主管

领导 审核 签字

- 1. [10 points] Design a DFA for $L = \{w \in \{0,1\}^* \mid w \text{ has exactly three 0s.}\}$
- 2. [10 points] Design an NFA for the language:

$$L = \{w \in \{a, b, c\}^* \mid w \text{ starts with } ac \text{ and ends with } cb.\}$$

- 3. [10 points] Design regular expressions for languages over $\Sigma = \{a, b\}$.
 - (1) All strings that do not end with aba.
 - (2) $L = \{w \mid w \text{ has no more than 5 } a\text{'s. }\}$
- 4. [10 points] Prove that the language $L = \{w \in \{a,b\}^* \mid w = w^R\}$ is not regular with pumping lemma.
- 5. [10 points] Consider the following ε -NFA.

- (1) Compute the ε -closure of each state.
- (2) Give all the strings of length three or less accepted by the automaton.
- (3) Convert the automaton to a DFA by subset construction. (diagram of transition function)
- 6. [10 points] Give a CFG for $L = \{a^i b^j c^k \mid i, j, k \ge 0 \text{ and } i = j + k\}.$
 - 7. [10 points] Find a grammar equivalent to

$$S \to AB \mid CA$$

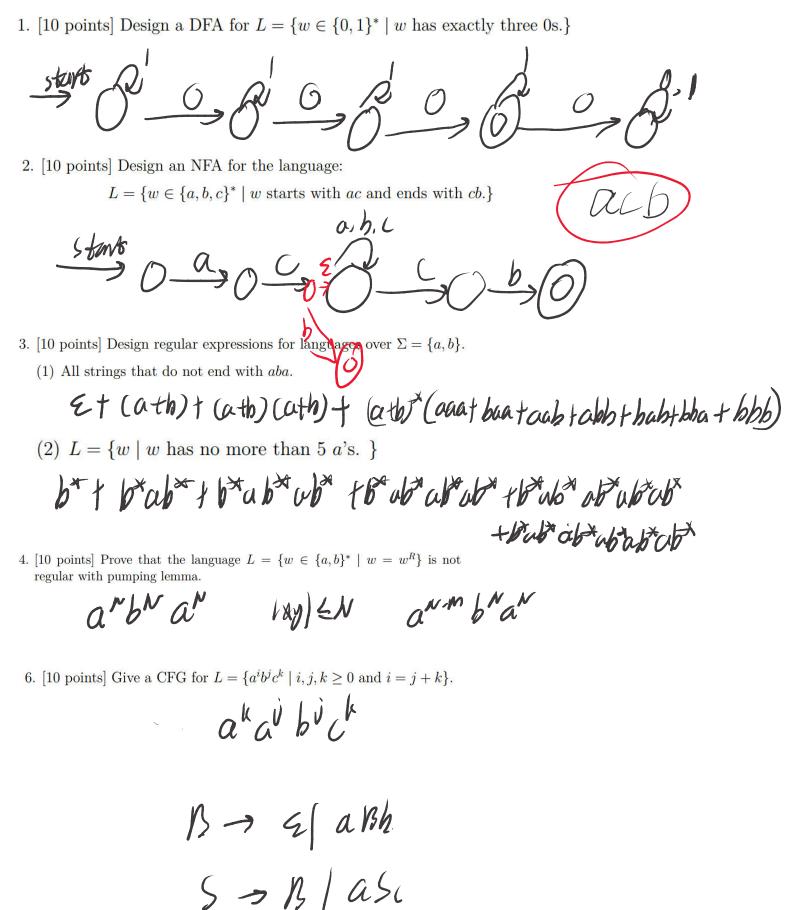
$$A \to a$$

$$B \to BC \mid AB$$

$$C \to aB \mid b$$

with no useless symbols.

- 8. [10 points] Design a PDA for $L_{eq} = \{w \in \{0,1\}^* \mid w \text{ contains the same}\}$ number of 0's and 1's \.
- 9. [10 points] Prove or disprove: if L_1 is CFL and $L_1 \cup L_2$ is also CFL, then L_2 must be CFL.
- [10. [10 points] Design Turing machine for the language $\{0^{2n}1^n \mid n \geq 0\}$.



5. [10 points] Consider the following ε -NFA.

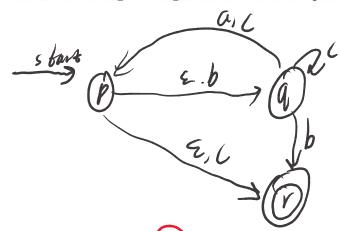
	ε	a	b	c
$\rightarrow p$	$\{q,r\}$	Ø	$\{q\}$	{r}
q	Ø	$\{p\}$	$\{r\}$	$\{p,q\}$
*r	Ø	Ø	Ø	Ø

(1) Compute the ε -closure of each state.

Eclose (p) = 4p, q, v}

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(2) Give all the strings of length three or less accepted by the automaton.



2+L+ (2+b)(c+a+c)(2+b) 5 (h+(atc/gy)

(3) Convert the automaton to a DFA by subset construction. (diagram of transition function)

$$\{p,q,r\}$$
 $\{p,q,r\}$ $\{q,r\}$ $\{p,q,r\}$ $\{p,q,r\}$ $\{p,q,r\}$ $\{p,q,r\}$ $\{p,q,r\}$ $\{p,q,r\}$ $\{p,q,r\}$ $\{p,q,r\}$ $\{q,r\}$ $\{q,r\}$

7. [10 points] Find a grammar equivalent to

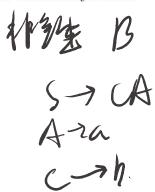
$$S \rightarrow AB \mid CA$$

$$A \rightarrow a$$

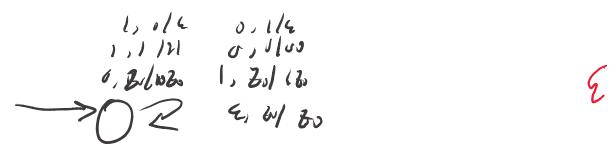
$$B \rightarrow BC \mid AB$$

$$C \rightarrow aB \mid b$$

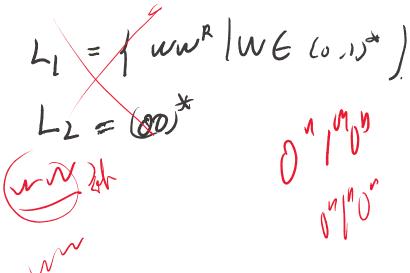
with no useless symbols.



8. [10 points] Design a PDA for $L_{eq} = \{w \in \{0,1\}^* \mid w \text{ contains the same number of 0's and 1's }\}.$



9. [10 points] Prove or disprove: if L_1 is CFL and $L_1 \cup L_2$ is also CFL, then L_2 must be CFL.



10. [10 points] Design Turing machine for the language $\{0^{2n}1^n \mid n \geq 0\}$.

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