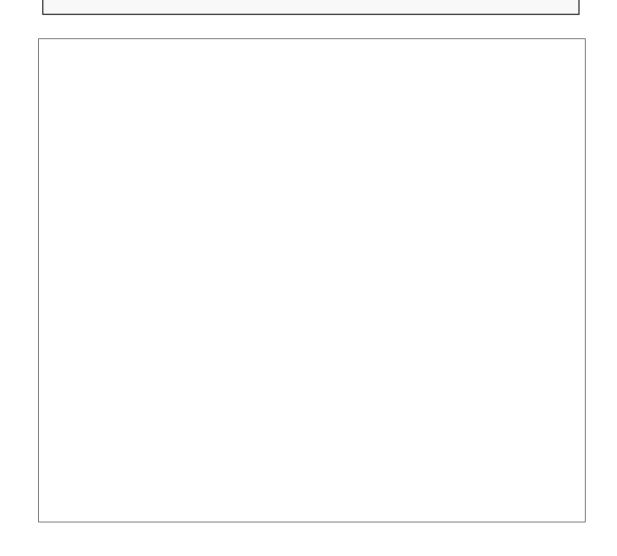


# **Ethernet and Wireless**

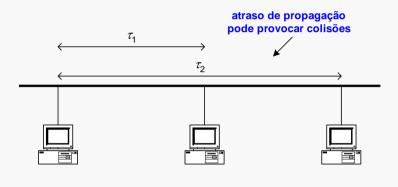
# Redes de Comunicações 1

#### Licenciatura em Engenharia de Computadores e Informática



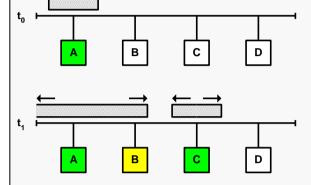
# **CSMA (Carrier Sense Multiple Access)**

- Stations transmit and receive in the same channel
- They sense the medium before transmission; only transmit if medium is free
- Number of collisions is minimized
- Collisions can occur because stations are distanced from each other



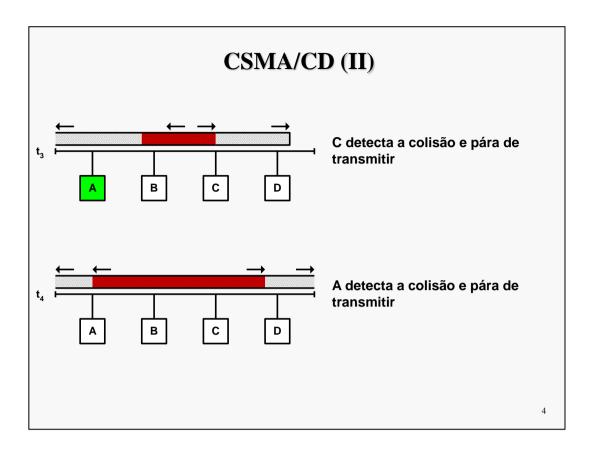
# CSMA/CD (CSMA with Collision Detection) (I)

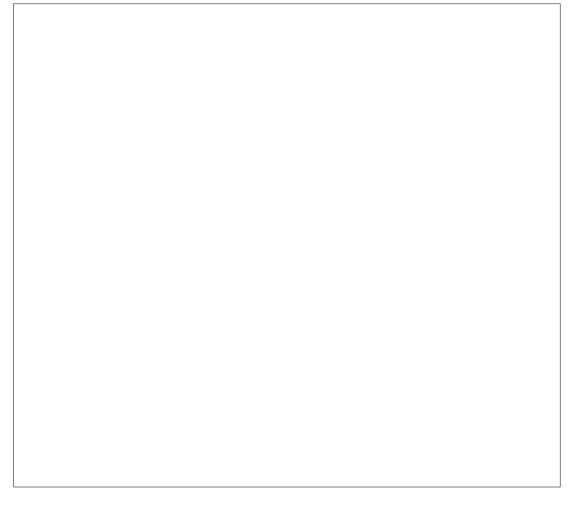
• Stations stop transmitting when they detect collisions

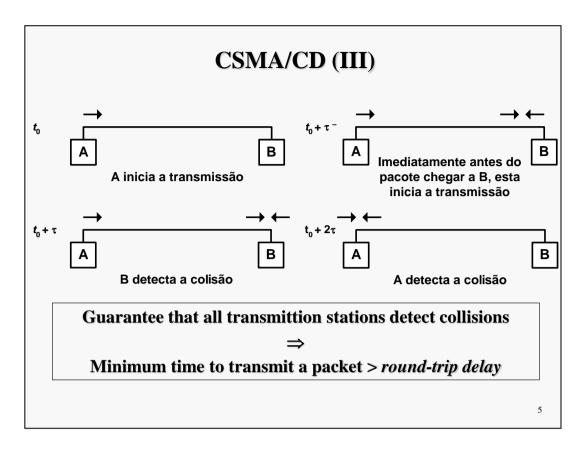


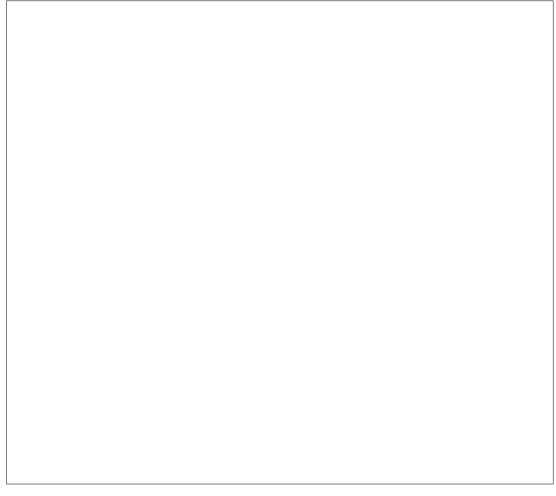
A estação A detecta o meio livre e inicia a sua transmissão

A estação B pretende transmitir mas não o faz porque detecta o meio ocupado; a estação C inicia a transmissão









### CSMA/CD (IV)

- The Carrier Sense Multiple Access with Collision Detection (CSMA/CD) protocol used on Ethernet works as follows: the medium access is ruled by carrier sense (the sending station first detects if the medium is being used by another station) and collision detection (the sending station checks if the medium has the same data being sent by it).
- When a station has an Ethernet frame to be sent, it first checks if the medium is busy with the transmission of a frame by another station. If the medium is free for an Inter Frame Spacing (IFS) time period, it starts sending its frame. If the medium is busy, it waits that the medium becomes free, waits another IFS time period and starts sending its frame (it is said that the protocol is 1-persistent since all stations waiting to transmit during a busy period will transmit their frames with 100% of probability as soon as the medium becomes free for a IFS time period).
- IFS is the minimum time interval required by all stations to accommodate one frame before being prepared to start receiving another frame. For example, in 10 Mbps Ethernet, the IFS is 9.6 μs.
- Note that it is possible that two (or more) stations start transmitting frames almost at the same time originating a collision. In a collision, multiple frames are being simultaneously transmitted and, therefore, will not be correctly received by any station. When a sending station detects a collision, it stops the frame transmission and sends a JAM signal (aimed to guarantee that all stations detect the collision). Then, it waits for a random period of time to send the frame again. This random period is defined by the Truncated Binary Exponential Backoff Algorithm described in the next slide.

#### **Ethernet**

The Carrier Sense Multiple Access with Collision Detection (CSMA/CD) protocol used on Ethernet works as follows: the medium access is ruled by carrier sense (the sending station first detects if the medium is being used by another station) and collision detection (the sending station checks if the medium has the same data being sent by it).

When a station has an Ethernet frame to be sent, it first checks if the medium is busy with the transmission of a frame by another station. If the medium is free for an Inter Frame Spacing (IFS) time period, it starts sending its frame. If the medium is busy, it waits that the medium becomes free, waits another IFS time period and starts sending its frame (it is said that the protocol is 1-persistent since all stations waiting to transmit during a busy period will transmit their frames with 100% of probability as soon as the medium becomes free for a IFS time period).

IFS is the minimum time interval required by all stations to accommodate one frame before being prepared to start receiving another frame. For example, in 10 Mbps Ethernet, the IFS is 9.6  $\mu$ s.

Note that it is possible that two (or more) stations start transmitting frames almost at the same time originating a collision. In a collision, multiple frames are being simultaneously transmitted and, therefore, will not be correctly received by any station. When a sending station detects a collision, it stops the frame transmission and sends a JAM signal (aimed to guarantee that all stations detect the collision). Then, it waits for a random period of time to send the frame again. This random period is defined by the Truncated Binary Exponential Backoff Algorithm described in the next slide.

FR - DETI/UA - 2016/2017

# CSMA/CD (V)

• Number of time slots of delay before the n<sup>th</sup> retry is a random variable uniformly distributed in the interval

$$0 \le r < 2^k$$
, with  $k = \min(n, 10)$ 

- Duration of the slot = 64 bytes = 512 bits =  $51.2 \mu s$  (10 Mbps)
- Example:

```
- n = 1 \Rightarrow r = 0 ou 1 (0 ou 51.2 μs)

- n = 2 \Rightarrow r = 0, 1, 2 ou 3 (0, 51.2, 102.4 ou 153.6 μs)

:
```

- n > 10, maximum delay fixed to  $2^{10}-1 = 1023$  slots
- Maximum number of retries = 16

#### Truncated Binary Exponential Backoff Algorithm

When a collision is detected by a sending station, the time that the station waits to send again its frame is given by an integer number r of time slots. In 10 Mbps and 100 Mbps Ethernet technologies, the time slot is defined by the time taken to transmit a minimum size frame of 64 bytes (64 Bytes =  $64 \times 8 = 512$  bits), which is  $51.2 \,\mu s$  (in 10 Mbps technologies) or  $5.12 \,\mu s$  (in 100 Mbps technologies).

A maximum of 16 retransmissions is allowed beyond which the frame is discarded by the sending station.

In the  $n^{th}$  retransmission of the same frame (with  $n \le 16$ ), the number of waiting time slots r is a uniform random value between  $0 \le r < 2^k$ , where k is the minimum value between n and 10.

Note that the average waiting time is short in the first retransmissions, grows exponentially with the number of retransmissions until the  $10^{th}$  one and remains the same above the  $10^{th}$  retransmission.

FR - DETI/UA - 2016/2017 7

# **CSMA-CD - performance**

Utilization of CSMA/CD is

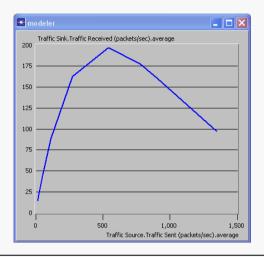
$$S \underset{N \to \infty}{\longrightarrow} \frac{1}{1 + 3.44a}$$

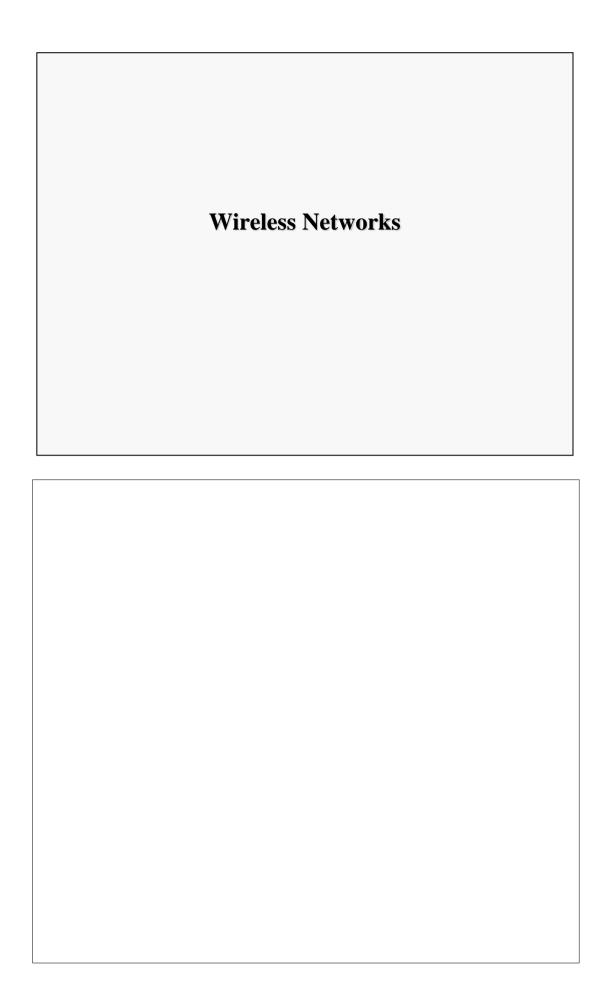
 $a = \tau/T$ , T – transmission time of a packet (useful time)

• a<1

# **CSMA-CD - performance**

- Increase of transmission traffic
  - Increase of collisions





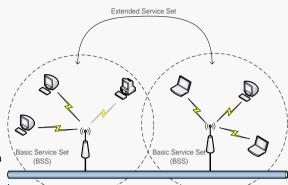
#### **Evolution of WLAN standards**

- WiFi 1 802.11b, 1999, 2.4 GHz band, 11 Mbps data rate
- WiFi 2 802.11a, 1999, 5 GHz band, 54 Mbps data rate
- WiFi 3 802.11g, 2003, 2.4 GHz band, 54 Mbps data rate
- WiFi 4 802.11n, 2009, 2.4 and 5 GHz bands, ~600 Mbps data rate
- WiFi 5 802.11ac, 2013, 5 GHz band, ~1.3 Gbps data rate
- WiFi 6 802.11ax, 2019 (2021 6E), 2.4, 5, (6 GHz 6E) bands, ~9.611 Gbps data rate
- WiFi 7 802.11be, 2024, 2.4, 5, 6 GHz, ~23 Gbps data rate
- WiFi 8 802.11bn, exp. 2028, 2.4, 5, 6 GHz, ~100 GHz data rate



### **Components**

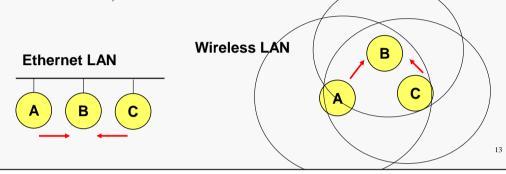
- Station (STA)
  - Mobile terminal
- Access Point (AP)
  - ◆STA connect to access points (infrastructured networks)
- Basic Service Set (BSS)
  - ◆STA and AP with same coverage form a BSS
  - ♦Group of IEEE 802.11 stations associated to an Access Point (AP)
  - Nown through the SSID
- Extended Service Set (ESS)
  - ♦Several BSSs interconnected by APs form
  - a ESS



#### Wired vs Wireless differences

- A and C sense the channel empty simultaneously
  - > Send traffic at the same time
- Ethernet: sender can detect collision
- Wireless: radios cannot detect collision (work in half-duplex)
  - Full-duplex: both can transmit and receive information between each other simultaneously

Half-duplex: transmission and reception of information must happen alternatively.
 While one point is transmitting, the other must only receive (avoid self-interference)



### **Wireless MAC**

- Wired MACs
  - Typical: CSMA/CD
  - Medium is free → send
  - Listen to sense collision
- What about wireless?
  - Signal power reduces with the square distance
  - Sender can apply CS and CD, but collisions occur in the receiver!
  - ➤ Sender may not listen the collision (CD does not work)
  - CS may not work either with hidden nodes

14

CS – carrier sense

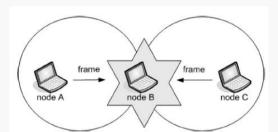
### **Hidden nodes**

#### Hidden terminals

- A and C do not ear each other
- Collision in B, if A and C send at the same time
- Nor A nor C understand that collision occured

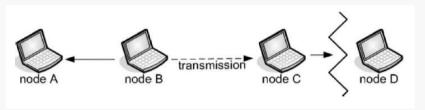
#### • Solution

- Detect collisions in the receiver
- "virtual carrier sensing": sender asks the receiver if he is receiving traffic;
   in the case of absense of answer, he assumes that the channel is busy



### **Exposed nodes**

- Exposed terminals
  - B sends to A; C wants to send to D
  - C senses the network and discovers that the medium is occupied
  - D is not in the range of B and A is not in the range of C, so the traffic could be transmitted
  - A and D are exposed terminals
- The transmissions could be done in parallel with no collision

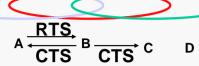


# MACA: Multiple Access with Collision Avoidance

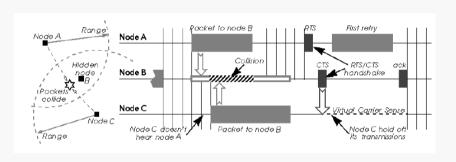
- MACA: avoids collisions using signalling packets
  - RTS (request to send)
    - A small packet is sent before transmitting
  - CTS (clear to send)
    - Receiver provides the right to transmit, when it is able to receive
- Signaling packets (RTS/CTS) contain
  - Sender address
  - Receiver address
  - Packet length (to be transmitted)
- Used in networks scenario with a large amount of traffic/collisions

### **MACA: Hidden Nodes**

- MACA and hidden nodes
  - A, C  $\rightarrow$  B (?)
  - A  $--RTS \rightarrow B$
  - B  $\rightarrow$  CTS  $\rightarrow$  A
  - C ears CTS of B



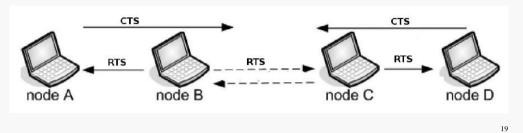
• C waits for the period announced in A transmission



# **MACA: Exposed Nodes**

- MACA and exposed nodes
  - B  $\rightarrow$  A, C  $\rightarrow$  D(?)
  - BRTS  $\rightarrow$  A
  - A CTS  $\rightarrow$  B
  - C ears RTS of B
  - C does not ear CTS of A
  - $CRTS \rightarrow D$

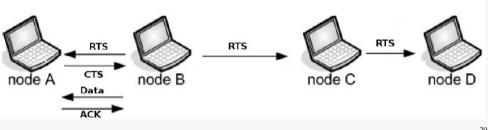


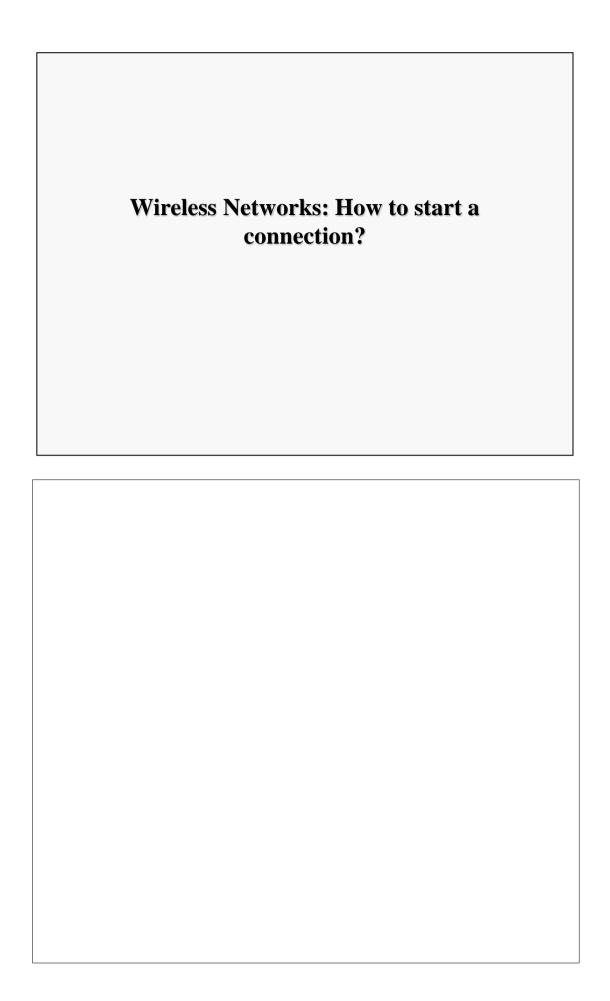


.

# **MAC** reliability

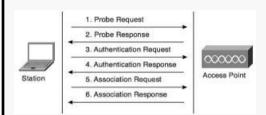
- Wireless connections are very prone to errors
  - Transport is not reliable
- Solution: use acknowledgements
  - When A receives DATA from B, answers with ACK.
  - If B does not receive ACK, B retransmits
  - C and D will not transmit until the ACK (to avoid collisions)
  - Total expected duration (including ACK) is included in the RTS/CTS packets

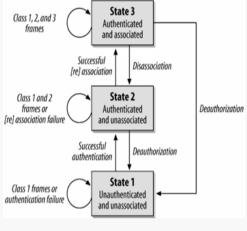




# Joining a BSS

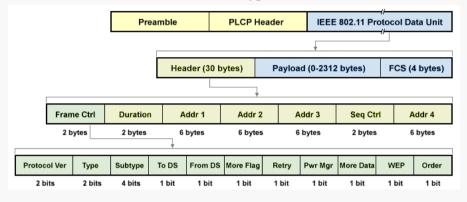
- Station finds BSS/AP by Scanning/Probing.
- BSS with AP: both Authentication and Association are necessary for joining a BSS.





### **WLAN Frames**

- Three types of frames
  - ◆Control: RTS, CTS, ACK
  - ◆Management
  - ◆Data
- Header is different for the different types of frames.



### **Joining BSS with AP: Scanning**

- A station willing to join a BSS must get in contact with the AP. This can happen through:
- 1. Passive scanning
  - ◆ The station scans the channels for a Beacon frame that is sent periodically from an AP to announce its presence and provide the SSID, and other parameters for WNICs within range
- 2. Active scanning (the station tries to find an AP)
  - ◆ The station sends a Probe Request frame Sent from a station when it requires information from another station
  - All AP's within reach reply with a Probe Response frame Sent from an AP containing capability information, supported data rates, etc., after receiving a probe request frame

#### **Beacon Frame**

### **Probe Request/Response Frames**

# Joining BSS with AP: Authentication

- Once an AP is found/selected, a station goes through authentication
- Open system authentication (default, 2-step process)
  - Station sends authentication frame with its identity
  - AP sends frame as an Ack / NAck
- Shared key authentication
  - Stations receive shared secret key through secure channel independent of 802.11
  - •After the WNIC sends its initial authentication request, it will receive an authentication frame from the AP containing a challenge text
  - The WNIC sends an authentication frame containing the encrypted version of the challenge text to the AP.
  - The AP ensures the text was encrypted with the correct key by decrypting it with its own key.
  - The result of this process determines the WNIC's authentication status.

#### **Authentication Frames**

- Nowadays, WPA\* secure networks use "Open System".
- Non-"Open System" authentication was used for WEP protected networks (unsecured and functionally deprecated).

```
TEEE 802.11 Authentication, Flags: ......
Type/Subtype: Authentication (0x000b)
Frame Control Field: 0xb000
.000 0001 0011 1010 = Duration: 314 microseconds
                                                                                                                       ← From Station
   Receiver address: Cisco_61:ee:d0 (00:1c:f6:61:ee:d0)
Destination address: Cisco_61:ee:d0 (00:1c:f6:61:ee:d0)
   Transmitter address: D-LinkIn 6a:cc:6e (84:c9:b2:6a:cc:6e)
   Source address: D-LinkIn_6a:cc:6e (84:c9:b2:6a:cc:6e)
- IEEE 802.11 Authentication, Flags: .......C
Type/Subtype: Authentication (0x000b)
Frame Control Field: 0xb000
.000 0001 0011 1010 = Duration: 314 microseconds
Receiver address: D-LinkIn_6a:cc:6e (84:c9:b2:6a:cc:6e)
   Fixed parameters (6 bytes)
                                                                                                Destination address: D-LinkIn_6a:cc:6e (84:c9:b2:6a:cc:6e)
Transmitter address: Cisco_61:ee:d0 (00:1c:f6:61:ee:d0)
Source address: Cisco_61:ee:d0 (00:1c:f6:61:ee:d0)
      Authentication Algorithm: Open System (0)
      Authentication SEQ: 0x0001
      Status code: Successful (0x0000)
                                                                                                BSS Id: Cisco 61:ee:d0 (00:lc:f6:61:ee:d0)
.................. 0000 = Fragment number: 0
1010 1001 0000 .... = Sequence number: 2704
Frame check sequence: 0x9f8350el [unverified]
[FCS Status: Unverified]
                                                               From AP \rightarrow
                                                                                             IEEE 802.11 wireless LAN
- Fixed parameters (6 bytes)
                                                                                                   Authentication Algorithm: Open System (0)
                                                                                                   Authentication SEQ: 0x0002
Status code: Successful (0x0000)
```

### Joining BSS with AP: Association

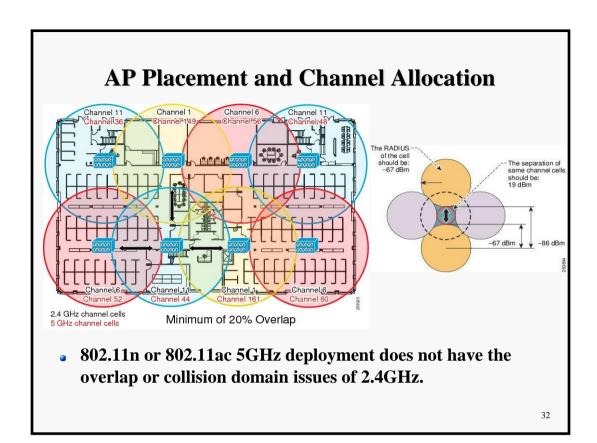
- Once a station is authenticated, it starts the association process, i.e., information exchange about the AP/station capabilities and roaming
- STA → AP: Associate Request frame
  - Enables the AP to allocate resources and synchronize. The frame carries information about the WNIC, including supported data rates and the SSID of the network the station wishes to associate with.
- **→** AP → STA: Association Response frame
  - Acceptance or rejection to an association request. If it is an acceptance, the frame will contain information such as association ID and supported data rates.
- New AP informs old AP (if it is a handover).
- Only after association is completed, a station can transmit and receive data frames.

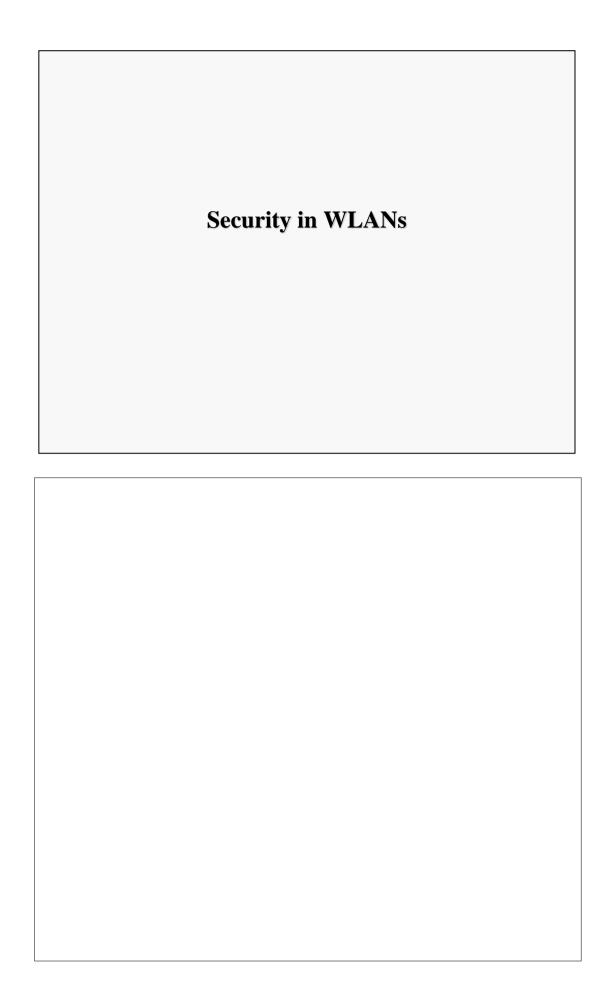
# **Association Request/Response Frames**

#### **Data Frame**

```
IEEE 802.11 QoS Data, Flags: .p....TC
Type/Subtype: QoS Data (0x0028)
Frame Control Field: 0x8841
  .000 0001 0011 1010 = Duration: 314 microseconds
                                                                 ← Node that will receive frame (AP)
  Receiver address: Cisco 61:ee:d1 (00:1c:f6:61:ee:d1)
 Transmitter address: IntelCor_e8:14:53 (b8:8a:60:e8:14:53) \leftarrow Node that sends frame
 Destination address: D-LinkIn_6a:cc:6e (84:c9:b2:6a:cc:6e) ← Station to receive data
  Source address: IntelCor e8:14:53 (b8:8a:60:e8:14:53)
                                                                 ← Station who sent data
 BSS Id: Cisco 61:ee:d1 (00:1c:f6:61:ee:d1)
 STA address: IntelCor_e8:14:53 (b8:8a:60:e8:14:53)
             .... 0000 = Fragment number: 0
  0000 0000 0011 .... = Sequence number: 3
  Frame check sequence: 0xc72771e8 [unverified]
  [FCS Status: Unverified]
Qos Control: 0x0000
CCMP parameters
Data (1244 bytes)
 Data: f8002648417037bc923106ead1717d4821fde0989beb08b1...
  [Length: 1244]
```

- Station "IntelCor\*" sending data to station "D-LinkIn\*" (via AP).
- Frame captured between station "IntelCor\*" and AP ("Cisco\*i").





### **Authentication and authorization mechanisms**

- Changing according to the organization and the security level
  - ◆Open network
  - ◆Open network + MAC authentication
  - →Open network + VPN-gateway
  - ◆Open network + web-gateway
  - **♦**SSID
  - ◆Shared key: WEP
  - ◆Wi-Fi Protected Access (WPA)
  - ▶IEEE 802.11i (WPA2)
  - **▶**IEEE 802.1X
  - ◆Virtual Private Networks (VPNs)

### **Open Network(s)**

#### Open network

- Network is open, providing IP addresses with DHCP
- There is no authentication and access is free
- Does not require specific software
- Access control is complicated
- →It is possible to 'see' all traffic in the network (sniffing)

#### Open network + MAC authentication

- ◆The control of the station MAC address is added
- Larger management load
- .... But MAC addresses can be falsified
- ... Difficult to support guests
- .... Impossible to use in public environments

#### **WEP Protocol**

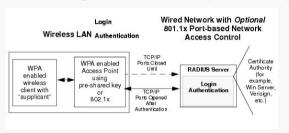
- Wired Equivalent Privacy  $\rightarrow$  shared key scheme.
- Part of basic 802.11 standard.
- Security protocol at link layer (L2).
- Designed to be computationally efficient and selfsynchronized.
- The station has to know the key (like a password) to access the AP.
- With passive monitoring, it can be broken (in seconds)
  - →Header is not ciphered, all destinations and origins are visible.
  - •Control frames are not ciphered, and then they can be changed.
  - AP is not authenticated and can be falsified.

### **WPA and 802.11i (WPA2)**

- IEEE 802.11i IEEE 802.11 task group "MAC enhancement for wireless security".
- Wi-Fi Protected Access (WiFi Alliance), WPA, is a subset internal in 802.11i.
  - Used in Eduroam
  - Defined to work in actual equipment.
    - Firmware update only.
  - Pass-phrase constant and shared, but keys are generated per session.
  - Used in the AP and station.
  - Uses "Open System" during authentication phase.
- WPA has two distinct components.
  - Authentication, based on 802.1X.
    - Login and password it is the user that is authenticated anywhere in the world
    - Does not need to know any password of the AP
    - Contact to a local server and a remote server where the user belongs
  - Ciphering based on TKIP (Temporal Key Integrity Protocol).

### **IEEE 802.1X**

- Layer 2 solution between station and AP.
  - Available in many equipments (e.g. IEEE 802.xx).
  - •Web systems frequently use 802.1X.
- Several authentication-mechanisms available (EAP-MD5, EAP-TLS, EAP-TTLS, PEAP)
- Multiple standard ciphering algorithms.
- Can cipher data with dynamic keys.





### **WPA\*** Key Exchange

Done during the Association process.

After Association Request/response frames.

```
205 595.669409767 IntelCor e8:14:53
206 595.671214291 Cisco 61:ee:d1
IntelCor e8:14:53
206 595.671214291 Cisco 61:ee:d1
IntelCor e8:14:53
208 595.676323781 (sico 61:ee:d1
IntelCor e8:14:53
208 595.6763333124 IntelCor e8:14:53
209 595.681295313 Cisco 61:ee:d1
IntelCor e8:14:53
218 595.683699439 IntelCor e8:14:53
218 595.683699439 IntelCor e8:14:53
218 595.683699439 IntelCor e8:14:53
218 77:211 bytes on wire (1688 bits), 211 bytes captured (1688 bits) on interface θ
Radiotap Header v0, Length 56
802.11 radio information
IEEE 802.11 005 Data, Flags: .....F.
Type/Subtype: QoS bat (axeo26)
Frame Control Field: 0x8082

Readiotap Header v0, Length 56
802.11 radio information
IEEE 802.11 0x8082
Frame Control Field: 0x8082
Frame Control Exception Control Excep
```

39