique homomorphism associated with  $f$ . It follows that

$$\ker \varphi = \{g = g_1^{e_1} \cdots g_n^{e_n} : g = 1\} = R$$

In other words,  $\ker \varphi$  is the set of all relations on  $G$ . Thus, we get that

$$G = F_n / \ker \varphi$$

□

Because of von Dyck's Theorem, we are able to define free groups in a very simple manner. By construction, the relations within a free group  $F_n$  are strictly those saying that  $x_i x_i^{-1} = 1$  for any  $i \in \{1, 2, \dots, n\}$ . Thus, free groups have the simplest possible presentation, and can be defined as just a set of generators with no relations:

**Definition 2.5.** The **free group of rank  $n$**  is the group  $F_n$ , the set of all reduced words in symbols  $\{x_1, \dots, x_n\}$ , with operation given by concatenation and reduction. Its presentation is

$$\langle a_1, a_2, \dots, a_n \rangle$$

and the identity is the empty word  $\emptyset$ .

Because of this result, we are able to translate the algebraic definition of a group, that which uses the universal property, to one using reduced words, which are much easier to work with.

## 2.2 The Rank of Subgroups of $F_2$

Rank is a group theoretic analogue of dimension for vector spaces.

**Definition 2.6.** The **rank** of a group  $G$  is the size of the smallest possible generating set for  $G$ .

$$\text{rank}(G) = \min\{|X| : X \subseteq G, \langle X \rangle = G\}$$

For a vector space  $V$ , any subspace  $W$  will have at most dimension  $\dim V$ , which makes intuitive sense. However, despite being an analogue of dimension for groups, this property does not hold for arbitrary groups; it holds for special types of groups, such as finite abelian groups.

The subgroups of free groups can oftentimes have larger rank than the free group itself. In particular,  $F_2$ , not only contains subgroups of larger rank, but they contain subgroups of every possible rank larger than 2, including subgroups of countably infinite rank!

**Theorem 2.7.** (a) For all  $n > 2$ ,  $F_n \subset F_2$ .

(b)  $F_\infty \subset F_2$ , where

$$F_\infty = \langle a_1, a_2, a_3, \dots \rangle$$

is the free group of countably infinite rank.

*Proof.* (a) We write  $F_2 = \langle a, b \rangle$ , and define

$$\begin{aligned} w_1 &= a \\ w_2 &= b^{-1}ab \\ w_3 &= b^{-2}ab^{-2} \\ &\vdots \\ w_n &= b^{-n}ab^n \end{aligned}$$

With the operations of concatenation and reduction, these form the generating set of a group. It thus suffices to show that there are no relations between these generators. First notice that for any  $w_i$  and any integer  $k$ :

$$\begin{aligned} w_i^k &= (b^{-i}ab^i)(b^{-i}ab^i)\cdots(b^{-i}ab^i) \\ &= b^{-i}a(b^i b^{-i})a(b^i b^{-i})a\cdots a(b^i b^{-i})ab^i \\ &= b^{-i}a^k b^i \end{aligned}$$

Now, let  $w = w_{i_1}^{k_1} \cdots w_{i_l}^{e_l}$  be an arbitrary word. We have that

$$\begin{aligned} w &= w_{i_1}^{k_1} \cdots w_{i_l}^{e_l} \\ &= (b^{-i_1}a^{e_1}b^{i_1})(b^{-i_2}a^{e_2}b^{i_2})\cdots(b^{-i_l}a^{e_l}b^{i_l}) \\ &= b^{-i_1}a^{e_1}b^{i_1-i_2}a^{e_2}b^{i_2-i_3}\cdots b^{i_{l-1}-i_l}a^{e_l}b^{i_l} \end{aligned}$$

and this is fully reduced. Moreover, it clearly cannot be expressed as the reduced product of any other set other product of the generators. Therefore, we conclude that there are no relations between the  $w_i$ 's, and hence they generate the free group  $F_n$ , so  $F_n \subset F_2$ .

(b) The proof is identical to (a), with the generating set of  $F_\infty$  being the set of  $w_i$  for all  $i \in \mathbb{N}$ . □

One might ask if this result generalizes to free groups of rank higher than 2? In particular, for any  $k \in \mathbb{N}$ , is it true that  $F_n \subset F_k$  for  $n > k$ ? If so, we'd get a chain of inclusions

$$F_2 \supset F_3 \supset F_4 \supset \cdots$$

The answer is no, and follows from a result of Schreier:

**Theorem 2.8** (Schreier Index Formula). *If  $H$  is a subgroup of  $F_n$  with index  $e$ , then  $H$  has rank*

$$1 + e(n - 1)$$

## Chapter 3

# A Group Acting Freely on a Tree is Free

We first prove the forwards direction.

**Theorem 3.1.** *If  $T$  is a tree and  $G$  is a group such that  $G \curvearrowright T$  freely, then  $G$  is isomorphic to a free group.*

To this end, we let  $T$  be a tree and we let  $G$  be a group such that  $G \curvearrowright T$  freely.

### 3.1 G-Tilings

**Definition 3.2.** A **tile** of a tree  $T$  is any subtree of  $T$ .

A **tiling** of  $T$  is a collection of tiles of  $T$  such that

1. The union of all tiles is  $T$ , and
2. Two tiles intersect, at most, at a vertex.

If  $G \curvearrowright T$ , and if

3. There is a tile  $T_0$  such that the set of all tiles is  $\{gT_0 : g \in G\}$

then we call the tiling a  **$G$ -tiling**.

It is not obvious a priori that  $T$  will always have a  $G$ -tiling. However, one can show that it always does.

**Theorem 3.3.** *Let  $G$  be a group and  $T$  a tree. If  $G \curvearrowright T$ , then  $T$  has a  $G$ -tiling.*

Before proving this we need two more objects.

**Definition 3.4.** For a tree  $T$ , its **barycentric subdivision** is the tree  $T'$  obtained by putting a new vertex into the middle of every edge in  $T$ .

Without loss of generality, we now assume that every non-leaf vertex in  $T$  has at least degree 2; if not, then use the barycentric subdivision  $T'$ , which has this property. The reason for this will become relevant later in Lemma 3.6.

Now, how exactly will we create our  $G$ -tiling? We can do this using our group action. Let  $v_0 \in V(T)$  be a vertex in the tree, and consider its orbit under the group action.

$$\mathcal{O}(v_0) = \{gv_0 : g \in G\}$$

For every vertex  $v \in V(T)$ , there is a  $gv_0 \in \mathcal{O}(v_0)$  that is closest to  $v$  under the path metric. From this, we define a subgraph  $T_g$  as follows:

$$\begin{aligned} V(T_g) &= \{v \in T : d(v, gv_0) \leq d(v, g'v_0) \ \forall g' \neq g\} \\ E(T_g) &= \{(u, v) \in E(T) : u, v \in V(T_g)\} \end{aligned}$$

In other words,  $T_g$  is a subgraph whose vertex set is all those vertices that are closest to  $gv$  in  $T$  under the path metric, and whose edge set is all those edges whose endpoints are both vertices in  $T_g$ .

**Lemma 3.5.** *Let  $G \curvearrowright T$ . Then  $T_g$  is a tile for all  $g \in G$ .*

*Proof.* It suffices to prove that  $T_g$  is connected, as every connected subgraph of a tree is a subtree. Let  $w \in V(T_g)$ . There is a unique edge path  $w$  to  $gv$ ; suppose it is of length  $n$ .

$$w_0 = w, w_1, w_2, \dots, w_{k-1}, w_n = gv$$

Then  $d(w_1, gv_0) = n - 1$ . We claim that  $w_i \in V(T_g)$  for all  $i$ . Starting with  $w_1$ , suppose by way of contradiction that  $w_1 \notin V(T_g)$ . Then there is a  $g'$  such that

$$d(w_1, g'v_0) = m < n - 1$$

But then going back on edge to  $w$ , we get that

$$d(w, g'v_0) \leq m + 1 < n$$

which contradicts the fact that  $w \in V(T_g)$ . Thus,  $w_1 \in V(T_g)$ . This logic can be repeated for all other vertices in the path, ensuring that all are in  $V(T_g)$ , as desired. This means that  $T_g$  is connected, hence a subtree of  $T$ , and thus, is a tile.  $\square$

Now that we know the  $T_g$ 's are in fact tiles, we should also check that they cover all of  $T$ .

**Lemma 3.6.** *The union of all tiles  $T_g$  for all  $g$  is equal to  $T$ .*

*Proof.* It is evident that  $\bigcup_{g \in G} V(T_g) = V(T)$ , since every vertex must be closest to some  $gv_0$ . It is also evident that

$$\bigcup_{g \in G} E(T_g) \subseteq E(T)$$

So it suffices to show the reverse inclusion: every edge in  $T$  belongs to some  $T_g$ . Let  $(u, w) \in E(T)$ . Then by construction of  $T$ , the distance to  $\mathcal{O}(v_0)$  from one of  $u, w$  is odd, while the distance from the other is even. Without loss of generality, suppose that

$$d(u, \mathcal{O}(v_0)) \text{ is even} \quad \text{and} \quad d(w, \mathcal{O}(v_0)) \text{ is odd}$$

that  $d(u, \mathcal{O}(v_0)) < d(w, \mathcal{O}(v_0))$ , and that  $u \in V(T_g)$ . By the Triangle Inequality, we have that

$$d(w, gv_0) \leq d(w, u) + d(u, gv_0) = d(u, gv_0) + 1$$

As the distance from  $w$  to  $\mathcal{O}(v_0)$  is an integer, and cannot be smaller than or equal to  $d(u, gv_0)$ , the only possibility is that it is equal to  $d(w, gv_0)$ . Thus,

$$d(w, gv_0) \leq d(w, g'v_0) \forall g' \neq g$$

so  $w \in V(T_g)$ . We conclude that  $(u, w) \in E(T_g)$ , as desired.  $\square$

The second condition, that two tiles share at most a vertex, is somewhat obvious. Indeed, if  $(u, w)$  is in both  $E(T_g)$  and  $E(T_{g'})$ , then  $u, w$  are of equal distance to  $gv_0$  and  $g'v_0$ . But assuming  $u$  is of equal distance to  $gv_0$  and  $g'v_0$ ,  $w$  must be closer to one of  $gv_0$  or  $g'v_0$ , since it is along the unique edge path from  $u$  to one of them.

The final condition can be proven using another lemma:

**Lemma 3.7.** *For  $g, h \in G$ ,  $gT_h = T_{gh}$ .*

*Proof.* It suffices to show that  $g$  maps vertices in  $T_h$  to those in  $T_{gh}$ , as they're completely determined by vertices. Let  $u \in V(T_h)$ . Then for all  $k \in G$ ,

$$d(u, hv_0) \leq d(u, kv_0)$$

Left multiplication by  $g^{-1}$  is a bijection of  $G$ , so we may write

$$d(u, hv_0) \leq d(u, (g^{-1}k)v_0)$$

For all  $k \in G$ . Now, as  $G \curvearrowright T$ , and  $T$  is a metric space using the path metric, the action of  $G$  preserves distances. Thus,

$$d(gu, (gh)v) \leq d(gu, kv)$$

so  $gu \in T_{gh}$ , as desired.  $\square$

*Proof of Theorem 3.3.* By Lemmas 3.5 and 3.6, the collection  $\{T_g : g \in G\}$  is a tiling of  $T$ . Setting  $h = 1$  and applying Lemma 3.7, all tiles are of the form  $gT_1$  for some  $g \in G$ .  $\square$

We have thus shown that any group acting on a tree induces a  $G$ -tiling on the tree.

### 3.2 Proof of Theorem 3.1

Now that we know  $G$  induces a  $G$ -tiling on our tree  $T$ , we can use it to create a free generating set.

We let  $\{T_g\}$  be the  $G$ -tiling found in Theorem 3.3. As  $G$  acts freely on  $T$ ,  $\text{Stab}_G(v_0) = \{1\}$ , so by the Orbit-Stabilizer Theorem,

$$|\mathcal{O}(v_0)| = 1 \cdot |G| = |G|$$

Thus the elements of  $G$  are in bijective correspondence with the elements of the orbit of  $v_0$ . In other words, every  $g \in G$  corresponds to a tile  $gT_1$ . Define

$$S = \{g \in G : (gT_1) \cap T_1 \neq \emptyset\}$$

to be the set of tiles that intersect with  $T_1$ .

**Lemma 3.8.** *S generates G*

*Proof.* First we show that if  $s \in S$ , so too is  $s^{-1}$ . Indeed let  $s \in S$ . Then there is a vertex  $w$  such that

$$\begin{aligned} (sT_1) \cap T_0 &= \{w\} \\ \implies s^{-1}(sT_1) \cap (s^{-1}T_1) &= \{s^{-1}w\} \\ \implies T_1 \cap (s^{-1}T_1) &= \{s^{-1}w\} \end{aligned}$$

so  $s^{-1} \in S$ .

Now let  $g \in G$  and consider  $gv_0$ . There is a unique path in  $T$  from  $gv_0$  to  $v_0$ , passing through tiles

$$T_{g_n}, T_{g_{n-1}}, \dots, T_{g_1}, T_{g_0}$$

We claim that  $g_{i-1}^{-1}g_i \in S$ . The above path travels from  $T_{g_{i+1}}$  to  $T_{g_i}$  without passing through a middle tile, so

$$T_{g_{i+1}} \cap T_{g_i} \neq \emptyset$$

Applying  $g_i^{-1}$ , we get

$$g_i^{-1}T_{g_{i+1}} \cap T_1 \neq \emptyset \implies T_{g_i^{-1}g_{i+1}} \cap T_1 \neq \emptyset$$

so  $g_i^{-1}g_{i+1} \in S$ , and we write it as  $s_i$ . Then, we have that

$$\begin{aligned}
g &= g_n \\
&= g_{n-1}g_{n-1}^{-1}g_n \\
&\vdots \\
&= g_1^{-1}g_1g_2^{-1}g_2 \cdots g_{n-1}^{-1}g_n \\
&= s_1s_2 \cdots s_n
\end{aligned}$$

so  $g$  may be written as a product of elements of  $S$ , meaning  $S$  generates  $G$ .  $\square$

With all these tools now in place, the last thing we need to do is show that there are no relations between elements of  $S$ .

*Proof of Theorem 3.1.* By Lemma 3.3 we have a  $G$ -tiling  $\{T_g\}$ , which we know generates  $G$  by Lemma 3.8. We claim that every  $g \in G$  may be written as a unique, reduced product of elements of  $S$ . Let

$$g = s_1s_2 \cdots s_k$$

with  $s_i \neq s_{i+1}^{-1}$ . We claim there is a unique path from  $gv_0$  to  $v_0$  that passes through, in order, the tiles

$$T_g = T_{s_1s_2 \cdots s_k}, T_{s_1s_2 \cdots s_{k-1}}, \dots, T_{s_1}, T_1$$

As  $T_{s_1} \cap T_1 = \{w_1\}$  for some  $w_1 \in V(T)$ ,  $T_{s_1} \cup T_1$  is a tree, where  $v_0 \in V(T_1)$  and  $s_1v_0 \in T_{s_1}$ . Thus, there is a unique path from  $s_1v_0$  to  $v_0$  in  $T_{s_1} \cup T_1$ . Similarly, we have

$$T_{s_1s_2} \cap T_{s_1} = s_1^{-1}(T_{s_2} \cap T_1) = \{s_1^{-1}w_2\}$$

for some  $w_2 \in V(T)$ , so  $T_{s_1s_2} \cup T_{s_1}$  is a tree. Again, this means there is a unique path from  $s_1v_0$  to  $s_1s_2v_0$  in  $T_{s_1s_2} \cup T_1$ . If we continue this process inductively, we get the desired unique path. This path completely determines the word  $s_1s_2 \cdots s_k$  representing  $g$ , and thus  $g$  can only be written as this freely reduced word in  $S$ . This tells us that  $S$  has no relations, and thus  $G$  is free, as desired.  $\square$

**Corollary 3.9** (Nielsen-Schreier Theorem). *Every subgroup of a free group is also a free group.*

*Proof.* A free action of a group  $G$  on a tree trivially restricts to an action of  $H \leq G$  on the tree. As  $1 \in H$ , if the action of  $G$  is free, so too must the action of  $H$ .  $\square$

Lorem ipsum dolor sit amet, consectetur adipiscing elit. Ut purus elit, vestibulum ut, placerat ac, adipiscing vitae, felis. Curabitur dictum gravida mauris. Nam arcu libero, nonummy eget, consectetur id, vulputate a, magna. Donec vehicula augue eu neque. Pellentesque habitant morbi tristique senectus et netus et malesuada fames ac turpis egestas. Mauris ut leo. Cras viverra metus rhoncus sem. Nulla et lectus vestibulum urna fringilla ultrices. Phasellus eu tellus sit amet tortor gravida placerat. Integer sapien est, iaculis in, pretium quis, viverra ac, nunc. Praesent eget sem vel leo ultrices bibendum. Aenean faucibus. Morbi dolor nulla, malesuada eu, pulvinar at, mollis ac, nulla. Curabitur auctor semper nulla. Donec varius orci eget risus. Duis nibh mi, congue eu, accumsan eleifend, sagittis quis, diam. Duis eget orci sit amet orci dignissim rutrum.

  Nam dui ligula, fringilla a, euismod sodales, sollicitudin vel, wisi. Morbi auctor lorem non justo. Nam lacus libero, pretium at, lobortis vitae, ultricies et, tellus. Donec aliquet, tortor sed accumsan bibendum, erat ligula aliquet magna, vitae ornare odio metus a mi. Morbi ac orci et nisl hendrerit mollis. Suspendisse ut massa. Cras nec ante. Pellentesque a nulla. Cum sociis natoque penatibus et magnis dis parturient montes, nascetur ridiculus mus. Aliquam tincidunt urna. Nulla ullamcorper vestibulum turpis. Pellentesque cursus luctus mauris.

  Nulla malesuada porttitor diam. Donec felis erat, congue non, volutpat at, tincidunt tristique, libero. Vivamus viverra fermentum felis. Donec nonummy pellentesque ante. Phasellus adipiscing semper elit. Proin fermentum massa ac quam. Sed diam turpis, molestie vitae, placerat a, molestie nec, leo. Maecenas lacinia. Nam ipsum ligula, eleifend at, accumsan nec, suscipit a, ipsum. Morbi blandit ligula feugiat magna. Nunc eleifend consequat lorem. Sed lacinia nulla vitae enim. Pellentesque tincidunt purus vel magna. Integer non enim. Praesent euismod nunc eu purus. Donec bibendum quam in tellus. Nullam cursus pulvinar lectus. Donec et mi. Nam vulputate metus eu enim. Vestibulum pellentesque felis eu massa.

  Quisque ullamcorper placerat ipsum. Cras nibh. Morbi vel justo vitae lacus tincidunt ultrices. Lorem ipsum dolor sit amet, consectetur adipiscing elit. In hac habitasse platea dictumst. Integer tempus convallis augue. Etiam facilisis. Nunc elementum fermentum wisi. Aenean placerat. Ut imperdiet, enim sed gravida sollicitudin, felis odio placerat quam, ac pulvinar elit purus eget enim. Nunc vitae tortor. Proin tempus nibh sit amet nisl. Vivamus quis tortor vitae risus porta vehicula.

  Fusce mauris. Vestibulum luctus nibh at lectus. Sed bibendum, nulla a faucibus semper, leo velit ultricies tellus, ac venenatis arcu wisi vel nisl. Vestibulum diam. Aliquam pellentesque, augue quis sagittis posuere, turpis lacus congue quam, in hendrerit risus eros eget felis. Maecenas eget erat in sapien mattis porttitor. Vestibulum porttitor. Nulla facilisi. Sed a turpis eu lacus commodo facilisis. Morbi fringilla, wisi in dignissim interdum, justo lectus sagittis dui, et vehicula libero dui cursus dui. Mauris tempor ligula sed lacus. Duis cursus enim ut augue. Cras ac magna. Cras nulla. Nulla egestas. Curabitur a leo. Quisque egestas wisi eget nunc. Nam feugiat lacus vel est. Curabitur consectetur.

  Suspendisse vel felis. Ut lorem lorem, interdum eu, tincidunt sit amet, laoreet vitae, arcu. Aenean faucibus pede eu ante. Praesent enim elit, rutrum at, molestie non, nonummy vel, nisl. Ut lectus eros, malesuada sit amet, fermentum eu, sodales cursus, magna. Donec eu purus. Quisque vehicula, urna sed ultricies auctor, pede lorem egestas dui, et convallis elit erat sed nulla. Donec luctus. Curabitur et nunc. Aliquam dolor odio, commodo pretium, ultricies non, pharetra in, velit. Integer arcu est, nonummy in, fermentum faucibus, egestas vel, odio.