COMP207 Database Development

Lecture 17

Review of query processing

and introduction of Distributed Databases

Chapter 4

Extensions of the simple relational database setup

Distributed databases

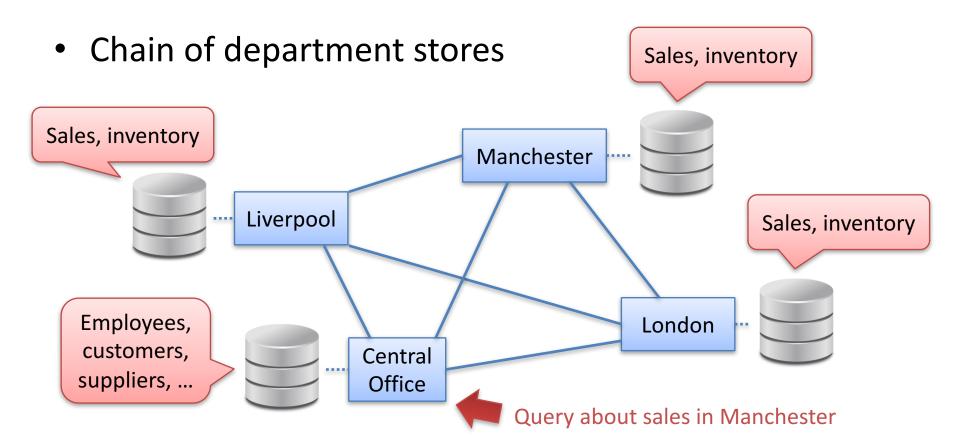
- Databases connected through a network
- More efficient by distributing tasks over several computers
- + other desirable properties

Digression: Map-reduce

- Framework for parallel processing of data
- Goes beyond relational data
- Next chapter: continues with non-relational data

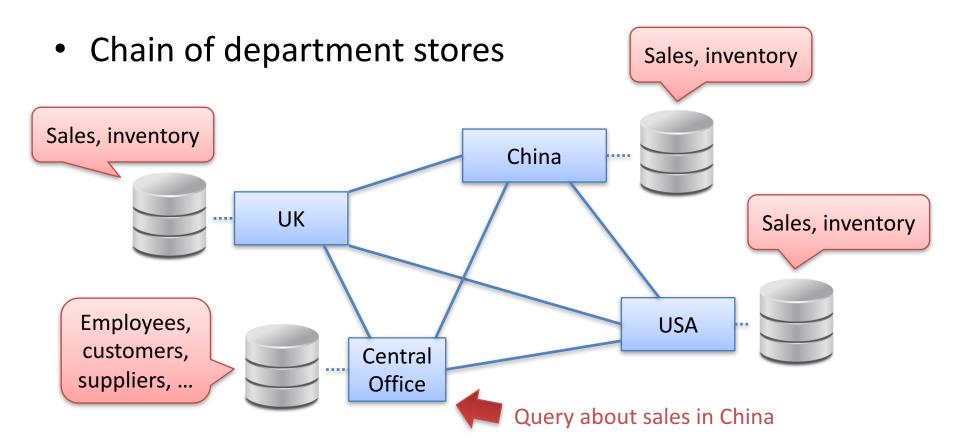
What is a Distributed Database?

Motivation



- Each site stores only data primarily relevant to it
- Distributed DBMS provide access to data at all sites

Motivation

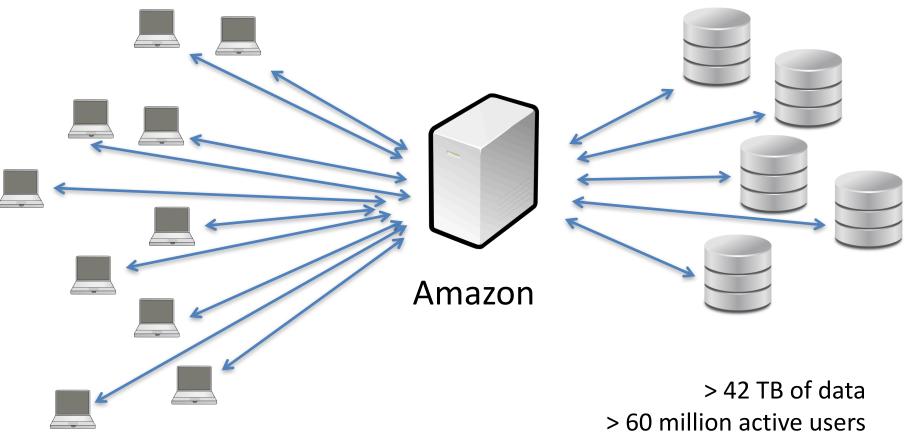


- Each site stores only data primarily relevant to it
- Distributed DBMS provide access to data at all sites

Other Applications

- More general: large organisations/companies
 - ...with different branches or offices
 - ...with different sub-companies
- Providing access to large datasets to many users (e.g., for an online store)
 - Distribute data over several computers
 - Computers could be at geographically separate locations
 - Possible advantages:
 - Balance workload & network traffic
 - Easier to extend capacity or scale to higher number of users

Recall Lecture 1...



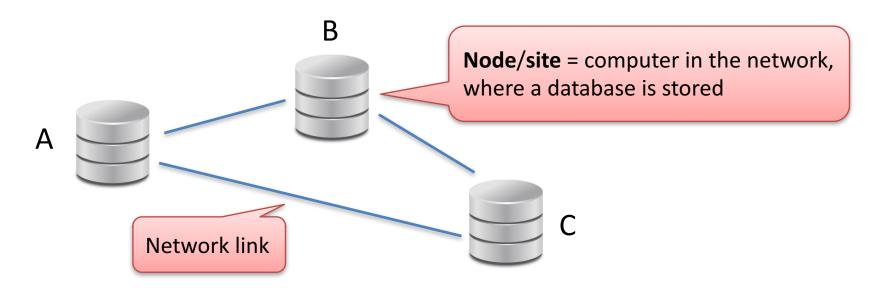
millions of users per day many at the same time

60 million active users millions of products

100s of servers

Distributed Databases

- Distributed Database:
 - Collection of multiple logically interrelated databases
 - Distributed over a computer network



Distributed DBMS: manages a distributed database

Advantages

Performance improvements

- Answer queries faster by distributing tasks over the nodes
- Reduces CPU time, disk accesses, communication cost, ...
 at individual nodes

Scalability

- Easier extension of the system: capacity, performance, ...
- Essentially just add a new node

Resilience

- Data can be replicated at geographically separate sites
- Catastrophic failures don't affect the entire system

How to Distribute Data?

Fragmentation

 Split database into different parts that can then be stored at different nodes

"Sharding"

Horizontal fragmentation

Vertical fragmentation

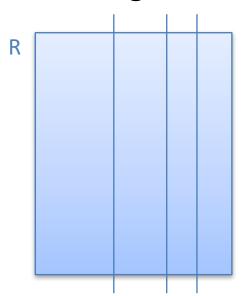
Fragment 1

Fragment 2

may overlap

Fragment 3

Fragment 4



Users don't see fragments, just the full relations

"Fragmentation transparency"

Horizontal Fragmentation Example

The chain of stores from our example might jointly

a+ aa	_		
store	a	relation	

item	date	price	purchaser	
Product A	17/10/2017	£29.99	Anna A.	← @Liverpool
Product B	19/10/2017	£199.99	Ben B.	← @Manchester
Product C	19/10/2017	£599.99	Chloe C.	← @Manchester
Product D	20/10/2017	£9.99	David D.	← @London
Product E	21/10/2017	£59.99	Emma E.	← @Liverpool
•••				

- Each site stores a relation that contains a subset of the tuples
- Entire relation = union of relations at the different sites

Does not exist

physically!

Horizontal Fragmentation Example

- Typically, tuples stored at different sites can be distinguished
 - by the value of one or a few attributes; or
 - by other conditions that are easy to test



store	item	date	price	purchaser
Liverpool	Product A	17/10/2017	£29.99	Anna A.
Manchester	Product B	19/10/2017	£199.99	Ben B.
Manchester	Product C	19/10/2017	£599.99	Chloe C.
London	Product D	20/10/2017	£9.99	David D.
Liverpool	Product E	21/10/2017	£59.99	Emma E.
•••	•••	•••	•••	

Vertical Fragmentation Example

The stores might also jointly store a relation

Does not exist physically!

item	date	purchaser	last_payment *
Product A	17/10/2017	Anna A.	30/09/2017
Product B	19/10/2017	Ben B.	30/09/2017
Product C	19/10/2017	Chloe C.	31/08/2017
			•••

@Central Office

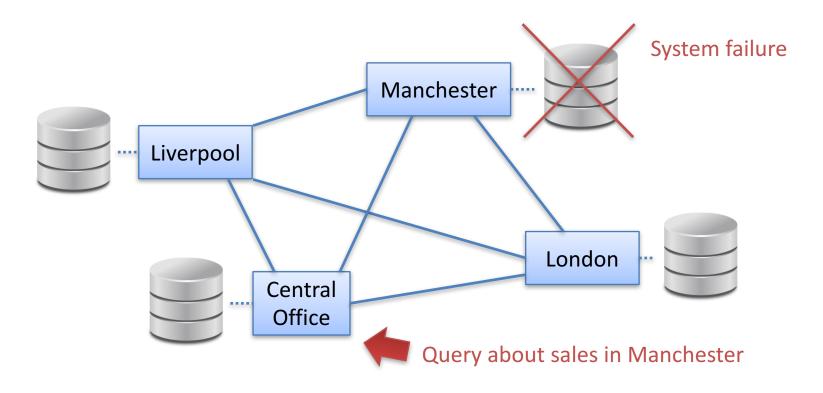
purchaser	last_payment	
Anna A.	30/09/2017	
Ben B.	30/09/2017	
Chloe C.	31/08/2017	

Stored at other sites

item	date	purchaser
Product A	17/10/2017	Anna A.
Product B	19/10/2017	Ben B.
Product C	19/10/2017	Chloe C.
	•••	

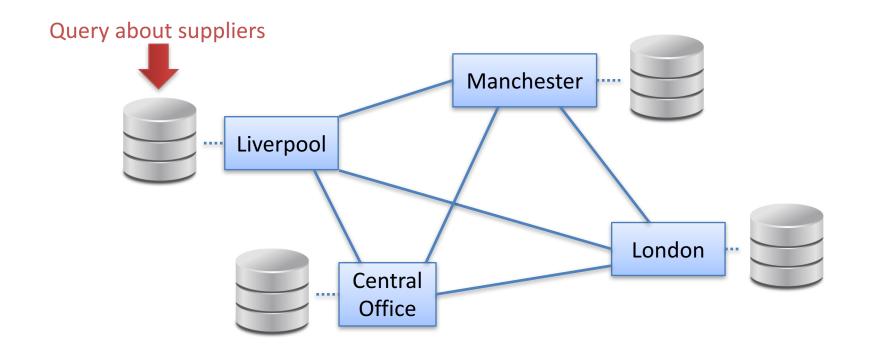
• Original relation \approx join of the fragments

Redundancy Improves Resilience



- Other sites keep copies of fragment stored at Manchester
- Allows us to answer queries involving data from Manchester

Redundancy Increases Efficiency



- Other sites keep copies of data about suppliers
- Allows stores to answer queries involving suppliers without establishing a connection to the central office

Replication

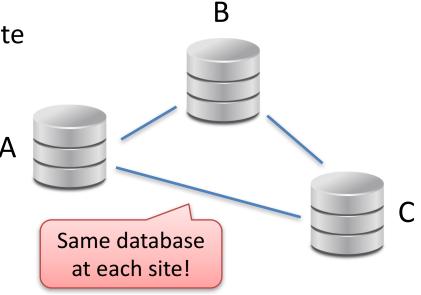
Controls how many sites keep a copy of a fragment

Full replication

- Each fragment stored at every site
 (→ there are no fragments)
- Faster query answering
- Very slow updates: consider every copy

No replication

Each fragment stored at a unique site



- Wide spectrum of partial replication
 - Limit number of copies of each fragment
 - Replicate only some fragments, ...

Transparency

- DBMSs hide how data is distributed over sites
- Transparency at different levels
 - Fragmentation transparency
 - Fragmentation is invisible to users
 - Users pose queries against the entire database
 - The distributed DBMS translates this into a query plan that fetches the required information from appropriate nodes
 - Replication transparency
 - Ability to store copies of data items / fragments at different sites
 - Replication is invisible to users
 - Distribution transparency
 - Naming transparency

Summary

- Distributed databases are...
 - Collection of multiple logically interrelated databases
 - Distributed over a computer network
- Advantages: Improved performance, scalability, resilience against failures
- Query answering more challenging
 - Goal: minimise communication
 - With joins: a strategy using semijoins to prefilter tuples may yield better communication costs
- Next lecture: Transaction management is more challenging, too!