



CPSC 457 Tutorial

Week 4

System calls

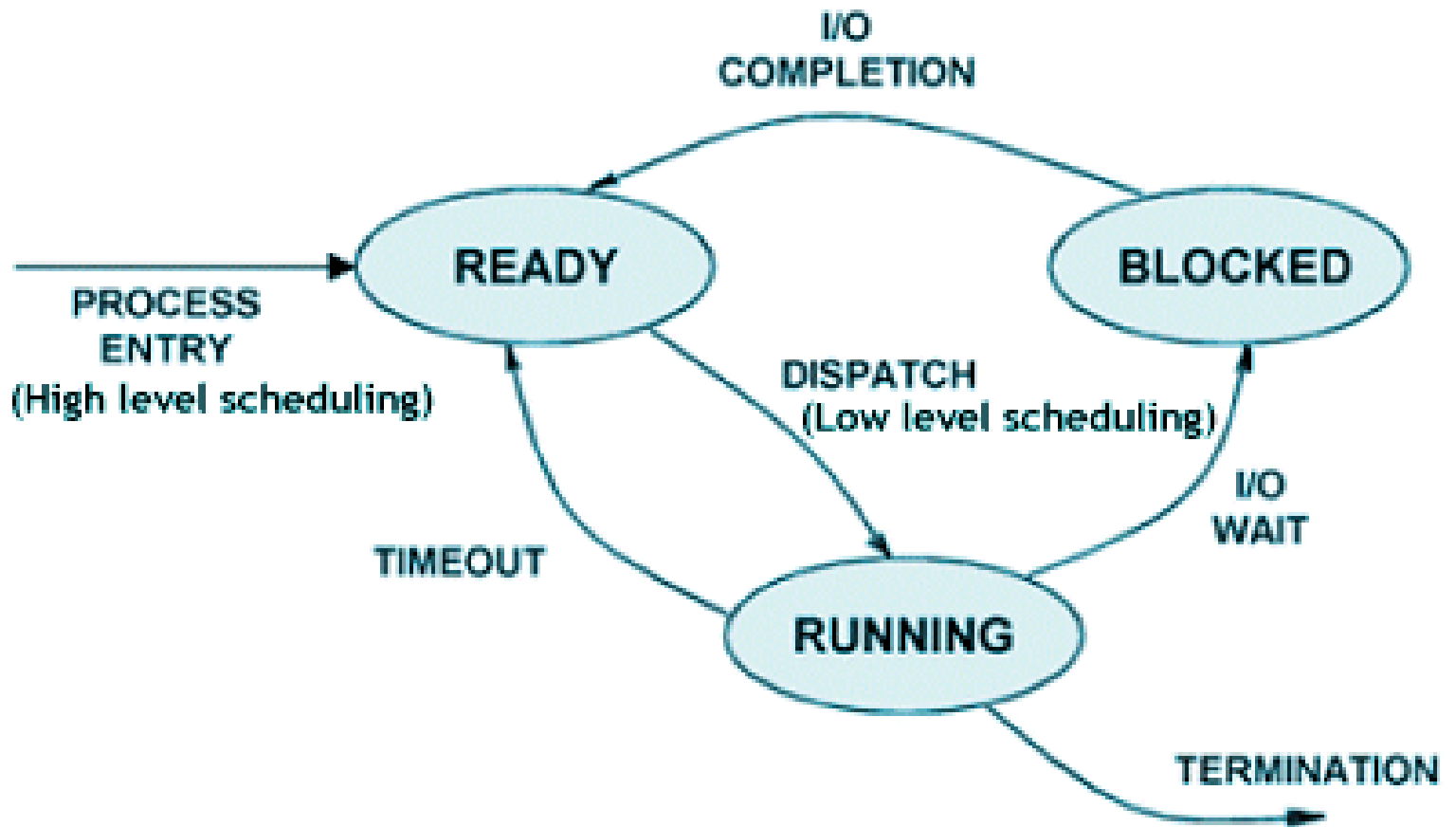


Figure 1: Process state diagram

Source: <https://www.technologyuk.net/computing/operating-systems/process-management.shtml>

Threads

- Basic unit of CPU utilization

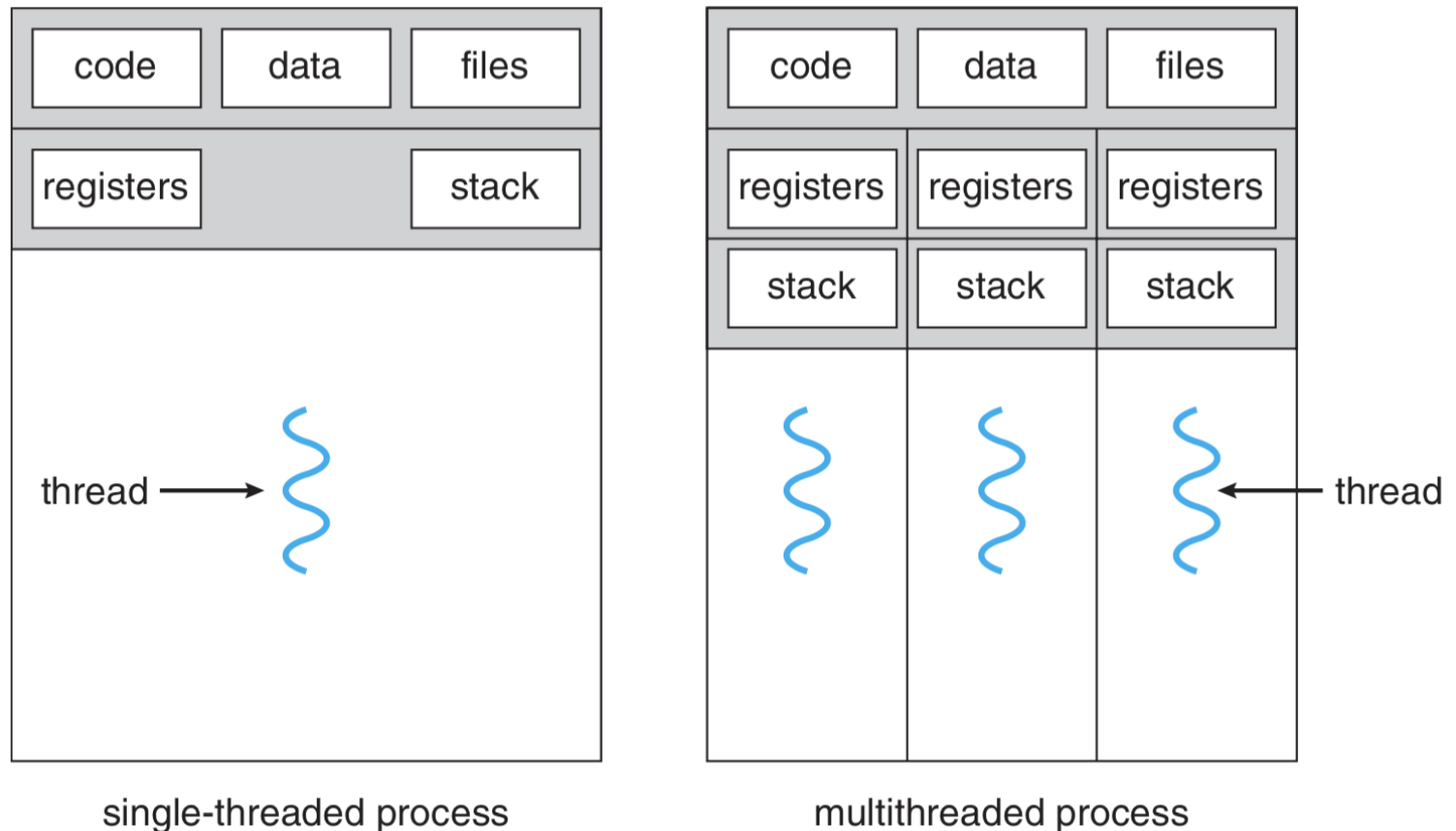


Figure 2: Single-threaded and multithreaded processes

Taken from Abraham Silberschatz, Peter Baer Galvin, and Greg Gagne. 2012. Operating System Concepts (9th ed.)

Threads (cont'd)

- Imagine an IDE program
- There will be multiple jobs performed simultaneously(!) while you are coding
 - Auto updater
 - Spell checker
 - Auto compilation
 - Text editor

Benefits

- **Responsiveness:** When one of the tasks is blocked, the others can continue
- **Resource sharing:** Sharing code and data allows running multiple threads on the same address space
- **Performance:** It is faster to create threads as they share resources of the process
- **Scalability:** While a single threaded process runs on one core, multiple threads may be running on different cores

POSIX threads

- `man pthreads` contains descriptions and examples
- `int pthread_create(pthread_t *thread, const pthread_attr_t *attr, void *(*start_routine) (void *), void *arg)`
- `void pthread_exit(void *value_ptr)`
- `int pthread_join(pthread_t thread, void **value_ptr);`

Thread concurrency

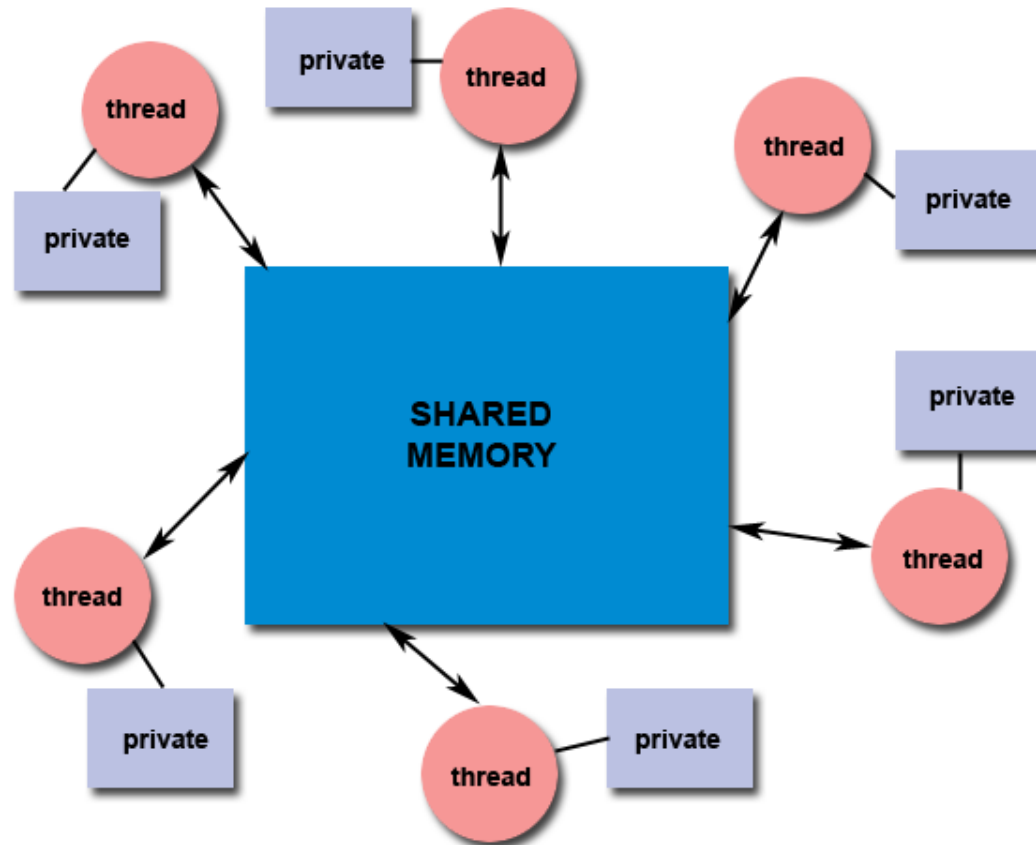
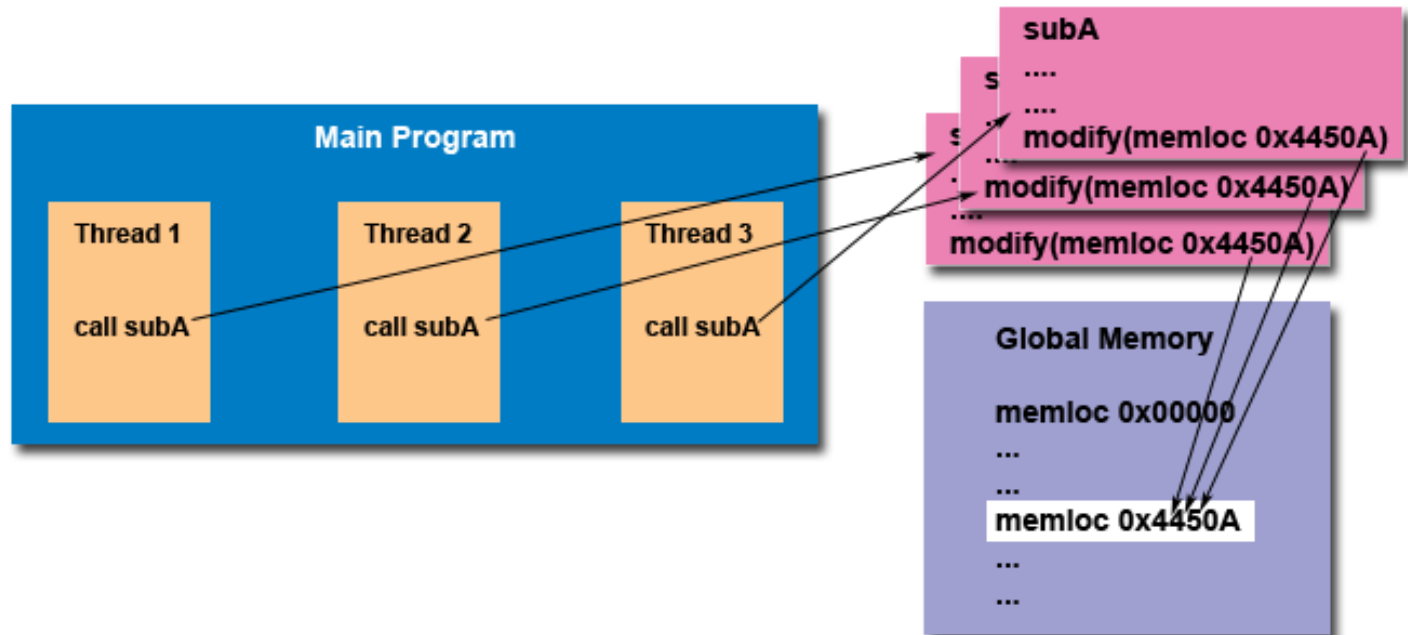


Figure 1: Multithread model

Taken from <https://computing.llnl.gov/tutorials/pthreads/>

Thread concurrency (cont'd)

- Accessing and modifying shared data should be secured in multithreaded processes
- Otherwise **race conditions** may occur in **critical sections**



Race condition example

```
int cnt = 0;
```

```
void* incr(void* args) {  
    cnt++;  
    return NULL;  
}
```

load cnt to reg0
increment reg0
store reg0 back to cnt



Thread 1	Thread 2	cnt
reg0 = cnt		0
reg0 = reg0 + 1		0
	reg1 = cnt	0
	reg1 = reg1 + 1	0
cnt = reg0		1
	cnt = reg1	1

Mutex

- Short of *mutual exclusion*
- Used to prevent race conditions
- There are two main operations:
 - lock()
 - unlock()
- These operations are atomic, meaning that only one thread can take the lock
- When a mutex is taken by one of the threads, the others wait for it to be released

POSIX Mutex

- `pthread_mutex_t` for type
- `pthread_mutex_init(pthread_mutex_t *)`
- `pthread_mutex_lock(pthread_mutex_t *)`
- `pthread_mutex_unlock(pthread_mutex_t *)`
- `pthread_mutex_destroy(pthread_mutex_t *)`