

Datapoint 2200

PROGRAMMERS' MANUAL

A product of COMPUTER TERMINAL CORPORATION
9725 Datapoint Drive
San Antonio, Texas 78229

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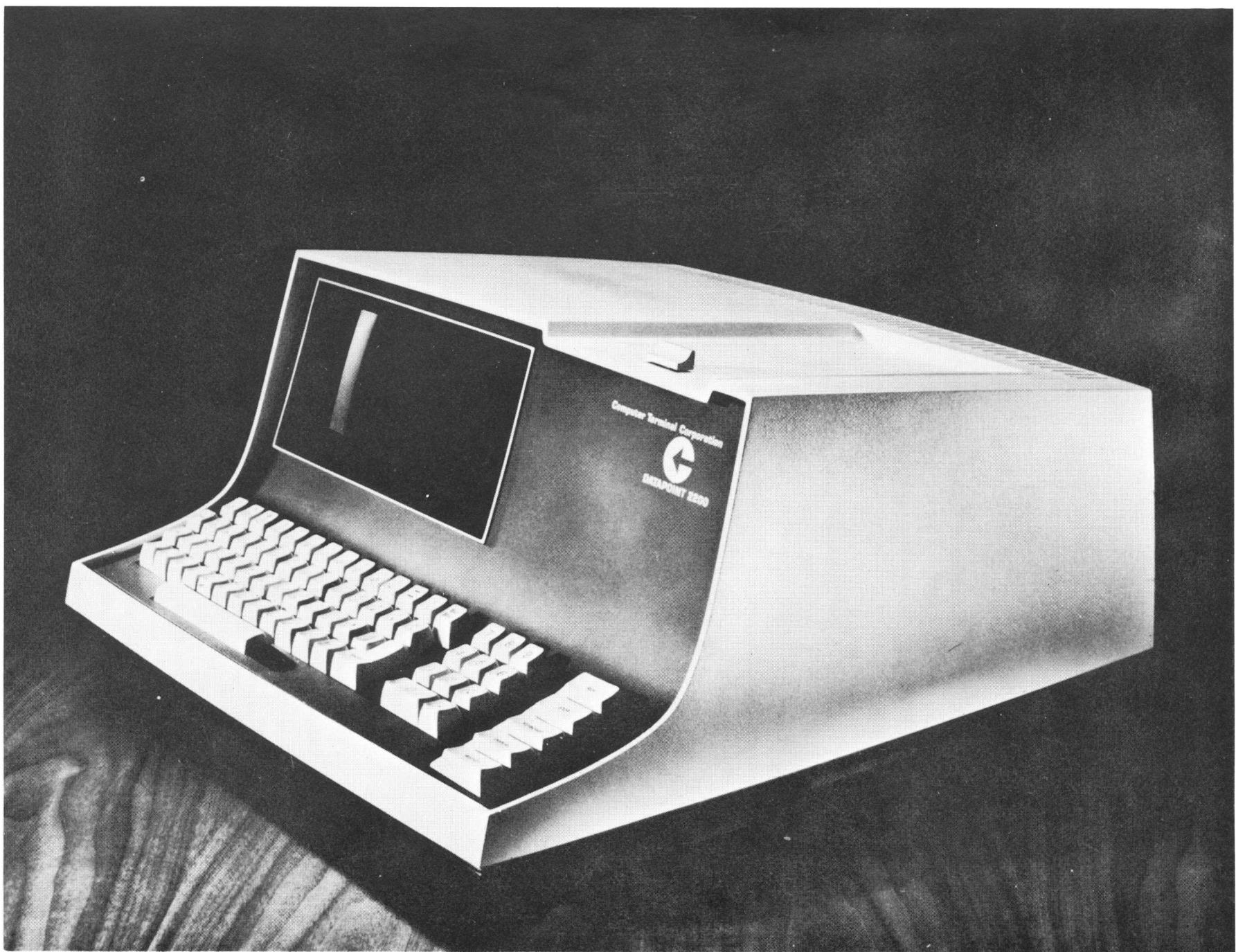
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SECTION 1

DATAPOINT 2200

REFERENCE MANUAL

PART 1

GENERAL FEATURES

1.1 INTRODUCTION

The Computer Terminal Corporation Datapoint 2200 is an integrated data system consisting of an alphanumeric keyboard for data entry, a cathode ray screen for data display, two digital cassette recorders for bulk data storage, a general purpose digital computer for control, and a communications capability for interface with external devices and communications facilities.

Through programming of the control computer the Datapoint 2200 may be used for an infinite variety of data processing applications.

The achievement of a small computer with integrated keyboard, display, storage and communications at such low cost now makes possible computer sophistication for applications not previously practical - particularly in the computer terminal/data entry/communications area.

This manual describes the specific hardware details of the Datapoint 2200. For information regarding specific applications the Datapoint 2200 Programmers' Guide and specific application manuals should be referred to.

1.2 SYSTEM ELEMENTS

There are four basic system elements in the basic Datapoint 2200 plus the capability of interface to a number of external peripheral devices.

This manual covers the basic elements (c.r.t., keyboard, processor, cassette tape decks) and one external device (communications adaptor).

1.3 CRT DISPLAY

The Datapoint 2200 CRT Display provides the following features:

- a. 7" x 2-1/2" viewing area;
- b. 960 characters;
- c. 80 character by 12 line format;
- d. 4/32" x 3/32" character size;
- e. Entire 94 character ASCII set;

- f. 60 frame per second refresh rate;
- g. 5 x 7 matrix character generation;
- h. 5 x 7 solid, blinking cursor, alternates with character, nondestructive;
- i. P31 green phosphor;
- j. Single control line erasure, frame erasure, and page roll-up; and
- k. Direct control of all c.r.t. functions by the 2200 processor, providing tab, editing, form control, etc.

1.4 KEYBOARD

The integral keyboard provides a basic 41 key alphanumeric key group, an 11 key numeric group and five system control keys.

The keyboard provides a unique multi-key roll-over characteristic providing maximum ease of typing. Transfer of characters from the keyboard is under control of the 2200 processor. An audible click providing an acoustical feedback to the typist is available under processor control.

A programmable audio "beep" is also provided when it is desired to gain a typist's attention.

1.5 PROCESSOR

The integral processor provides all control functions and includes:

- a. 28 different instruction types;
- b. 7 addressable registers;
- c. 7 deep pushdown stack;
- d. 8 bit memory word length;
- e. Up to 8192 word memory;
- f. Complete parallel I/O system;
- g. Automatic power-up restart.

1.6 CASSETTE TAPE DECKS

Two read-write tape decks are provided for program and data storage. The deck accepts Norelco-type cassettes and provides:

- a. 47 characters per inch density;
- b. Dual-capstan forward-reverse operation;
- c. Processor controlled data transfer, direction control, and high-speed rewind.

1.7 COMMUNICATIONS ADAPTOR

The communications adaptor is a unique feature of the Data-point 2200 system. There are four versions of the adaptor:

- a. EIA RS-232 interface for use with external data sets or peripherals;
- b. High-level keying interface for connection to telegraph-type communications channels or equipment;
- c. 103-type data set interface for direct connection to common carrier lines, and including automatic dialing and answering;
- d. 202-type data set interface with 150 bit/sec supervisory channel operation for direct connection to common carrier lines, and including automatic dialing and answering.

The adaptor permits program selection of the desired bit rate, character length, and character set providing the most versatile communications capability yet provided for a remote terminal.

1.8 GENERAL SPECIFICATIONS

The Datapoint 2200 has the following general characteristics:

- a. 105-135 v.a.c., 60 cycle, 180 watts, power input;
- b. 47 pounds weight;
- c. 9-5/8" high, 18-1/2" wide, by 19-5/8" deep outside dimensions;
- d. 0° to 50°C (32° to 122°F), 10 to 90 percent relative humidity operation environment.

1.9 OPTIONAL PERIPHERALS

A number of optional peripherals are available (in addition to the communications adaptor) for use with the Datapoint 2200 including a:

a. 132 column, 30 c.p.s. impact page printer; and a

b. IBM compatible magnetic tape deck.

For further information on these devices reference should be made to their respective reference manuals.

PART 2

BASIC PROCESSOR

2.1 PROCESSOR ORGANIZATION

The processor contained in the Datapoint 2200 is comprised of the Arithmetic/Logic Unit, 7 program accessible registers, 2K to 8K words of read/write memory, an instruction decoder and a seven-level hardware pushdown stack used in subroutine type operations.

2.2 ARITHMETIC/LOGIC UNIT

The Arithmetic/Logic Unit is capable of processing both binary integers and logical operands. All arithmetic and logical operations may take place between the A-register and any of the 7 program accessible registers (or between the A-register and memory). The A-register always contains the result of an arithmetic or logic operation, with the other register (or memory cell) being unaffected. Arithmetic and logic operations affect the Sign, Carry, Zero and Parity Flip-flops.

2.3 PROCESSOR REGISTERS

A - The Accumulator register is used to hold the result of all arithmetic and logical instructions. All data transfers into or out of the computer take place through this register.

7	6	5	4	3	2	1	0
---	---	---	---	---	---	---	---

A Register

B, C, D, E - These are general purpose registers which may be used in conjunction with the Accumulator in arithmetic and logical operations. Each register may be loaded from or stored into memory or another register. When used in conjunction with the A and H, L registers, the B, C, D and E registers may function as indexes.

7	6	5	4	3	2	1	0
---	---	---	---	---	---	---	---

B, C, D, E Registers

H, L - The H and L Registers are utilized to contain respectively the most significant portion (MSP) and least significant portion (LSP) of the address of a memory location being referenced. All memory reference instructions utilize these registers with the exception of CALL and JUMP commands. However, the H and L Registers may be used as general purpose registers when not being used as above.

H (8 bits)	L (8 bits)
MSP of address being referenced	LSP of address being referenced

P - The program register or "Location Counter" contains the address of the next instruction to be executed. This

register is stored in the pushdown stack upon the execution of a "CALL" instruction and is loaded with the effective address upon execution of a "JUMP", "CALL" or "RETURN" instruction. The P register is 13 bits in length and is capable of addressing up to 8K of memory.

12	P (13 bits)	0
----	-------------	---

I - The I register is the register which holds the "operation code" of the instruction currently being executed. The contents of I are gated through a decoding network to determine what operation, internal or external, is to be performed.

2.4 MEMORY

The basic Datapoint 2200 is supplied with 2048 eight-bit words of memory. Additional modules of 2048 words each may be incorporated with the total memory capacity of the processor being 8192 words. Each 2K memory is made up of 32 individual MOS shift registers with each one having a capacity of 512 bits or 64 eight-bit words. These registers are clocked at a rate of 1.2 MHz. Data is read out in bit serial fashion with one word taking 8 microseconds. During this period of time, two clock pulse times are available for the processor to perform any necessary gating or testing functions.

The Datapoint 2200 memory might be likened to a drum type memory in some respects. It takes approximately 1/2 millisecond for the memory to completely circulate. Thus, if the current instruction referenced a memory location for data access, there would be a 1/2 millisecond delay before that instruction could be completed. However, unlike a drum memory the MOS memory may be stopped during instruction execution so that each succeeding instruction may be read from memory without delay (in 8 usec.).

Physically, instructions require a variable number of cycles for completion. In the first cycle, the instruction is fetched from memory and decoded. If the instruction involves no memory reference, it is then executed within 8 microseconds for a total completion time of 16 microseconds.

"Immediate" type instructions are the same as instructions requiring no memory reference and require a 16 usec interval for the operand fetch and execute cycle. Jump and Call type instructions require a variable amount of time for execution, depending on the difference between the old and new locations.

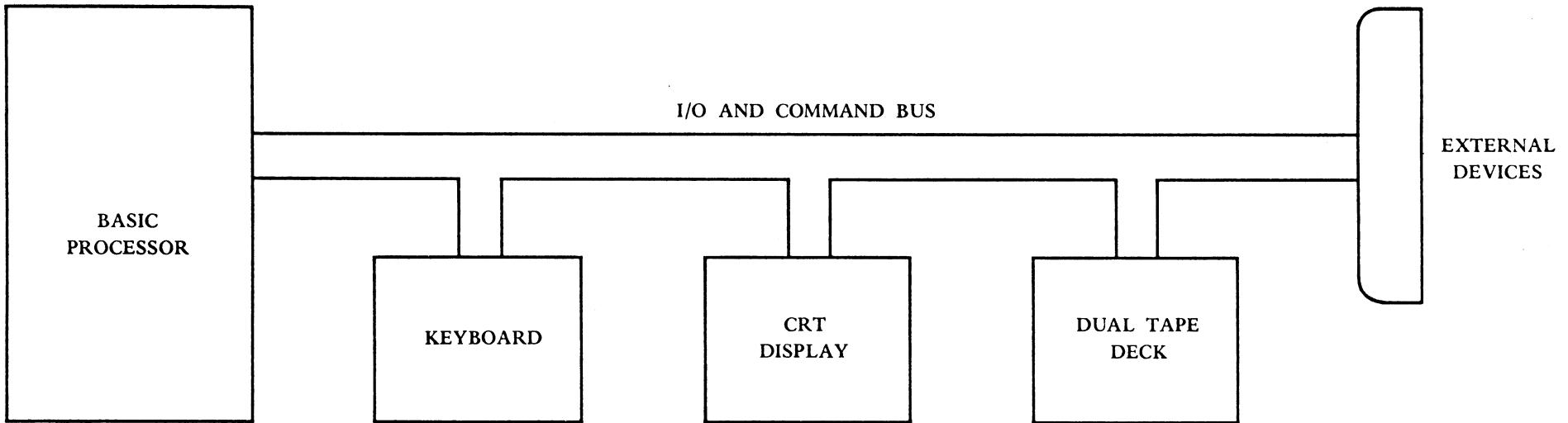
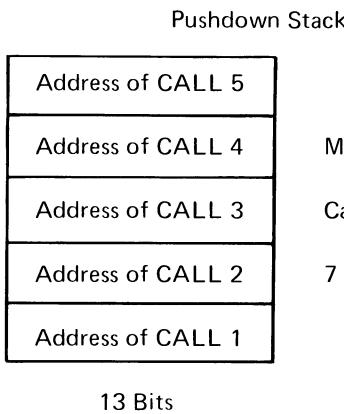


FIGURE 2-1

**DATAPOINT 2200
BLOCK DIAGRAM**

2.5 PUSHDOWN STACK

A unique feature of a machine this size is the incorporation into the processor's structure of a pushdown stack which is useful in any type of application which requires program subroutines. The stack automatically stores the contents of the P register upon execution of a "CALL" instruction and automatically restores P upon execution of a "RETURN". The stack is a group of "last-in/first-out" registers and has a capacity of 7 CALLS. Note that "CALLS" may be "nested", that is more than one CALL may be made before the execution of a RETURN. The execution of a "RETURN" will cause processor control to be given to the next instruction following the last executed CALL.



2.6 CONTROL FLIP-FLOPS

Also contained in the basic processor are four control flip-flops which reflect the state of the arithmetic logic unit and which may be tested through the execution of a conditional jump, call or return instruction. The flip-flop mnemonics with their associated functions are as follows:

C_f-Carry Flip-flop. Set when an arithmetic operation results in either a carry (add) or borrow (subtract). * The Carry Flip-flop also reflects the state of the most significant bit in the accumulator after completion of a shift right instruction. Likewise, it reflects the state of the accumulator least significant bit after completion of a shift left instruction.

Z_f-Zero Flip-flop. Set when the result of an arithmetic or logical operation is equal to zero.*

S_f-Sign Flip-flop. This flip-flop reflects the state of bit 7 in the accumulator after an arithmetic type operation.*

P_f-Parity Flip-flop. Indicates the parity or "number of one bits" contained in the accumulator. If this flip-flop is set (true), the A register contains an odd number of one bits; if it is reset (false), the A register contains an even number of one bits.*

*In the event of a compare instruction the contents of the accumulator are not changed; however, the control flip-flops reflect the equivalent of a subtract instruction.

2.7 DATA FORMAT

Data is represented in the Datapoint 2200 in the form of 8-bit binary integers.

7	6	5	4	3	2	1	0
---	---	---	---	---	---	---	---

DATA WORD

2.8 INSTRUCTION FORMATS (GENERAL)

Instruction formats, dependent upon the operation to be performed, may be eight, sixteen or twenty-four bits in length.

Type-1- register to register, memory reference, arithmetic, logical, shift instructions

OP CODE 8 bits

Type-2- immediate mode instructions

OP CODE 8 bits	OPERAND 8 bits
-------------------	-------------------

Type-3- JUMP & CALL instructions

OP CODE 8 bits	ADDRESS 16 bits
-------------------	--------------------

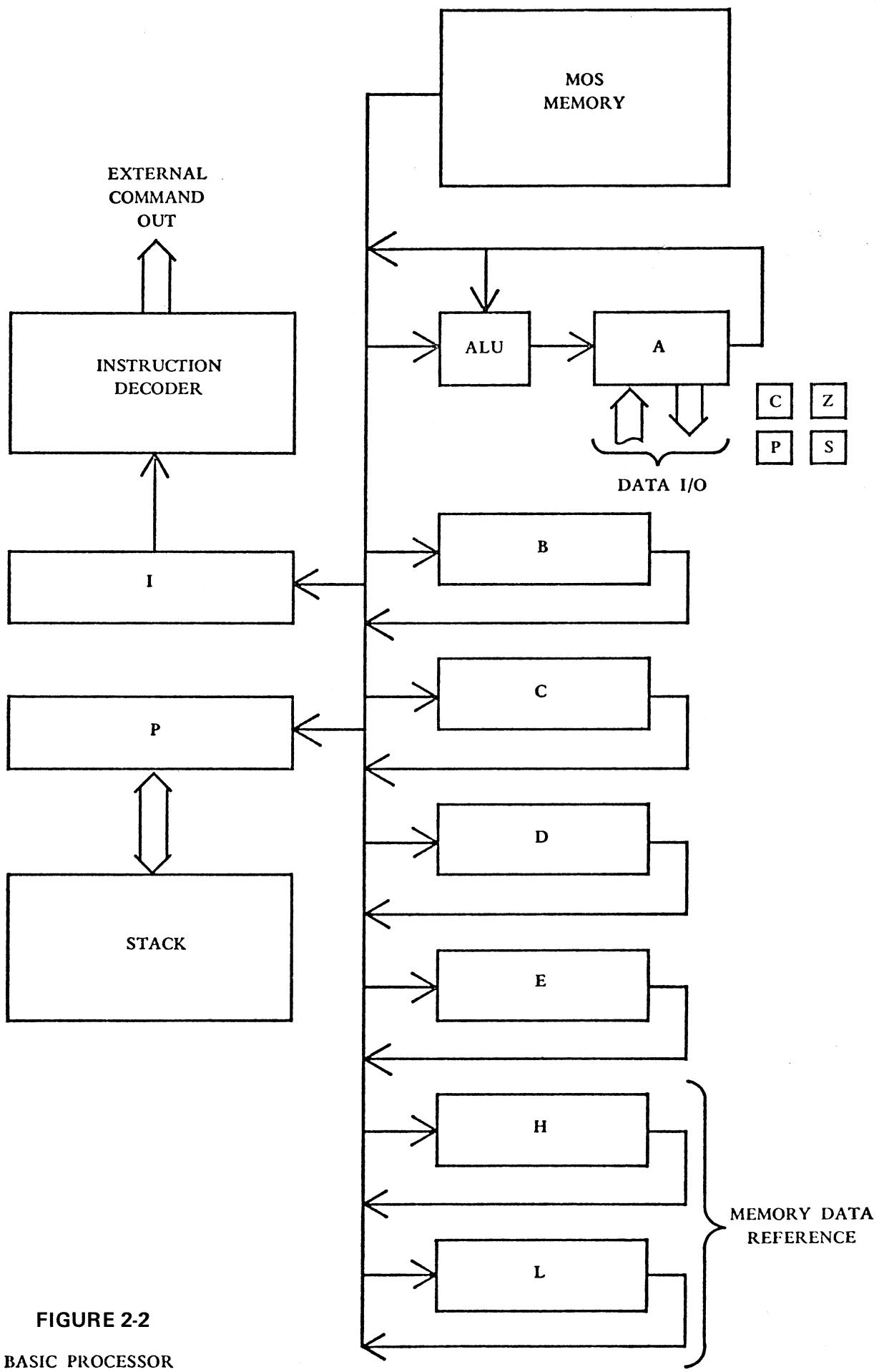


FIGURE 2-2
BASIC PROCESSOR

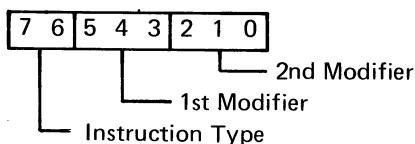
PART 3

INSTRUCTION REPERTOIRE

3.1 PRESENTATION FORMAT

This section gives a detailed description of each of the Data-point 2200 instructions. The use and operations of each instruction is presented as follows:

- FUNCTION:** Mnemonic Code
OPERATION: Symbolic representation of instruction description.
OP CODE: Operation Code, expressed in octal.
TIMING: Execution time. (Times are approximate).
DESCRIPTION: Definition of function of the instruction.



INSTRUCTION FORMAT: Explanation of the function of each part of the instruction word.

NOTE

Considerations in instruction use and further definition of function.

Symbols and Abbreviations

The following symbols and abbreviations are used in the instruction format:

()	the contents of
←	is replaced by
→	is transferred to
:	is compared with
A	8 bit arithmetic register (accumulator)
B	
C	8 bit general purpose registers
D	
E	
H	8 bit register used to specify most significant portion of operand address
L	8 bit register used to specify least significant portion of operand address
M	memory location designated by contents of H, L
r	one of the following register designators: A, B, C, D, E, H, L
r _s	designates operand source register (s=0-7)
r _d	designates operand destination register (d=0-7)
V	Logical "OR" operation
⊕	Logical "exclusive-OR" operation

Λ	Logical "AND" operation
STACK	Instruction counter (P) pushdown queue
P	Program counter
f _c	Flag flip-flop codes: C _f , Z _f , S _f , P _f
RR	Register to Register
IM	Immediate (from P+1)
MR	Memory Reference (Contents of memory location designated by H, L)
I	Instruction Register

TABLE 3-1

SOURCE AND DESTINATION CODES (s and d)

	s/d	SYMBOLIC CODE	
r _s /r _d	0	A	A Register
	1	B	B Register
	2	C	C Register
	3	D	D Register
	4	E	E Register
	5	H	H Register
	6	L	L Register
M	7	M	Memory location specified by contents of H&L

TABLE 3-2

FLIP-FLOP CODE (f_c)

c	SYMBOLIC CODE	NAME
0	C _f	Carry
1	Z _f	Zero
2	S _f	Sign
3	P _f	Parity

LOAD IMMEDIATE: Lr_d
 OP CODE: 0d6 TIMING: 16 usec.
 OPERATION: $(P+1) \rightarrow r_d, P+2 \rightarrow P$
 DESCRIPTION: Transfers the contents of the memory location immediately following the instruction, to the register specified by bits 3-5(d) of the instruction.

INSTRUCTION FORMAT:

P				P+1	
7	6	5	4	3	2 1 0
0		d		6	OPERAND

d: is the destination designator
 d=7: is not allowed

NOTE

1. The contents of P+1 are unchanged.
2. None of the Flag Flip-flops are affected.
3. Refer to Table 3-1 for destination codes.

LOAD: Lr_dM, Lr_drs, LMr_s
 OP CODE: 3ds TIMING: 16 usec. for register to register transfers, 520 usec. for memory reference.

OPERATION: $(M) \rightarrow r_d, s=7, d \leq 6 (Lr_dM)$
 $(r_s) \rightarrow r_d, s \leq 6, d \leq 6 (Lr_drs)$
 $(r_s) \rightarrow M, s \leq 6, d=7 (LMr_s)$

DESCRIPTION: Transfers the operand from the source specified by bits 0-2 of the instruction word to the destination specified by bits 3-5 of the instruction word.

INSTRUCTION FORMAT:

P					
7	6	5	4	3	2 1 0
3		d		s	

d: designates the destination of data.
 s: designates the source. If either s or d=7 a memory reference is indicated and the contents of registers H&L specify the address of the memory location.

NOTE

1. The data source is unaffected.
2. s & d both = 7 results in a Halt instruction.
3. None of the Flag Flip-flops are affected by execution of this instruction.
4. s=d results in a NOP, except as stated in Note 2.

ADD IMMEDIATE: AD

OPERATION: $(A) + (P+1) \rightarrow A \ P+2 \rightarrow P$

DESCRIPTION: Adds to the contents of the A register the contents of the memory location immediately following the instruction, and retains the sum in the A register. Sets the Cf flip-flop if ADD overflow occurs, otherwise resets Cf.

INSTRUCTION FORMAT:

P						P+1	
7	6	5	4	3	2	1	0
0	0		4			OPERAND	0

NOTE

1. The Sign, Zero and Parity Flip-flops will indicate the status of the A register at completion.
 2. The contents of P+1 are unchanged.
 3. The Carry Flip-flop is cleared at the beginning of this instruction.

ADD:

OP CODE: 20s

ADr_s ADM

TIMING: 16 usec. if RR, 520
usec. if MR

OPERATION: $(A) + (r_s) \rightarrow A$ or $(A) + (M) \rightarrow A$

DESCRIPTION: This instruction is identical to ADD IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

P						
7	6	5	4	3	2	1
2		0				s

s: specifies the operand source.
Refer to Table 3-1 for source codes.

ADD WITH CARRY

OP CODE: 014 TIMING: 16 usec

OPERATION: $(A) + (P+1) + (C_f) \rightarrow A$, $P+2 \rightarrow P$
DESCRIPTION: Adds the C_f bit and the contents of the location immediately following the instruction to the contents of the A register, and retains the sum in the A register. If add overflow occurs, the C_f Flip-flop is set, otherwise C_f is reset.

INSTRUCTION FORMAT:

P					P+1	
7	6	5	4	3	2	1 0 7
0		1		4		OPERAND 0

NOTE

1. The Sign, Zero and Parity Flip-flops will indicate the status of the A register at completion.
 2. The contents of P+1 remain unchanged.

ADD WITH CARRY: ACr_s ACM

OP CODE: 21s

TIMING: 16

520 usec. if MR

DESCRIPTION: This instruction is identical to ADD WITH CARRY IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

	P	.
7 6	5 4 3	2 1 0
2	1	s

s: specifies the operand source.
Refer to Table 3-1 for source codes.

SUBTRACT**IMMEDIATE: SU**

OP CODE: 024

TIMING: 16 usec.

OPERATION: $(A) - (P+1) \rightarrow A, P+2 \rightarrow P$ DESCRIPTION: Subtracts the contents of the memory location immediately following the instruction from the contents of the A register, and retains the difference in the A register. The C_f Flip-flop is set if underflow occurs.**INSTRUCTION FORMAT:**

P				P+1	
7	6	5	4	3	2 1 0 7 0
0	2	4			OPERAND

NOTE

1. The contents of P+1 is unchanged.
2. The Zero, Sign, and Parity Flip-flops represent the status of the A register at the completion of this instruction.

SUBTRACT:**OP CODE: 22s****SUr_s SUM**

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: $(A) - (r_s) \rightarrow A$ or $(A) - (M) \rightarrow A$

DESCRIPTION: This instruction is identical to SUBTRACT IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

P			
7	6	5	4 3 2 1 0
2	2		s

s: specifies the operand source.
Refer to Table 3-1 for source codes.

SUBTRACT WITH**BORROW IMMEDIATE:****SB**

OP CODE: 034

TIMING: 16 usec.

OPERATION: $(A) - (P+1) - (C_f) \rightarrow A, P+2 \rightarrow P$ DESCRIPTION: Subtracts the contents of the memory location immediately following the instruction and the C_f bit, from the contents of the A register. Sets the C_f bit if underflow occurs, otherwise resets C_f .**INSTRUCTION FORMAT:**

P				P+1	
7	6	5	4	3	2 1 0 7 0
0	3	4			OPERAND

NOTE

1. The contents of P+1 are unchanged.
2. The Zero, Sign, and Parity Flip-flops represent the status of the A register at the completion of this instruction.

SUBTRACT:**OP CODE: 23s****SBr_s SBM**

TIMING: 16 usec. if

RR, 520 usec. if MR

OPERATION: $(A) - (r_s) - (C_f) \rightarrow A$ or $(A) - (M) - (C_f) \rightarrow A$

DESCRIPTION: This instruction is identical to SUBTRACT WITH BORROW IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

P			
7	6	5	4 3 2 1 0
2	3		s

s: specifies the operand source.
Refer to Table 3-1 for source codes.

AND IMMEDIATE: ND

OP CODE: 044 TIMING: 16 usec.

OPERATION: $(P+1) \Delta (A) \rightarrow A, P+2 \rightarrow P$

DESCRIPTION: Forms the logical product of the contents of the A register with the contents of the memory location immediately following the instruction, and places the results in the A register.

INSTRUCTION FORMAT:

P				P+1		
7	6	5	4	3	2	1
0	4	4	OPERAND			

NOTE

1. The Carry Flip-flop will be reset upon completion of the operation.
2. The Zero, Sign, and Parity Flip-flops will represent the status of the A register upon completion of the operation.

SAMPLE OPERATION:

(A Reg)	0	0	1	1
(P+1)	0	1	0	1
(A Reg)	0	0	0	1

OR IMMEDIATE: OR

OP CODE: 064 TIMING: 16 usec.

OPERATION: $(A) V (P+1) \rightarrow A, P+2 \rightarrow P$

DESCRIPTION: Forms the logical sum of the contents of the A register and the contents of the memory location immediately following the instruction, and places the result in the A register.

INSTRUCTION FORMAT:

P				P+1		
7	6	5	4	3	2	1
0	6	4	OPERAND			

NOTE

1. The Carry Flip-flop will be reset at conclusion.
2. The Zero, Sign, and Parity Flip-flops will represent the status of the A register at completion of the operation.

SAMPLE OPERATION:

(A Reg)	0	0	1	1
(P+1)	0	1	0	1
(A Reg)	0	1	1	1

AND:

OP CODE: 24s

ND_{r_s}, NDM

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: $(A) \Delta (r_s) \rightarrow A$, or $(A) \Delta (M) \rightarrow A$

DESCRIPTION: This instruction is identical to AND IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

7	6	5	4	3	2	1	0
2	4	_s					

_s: specifies the operand source.
Refer to Table 3-1 for source codes.

OR:

OP CODE: 26s

OR_{r_s}, ORM

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: $(A) V (r_s) \rightarrow A$, or $(A) V (M) \rightarrow A$

DESCRIPTION: This instruction is identical to OR IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

7	6	5	4	3	2	1	0
2	6	_s					

_s: specifies the operand source.
Refer to Table 3-1 for source codes.

EXCLUSIVE OR**IMMEDIATE:** XR

OP CODE: 054 TIMING: 16 usec.

OPERATION: $(A) \vee (P+1) \rightarrow A, P+2 \rightarrow P$

DESCRIPTION: The logical difference of the contents of the A register and the contents of the memory location immediately following the instruction is formed, and the result replaces the contents of the A register.

INSTRUCTION FORMAT:

P				P+1	
7	6	5	4	3	2 1 0
0	5		4		OPERAND

NOTE

1. The Carry Flip-flop will be reset at conclusion.
2. The Zero, Sign and Parity Flip-flops will represent the status of the A register upon completion of the operation.

SAMPLE OPERATION:

(A Reg)	0	0	1	1
(P+1)	0	1	0	1
(A Reg)	0	1	1	0

EXCLUSIVE OR:**XR_{r_s} XRM**

OP CODE: 25s

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: $(A) \vee (r_s) \rightarrow A, (A) \vee (M) \rightarrow A$

DESCRIPTION: This instruction is identical to EXCLUSIVE OR IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

P			
7	6	5	4
4	3	2	1
s			

s: specifies the operand source.
Refer to Table 3-1 for source codes.

COMPARE**IMMEDIATE:** CP

OP CODE: 074 TIMING: 16 usec.

OPERATION: $(A) : (P+1), P+2 \rightarrow P$

DESCRIPTION: Compares the contents of the A register with the contents of the memory location immediately following the instruction. The flag flip-flops assume the same state as they would for a Subtract instruction.

INSTRUCTION FORMAT:

P				P+1	
7	6	5	4	3	2 1 0
0	7		4		OPERAND

NOTE

1. The contents of the A register are unaffected.

EXCLUSIVE OR:**XR_{r_s} XRM**

OP CODE: 25s

TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: $(A) \vee (r_s) \rightarrow A, (A) \vee (M) \rightarrow A$

DESCRIPTION: This instruction is identical to EXCLUSIVE OR IMMEDIATE with the exception of operand source.

COMPARE:**CPr_s CPM**

OP CODE: 27s TIMING: 16 usec. if RR, 520 usec. if MR

OPERATION: $(A) : (r_s)$ or $(A) : (M)$

DESCRIPTION: This instruction is identical to COMPARE IMMEDIATE with the exception of operand source.

INSTRUCTION FORMAT:

P			
7	6	5	4
4	3	2	1
s			

s: specifies the operand source.
Refer to Table 3-1 for source codes.

UNCONDITIONAL

JUMP:

JMP

OP CODE: 104

TIMING: Variable*

OPERATION: $(P+1, P+2) \rightarrow P$

DESCRIPTION: An unconditional transfer of control. The contents of P+1 represent the least significant portion of the address, while the contents of P+2 represent the most significant portion.

INSTRUCTION FORMAT:

P	P+1	P+2
OP CODE	ADDRESS	
7 6 5 4 3 2 1 0 7	0 7	0
1 0 4 LSP MSP		

The three high order bits in the address are ignored, the remaining 13 bits specify the address to which control is to be transferred.

NOTE

*Timing is variable dependent upon cyclic difference between instruction and effective address locations.

JUMP IF CONDITION

TRUE:

JTf_c

OP CODE: 1(c+4)0

TIMING: Variable if condition true, 24 usec if condition false

OPERATION: If ($f_c = \text{TRUE}$), $(P+1, P+2) \rightarrow P$. Otherwise, $P+3 \rightarrow P$

DESCRIPTION: Examines the designated flip-flop. If set, transfers control to the address designated by the contents of the two memory locations immediately following the instruction. If the selected flip-flop is reset, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

P	P+1	P+2
OP CODE	ADDRESS	
7 6 5 4 3 2 1 0 7 0 7 0		
1 c+4 0 LSP MSP		

c: designates which flip-flop condition is to be tested. Refer to Table 3-2 for list of Flip-flop codes

NOTE

1. The condition of the selected Flip-flop is unchanged by this instruction.

JUMP IF CONDITION

FALSE:

OP CODE: 1c0

JFf_C

TIMING: Variable if condition false, 24 usec. if condition true.

OPERATION: If ($f_c = \text{FALSE}$), $(P+1, P+2) \rightarrow P$. Otherwise $P+3 \rightarrow P$.

DESCRIPTION: Examines the designated flip-flop. If reset, transfers control to the address designated by the contents of the two memory locations immediately following the instruction. If the selected flip-flop is set, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

P	P+1	P+2
OP CODE	ADDRESS	
7 6 5 4 3 2 1 0 7 0 7 0	LSP MSP	
1 c 0		

c: designates which flip-flop (condition) is to be tested. Refer to Table 3-2 for list of flip-flop codes.

NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.

SUBROUTINE CALL: CALL

OP CODE: 106

TIMING: Variable

OPERATION: $P+3 \rightarrow \text{STACK}$, $(P+1, P+2) \rightarrow P$

DESCRIPTION: Transfers the address of the next sequentially available instruction to the Pushdown Stack, and transfer control to the address specified by the contents of the two memory locations immediately following the Op Code.

INSTRUCTION FORMAT:

P	P+1			P+2
7 6	5 4 3	2 1 0	7	0 0 7 0
1	0	6	LSP	MSP

ADDRESS

NOTE

1. The Stack is open-ended in operation. If it is overfilled, the deepest address will be lost.

CONDITIONAL SUBROUTINE CALL

IF CONDITION TRUE: CTf_c

OP CODE: 1(c+4)2

TIMING: Variable if condition true, 24 usec. if condition false.

OPERATION: If (f_c =TRUE), $P+3 \rightarrow$ STACK, $(P+1, P+2) \rightarrow P$. Otherwise, $P+3 \rightarrow P$.

DESCRIPTION: Examines the designated flip-flop. If set, transfers the address of the next sequentially available instruction to the pushdown stack, and transfers control to the address of the two memory locations immediately following the Op Code. If the selected flip-flop is reset, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

ADDRESS							
				P+1	P+2		
7	6	5 4 3	2 1 0	7	0	7	0
1	c+4		2		LSP		MSP

c: designates which flip-flop (condition) is to be tested.

NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The stack is open-ended in operation. If it is overfilled, the deepest address will be lost.
3. Refer to Table 3-2 for list of flip-flop codes.

CONDITIONAL SUBROUTINE CALL

IF CONDITION FALSE: CFF_c

OP CODE: 1c2

TIMING: Variable if condition false, 24 usec. if condition true.

OPERATION: If ($f_c = \text{FALSE}$), $P+3 \rightarrow \text{STACK}$, $(P+1, P+2) \rightarrow P$.

DESCRIPTION: Examines the designated flip-flop. If reset, transfers the address of the next sequentially available instruction to the pushdown stack, and transfers control to the address of the two memory locations immediately following the Op Code. If the selected flip-flop is set, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

ADDRESS								P+1	P+2
7	6	5	4	3	2	1	0	7	0 7 0

c: designates which flip-flop (condition) is to be tested.

NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The stack is open-ended in operation. If it is overfilled, the deepest address will be lost.
3. Refer to Table 3-2 for list of flip-flop codes.

SUBROUTINE**RETURN:**

OP CODE: 007

RETURN

TIMING: Variable

OPERATION: (STACK) \rightarrow P

DESCRIPTION: Transfer control to the address specified by the most recent entry in the Pushdown Stack. Deletes the most recent entry from the Stack.

INSTRUCTION FORMAT:

P
7 6 5 4 3 2 1 0
0 0 7

NOTE

1. The effect of attempting more "RETURN" than the Stack is capable of handling is undefined.

CONDITIONAL SUBROUTINE RETURN**IF CONDITION FALSE: RFf_c**

OP CODE: 0c3

TIMING: Variable if condition false, 16 usec. if condition true.

OPERATION: If ($f_c = \text{FALSE}$), Stack \rightarrow P. Otherwise, P+1 \rightarrow P

DESCRIPTION: Examines the designated flip-flop. If reset, transfers control to the address specified by the most recent entry in the stack. If the selected flip-flop is set, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

7 6 5 4 3 2 1 0
1
c 3

c: designates which flip-flop (condition) is to be tested.

NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The effect of attempting more "RETURN" than the stack is capable of handling is undefined.
3. Refer to Table 3-2 for list of flip-flop codes.

CONDITIONAL SUBROUTINE RETURN**IF CONDITION TRUE: RTf_c**

OP CODE: 0(c+4)3

TIMING: Variable if condition true, 16 usec. if condition false.

OPERATION: If ($f_c = \text{TRUE}$), Stack \rightarrow P. Otherwise P+1 \rightarrow P

DESCRIPTION: Examines the designated flip-flop. If set, transfers control to the address specified by the most recent entry in the pushdown stack. Deletes the most recent entry in the stack. If the selected flip-flop is reset, executes the next sequentially available instruction.

INSTRUCTION FORMAT:

7 6 5 4 3 2 1 0
1
c+4 3

c: designates which flip-flop (condition) is to be tested.

NOTE

1. The condition of the selected flip-flop is unchanged by this instruction.
2. The effect of attempting more "RETURN" than the stack is capable of handling is undefined.
3. Refer to Table 3-2 for list of flip-flop codes.

SHIFT RIGHT**CIRCULAR:** **SRC**

OP CODE: 012 TIMING: 16 usec.

OPERATION: $A_m \rightarrow A_{m-1}$, $A_0 \rightarrow A_7$, $A_0 \rightarrow C_f$

DESCRIPTION: Shifts the contents of the A register right in a circular fashion. Shifts the least significant bit into the most significant bit position. Upon completion of the operation, the Carry Flip-flop is equal to the most significant bit.

INSTRUCTION FORMAT:

P	7	6	5	4	3	2	1	0
	0	1	2					

NOTE

None of the flag flip-flops other than C_f is affected by this instruction.

SHIFT LEFT**CIRCULAR:** **SLC**

OP CODE: 002 TIMING: 16 usec.

OPERATION: $A_m \rightarrow A_{m+1}$, $A_7 \rightarrow A_0$, $A_7 \rightarrow C_f$

DESCRIPTION: Shifts the contents of the A register left in a circular fashion. Shifts the most significant bit into the least significant bit position. Upon completion of the operation, the Carry Flip-flop is equal to the least significant bit.

INSTRUCTION FORMAT:

P	7	6	5	4	3	2	1	0
	0	0	2					

NOTE

None of the flag flip-flops other than C_f is affected by this instruction.

NO OPERATION: **NOP**

OP CODE: 300 TIMING: 16 usec.

OPERATION: $P+1 \rightarrow P$

DESCRIPTION: No instruction is executed.

INSTRUCTION FORMAT:

P	7	6	5	4	3	2	1	0
	3	0	0					

HALT:

OP CODE: 000,001,377

OPERATION:

DESCRIPTION: The computer halts. When the START button on the console is depressed, operation resumes at $P+1$.**HALT**

TIMING: Execution Stops

INSTRUCTION FORMAT:

P	7	6	5	4	3	2	1	0
	0	0	0					
	0	0	1					
	3	3	7					

INPUT:

OP CODE: 101

INPUT

TIMING: 16 usec.

OPERATION: (I/O Bus) \rightarrow A

DESCRIPTION: Transfers the contents of the I/O Bus to the A register.

INSTRUCTION FORMAT:

P	7	6	5	4	3	2	1	0
	1	0	1					

EXTERNAL COMMAND: EX (exp)
OP CODE: 121-177 depending
on the specific command being
executed.
OPERATION: Performs I/O control functions according
to (exp)
DESCRIPTION: These instructions perform the functions
necessary for control of the I/O system and external de-
vices. Many of these functions are specifically related to
operation of particular devices. The device oriented com-
mands for the Keyboard, CRT Display, Cassette Tapes,
and Communications Interface are explained in the sections
covering these devices.

INSTRUCTION FORMAT:

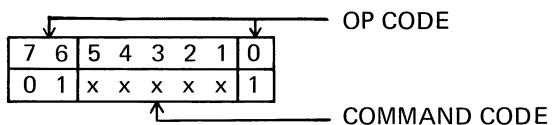


Table 3-3 is a list of External Commands used. For a detailed discussion of their use, reference should be made to Part 4 (Input/Output Operations) and to descriptions of the separate external devices.

TABLE 3-3
EXTERNAL COMMANDS

EX (exp)

COMMAND NUMBER	(exp)	OCTAL CODE	COMMAND	DESCRIPTION	DEVICE ADDRESS
1	ADR	121	Address	Selects device specified by A-register	ALL
2	STATUS	123	Sense Status	Connects selected device status to input lines	
3	DATA	125	Sense Data	Connects selected device data to input lines	
4	WRITE	127	Write Strobe	Signals selected device that output data word is on output lines	
5	COM1	131	Command 1	Outputs a control function to selected device	
6	COM2	133	Command 2	Outputs a control function to selected device	
7	COM3	135	Command 3	Outputs a control function to selected device	
8	COM4	137	Command 4	Outputs a control function to selected device	ALL
9	--	141	(Unassigned)	--	--
10	--	143	(Unassigned)	--	--
11	--	145	(Unassigned)	--	--
12	--	147	(Unassigned)	--	--

TABLE 3-3
EXTERNAL COMMANDS
EX (exp)
(Continued)

COMMAND NUMBER	(exp)	OCTAL CODE	COMMAND	DESCRIPTION	DEVICE ADDRESS
13	BEEP	151	Beep	Activates tone producing mechanism	341
14	CLICK	153	Click	Activates audible click producing mechanism	341
15	DECK1	155	Select Deck 1	Connects deck 1 to I/O bus	360
16	DECK2	157	Select Deck 2	Connects deck 2 to I/O bus	
17	RBK	161	Read Block	Enables read circuitry and sets tape in forward motion	
18	WBK	163	Write Block	Enables write circuitry and sets tape in forward motion	360
19	--	165	(Unassigned)	--	--
20	BSP	167	Backspace One Block	Backs up the selected tape one record	360
21	SF	171	Slew Forward	Sets selected tape deck in forward motion	
22	SB	173	Slew Backward	Sets selected tape deck in backward motion	
23	REWIND	175	Rewind	Rewinds the selected deck to beginning of tape	
24	TSTOP	177	Stop Tape	Halts motion of the selected tape deck	360

TABLE 3-4
INSTRUCTION REPERTOIRE

OP CODE	MNEMONIC	OP CODE	MNEMONIC	OP CODE	MNEMONIC
000	HALT	050		120	JFS
001	HALT	051		121	EX ADR
002	SLC	052		122	CFS
003	RFC	053	RTZ	123	EX STATUS
004	AD	054	XR	124	
005		055		125	EX DATA
006	LA	056	LH	126	
007	RETURN	057		127	EX WRITE
010		060		130	JFP
011		061		131	EX COM1
012	SRC	062		132	CFP
013	RFZ	063	RTS	133	EX COM2
014	AC	064	OR	134	
015		065		135	EX COM3
016	LB	066	LL	136	
017		067		137	EX COM4
020		070		140	JTC
021		071		141	
022		072		142	CTC
023	RFS	073	RTP	143	
024	SU	074	CP	144	
025		075		145	
026	LC	076		146	
027		077		147	
030		100	JFC	150	JTZ
031		101	INPUT	151	EX BEEP
032		102	CFC	152	CTZ
033	RFP	103		153	EX CLICK
034	SB	104	JMP	154	
035		105		155	EX DECK1
036	LD	106	CALL	156	
037		107		157	EX DECK2
040		110	JFZ	160	JTS
041		111		161	EX RBK
042		112	CFZ	162	CTS
043	RTC	113		163	EX WBK
044	ND	114		164	
045		115		165	
046	LE	116		166	
047		117		167	EX BSP

TABLE 3-4
INSTRUCTION REPERTOIRE
(Continued)

OP CODE	MNEMONIC	OP CODE	MNEMONIC	OP CODE	MNEMONIC
170	JTP	240	NDA	310	LBA
171	EX SF	241	NDB	311	
172	CTP	242	NDC	312	LBC
173	EX SB	243	NDD	313	LBD
174		244	NDE	314	LBE
175	EX REWND	245	NDH	315	LBH
176		246	NDL	316	LBL
177	EX TSTOP	247	NDM	317	LBM
200	ADA	250	XRA	320	LCA
201	ADB	251	XRB	321	LCB
202	ADC	252	XRC	322	
203	ADD	253	XRD	323	LCD
204	ADE	254	XRE	324	LCE
205	ADH	255	XRH	325	LCH
206	ADL	256	XRL	326	LCL
207	ADM	257	XRM	327	LCM
210	ACA	260	ORA	330	LDA
211	ACB	261	ORB	331	LDB
212	ACC	262	ORC	332	LDC
213	ACD	263	ORD	333	
214	ACE	264	ORE	334	LDE
215	ACH	265	ORH	335	LDH
216	ACL	266	ORL	336	LDL
217	ACM	267	ORM	337	LDM
220	SUA	270	CPA	340	LEA
221	SUB	271	CPB	341	LEB
222	SUC	272	CPC	342	LEC
223	SUD	273	CPD	343	LED
224	SUE	274	CPE	344	
225	SUH	275	CPH	345	LEH
226	SUL	276	CPL	346	LEL
227	SUM	277	CPM	347	LEM
230	SBA	300	NOP	350	LHA
231	SBB	301	LAB	351	LHB
232	SBC	302	LAC	352	LHC
233	SBD	303	LAD	353	LHD
234	SBE	304	LAE	354	LHE
235	SBH	305	LAH	355	
236	SBL	306	LAL	356	LHL
237	SBM	307	LAM	357	LHM

TABLE 3-4
INSTRUCTION REPERTOIRE
(Continued)

OP CODE	MNEMONIC	OP CODE	MNEMONIC	OP CODE	MNEMONIC
360	LLA	370	LMA		
361	LLB	371	LMB		
362	LLC	372	LMC		
363	LLD	373	LMD		
364	LLE	374	LME		
365	LLH	375	LMH		
366		376	LML		
367	LLM	377	HALT		

NOTE

OP Codes shown without Mnemonics are undefined.

PART 4

INPUT/OUTPUT OPERATIONS

4.1 GENERAL

The versatile input/output capability of the Datapoint 2200 permits it to communicate with external devices (such as the 2200 communications adaptor) through a parallel I/O system. The keyboard, c.r.t. and tape decks that are an integral part of the Model 2200 perform all operations over the same I/O system as external devices.

4.2 INPUT/OUTPUT INSTRUCTIONS

Two types of instructions provide for I/O operations. One is the INPUT command (see section 3) which, upon execution, transfers whatever is on the input bus to the A-register. The second is the EXTERNAL command which is sub-divided into 24 separate command operations (8 of which are available to devices physically external to the Model 2200). Each EXTERNAL command produces a strobe pulse which may be used for control external to the processor. The actual control functions assigned to each external command are listed in Table 3-3.

4.3 INPUT/OUTPUT CABLE

The parallel I/O cable carries data, input strobe, external commands, and power between the 2200 processor and external devices connected to it. A complete I/O system is structured by connecting external devices in partyline fashion as shown in Figure 2-1. The I/O cable contains 8 input data lines, 8 output data lines, 1 input strobe line, 8 (of the 24) external command lines, 1 clock line, and 7 power and ground lines.

4.4 I/O DATA LINES

The data lines are broken into two groups. 8 lines are used for output and 8 lines are used for input.

The data output lines are connected (at all times) to the A-register in the processor and are used to perform three basic functions:

- a. To transfer an address to select an external device (including the keyboard, c.r.t. and tape decks);
- b. To transfer commands to an addressed device; and
- c. To transfer data to an addressed device.

The data input lines are strobed into the A-register upon execution of the INPUT instruction and used to perform two basic functions.

- a. To transfer status information from an addressed external device; and
- b. To transfer data from an addressed external device.

As shown in Figure 4-1, input data or status from the data input lines is processed through input receivers and gated into the A-register. Once in the A-register data can be manipulated or stored as desired. Addresses, commands, or data that is to be transferred to an external device must first be loaded into the A-register. From the A-register it is transmitted through output devices onto the data output lines. The A-register, then, is used as a buffer register between the 2200 processor and external devices for all input and output data transfers.

4.5 INPUT STROBE

The INPUT STROBE carries a signal (8 usec. pulse) from the processor to the external device to indicate that whatever data is on the data input lines has been sampled and transferred into the A-register. The trailing edge of the pulse may be used by an external device to remove data from the data input line or to clear a status bit. The INPUT strobe is generated upon execution of the INPUT instruction.

4.6 EXTERNAL COMMAND STROBES

The eight EXTERNAL commands used by devices physically external to the Model 2200 are given function assignments as follows:

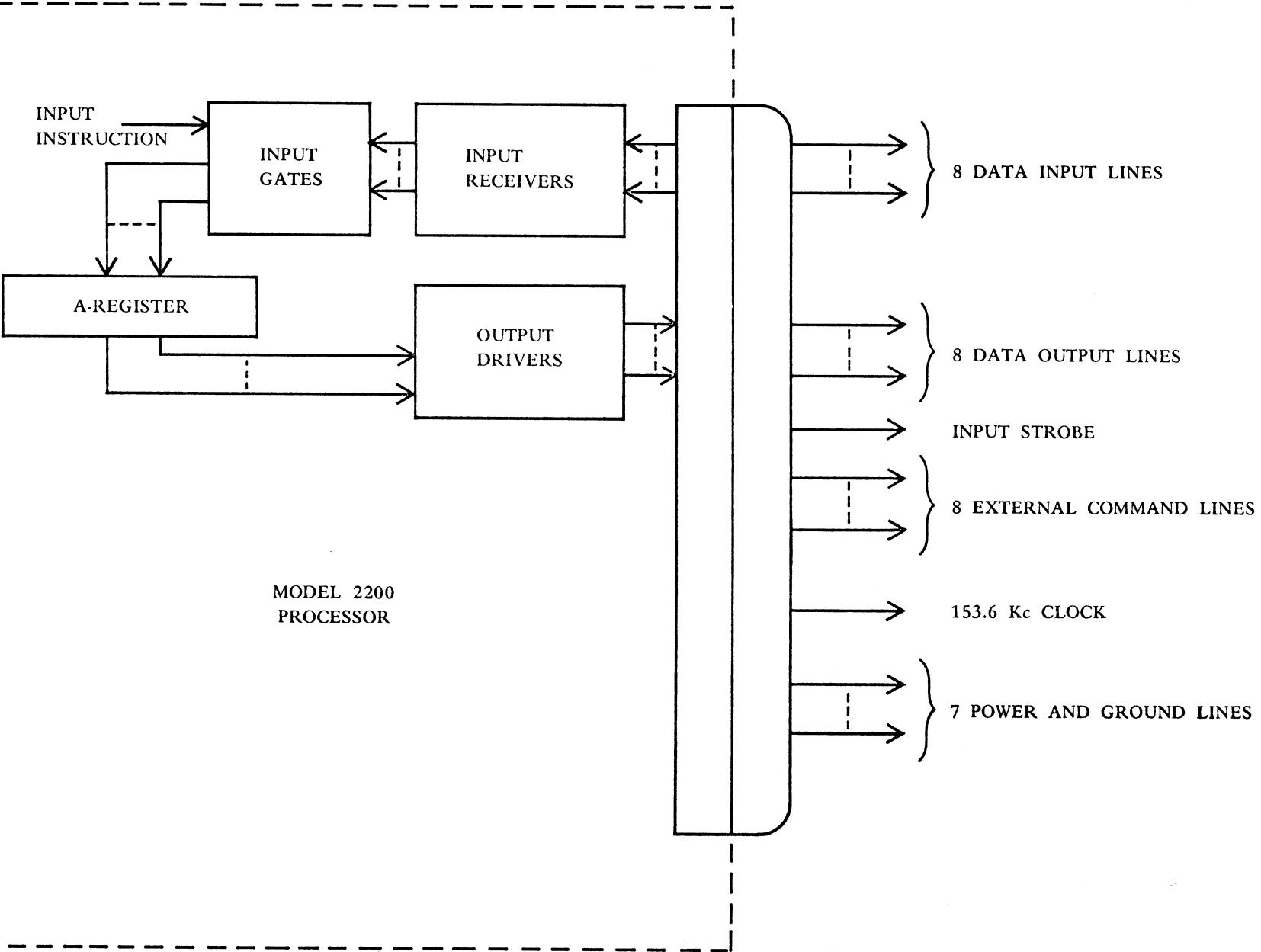


FIGURE 4-1
I/O SYSTEM, FUNCTIONAL DIAGRAM

TABLE 4-1
EXTERNAL COMMANDS
EX (exp)

COMMAND NUMBER	(exp)	OCTAL CODE	COMMAND	DESCRIPTION
1	ADR	121	Address	Selects device specified by A-register
2	STATUS	123	Sense Status	Connects selected device status lines to data input bus
3	DATA	125	Sense Data	Connects selected device data lines to data input bus
4	WRITE	127	Write Strobe	Signals selected device that output data is on data output lines
5	COM1	131	Command 1	Signals selected device that a control word is on data output lines
6	COM2	133	Command 2	Signals selected device that a control word is on data output lines
7	COM3	135	Command 3	Signals selected device that a control word is on data output lines
8	COM4	137	Command 4	Signals selected device that a control word is on data output lines

Execution of an EXTERNAL instruction provides a pulse 8 microseconds long. No functions are performed within the 2200 processor during execution of an EXTERNAL instruction. The interpretation of each of the EXTERNAL instructions is as follows:

- a. Address. The address command (EX ADR) is a signal from the processor to all external devices to indicate that the information on the data output bus is to be interpreted as an external device address. Whenever an address command is executed all external devices should be disconnected from the I/O system except the device whose address appears in the A-register. (See paragraph 4.10 for discussion of address assignments).
- b. Sense Status. The sense status (EX STATUS) command is a signal from the processor to the selected external device to place status information on the data input lines. (Note: External devices should be configured such that status is connected to the data input line whenever the device is first addressed. It is only necessary to use the EX STATUS instruction when it is desired to sense status after an EX DATA instruction has been used and a new address sequence has not been executed).
- c. Sense Data. The sense data (EX DATA) command is a signal from the processor to the selected external device to place its data on the data input lines.

- d. Write Strobe. The write strobe (EX WRITE) command is a signal from the processor that data is present on the data output lines for the selected external device.
- e. Command 1 through Command 4. Command 1 through Command 4 (EX COM1, etc.) have meaning appropriate to the device selected. Reference should be made to a description of each device for specific function assignments.

4.7 CLOCK LINE

The clock line is crystal controlled 153.6 kilohertz square-wave that is available to external devices for timing purposes.

4.8 I/O BUS ELECTRICAL SPECIFICATIONS

All signals in the I/O System operate with a voltage swing of zero to +5 volts. Line drivers have a source impedance of approximately 470 ohms and line receivers have an input impedance in excess of 18,000 ohms and a decision threshold of +1.7 volts. Figure 4-2 illustrates a typical output line circuit.

All logic levels are True (logical 1) for zero (less than +1.7) volts and False (logical 0) for +5 (greater than 1.7) volts.

4.9 DATA TRANSFER OPERATION

- a. Data Output. Figure 4-3 illustrates the sequence of events that occur when data is transferred from the

2200 processor to an external device. A typical program sequence to execute a transfer is as follows:

WDATA	LA 0322	Load device address into A-register
EX ADR		
INPUT		Load device status into A-register
SRC		Shift desired status bit into C flip-flop
JFC EXIT		Exit if device not ready
LAM		Load A-register with DATA
EX WRITE		Write Data to device

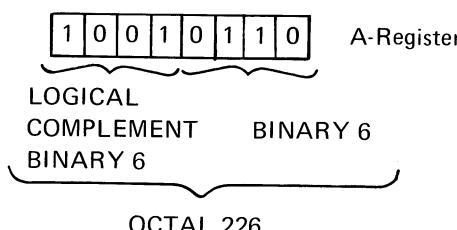
Once a device is addressed it remains addressed until another device is addressed so that succeeding commands may be transmitted to a device without re-addressing the device. Transmitting a command to a device would follow a program sequence similar to a data transfer except that EX COM_n would replace EX WRITE.

- b. Data Input. Figure 4-4 illustrates a sequence of events that occur when data is transferred from an external device to the 2200 processor. A typical program sequence is as follows:

RDATA	LA 0322	Load A-register with device address
EX ADR		
INPUT		Load device status into A-register
SRC		Shift status bit into C flip-flop
SRC		
JFC EXIT		Exit if device not ready
EX DATA		Place data on input lines
INPUT		Load A-register with data

4.10 DEVICE ADDRESS NUMBERING

Address assignments in the I/O system provides for up to 16 devices external to the 2200 processor. The address word is formulated such that the low-order four bits form the binary value for the address and the high-order four bits form the logical complement of the low order bits. For example device number 6 would have an address word as follows:



This addressing system permits any device to be coded for its particular address with only a four-input gate strapped to those output lines that are set to one during the address command. In addition, all devices can be cleared by setting the A-register to all zeros and executing an EX ADR instruction.

Device addresses used in the Model 2200 are given in the following table:

TABLE 4-2

DEVICE ADDRESS ASSIGNMENTS

DEVICE	NUMBER	BINARY	OCTAL
Cassette Tape Decks	0	11110000	360
CRT/Keyboard	1	11100001	341
Communications Adaptor	2	11010010	322
2200P Printer	3	11000011	303
2200T Tape Transport	4	10110100	264
Unassigned	5	10100101	245
"	6	10010110	226
"	7	10000111	207
"	8	01111000	170
"	9	01101001	151
"	10	01011010	132
"	11	01001011	113
"	12	00111100	074
"	13	00101101	055
"	14	00011110	036
"	15	00001111	017

4.11 I/O POWER AND GROUND LINES

The Model 2200 provides several power supply voltages for use by external devices. Table 4-3 below lists the characteristics of each power and ground line.

4.12 I/O SYSTEM CONNECTOR

Connection to the I/O system is made through an Amphenol 17-20500-1 connector. The mating I/O cable should have a 50-pin Amphenol 17-10500-1 connector.

Table 5-5 lists the pin assignments.

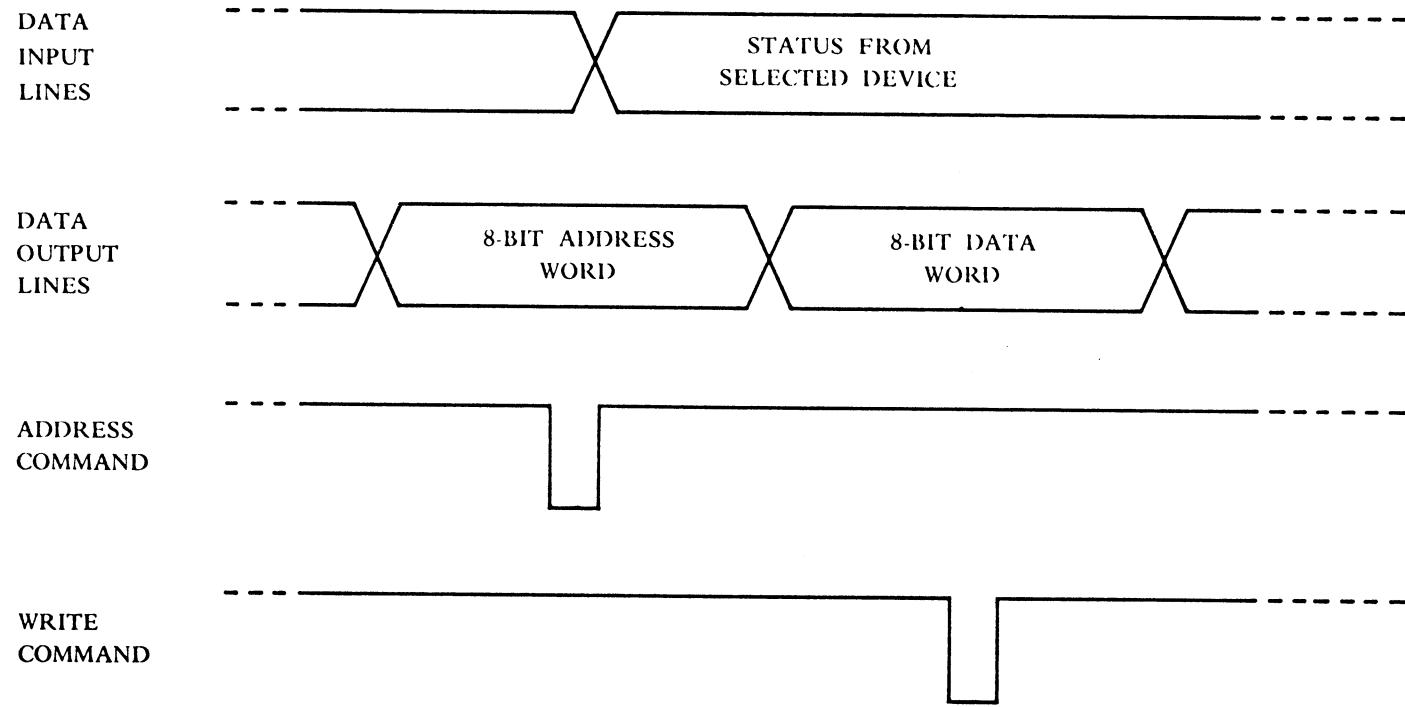


FIGURE 4-3
TYPICAL DATA OUTPUT
SEQUENCE

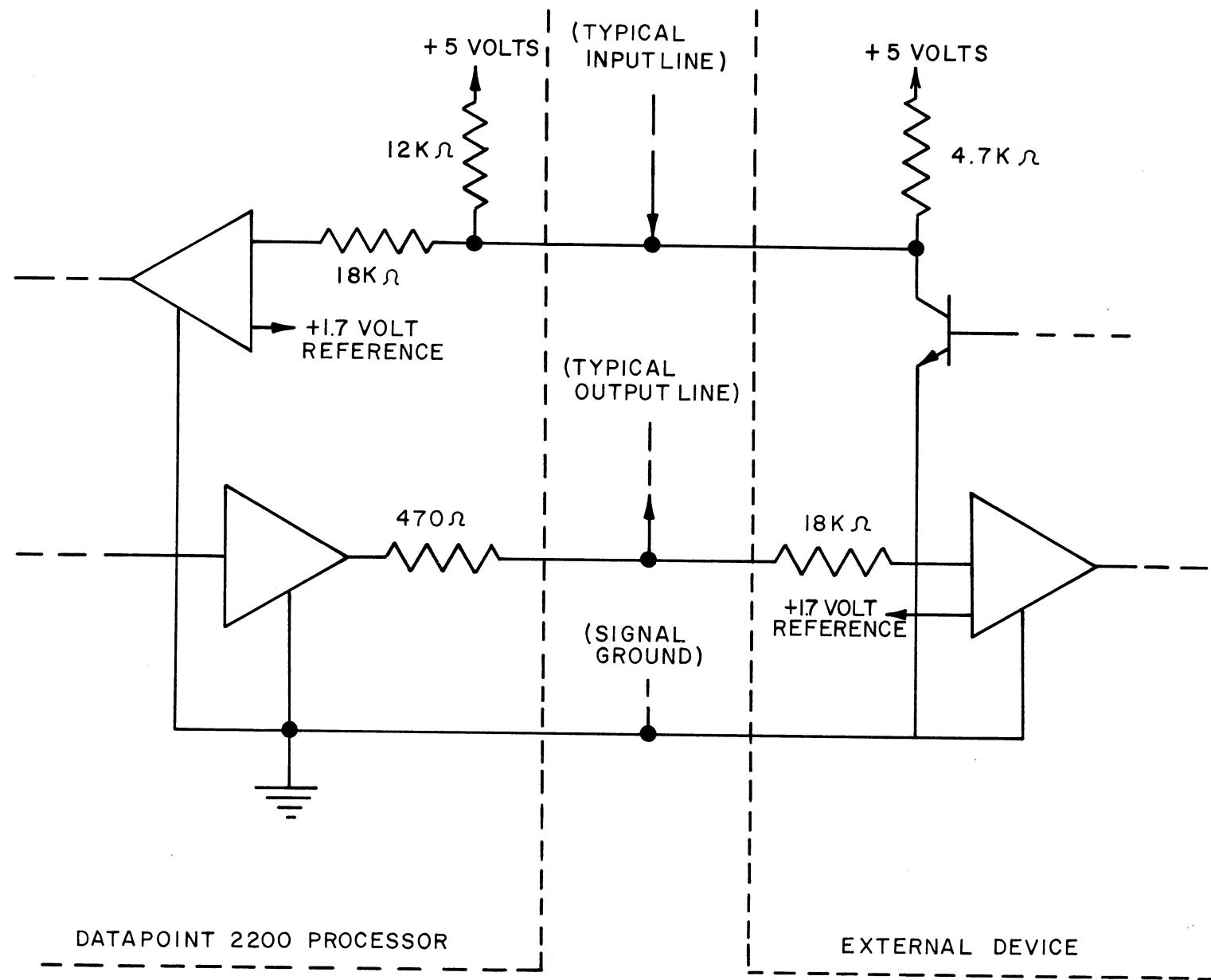


FIGURE 4-2

I/O CABLE, ELECTRICAL CHARACTERISTICS

TABLE 4-3
I/O POWER AND GROUND LINES

VOLTAGE	MAX. CURRENT	REGULATION
-12 Volts	0.5 amps	±10%
- 5 Volts	0.1 amps	±5%
+ 5 Volts	3.4 amps	±5%
+12 Volts	0.5 amps	±10%
+24 Volts	0.1 amps	±5%
Power Ground	—	—
Signal Ground	—	—

TABLE 4-4
I/O CONNECTOR PIN ASSIGNMENTS

ASSIGNMENT	PIN NUMBER
Data output 0	44
1	45
(A Bus Outputs) 2	46
3	29
4	30
5	31
6	32
7	33
Data Input 0	1
1	2
(A Bus Inputs) 2	3
3	4
4	5
5	6
6	7
7	18
Input Strobe (Read)	12
Address Command	15
Sense Status Command	13
Sense Data Command	14
Write Command	19
Command 1	20
Command 2	21
Command 3	22
Command 4	23
153.6 KHz Clock	39
-12v	24
-5v	27
+5v	8, 9, 10, 11
+12v	25
+24v	26
Ground (Power & Signal)	40, 41, 42, 43

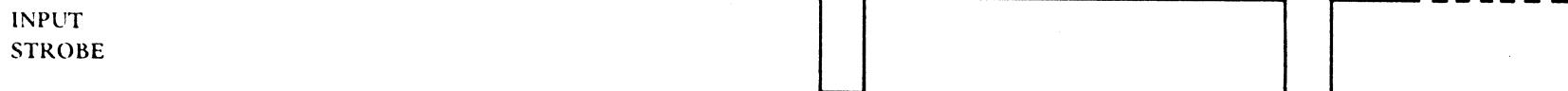
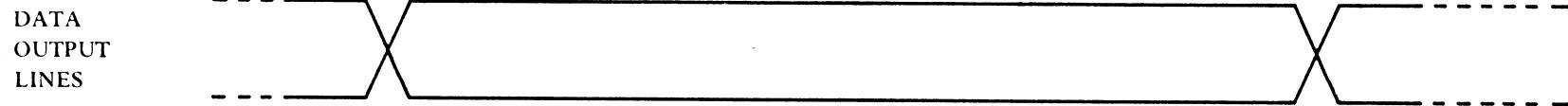


FIGURE 4-4
TYPICAL DATA INPUT
SEQUENCE

PART 5

KEYBOARD

5.1 GENERAL DESCRIPTION

The keyboard on the Datapoint 2200 performs the functions of data entry and processor control. The keys are divided into three sections, each of which has its own function.

Section 1 consists of 41 standard alphabetic, numeric and special character keys found in the ASCII character set. Figure 5-1 illustrates the keyboard layout.

Section 2 consists of an 11 key matrix which is identical to a standard adding machine keyboard with the addition of a decimal point (period). The keys in this section are duplicates of certain keys found in Section 1 and are provided to facilitate entry of large amounts of numeric data.

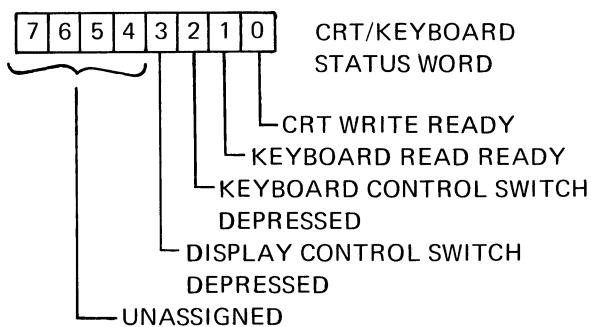
The keys in Section 3 are special function keys which exert control over the processor. Their names and associated functions are as follows:

RUN	Momentary contact switch, which when depressed, causes the processor to begin execution of the instruction located at the address in memory currently addressed by the program counter.
STOP	Momentary contact switch which, when depressed, causes instruction execution to halt at the completion of the current instruction. Care should be taken when using this switch, because any tape operation which may be in progress will be aborted.
KEYBOARD	Momentary contact switch which sets a status bit that may be tested at any time by the processor.
DISPLAY	Momentary contact switch with a function similar to that of KEYBOARD switch. Either one or both of these switches may be depressed.
RESTART	Momentary contact switch which causes the processor to halt, rewind the system or program tape mounted on Deck 1, load and execute the first record found on tape.

5.2 OPERATION

The keyboard is addressed by the processor by loading the A-register with 3418 and executing an EX ADR command. (The crt display also uses this address. Data transfers to the

processor are from the keyboard and transfers from the processor are to the display). Following the address sequence the c.r.t./keyboard status word can be loaded into the A-register by executing an INPUT instruction. Bit 1 of the A-register may be tested by the program to determine if a character is ready for transfer from the keyboard. Bits 2 and 3 will indicate if either the KEYBOARD or DISPLAY control switch is pressed.



The External Commands associated with the operation of the keyboard are as follows:

- a. EX BEEP. This command produces a 1500 Hertz tone for a duration of about 100 msec. The tone could be used as an error or ready signal to the keyboard operator.
- b. EX CLICK. This command produces an audible click which could be used to acknowledge receipt of a valid character when a key is depressed.
- c. EX COM1 (Command 1). Presents a control word contained in the A-register to the keyboard. Bit 5 of the control word controls the KEYBOARD switch light and bit 6 controls the DISPLAY switch light as follows:

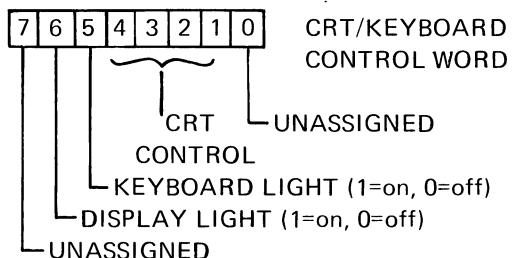


TABLE 5-1
KEYBOARD CODING (ASCII)

A-101	a -141	0-060	:	-072
B-102	b -142	1-061	;	-073
C-103	c -143	2-062	<	-074
D-104	d -144	3-063	=	-075
E-105	e -145	4-064	>	-076
F-106	f -146	5-065	?	-077
G-107	g -147	6-066	[-133
H-110	h -150	7-067	~	-176
I -111	i -151	8-070]	-135
J -112	j -152	9-071	Λ	-136
		Space-040	—	-137
K-113	k -153			
L-114	l -154	!-041	@	-100
M-115	m -155	"-042	{	-173
N-116	n -156	#-043	\	-134
O-117	o -157	\$-044	'	-140
P-120	p -160	%-045		-174
Q-121	q -161	&-046	}	-175
R-122	r -162	'-047	Enter	-015
S-123	s -163	(-050	Cancel	-030
T-124	t -164)-051	Backspace	-010
U-125	u -165	*-052	Rubout (R.O.)	-177
V-126	v -166	+053		
W-127	w -167	, -054		
X-130	x -170	-055		
Y-131	y -171	. -056		
Z-132	z -172	/ -057		

PART 6

CRT DISPLAY

6.1 GENERAL DESCRIPTION

The display unit on the Datapoint 2200 consists of a CRT capable of displaying 12 lines of 80 characters each, a character generator, 960 cells of refresh memory (refresh rate 60 Hz), and a group of registers utilized to position the cursor. Maximum character transfer rate to the CRT is 60 characters per second.

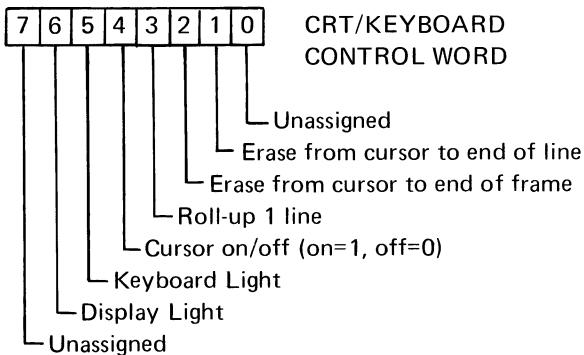
The character set utilized by the CRT display consists of the full ASCII set with both upper and lower case alphabetics and all numeric and special characters.

6.2 OPERATION

The CRT is addressed and status tested in the same manner as the keyboard (see paragraph 5.2). Bit 0 of the status word indicates that the CRT is ready to accept data or commands. Characters are transferred to the screen by loading the A-register with the character to be displayed and executing an EX WRITE. The character will be displayed at the current cursor location.

Control of the CRT is accomplished through the use of the three external commands - Command 1, Command 2, and Command 3. The functions performed by these commands are as follows:

- a. EX COM1 (Command 1) Transfers a control word contained in the A-register to the CRT. The applicable bit assignments and their functions are as follows:



The erase functions permit selective erasures on the screen by limiting erasures to those character positions following the current cursor position to the end of the line (or page).

The roll-up function causes all displayed characters (not the cursor) to move up one line. The top line on the screen is lost.

The cursor image may be turned on or off through the control word. The cursor position is the same in either case. The cursor image is automatically turned off whenever the processor is in the HALT state.

- b. EX COM2 (Command 2) Positions the cursor to the horizontal character slot designated by the contents of the A-register. Character position $0-79_{10}(0-117_8)$ are valid.
- c. EX COM3 (Command 3) Positions the cursor to the line designated by the contents of the A-register. Line number $0-11_{10}$ ($0-13_8$) are valid.

In order to write a new character, the cursor must occupy that character's position on the screen. After the character has been written, the cursor should then be moved to the next horizontal (or vertical) position desired. The CRT Write Ready status bit must be true before positioning the cursor or displaying a character.

Both the CRT and keyboard utilize the standard ASCII character set. (See Table 5-1). Any invalid character code will appear as a blank space on the CRT screen.

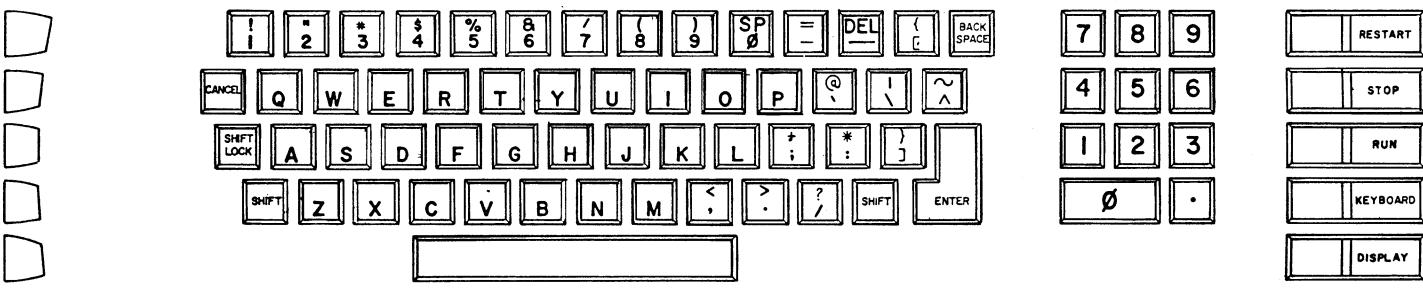


FIGURE 6-1
KEYBOARD LAYOUT

PART 7

CASSETTE TAPES

7.1 GENERAL DESCRIPTION

The Datapoint 2200 contains two cassette tape recording devices for storage of programs and data. Since the hardware RESTART (section 5.1) uses the rear deck (number one), programs will typically be on it while data areas will be the front deck (number two). However, once the machine is initially loaded, either deck may be used for both purposes.

Data on the Tape is organized by record (of any length). Records are written and read at 350 eight-bit characters per second with a tape speed of approximately 7.5 inches per second. See Table 7.1 for a list of the physical specifications.

7.2 OPERATION

Data is recorded or read in bit serial fashion on one track. Each eight bit character is framed by three sync bits on either side of the character:

-	0 1 0	x x x x x x x x	0 1 0	x x x x x x x x	0 1 0	-
Sync Code	Character 1 Code		Sync Code	Character 2 Code		Sync Code

The appearance of the correct sync code indicates that the character is valid. Any other sync code causes special action to be taken on data reads. Note that the sync codes are valid for tape motion in either direction so the tape may be read backwards although in the reverse direction the data bits will appear reverse d (bit 0 will be bit 7, 1 will be 6 etc.)

A record is a group of successive valid characters. An inter-record gap is indicated by the failure of the sync code to be zero one zero and all mark code. (ones):

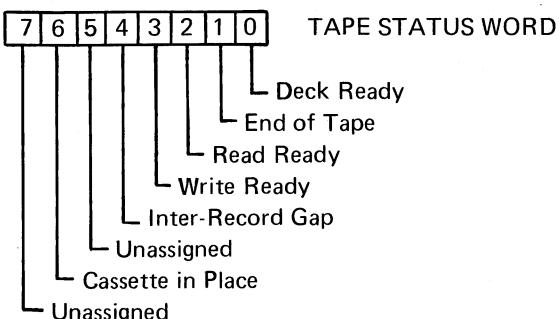
-	0 1 0	x x x x x x x x	0 1 0	x x x x x x x x	1 1 1	x x x	-
Sync Code	Valid Character		Sync Code	Invalid Character	Sync Code	Inter Record Gap	

Only valid characters will be presented as data from the tape unit.

7.3 STATUS

The cassette tape unit is addressed by the processor by loading the A-register with 3608 and executing the EX ADR instruction. Following this sequence, the tape unit

status can be loaded into the A-register by executing an INPUT instruction. The bit assignments are as follows:



- | | |
|-------------------|---|
| DECK READY | Deck ready will be set whenever the tape unit is ready to accept another command. (Only the TSTOP command should be issued if this bit is false). The tape will be stopped, a cassette in the selected deck and not wound to the clear leader at either end, and the head engaged when this bit is true. This bit should be checked after selecting a deck. |
| END OF TAPE | End of Tape indicates that the cassette has run onto leader (in either direction). |
| READ READY | Read Ready indicates that the selected deck has read another character. |
| WRITE READY | Write Ready indicates that the selected deck is ready to write another character. |
| INTER-RECORD GAP | Inter-Record Gap indicates the selected deck has come across an inter-record gap (invalid sync code). |
| CASSETTE IN PLACE | Cassette in Place indicates that a cassette is physically in place in the selected deck. |

7.4 CONTROL

When the cassette tape unit is addressed the following instructions will control the action of the tape:

- a. EX TSTOP causes any motion of either deck to be stopped, any read or write operations to be terminated. When everything has settled, the ready status bit will come true and operations may be resumed.
- b. EX DECK1 causes deck one (rear) to be the currently selected deck. Before commanding a deck selection, care should be taken that the currently selected deck has completed all operations.

- c. EX DECK2 causes deck two (front) to be the currently selected deck. Note the precaution in (b).
- d. EX RBK causes the currently selected deck to be set in forward motion and, after 70 msec, for the read circuitry to be enabled. The read ready status bit will come true upon appearance on the tape of the first valid character. Upon appearance of an invalid sync code, the inter-record gap status bit comes true and tape motion is automatically stopped. Note that this will happen only after at least one valid character has been found. Once the read ready status bit comes true, the character must be taken within 2.8 milliseconds or it will be overwritten with the next one. The tape read hardware double-buffers incoming characters to allow the 2.8 msec character availability.
- e. EX BSP is similar to EX RBK except that tape motion is in the reverse direction so the data bits will be reversed.
- f. EX SF is similar to EX RBK except the tape is not stopped upon appearance of an inter-record gap, and if allowed to continue will start to read the next record on the tape. In this case, the read ready status bit will come true again after the first character of the next record is read. Only an EX TSTOP will stop the motion initiated by EX SF.
- g. EX WBK causes the currently selected deck to be set in forward motion and for all status bits except the write ready to go false. A character must then be presented within 2.8 milliseconds (the first character will be accepted at once due to the buffering in the tape hardware and then there will be a pause while the tape comes up to speed), at which time the write ready will go false until the writing circuitry is ready to accept another character. An end of record is signalled to the hardware by withholding a character for a period of time longer than 2.8 milliseconds specified above. When this is done, the write ready will go false, an inter-record gap will be written, the tape motion will cease, and the deck ready status bit will come true again.
- h. EX REWIND causes the tape to be rewound to the beginning on the selected deck. Worst case rewind time is approximately 40 seconds.
- i. PUNCH TABS, on the Cassette Cartridge are used for "write protect" and "automatic restart". The punch tab on the left (as you face the terminal) inhibits the ability to write on tape, when punched. When the tab on the right is punched, it causes an automatic restart whenever a halt or power-up occurs.

TABLE 7-1
TAPE UNIT PHYSICAL SPECIFICATIONS

Density	47 characters/inch
Speed	7.5 ips
Recording Rate	350 c.p.s.
Capacity	130,000 characters (typical)
Start/Stop Time (Inter-Record Gap)	280 msec.
Start/Stop Distance (Inter-Record Gap)	2 inches
Rewind Speed	90 ips
Rewind Time (max 300 ft.)	40 sec.
Character Transfer Time	2.8 msec.

PART 8

COMMUNICATIONS ADAPTOR

8.1 GENERAL DESCRIPTION

The 2200 Communications Adaptor is an external device, which when connected to the Datapoint 2200 Input/Output System permits asynchronous serial data interchange to other remote systems or devices.

The Communications Adaptor consists of three basic parts:

- a. The serial data transmitter and time base;
- b. The serial data receiver and time base; and
- c. The communications channel interface.

The communications channel interface may be one of four types:

- a. An EIA RS-232 type interface;
- b. An isolated high-level neutral or polar telegraph loop interface;
- c. A modem compatible with the Bell System 103 type modems;
- d. A modem compatible with the Bell System 202 type modems.

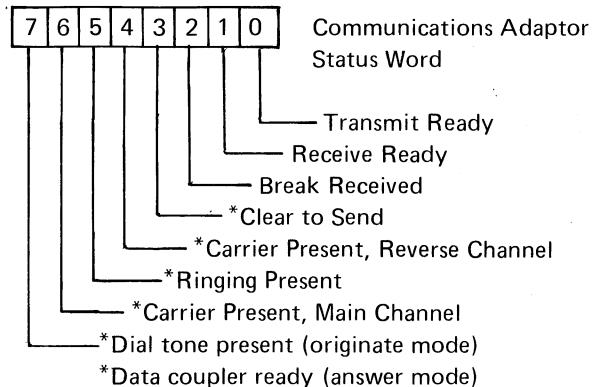
8.2 OPERATION

The serial data transmitter and receiver are addressed at the same time (the address of the first used communications adaptor is 3228 - see Table 4-2). Additional adaptors may be given previously unassigned addresses.

To set the bit rate desired for the transmitter time base two successive EX COM3 instructions are used to transfer two 8-bit masks from the A-register (See paragraph 8.6 for a discussion of time base mask words). For the receiver EX COM2 is used.

To set the character length for the transmitter and receiver an EX COM4 command is executed with a character length mask from the A-register (see paragraph 8.7 for a discussion of character length mask words).

The status of the communications adaptor is transmitted to the A-register with the following bit assignments:



*Used with data set options.

Communications Adaptor Status Bit Description

Bit 0, Transmit Ready

The "true" condition of this bit indicates that the serial transmitter is ready to accept a new character for transmission. Should another write command be issued to the Communications Adaptor while this bit is "false", i.e. transmitter NOT ready, the previous character will be written over.

Bit 1, Receiver Ready

The Receive Ready bit, in the true state, indicates the presence of a new received character. A read command to the Communications Adaptor returns this bit to the false state. If a read command is not issued before another new character is received, the new character will replace the existing character and the status will remain true.

Bit 2, Break Received

The Break Received status bit simply indicates that the received data is in the "space" or "zero" condition for longer than one character time.

Bit 3, Clear to Send

The true state of Clear to Send status indicates that the data set (internal or external) is prepared to accept data for transmission. This bit has meaning only when an internal or external data set is in used.

Bit 4, Carrier Present - Reverse Channel

This status bit has significance only when operating half-duplex with either an internal or external 202 type data set (modem). The true condition indicates that the reverse (supervisory) channel carrier is being received.

Bit 5, Ringing Present

The true condition of Ringing Present indicates that the ringing of an incoming call has been detected. This bit has significance only when used with an internal or external (with proper options) data set.

Bit 6, Carrier Present - Main Channel

The true condition of this status bit indicates that the main channel carrier is being received. This status bit has meaning only when used with an internal or external data set.

Bit 7, (1) Dial Tone Present (Originate Mode)

(2) Data Coupler Ready (Answer Mode)

(1) When originating a call, the true condition of this status bit indicates that a dial tone is present and dialing may proceed; during dialing, the status will become false. Following dialing, and a 2 to 5 second delay, this bit will return to the true condition indicating connection to the telecommunication network (but does not indicate the called number has answered).

(2) When answering a call, the true condition of this status bit indicates that the data coupler is connected to the telecommunications network.

8.3 DATA OUTPUT

After addressing the communications adaptor transmission of each character is accomplished in the following manner:

- a. Input the status word and verify that status bit 0, Transmit Ready, is set to 1 indicating that the adaptor can accept another character.
- b. Load the A-register with the byte to be transmitted.
- c. Apply a write strobe (EX WRITE). Data present on the A-bus will be loaded into the data transmitter and data will be serially transmitted at the selected code length and bit rate.

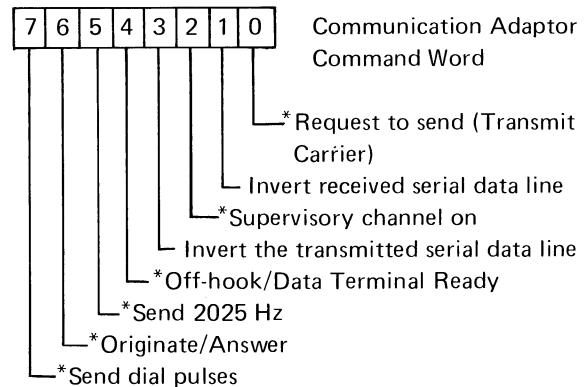
8.4 DATA INPUT

After addressing the communications adaptor, reception of each character is accomplished in the following manner:

- a. Input the status word and verify that status bit 1, Receiver Ready, is set to 1, indicating that a character has been received.
- b. Execute an EX DATA instruction.
- c. Execute an INPUT command, transferring the received character to the A-register.

8.5 COMMAND WORD

Control of the communications adaptor is accomplished through the use of a command word. The command word is transmitted to the adaptor by executing EX COM1.



*Used with data set options

Communication Adaptor Command Word Description

Bit 0 - Request to Send

This command bit controls the transmit carrier of an internal or external data set. A "one" in this position turns on the transmit carrier and indicates to the data set that it must prepare for data transmission.

Bit 1 - Invert Received Serial Data Line

A "one" in this position permits data to be received normally when the received serial data line is inverted.

Bit 2 - Supervisory Channel On

This command is used only with a 202 type modem in half-duplex operation. A "one" in this command indicates to the modem that the supervisory (or reverse) channel will be operative, transmit or receive.

Bit 3 - Invert Transmitted Serial Data

A "one" in this command inverts the transmitted serial data.

Bit 4 - Off-Hook

A "one" must be placed in this bit position any time a telecommunication call is to be originated or answered. This command allows connection to be made to the telecommunication network with an internal modem and a Bell System Data Access Arrangement. When using an external modem, this command provides "Data Terminal Ready" to the external modem, i.e., the system is prepared for on-line communications. This command is used only for the cases described above.

Bit 5 - Send 2025

This command is used only with an internal 202 type modem, half-duplex operation and "answer" mode. The only use of this command is described as follows:

- 1.) following receipt of Ringing Present, Status Bit 5, the Off Hook Command, Command Bit 4, is set to a "one".
- 2.) Next, Status Bit 7, Data Coupler Ready must become "true".
- 3.) Send 2025 command must now be set to a "one" only for a period of 1/2 second to 3 seconds to inform the calling data set of our response.

Bit 6 - Originate

This command is used only with internal data sets (modems). A "one" in this command instructs the modem that the system will originate a telecommunication call, A "zero" tells the modem the system is prepared to answer a telecommunication call.

Bit 7 - Send Dial Pulses

This command is used only with internal data sets (modems) and is set to "one" only when dialing. Its use is described as follows:

1.) Off-Hook Command (Bit 4) is set to "one".

2.) Status Bit 7 - Dial Tone Present becomes "true".

3.) Bit 7, Bit 4 and Bit 3 (invert xmit), are now set to "one".

4.) When the last dial pulse is completely transmitted, Bit 7 and Bit 3 must be returned to "zero".

8.6 TIME BASE MASK WORDS

Both time base generators are programmed for their respective bit rates by the processor. Each time base is independently controlled to allow transmission and reception at different rates.

After addressing the communication interface, two eight-bit mask words are loaded into the time base registers to synthesize the selected bit rates. As each respective byte is presented, a corresponding EX COM2 instruction must be executed to load the receive time base and an EX COM3 instruction to load the transmit time base.

These two bytes are combined to form a 16 bit word which is placed in a holding register. A counter is then set to the value in the holding register. This counter is incremented at the rate of 153,600 Hz. Each time the counter overflows, i.e., goes from all one to all zeroes, a pulse is generated and the counter is reset to the value in the holding register. The time between pulses represents 1/2 clock period or 1/2 bit time. Given a bit rate (bps), the following formula can be used to determine the number N to be entered into the holding register:

$$N = 65,536 \cdot \left(\frac{76,800}{\text{bps}} \right)$$

This number N may then be converted to a 16 bit binary number and separated into the two 8-bit mask words.

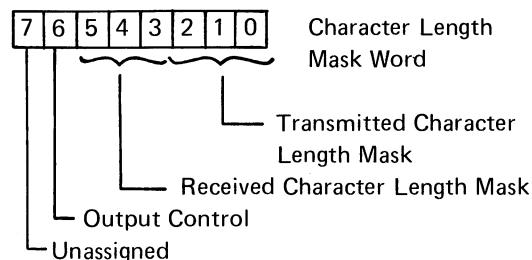
The octal codes for some of the more frequently used rates are listed below:

BIT RATE	1ST MASK WORD	2ND MASK WORD
1) 100*	375	000
2) 110	375	106
3) 220	376	243
4) 440	377	121
5) 150	376	000
6) 300	377	000
7) 600	377	200
8) 1200	377	300
9) 2400	377	340
10) 4800	377	360
11) 9600	377	370

*(Dialing)

8.7 CHARACTER LENGTHS

Character lengths for the transmit and receive sections and its output control bit are determined by a character length mask word which is transmitted to the communications adaptor with an EX COM4 instruction.



The functions of the mask bits are given in the following tables:

TABLE 8-1
TRANSMITTED CHARACTER LENGTH MASK BITS

MASK BIT POSITION 210	START UNITS	INFORMATION UNITS	STOP UNITS	CODE BIT POSITIONS 76543210
000	1	8	1	87654321
001	1	8	2	87654321
010	1	7	1	7654321
011	1	6	1	654321
100	1	5	1	54321
101	—	—	—	—
110	—	—	—	—
111	—	—	—	—

When codes having 5, 6, or 7 information units are to be transmitted, the remaining high-order bits in the character byte must be coded to "1".

When a two-unit stop pulse is required for characters having 5, 6, or 7 information bits, the next larger character length is used; the remaining high-order bits (all coded 1) form the stop pulses.

When received characters contain 5, 6 or 7 information bits, the remaining low-order bits (as shown above) must be disregarded.

One additional command bit, Bit 6 (output control) of the Character Length Mask Word, is used to control the EIA RS-232 Transmitted Data and the High-Level Keyer Transmitted Data. A "one" in this command bit enables serial data to be transmitted only to the EIA RS-232 output or to

TABLE 8-2
RECEIVED CHARACTER LENGTH MASK BITS

MASK BIT POSITION 543	START UNITS	INFORMATION UNITS	STOP UNITS	CODE BIT POSITIONS 76543210
000	1	8	1 or more	87654321
001	1	8	1 or more	87654321
010	1	7	1 or more	7654321x
011	1	6	1 or more	654321xx
100	1	5	1 or more	54321xxx
101	—	—	—	—
110	—	—	—	—
111	—	—	—	—

the High Level Keyer. A "zero" in this command bit allows serial data to be transmitted only to an internal data set (modem).

8.8 INTERFACE CONNECTOR

This interface is provided through an Amphenol 17-10500-1 connector. Pin assignments are as follows:

LEAD	FUNCTION	INPUT/OUTPUT
------	----------	--------------

1	Protective Ground	-----
2	Protective Ground	-----
3	OH (Off Hook)	Output
4	+25v	-----
5	DA (Transmission Path Request)	Output
6	R (Ring Indicator)	Input
7	CCT (Data Coupler Ready)	Input
9	DT (4 wire)	Direct Private Line Connection
10	DT (2 wire)	
11	DR (2 wire)	
12	DR (4 wire)	
23	Clear to Send (RS-232)	Input
24	Transmitted Data (RS-232)	Output
28	Signal Ground	-----
29	Signal Ground	-----
32	+5v	-----
33	+5v	-----
40	Request to Send (RS-232)	Output
41	Received Data (RS-232)	Input
42	Data Terminal Ready (RS-232)	Output

LEAD FUNCTION INPUT/OUTPUT

44	Supervisory Transmitted Data (RS-232)	Output
45	Data Carrier Detector (RS-232)	Input
46	Supervisory Received Data (RS-232)	Input
49	Clock for 3300P	Output
50	Transmit Bit Rate Clock	Output

8.9 HIGH LEVEL OPTION

Interface with telegraph-type current loops is provided with the high level option. This option provides for completely isolated electronic neutral/polar output relay and a completely isolated neutral/polar input relay. Loop voltage may be as high as 400 volts across the relay and as high as 1000 with respect to ground.

Loop resistance and power is not included with the option.

For further information, refer to the Datapoint 2200 Installation Manual.

8.10 103-DATA SET OPTION CHARACTERISTICS

The 103-Data Set option provides for full duplex data transmission for rates up to 300 bits per second with a signalling system that is compatible with the Bell System 103 series Dataphones. Connection to the common carrier lines would normally be made through a Bell System Access Arrangement type F-58118, CBT, or 1001B. Other connections are also possible where automatic dialing or answering is not required.

The data set may be placed in either the answer mode or originator mode through the use of bit 6 of the communications adaptor command word (see paragraph 8.5). Bit 6 is set to 0 for answer mode and 1 for originator mode. The request to send command bit (bit 0) is normally set to 1 with the 103 option to maintain the transmit carrier on.

Operation of the automatic dialing and answering features discussed in paragraph 8.12 and 8.13.

Table 8-3 provides a summary of characteristics of the 103 Data Set option.

TABLE 8-3

103 DATA SET OPTION CHARACTERISTICS

Originate Mode

Carrier Frequencies:	Transmit:	Mark: 1270 Hz Space: 1070 Hz
	Receive:	Mark: 2225 Hz Space: 2025 Hz
Answer Mode		
Carrier Frequencies:	Transmit:	Mark: 2225 Hz Space: 2025 Hz
	Receive:	Mark: 1270 Hz Space: 1070 Hz
Keying Rate:	Up to 300 bits per second	
Transmit Level:	0 to -10 dbm.	
Impedance:	600 ohms nominal	
Receive Sensitivity:	+5 to -30 dbm.	

8.11 202-DATA SET OPTION

The 202 Data Set option provides for either full or half duplex data transmission for rates up to 1200 bits per second (1800 bits per second on conditioned private lines). This option is compatible with Bell System 202 series Dataphones (including supervisory channel operation) and in addition provides a 150 bit per second supervisory channel when used with another Datapoint 2200 Data Set option of the same type. Connection may be directly to private lines or to common carrier lines through a Bell System Access Assignment type F-58118, CBT, or 1001B where access to the telephone switched network is desired.

Operation of the automatic dialing and answering are discussed in paragraphs 8.12 and 8.13 respectively.

Table 8-4 provides a summary of characteristics of the 202 Data Set option.

TABLE 8-4

202 DATA SET OPTION CHARACTERISTICS

Main Channel Frequencies:	Mark: 1200 Hz Space: 2200 Hz Soft Turn-Off: 880 Hz
Supervisory Channel Frequencies:	Mark: 387 Hz Space: 470 Hz Soft Turn-Off: 330 Hz
Special Command Frequency:	2025 Hz
Main Channel Keying Rate:	Up to 1200 baud (1800 baud on conditioned private lines.)
Supervisory Channel Keying Rate:	Up to 150 baud
Transmit Level:	0 to -10 dbm
Impedance:	600 ohms nominal
Receive Sensitivity:	+5 to -30 dbm

8.12 AUTOMATIC DIALING OPERATION

When using the Datapoint 103 or 202 data set options with the Bell System Access Arrangement type F-58118, CBT, or 1001B it is possible to automatically originate a call into the telephone switches network. The procedure for this function is as follows:

- a. Set bits 4 and 6 of the communications adaptor command word to 1 to provide an off-hook signal to the telephone network and to prepare the modem for originate operation.
- b. Test bit 7 of the communications adaptor status word for a 1 indicating dial tone present.
- c. Set the transmitter time base to 100 bits per second (see paragraph 8.6).
- d. Set the character length mask word to all zeros (ten bit length-see paragraph 8.7).
- e. Set bits 3 and 7 of the command word to 1 thus inverting the serial transmitter output and transferring this output to the dial pulse keyer.
- f. Sequentially transmit the octal byte 360 for each dial pulse required for each number (see paragraph 8.3-Data Output).
- g. Program approximately 1 second delay between each number and at the end of the last number transmitted.

-
-
-
-
-
-
-
- h. Re-establish the correct code length and bit rate for data transmission and set command word bits 3 and 7 to zero to restore the normal transmitter output.

8.13 AUTOMATIC ANSWERING OPERATION

When using the Datapoint 103 or 202 data set options with the Bell System Access Arrangement type F-58118, CBT, or 1001B it is possible to automatically answer a call from the telephone switched network.

Ringing is detected simply by testing bit 5 of the Communications Adaptor Status Word, Response to ringing would be to set bit 4 of the Communications Adaptor Command Word to 1 to provide an off-hook signal to the telephone network.

If the 103 Option is used Command Word bit 0 is set to 1 and bit 6 is set to 0 turning on the transmit carrier and selecting the answer-mode carrier frequencies.

If the 202 Option is used bit 5 of the command word is set to 1 for 1/2 to 3 seconds to transmit a 2025 Hz tone to disable echo suppressors and to inform the calling data set of out sequence in the telephone network, after which normal data transmissions occurs.

SECTION 2

DATAPoint 2200

OPERATING SYSTEM

SECTION 2

THE OPERATING SYSTEM

The operating system is a conversational mode program for the Datapoint 2200 user to have a means to catalog, load, debug and run user programs and to provide other utilities important to the use of the 2200. All other programs discussed in the "Programmers Manual" such as the Program Editor and the Assembler are programs which the user may catalog onto an operating system tape and call into use as required.

The operating system itself is a relatively long program which is generally overlayed when user programs are called in from tape (unless they are less than 2K bytes in length and properly located). However, a family of resident utility routines is loaded with the operating system that may be used by user programs to simplify frequently used functions such as reading from the keyboard, writing to the CRT screen, reading and writing tape records, etc. The detailed use of the routines and the makeup of the operating system are described in Section 5—Advanced operating system command and subroutine usage.

Section 1 will describe the command language for the operating system and does not require any particular programming skills.

Start-Up Procedures

When power is first applied to the Datapoint 2200, it is incapable of performing any useful function except to load a block of data from the rear tape cassette deck into the processor's memory and transferring control to it. In the operating system, this first block of data is called a LOADER and when control is transferred to it, it proceeds to the first check itself to see if it was loaded properly and then to load the next file on the same tape which is the rest of the operating system program. This process can be executed at any time (assuming a proper program tape is in the rear deck) by pressing the RESTART key on the right hand side of the 2200.

This first block of data can be used to load programs other than the operating system and is generally useful for all applications of the 2200. In order to use the operating system, a full 8K bytes of memory must be provided in the 2200, but the loader alone can be used with any size memory.

When an operating system program tape is loaded, the first thing that appears on the screen is:

COMPUTER TERMINAL OPERATING SYSTEM

READY

At this point, any of the operating system commands discussed below may be typed into the 2200.

Each command has to be in the form of a word followed by the ENTER key, or a word followed by a space, or a dash and a name, or a modifier. Each of the following are valid operating system commands:

CATALOG
REPLACE RST4
REPLACE-RST4
RUN*

Only the first three letters of a command are actually decoded so that the following are valid commands:

CAT
REP RST4

If an invalid command is typed the system responds with:

WHAT?

Operating System Commands:

CATALOG

The CATALOG command will print out a list of programs currently available on that particular operating system tape. Up to 14 programs may be cataloged on a tape under this system and the order that they appear on the tape is the order their names appear on the screen when the catalog command is given. A typical catalog response might be:

CATALOG
RST4 CODER BANDIT ANNUIT
READY

Which would indicate that four programs are logged onto the tape.

NAME

The NAME command allows any program in the system to have its name changed. For example, the program named BANDIT in the above example can be changed to GAME with the command:

NAME-BANDIT, GAME

Names may have any combination of letters and digits up to six characters in length and beginning with a letter. All of the following are valid program names:

BOB
R12345
A
NAME

RUN

The RUN command causes the operating system to position the cassette tape to the program named in the run command, load the program into the 2200 memory and transfer control to it. The program being run may overlay part or all of the operating system. If it does, returning to the operating system can only be done by reloading it. This can be done by using the restart switch or by program control (See Section 5).

A run command would appear like this:

RUN-BANDIT

A program that has not been logged onto the system tape may be run by placing an assembled form of the program on the front deck and typing:

RUN*

IN

The IN command causes a program to be cataloged onto the operating system tape. The program must be assembled and the assembled program tape placed on the front deck. The IN command is typed giving the name to be assigned to the program as shown here:

IN RST4

The operating system tape (on the rear deck) will position itself to the end of its program library and will then copy the program from the front deck into place and add its name and position to the system catalog.

DELETE

The DELETE command causes the program named in the command to be removed from the system library. Unless the program being deleted happens to be the last program in the library, a SCRATCH tape will be required in the front deck to copy part of the library out and back to the system tape to CLOSE-UP the space. In this case, when the command is entered, the system will write a message back:

FRONT TAPE SCRATCH?

Then the processor will stop. The stop key on the right hand side of the keyboard will be lighted. If there is a tape in the front deck that can be recorded on, press the run key on the right hand side and the system will proceed to delete the named program for you.

This will generally take a little time. When the system is through it will write READY on the screen.

REPLACE

The REPLACE command allows a program already in the system catalog to be deleted and a new program to be put in its place in the same order on the tape. The new program does not have to be the same length as the old one. Again, this command takes time to execute due to the amount of "shuffling" of tape files to get everything in place.

AUTO

The AUTO command allows one of the programs in the system library to be marked for an automatic RUN whenever the operating system is restarted. Once a program has been named in an auto command, it may be cleared by typing a MANUAL command. The automatic feature may be overridden during a restart by holding down the KEYBOARD key on the right hand side of the keyboard.

The automatic program calling feature is particularly valuable when the program is to be run in an unattended situation. If the knock-out tab on the back of the operating system tape is removed then whenever power is reapplied or the processor is halted for any reason (including a programming halt), an automatic restart is executed and, of course, the program named in the auto command is reloaded and given control.

If the command AUTO is typed without a name then the system will respond with:

NAME REQUIRED

If a program is already named in an auto command then the system will respond with:

AUTO SET TO (PROGRAM NAME)

MANUAL

The MANUAL command will delete any program from the auto mode.

OUT

The OUT command causes any named program to be copied to the front tape causing any data already on the front tape to be lost. This copy may then be cataloged onto some other

system tape or be saved for some future use. When using the OUT command the system will write FRONT TAPE SCRATCH? onto the screen and the processor will halt. If you have a usable tape in the front deck, then press the run key on the right hand side of the keyboard and the system will continue.

If the command OUT \$ is typed, then the entire operating system including the library is copied to the front deck. If the command OUT * is typed, then the loader and the library is copied to the front deck but the operating system is deleted. This permits a program or family of programs to be used on a Datapoint 2200 with less than the full 8K of memory. (See Section 5 for details).

PREP

The PREP command causes the tape in the front deck to be rewound and a NULL file to be written at the beginning of the tape, effectively ERASING the tape and making it ready for use by operating system tape routines. The operating system commands that write on the front deck execute the PREP function automatically, however, USER PROGRAMS WRITING DATA TO THE TAPE MAY REQUIRE THE TAPE TO BE "PREPed" IN ADVANCE.

HEX

The HEX command allows programs generated on other machines that follow a specified hexadecimal format, to be loaded into the Datapoint 2200. Users will not normally be concerned with the HEX command. (See Section 5 for details).

DEBUG

The DEBUG command transfers program control to a small sub-program within the operating system that is used as a programming aid to debug and modify programs that are loaded into the 2200 memory. The debug sub-program allows you to write the contents of memory locations to the screen, modify memory locations, load programs into memory from the library or from the front deck, transfer control to parts of a program in memory and TRAP register values upon return to the debug program. Instructions on the use of the debug sub-program are given in Section 5.

SECTION 3

DATAPOINT 2200

SOURCE CODE EDITOR

SECTION 3

ASSEMBLER SOURCE CODE EDITOR

The assembler source code editor program provides for the preparation and editing of source data tapes in an assembler compatible format.

The editor program is called by the Computer Terminal Operating System (CTOS), if it is cataloged therein, by typing the following command:

STEP 1

RUN EDIT

STEP 2

- a) When the editor has been loaded, the following message will appear on the screen:

Compressed Source Code Editor

Edit (E) or Convert (C)?

Type 'C' only if you have a tape generated by EDIT (1.1). This version of the Editor generated unblocked string records (using SSFW\$). The present version generates "compressed source" records (using SNFW\$).

If conversion is required, type 'C'. The following message will appear:

**PLACE SYMBOLIC TAPE IN FRONT DECK-
WHEN READY PRESS RUN**

Place the old tape in the front deck and push RUN. A converted file will be generated in the scratch area of the CTOS tape. When the 1st pass of the conversion is complete, the following message will appear:

**PLACE SCRATCH TAPE IN FRONT DECK-
WHEN READY PRESS RUN**

To protect your original tape, use a new tape to record the new compressed source code. Place the new tape in the front deck and press run.

The conversion process may be repeated for several tapes.

The Editor will now accept only compressed source tapes.

- b) Type 'E' to EDIT, this message will appear on screen:

TYPE (:NEW,:OLD, OR:DUPLICATE)?

The meaning of the possible responses are as follows:

(NOTE:) All commands to the editor must begin with colons.

:NEW	Indicates that the tape on the front deck is to be treated as a new source data tape. Any old data on this tape will be written over by the editor program.
:OLD	Indicates that the front deck contains a tape with assembler source data on it. The operator will be allowed to edit this tape, changing only those lines which the operator specifies.
:DUPLICATE	Indicates to the editor that it should copy the contents of the scratch file to the source data tape. This provides copies of a single source data file. It also provides for recovery capability should a system failure occur during the editing process. See Step 4, Recovery Procedures.

STEP 3

- a) If the response in Step 2 was :NEW the following question will appear on the screen:

NEW NAME?

The operator may now enter a character string of up to 40 characters. This text will become the first record of the source data tape and will appear as part of a comment line in the assembled text.

- b) If the response in Step 2 was :OLD the editor will read the header record from the source data tape (front deck) and display that header on the screen in the following format:

OLD NAME IS XXXX XXXX...XXXX

This old name header is also written to the scratch file and is retained as the header record for the source tape.

- c) If the response in Step 2b was :DUPLICATE the editor will return to Step 2b after the copying operation has been completed. See Step 4, Recovery Procedures.

After a) or b) above has been completed, a "READY" message will be written on the screen and the cursor will appear at the beginning of the bottom line. The editor is now ready to accept new text data or a command. In order to enter text, simply type the desired text. Upon pressing the enter key, the typed line will be rolled up on the screen one line and the cursor will reappear on the bottom line and accept another text line or a command. When a line of text rolls off

the top of the screen, it is written to the scratch file. Lines will be written to the scratch file in the same order as they appear on the screen, the top line being first in the file.

Commands which can be entered, and their respective functions, are listed below. Command lines are distinguished from text lines by a leading colon; therefore, it is necessary to begin any command by typing a leading colon.

The "pointed line" referred to by some command descriptions below, is the line currently being pointed to by the visible pointer at the left side of the screen (col. 0). The "point's" vertical position is controlled by the keyboard/display keys. Specifically, pressing KEYBOARD causes the pointer to move up one line, pressing DISPLAY moves it down one line. Motion in either direction is circular around the screen. (It wraps around).

Commands may be entered from the bottom line only and must be preceded by a colon (:).

The functions available and their respective descriptions are listed below:

MANUAL SEARCH	(Not a typed command). MANUAL SEARCH is like a continuing find of the very next line. That is, searching continues line at a time, with the next new line going to the eleventh line and the screen rolling up. It is unlike a find in that the screen isn't cleared for each new line acquired. It is useful for manually scanning through the data to bring to the screen and, therefore, into a position to edit the data of interest. MANUAL SEARCH is achieved by holding down the KEYBOARD and DISPLAY keys simultaneously. While held down, the search will proceed until end of file at which point the keys will become inoperable.	:COPY or :COP	The pointed line is copied to the bottom line and is simultaneously deleted from its previous location. The cursor will occupy the pointed line and will accept new text at this position. Striking the <enter> key returns the cursor to the bottom line and rolls the screen up one line.
:FIND <TEXT>	Where <text> represents n characters of text data. The editor searches the source data tape for the first match with the desired n-<text> characters. When the desired line is found it is displayed on the bottom line of the screen. The search is circular through the data files. If no match is found the text which occupied the bottom line at the time the command was issued is restored. Leading blanks on data-lines are ignored during the search. A FIND or EOF in progress can be stopped by the manual search operation.	:DELETE or :DEL	The pointed line is deleted from the screen and the scratch area. The cursor will occupy the pointed line and will accept new text at this position. Striking the <enter> key returns the cursor to the bottom line. The screen is not rolled up.
		:INSERT or :INS	The pointed line and all lines above it are rolled up one line. The cursor will now occupy the blank line created and is ready to accept the new text. Striking the <enter> key returns the cursor to the bottom line.
		:EOF	The editor will search the source data tape for an end-of-file. Upon finding it, the last 11 text lines are on the screen and the cursor occupies the bottom line ready to accept new text or commands. A FIND or EOF in progress can be stopped by the manual search operation.
		:END	Causes the screen and the remaining source data to be copied to the scratch file. The scratch file is then copied to the source data tape.
		:END/DEL	The same as :END except that all data on the source tape which follows the current screen data is deleted.
		:SCRATCH or :SCR	Cause all lines between the top of the screen and the pointer, inclusive, to be deleted from the screen and the scratch area. The cursor will occupy the pointed line and will accept new text at this position. Striking the <enter> key returns the cursor to the bottom line. The screen is not rolled up.

STEP 4

Recovery procedures using the :DUPLICATE command can be implemented should the edit program be aborted during its execution without benefit of having completed all necessary copying and end of file writing.

Causes of difficulties which could require such action are:

1. Power failure during execution of the program
2. Turning off the power execution of the program
3. Striking the restart switch during execution
4. Encountering unrecoverable tape errors during execution
5. Removal of the tape cassette during execution

The EDITOR edits from one tape [the "Source Tape," containing old data] through the screen to the "scratch tape." The identity of the "source" and "scratch" tapes or files are reversed each time the current "source" tape reaches a file marker indicating end of file. Therefore, if an even number of passes have been completed, the updated file is on the front deck. If an odd number of passes have been made, the CTOS scratch file, file 40g, is the updated file. New programs are written on the front deck.

Understanding the activity of the data with respect to the tapes will allow the operator to determine the course of action should difficulty occur.

Should the operator encounter difficulty and be rather vague as to which action to take, it is recommended that the operator choose the most valuable data tape by the following:

When unrecoverable difficulty arises, remove the data tape and replace it with a new tape which will be written over by the scratch area of the operating system tape.

Run the edit program and type :DUPLICATE in response to the original questions.

The scratch area will be deposited on the new tape.

Then the operator can manually search through the two tapes as separate data tapes and make a decision as to the most valuable to keep.

Special situations:

It is possible that the scratch area has the most valuable data on it but it is missing any file termination record. This can occur when the program was interrupted during data entry and the previous scratch-to-data tape copy (if any) was so far back as to render the scratch area the only desirable data. Should this occur, the copy will proceed from scratch to source until it runs out of data, in which case the operating system will encounter garbage on the tape. It will then write end-of-file marks on the data tape at that point. This will give a clean data tape suitable for beginning again. A note of caution: it is usually advisable, when attempting to recover valuable data, to use the two data tape approach and visually compare the two to make a value judgment as to which has the most desirable data.

SECTION 4

DATAPOINT 2200

ASSEMBLER

SECTION 4

THE ASSEMBLER

The 2200 assembly system consists of the ASSEMBLER, EDITOR and the OPERATING SYSTEM.

The ASSEMBLER generates a block of absolute object code which can be loaded by the operating system loader and cataloged by the operating system. It generates the object code from the symbolic source code which was generated by the editor.

The ASSEMBLER makes two passes over the source code.

The first pass generates a symbol table from the labels in the source code and checks for certain error conditions, primarily syntax and form. The symbol table is maintained in memory.

The second pass generates the program listing and the object code on the tape. It also produces further diagnostics of a more subtle nature.

Basically, the ASSEMBLER is a program that assigns numerical values to symbols and outputs these values upon input of the associated symbols. Symbols in certain fields have pre-assigned values such as the opcode mnemonics. The value assigned to an instruction mnemonic is the binary bit configuration recognized by the 2200 processor for that instruction.

For example, the following instruction mnemonics have the following octal values:

MNEMONIC	VALUE
ADB	201
RETURN	007
SUB	221

Symbols in fields other than the opcode field may be defined by the user. Pre-defined and user-defined symbols are kept separately by the ASSEMBLER so that the user may define symbols that are the same as the pre-defined symbols without encountering any difficulties.

Along with relating symbols with numbers, another major function of the ASSEMBLER is to enable one to reference a symbol that is defined later in the program. This is called FORWARD REFERENCING, and may be handled in a variety of ways. One of the simplest is to look at the source code twice. The first time determines the definitions of all the symbols and the second time uses the symbols to produce the object code. Each "look" at the source code is called a "PASS". Therefore, we end up with a two pass assembly process.

Statements

A 2200 assembly code statement consists of a label field, an instruction field, an expression field and a comment field. An example:

1	2	3	4
LABEL1	JTC	START	THIS IS THE COMMENT FIELD

Field 1 is the label field

Field 2 is the instruction field

Field 3 is the expression or operand field

Field 4 is the comment field

The 2200 editor provides automatic formatting so that the fields always are justified to begin in a certain column with tabbing to that field automatic. However, the ASSEMBLER only requires the following:

A non-space in the first column means that the first field is a label, except for a leading period which designates the entire line as a comment line.

A space in the first field means a null label and the first field is an instruction.

Scanning proceeds from left to right with one or more spaces serving as field delimiters.

Terminating fields by other than a space or a line termination will result in E-flags during the assembly.

THE LABEL FIELD may consist of up to 6 characters. An excess of 6 characters will be truncated. The first character may be any alphabetic character or a \$ sign. The other characters may be any alphanumeric character or a \$ sign. For example:

LEGAL	ILLEGAL
LABEL1	1LABEL (starts with numeric)
LABEL2	LABEL* (non-alphanumeric character)
LABEL\$	LABEL. (non-alphanumeric character)
L1B2L3	

THE INSTRUCTION FIELD may be any of the instruction mnemonics listed in the Datapoint 2200 Reference Manual, compound instruction (described later) or assembler directives.

The Instruction Field may be from two to four characters. However, only the first three are scanned and consequently the user may abbreviate. For example:

LEGAL	ILLEGAL	
CALL	CALL2 (instructions have no more than four characters or numeric characters in the field)	
JTZ		
SET		
TP		

THE EXPRESSION or OPERAND FIELD consists of any number of strings, numbers or symbols with operators between them. If a space or line end terminates a number or a symbol, the expression is assumed to be ended. Numbers are assumed to be decimal (base 10) unless they have one or more leading zeros, in which case they are taken to be octal. That is, 123 is 123 decimal, whereas 0123 or 00123 (the octal number 123) is really 83 decimal.

String quantities are delimited (preceded and followed) by apostrophes. In expressions, only the last character of a string is used if more than one appears. If a string were to be added to a number, only the last character of the string would be added. The character value is the ASCII binary number with the parity bit always a zero. A null string is legal ('') and results in a zero value. The forcing character, #, is used in strings to indicate that the next character should be taken as ASCII no matter what it is. This is useful for getting the characters (') and (#) themselves into the string. For example:

' #' #' is the character string '#

There are three operators allowed in the expressions:

1. + This means addition
2. - This means subtraction
3. >8 This means shift right by 8. Use this to get the MSP of an expression.

Expressions are evaluated from left to right and all operations are assumed to have the same priority.

The operand or expression is a symbolic expression which is evaluated at assembly time and the value is used in whatever manner is required by the opcode.

THE COMMENT FIELD begins immediately after the first delimiter space after the operand. The comment field may have any character including punctuation within it. It is terminated by the end of the line which was written by the editor. Comments may take over the entire line, in which case that line must begin with a period.

ASSEMBLER DIRECTIVES are available for setting label and location counter values to other than the normal sequential location assignments and for defining constants.

There are seven:

1. EQU EQUALS. Sets the value of the label on the statement to the value of the operand expression.
2. SET SET. Changes the value of the location counter to the value of the operand expression.
3. SK SKIP. Increments the value of the location counter by the value of the operand expression.
4. TP TABULATE PAGE. Increments the value of the location counter until it is a multiple of 256. This is useful for minimizing execution time and for blocking out data areas addressable by single precision.
5. DC DEFINE CONSTANT. Generates eight bit object words from one or more expressions or strings following the opcode. If the expression is terminated by a space, the DC directive returns control to the main assembly process loop which obtains another instruction. If it is terminated by a comma, another expression or string is looked for. Another special exception is made for string items found in the DC directive. All the characters of a string item are significant and as many words as necessary are generated to accomodate all the characters of the given string. Again, a comma is looked for after the closing apostrophe in a string item to see if more expressions follow. This special string item is in effect only if the expression opened with an apostrophe. String items in expressions still have only one character of significance. For example:
DC 1,2+3,2+'A','ABC'
generates the following octal values:
1,5,103,101,102,103
6. DA DEFINE ADDRESS. Generates a two byte constant which is the address, LSP first, of the expression.
7. RP REPEAT. Will cause the following line to be processed, the number of times, indicated by the operand value. For example:
RP 5
LDA 0123
would produce the same code as:

LDA 0123
LDA 0123
LDA 0123
LDA 0123
LDA 0123

NOTE

Repeated statements which have a label on them will result in multiple definition of that label and all that entails, including the "D" error flags.

FORWARD referencing in the expression field in assembler directives only is not permitted.

8. END END. Indicates that there is no more input data to be processed and that the ASSEMBLER should complete generating the output. The operand field has special significance in the END statement. The value of the expression in the operand of the END statement is the starting value of the execution of the program. That is the starting address. This is, of course, optional. When no operand is specified, the results are indeterminate. It should only be left vacant when the program is to be loaded without direct transfer of execution to the program such as an overlay.

Compound instructions are instructions which directly result in the assembly producing a sequence of source code. In this case, the 2200 ASSEMBLER has two: The HL instruction and the DE instruction:

1. HL LABEL The HL compound instruction generates the load H-REGISTER and load L-REGISTER instruction necessary to place the address of the label LABEL in the H-REGISTER and L-REGISTER properly so that the load to and from memory will operate to that address. In doing the HL, it loads the most significant byte of the value of LABEL into the H-REGISTER and the least into L.
2. DE LABEL The same as with HL except loads into the D and E registers.

THE ERROR FLAGS produced by the 2200 ASSEMBLER are as follows:

The error flags can occur during either pass of the ASSEMBLER in response to bad statements.

They are:

1. D The D flag means DIFFERENT DEFINITION. It is flagged if the label has been redefined to

a different value during the assembly. In that case, it has the second value.

2. I The I flag means INSTRUCTION MNEMONIC UNKNOWN. The instruction was not an accepted instruction in which case a zero is inserted for this instruction.

3. E The E flag means that an error has occurred in an expression or some unrecognizable character appeared in the wrong place. In this case a zero is substituted for the expression or in whatever was unrecognizable.

4. U The U flag means UNDEFINED LABEL. It is used whenever a label is referenced and is not defined. This can occur in pass 1 when an assembly directive is operating on an expression containing a forward reference.

EXTERNAL COMMANDS & REFERENCES can be taken care of in two ways:

1. Directly produce the numeric value in the expression field corresponding to the reference external address (such as an operating system subroutine resident in memory) or the external command operand such as EX 1 instead of writing EX ADR.
2. Equating labels to these referenced locations using the EQU assembler directive and then referencing the labels. This can be done for external references to operating system subroutines by duplicating the operating system subroutine entry point label in your program and equating it to that address. i.e. instead of:

CALL 017000

to get the operating system keyboard string input routine, a more meaningful listing can be obtained if, at the beginning of the program, this was entered:

KEYIN\$ EQU 017000

and then all references to this routine can be this way:

CALL KEYIN\$

The same is true of the external commands used in the 2200. Rather than say:

EX 1

it is more meaningful to say:

EX ADR

Since it is an external command address that is desired.

A list of the external commands and the operands which the ASSEMBLER incorporates into the proper EX coding are below.

The ASSEMBLER treats external command labels differently to produce the octal command byte. For the commands, the operands are as follows:

ADR	1
STATUS	2
DATA	3
WRITE	4
COM1	5
COM2	6
COM3	7
COM4	8
UNUSED	9
UNUSED	10
UNUSED	11
UNUSED	12
BEEP	13
CLICK	14
DECK 1	15
DECK 2	16
RBK	17
WBK	18
UNUSED	19
BSP	20
SF	21
SB	22
REWIND	23
TSTOP	24

It is recommended that for those external commands used, the EQU to the table number is done at the start of the source program and then the external command references are done to the label.

Operating The Assembler

The ASSEMBLER must have a symbolic source tape generated by the 2200 editor.

Place this tape in the front tape deck.

Run the ASSEMBLER.

It will ask for printer speed. For the Datapoint 3300P, state 300. For a model 33 or 35 Teletype, state 110. For a model 37 Teletype, state 150. For no printer or no listing desired, state 0.

The source deck will rewind and begin to read in.

At the end of the first pass the ASSEMBLER asks if the second pass should proceed. It only requires a YES or NO.

This is a convenience, since many times many errors will be uncovered by the ASSEMBLER already after the first pass and the user will desire to correct those errors before proceeding to the second pass and the listing.

If the second pass begins, the tape will rewind and begin accepting data again from the source tape, printing the listing and writing the object file on the scratch area of the rear tape.

When the tape has reached the end of the source the second time, assembly is complete and it only needs to copy the object code block on the rear tape to the area on the front tape just after the source code. This results in the rear tape being backspaced to the beginning of the block of code and then copying proceeding forward reading a block from the rear deck and writing it on the front deck.

At the end of ASSEMBLY, the operating system will be reloaded and come up running.

The front tape can be loaded into the machine to test using the operating system command RUN*, inputted into operating system catalog or loaded using the Debug program by using the F command.

SECTION 5

DATAPOINT 2200

ADVANCED OPERATING SYSTEM COMMAND and SUBROUTINE USAGE

SECTION 5

1. INTRODUCTION

The primary function of CTOS is to provide the user with an easily accessible data environment which will greatly facilitate program generation. This function is fulfilled through the use of a file handling system which is available both directly from the keyboard in the form of system commands and through program calls to file handling input/output subroutines. Note that the keyboard facility deals mainly with the system (rear) tape (using the data (front) tape mainly for input/output and scratch space) but that the program routines are generalized to allow use of either tape.

1.1 KEYBOARD FACILITIES

The keyboard accessible facility allows the user to fetch and execute object files, which may be either system packages, such as the editor and assembler, or files the user has generated with either the assembler or other code generating programs. This facility also allows the user to create new files, alter or delete old ones, or perform certain utility functions. The system tracks the files on the system tape in a symbolic catalog which may be manipulated by the operator at the keyboard or used in program linking.

1.2 PROGRAM FACILITIES

The program routines perform basic operations such as reading and writing records with all parity checking and generation handled for the user. Other operations such as positioning to the beginning or end of a file, backspacing over records, or rewinding the tape are also provided. Parameterization is handled in a generalized way to make subroutine usage easy and consistent.

1.3 PHYSICAL LAYOUT

The memory layout of the operating system is shown below. The OS FILE HANDLER is the program accessible facility mentioned above while the OS COMMAND HANDLER is the keyboard accessible facility. Note that only 017400 and up need be in memory if only the symbol linker (which calls in an overlay by name so that its physical file number may be changed without having to rewrite the program calling in the overlay) is to be used, only 016200 and up need be in memory if only the debugging tool is to be used, and only 014000 and up need to be in memory if the keyboard facilities are not to be used (of course, 0-0777 is always reserved by the system). Also note that the user may load a program designed to fit into a 2K machine without overlaying any part of the full operating system.

CTOS MEMORY USAGE MAP:

SYMBOLIC LINKER	017777
CATALOG	017600
KEYBOARD DISPLAY	017400
DEBUG	017000
	016200
OS FILE HANDLER	014000
OS BOOTBLOCK COPY	013000
OS COMMAND HANDLER	05000
2K UNUSED	01000
LOADER	0

2. THE LOADER

The loader is the heart of CTOS. It enables other programs to load files from the tapes into memory without the tape having to be at the beginning of the desired file and provides extensive error protection. It is the routine used by the bootstrap mechanism (indeed, it is part of the bootblock) to load the initial program and is also the routine used in overlay and linkup operations both by CTOS and utility packages.

2.1 BOOTSTRAP ACTIONS

When a restart occurs, the rear deck is rewound and the first block on the tape (called the bootblock) is loaded into memory starting at location zero. The first 512 bytes of memory (0 to 0777 octal) have been reserved for a permanently resident program which is loaded from the bootblock. The first 40 bytes of this block constitute a program which runs a parity check on the rest of the block that should have been loaded. The processor is halted (note auto-restart implications if the auto-restart tab on the cassette

is punched out) if this routine finds a fault in the check. Otherwise, zeros are stored in the memory locations used in the parity check routine. This will cause a halt if an early data drop-out from the tape machine occurs during the next bootstrap load (typically only one or two bytes get loaded in this failure mode). After the low memory has been cleared, a routine calls the loader, which has been loaded in the bootstrapping operation, asking for file zero to be loaded from the rear deck. If file zero cannot be loaded for some reason, the program halts the processor without a whimper (no bells or whistles in any of the bootstrap operation), otherwise, it jumps to the starting address supplied with file zero. Note that if the auto-restart tab is punched out of the rear cassette, any failure along the road of bootstrapping will cause the whole process to be tried again.

2.2 FILE ORGANIZATION

Once file zero has been loaded from the system tape, the bootstrap program (locations 0 through 074) is never used again until the next restart operation which will overlay it. The loader, however, will be used many times. The physical layout of information on the system tape is as follows:

BOOTBLOCK/FILE0(FILE1/.../FILE15/FILE32/FILE127

File 0 is the one executed by the bootstrap and is typically followed by a sequential (required to be sequential by the loader) set of minimally increasing (file numbers go up by only one at a time) files up to 15 (a CTOS catalog size limitation, although the loader will load a file with any positive number), followed by a file 32, which is a system scratch file, followed by a file 127 (largest positive eight bit number), which is a dummy to mark the logical end of the tape.

2.3 FILE LAYOUT AND RECORD FORMAT

Each file is a group of records starting with a very special four byte record. Every record used by CTOS starts with two special bytes to indicate that it is one of three types: file marker, numeric data, or symbolic data. The file marker, which is the special four-byte record at the beginning of a file, contains two additional bytes that denote the file number. The use of two bytes for both the record type and file number provides redundancy for error control, since the second byte is simply the one's complement of the first. The record types are denoted by 0201 for file marker, 0303 for numeric data, and 0347 for symbolic data. The following table summarizes all of the various data formats used by the system. XP and CP denote the two longitudinal parity checks and will be described later. FN denotes the file number and -FN its one's complement.

FILE MARKER RECORD: 0201 / 0176 / FN / -FN
NUMERIC DATA RECORD: 0303 / 074 / XP / CP /
DATA
SYMBOLIC DATA RECORD: 0347 / 030 / XP / CP /
DATA (with VRC)
FILE: FILE MARKER / DATA RECORD / DATA
RECORD / ...
SYSTEM TAPE: BOOTBLOCK / FILE 0 / FILE 1 /
... / FILE 15 / FILE 32 / FILE 127
DATA TAPE: FILE 0 / FILE 1 / ... / FILE 127

2.4 LOADER ACTION

When the loader is told to load a given file, it begins searching the tape (the loader can load files from either deck, depending upon which entry point is used) forward until it finds a file marker record. Note that all records passed over must have a valid type number pair or an error recovery procedure will be initiated which will try up to three times to read the record correctly and then make an error exit if failure occurred all three times. Upon finding a file marker, the loader determines, from the number in that marker, whether the tape is positioned to the correct place (the number is equal to that requested), is not positioned far enough forward (the number is greater than that requested). If the tape is positioned to the correct place, the loader proceeds to load all of the numeric records it finds, obtaining the memory address of where it is to put the data from the beginning of each record, (symbolic records are ignored) until it runs across another file marker. At this point it stops the tape (which was in slew forward mode) and backs up over the file marker so a succeeding call on the loader would cause a file marker to be found immediately. If there were no numeric records in the file, an error return is made. If the tape is not positioned far enough forward, the loader searches forward for the next file marker. If the tape is positioned too far forward, the loader enters a reverse search mode. If, in this mode, the loader finds a file marker that indicates that the tape is now positioned to the correct place, tape motion is reversed and the file is loaded as in the forward search case. If it finds a file marker which indicates that the tape is not positioned backward far enough, the loader continues searching in the reverse mode for the next file. If, however, a file marker is found that indicates that the tape has been positioned too far backward, the loader decides that the file is not on the tape and makes an error return. Error returns are also made if a record can not be read without a parity failure or type indicator discrepancy (the two characters are either not the one's complement of each other or are not one of the three special numbers) occurring in all three trials or if loading the record would overstore the loader. In all of these cases, the carry condition will be true (a satisfactory load always rendering the carry condition false) and the tape will be positioned after any offending record.

2.5 PARITY CHECKING

The third and fourth bytes of every data record contain longitudinal parity checks. These bytes are set up by the record generation program such that the following exclusive OR sums will yield zeros: the first byte with all the data characters (data characters start with the fifth byte of the record and proceed to the end) and the second byte with the same characters except the sum is shifted right circularly one place after each exclusive OR. In the case of symbolic records, the additional condition of the vertical parity of each character being odd must also be met. One thing not mentioned in the discussion of the loader was that the first four data characters (fifth through eighth bytes in the record) are not really data but are the MSP and LSP of the starting memory address followed by the one's complement of the MSP and LSP of the starting memory address of where the data is to be loaded.

3. THE CATALOG, SYMBOLIC LOADER, BASIC I/O, AND DEBUG

As mentioned above, the operating system maintains a catalog of names which correspond to the files on the system tape. This catalog may be used in manipulating the files from the keyboard or in symbolically calling in overlays using the symbolic loader from a user program.

3.1 CATALOG CHARACTERISTICS

Each name in the catalog must start with a letter and may additionally contain from one to five alpha-numeric characters. There is room in the catalog for up to fourteen names so there is a limit of fourteen catalogued files on one system tape. The symbolic loader contains routines which will look up a given name up in the catalog and load the corresponding file. This same lookup routine is used by the command handler and is labeled LOOKUP.

3.2 UTILITY ROUTINES IN THE SYMBOLIC LOADER

Other utility routines in the symbolic loader area are a block transfer, labeled BLKTFR, and a routine, labeled INCSWP, which increments the H and L register pair and then swaps it with the D and E register pair. The block transfer will move the number of characters specified by the entry value in the C register from a memory address starting with the entry values in the H and L registers to a memory address starting with the entry values in the D and E registers.

3.3 LOADING ROUTINES

To use the symbolic loading routine, one loads into the D and E registers the address of the six characters of the desired name (trailing blanks must be included) and calls

MLOAD\$. If the zero condition is false upon return, then the given name was not in the catalog. If the zero condition is true but the carry condition is false upon return the loader could not either find or correctly load the file requested. Note that one must be certain to place the call to MLOAD\$ in a place that will not be overlayed since execution will resume following the CALL instruction after the file has been loaded.

3.4 OTHER SYMBOLIC LOADER FACILITIES

Another facility in the symbolic loader area will load and execute a file whose number (not name) is in the B register upon call. Calling MAUTO\$ will load the file from the system tape and calling MAUT2\$ will load the file from the data tape. If the loader could not either find or load the file, the operating system is automatically reloaded.

3.5 KEYBOARD AND DISPLAY ROUTINES

The operating system contains facilities to ease the burden of communicating with the operator. Two routines exist. The first accepts the characters from the keyboard, displays them on the screen, and stores them into a memory buffer. The second writes a string of characters from a memory buffer onto the screen.

3.5.1 KEYBOARD INPUT

The keyboard input routine, labeled KEYIN\$, accepts a specified maximum number of characters, given by the entry value of the C register, from the keyboard and puts them into memory starting at the entry value of the H and L registers and onto the screen at a starting horizontal cursor position of the entry value of the D register and vertical cursor position (which cannot be changed during the course of one input) of the entry value of the E register. Note that if the cursor collides with the right edge of the screen during entry, characters other than backspace, cancel, and ENTER will not be accepted, although they will print over each other in the last display position. The ENTER character (015) terminates input and is stored in the memory buffer to specify the end of data but is not written to the screen. Hitting the backspace key will delete the last character entered and move the cursor appropriately while hitting the delete key will delete all characters entered and also move the cursor appropriately. These two keys also back up the buffer memory pointer appropriately. Note that if one has typed a character at either the screen limit or at the maximum character count limit, hitting a backspace will cause the previous character to be erased and leave the last character still on the screen, although it will either not appear on the memory buffer or be after the 015.

3.5.2 DISPLAY OUTPUT

The display routine, labeled DSPLY\$, will display the string of characters stored in memory starting at the address which is the entry value of the H and L registers and terminating with a character whose numerical value is either a 3 (ETX) or 015 (ENTER). The cursor starts at the entry values in the D (horizontal) and E (vertical) registers (a cursor position that is off the screen will not be sent to the CRT) and stops after the last character printed if the terminating character was a 3 or at the beginning of the following line if the terminating character was an 015. Note that, as in KEYIN\$, the cursor stops at the right edge of the screen and the characters overwrite each other if more are available after collision. Also note that if display was occurring on the bottom line and the terminating character is an 015, then the whole screen is rolled up to force the existence of a following line and the information that was at the top of the screen is lost. After return from the display routine, the H and L registers will point to the location after the terminating character and the D and E registers will reflect the current cursor position. The cursor will be off while the display routine is writing, but it is turned back on upon exit even if it was off upon entry. Other special control characters can cause cursor positioning, line/frame erasure, and screen roll-up:

011 - a new horizontal position (0 to 79) follows
013 - a new vertical position (0 to 11) follows
021 - erase to the end of the frame
022 - erase to the end of the line
023 - roll the screen up one line

3.6 THE DEBUGGING TOOL

The debugging program allows the user to observe and modify any location in memory, to load files from either the system or data tapes, and to start execution at any place in memory. This allows him to load and debug programs with surprising ease. The major debugging technique is to insert RETURN instructions in critical places in memory so one routine at a time may be checked using the CALL command. All but two (user specifiable) of the registers may be saved upon return from the program being tested, allowing the user to determine if the proper actions are taking place by observing critical register and memory values. The registers A, B, C, D, E (subject to the H and L commands in Section 3.6.3) are stored in locations 16770, 16771, 16772, 16773 and 16774 respectively upon a return to Debug from a program which was called from Debug.

3.6.1 INPUT SYNTAX AND ERROR ACTION

The debugging program is entered from the command handler as explained in a later section or by processor execution control being passed to the location labeled DEBUG\$. At

this time the bottom line of the display will be erased and the current location and its contents will be displayed there. The program is now ready to accept input in the format <number><command>. The number is assumed to be octal and the absence of any digits between zero and seven implies a value of zero for the number. Only sixteen bits of significance are kept for the input value. If more are entered, the first digits entered are lost. Some commands use only the lower order eight bits. The number is terminated by the first character that is not between zero and seven and this character is taken to be the command. Note that leading spaces are not permitted. This line is read in using the KEYIN\$ routine previously discussed, thus enabling the use of the backspace and cancel keys but requiring the ENTER key to be struck to obtain a response. In one case the ENTER character is the command and in some others the number is disregarded. If the command is not recognized, the program simply ignores it and the old current address and its contents are displayed again. After every command, control is returned to the entry point of the debugging program which will display the now current address and its contents.

3.6.2 THE CURRENT ADDRESS

Two memory locations in the debug contain an address (initialized to zero upon loading) which points to a memory location which is the current center of interest. Available commands allow one to change the contents of this memory location and move the pointer as well as perform other functions.

3.6.3 COMMAND MEANINGS

The following is a list of each command character with its effect and the number (in parenthesis) of bits of the given number used:

ENTER - set the current address to the number given (16)
I - increment the current address by one (0)
D - decrement the current address by one (0)
M - change the current address contents to the number given (8)
. - do the M followed by the I command (8)
L - upon return from a C command, cause the L register to be stored into the register whose number is given (3)
H - same as the L command but for the H register (3)
G - load from the system tape the file whose number is given (8)
F - load file one from the data tape (0)
O - return to the operating system command handler (be sure it is there) (0)
C - execute a CALL instruction to the location whose number is given (16)

4. KEYBOARD FACILITIES (OS COMMAND HANDLER)

The operating system contains a program which will interpret user commands given at the keyboard and perform the tasks indicated. These tasks mainly involve copying new files from the data tape onto the system tape, copying files from the system tape onto the data tape, deleting and updating files on the system tape, and executing programs kept in these files, as well as several other functions.

4.1 SYNTAX RULES AND ENTRY ERROR ACTIONS

The command input format is purposely made quite strict to reduce the chance of causing unwanted action which could be catastrophic to the user's data. The command must start with the first character entered (leading spaces are illegal) and any alphabetic after the third character is ignored (thus DEBACLE will be interpreted as the DEB command just as well as DEBUG). The first non-alphabetic character must be either an ENTER, a space, or a dash (minus sign). Some commands will not allow the ENTER but typing a non-alphabetic other than these three will always net you an error message of WHAT?. This will also appear if a command that has legal syntax but is not one of those defined is entered. If the command is to be parameterized, the first name must follow the dash or space immediately and must be terminated with an ENTER if that is the only parameter. The name must start with an alphabetic but may contain any number of alpha-numeric characters even though all after the six will be ignored. If the command has two parameters, the first must be terminated by an ENTER. If a name is terminated by characters other than those specified, the error message BAD NAME will be displayed. If a name is not supplied but the command requires one, the error message NAME REQUIRED will be displayed. If the name given is required to be in the file catalog but is not, the error message NO SUCH NAME will be displayed. If the inverse is true, the error message NAME IN USE will be displayed.

4.2 OPERATING SYSTEM COMMAND INSTRUCTIONS

The following paragraphs describe the usage and effect of each command in the system. Each paragraph is titled by what must be entered to use the corresponding routine. Note that, for clarity, more than just the necessary three characters have been shown.

4.2.1 CATALOG

The CATALOG command lists the names of all files that are currently on the system tape. They are listed across the screen in the physical order in which the files appear on the tape. Any parameters supplied are ignored.

4.2.2 NAME (old), (new)

The NAME command will change the name of the file specified by the first name given to the second name given. This command requires that the first and not the second name be in the catalog. The catalog file on the system tape (a one record file (number one) that immediately follows the operating system file) will be overwritten with the new catalog. Note that this leaves the system tape positioned before the file marker of any existing first cataloged file. This operation is performed by all commands that change the catalog.

4.2.3 RUN (name)

The RUN command uses the loader to load the file specified by the name given and then transfers processor control to the starting address indicated to the loader by the file information. Note that it is the responsibility of the loaded program to return to or reload the operating system if this is desired. There is a special case to the RUN command that breaks the general syntax rules. If the name consists of exactly one asterisk terminated by an ENTER (RUN-*), the loader will be directed to load physical file 1 from the data tape. This provision is made to allow the user to run a program he has generated without having to load it onto the system tape. This, along with the F command in the debugging tool, eliminates a lot of tape movement when debugging programs.

4.2.4 IN (name)

Note that exactly the characters shown must be typed to execute this command since the space which must be the third command character will also terminate the command. This command will position the system tape after the last cataloged file and the data tape to the beginning of physical file 1. (The data tape convention is that physical file 0 will be the first piece of information on the tape, containing the users symbolic data for a given program, and that physical file 1 will be the second piece of information on the tape and will contain the users object data for a given program, and all tape after this is to be considered a scratch area which is properly terminated by physical file 127 to indicate the logical end of the tape.) The command then copies all records in the file from the data tape onto the system tape creating a file on the system tape (a file marker being written before the data was copied) which has the next available physical file number. Following this new file, file markers 32 and 127 are written on the system tape to indicate the new start of system scratch and logical end of tape. If the system tape contained no cataloged files before this command was issued, the file entered will be physical file 2 and immediately follow the catalog file. After the new file has been written, the new name is entered into the catalog and the catalog file is updated as in the NAME command. Note that if the catalog was full when the command was entered, the error message LIBRARY FULL will be displayed and no

other action will occur. The name supplied must not already be in the catalog.

4.2.5 OUT (name)

The OUT command first executes the PREPARE command to provide itself with a null data tape which can be handled by the file handling routines. It then positions the system tape to the beginning of the given file (the name must have been in the catalog) and the data tape to the beginning of physical file one and copies all the records in the file on the system tape onto the data tape. It then places a file marker 127 on the data tape and quits. Note that the catalog file is not updated for this command. This command is provided to allow moving a file from one system tape to another through the associated use of the IN command.

There are two special cases to the OUT command that break the general syntax rules. If the name consists of exactly one dollar sign terminated by an ENTER (OUT-\$) then an exact copy is made of the system tape up to file marker 32 at which time the copy is terminated by file markers 32 and 127 (which causes any scratch data on the old system tape to be removed). If the name is exactly one asterisk terminated by an ENTER (OUT-*), the action is similar to the previous case except physical files 0 and 1 (namely, the operating system) are deleted and the file numbers of all following data files (not file 32 or 127) are lowered by two. Note that if this tape is now bootstrap loaded, the first program loaded will be what was the first file catalogued in the operating system. This is most useful in preparing bootstrap tapes that will be used in machines with less than 8K of memory.

4.2.6 DELETE (name)

The DELETE command takes two different courses of action depending on whether or not the file deleted is the last one catalogued. If it is, the system tape is moved to the end of the next to the last catalogued file and file markers 32 and 127 are written, thus logically destroying the last file. The name is then deleted from the catalog and the catalog file is updated. If the file is not the last one catalogued, the PREPARE command is called to obtain a fresh data tape, as in the OUT command, and the system tape is positioned to the end of the named file. The rest of the system tape (up to the file 32 marker) is then copied onto the data tape and the data tape is terminated with a file marker 127. Note that the data tape file numbers start out at one and increase by one for each succeeding file copied onto the data tape. These numbers are not used since all the copy back part needs to know is file delimitation since it is getting its file number information from catalog positions. The copying onto the data tape is followed by the system tape being positioned to the end of the file before the one named and the data tape being positioned to the beginning of file one. A file marker having a value one greater than the previous

marker is then written on the system tape and then the data tape is copied back onto the system tape with every file marker encountered on the data tape causing a file marker of value one greater than the previous marker to be written on the system tape. This process terminates when a file marker 127 is encountered on the data tape which causes file markers 32 and 127 to be written on the system tape. The given name is deleted from the catalog, all following entries are dropped down one place to correspond to the similar shift in file numbers that took place, and the catalog file is updated.

4.2.7 REPLACE (name)

The REPLACE command is quite similar to the DELETE command except that instead of preparing the data tape with the PREPARE command, it positions it to the end of file 1 and then writes a file marker 2. Now, copying all the files after the named one onto the data tape in a fashion similar to the DELETE command and copying the data tape back onto the system tape in exactly the same fashion as in the DELETE command will replace the named file by file 1 on the data tape, with any necessary physical expansion or contraction taking place. Even though the catalog is not changed in this operation, it is updated anyway since this is an easy way to position the system tape to a place before file marker 127. Without this, a succeeding call on the loader would run into trouble since the system tape would be left positioned after file marker 127 and the loader always starts by searching a tape forward which in this case would be off the logical end of the tape. The loader starts with a forward search because the very first time it is used, the tape is positioned just after the boot-block and a backward search for a file marker would cause trouble. The operating system routine which searches for files can start with a reverse search to avoid the problem since the tape will never be resting before file zero.

4.2.8 AUTO or AUTO (name)

There is a word in the catalog which contains the physical file number of a file which should be loaded and executed immediately upon loading and execution of the operating system. This enables a user program to be run after restart without interaction with the operation system being required. If this word is a zero or the keyboard switch is being depressed upon initial execution of the operating system, the normal entry is made into the operating system and the start up message and response request are displayed.

If the AUTO command is given with no name and the auto pointer is zero then the error message NAME REQUIRED will be displayed. Otherwise the name of the file being pointed to will be displayed in the message AUTO SET TO (name). If the auto command is given with a name (which must be in the catalog) and the auto pointer is a zero, the

pointer will be changed to the corresponding file number and the catalog (which contains the pointer) will be updated. If the auto pointer is non-zero, the name is ignored and the AUTO SET TO (name) will be displayed as in the no-name case.

4.2.9 MANUAL

The MANUAL command will zero the auto pointer and update the catalog if the auto pointer was non-zero. Otherwise, the message AUTO NOT SET will be displayed.

4.2.10 PREPARE

The PREPARE command first asks the operator if the data tape contains anything of value and then halts. (Note that the auto-restart tab should not be broken out of the operating system tape because it will prevent use of the OUT, DELETE, or PREPARE commands since halting the processor will cause an auto-restart.) After the operator hits the RUN button as a response, it is assumed that the data tape is of no value as it is rewound and file markers 0, 1, and 127 are written on it. This is needed since the operating system routines require file markers for which they can search in using the data tape.

4.2.11 HEX (name)

The HEX command is similar to the IN command except that the data tape is formatted in symbolic records with no parity checking. This is useful in loading onto the system tape data produced by sources other than the 2200. There are four types of records accepted. The type is determined by the second character (the first must always be an 012 (LF)): asterisk means ignore the record; pound sign denotes the logical end of the tape; plus sign means the following four hexadecimal characters are a new starting address (these must be terminated by an 023 (XOFF)); and a hexadecimal character denotes a data record. All other cases are assumed to be data read errors. A data record must always contain an even number of only hexadecimal (0 through 9 and A through F) characters terminated by either an 023 or a plus sign. The characters are paired up to form successive bytes of eight bit data. If the terminating character is an 023 then the block of data bytes is written out in loader format and the starting address is incremented by the number of data bytes. If the terminating character is a plus sign then the data remains in the buffer and the following record will be appended to it. This allows blocks of larger than 36 bytes (128 is the upper limit) to be written when the device which writes the tape is limited to lines of 72 characters. Note that there is no buffer overflow protection and it is the responsibility of the program generating the symbolic data to keep the total number of continued bytes to 128 or less (128 hexadecimal character pairs). Also note that if a continuation line is followed by a new address line, the data will remain in the buffer but the starting address will change. This combination will cause incorrect results

since even if the buffer did not overflow will also overwrite critical pointers which will cause the operating system to produce an error message (because it will be called with incorrect parameters when the critical pointers are overwritten) and be reloaded. If a read error is detected, the data tape is backspaced one record and read again. This will go on until the data appears correctly or the keyboard switch is depressed. Depression of the keyboard switch causes the same action as reading from the data tape a record starting with a pound sign.

4.2.12 DEBUG

The DEBUG command causes the debugging tool described earlier to be entered.

4.3 SYMBOLIC OPERATING SYSTEM AND EXTENDED COMMAND INSTRUCTIONS

The overlay program SOSX is available to extend the operating system command set. The following paragraphs describe the usage and effect of each new command. Each paragraph is titled by what must be entered to use the corresponding routine. Note that, for clarity, more than just the necessary three characters have been shown.

4.3.1 CHOP (name)

The CHOP command deletes the named file and all subsequent files.

4.3.2 INSERT (new, (old)

The INSERT command proceeds like a REPLACE command except it includes the old named file as one of the files written after file 1 on the front deck. When the front deck is copied back onto the CTOS tape a new object file has been inserted.

4.3.3 APPEND (name)

The APPEND command appends the object file from deck 2 onto the end of the named file on the CTOS tape. Like the DELETE command, it has two possible courses of action, depending on whether or not the file being appended is the last cataloged file. If it is, the tape is positioned to the end of the cataloged file and a new object file is copied from the front deck. New file 32 and 127 markers are written. If the named file is not the last cataloged file, the operation proceeds like REPLACE except that the CTOS tape is positioned to the end of the named file before the copy is performed.

4.3.4 LGO (name [, name, name . . .])

The LGO command makes a tape with a loader and the named file(s) in the sequence named in the command. The files will have sequential file markers starting with 0. There is a limitation

of 23 characters on the command length, thus to name many files in the LGO command it may be necessary to temporarily rename the files with one character labels. LGO * is not permitted. OUT * has the desired effect of generating a load and go tape of all catalogued files.

4.3.5 SYMBOLIC (name)

The SYMBOLIC command adds a compressed source file (file #0) to the CTOS tape (in a fashion similar to IN). The name in the internal catalog will have an 'S' in the seventh (not displayed) position to identify the file as symbolic.

4.3.6 SREPLACE (name)

The SREPLACE, symbolic replace, command is performed exactly as the REPLACE except the compressed source file (file #0) is used instead of the object file. File 1 may be overwritten.

4.3.7 SINSERT (new), (old)

The SINSERT, symbolic insert, command is performed exactly as the INSERT except the compressed source file is inserted instead of the object file. File 1 will be overwritten.

4.3.8 ATTACH (name [, name, name . . .])

The ATTACH command positions the front deck to the end of file 0 and (without file markers) copies specified file(s) from the CTOS tape to the front deck. When all specified files are copied, the question 'END (LABEL OR :)?' will appear. A six character label may be entered. If ':' is typed no end statement will be added. The ATTACH * form of the command will attach, in catalogued sequence, all symbolic files to file 0 on the front deck.

5. PROGRAM FACILITIES (OS FILE HANDLER)

The operating system contains a set of routines which will perform all of the various input/output functions needed to maintain the files of data on the tapes. These routines are packed in the upper 2K of memory and are made available to the user if he wishes to handle his mass storage problems in conformance with the conventions of the operating system. All routines are uniformly parameterized and are accessed through an entry point table (a group of JUMP instructions to the actual routine locations) so any updates to the operating system will not have any effect upon the user's code.

5.1 ROUTINE PARAMETERIZATION

Routine parameterization consists of a memory location in the D (MSB) and E (LSB) registers of the first byte of a group of four bytes (called a packet) which parameterizes the call

more explicitly. This method reduces the number of memory locations required to perform a routine call since, in a typical program, one needs only a few different packets but will have many different calls. The parameterization of some routines is not as extensive as that of others, but the same packet can generally be used for the different calls when they are affecting the same file.

5.1.1 LOGICAL FILE NUMBERS

The first byte in the packet is the logical file number and must be between zero and seven or an internal error H will occur upon calling any routine using this packet. This error condition usually occurs when the user has either failed to load the D and E registers at all or has loaded them with an erroneous value before calling the routine. The second and third bytes in the packet contain the LSB and MSB (respectively) of the first location in memory to be used as a data buffer. Actually, the two bytes previous to this location will be used by some of the routines as discussed later. This data buffer may be located anywhere in memory. The fourth byte in the packet specifies the length of the data buffer when numeric data is being handled. Note that using only one byte for the length implies that numeric records may not contain more than 256 data bytes. Actually, the maximum number of data bytes specified may not be greater than 254 for reasons that are made clear in the numeric routine instructions. The four bytes of the packet may be located anywhere in memory.

5.1.2 PHYSICAL DEVICE AND FILE NUMBERS

The logical file number specified in the packet is converted by each routine, via an internal transformation table, into physical file and device numbers. The physical device number specifies whether the operation is to be performed on deck one (rear) or deck two (front) and the physical file number specifies which file is to be treated on the given deck. Actually, not all routines use all of this information since, for instance, when one is reading records from a file he assumes that he is using the file to which the tape was last positioned. The internal transformation table is initialized at load time to the following values:

LOGICAL FILE	PHYSICAL FILE	PHYSICAL DEVICE	GENERAL USE
0	0	0	Unassigned
1	0	1	General deck one
2	0	2	General deck two
3	1	1	CTOS catalog
4	0	2	Symbolic data
5	1	2	Object data
6	0	0	Unassigned
7	32	1	System scratch

It is shown that logical files 1 and 2 are specified for use of any physical file, even though 0 is shown in the table. This can be done by use of a routine that will change the physical file number of a given logical file. A routine also exists to allow the physical device number to be changed, thus allowing the user to set up logical files in any physical configuration needed. Note, however, that one must have logical files 1 through 5 and 7 in the state shown (except for the physical device numbers of logical files 1 and 2) if one returns control to the operating system command handler, since the loaded values are assumed by this program. Logical files 0 and 6 may be used freely but must be set before the first call utilizing them. The following is an example of a packet usage as it would be expressed in the assembler: (Note all calls to CTOS tape routines must, as in the following example, be preceded by a DE to the first byte of the packet. Note also that the packet consists of 4 bytes: Logical file number, LSP of buffer, MSP of buffer and length of buffer)

	LA	2	Set up logical file six to be used as physical file 3 on the front deck
	DE	PACKET	
	CALL	CPDN\$	
	LA	3	
	DE	PACKET	
	CALL	CPFN\$	
	DE	PACKET	Position to the beginning of the file
	CALL	PBOF\$	
LOOP	DE	PACKET	Read a record of symbolic from it into BUFFER
	CALL	SSFR\$	
	JTC	DONE	Quit if to the next file marker
	JFZ	TERR	Exit if type error
	.	.	Action taken for each record
	JMP	LOOP	
DONE	.	.	Action taken when file completely in
	.	.	
TERR	.	.	Type error action
	.	.	
PACKET	DC	6	Logical file 6
	DA	BUFFER	Buffer address
	DC	0	Length (not used)
	DC	0,0	Room for parity check generation
BUFFER	SKIP	128	Buffer area

5.2 ROUTINE USAGE INSTRUCTIONS

To use a routine, one sets up whatever is required for proper parameterization and then calls the desired location in the entry point table. The locations are labeled with the first word in the following paragraph titles followed by a \$. For example, to call the serial numeric file read, one would say CALL SNFR\$. The routine will either perform the requested task or take one of two error exit paths. The first path is

taken in the case of fatal errors, for which it is decided that the only recourse is to reload the operating system. This is called an internal error and the message INTERNAL ERROR (letter) is written on the bottom line of the display before the system is reloaded. The various letters which may appear are the following:

- A - Illegal device specification ..
- B - Illegal record format
- D - Unrecoverable parity error
- G - Unfindable file
- H - Illegal logical file specification

The other path is non-fatal and simply returns with certain condition flags in states other than normal to indicate that something unusual happened. Since every routine uses a common subroutine (labeled GETPKT) to get the parameters from the specified packet, common internal errors can occur. If the logical file number is not between zero and seven, an

internal error H occurs. If the physical device number is not either a 1 or a 2, an internal error A occurs. Other than for these error actions, the following paragraphs described the effects of and the exact parameterization needed for each routine.

5.2.1 SNFR - SERIAL NUMERIC FILE READ

This routine reads the next record from the specified device. If the record is of type symbolic, the zero and carry conditions

are set false and return occurs with no parity checking or data storage being performed. If the record is a file marker, the carry condition is set true and the tape is backed up to where it was before the routine was called. Again, return occurs with no data storage being performed. If the type is numeric, the two parity bytes followed by the data are read into the buffer. If the parity checking fails or the record type is bad, three efforts are made at reading the record by backing up to its beginning and starting over. If recovery is not made in one of these efforts, an internal error D occurs. If the record is read successfully, return occurs with the zero condition true, the carry condition false, and the H and L registers containing the memory location of the byte following the last one loaded from the tape. To calculate the length of the buffer area used, one must subtract the buffer starting address from returned values in the H and L registers. Remember that the first two characters in the buffer are not data characters but are the two longitudinal check sums. To obtain the number of data characters loaded, one must subtract the buffer starting address plus two from the returned values in the H and L registers. The parity checks are stored because the SBFW routine uses them instead of regenerating them from the data, thus shortening the time required to copy numeric records from one deck to the other.

5.2.2 SSFR - SERIAL SYMBOLIC FILE READ

This routine reads the next record from the specified device. If the record is of type numeric or file marker, the action taken will be the same as when SNFR reads a symbolic or file marker record. Action similar to that taken by SNFR is also taken if parity or type faults occur. If the record is read satisfactorily, only data characters will be in the buffer starting at the address specified. An 015 will mark the end of the data string and all vertical parity bits will be zero. The same normal exit conditions as in SNFR will occur.

5.2.3 SBFW - SERIAL BLOCK FILE WRITE

This routine writes a record of type numeric on the specified device. The total number of bytes, including the parity initialization sums as the first two, must be in the fourth byte of the packet. Note that inclusion of the parity initialization sums implies that the total number of actual data bytes cannot exceed 254. This routine assumes that the first two bytes in the buffer are the correct parity initialization sums since it does not generate them from the data. There are no error exits from this routine which implies that writing off the end of the tape will not be caught and that read-after-write checking is not performed.

5.2.4 SNFW - SERIAL NUMERIC FILE WRITE

This routine performs in a fashion similar to the SBFW routine except the two parity bytes are not included

in the data buffer and the length specifies the number of actual data characters. The routine generates the two longitudinal parity sum initialization values and inserts them in the two locations preceding the buffer. It then writes on the specified device a record of type numeric containing the two parity bytes generated, followed by the number of data bytes specified. Note that the length is adjusted to accommodate the two parity bytes so, as in the SBFW routine, only 254 actual data bytes may be written. If one specifies a length of 255 or 0, the only bytes (besides the record type) written on the tape will be respectively the first or both parity initialization sums. No error exits are made from this routine.

5.2.5 SSFW - SERIAL SYMBOLIC FILE WRITE

This routine performs in a fashion similar to the SNFW routine except that an 015 character in the data string rather than a specified value is used to determine the buffer length, vertical (in addition to longitudinal) parity generation is performed, and a record of type symbolic rather than numeric is written. The terminating 015 character is not included in the set of characters written to the tape, but remember that it will appear again if the SSFR routine is used to read the record.

5.2.6 PEOF - POSITION TO END OF FILE

This routine searches forward on the specified device until it finds a file marker. It then backspaces the tape until it is between the next to the last and the last record in the file. It then forward spaces the tape one record which puts it at the end of the file, having arrived there via forward tape motion. This forward arrival is important to observe when one plans to append one record after another and still maintain physical interrecord gap integrity. Note that every record passed over by the PEOF routine must have a valid record type or it will be read again in action similar to parity failure action in the SNFR routine.

5.2.7 PBOF - POSITION TO BEGINNING OF FILE

This routine searches for a file marker in a fashion similar to the loader except it starts by searching backwards. The file number searched for is specified by the physical file number supplied by the generalized parameterization. Note that since this routine starts by searching backwards, it will not decide that the requested file is not on the tape until it has found in the search forward mode a file marker that specifies a file number greater than the one desired, if indeed the file is not on the tape. Also note that if the leader is found in the search backward mode, the tape is positioned forward past the first record and the backward search is continued. If the first record is not a file marker (operating system convention requires it to be) or is a file marker whose value is greater than the one desired, the first record on the tape will be passed over back and forth until external intervention is imposed.

Otherwise, all search rules and error exit conditions of the loader routine apply here. If, upon return, the carry condition is true, then the file was not found. Otherwise, the tape will be positioned at the interrecord gap following the file marker, having approached that point with forward tape motion for the reasons expressed in the PEOF routine instructions.

5.2.8 BSP - BACKSPACE

This routine simply backspaces the tape one record using the hardware backspace function. No checking is made to see if the record was of proper type or if the tape ran onto the leader.

5.2.9 CPDN - CHANGE PHYSICAL DEVICE NUMBER

This routine stores the entry value of the A register (note the break from generalized parameterization) into the physical device number entry for the specified logical file in the internal transformation table.

5.2.10 CPFN - CHANGE PHYSICAL FILE NUMBER

This routine stores the entry value of the A register (note the break from generalized parameterization) into the physical file number entry for the specified logical file in the internal transformation table.

5.2.11 TRW - TAPE REWIND

This routine performs a hardware high speed rewind of only the front deck. If the rear deck (physical device 1) is specified, an internal error A will occur. Upon exit from this routine, the tape will be positioned to the clear leader.

5.2.12 TFNR - TAPE FILE NUMBER READ

This routine acts in a fashion similar to PEOF until it finds the file marker. At this point, it simply reads the value of that marker and leaves the tape positioned after the marker record. The value read is returned in the C register. Error exits similar to the PEOF routine can occur.

5.2.13 TFNW - TAPE FILE NUMBER WRITE

This routine will write on the specified deck the special four byte file marker record containing the physical file number specified. No error exits will occur.

SECTION 6

DATAPOINT 2200

TRACE

SECTION 6

DATAPOINT 2200 TRACE PROGRAM

Introduction

TRACE is an interactive octal debugging aid for the Data-point 2200. It operates under the Computer Terminal Operating System (CTOS) and occupies memory space between 5600₈ and 13777₈. The normally resident operating system subroutines are not overlayed and are callable by the program being traced.

TRACE accepts commands from the keyboard and displays its results in the rightmost eight columns of the CRT display. It allows a user to trace the execution of a program, to examine and change the contents of the registers and memory.

Entering Commands

TRACE commands consist of up to two octal operands followed by a single letter operation. If there are two operands, a comma shall separate the two. An operand may be a 13-bit address or an 8-bit byte value, either expressed in octal. If the operand is an address, it may be given in two parts, separated by a blank. The first part consisting of the five most significant bits, and the second part consisting of the remaining eight bits. An address may also be given as a single octal number, but it is displayed in two parts as described. Leading zeros need not to be entered.

Examples:

TYPED VALUE	DISPLAYED VALUE	
	FIRST PART	SECOND PART
101A	—	101A
707J	01	307 J
0331, 1477W	{ — 03	331, 077 W

The command being entered is displayed on lines 10 and 11 of the CRT display, as shown in Figure 1. Line 10 shows the first operand. Line 11 shows the second operand and operation. If there is only one operand, line 11 will be blank. If there are no operands, line 11 shows only the operation and line 10 will be blank. NOTE: In all ensuing examples, the display format is used to exemplify the referenced operation.

If an illegal character is typed, the beep signal will sound and the character will be ignore. The CANCEL key will cause the command being entered to be discarded and another command can be entered.

The **ENTER** key will cause the command just entered to be executed. The command may be CANCELED at any time before the **ENTER** key is depressed.

Command To Modify Registers Or Memory

A, B, C, D, E, H, L

The operations A, B, C, D, E, H or L take a single byte value operand. The register specified by the operation is set to the operand value.

Examples: 1 7 3 B sets B to 173
0 0 1 A sets A to 001
3 7 5 L sets L to 375

Operation F

The operation F takes a single address operand. The Zero, Sign, and Parity flip-flops are set as if the lower 8 bits of the address were the result in A of some arithmetic instruction. The Carry flip-flop is set to the rightmost bit of the first part of the address.

Examples: 0 0 0 F sets Zero
 resets Sign, Carry, and Parity

0 1 0 0 0 F sets Zero and Carry
 resets Sign, and Parity

200 F sets Parity and Sign
resets Zero and Carry

Operation O

The operation O takes a single address operand. It opens the specified location for possible modification. The contents of the location are shown on line 12 of the CRT display. A byte value can be entered followed by ; , <,or> followed then by ENTER. The location is set to that value. If the terminating character is ; , TRACE will accept another command. If it is <, the previous location is now opened. If it is >, the next location is opened. If the CANCEL key is used, the currently open location remains open and any modification for it is discarded. The modifying byte value is shown as it is entered, following the contents of the open location on the CRT display.

Examples:

01 115 O opens 01 115
 < now opens 01 114
 > reopens 01 115
 57; sets 01 115 to 057

00 017 O opens 00 017
5> sets 0 017 to 005,
opens 0 020
20; sets 0 020 to 20

Command To Displayed Memory

Operation M

The operation M takes two operands, both addresses. They are the lower and upper bounds respectively of the region of memory to be displayed. Sixteen bytes are displayed across the entire width of the CRT display. The address is given on one line followed by the memory contents on a second line. The display continues being built and rolled up unless the DISPLAY key is depressed. The display then stops until the DISPLAY key is depressed again. The KEYBOARD key terminates the memory display.

Example: 000, displays the first 256 bytes of
00 377 M memory

Transfer Of Control Commands

Operation K

The operation K takes one address operand. It causes a Call instruction to be performed to the address given as the operand. Return is to TRACE.

Example: 02 000 K calls routine at 2 000

Operation J

The operation J optionally takes one address operand. If the operand is absent, the content of the P register is used. It causes a Jump instruction to be performed to the address given as the operand or in P.

Examples: 03 101 J jumps to 3 101
(octal 1501)
J jumps to address in P

Mode Setting Commands

Operation X

The operation X optionally takes two address operands. They are the lower and upper bounds respectively of a region in memory. Any Call instructions into this memory region are actually executed rather than being simulated. Since TRACE loses control at this point, it is imperative that routines called in this region return. All registers are properly set when the Call is performed. The contents of H and L are lost on the Return. If the operation X is given without operands, any region in effect is removed.

Example: 0 1 000 , sets the special CALL region to
to 1 000 to 3 377
03 377 X (octal 400 to 1777)

Operation W

The operation W optionally takes two address operands. They are the lower and upper bounds respectively of a region in memory. At the completion of any instruction in this region, the registers contents are shown in the first nine (9) lines of the CRT display as shown in Figure 1. The contents of the Carry, Zero, Sign, and Parity flip-flops are shown as the letters C, Z, S, and P respectively on the right-hand side of the CRT display if set and blank is reset. If the operation W is given without operands, any region in effect is removed. If the regions for the X operation and W operation overlap, the X operation takes precedence.

Example: 12 000 , display the register contents after
12 377 W every instruction in 12 000 to
12 377 (octal 5000 to 5377)

Operation S

The operation S optionally takes one address operand. Before TRACE executes the instruction at the given address, the register contents are shown as in a W operation region and TRACE stops to accept commands. A J operation with no operands will restart the program. If the S command is given without an operand, any stop address in effect is removed. If the stop address falls within an X operation region, the X operation takes precedence.

Example: 7 011 S the program will stop before the instruction at 7 011 is executed

To stop a program when TRACE is in control, depress the KEYBOARD and DISPLAY Keys at the same time. This has the same effect as an S operation for the current program address.

If a Halt instruction is executed by TRACE, the result is the same, as if an S operation was set for the Halt otherwise ignored.

Starting Trace

TRACE is loaded like any other program from the operating system library (RUN TRACE). Once started, it will request the name of the program desired to TRACE. The name of the program must be typed, followed by ENTER. CANCEL will cause TRACE to ask again. The named program will be loaded using the symbolic linker and loader in the operating system, and TRACE will show the register contents as if the program had been stopped at its entry point. P and I will be the only registers with non-zero contents.

To TRACE a program already in memory, simply depress the ENTER key without entering a program name, then jump to the entry point of the traced program, using the J command.

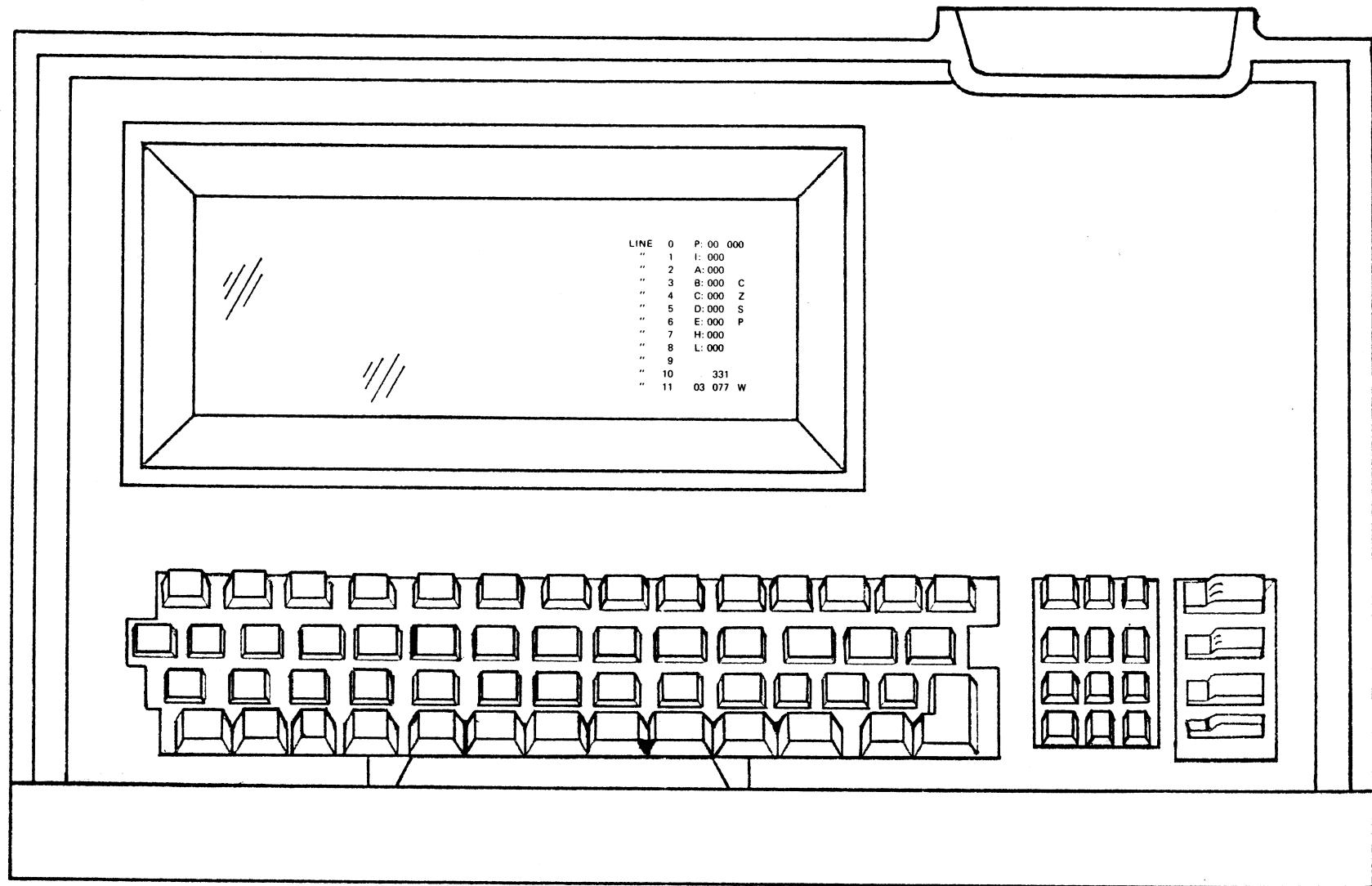
To TRACE a program located on the front deck type * and depress the ENTER key. The object file (file #1) from the front tape will be loaded using the operating system loader, and TRACE will show the register contents as if the program had been stopped at its entry point. P and I will be the only registers with non-zero contents.

Operational Summary

A	Set A to operand
B	Set B to operand
C	Set C to operand
D	Set D to operand
E	Set E to operand
H	Set H to operand
L	Set L to operand
F	Set flip-flops from operands
O	Open location
J	Jump
K	Call
M	Display memory
X	Set special CALL region
W	Set register display region
S	Set stop address

Figure 1. Operation W

6-4



SECTION 7

DATAPOINT 2200

ARITHMETIC SUBROUTINES

SECTION 7

7.1.1 INTRODUCTION

STATH is a subroutine package specifically designed to provide formatted keyboard input, screen display, checksum and arithmetic operations on numeric strings. Each function of STATH is obtained by calling the entry point associated with that function.

Following is a list of the functions available through STATH. The labels given to their entry points and the sections incorporating their usage parameters:

Entry Point	Function
ADD\$	Addition
COM\$	Compare Magnitude
DIV\$	Division
DSP\$	Display on screen
KEY\$	Keyboard formatted Input
MOD10\$	MOD 10 checksum calculation
MOD11\$	MOD 11 checksum calculation
MOV\$	Move string
MUL\$	Multiply
SUB\$	Subtract

7.1.1.1 INTRODUCTION TO STRINGS - NUMERIC AND OTHERWISE

The purpose of a 'string' is to carry around a 'package' of text. A string is an individual block of text and just like a string it has a definite beginning and end. The composition of the string is an uninterrupted sequence of ASCII characters. That is, between the beginning and end of the string only ASCII characters are allowed. The ASCII character may be any of the 95 plus space (blank) characters listed in Section 1 of the Programmers' Manual (2200 Reference Manual).

The string is bounded at the beginning and end in different ways. The end is determined by the first occurrence of the ASCII 'ETX' which is equal to (003₈) in the sequence of characters called the string. The 003 tells STATH that the string is ended. The CTOS will also accept a carriage return character (015₈) in place of the 003 but STATH only accepts the 003.

The following are valid strings. The contents of the parentheses are intended to be the byte value of the ASCII character for single character values or the octal value of the octal triple such as 003.

(N) (O) (W) () (I) (S) () (T) (H) (E) () (T) (I) (M) (E) (003)

(0) (1) (2) (3) (4) (5) (6) (7) (8) (9) (0) (003)

Which are in octal:

116,117,127,040,111,123,040,124,110,105,040,124,111,
115,105,003

and

060,061,062,063,064,065,066,067,070,071,003

Although a string has an inherent end built into itself, the 003, there is no beginning. At least no beginning which itself is part of the string of characters in memory. The beginning is combined with the pointer to the string itself. That is, a string is referred to by calling out a location in memory. That location is the first character of the string. In the above samples, for 'now is the time' to be referred to beginning with the word 'now' the location of the letter 'N' would be specified. It is clear that specifying only the 'N' yields a complete description which is:

'Begin with the character in the location specified and continue until a 003 is reached.'

Beginning with 'N' and continuing to the 003 gives: 'Now is the time'. If the location of the letter 'W' in now were specified, the string resulting would be 'w is the time'.

Therefore, to specify a string to a routine (like STATH) which is going to use the string, the user must only transfer the address of the first character of the string or the character in the string the user wants to begin the string (it may not be the first) to the routine. Also, if the user created the string, he must be assured that there is a terminating 003 byte immediately following the last character of the string in memory.

STATH differentiates between two categories of strings:

1) Numeric strings

and

2) Non-numeric strings

Where numeric strings are only regular strings with the character set restricted to the characters 012345689 with an optional single period representing the decimal point and/or a single hyphen leading the string representing a minus sign.

A non-numeric string is any string which is not numeric by the above definition.

A numeric string (omitting temporarily the 003) can look like:

00000034567788888777.99999999999999

or

-123.45

or

34.5000000000

There is a size limit as to the number of characters a string may have in STATH. This is not true of ordinary text strings in CTOS where a string may, for some strange purpose, have thousands of characters in it. STATH is a mathematic package and the numeric strings represent numbers. The largest number of digits, therefore, is limited in STATH and that limit is 126.

7.1.1.2 INTRODUCTION TO THE FUNCTIONS OF STATH

STATH functions fall in the following four categories. The categories are listed with their appropriate functions below.

Arithmetic Analysis Manipulative Input/Output

Addition	Compare	Move	Keyboard formatted input
Subtraction	MOD10 Check		Display on screen
Division	MOD11 Check		
Multiplication			

The arithmetic functions are the normal functions with which everyone is familiar.

The analysis functions permit decisions to be made on the content of a number. MOD10 and MOD11 verify the checksum Modulo 10 or 11 as is used in many business applications. Compare will permit comparing two numbers to determine equality or relative magnitude.

The move function is necessary as a preparation for using the multiplication and division functions in STATH, applicable for general use in the user's program to move numeric strings from one location to another and to format and round them in the process.

The input/output functions provide the user with simple techniques for bringing numbers into memory from the keyboard and displaying string numbers in memory onto the screen.

7.1.2 STATH FUNCTIONS AND ARGUMENTS

Each routine takes one or two arguments. An argument consists of a CTOS-compatible string. The argument strings are bounded at the end by an ASCII ETX (=003), and the beginning boundary is determined by the address contained in the register-pair associated with that argument. The maximum size for any STATH string is 126 characters. This means arguments and results are limited each to 126 digits.

Except for the routine DSP\$, all strings must be 'numbers' which means a sequence only of ASCII numeric digits (0123456789) with an optional decimal point. Optional leading minus sign and optional leading blanks (an octal 040). The number must be right justified in the argument string. All strings except for DSP\$ set the condition flags as follows:

Flag	Indication
Zero	The result was zero
Sign	The result was negative
Carry	An overflow occurred
Parity	One or both arguments were improperly formatted

If parity is not set at the end of an operation, HL and DE contain the addresses of the location in memory past their respective ETX's. In the case of KEY\$ and DSP\$, D contains the column and E contains the row of the position immediately beyond the display area used. MUL\$ and DIV\$ leave D and E with junk in them. MOD10 and MOD11 leave H and L containing the address of the check digit position.

7.1.2.1 EXAMPLES ON THE USE OF STATH

Following is a 488-byte program which is a useful desk calculator using STATH. It is included as an example of a program calling STATH functions.

'DCLAC', the desk calculator, inputs a numeric string and provides addition, subtraction, multiplication or division of that inputted string against an accumulator. 'DCALC' always inputs the string from the keyboard into a string labeled 'input'. The accumulator is in a string labeled 'accum'.

The four arithmetic operations performed in the program are routines labeled as 'ADDOP', 'SUBOP', 'MULOP', and 'DIVOP'. The routines are very short but demonstrate the use of STATH.

	SET	01000		LBR	
BOOTS	EQU	064		LA	DIVID+20
MOV\$	EQU	010000		SUB	
ADD\$	EQU	010003		LLA	
SUB\$	EQU	010006		LMC	
MUL\$	EQU	06000		LD	28
DIV\$	EQU	06003		LE	2
KEY\$	EQU	010014		HL	CLEAR
DSP\$	EQU	010017		CALL	DSPLY\$
KEYIN\$	EQU	017000		DE	ACCUM
DSPLY\$	EQU	017151		HL	ACCUM
MLOAD\$	EQU	017620		CALL	SUB\$
BEEP	EQU	13		DCALCL	LD 38
HEADING	DC	021,811,20,2 2 0 0		LE	5
	DC	013,5,011,31,'Total'		HL	ACCUM
	DC	013,7,011,28,'Keyboard'		CALL	DSP\$
	DC	013,2,011,28,'0 To 9'		LD	50
	DC	'Decimal Places?'		LE	7
DECPL	DC	'0',3		HL	BLANK+6
OVFMSG	DC	'Overflow',3		CALL	DSP\$
BLANK	DC	' ',3		LE	38
CLEAR	DC	022,3		LE	7
INPUT	DC	'0000000000',3		HL	INPUT
ACCUM	DC	'000000000000',3		CALL	KEY\$
DIVID	DC	'00000000000000000000000000000000',3		LC	1
NAME1	DC	'Stath'		LE	50
OPCODE	DC	' ',815		LE	7
DCALC	DE	NAME1		HL	OPCODE
	CALL	MLOAD\$		CALL	KEYIN\$
	JFZ	BOOT\$		HL	OPCODE
DCALCH	DE	0		LAM	
	HL	HEADING		CP	015
	CALL	DSPLY\$		JTZ	ADDOP
	LD	51		CP	'A'
	LE	2		JTZ	ADDOP
	HL	DECPL		CP	'S'
	CALL	KEY\$		JTZ	SUBOP
	LL	INPUT		CP	'M'
	CALL	FILLIN		JTZ	MULOP
	LL	ACCUM		CP	'D'
	CALL	FILLIN		JTZ	DIVOP
	LL	DIVID		CP	'E'
	CALL	FILLIN		JTZ	MOVOP
	LL	DECPL		CP	'R'
	LAM			JTZ	DCALCH
	SU	'0'		EX	BEEP
	LBA			JMP	DCALCL
	LC	'.'		ADDOP	DE INPUT
	LA	INPUT+10			HL ACCUM
	SUB			CALL	ADD\$
	LLA			OVFTST	JFC NOOVF
	LMC				LD 36
	LA	ACCUM+10			LE 3
	SUB				HL OVFMSG
	LLA				CALL DSP\$
	LMC				EX BEEP
	LAB				JMP DCALCL
	SLC			NOOVF	LE 36

	LE	3
	HL	BLANK
	CALL	DSP\$
	JMP	DCALCL
SUBOP	DE	INPUT
	HL	ACCUM
	CALL	SUB\$
	JMP	OVFTST
MULOP	DE	ACCUM
	HL	ACCUM
	CALL	MOV\$
	DE	INPUT
	HL	ACCUM
	CALL	MUL\$
	JMP	OVFTST
MOVOP	DE	INPUT
	HL	ACCUM
	CALL	MOV\$
	JMP	OVFTST
DIVOP	DE	ACCUM
	HL	DIVID
	CALL	MOV\$
	DE	INPUT
	HL	ACCUM
	CALL	DIV\$
	JMP	OVFTST
FILLIN	LAM	
	CP	3
	RTZ	
	LA	'0'
	LMA	
	LAL	
	AD	1
	LLA	
	JMP	FILLIN
	END	DCALC

Observe the addition routine, 'ADDOP'. To add together the inputted string 'input' to the accumulator 'accum' the user only writes the following code as found at 'ADDOP'.

ADDOP	DE	INPUT
	HL	ACCUM
	CALL	ADD\$

Executing this code will cause string 'input' to be added to the string 'accum' with the result in the string 'accum'. The accumulator, it must be realized, is simply a string which the writer of 'DCALC' is using as his result string and he preferred to call it an accumulator.

Note that after each operation there is a jump to 'OVFTST' or as in 'ADDOP', the code is immediately after and executed right after 'ADDOP'. Observe that the first instruction

OVFTST JFC NOOVF

of the overflow test is the actual test: If the carry isn't set then there was no overflow resulting from the operation. If the carry was set, in 'DCALC' the message 'overflow' is printed on the screen as is seen from the code following the 'JFC NOOVF'.

Subtraction behaves the same as addition except for the CALL to SUB\$.

Multiplication and division are slightly different from addition and subtraction but operate similar to each other. Observe the following code as taken from 'DCALC'.

MULOP	DE	ACCUM
	HL	ACCUM
	CALL	MOV\$
	DE	INPUT
	HL	ACCUM
	CALL	MUL\$
	JMP	OVFTST

This demonstrates the requirement, as stated in 7.1.5, that, in MUL\$ and DIV\$, the argument #2 must be the result of the previous move. The reason for this is that multiplication and division really require three 'registers' or strings: The two strings being multiplied and the result. The 'MOV\$' move operation makes a copy of whatever is being moved, during the move, in an internal STATH 'register' string. Therefore, note that the first three instructions in 'MULOP' cause the accumulator to be 'MOV\$' moved to itself. Frequently the user can save time by utilizing this fact in making the last move before calling 'MUL\$' a move of a string involving argument #2. (Again, argument #2 is the argument associated with the H and L registers).

Also note that 'MULOP' tests overflow using the same routine that is used for the other three arithmetic routines 'OVFTST' as described above.

7.1.3 LOADING STATH

STATH may be loaded in memory in either of two ways:

- 1) Incorporating the source code of STATH into the problem source code.
- 2) Catalog STATH as an object file and call it in through the operating system.

The second is preferred and simpler, as is done in 'DCALC'. Once cataloged, the following calls STATH into memory:

NAME1	DC	'STATH'
	DE	NAME1
	CALL	BOOT\$

7.1.3 ADDITION

Entry Point Name	ADD\$
Entry Point Address	10003 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) + (Argument #1)

Action:

Adds two numeric string numbers, rounds, and installs leading blanks and trailing zeros when needed in the result.

Typical calling sequence:

ADD\$	EQU	010003
DE	ARG1	
HL	ARG2	
CALL	ADD\$	

Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the sum of arguments 2 and 1 and will have leading blanks and trailing zeros when needed.

Changes:

The contents of argument 2 are changed to contain the result.

Errors Recognized:

Improper argument format (parity bit set)
Overflow occurrence (carry bit set)

Comparison Flags:

Result was zero (zero bit set)
Result was negative (sign bit set)

7.1.4 SUBTRACTION

Entry Point Name	SUB\$
Entry Point Address	10006 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) - (Argument #1)

Action:

Subtracts one numeric string number from another, rounds and installs leading blanks and trailing zeros when needed in the result.

Typical calling sequence:

SUB\$	EQU	010006
DE	ARG1	
HL	ARG2	
CALL	SUB\$	

Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the difference of arguments 2 and 1 and will have leading blanks and trailing zeros when needed.

Changes:

The contents of argument 2 are changed to contain the result.

Errors Recognized:

Improper argument format (parity bit set)
Overflow occurrence (carry bit is set)

Comparison Flags:

Result was zero (zero bit is set)
Result was negative (sign bit is set)

7.1.5 MULTIPLICATION

Entry Point Name	MUL\$
Entry Point Address	6000 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) X (Argument #1)
Argument Restrictions	Argument #2 must be result of last MOV\$ call

Action:

Multiplies two numeric string numbers, rounds and installs leading blanks and trailing zeros when needed in the result.

Typical calling sequence:

```
MUL$ EQU 06000
      DE   ARG1
      HL   ARG2
      CALL MOV$  
  
      DE   ARG1,
      HL   ARG2
      CALL MUL$
```

Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result. Argument 2 must have been involved in the previous move operation.

Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the product of arguments 2 and 1 and will have leading blanks and trailing zeros when needed.

Changes:

The contents of argument 2 are changed to contain the result.

Errors Recognized:

Improper argument format (parity bit set)
Overflow occurrence (carry bit set)

Comparison Flags:

Result was zero (zero bit set)
Result was negative (sign bit set)

7.1.6 DIVISION

Entry Point Name	DIV\$
Entry Point Address	6003 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #2) / (Argument #1)
Argument Restrictions	Argument #2 must be result of last MOV\$ call

Action:

Divides one numeric string number into another, rounds and installs leading blanks and trailing zeros when needed in the result.

Typical calling sequence:

```
MOV$ EQU 010000
      DE   ARG1
      HL   ARG2
      CALL  MOV$

DIV$ EQU 06003
      DE   ARG1
      HL   ARG2
      CALL  DIV$
```

Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result. Argument 2 must have been involved in the previous move operation..

Result:

The contents of argument 1 (D and E) will remain unchanged.

The contents of argument 2 (H and L) will contain the result of the division of argument 1 into argument 2 and will have leading blanks and trailing zeros when needed.

The number of decimal places in the result is equal to the number of decimal places in the dividend less the number of decimal places in the divisor. This number may not be negative and if it is, the number of decimal places is extended to make the difference zero.

The size of the result equals the size of the extended dividend less the size of the divisor.

Note that the string '10.0' divided by the string '3.0' is the string '3'. It is rounded to ZERO decimal places.

Changes:

The contents of argument 2 are changed to contain the result.

Errors Recognized:

Improper argument format (parity bit set)
Overflow occurrence (carry bit set)

Comparison Flags:

Result was zero (zero bit set)
Result was negative (sign bit set)

7.1.7 COMPARE

Entry Point Name	COM\$
Entry Point Address	10011 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Arguments unchanged. Only sets condition code
Arithmetic Function	(cond-code) = (cond [(Argument #2) - (Argument #1)])

Action:

Compares two numeric string numbers as to magnitude. No change to arguments results. Changes are only made to the condition flags.

Typical calling sequence:

```
COM$ EQU 010011  
DE     ARG1  
HL     ARG2  
CALL   COM$
```

Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

Result:

The contents of both arguments will remain unchanged. Only the condition code will change and will obtain the exact same condition as if a call to SUB\$ were done. Therefore, the resultant condition flags will behave as if the result were to be rounded.

Changes:

The contents of both arguments remain unchanged. Only the condition flags are changed.

Errors Recognized:

Improper argument format (parity bit set)
Overflow occurrence (carry bit set)

Comparison Flags:

Result was zero (zero bit set)
Result was negative (sign bit set)

7.1.8 MOVE

Entry Point Name	MOV\$
Entry Point Address	10000 Octal
Argument #1 Address	D-E Registers
Argument #2 Address	H-L Registers
Result Location	Argument #2
Arithmetic Function	(Argument #2) = (Argument #1)

Action:

Replaces the numeric string number in argument 2 with that of argument 1, rounds and installs leading blanks and trailing zeros when needed in the result.

Typical calling sequence:

```
MOV$ EQU 010000  
DE     ARG1  
HL     ARG2  
CALL   MOV$
```

Arguments:

Arguments must be each numeric strings of less than 126 characters in length. Argument 1 is addressed by the D and E Registers. Argument 2 is addressed by the H and L Registers and will contain the result.

Result:

The contents of argument 1 (D and E) will remain unchanged.
The contents of argument 2 (H and L) will contain the number of argument 1 rounded and reformatted if necessary.

Changes:

The contents of argument 2 are changed to contain the result.

Errors Recognized:

Improper argument format (parity bit set)
Overflow occurrence (carry bit set)
[Note that overflow can occur in a MOV\$ if a move from a larger to smaller field is attempted]

Comparison Flags:

Result was zero (zero bit set)
Result was negative (sign bit set)

7.1.9 MOD10 CHECKSUM CALCULATION

Entry Point Name	MOD10\$
Entry Point Address	6006 Octal
Argument #1 Address	H-L Registers
Result Location	A-Register (no reformatting of argument)
Arithmetic Function	(A Reg) = Check-MOD-10 (Argument #1)

Action:

Checks validity of Modulo 10 checksum of a numeric string number.

Typical calling sequence:

```
MOD10$ EQU 06006
        HL    ARG1
        CALL  MOD10$
```

Arguments:

The argument must be a numeric string of less than 126 characters in length. Argument 1 is addressed by the H and L Registers.

Result:

- The contents of the argument remains unchanged.
- The carry bit is set if the check digit is 10.
- The zero bit is set if the check digit is not 10.
- The check digit is in the A-Register upon return.

Changes:

- The contents of the argument remain unchanged.

Errors Recognized:

- Improper argument format (parity bit set)

Comparison Flags:

- Check digit was 10 (carry bit set)
- Check digit was not 10 (zero bit set)

7.1.10 MOD11 CHECKSUM CALCULATION

Entry Point Name	MOD11\$
Entry Point Address	6011 Octal
Argument #1 Address	H-L Registers
Result Location	A-Register (no reformatting of argument)
Arithmetic Function	(A Reg) = Check-MOD-10 (Argument #1)

Action:

Verifies the Modulo 11 checksum of the numeric string number.

Typical calling sequence:

```
MOD11$ EQU 06011
        HL    ARG1
        CALL  MOD11$
```

Arguments:

The argument must be a numeric string of less than 126 characters in length. The argument is addressed by the H and L Registers.

Result:

- The contents of the argument remains unchanged.
- The carry bit is set if the check digit is 11.
- The zero bit is set if the check digit is not 11.
- The A-Register contains the check digit.

Changes:

- The contents of the argument remain unchanged.

Errors Recognized:

- Improper argument format (parity bit set)

Comparison Flags:

- Check digit was 11 (carry bit set)
- Check digit was not 11 (zero bit set)

7.1.11 KEYBOARD FORMATTED INPUT

Entry Point Name	KEY\$
Entry Point Address	10014 Octal
Argument #1 Address	H-L Registers
Extra Parameters	(D Reg) = Column. (E Reg) = Row for cursor
Input Function	(Argument #1) = (Keyed in number)
Input Restrictions	Screen format and, therefore, keyed in number has same format as originally in Argument #1

Action:

Provides formatted input from the keyboard into a numeric string. The format is maintained on the screen and only a number fitting the format can be entered. The inputted numeric string is placed in argument 1.

Typical calling sequence:

```
KEY$ EQU 010014
LD    COLUMN
LE    ROW
HL    ARG1
CALL  KEY$
```

Arguments:

The argument must be a formatted numeric string. The D and E Registers must contain the column and row of the cursor position of the first character to be typed in.

Result:

The contents of argument 1 are replaced by the inputted number.
Striking the enter key with no input will cause the argument to be replaced with a zero.
The H and L Registers are pointing immediately after the ETX.

Changes:

The contents of the argument are replaced with the inputted string

Errors Recognized:

Improper argument format (parity bit set)

Comparison Flags:

Result was zero (zero bit set)

Result was negative (sign bit set)

7.1.12 DISPLAY STRING

Entry Point Name	DSP\$
Entry Point Address	10017 Octal
Argument #1 Address	H-L Registers
Extra Parameters	(D Reg) = Column. (E Reg) = Row for cursor
Input Functions	(Display starting at D,E) = (Argument #1)
Input Restrictions	None. May even be non-numeric string

Action:

Displays a string onto the screen. String may be non-numeric.

Typical calling sequence:

```
DSP$ EQU 010017
LD    COLUMN
LE    ROW
HL    ARG1
CALL  DSP$
```

Arguments:

The argument may be a numeric or non-numeric string as long as it terminates with an ETX. The D and E Registers contain the column and row of the location of the first character of the string.

Result:

The string in argument 1 is displayed on the screen starting at the cursor location beginning with the column and row specified by the D and E Registers. The H and L Registers point the location immediately after the ETX in argument 1.

Changes:

The contents of the argument remain unchanged.

Errors Recognized:

None

Comparison Flags:

None

7.2.1 INTRODUCTION

FPAK is a subroutine package which gives the Datapoint 2200 the capability of performing numerical operations with numbers in the range of -10^{38} to 10^{37} . This is accomplished by representing all numbers in a form called "floating point." Floating point notation is a shorthand method of number representation and is very similar to the familiar "scientific notation" used in technical work.

FPAK also provides conversion of floating point numbers to and from 16 bit binary integers, particularly attractive for analyzing binary data gathered by the 2200 from instrumentation systems.

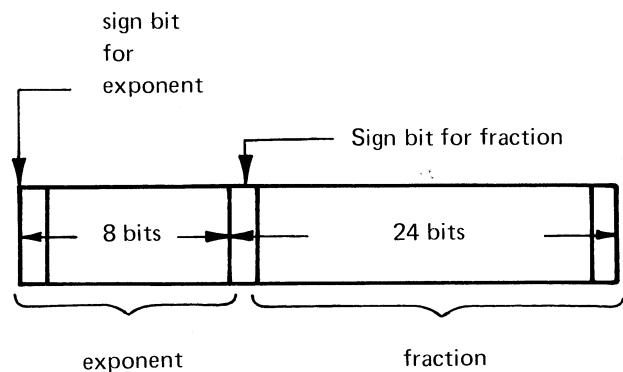
Supplied with FPAK is FCON which supplies the user with simple conversion of ASCII numeric string representation (suitable for displaying or printing) to and from the internal floating point representation.

The CTC 2200 floating point software consists of two main sections: FCON, the conversion section and FPAK, the arithmetic section. The conversion section converts a floating point number to an ASCII string, and visa versa. ASCII is the character code used by the 2200 keyboard and CRT display, and this section of the floating point software allows the user to interface with the computer. The user can enter numbers in a form familiar to him and read the results in a similar form. The ASCII string which the user enters through the keyboard (or from tape or some other means) is converted into the internal floating point form outlined in Section 7.2.1.1. When all arithmetic operations are completed, the user can request that the result be converted back into ASCII for display on the CRT (or for output to tape or for some other use).

7.2.1.1 INTRODUCTION TO FLOATING POINT REPRESENTATION

A number (N) in floating point form consists of two parts within the computer's memory--the "exponent" (e) and the "fraction" (f) --such that: $N = f \cdot 2^e$ (where the \cdot denotes the multiplication operation). The exponent occupies one byte (word) in the 2200 and is an 8 bit signed integer. Thus, exponents on the 2200 can have a range of 127 to -128.

The fraction (sometimes called the "mantissa") on the 2200 occupies three bytes and is a 24 bit signed quantity. Like a decimal fraction (such as .5 or .0001), the fraction in a floating point number has a "decimal point," although "decimal point" is not what it is called. Its proper name, in a binary fraction, is a "binary point." In the notation used on the 2200, the binary point is located immediately to the right of the sign bit (high order bit) of the fraction. Thus, a picture of a 2200 floating point number would look like this:



The exponent and the fraction are separate parts of the number, and one can be positive while the other is negative. On the 2200, negative numbers are represented in their 2's complement form. Since the floating point representation requires more than one byte (word) on the 2200, a convention is used to address a floating point number in memory. The address of a floating point number is the byte (word) address of the exponent byte of the number. The software will use the addressed byte, and the three bytes immediately following, in whatever operation is being performed. Later in the documentation, reference is made to addressing floating point numbers. In such cases, this means that the MSP of the address of the exponent should be in the H or D register and the LSP of the address should be in the L or E register, depending upon whether the HL or DE pair is being used.

7.2.2 FCON - FLOATING POINT/STRING CONVERSION

7.2.2.1 INTRODUCTION TO FCON - FLOATING POINT/STRING CONVERSION

LOCATION	FUNCTION
FISC (0441)	Floating Internal to String Conversion entry point
FSIC (0444)	Floating String to Internal Conversion entry point
FSCE (04460)	Floating Set Conversion Error Branch entry point
OPER (013403)	Location of Floating point number to be converted to or from ASCII

7.2.2.2 FISC - FLOATING INTERNAL TO STRING CONVERSION

The Floating Internal to String Conversion routine has been designed so that the user need not specify the type of number he is going to supply. That is, as long as the ASCII characters being converted represent a valid, decimal number, the conversion routine can decide what type of number it is (i.e., integer, fraction) and perform the proper conversion without any further instructions. This type of input is referred to as "free form" input.

The result of all Floating Point Arithmetic routines end up at location 013403, labeled 'OPER'. The conversion routine, FISC, converts floating point numbers at OPER into a string beginning at the location specified by the H and L Registers upon execution of the CALL to FISC.

For example, should OPER (and the subsequent 3 bytes) contain the floating point number represented by 123,450,000,000,000,000 the string resulting from a call to FISC would look like this: 1.2345E20 where the ASCII number 1, an octal 61, would appear in the location specified by H and L and the period (an octal 056) in H and L plus 1 etc. A note of caution, FISC does not put a terminating 003 or 015 after the string. To be compatible with the CTOS string routines, the string must be terminated with either 003 or 015. However, FISC, upon return from being called, leaves the H and L registers pointing to the location immediately after the last character in the string. This enables the user to immediately store the terminating character of his choice (003 or 015) in that location upon a return. The following call to FISC will illustrate:

FISC	EQU	04441
HL	String	
CALL	FISC	
LA	015	
LMA		

Note the LA 015 and LMA will install a 015 as the terminating character to the resultant string which is the ASCII representation of the floating point number in OPER.

Name: Floating Internal to String Conversion (FISC)

Action:

Converts a floating point number to its ASCII character representation.

Calling Sequence:

FISC	EQU	04441
	HL	String
	CALL	FISC
	LA	015
	LMA	

Arguments:

OPER contains the number to be converted to ASCII. The H and L registers contain the address of the location, in memory, where the first (leftmost) ASCII character should be placed.

Result:

The floating point number in OPER is converted to its ASCII representation, and the ASCII characters comprising this representation are placed in memory, beginning at the address specified by the contents of the H and L registers and continuing in sequential memory locations. H and L end up pointing to the next location after the last string character enabling the user to store the string termination character of his choice up on the return from FISC.

Changes:

The contents of OPER are destroyed; the previous contents of the output string are destroyed. At the end of the execution of this routine, H and L contain the address of the memory location immediately after the last ASCII character in the converted number.

Errors Recognized:

None.

Comments:

Numbers are represented to six significant (decimal) digits and are rounded where appropriate. The format of the output is "free," with small integer in FORTRAN I format, floating point numbers with decimal exponents between -6 and 6 in FORTRAN F format, and other numbers in FORTRAN E format.

7.2.2.3 FSCI - FLOATING STRING TO INTERNAL CONVERSION

The Floating String to Internal Conversion routine has been designed to convert floating point numbers into the proper ASCII representation. If the floating point numeric string is a small integer, it will be converted to an integer, with no decimal point in the representation. If the numeric string is a large integer, or a noninteger, it will be converted into scientific notation, or more precisely what is known as the FORTRAN E format, such as 1.3456E17.

FSIC converts to internal floating point representation an ASCII numeric string with optional leading minus sign, optional decimal point, and optional trailing FORTRAN E, type exponent, i.e. -1.2345E20. The H and L registers must point to the first character of the string. The result goes into the FPAK 'register' called OPER starting at 013403, ready to be used by FPAK. FSCE, Floating Set Conversion Error Branch, should be set first to cover format problems in the string being converted. A simple call to FSCE with the D and E registers specifying the location of your error recovery routine will set the error branch.

Name: Floating String to Internal Conversion (FSIC)

Action:

Converts an ASCII string, which represents a decimal number, into that number's floating point form.

Calling Sequence:

FSIC	EQU	04444
FSCE	EQU	04460

DE	ERROU	Location of error routine
CALL	FSCE	
HL	String	
CALL	FSIC	

Arguments:

The H and L registers contain the address of the first byte (character) of the ASCII string which represents the number to be converted.

Result:

The character string, if it represents a valid number, is converted to a floating point number, and that value is left in OPER. The result in OPER is normalized and rounded.

Changes:

The original contents of OPER are destroyed; the ASCII string is left unchanged, and upon successful conversion, the H and L registers contain an address of the character which caused termination of the number (i.e., was a character not allowed in the ASCII representation of a number).

Errors Recognized:

Invalid character found while converting from ASCII to floating point.

Comments:

The ASCII string may be in free form, that is, in FORTRAN I, F, or E format. All of those formats will be properly converted by this routine. Conversion stops when an invalid character (something other than a digit, "E", "+", "-", or ".") is encountered after a valid number has been found. An invalid character encountered before a valid number has been found will generate an error. Some of the above characters can be considered invalid if used incorrectly (i.e., a "." in an exponent, such as 1.333E1.5, is an error) and will generate an error condition.

7.2.2.4 FSCE - FLOATING SET CONVERSION ERROR BRANCH

Name: Floating Set Conversion Error Branch (FSCE)

Action:

Specifies the location of the user's routine to be branched to in the event an invalid character is encountered while converting an ASCII representation of a number to the floating point representation of that number.

Calling Sequence:

DE	ERROU	Location of error routine
CALL	FSCE	

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Errors Recognized:

None.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the contents of the A register will be non-zero, and the result in OPER will, in general, be erroneous. The user's error routine may end with a return if the user wishes to continue execution immediately after the call to the routine which generated the error.

7.2.3 FPAK - FLOATING POINT ARITHMETIC PACKAGE

7.2.3.1 INTRODUCTION TO FPAK - FLOATING POINT ARITHMETIC PACKAGE

LOCATION	FUNCTION
FCMP (04422)	Floating Point Compare
FADD (04400)	Floating Point Addition
FSUB (04403)	Floating Point Subtraction
FMUL (04406)	Floating Point Multiplication
FDIV (04411)	Floating Point Division
FLOD (04414)	Floating Point Load [memory to 'OPER']
FSTO (04417)	Floating Point Store ['OPER' to memory]
FNEG (04425)	Floating Point Negate [Two's complement]
FABS (04430)	Floating Point Absolute Value
FSTL (04463)	Floating Point Set Tolerance [For Equal Flag]
FFIX (04433)	Floating Point Fix [to 16 bit integer]
FFLT (04436)	Floating Point Float Conversion from 16 bit integer
FSOV (04447)	Floating Point Set Overflow Error Branch
FSUN (04452)	Floating Point Set Underflow Error Branch
FSDV (04455)	Floating Point Set Divide Check Error Branch

The second section of the floating point software is the arithmetic part. This section contains the routines for performing the common arithmetic operations of add, subtract, multiply, divide, compare, negate, and absolute value, and two routines for converting between integer and floating point formats (an integer, in the floating point software, is a 16 bit (2-byte) signed quantity which is addressed by specifying the address of the high order byte).

Within the floating point software package is a 4-byte area called OPER. OPER is to the floating point software what the A register is to the 2200 processor. Floating point operations are performed on numbers in OPER, or on pairs of numbers, one of which is in OPER and the other in memory. The software supplies two routines, FLOD and FSTO which provide the user with the capability of copying numbers from memory to OPER and from OPER to memory.

With two exceptions, all of the routines in the arithmetic part of the floating point software, which take floating point numbers as their arguments, expect their operands to be "normalized." Normalization is nothing more than an agreed upon standard for writing a floating point num-

ber. A number is considered normalized if the sign bit of the fraction and the bit immediately to the right of the sign bit (the high order bit of the fraction) are unequal. Thus, a positive fraction (sign bit 0) has a 1 as its high order bit, and a negative fraction (sign bit 1) has a 0 as its high order bit. This convention makes sure that the maximum precision possible is maintained in all floating point operations.

As a rule, all routines expect their floating point operands to be normalized. The significant exceptions to this rule are the add and subtract routines, FADD and FSUB. If the user is adding or subtracting two numbers, the numbers should be normalized for a result with the greatest accuracy possible. However, if the user has a floating point number which is not normalized, he can convert the number to its normalized form by adding or subtracting a "normal" 0 to or from the unnormalized number. A normal 0 has a fraction equal to 0 and an exponent of -128 (200 octal). Except in this case, it is not recommended that the user perform operations on unnormalized numbers.

7.2.3.2 ERROR CONDITIONS

There are several error conditions that can arise during the course of executing routines in the floating point software package. These errors are:

exponent overflow
exponent underflow
divisor of 0 (in FDIV)

For these errors, a flag (see below) is set to 1 when the error is detected. For all of these errors, an "error branch" is provided. When the error condition arises, the appropriate flag (or A register) is set, and a jump is made to a location in the floating point software package. This location contains a jump to the address of either a user-specified error routine or a return instruction (the default case if the user does not supply an error routine). There is a separate location for each error condition, and there are three routines -- FSOV, FSUN, and FSDV -- which are used to set or change the address of the error routines.

The error conditions and their respective flags are:

Exponent Underflow	UNFLO	Location 013400
Exponent Overflow	OVFLO	Location 013401
Divide by 0	DVDCK	Location 013402

If an error condition arises, the flag is set to 1 and a branch is made to the error routine address. If no error condition arises, the flag is set to 0, and a normal return from the routine occurs.

7.2.3.2.1 FSOV - FLOATING SET OVERFLOW ERROR BRANCH

Name: Floating Set Overflow Error Branch (FSOV)

Action:

Specifies the location of the user's routine to be branched to in the event an operation causes exponent overflow (the value of the binary exponent in the result is greater than 127).

Calling Sequence:

Execute CALL instruction location 04447.
See 7.2.3.2.

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Errors Recognized:

None.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the appropriate error flag will be set to 1, and the result in OPER will, in general, be erroneous. The user's error routine should not end with a return since that would cause processing to continue in the floating point software with incorrect values in the machine registers.

7.2.3.2.2 FSUN - FLOATING SET UNDERFLOW ERROR BRANCH

Name: Floating Set Underflow Error Branch (FSUN)

Action:

Specifies the location of the user's routine to be branched to in the event an operation causes exponent underflow (the value of the binary exponent in the result is less than -128).

Calling Sequence:

Execute CALL instruction to location 04452.
See 7.2.3.4.

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the appropriate error flag will be set to 1, and the result in OPER will, in general, be erroneous. The user's error routine should not end with a return since that would cause processing to continue in the floating point software with incorrect values in the machine registers.

7.2.3.2.3 FSDV - FLOATING SET DIVIDE CHECK ERROR BRANCH

Name: Floating Set Divide Check Error Branch (FSDV)

Action:

Specifies the location of the user's routine to be branched to in the event the divisor in a floating divide operation is 0.

Calling Sequence:

Execute CALL instruction to location 04455.
See 7.2.3.7.

Arguments:

The D and E registers contain the address of the error routine.

Result:

The location in the floating point software which specifies the location of the error routine is set to the address provided by the user in the D and E registers.

Changes:

The previous error routine address is destroyed.

Errors Recognized:

None.

Comments:

In the event the user does not specify an error routine location, the floating point software will execute a return (RET instruction) if the error condition arises, and the arithmetic routine called by the user will continue to completion. At the conclusion of that routine, the appropriate error flag will be set to 1, and the result in OPER will, in general, be erroneous. The user's error routine may end with a return if the user wishes to continue execution immediately after the call to the routine which generated the error.

7.2.3.3 FLOATING COMPARE

Name: Floating Compare (FCMP)

Action:

Compares, algebraically, two floating point numbers.

Calling Sequence:

FSTL	EQU	04463	Only necessary to EQU once per program
FCMP	EQU	04422	
	LA	TLRNC	Where TLRNC is the comparison tolerance only necessary once per program if tolerance doesn't change
	CALL	FSTL	
	HL	NUMBER	Number will be compared with OPER
	CALL	FCMP	

Arguments:

OPER contains one of the floating point numbers being compared, and the contents of the H and L registers address the other floating point number being compared.

Result:

Floating Compare sets the Sign and Zero flip-flops as if a subtraction of the floating point number addressed by the contents of the H and L registers from the floating point number in OPER had taken place. However, if the absolute value of the difference is less than or equal to the tolerance specified (see the description of the routine FSTL for an explanation of how the tolerance is specified), then the Sign and Zero flip-flops are set as if both floating point numbers were found to be equal.

Changes:

Neither operand is altered by the Floating Compare operation.

Errors Recognized:

None.

Comments:

Since representations of decimal fractions in a binary machine are approximate, the Floating Compare operation allows for an "approximate" compare by allowing the user to specify how close two numbers may be before they are considered equal.

7.2.3.4 FLOATING ADD

Name: Floating Add (FADD)

Action:

Adds two floating point numbers, rounds and normalizes the result.

Calling Sequence:

FSOV	EQU	04447	Only necessary to EQU once per program
FSUN	EQU	04452	
FADD	EQU	04400	
DE	OVERR		Only necessary to set these once per program or until it is desired to change.
CALL	FSOV		
DE	UNERR		Where OVERR and UNERR are addresses of user and recovery routines.
CALL	FSUN		
HL	NUMBER		
CALL	FADD		Number will be added to OPER

Arguments:

OPER contains one of the floating point numbers, and the contents of the H and L registers address the other floating point number.

Result:

The contents of OPER and the floating point number addressed by the contents of the H and L registers are added together with the result left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

Maximum precision is maintained by having both operands normalized; however, an unnormalized number may be converted to its normalized form by using this routine to subtract a "normal" 0 from the unnormalized number.

7.2.3.5 FLOATING SUBTRACT

Name: Floating Subtract (FSUB)

Action:

Subtracts two floating point numbers, rounds and normalizes the result.

Calling Sequence:

FSUB is identical to FADD except the program must now contain a FSUB EQU 04403 and the last statement in calling sequence is:

CALL	FSUB	Number will be subtracted from OPER
------	------	-------------------------------------

Arguments:

OPER contains the minuend and the contents of the H and L registers address the subtrahend.

Result:

The floating point number addressed by the contents of the H and L registers is subtracted from the floating point number in OPER, and the result is left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

Maximum precision is maintained by having both operands normalized; however, an unnormalized number may be converted to its normalized form by using this routine to subtract a "normal" 0 from the unnormalized number.

7.2.3.6 FLOATING MULTIPLY

Name: Floating Multiply (FMUL)

Action:

Multiplies two floating point numbers, rounds and normalizes the result.

Calling Sequence:

FMUL is identical to FADD except the program must now contain a FMUL EQU 04406 and the last statement in the calling sequence is:

CALL	FMUL	Number will multiply OPER
------	------	---------------------------

Arguments:

OPER contains the multiplicand, and the H and L registers contain the address of the multiplier.

Result:

The floating point of OPER and the floating point number addressed by the contents of the H and L registers are multiplied together with the result left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

This routine expects both operands to be normalized. If one or both of the operands is not normalized, erroneous results may occur.

7.2.3.7 FLOATING DIVIDE

Name: Floating Divide (FDIV)

Action:

Forms the quotient of two floating point numbers, rounds and normalizes the result.

Calling Sequence:

<table border="0"> <tr> <td style="text-align: center;">FSOV</td> <td style="text-align: center;">EQU</td> <td style="text-align: center;">04447</td> <td style="text-align: center;">Only necessary to EQU these once per program</td> </tr> <tr> <td style="text-align: center;">FSUN</td> <td style="text-align: center;">EQU</td> <td style="text-align: center;">04452</td> <td></td> </tr> <tr> <td style="text-align: center;">FSDV</td> <td style="text-align: center;">EQU</td> <td style="text-align: center;">04455</td> <td></td> </tr> </table> <table border="0"> <tr> <td style="text-align: center;">FDIV</td> <td style="text-align: center;">EQU</td> <td style="text-align: center;">04411</td> <td></td> </tr> <tr> <td style="text-align: center;">DE</td> <td style="text-align: center;">OVERR</td> <td style="text-align: center;">Only necessary to set these once per program or when it is desired to change recover routine.</td> <td></td> </tr> <tr> <td style="text-align: center;">CALL</td> <td style="text-align: center;">FSOV</td> <td></td> <td>Where OVERR, UNERR, and CKERR are addresses of user error recovery routines.</td> </tr> <tr> <td style="text-align: center;">DE</td> <td style="text-align: center;">UNERR</td> <td></td> <td></td> </tr> <tr> <td style="text-align: center;">CALL</td> <td style="text-align: center;">FSUN</td> <td></td> <td></td> </tr> <tr> <td style="text-align: center;">DE</td> <td style="text-align: center;">CKERR</td> <td></td> <td></td> </tr> <tr> <td style="text-align: center;">CALL</td> <td style="text-align: center;">FSDV</td> <td></td> <td></td> </tr> </table> <table border="0"> <tr> <td style="text-align: center;">HC</td> <td style="text-align: center;">NUMBER</td> <td></td> <td></td> </tr> <tr> <td style="text-align: center;">CALL</td> <td style="text-align: center;">FDIV</td> <td style="text-align: center;">Number divides OPER</td> <td></td> </tr> </table>	FSOV	EQU	04447	Only necessary to EQU these once per program	FSUN	EQU	04452		FSDV	EQU	04455		FDIV	EQU	04411		DE	OVERR	Only necessary to set these once per program or when it is desired to change recover routine.		CALL	FSOV		Where OVERR, UNERR, and CKERR are addresses of user error recovery routines.	DE	UNERR			CALL	FSUN			DE	CKERR			CALL	FSDV			HC	NUMBER			CALL	FDIV	Number divides OPER		<table border="0"> <tr> <td style="text-align: center;">DE</td> <td style="text-align: center;">OVERR</td> <td style="text-align: center;">Only necessary to set these once per program or when it is desired to change recover routine.</td> <td></td> </tr> <tr> <td style="text-align: center;">CALL</td> <td style="text-align: center;">FSOV</td> <td></td> <td>Where OVERR, UNERR, and CKERR are addresses of user error recovery routines.</td> </tr> <tr> <td style="text-align: center;">DE</td> <td style="text-align: center;">UNERR</td> <td></td> <td></td> </tr> <tr> <td style="text-align: center;">CALL</td> <td style="text-align: center;">FSUN</td> <td></td> <td></td> </tr> <tr> <td style="text-align: center;">DE</td> <td style="text-align: center;">CKERR</td> <td></td> <td></td> </tr> <tr> <td style="text-align: center;">CALL</td> <td style="text-align: center;">FSDV</td> <td></td> <td></td> </tr> </table>	DE	OVERR	Only necessary to set these once per program or when it is desired to change recover routine.		CALL	FSOV		Where OVERR, UNERR, and CKERR are addresses of user error recovery routines.	DE	UNERR			CALL	FSUN			DE	CKERR			CALL	FSDV		
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CALL	FSUN																																																																								
DE	CKERR																																																																								
CALL	FSDV																																																																								

Arguments:

OPER contains the dividend, and the H and L register contain the address of the divisor.

Result:

The floating point number in OPER is divided by the floating point number addressed by the contents of the H and L registers with the result left in OPER.

Changes:

The contents of OPER are altered; the floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

Exponent overflow, exponent underflow, divisor equal to 0.

Comments:

This routine expects both operands to be normalized. If one or both of the operands is not normalized, erroneous results may occur.

7.2.3.8 FLOATING LOAD

Name: Floating Load (FLOD)

Action:

Copies a floating point number from its location in memory to OPER.

Calling Sequence:

FLOD	EQU	04414	Only necessary to EQU this once per program.
HC CALL	NUMBER FLOD		Number is loaded into OPER

Arguments:

The H and L registers contain the address of the floating point number that is to be copied into OPER.

Result:

The floating point number addressed by the H and L registers is copied into OPER.

Changes:

The original contents of OPER are destroyed. The floating point number addressed by the contents of the H and L registers is unchanged.

Errors Recognized:

None.

Comments:

None.

7.2.3.9 FLOATING STORE

Name: Floating Store (FSTO)

Action:

Copies a floating point number from OPER to memory.

Calling Sequence:

FSTO	EQU	04417	Only necessary to EQU this once per program.
HL CALL	NUMBER FSTO		Number is loaded from OPER

Arguments:

The H and L registers contain the address of the location, in memory, to which the floating point number is to be copied.

Result:

The floating point number is copied into the location addressed by the contents of the H and L registers.

Changes:

The original contents of memory (4 bytes) addressed by the H and L registers are destroyed. The contents of OPER are unchanged.

Errors Recognized:

None.

Comments:

None.

7.2.3.10 FLOATING NEGATE

Name: Floating Negate (FNEG)

Action:

Forms the two's complement of the floating point number in OPER.

Calling Sequence:

FNEG	EQU	04425	Only necessary to EQU this once per program.
CALL	FNEG		OPER is negated

Arguments:

OPER contains the floating point number to be negated.

Result:

The number in OPER is converted to two's complement form and then this result is normalized. The final result is left in OPER.

Changes:

The original contents of OPER are destroyed.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

None.

7.2.3.11 FLOATING ABSOLUTE VALUE

Name: Floating Absolute Value (FABS)

Action:

Forms the absolute value of a floating point number.

Calling Sequence:

FABS	EQU	04430	Only necessary to EQU this once per program
CALL	FABS		OPER becomes the absolute value of OPER

Arguments:

OPER contains the floating point number whose absolute value is to be computed.

Result:

If the contents of OPER are greater than or equal to zero, then they are left unchanged. Otherwise, the contents of OPER are negated (see the description of FNEG). In the latter case, the original contents of OPER are destroyed.

Changes:

Contents of OPER are destroyed if they are less than zero; otherwise, the contents of OPER are unchanged.

Errors Recognized:

Exponent overflow, exponent underflow.

Comments:

None.

7.2.3.12 FLOATING FIX

Name: Floating Fix (FFIX)

Action:

Converts a floating point number into a 16 bit integer.

Calling Sequence:

FFIX	EQU	04433	Only necessary to EQU this once per program
HL CALL	NUMBER FFIX		Number and number+1 will contain the 16 bit integer made from OPER

Arguments:

OPER contains the floating point number to be fixed (converted to an integer), and the H and L registers contain the address, in memory, of the high order byte (upper eight bits of the integer) where the integer is to be placed.

Result:

The floating point number is converted to a 16 bit integer. If the number has a fractional part, that part is lost. The 16 bit integer is stored in memory beginning at the byte addressed by the contents of the H and L registers.

Changes:

The original contents of the 16 bits addressed by the contents of the H and L registers are destroyed. The contents of OPER are unchanged.

Errors Recognized:

None.

Comments:

If the number in OPER is such that it cannot be represented in 16 bits, only the low order 16 bits are stored in memory. Any higher order bits are lost.

7.2.3.13 FLOAT

Name: Float (FFLT)

Action:

Converts a 16 bit integer into a normalized floating point number.

Calling Sequence:

FFLT	EQU	04436	Only necessary to EQU this once per program
HL CALL	NUMBER FFIX		The 16 bit integer in number and number+1 will be converted to floating point in OPER

Arguments:

The H and L registers contain the address, in memory, of the high order byte (high order eight bits) of the 16 bit integer that is to be converted.

Results:

The 16 bit integer is converted from its integer form to the floating point form, and the result is normalized and left in OPER.

Changes:

The original contents of OPER are destroyed. The 16 bit integer addressed by the H and L registers is unchanged.

Errors Recognized:

None.

Comments:

None.

7.2.3.14 FLOATING SET COMPARE TOLERANCE

Name: Floating Set Compare Tolerance (FSTL)

Action:

Specifies a range in which the difference of two floating point numbers must lie for the two numbers to be considered equal.

Calling Sequence:

See 7.2.3.3.

Arguments:

The A register contains the tolerance as a positive eight bit integer (the high order bit of the A register must be 0).

Result:

The location in the floating point software which specifies the floating point compare tolerance is set to reflect the value provided by the user in the A register.

Changes:

The previous value of the tolerance is destroyed.

Errors Recognized:

None.

Comments:

When the floating point package is initialized, the tolerance is set as if the user had called FSTL with a 2 in the A register. If the value in the A register is less than .0 when FSTL is called, erroneous results may occur when using the floating compare routine, FCMP.

SECTION 8

DATAPOINT 2200

COMMUNICATIONS SUBROUTINES

SECTION 8

1. INTRODUCTION

Interfacing the Datapoint 2200 with a wide range of communication facilities is a simple task. All that is needed is the 2210 ACA Communications Adaptor with the required data set or keyer option and the necessary software subroutines to drive it. The software subroutines may or may not have been written for a particular application. However, it seems likely that most users will choose to develop their own to fit their particular needs. This chapter is devoted to aiding the user in fulfilling this goal.

Understanding communications subroutines is useful for many reasons. It enables use of communication disciplines not previously used to fill a specialized need. It enables a user to develop routines that are most efficient for his particular application, it permits a user to modify previously written routines for special purposes and provides greater insight into how the communications system functions.

There is nothing difficult about the communications routines. They are just another part of the user's applications program. The routines are given special treatment here because they are used so frequently and because the terminology and hardware used for communication is foreign to many users.

In addition to the material covered in this chapter the user should be familiar with material covered in other publications on the subject of data communications. Two references that are highly recommended before embarking on any communications oriented 2200 applications are:

Bell System Data Communications Technical Reference Manual*
Martin, James; Teleprocessing Network Organization;
Prentice - Hall, 1970

2. TYPES OF SUBROUTINES

As in most modern computers, the input/output devices used with the Datapoint 2200 are much slower than the 2200 processor. In order for an input/output (I/O) routine to be efficient it must be possible for the processor to perform other tasks (including other I/O operations) while a given I/O routine is active. One approach is to use an interrupt system in the processor to stop one routine and give control to another when an I/O operation is needed. The Datapoint 2200 does not have an interrupt system but in its place it has a very powerful subroutine calling mechanism

which permits many separate I/O routines to be "scanned" during normal execution of a program so that several I/O or other subroutines can be active at the same time.

This leads to the two possible types of communications subroutines: "in-line" and "interleaved". In-line subroutines are those routines which are written in such a way that whenever they are called they "capture" the processor until their function is complete and hence do not permit any other subroutine to be active at the same time. In many situations in-line subroutines are all that is required (such as during an automatic dialing operations when the 2200 has no other functions to perform). Interleaved subroutines are written in such a way that they return to the calling routine at regular intervals while they are active - to be called again to complete their work. Return points in communications subroutines frequently occur following status checks of external devices so that the communications subroutine does not sit in a "tight-loop" waiting for some external operation to be completed.

All of the I/O routines in the CTOS (Operating System) are in-line and would not be used during interleaved operations.

3. INPUT/OUTPUT OPERATIONS

In order to write any type of I/O subroutine for the Datapoint 2200 it is necessary to have a working knowledge of the input/output section of the processor. All 2200 I/O devices (including the CRT, keyboard and tape cassette decks as well as the Communications Adaptor) operate alike and have the same general I/O structure.

The basic physical details of the I/O structure are given in Part 4 of the Datapoint 2200 Reference Manual. We will deal with this system here from a programmer's point of view.

3.1 Data Buses

Data flow to and from the processor takes place over a set of I/O data lines connected to the A-register in the processor. Output data is transmitted from the A-register by eight wires which at all times reflect the contents of the A-register. Whenever the content of the A-register is to be transmitted to an I/O device, one of the external command instructions is executed, which causes one of the External Command Strobes to pulse a signal to the I/O device, informing it that the data on the output bus is for it, and should be read.

*Obtained through Engineering Director—Data Communications, American Telephone and Telegraph Co., 195 Broadway, N.Y., N.Y. 10007

Input data is transmitted to the A-register in the processor by eight wires which form a bus connected to all I/O devices. Each I/O device is so arranged that only the one currently addressed will have access to this bus. Normally, when an I/O device is first addressed, a status word is placed on this bus. The status word (or whatever is placed on the bus) is loaded into the A-register whenever an INPUT instruction is executed.

3.2 External Command Strobes

The Datapoint 2200 processor has 24 External Command Strobes in its I/O structure, only eight of which are brought to devices outside of the 2200 proper (e.g. the Communications Adaptor) and need be considered here.

These eight command lines are physically identical, and their functions are pre-assigned in the table below for the sake of consistency between I/O devices.

ADR	EQU	1
STATUS	EQU	2
DATA	EQU	3
WRITE	EQU	4
COM1	EQU	5
COM2	EQU	6
COM3	EQU	7
COM4	EQU	8

(In all examples following in this chapter it is assumed that all External Command labels have been defined.)

When an External Command is executed, physically all that occurs in the processor is a pulse (or strobe) on the indicated command line. All other action occurs in one of the I/O devices.

- a. EX ADR is the only command strobe acted upon by all I/O devices at the same time. All other command strobes affect only the I/O device that is currently addressed.

EXTERNAL COMMAND

COMMAND NUMBER	(exp)	OCTAL CODE	COMMAND	DESCRIPTION
1	ADR	121	Address	Selects device specified by A-register
2	STATUS	123	Sense Status	Connects selected device data lines to data input bus
3	DATA	125	Sense Data	Connects selected device data lines to data input bus
4	WRITE	127	Write Strobe	Signals selected device that output data is on data output lines
5	COM1	131	Command 1	Signals selected device that a control word is on data output lines
6	COM2	133	Command 2	Signals selected device that a control word is on data output lines
7	COM3	135	Command 3	Signals selected device that a control word is on data output lines
8	COM4	137	Command 4	Signals selected device that a control word is on data output lines

When external commands are to be used in a program the names or labels for the commands used should be defined to the assembler at the beginning of the source code listing as in the following example:

- b. EX STATUS causes the selected device to place its status word on the input bus (it may already be on the bus in which case the EX STATUS does nothing).

c. EX DATA causes the selected device to place its data word on the input bus. This data will remain there until an EX STATUS or an EX ADR is executed.

LA 0322
EX ADR

d. EX WRITE- The write strobe command is a signal from the processor that data is present on the data output lines for the selected external device.

e. EX COM1 thru EX COM4 are used generally to load command words into I/O device command word registers. Depending on the device, however, they may be used for any purpose.

Device addressing in the Datapoint 2200 follows an unusual convention which the programmer should be aware. Up to 16 devices may be addressed, and the first four (low order) bits of the address word indicate which address is selected (zero through fifteen). The second four (high order) bits of the address word must contain the binary complement of the first four bits. Some of the sixteen possible addresses are reserved for specific devices. The remaining ones may be assigned as needed for a particular application.

DEVICE ADDRESS ASSIGNMENTS

DEVICE	NUMBER	BINARY	OCTAL
Cassette Tape Decks	0	11110000	360
CRT/Keyboard	1	11100001	341
Communications Adaptor	2	11010010	322
2200P Printer	3	11000011	303
2200T Tape Transport	4	10110100	264
Unassigned	5	10100101	245
"	6	10010110	226
"	7	10000111	207
"	8	01111000	170
"	9	01101001	151
"	10	01011010	132
"	11	01001011	113
"	12	00111100	074
"	13	00101101	055
"	14	00011110	036
"	15	00001111	017

By way of example, to address (or select) the Communications Adaptor (and de-address all other devices) the following instructions are all that is required:

LA 0322
EX ADR

3.3 The Input Command

In order to load the A-register with whatever is on the input bus an INPUT instruction is executed. In addition to loading the A-register with a new value, it transmits a strobe to the selected external device to inform it that the input bus has been read. Generally, if the status word is on the input bus, the input strobe is of no interest to the I/O device. However, if the data word from the I/O device is on the input bus, then the input strobe informs the I/O device that it has been read by the processor and the device then clears the read ready status bit.

3.4 Command Words

Through the use of the EX COM1 through EX COM4 strobes, it is possible to load command words in an I/O device, which causes the device to carry out specific instructions, or to assume some specific configuration. An excellent example of a command word structure is shown in paragraph 8.5 of the Datapoint 2200 Reference Manual. Each bit of the command word affects some aspect of the Communications Adaptor configuration and the entire operating mode of the adaptor is determined by the content of the Command Word Register at any given time.

When an EX COMn is executed, all of the bits in the affected command word register are loaded from the A-register so care must be taken that all eight bits are accounted for whenever a change is made in a command word register. Generally, when a word is loaded into a command word register, it remains there until another one replaces it. In some devices, (not the Communications Adaptor) a bit set to one will return to zero automatically when some function is carried out.

To give an example; suppose it was desired to instruct the Communications Adaptor to go "off-hook" and to "send 2025 Hz". The device would be addressed:

and then a command word loaded

LA	060
EX	COM1
	,

Where 060 is the octal value of the command word.

3.5 The Status Words

The status word provides a means of communicating to the processor the state of an I/O device at any given time. The status word is placed on the input bus whenever an I/O device is addressed, and remains there until the device is de-addressed or an EX DATA is executed. If the status of the device changes while it is selected, the value on the input bus changes with it, and may be read into the A-register without re-addressing the device. Paragraph 8.2 of the 2200 Reference Manual provides a detailed example of the status word structure used in the Communications Adaptor.

If it were desired to jump to a subroutine if the "Ring-
ing Present" bit of this status word were to come true
it could be coded as follows:

LA	0322	ADDRESS DEVICE
EX	ADR	
INPUT		INPUT STATUS WORD
ND	040	MASK OFF ALL OTHER BITS
JFZ	SUBR	JUMP TO SUBROUTINE IF A CONTAINS A ONE
	,	

3.6 Character Buffers

An I/O device generally has one or more registers or buffers used to hold characters (also called "data") which are being transmitted or received by the device. Slow devices such as the keyboard usually have only one character buffer since the processor has plenty of time to read a character from the buffer before another is loaded.

Faster devices such as the Communications Adaptor have a double character buffer for transmitting or receiving data so that the processor may be reading (or writing) from one buffer while a data set (or some other external equipment) is writing (or reading) to the other buffer. This means that the processor always has at least one full character time in which to service the Communications Adaptor between data transfers.

Some even faster devices (such as the 2200T IBM Compatible Tape Deck) buffer an entire string of characters (up to 1024 in this example).

4. SIMPLE COMMUNICATIONS ADAPTOR ROUTINES

In writing any routine for the Communications Adaptor some simple rules must be followed. Reference should be made to Section 8 of the Datapoint 2200 Programmer's Manual in order to understand the following discussion.

Whenever data is to be transmitted or received through the Communications Adaptor, the device must first be configured for the mode of operation to be used. This is generally done with a prep subroutine which sets the Communications Adaptor Command Word (EX COM1). The transmit and receive time base registers (EX COM3 and EX COM2) and the Character Length Mask Word (EX COM4).

4.1 External Printers

Suppose it is desired to drive an external printer such as the Datapoint 3300P from the Communications Adaptor (The Datapoint 2200P connects directly to the I/O bus and does not use the Communications Adaptor). The 3300P is an EIA RS-232 interface serial printer, operates at 300 baud (bits/second)

uses an 8-information-bit code, and works best with two stop units.

Referring to paragraph 8.5 of the Reference Manual we see that the Command Word can be all zeros. (No data set is involved and neither transmit or received data is inverted).

Referring to paragraph 8.6 we see that to transmit 300 baud the transmit time base must be loaded with 377 followed by 000. The receive time base need not be set since we are only transmitting to a printer.

Referring to paragraph 8.7 we see that the transmitted character length mask must be 001 (binary) and the receive character length can be 000 (binary) since we are not going to receive anything. Bit 6 must be 1 since we are using the EIA-RS-232 output. The binary value of this word then is 01000001 (binary) or 101 (octal).

The following subroutine will therefore configure the Communications Adaptor for the 3300P printer:

PREP1	LA	0322	
	EX	ADR	ADDRESS DEVICE
	LA	0	
	EX	COM1	SET COMMAND WORD
	LA	0377	
	EX	COM3	SET TRANSMIT TIME BASE
	LA	0	
	EX	COM3	
	LA	0101	
	EX	COM4	SET CHARACTER LENGTH
		RET	

This routine need only be executed once at the start of the use of the printer.

Once the Communications Adaptor is configured a subroutine must be called to transmit data to the printer. An in-line subroutine could look like this:

PRINT1	HL	MSG	LOAD H AND L WITH BUFFER ADDRESS OF MESSAGE TO BE TRANSMITTED
LOOP	LA	0322	
	EX	ADR	
	INPUT		ADDRESS DEVICE
	ND	1	MASK FOR TRANSMIT READY
	JTZ	LOOP	LOOP BACK IF NOT READY
	LAM		LOAD A FROM MEMORY IF READY
	EX	WRITE	TRANSMIT TO COMMUNICATIONS ADAPTOR
	CP	015	COMPARE WITH END OF MESSAGE CHARACTER
	RTZ		RETURN IF END OF MESSAGE
	CALL	INCHL*	INCREMENT H AND L
	JMP	LOOP	LOOP BACK IF NOT END OF MESSAGE

The above example assumes that a message has been stored in a buffer area in memory and is transmitted to the printer to the exclusion of all other activity.

A more general routine might be to transmit a single character to the printer and the return to the calling program for other activity while the printer is printing.

An example of this might be as follows:

PRINT	LA	0322	
	EX	ADR	ADDRESS DEVICE
	INPUT		
	ND	1	MASK FOR TRANSMIT READY
	RTZ		RETURN TO CALLING PROGRAM WITH A ZERO IN A REGISTER IF PRINTER NOT READY
	LAB		IF PRINTER READY, LOAD A FROM B WITH CHARACTER TO BE PRINTED
	EX	WRITE	
	OR	1	MAKE SURE Z-FLAG IS SET TO 0
	RET		RETURN TO CALLING PROGRAM

*See end of chapter for frequently used utility routines

Before calling this subroutine, B is loaded with the character to be printed and when the subroutine returns the Z-flag can be tested to see if the printer accepted the character.

4.2 Non-Automatic Data Sets

Data sets that are not automatically controlled from the software such as acoustic couplers or external data sets using private line connections are generally the easiest to program and will be used as our first examples of programming for data sets.

For an example, let us program a Datapoint 2200 to interface with an acoustic coupler which will be used to call a time-sharing service and operate full-duplex at 110 baud. (This program will make the 2200 look like a typical KSR teletype machine). The main program might be written like this:

START1	CALL LA	PREP2 012	CONFIGURE COMM ADAPTOR (LINE FEED)
SCAN1	CALL CALL	DISPLAY READ1	CLEAR BOTTOM LINE OF CRT INPUTS CHAR FROM COMM ADAPTOR IF ONE READY
HDX	JTZ CALL	SCAN2 DISPLAY	GO TO KYBD CHECK WRITE CHAR IN A-REG TO CRT
SCAN2	CALL	KEYIN	INPUTS CHAR FROM KYBD IF ONE READY
	JTZ CALL	SCAN1 WRITE1	CHECK COMM ADAPTOR OUTPUT KYBD CHAR TO COMM ADAPTOR
	JMP	SCAN1	

This is all there needs to be to the main program. When starting, a prep subroutine is called to configure the Communications Adaptor. A scanning loop is then entered which looks for characters from the Communications Adaptor or the keyboard and transmits them to their respective destinations.

If it were desired to operate the program in a half-duplex mode where the characters are displayed directly on the CRT rather than full-duplex where the characters are transmitted back from the remote computer then the last instruction in the main program should be JMP HDX rather than JMP SCAN1.

In this particular mode of operation the Command Word would have bits 0 and 4 set to one and all others set to zero (Paragraph 8.5, 2200 Reference Manual). The time base mask words would be 375 and 106 for both transmit and receive, and the Character Length Mask word would be 111 (octal). (Bit 6 is set to one since the acoustic coupler is an external data set and uses the EIA-RS-232 interface.)

The PREP2 subroutines would therefore be coded as follows:

PREP2	LA	0322	
	EX	ADR	ADDRESS DEVICE
	LA	021	OUTPUT COMMAND WORD
	EX	COM1	'
	LA	0375	SET TRANSMIT AND RECEIVE
	EX	COM2	TIME BASES TO 110 BAUD
	EX	COM3	'
	LA	0106	'
	EX	COM2	'
	EX	COM3	'
	LA	0111	SET CHAR LENGTH MASK
	EX	COM4	TO 11-UNIT CODE
	RET		

To input characters from the Communications Adaptor a subroutine READ1 is written. It will test the Communications Adaptor to see if a character is ready, and if so, read it. If no character is found the Z-flag is returned set to 1 and if a character is read it is returned set to zero. The code is as follows:

READ1	LA	0322	
	EX	ADR	ADDRESS DEVICE
	IN		CHECK READ READY AND
	ND	2	RETURN IF NOT READY
	RTZ		
	EX	DATA	PUT DATA ON INPUT BUS
	IN		TRANSFER CHAR TO A-REG
	ORA		SET Z-FLAG IF CHAR = 0
	RET		

To output characters to the Communications Adaptor a subroutine WRITE 1 is written. It will accept a character in the A-register, transmit it to the Communications Adaptor, and return to the main program when the task is finished with the character remaining in the A-register. It is coded as follows:

WRITE1	LBA	SAVE A IN B	
	LA	0322	
	EX	ADR	ADDRESS DEVICE
RETRY	IN		TEST FOR TRANSMIT READY
	ND	1	AND RETRY IF BUSY
	JTZ	RETRY	
	LAB		
	EX	WRITE	SEND CHAR OUT
	RET		

The subroutines DISPLAY and KEYIN are shown at the end of this chapter for information purposes. Since they do not involve the Communications Adaptor they will not be discussed here.

4.3 The High Level Keyer

When the high level keyer is used it operates in every respect like an external data set except that the Command Word is set to all zeros. Bit 6 of the Character Length Mask is set to one.

5. AUTOMATIC DATA SET OPERATION

One of the major features of the Datapoint 2200 is its ability to operate with the telephone network, providing completely automatic call origination and answering.

5.1 Automatic DDD Network Call Origination.

Automatic Call origination requires the Communications Adaptor to be provided with either a 103 or 202 internal data set option. These data sets interface with the telephone network through a Bell System Direct Access Arrangement (DAA). (See the Datapoint 2200 Installation Manual for specific details).

To automatically originate a call the following events must occur:

- a. The DAA must have been on-hook long enough to assure complete termination of any previous call.
- b. The Communications Adaptor must be configured for an automatic dialing mode.
- c. The DAA must be set "Off-hook" and the dial tone present bit tested for ready (one).
- d. The desired number transmitted.
- e. The Communications Adaptor configured for the type of data set used and the connection confirmed (answered by another data set). (If the call is not confirmed within a reasonable time, usually about 30 seconds, a retry is probably indicated about 3 to 5 times).
- f. Normal data transmission occurs.
- g. The DAA is set to "on-hook" as soon as the connection is no longer desired.

The following code (page 8-9) provides an example of a complete automatic call origination sequence up to the point of reconfiguring the Communications Adaptor for the particular data set used (Step e. above). The number to be dialed is assumed to have been previously stored in an ASCII character sequence in a buffer area in memory beginning at NUMBER. An ASCII '*' (052) between digits results in an extra delay between dial pulses when such might be required to obtain an outside line in a private exchange or for some other reason. The end of the number is indicated by an ASCII return (015). For example:

NUMBER DC '9*5125551234',015

would cause 9 to be dialed, then a pause, then 512-555-1234 to be dialed then control transferred to the calling program.

All other characters in the buffer area are ignored.

DIAL	HL	PHNUMB	BUFFER POINTER
	LA	0322	
	EX	ADR	ADDRESS DEVICE
	SUA		
	EX	COM1	SET DAA ON-HOOK
	DE	10000	
	CALL	DELAY	DELAY 5 SECONDS
	LA	0330	CONFIGURE FOR DIALING; OFF-HOOK,
	EX	COM1	INVERT DATA, SEND DIAL PULSES.
	LA	0375	SET 100 BAUD (10 CPS) DIAL RATE
	EX	COM3	
	LA	0	
	EX	COM3	
	EX	COM4	
DTONE	IN		WAIT FOR DIAL TONE
	ND	0200	'
	JTZ	DTONE	'
	JMP	LDIG	GET FIRST DIGIT
NEXDIG	CALL	INCHL	INCREMENT H AND L
LDIG	LAM		
	CP	**	IF THE A-REG CONTAINS '** THEN
	JFZ	CMPR	CALL 5 SECOND DELAY
	DE	10000	
	CALL	DELAY	'
	JMP	NEXDIG	
CMPR	CP	015	
	RTZ		RETURN IF END OF NUMBER
	CP	'9'+1	TEST FOR VALID DIGIT
	JFS	ERR1	'
	CP	'0'	'
	JTS	ERR1	'
	JFZ	MASK	
	LA	10	CHANGE ZERO TO TEN
MASK	ND	017	MASK-OFF HIGH ORDER BITS
	LBA		SAVE A IN B
PLOOP	IN		WAIT FOR TRANSMIT READY
	ND	1	
	JTZ	PLOOP	'
	LA	0360	
	EX	WRITE	SEND DIAL PULSE
	LAB		DECREMENT PULSE COUNTER
	SU	1	'
	LBA		
	DE	2000	
	JFZ	PLOOP	
	CALL	DELAY*	DELAY ONE SECOND
	JMP	NEXDIG	

Upon returning from the DIAL subroutine the Communications Adaptor should be reconfigured for the type of data set used and the status bit tested for main channel carrier present. If it is not received within 30 seconds the call should be terminated and retried. The following code shows how this could be done for a 103 type data set operating at 150 baud.

*See end of chapter for frequently used utility routines

CNFIRM	LA	0322	
	EX	ADR	
	LA	0121	
	EX	COM1	
	LA	0376	
	EX	COM2	RE-CONFIGURE COMMUNICATIONS ADAPTOR
	EX	COM3	
	LA	0	
	EX	COM2	
	EX	COM3	
	DE	60000	SET TIME COUNTER
MCCDET	INPUT		
	ND	0100	RETURN IF MAIN CHANNEL CARRIER PRESENT
	RFZ		
	LAE		DECREMENT TIME COUNTER
	SU	1	
	LEA		
	LAD		
	SB	0	
	LDA		
	JFZ	MCCDET	
	RET		RETURN IF TIME OUT

This subroutine returns with the Z-flag 0 if the carrier has been detected and 1 if it has not after 30 seconds.

The first part of a main program using these routines might be coded as follows:

START	LC	5	SET NUMBER OF CALL TRIES
RSTART	LAC		
	SU	1	DECREMENT CALL COUNTER
	LCA		
	JTZ	QUIT	EXIT AFTER 5 TRIES
	CALL	DIAL	
	CALL	CONFIRM	
	JTZ	RSTART	
	'		
	'		REMAINDER OF MAIN PROGRAM
	'		
	'		
NUMBER	DC	'9*5125551234',015	

- 5.2 Automatic DDD Network Call Answering.
 Answering a call from the DDD network is very simple and we will not repeat coding examples for this function. The procedure is as follows:

- Make sure the Command Word bit 4 is zero, maintaining the DAA "on-hook".
- At regular intervals test bit 5 of the Status Word for ringing present.

- c. If ringing is detected configure the Communications Adaptor for the type of data set used and set Command Word bit 5 for off-hook (1).
- d. Depending on the type of data set, test for a received carrier (main channel in a type 103, main or supervisory channel depending on initial direction of communications in a 202 type). If no carrier is received after 30 seconds return to step a. above). If normal carrier is received then continue with normal communications.

6. FREQUENTLY USED SUBROUTINES

6.1 INCHL

This subroutine is used to increment the value stored in the H and L register as a double precision (16-bit) number.

```

INCHL    LAL
         AD      1
         LLA
LA       LAH
         AC      0
         LHA
         RET

```

6.2 DELAY

This subroutine provides a means for a time delay up to 30 seconds. Before calling the routine a double precision number is loaded into the D and E registers using the DE macro. This number is decremented at a rate of 2000 counts per second until D and E are zero and then the subroutine returns to the calling program.

```

DELAY    LAE
         SU      1
         LEA
         LAD
         SB      0
         LDA
         JFZ     DELAY
         ADE
         RTZ
         JMP     DELAY

```

6.3 DISPLAY

This routine accepts a character in the A-register and displays it on the CRT screen at the current cursor position and then increments the cursor to the next position. Characters are always entered on the bottom line of the screen and the screen is rolled up one line whenever an ASCII line-feed is received (012). The character displayed is in the A-register when the routine returns.

DISPLAY	LBA		SAVE A IN B
	LA	0341	ADDRESS DEVICE
	EX	ADR	
	LAB		LOAD A FROM B AND
	ND	0177	MASK PARITY BIT
	CP	015	TEST FOR CR
	JTZ	CRDET	
	CP	012	TEST FOR LF
	JTZ	LFDET	
	CP	040	TEST FOR VALID
	RTS		ASCII CHARACTER
	CP	0177	(RUBOUT)
	RTZ		
	EX	WRITE	
	HL	CURPOS	INCREMENT CURSOR POS
	LAM		
	AD	1	
	CP	80	
	JFS	OFDET	
	LMA		SAVE CURSOR POS
	LCA		SAVE A IN C
WCOMP	IN		TEST FOR WRITE DONE
	ND	1	
	JTZ	WCOMP	
	HL	CMDWRD	GET COMMAND WORD
	LAM		
	EX	COM1	
	LA	020	
	LMA		RESTORE COMMAND WORD
	LAC		
	EX	COM2	WRITE NEW CURSOR POS
	LA	11	MAINTAIN CURSOR ON
	EX	COM3	BOTTOM LINE
	LAB		RESTORE CHAR TO A
	RET		
LFDET	HL	CMDWRD	
	LA	030	SET NEW
	LMA		COMMAND WORD
	HL	CURPOS	
	LCM		LOAD CURSOR POS
	JMP	WCOMP	
LFDET	HL	CMDWRD	
	LA	030	SET NEW
	LMA	030	COMMAND WORD
	LC	0	
	HL	CURPOS	SET NEW CURSOR
	LMC		POSITION AND STORE
	JMP	WCOMP	
CURPOS	DC	0	
CMDWRD	DC	020	
CRDET	EQU	LFDET	

6.4 KEYIN

This subroutine is used to scan the keyboard and if a character is present return it to the calling program in the A-register. If the keyboard switch is held down during a keyboard entry, bit six of the data word is set to 0 allowing upper case ASCII characters to be converted to ASCII control characters (e.g., upper case J is converted to ASCII line-feed). The subroutines exits with the Z-flag set to one if no character is input and set to zero if a character is present.

KEYIN	LA	0341	ADDRESS DEVICE
EX		ADR	
IN			INPUT STATUS
LBA			SAVE STATUS IN B
ND	2		
RTZ			RETURN IF READ NOT READY
LAB			RESTORE STATUS
ND	4		MASK FOR KYBD SENSE SW
EX		DATA	
IN			READ DATA FROM KYBD
JFZ		CCONT	JUMP IF KYBD SW SET
ORA			RESET Z-FLAG
RET			
CCONT	ND	077	MASK BIT 6
	RET		

SECTION 9

DATAPOINT 2200

OPERATING SYSTEM LISTING

; PARITY CHECK THE BOOTSTRAPED DATA

000000	066 050 056 000	OKLOAD: HL	PSTART	
000004	036 000	LD	\$-\$	INITIALIZE XOR CHECK
000006	046 000	LE	\$-\$	INITIALIZE CIRCLE CHECK
000010	307	OKLOOP: LAM		GET A BYTE
000011	320	LCA		SAVE IT
000012	253	XRD		ACCUMULATE THE XOR PARITY
000013	330	LDA		
000014	302	LAC		
000015	254	XRE		ACCUMULATE THE CIRCLE PARITY
000016	012	SRC		
000017	340	LEA		
000020	306	LAL		INCREMENT HL
000021	004 001	HALT: AD	1	
000023	360	LLA		
000024	305	LAH		
000025	014 000	AC	0	
000027	350	LHA		
000030	074 002	OP	PEND>8	STOP WHEN PAST END
000032	110 010 000	JFZ	OKLOOP	
000035	306	LAL		
000036	074 000	OP	PEND	
000040	110 010 000	JFZ	OKLOOP	
000043	303	LAD		CHECK THE PARITY ACCUMULATIONS
000044	264	ORE		
000045	110 022 000	JFZ	HALT+1	
000050	066 054 056 000	PSTART: HL	SCLOOP	CLEAR LOW CORE TO HALT SHORT LOADS
000054	306	SCLOOP: LAL		DECREMENT MEMORY POINTER
000055	024 001	SU	1	
000057	360	LLA		
000060	373	LMD		CLEAR THE LOCATION
000061	110 054 000	JFZ	SCLOOP	GO UNTIL LOCATION ZERO CLEAR

; BOOTSTRAP LOADS THE ZEROTH FILE

000064	016 000	BOOT\$# LB	0	LOAD FILE ZERO
000066	106 100 000	CALL	LOAD\$	
000071	100 075 000	JFC	RUN\$	EXECUTE IF LOAD WAS OKAY
000074	377	HALT		
000075	104 064 000	RUN\$# JMP	BOOT\$	OVERSTORED WITH STARTING ADDRESS

- ; 2200 BINARY IMAGE FILE LOADER
- ; UPON ENTRY THE B REGISTER SHOULD CONTAIN
- ; THE DESIRED FILE NUMBER (POSITIVE)
- ; FILE LABEL RECORD FORMAT: 0201/0176/N/-N
- ; DATA RECORD FORMAT: 0303/074/XP/OP/H/L/-H/-L/DATA...
- ; THE 0303/074 INDICATES NUMERIC TYPE DATA
- ; H AND L DEFINE THE STARTING ADDRESS
- ; XP IS THE XOR PARITY AND OP IS THE CIRCULAR PARITY
- ; FOR THE CHARACTERS FOLLOWING THE OP

001000	006 360	LOAD\$# LA	0360	ADDRESS THE CASSETTE MECHANISM
--------	---------	------------	------	--------------------------------

00102	121	EX	ADR	
00103	106 322 001	CALL	STOP	STOP ANY TAPE MOTION
00106	155	EX	DECK1	SELECT THE SYSTEM DECK
00107	104 121 000	JMP	LOAD	
00112	006 360	LOAD\$*	LA 0360	ADDRESS THE CASSETTE MECHANISM
00114	121	EX	ADR	
00115	106 322 001	CALL	STOP	STOP ANY TAPE MOTION
00120	157	EX	DECK2	SELECT THE DATA DECK
00121	106 323 001	LOAD:	CALL DWAIT	WAIT FOR DECK SELECTION
00124	301	LAB		THE REQUESTED FILE NUMBER MUST BE
00125	260	ORA		POSITIVE
00126	106 270 001	JTS	ARGH	
00131	066 077 056 000	HL	RUN\$+2	INITIALIZE THE STARTING LOCATION MSB
00135	250	XRA		FOR 'NOTHING LOADED' FLAG
00136	370	LMA		
00137	104 012 001	JMP	FSTART	
: SEARCH FOR THE DESIRED FILE				
00142	106 360 001	FWAIT:	CALL GETCH	WAIT FOR END OF RECORD
00145	100 142 000	JFC	FWAIT	
00150	106 350 001	FNEXT:	CALL RTINIT	INITIALIZE THE RE-TRY COUNT
00153	026 006	FREAD:	LC 6	WAIT FOR DATA OR LEADER
00155	106 325 001	CALL	TWAIT	
00160	044 002	ND	2	QUIT IF LEADER
00162	110 270 001	JFZ	ARGH	
00165	106 360 001	CALL	GETCH	GET THE RECORD TYPE
00170	330	LDA		SAVE IT
00171	106 360 001	CALL	GETCH	GET THE RECORD TYPE COMPLEMENTED
00174	054 377	XR	0377	UN-COMPLEMENT IT
00176	273	CPD		THE TWO MUST MATCH
00177	110 244 000	JFZ	FSTOP	
00202	074 303	OP	0303	IGNORE NUMERIC RECORDS
00204	150 142 000	JTZ	FWAIT	
00207	074 347	OP	0347	IGNORE SYMBOLIC RECORDS
00211	150 142 000	JTZ	FWAIT	
00214	074 201	OP	0201	ELSE IT MUST BE AN EOF RECORD
00216	110 244 000	JFZ	FSTOP	
00221	106 360 001	CALL	GETCH	GET THE FILE NUMBER
00224	330	LDA		SAVE IT
00225	106 360 001	CALL	GETCH	GET THE FILE NUMBER COMPLEMENTED
00230	054 377	XR	0377	UN-COMPLEMENT IT
00232	273	CPD		MAKE SURE THE TWO MATCH
00233	110 244 000	JFZ	FSTOP	
00236	106 360 001	CALL	GETCH	MAKE SURE THIS IS THE END OF THE RECORD
00241	140 262 000	JTC	WCHWAY	
00244	106 322 001	CALL	STOP	STOP THE TAPE
00247	167	EX	BSP	BACK UP OVER THE RECORD
00250	106 334 001	CALL	DECRTC	DECREMENT THE RE-TRY COUNT
00253	160 270 001	JTS	ARGH	QUIT IF TOO MANY RE-TRIES
00256	171	EX	SF	RE-INITIATE FORWARD MOTION
00257	104 153 000	JMP	FREAD	
00262	303	WCHWAY:	LAD	SEE IF WE ARE THERE YET
00263	271		OPB	

00264	160 150 000	JTS	FNEXT	KEEP GOING IF NOT FAR ENOUGH
00267	150 037 001	JTZ	NXTREC	START LOADING IF THERE
00272	106 322 001	CALL	STOP	ELSE STOP THE TAPE
00275	173	EX	SB	AND START SEARCHING BACKWARD
00276	106 350 001	BWAIT:	RTINIT	INITIATE THE RE-TRY COUNT
00301	026 006	CALL	6	WAIT FOR DATA OR LEADER
00303	106 325 001	CALL	TWAIT	QUIT IF LEADER
00306	044 002	ND	2	
00310	110 270 001	JFZ	ARGH	PUSH THE CHAR ONTO THE STACK
00313	365	BREAD:	LLH	
00314	354		LHE	
00315	343		LED	
00316	330		LDA	
00317	106 360 001	CALL	GETCH	GET THE NEXT RECORD CHARACTER
00322	100 313 000	JFC	BREAD	
00325	304		LAE	GET THE RECORD TYPE COMPLEMENTED
00326	054 377		XR	UN-COMPLEMENT IT
00330	273		CPD	IT MUST MATCH THE TYPE
00331	110 021 001	JFZ	BSTOP	
00334	074 303	OP	0303	IGNORE NUMERIC RECORDS
00336	150 276 000	JTZ	BWAIT	
00341	074 347	OP	0347	IGNORE SYMBOLIC RECORDS
00343	150 276 000	JTZ	BWAIT	
00346	074 201	OP	0201	ELSE IT MUST BE AN EOF RECORD
00350	110 021 001	JFZ	BSTOP	
00353	306		LAL	GET THE FILE NUMBER COMPLEMENTED
00354	054 377		XR	UN-COMPLEMENT IT
00356	225		SUH	MAKE SURE IT MATCHES THE FILE NUMBER
00357	110 021 001	JFZ	BSTOP	
00362	340		LEA	FLIP OVER THE FILE NUMBER
00363	026 010		LC	
00365	305	FLIP:	LAH	
00366	012		SRC	
00367	350		LHA	
00370	304		LAE	
00371	210		ACA	
00372	340		LEA	
00373	302		LAC	
00374	024 001		SU	1
00376	320		LOA	
00377	110 365 000	JFZ	FLIP	COMPARE IT TO THE DESIRED FILE NUMBER
00402	304		LAE	
00403	271		CPB	
00404	160 270 001	JTS	ARGH	IT AINT THERE
00407	110 276 000	JFZ	BWAIT	WE HAVEN'T GONE BACK FAR ENOUGH
00412	106 322 001	FSTART:	CALL	ELSE STOP THE TAPE
00415	171		STOP	AND START GOING FORWARD AGAIN
00416	104 150 000		EX	
00421	106 322 001		JMP	TRY THAT RECORD IN REVERSE AGAIN
00424	161	BSTOP:	CALL	
00425	106 334 001		STOP	
00430	160 270 001		EX	DECREMENT THE RE-TRY COUNT
00433	173		RBK	QUIT IF TOO MANY RE-TRIES
00434	104 313 000		JTS	RE-INITIATE BACKWARD MOTION
			EX	
			JMP	
			BREAD	

; READ IN A DATA RECORD HEADER

00437	106 350 001	NXTREC: CALL RTINIT	INITIALIZE THE RE-TRY COUNT
00442	026 020	NXTWAT: LC 020	WAIT FOR IRG
00444	106 325 001	CALL THAIT	
00447	106 360 001	NEXTRY: CALL GETCH	GET THE RECORD TYPE
00452	140 047 001	JTC NEXTRY	WAIT FOR DATA
00455	330	LDA	SAVE THE RECORD TYPE
00456	106 360 001	CALL GETCH	GET THE RECORD TYPE COMPLEMENTED
00461	054 377	XR 0377	UN-COMPLEMENT IT
00463	273	CPD	THE TWO MUST MATCH
00464	110 304 001	JFZ AGAIN	
00467	074 347	OP 0347	IGNORE SYMBOLIC RECORDS
00471	150 042 001	JTZ NXTWAT	
00474	074 303	OP 0303	LOAD NUMERIC RECORDS
00476	150 130 001	JTZ NXTONE	
00501	074 201	OP 0201	QUIT ON EOF MARKER
00503	110 304 001	JFZ AGAIN	
00506	106 322 001	CALL STOP	STOP THE TAPE
00511	167	EX BSP	BACK UP TO THE END OF THE FILE
00512	106 323 001	CALL DWAIT	
00515	066 077 056 000	HL RUN\$+2	MAKE SURE SOMETHING WAS LOADED
00521	307	LAM	
00522	260	ORA	
00523	150 270 001	JTZ ARGH	
00526	250	XRA	
00527	007	RET	
00530	106 360 001	NXTONE: CALL GETCH	ERROR EXIT IF NOT
00533	350	LHA	ELSE SET THE ZERO CONDITION
00534	106 360 001	CALL GETCH	AND QUIT
00537	360	LLA	GET THE PARITY INITIALIZATION VALUES
00540	106 360 001	CALL GETCH	IN H (XP) AND L (OP)
00543	330	LDA	
00544	106 360 001	CALL GETCH	GET THE STARTING ADDRESS IN DE
00547	340	LEA	
00550	106 360 001	CALL GETCH	
00553	054 377	XR 0377	GET IT AGAIN FOR A CHECK
00555	273	CPD	IT IS COMPLEMENTED THIS TIME
00556	110 304 001	JFZ AGAIN	
00561	106 360 001	CALL GETCH	
00564	140 304 001	JTC AGAIN	CATCH THE RECORD BEING OVER ALREADY
00567	054 377	XR 0377	UN-COMPLEMENT
00571	274	CPE	
00572	110 304 001	JFZ AGAIN	
00575	306	LAL	SAVE THE PARITY ACCUMULATORS
00576	325	LCH	
00577	066 076 056 000	HL RUN\$+1	STORE THE STARTING ADDRESS IN RUN\$ JUMP
00603	374	LME	
00604	066 077	LL RUN\$+2	
00606	373	LMD	SET STORAGE POINTER TO STARTING ADDRESS
00607	353	LHD	
00610	364	LLE	
00611	332	LDC	RESTORE THE PARITY ACCUMULATORS

00612	255	XRH	ACCUMULATE IN THE STARTING ADDRESS
00613	012	SRC	
00614	256	XRL	
00615	012	SRC	
00616	255	XRH	
00617	012	SRC	
00620	256	XRL	
00621	012	SRC	
00622	340	LEA	
00623	305	LAH	

: LOAD A RECORD ACCUMULATING PARITY

00624	106 360 001	NXTBYT: CALL	GETCH	GET A BYTE OF DATA
00627	140 277 001	JTC	EOR	CATCH END OF RECORD
00632	320	LCA		ELSE SAVE IT
00633	253	XRD		ACCUMULATE THE PARITIES
00634	330	LDA		
00635	302	LAC		
00636	254	XRE		
00637	012	SRC		
00640	340	LEA		
00641	306	LAL		PREVENT LOADING INTO THE LOADER
00642	024 000	SU	PEND	
00644	305	LAH		
00645	034 002	SB	PEND>8	
00647	160 270 001	JTS	ARGH	
00652	372	LMC		STORE THE DATA IF ADDRESS OKAY
00653	306	LAL		INCREMENT THE MEMORY ADDRESS
00654	004 001	AD	1	
00656	360	LLA		
00657	305	LAH		
00660	014 000	AC	0	
00662	044 037	ND	037	DO MEMORY WRAP-AROUND
00664	350	LHA		
00665	104 224 001	JMP	NXTBYT	GET THE NEXT DATA BYTE
00670	106 322 001	ARGH:	CALL	STOP
00673	064 001		OR	1
00675	012	SRC		INDICATE ABORTIVE EXIT WITH CARRY TOGGLE
00676	007	RET		
00677	303	EOR:	LAD	CHECK PARITY ACCUMULATIONS
00700	264		ORE	
00701	150 037 001	JTZ	NXTREC	
00704	106 322 001	AGAIN:	CALL	STOP
00707	167		EX	BSP
00710	106 334 001		CALL	DECRTC
00713	160 270 001		JTS	ARGH
00716	171		EX	SF
00717	104 047 001		JMP	NEXTRY
: UTILITY ROUTINES				
00722	177	STOP:	EX	TSTOP
00723	026 001	DWAIT:	LC	1
STOP THE TAPE				
WAIT FOR DECK READY				

00725	123	TWAIT:	EX	STATUS	
00726	101	WAITL:	IN		
00727	242		NDC		
00730	150 326 001	JTZ	WAITL		
00733	007	RET			WAIT FOR SPECIFIED STATUS
00734	106 323 001	DECRTC:	CALL	DWAIT	
00737	066 377 056 001		HL	RTC	WAIT FOR I/O OPERATION
00743	307		LAM		DECREMENT THE RE-TRY COUNT
00744	024 001		SU	1	
00746	370		LMA		
00747	007		RET		
00750	066 377 056 001	RTINIT:	HL	RTC	
00754	006 003		LA	3	INITIATE THE RE-TRY COUNT
00756	370		LMA		TO TRY FOUR TIMES
00757	007		RET		
00760	123	GETCH:	EX	STATUS	
00761	101		IN		GET A CHARACTER
00762	044 024		ND	024	
00764	150 360 001	JTZ	GETCH		WAIT FOR DATA OR IRG
00767	002		SLC		
00770	002		SLC		
00771	002		SLC		
00772	002		SLC		
00773	043		RTC		END OF RECORD
00774	125		EX	DATA	ELSE GET THE CHARACTER
00775	101		IN		
00776	007		RET		
; SYSTEM STORAGE					
00777	000	RTC:	DC	0	RE-TRY COUNT
01000		PEND:	EQU	\$	END OF LOADER LOCATION

DONE

OPERATING SYSTEM COMMAND DECODER

05000		SET	05000	
05000	016 001	LB	1	LOAD THE TAPE DIRECTORY
05002	106 100 000	CALL	LOAD\$	
05005	100 017 012	JFC	GOODL	IT LOADED OKAY
05010	066 001 056 013	HL	BDOMSG	ELSE PRINT CAT UN-LOADABLE MSG
05014	104 046 012	JMP	NOCAT	
05017	006 341	GOODL:	LA 0341	KEYBOARD SWITCH OVERRIDES AUTO-LOAD
05021	121	EX	ADR	
05022	101	IN		
05023	044 004	ND	4	
05025	110 042 012	JFZ	OS\$	
05030	066 171 056 037	HL	ALPFN	RUN ANY AUTO-LOAD PROGRAM
05034	307	LAM		
05035	260	ORA		
05036	310	LBA		
05037	110 201 037	JFZ	MAUTO\$	
05042	066 303 056 012	OS\$#	HL OSMSG	PRINT THE START-UP MESSAGE
05046	106 151 036	NOCAT:	CALL DSPLY\$	
05051	066 367 056 012	NXTCMD:	HL RDYMSG	PRINT 'READY'
05055	106 151 036	CALL	DSPLY\$	
05060	066 151 056 015	HL	CMDBUF	
05064	046 013 036 000	DE	11	INPUT THE COMMAND
05070	026 024	LC	20	POSITION THE CURSOR FOR ENTRY
05072	106 000 036	CALL	KEYIN\$	ONLY ACCEPT 20 CHARACTER
05075	066 362 056 012	HL	CRLF	
05101	106 151 036	CALL	DSPLY\$	
05104	250	XRA		
05105	131	EX	COM1	
05106	066 150 056 015	HL	INPTR	INITIALIZE THE SCANNER POINTER
05112	006 151	LA	CMDBUF	
05114	370	LMA		
05115	106 316 013	CALL	GETSYM	
05120	066 200 056 037	HL	SYMBOL+6	CHECK THE TERMINATING CHARACTER
05124	307	LAM		
05125	074 015	CP	015	
05127	150 144 012	JTZ	FNDCMD	IT MUST BE AN ENTER
05132	074 055	CP	' '	A DASH
05134	150 144 012	JTZ	FNDCMD	
05137	074 040	CP	' '	OR A SPACE
05141	110 264 012	JFZ	BADCMD	
05144	066 175 056 037	FNDCMD:	HL SYMBOL+3	USE ONLY THE FIRST THREE CHARACTERS
05150	016 040	LB	' '	
05152	026 003	LC	3	
05154	106 040 014	CALL	BLKSET	
05157	046 000 036 015	DE	CMDLST	LOOK IT UP IN THE COMMAND LIST
05163	106 264 037	CALL	LOOKUP	
05166	306	LAL		
05167	044 370	ND	0370	POINT THE MEMORY POINTER TO THE
05171	004 006	AD	6	BRANCH ADDRESS
05173	360	LLA		
05174	347	LEM		

05175	106 353 036	CALL INCHL	
05200	337	LDM	
05201	066 214 056 012	HL CBI+1	PUT THE ADDRESS IN THE JUMP INSTRUCTION
05205	374	LME	
05206	066 215 056 012	HL CBI+2	
05212	373	LMD	
05213	106 264 012	OBI: CALL BADCMD	
05216	104 051 012	JMP NXTCMD	

. ERROR MESSAGES

05221	066 073 056 013	NAMREQ: HL NRQMSG	
05225	104 270 012	JMP BADSPL	
05230	066 111 056 013	NONNAME: HL NONMSG	
05234	104 270 012	JMP BADSPL	
05237	066 031 056 013	BADNAM: HL BDNMSG	
05243	104 270 012	JMP BADSPL	
05246	066 042 056 013	CATFUL: HL CFLMSG	
05252	104 270 012	JMP BADSPL	
05255	066 057 056 013	DUPNAM: HL DUPMSG	
05261	104 270 012	JMP BADSPL	
05264	066 217 056 013	BADCMD: HL BOMSG	
05270	036 000	BADSPL: LD Ø	
05272	046 013	LE 11	
05274	106 151 036	CALL DSPLY\$	
05277	151	EX BEEP	
05300	104 051 012	JMP NXTCMD	

05303	011 000 013 000	OSMSG: DC 011.0.013.0.021.011.23.013.11	
05314	103 117 115 120	DC 'COMPUTER TERMINAL OPERATING SYSTEM'.023.023.023.015	
05362	011 000 013 013	CRLF: DC 011.0.013.11.015	
05367	011 000 013 013	RDYMSG: DC 011.0.013.11. 'READY'.015	
05401	011 000 013 013	BDCMSG: DC 011.0.013.11.022. 'CATALOG UNLOADABLE'.015	
05431	102 101 104 040	BDNMSG: DC 'BAD NAME'.015	
05442	114 111 102 122	CFLMSG: DC 'LIBRARY FULL'.015	
05457	116 101 115 105	DUPMSG: DC 'NAME IN USE'.015	
05473	116 101 115 105	NRQMSG: DC 'NAME REQUIRED'.015	
05511	116 117 040 123	NONMSG: DC 'NO SUCH NAME'.015	
05526	101 125 124 117	NOAMSG: DC 'AUTO NOT SET'.015	
05543	101 125 124 117	AUTMSG: DC 'AUTO SET TO '	
05557	040 040 040 040	AUTENT: DC ' '.015	
05566	011 000 013 013	CBTMSG: DC 011.0.013.11.022. 'FRONT TAPE SCRATCH?'.015	
05617	127 110 101 124	BOMSG: DC 'WHAT?'.015	

05625	001	D1PKT: DC 1	DECK ONE IS LOGICAL FILE ONE
05626	002 016	DA TFRBUF	
05630	000	DC Ø	
05631	002	D2PKT: DC 2	DECK TWO IS LOGICAL FILE TWO
05632	002 016	DA TFRBUF	
05634	000	DC Ø	
05635	003	CATPAK: DC 3	CATALOG IS LOGICAL FILE THREE
05636	004 037	DA CATH	
05640	166	DC ALPFN-CATHW+1	
05641	005	OBJPKT: DC 5	OBJECT FILE IS LOGICAL FILE FIVE

05642 002 016
05644 000

DA TFRBUF
DC 0

: CALCULATE A PHYSICAL FILE NUMBER FROM CATALOG ADDRESS

05645 024 010
05647 012
05650 012
05651 012
05652 004 002
05654 007

NOCALC: SU DAT
SRC
SRC
SRC
AD 2
RET,

: SCAN OFF A NAME AND LOOK IT UP

05655 106 316 013
05660 074 015
05662 110 237 012
05665 066 172
05667 307
05670 074 040
05672 150 221 012
05675 046 010 036 037
05701 106 264 037
05704 306
05705 044 007
05707 150 230 012
05712 306
05713 044 370
05715 007

GETNAM# CALL GETSYM GET THE NAME
GETNAN: CP 015
JFZ BADNAM TERMINATING CHARACTER MUST BE AN 015
LL SYMBOL GET THE FIRST CHARACTER
LAM
CP
JTZ NAMREQ THERE MUST BE A NAME
GETNAX: DE DAT LOOK IT UP IN THE CATALOG
CALL LOOKUP
LAL
ND ?
JTZ NONAME IT ISN'T THERE
LAL SET TABLE POINTER TO BEGINNING OF ENTRY
ND 0370
RET

: OPERATING SYSTEM LEXICAL SCANNING SUBROUTINES

05716 016 040
05720 026 007
05722 066 172 056 037
05726 106 040 014
05731 026 172
05733 106 010 014
05736 074 101
05740 160 002 014
05743 074 133
05745 120 002 014
05750 056 037
05752 362
05753 370
05754 302
05755 074 200
05757 014 000
05761 320
05762 106 010 014
05765 074 060
05767 160 002 014
05772 074 072
05774 160 350 013

GETSYM: LB BLANK THE SYMBOL STORAGE
LC ?
HL SYMBOL
CALL BLKSET
LC SYMBOL
CALL GETCH INITIALIZE THE SYMBOL STORAGE POINTER
GETLTR: OP 'A'
JTS GETERM
CP 'Z'+1
JFS GETERM
GETNBR: LH SYMBOL>8 STORE THE CHARACTER
LLC
LMA
LAC
CP SYMBOL+6 BUMP THE STORAGE INDEX
AC 0 UNLESS IT IS AT THE END OF THE STORAGE
LOA
CALL GETCH
CP '0'
JTS GETERM
CP '9'+1
JTS GETNBR
GET THE NEXT CHARACTER
CHECK IT'S RANGE BETWEEN 0 AND 9

05777 104 336 013 JMP GETLTR
 06002 066 200 056 037 GETERM: HL SYMBOL+6 STORE THE TERMINATING CHARACTER
 06006 370 LMA
 06007 007 RET

; GET THE NEXT CHARACTER

06010 066 150 056 015 GETCH: HL INPTR GET THE INPUT POINTER
 06014 307 LAM
 06015 310 LBA
 06016 004 001 AD 1 SAVE IT
 06020 370 LMA
 06021 361 LLB
 06022 307 LAM
 06023 074 015 OP 015 BUMP IT TO THE NEXT CHARACTER
 06025 013 RFZ
 06026 066 150 LL INPTR GET THE CHARACTER POINTED TO
 06030 307 LAM
 06031 024 001 SU 1 ELSE DECREMENT THE CHARACTER POINTER
 06033 370 LMA
 06034 250 XRA
 06035 006 015 LA 015 AND EXIT WITH ZERO CONDITION TRUE
 06037 007 RET AND WITH A 015

; SET A BLOCK OF CORE TO THE B REGISTER CONTENTS
; STARTING ADDRESS IN HL: NUMBER OF POSITIONS IN C

06040 371 BLKSET: LMB
 06041 106 353 036 CALL INCHL
 06044 302 LAC
 06045 024 001 SU 1
 06047 320 LOA
 06050 110 040 014 JFZ BLKSET
 06053 007 RET

; STORAGE

06400	TP			
06400	103 101 124 040	CMDLST: DC	'CAT	COMMAND LIST
06406	000 017	DA	'CATCMD	
06410	116 101 115 040	DC	'NAM	
06416	127 017	DA	'NAMCMD	
06420	122 125 116 040	DC	'RUN	
06426	267 022	DA	'RUNCMD	
06430	111 116 040 040	DC	'IN	
06436	263 017	DA	'INCMD	
06440	117 125 124 040	DC	'OUT	
06446	036 020	DA	'OUTCMD	
06450	104 105 114 040	DC	'DEL	
06456	147 021	DA	'DELCMD	
06460	122 105 120 040	DC	'REP	
06466	341 020	DA	'REPCMD	
06470	101 125 124 040	DC	'AUT	
06476	344 022	DA	'AUTCMD	

06500	115 101 116 040	DC	'MAN	
06506	031 023	DA	MANCMD	
06510	120 122 105 040	DC	'PRE	
06516	073 023	DA	PRECMD	
06520	110 105 130 040	DC	'HEX	
06526	250 023	DA	HEXCMD	
06530	104 105 102 040	DC	'DEB	
06536	200 034	DA	DEBUG\$	
06540	040 040 040 040	DC		
06546	264 012	DA	BADCMD	
06550	000	INPTR:	DC 0	INPUT SCANNER INDEX
06551		CMDBUF:	SKIP 22	
06577	000	CATPTR:	DC 0	LIBRARY CATALOG POINTER
06600	000	CSCPTR:	DC 0	CATALOG SCREEN POINTER
06601	040 040 040 040	CATSPS:	DC 0	CATALOG NAME PRINT STRING
06612	000	ENTSAV:	DC 0	CATALOG ENTRY ADDRESS STORAGE
06613	000	PFNSEL:	DC 0	PHYSICAL FILE NUMBER SELECTED
06614	000	PFNCTR:	DC 0	PHYSICAL FILE NUMBER COUNTER
07000		TP		
07000	000 000	DC	0,0	PARITY STORAGE FOR I/O ROUTINES
07002		TFRBUF:	SKIP 254	I/O TRANSFER BUFFER
 : LIST THE CATALOG				
07400	066 177 056 015	CATCMD:	HL CATPTR	INITIALIZE THE CATALOG POINTER
07404	006 010		LA CAT	
07406	370		LMA	
07407	066 200		LL CSCPTR	INITIALIZE THE SCREEN POSITION
07411	250		XRA	
07412	370		LMA	
07413	066 177 056 015	CATLOP:	HL CATPTR	GET THE ADDRESS OF THE NEXT CAT ENTRY
07417	367		LLM	
07420	056 037		LH CAT>8	
07422	307		LAM	
07423	074 040		CP	GET THE FIRST CHARACTER
07425	150 117 017		JTZ CATEND	LISTING IS FINISHED IF IT IS A SPACE
07430	074 052		CP *	OR AN ASTERISK
07432	150 117 017		JTZ CATEND	TRANSFER NAME INTO PRINT STRING
07435	046 201 056 015		DE CATSPS	
07441	026 006		LC 6	
07443	106 345 037		CALL BLKTFR	
07446	066 200 056 015		HL CSCPTR	GET THE CURSOR POSITION
07452	337		LDM	
07453	303		LAD	SEE IF WE NEED TO GO TO A NEW LINE
07454	074 111		OP 73	
07456	160 070 017		JTS CATMOR	
07461	066 362 056 012		HL CRLF	PUT OUT CR LF IF SO
07465	106 151 036		CALL DSPLY\$	
07470	046 013		LE 11	ALWAYS PRINT ON LINE 11
07472	066 201 056 015		HL CATSPS	PRINT THE NAME
07476	106 151 036		CALL DSPLY\$	
07501	066 200 056 015		HL CSCPTR	UPDATE THE CURSOR POSITION
07505	373		LMD	
07506	066 177		LL CATPTR	UPDATE THE CATALOG ENTRY POSITION

07510	307	LAM		
07511	004 010	AD	8	
07513	370	LMA		
07514	104 013 017	JMP	CATLOP	DO NEXT ENTRY
07517	066 362 056 012	HL	CRLF	MAKE ROOM FOR NEXT COMMAND
07523	106 151 036	CALL	DSPLY\$	
07526	007	RET		

: CHANGE THE FILE NAME

07527	106 316 013	NAMCMD:	CALL	GETSYM	GET THE OLD NAME
07532	074 054	CP	'	'	
07534	110 237 012	JFZ	BADNAM	IT MUST BE TERMINATED BY A COMMA	
07537	066 172	LL	SYMBOL		
07541	307	LAM			
07542	074 040	DP	'		
07544	150 221 012	JTZ	NAMREQ	THERE MUST BE A NAME	
07547	046 010 036 037	DE	CAT	LOOK IT UP	
07553	106 264 037	CALL	LOOKUP		
07556	306	LAL			
07557	044 007	ND	7		
07561	150 230 012	JTZ	NONAME	IT MUST BE IN CATALOG	
07564	335	LDH		SAVE THE CATALOG POINTER	
07565	306	LAL			
07566	044 370	ND	0370		
07570	066 177 056 015	HL	CATPTR		
07574	370	LMA			
07575	066 200	LL	CSCPTR		
07577	373	LMD			
07600	106 316 013	CALL	GETSYM	GET THE NEW NAME	
07603	074 015	CP	015		
07605	110 237 012	JFZ	BADNAM	THE NEW NAME MUST BE TERMINATED BY 015	
07610	066 172 056 037	HL	SYMBOL		
07614	307	LAM			
07615	074 040	CP	'		
07617	150 221 012	JTZ	NAMREQ	THERE MUST BE A NEW NAME	
07622	046 010 036 037	DE	CAT	IT MUST NOT ALREADY BE IN THE CATALOG	
07626	106 264 037	CALL	LOOKUP		
07631	306	LAL			
07632	044 007	ND	7		
07634	110 255 012	JFZ	DUPNAM		
07637	066 177 056 015	HL	CATPTR	RESTORE THE CATALOG POINTER	
07643	347	LEM			
07644	066 200	LL	CSCPTR		
07646	337	LDM			
07647	066 172 056 037	HL	SYMBOL	TRANSFER THE SYMBOL INTO THE CATALOG	
07653	026 006	LC	6		
07655	106 345 037	CALL	BLKTFR		
07660	104 054 023	JMP	UPCAT	UPDATE THE CATALOG FILE	

: BRING A NEW OBJECT FILE INTO THE SYSTEM

07663	106 300 017	INCMD:	CALL	INGET	DO THE PART COMMON WITH HEXCMD
07666	046 241 036 013	DE		OBJPKT	GET TO THE BEGINNING OF THE INPUT FILE

07672	106 022 030	CALL	PBOF\$		
07675	104 076 021	JMP	REPFILE		
07700	106 316 013	INGET:	CALL	GETSYM	GET THE NAME SYMBOL
07703	074 015		CP	'15	
07705	110 237 012		JFZ	BADNAM	TERMINATING CHARACTER MUST BE 015
07710	066 172		LL	SYMBOL	GET THE FIRST CHARACTER
07712	307		LAM		
07713	074 040		OP		
07715	150 221 012		JTZ	NAMREQ	
07720	046 010 036 037		DE	CAT	
07724	106 264 037	INEXT:	CALL	LOOKUP	
07727	074 052		OP	'*	
07731	150 246 012		JTZ	CATFUL	CATALOG FULL IF FIRST CHARACTER IS *
07734	306		LAL		
07735	044 007		ND	7	
07737	110 255 012		JFZ	DUPNAM	ENTRY MUST NOT BE IN THE TABLE
07742	335		LDH		PUT THE NEW NAME IN CATALOG
07743	306		LAL		BUMP MEMORY POINTER TO START OF ENTRY
07744	044 370		ND	0370	
07746	340		LEA		
07747	066 212 056 015		HL	ENTSAV	SAVE THE CATALOG ADDRESS
07753	370		LMA		
07754	066 172 056 037		HL	SYMBOL	
07760	026 006		LC	6	
07762	106 345 037		CALL	BLKTFR	
07765	066 212 056 015		HL	ENTSAV	CALCULATE THE SELECTED FILE NUMBER - 1
07771	307		LAM		
07772	106 245 013		CALL	NOALC	
07775	370		LMA		
07776	024 001		SU	1	
10000	046 225 036 013		DE	D1PKT	
10004	106 033 030		CALL	CPFN\$	
10007	046 225 036 013		DE	D1PKT	
10013	106 022 030		CALL	PBOF\$	
10016	046 225 036 013		DE	D1PKT	
10022	106 017 030		CALL	PEOF\$	
10025	066 212 056 015		HL	ENTSAV	
10031	307		LAM		
10032	106 174 023		CALL	D1FNW	
10035	007			RET	
: OUTPUT AN ELEMENT					
10036	106 316 013	OUTCMD:	CALL	GETSYM	GET THE ELEMENT NAME
10041	074 052		CP	'*	CHECK THE TERMINATING CHAR
10043	150 166 020		JTZ	OUTALL	COPY WHOLE SYSTEM TAPE IF *
10046	074 044		OP	'\$'	
10050	150 166 020		JTZ	OUTALL	
10053	106 260 013		CALL	GETNAM	COPY ALL BUT OS AND CAT IF \$
10056	106 245 013		CALL	NOALC	ELSE DO THE REST OF GETNAM
10061	046 225 036 013		DE	D1PKT	CALCULATE THE PHYSICAL FILE NUMBER
10065	106 033 030		CALL	CPFN\$	POSITION SYSTEM TAPE TO THAT FILE

10070	046 225 036 013	DE D1PKT	
10074	106 022 030	CALL PBOF\$	
10077	106 073 023	CALL PRECMD	PREP THE DATA TAPE
10102	046 241 036 013	DE OBJPKT	POSITION TO THE OUTPUT FILE
10106	106 022 030	CALL PBOF\$	
10111	046 225 036 013	OUTTFR: DE D1PKT	PUT OUT THE FILE
10115	106 000 030	CALL SNFR\$	READ A RECORD FROM THE SYSTEM TAPE
10120	140 145 020	JTC OUTEND	CATCH END OF FILE
10123	506	LAL	CALCULATE THE LENGTH
10124	024 002	SU TFRBUF	
10126	066 244 056 013	HL OBJPKT+3	PUT IT IN THE OUTPUT FILE LENGTH
10132	370	LMA	
10133	046 241 036 013	DE OBJPKT	WRITE OUT THE RECORD
10137	106 006 030	CALL SBFW\$	
10142	104 111 020	JMP OUTTFR	DO THE NEXT RECORD
10145	046 231 036 013	OUTEND: DE D2PKT	PUT FILE MARKER 127 ON OUTPUT FILE
10151	006 177	LA 127	
10153	106 033 030	CALL CPFN\$	
10156	046 231 036 013	DE D2PKT	
10162	106 044 030	CALL TFNW\$	
10165	007	RET	
10166	066 172	OUTALL: LL SYMBOL	THERE MUST NOT HAVE BEEN A NAME
10170	307	LAM	
10171	074 040	CP	
10173	110 237 012	JFZ BADNAM	
10176	066 166 056 013	HL CBTMSG	MAKE SURE THE FRONT TAPE IS SCRATCH
10202	106 151 036	CALL DSPLY\$	
10205	151	EX BEEP	
10206	377	HALT	
10207	006 360	LA 0360	ADDRESS DECK 2
10211	121	EX ADR	
10212	106 146 024	CALL DWAIT	
10215	157	EX DECK2	
10216	106 146 024	CALL DWAIT	
10221	175	EX REWND	REWIND THE TAPE
10222	106 146 024	CALL DWAIT	
10225	066 000 056 026	HL BOOTS	WRITE THE BOOT BLOCK
10231	046 000 036 030	DE BOOTE	
10235	106 213 023	CALL WBLOK	
10240	106 146 024	CALL DWAIT	
10243	066 200 056 037	HL SYMBOL+6	SEE IF THIS IS A FULL COPY
10247	307	LAM	OR JUST FILES 2 TO THE END
10250	024 044	SU '\$'	
10252	150 257 020	JTZ OUTSYS	START COPYING FROM FILE ZERO
10255	006 002	LA 2	START COPYING FROM FILE TWO
10257	046 225 036 013	OUTSYS: DE D1PKT	
10263	106 033 030	CALL CPFN\$	
10266	046 225 036 013	DE D1PKT	
10272	106 022 030	CALL PBOF\$	
10275	066 214 056 015	HL PFNCTR	COPY THE TAPE USING FIRST HALF OF UPDATE
10301	006 377	LA -1	SET UP TO START WRITING FILE MARKERS AT 2
10303	370	LMA	
10304	106 346 021	CALL UPDATA0	
10307	006 177	LA 127	TERMINATE THE DATA TAPE

10311 046 231 036 013	DE D2PKT	WITH FILE MARKER 127
10315 106 033 030	CALL CPFN\$	
10320 046 231 036 013	DE D2PKT	
10324 106 044 030	CALL TFNW\$	
10327 106 146 024	CALL DWAIT	
10332 175	EX REWND	REWIND DECK 2
10333 106 146 024	CALL DWAIT	
10336 104 054 023	JMP UPDAT	

REPLACE THE NAMED FILE

10341 106 255 013	REPCMD: CALL GETNAM	GET THE FILE NAME
10344 066 212 056 015	HL ENTSAV	SAVE THE CATALOG ENTRY ADDRESS
10350 004 010	AD 8	
10352 370	LMA	
10353 024 010	SU 8	
10355 106 245 013	CALL NCALC	CALCULATE THE PHYSICAL FILE NUMBER
10360 066 213 056 015	HL PFNSEL	SAVE IT
10364 370	LMA	
10365 046 241 036 013	DE OBJPKT	POSITION TO THE INPUT FILE
10371 106 022 030	CALL PEOF\$	
10374 066 212 056 015	HL ENTSAV	SEE IF THIS IS THE LAST ENTRY IN THE CATA
10400 367	LLM	
10401 056 037	LH CAT>8	
10403 307	LAM	
10404 074 040	CP	
10406 150 053 021	JTZ REPUP	DO SPECIAL UPDATE IF IT IS
10411 074 052	CP '*'	
10413 150 053 021	JTZ REPUP	POSITION TO THE END OF THE INPUT FILE
10416 046 241 036 013	DE OBJPKT	
10422 106 017 030	CALL PEOF\$	PUT OUT A FILE MARKER AFTER IT
10425 066 213 056 015	HL PFNSEL	
10431 307	LAM	
10432 046 231 036 013	DE D2PKT	
10436 106 033 030	CALL CPFN\$	
10441 046 231 036 013	DE D2PKT	
10445 106 044 030	CALL TFNW\$	
10450 104 066 022	JMP UPDATE	AND THEN DO THE NORMAL UPDATE
10453 066 213 056 015	HL PFNSEL	GET SELECTED FILE NUMBER
10457 307	LAM	
10460 046 225 036 013	DE D1PKT	POSITION SYSTEM TAPE TO THAT FILE
10464 106 033 030	CALL CPFN\$	
10467 046 225 036 013	DE D1PKT	
10473 106 022 030	CALL PEOF\$	
10476 046 241 036 013	DE OBJPKT	READ AN INPUT RECORD
10502 106 000 030	CALL SNFR\$	
10505 140 132 021	JTC REPEND	CATCH END OF FILE
10510 306	LAL	CALCULATE THE LENGTH
10511 024 002	SU TFRBUF	
10513 066 230 056 013	HL D1PKT+3	
10517 370	LMA	
10520 046 225 036 013	DE D1PKT	
10524 106 006 030	CALL SBFW\$	
10527 104 076 021	JMP REPFL	WRITE THE RECORD DO THE NEXT RECORD

10532	006 040	REPEND: LA 32	FOLLOW THE FILE BY FILE MARKERS
10534	106 174 023	CALL D1FNW	32 AND 127
10537	006 177	LA 127	
10541	106 174 023	CALL D1FNW	
10544	104 054 023	JMP UPDAT	UPDATE THE CATALOG FILE
; DELETE A NAMED FILE			
10547	106 255 013	DELOMD: CALL GETNAM	GET THE NAMED FILE
10552	340	LEA	SAVE IT
10553	066 212 056 015	HL ENTSAV	SAVE THE CATALOG ENTRY ADDRESS
10557	370	LMA	
10560	106 245 013	CALL NCALC	CALCULATE THE PHYSICAL FILE NUMBER
10563	066 213 056 015	HL PFNSEL	SAVE IT
10567	370	LMA	
10570	066 171 056 037	HL ALPFN	KILL AUTO PTR IF IT IS POINTING
10574	227	SUM	TO THE FILE TO BE DELETED
10575	110 204 021	JFZ DELDEC	
10600	370	LMA	
10601	104 213 021	JMP DELAUT	
10604	120 213 021	DELDEC: JFS DELAUT	DELETED FILE AFTER AUTO-POINTED FILE
10607	307	LAM	ELSE BUMP DOWN THE AUTO POINTER
10610	024 001	SU 1	TO CORRESPOND TO CATALOG SHIFT
10612	370	LMA	
10613	304	DELAUT: LAE	SEE IF AN ENTRY FOLLOWS
10614	004 010	AD 8	
10616	056 037	LH DAT>8	
10620	360	LLA	
10621	307	LAM	
10622	074 040	CP	
10624	150 275 021	JTZ DELAST	TAKE SPECIAL ACTION IF NOT
10627	074 052	CP '*'	
10631	150 275 021	JTZ DELAST	
10634	026 010	DELMOV: LC 8	SHIFT DOWN THE CATALOG
10636	106 345 037	CALL BLKTFR	
10641	307	LAM	
10642	074 040	CP	
10644	150 254 021	JTZ DELEND	DONE WHEN NO NEXT ENTRY
10647	074 052	OP '*'	OR AT CATALOG STOP ENTRY
10651	110 234 021	JFZ DELMOV	CLEAR THE LAST ENTRY VACATED
10654	364	DELEND: LLE	BY THE MOVE
10655	006 040	LA	
10657	370	LMA	
10660	106 073 023	CALL PRECMD	PREP THE DATA TAPE
10663	046 241 036 013	DE OBJPKT	POSITION FRONT DECK TO OBJECT FILE
10667	106 022 030	CALL PBOF\$	
10672	104 066 022	JMP UPDATE	
10675	066 213 056 015	DELAST: HL PFNSEL	ANN DO THE NORMAL UPDATE
10701	307	LAM	SORTS THE LAST FILE
10702	024 001	SU 1	POSITION THE SYSTEM TAPE TO THE
10704	046 225 036 013	DE D1PKT	SELECTED FILE MINUS ONE
10710	106 033 030	CALL CPFN\$	
10713	046 225 036 013	DE D1PKT	
10717	106 022 030	CALL PBOF\$	

10722	046 225 036 013	DE	D1PKT	POSITION TO THE END OF THE FILE
10726	106 017 030	CALL	PEOF\$	
10731	066 212 056 015	HL	ENTSAV	DELETE THE ENTRY FROM THE CATALOG
10735	367	LLM		
10736	056 037	LH	CAT>8	
10740	006 040	LA		
10742	370	LMA		
10743	104 132 021	JMP	REPEND	TERMINATE TAPE AND UPDATE CATALOG
 : UPDATE THE SYSTEM TAPE				
10746	066 214 056 015	UPDATE:	HL PFNCTR	WRITE THE CURRENT PFN ON DECK TWO
10752	307		LAM	INCREMENT THE CURRENT PFN
10753	004 001		AD 1	
10755	370		LMA	
10756	046 231 036 013	DE	D2PKT	
10762	106 033 030	CALL	CPFN\$	
10765	046 231 036 013	DE	D2PKT	WRITE IT ON DECK 2
10771	106 044 030	CALL	TFNW\$	
10774	046 225 036 013	UPDAT1:	DE D1PKT	READ A RECORD FROM DECK 1
11000	106 000 030		CALL SNFR\$	
11003	140 030 022	JTC	UPDAT2	
11006	306	LAL		CATCH EOF
11007	024 002	SU	TFRBUF	CALCULATE ITS LENGTH
11011	066 234 056 013	HL	DEPKT+3	AND PUT IT IN THE WRITE PACKET
11015	370	LMA		
11016	046 231 036 013	DE	D2PKT	
11022	106 006 030	CALL	SBFW\$	WRITE THE RECORD INCLUDING PARITIES
11025	104 374 021	JMP	UPDAT1	DO THE NEXT RECORD
11030	046 225 036 013	DE	D1PKT	READ FILE NUMBER FROM DECK 1
11034	106 041 030	CALL	TFNR\$	
11037	302	LAC		
11040	074 040	CP	32	
11042	160 346 021	JTS	UPDAT0	MORE TO GO IF LESS THAN 32
11045	006 040	LA	32	ELSE PUT FILE MARKER 32 ON DECK 2
11047	046 231 036 013	DE	D2PKT	
11053	106 033 030	CALL	CPFN\$	
11056	046 231 036 013	DE	D2PKT	
11062	106 044 030	CALL	TFNW\$	
11065	007	RET		
 11066	066 213 056 015	UPDATE:	HL PFNSEL	GET THE SELECTED PHYSICAL FILE NUMBER
11072	307	LAM		
11073	066 214 056 015	HL	PFNCTR	INITIALIZE THE PFN COUNTER
11077	370	LMA		
11100	004 001	AD 1		
11102	046 225 036 013	DE	D1PKT	POSITION TO THE FILE AFTER THE ONE SELECT
11106	106 033 030	CALL	CPFN\$	
11111	046 225 036 013	DE	D1PKT	
11115	106 022 030	CALL	PBOF\$	
11120	106 374 021	CALL	UPDAT1	
11123	046 241 036 013	DE	OBJPKT	
11117	106 022 030	CALL	PBOF\$	COPY SYSTEM TAPE TO DATA TAPE POSITION DATA TAPE TO THE OBJECT FILE

11132	066 213 056 015	HL	PFNSEL	RE-INITIALIZE THE FILE COUNTER
11136	307	LAM		
11137	066 214 056 015	HL	PFNCTR	
11143	370	LMA		
11144	046 225 036 013	DE	D1PKT	POSITION DECK 1 TO SELECTED FILE
11150	106 033 030	CALL	CPFN\$	
11153	046 225 036 013	DE	D1PKT	
11157	106 022 030	CALL	PBOF\$	
11162	104 174 022	JMP	UPDAT4	
11165	046 225 036 013	UPDAT3:	DE	WRITE A FILE NUMBER ON DECK 1
11171	106 044 030	CALL	D1PKT	
11174	046 231 036 013	UPDAT4:	DE	READ A RECORD FROM DECK 2
11200	106 000 030	CALL	D2PKT	
11203	140 230 022	JTC	SNFR\$	
11206	306	LAL	UPDAT6	CATCH EOF
11207	024 002	SU	TFRBUF	CALCULATE IT'S LENGTH
11211	066 230 056 013	HL	D1PKT+3	PUT IT IN THE WRITE PACKET
11215	370	LMA		
11216	046 225 036 013	DE	D1PKT	WRITE THE FILE
11222	106 006 030	CALL	SBFW\$	INCLUDING THE PARITY CHARACTERS
11225	104 174 022	JMP	UPDAT4	DO THE NEXT RECORD
11230	066 214 056 015	UPDAT6:	HL	INCREMENT THE CURRENT PFN COUNTER
11234	307	LAM		
11235	004 001	AD	1	
11237	370	LMA		
11240	046 225 036 013	DE	D1PKT	CHANGE THE PACKET NUMBER
11244	106 033 030	CALL	CPFN\$	
11247	046 231 036 013	DE	D2PKT	READ THE NEXT FILE NUMBER FROM DECK 2
11253	106 041 030	CALL	TFNR\$	
11256	302	LAC		
11257	074 040	CP	32	
11261	160 165 022	JTS	UPDAT3	DO THE NEXT FILE IF IT IS LESS THAN 32
11264	104 132 021	JMP	REPEND	ELSE TERMINATE TAPE AND UPDATE CATALOG
 : LOAD AND EXECUTE A FILE				
11267	106 316 013	RUNCMD:	CALL GETSYM	GET THE FILE NAME
11272	074 052	CP	'*'	LOAD OBJECT FILE IF *
11274	150 311 022	JTZ	RUNOBJ	
11277	106 260 013	CALL	GETNAN	
11302	106 245 013	CALL	NCALC	
11305	310	LBA		
11306	104 201 037	JMP	MAUTO\$	
11311	066 172	RUNOBJ:	LL SYMBOL	MAKE SURE THERE
11313	307	LAM		WAS NO NAME BESIDES *
11314	074 040	CP	' '	
11316	110 237 012	JFZ	BADNAM	
11321	046 241 036 013	DE	OBJPKT	POSITION THE FILE FOR THE LOADER
11325	106 022 030	CALL	PBOF\$	
11330	046 231 036 013	DE	D2PKT	
11334	106 025 030	CALL	BSP\$	
11337	016 001	LB	1	RUN THE OBJECT FILE
11341	104 212 037	JMP	MAUT2\$	ON THE FRONT DECK

SET THE AUTO-LOAD POINTER

11344 066 171 056 037	AUTCMD: HL ALPFN	GET THE POINTER
11350 307	LAM	
11351 260	ORA	
11352 110 377 022	JFZ AUTDUP	ERROR IF ALREADY SET
11355 106 255 013	CALL GETNAM	ELSE GET THE NAME
11360 024 010	SU CAT	CALCULATE THE FILE NUMBER
11362 012	SRC	
11363 012	SRC	
11364 012	SRC	
11365 004 002	AD Z	
11367 066 171 056 037	HL ALPFN	AND SET THE POINTER
11373 370	LMA	
11374 104 054 023	JMP UPDAT	AND UPDATE THE CATALOG FILE
11377 024 002	AUTDUP: SU Z	CALCULATE TABLE ADDRESS
11401 002	SLC	
11402 002	SLC	
11403 002	SLC	
11404 004 010	AD CAT	
11406 360	LLA	
11407 056 037	LH CAT>8	
11411 046 157 036 013	DE AUTENT	
11415 026 006	LC 6	
11417 106 345 037	CALL BLKTFR	PUT TABLE ENTRY IN STRING
11422 066 143 056 013	HL AUTMSG	
11426 104 270 012	JMP BADSPL	AND PRINT IT

RESET THE AUTO-LOAD POINTER

11431 066 171 056 037	MANCMD: HL ALPFN	
11435 307	LAM	
11436 260	ORA	
11437 066 126 056 013	HL NOAMSG	
11443 150 270 012	JT2 BADSPL	AUTO IS NOT SET
11446 066 171 056 037	HL ALPFN	
11452 250	XRA	
11453 370	LMA	

UPDATE THE CATALOG FILE

11454 046 235 036 013	UPCAT: DE CATPAK	
11460 106 022 030	CALL PBOF\$	
11463 046 235 036 013	DE CATPAK	
11467 106 011 030	CALL SNFW\$	
11472 007	RET	

PREPARE A BLANK DATA TAPE

11473 066 166 056 013	PRECMD: HL OBTMSG	WAIT FOR BLANK TAPE
11477 106 151 036	CALL DSPLY\$	
11502 151	EX BEEP	
11503 377	HALT	
11504 046 231 036 013	DE D2PKT	REWIND THE DATA TAPE

11510 106 036 030 CALL TRW\$
 11513 046 231 036 013 DE D2PKT WRITE A FILE NUMBER 0 ON IT
 11517 006 000 LA 0
 11521 106 033 030 CALL CPFN\$
 11524 046 231 036 013 DE D2PKT
 11530 106 044 030 CALL TFNW\$
 11533 046 231 036 013 DE D2PKT WRITE A FILE NUMBER 1 ON IT
 11537 006 001 LA 1
 11541 106 033 030 CALL CPFN\$
 11544 046 231 036 013 DE D2PKT
 11550 106 044 030 CALL TFNW\$
 11553 006 177 LA 127
 11555 046 231 036 013 DE D2PKT WRITE A FILE NUMBER 127 ON IT
 11561 106 033 030 CALL CPFN\$
 11564 046 231 036 013 DE D2PKT
 11570 106 044 030 CALL TFNW\$
 11573 007 RET

: WRITE A FILE MARKER ON DECK 1

11574 046 225 036 013 D1FNW: DE D1PKT
 11600 106 033 030 CALL CPFN\$
 11603 046 225 036 013 DE D1PKT
 11607 106 044 030 CALL TFNW\$
 11612 007 RET

: WRITE A BLOCK TO TAPE

11613 163 WBLOK: EX WBK FIRE UP THE WRITE
 11614 317 WNEXT: LBM GET THE DATA CHARACTER
 11615 123 WWAIT: EX STATUS WAIT FOR WRITE READY
 11616 101 IN
 11617 044 010 ND 010
 11621 150 215 023 JTZ WHAIT
 11624 301 LAB WRITE THE DATA CHARACTER
 11625 127 EX WRITE
 11626 306 LAL BUMP THE MEMORY POINTER
 11627 004 001 AD 1
 11631 360 LLA
 11632 305 LAH
 11633 014 000 AC 0
 11635 350 LHA
 11636 273 OPD SEE IF AT END OF BLOCK YET
 11637 110 214 023 JFZ WNEXT NO CHANCE
 11642 306 LAL
 11643 274 CPE TRY LSB
 11644 110 214 023 JFZ WNEXT ELSE WE ARE DONE
 11647 007 RET

: PUT A TSB TAPE INTO THE LIBRARY

11650 106 300 017 HEXCMD: CALL INGET DO THE PART THAT IS LIKE INCMD
 11653 046 231 036 013 DE D2PKT
 11657 106 036 030 CALL TRW\$

11662	106 157 024	HEXASR:	CALL HEXRBK	SEARCH FOR THE FIRST STARTING ADDRESS
11665	066 007		LL HEXBUF+1	
11667	307		LAM	
11670	074 053		CP '+'	THE FIRST CHARACTER MUST BE A +
11672	110 262 023		JFZ HEXASR	
11675	066 010	HEXGAD:	LL HEXBUF+2	GET THE STARTING ADDRESS
11677	106 256 024		CALL HEXCON	
11702	140 123 024		JTC HEXERR	IT MUST BE FOUR GOOD HEX CHARACTERS
11705	321		LOB	SAVE MSB
11706	106 256 024		CALL HEXCON	
11711	140 123 024		JTC HEXERR	
11714	066 004		LL HEXADR	SAVE THE ADDRESS
11716	372		LMC	
11717	066 005		LL HEXADR+1	
11721	371		LMB	
11722	106 157 024	HEXREC:	CALL HEXRBK	LOAD A RECORD
11725	066 006		LL HEXBUF	GET THE FIRST CHARACTER
11727	307		LAM	
11730	074 012		CP 012	IT MUST BE A LINE FEED
11732	110 123 024		JFZ HEXERR	
11735	066 007		LL HEXBUF+1	GET THE SECOND CHARACTER
11737	307		LAM	
11740	074 052		CP '#'	IGNORE RECORD IF #
11742	150 322 023		JTZ HEXREC	
11745	074 053		CP '+'	GET ADDRESS IF +
11747	150 275 023		JTZ HEXGAD	
11752	074 043		CP '##'	END OF FILE IF #
11754	150 132 021		JTZ REPEND	
11757	066 114		LL HEXWBP	CONVERT THE HEX IN HEXBUF
11761	347		LEM	TO BINARY IN HEXWBF
11762	066 007		LL HEXBUF+1	
11764	106 256 024	HEXCL:	CALL HEXCON	QUIT IF NON-HEX CHARACTER
11767	140 006 024		JTC HEXCL	SWAP E AND L
11772	306		LAL	
11773	364		LLE	
11774	340		LEA	
11775	371		LMB	
11776	306		LAL	STORE BINARY NUMBER
11777	004 001		AD 1	INCREMENT AND SWAP L AND E
12001	364		LLE	
12002	340		LEA	
12003	104 364 023	HEXCL:	JMP HEXCL	DO NEXT HEX PAIR
12005	307		LAM	TERMINATING CHAR MUST BE 023
12007	074 023		CP 023	
12011	150 027 024		JTZ HEXWRT	UNLESS THIS BLOCK IS TO BE CONTINUED
12014	074 053		CP '+'	
12016	110 123 024		JFZ HEXERR	
12021	066 114		LL HEXWBP	
12023	374		LME	
12024	104 322 023	HEXWRT:	JMP LL	
12027	066 114		HEXWBP	
12031	036 123		LD HEXWBF+4	ELSE RESET THE WRITE BUFFER PTR

12033	373	LMD			
12034	066 004	LL	HEXADR	PUT THE STARTING ADDRESS IN BUFFER	
12036	307	LAM			
12037	066 117	LL	HEXWBF		
12041	370	LMA			
12042	054 377	XR	0377		
12044	066 121	LL	HEXWBF+2		
12046	370	LMA			
12047	066 005	LL	HEXADR+1		
12051	307	LAM			
12052	066 120	LL	HEXWBF+1		
12054	370	LMA			
12055	054 377	XR	0377		
12057	066 122	LL	HEXWBF+3		
12061	370	LMA			
12062	304	LAE		CALCULATE THE CORE BLOCK LENGTH	
12063	024 123	SU	HEXWBF+4		
12065	340	LEA			
12066	066 005	LL	HEXADR+1	UPDATE THE CORE ADDRESS	
12070	307	LAM			
12071	204	ADE			
12072	370	LMA			
12073	066 004	LL	HEXADR		
12075	307	LAM			
12076	014 000	AC	0		
12100	370	LMA			
12101	304	LAE		CALCULATE THE WRITE BLOCK LENGTH	
12102	004 004	AD	4	COMPENSATE FOR HL GIVEN TWICE	
12104	066 003 056 025	HL	HEXPKT+3	PUT THE LENGTH IN THE PACKET	
12110	370	LMA			
12111	046 000 036 025	DE	HEXPKT	WRITE THE BUFFER	
12115	106 011 030	CALL	SNFW\$		
12120	104 322 023	JMP	HEXREC	AND DO THE NEXT RECORD	
12123	106 146 024	HEXERR:	CALL	DWAIT	TRY THAT RECORD AGAIN
12126	167		EX	BSP	
12127	106 146 024		CALL	DWAIT	
12132	006 341		LA	0341	UNLESS KEYBOARD SWITCH DEPRESSED
12134	121		EX	ADR	
12135	101		IN		
12136	044 004		ND	4	
12140	150 322 023		JTZ	HEXREC	
12143	104 132 021		JMP	REPEND	
12146	026 001	DWAIT:	LC	1	DECK WAIT LOOP
12150	123	TWAIT:	EX	STATUS	
12151	101		IN		
12152	242		NDC		
12153	150 150 024		JTZ	THAIT	
12156	007		RET		
12157	006 360	HEXRBK:	LA	0360	MAKE SURE THE CASSETTE IS ADDRESSED
12161	121		EX	ADR	
12162	106 146 024		CALL	DWAIT	READ A BLOCK

12165	157		EX	DECK2	FROM DECK 2
12166	106 146 024		CALL	DWAIT	
12171	066 006 056 025		HL	HEXBUF	INTO HEXBUF
12175	161		EX	RBK	
12176	026 024	HEXRNX:	LC	024	WAIT FOR IRG OR DATA
12200	106 150 024		CALL	TWAIT	
12203	044 020		ND	020	
12205	110 146 024		JFZ	DWAIT	QUIT IF IRG
12210	125		EX	DATA	ELSE PUT DATA INTO BUFFER
12211	101		IN		
12212	044 177		ND	0177	STRIP THE PARITY
12214	370		LMA		
12215	306		LAL		BUMP THE MEMORY POINTER
12216	004 001		AD	1	
12220	360		LLA		
12221	104 176 024		JMP	HEXRNX	
12224	024 060	HEXGET:	SU	'0'	CONVERT HEX TO 4-BIT BINARY
12226	160 252 024		JTS	HEXCEN	
12231	074 012		CP	10	
12233	160 250 024		JTS	HEXLOW	
12236	024 007		SU	7	
12240	160 252 024		JTS	HEXCEN	
12243	074 020		CP	16	
12245	120 252 024		JFS	HEXCEN	
12250	260	HEXLOW:	ORA		CLEAR THE CARRY TIGGLE
12251	007		RET		
12252	064 001	HEXCEN:	OR	1	SET THE CARRY TIGGLE
12254	012		SRC		
12255	007		RET		
12256	307	HEXCON:	LAM		GET THE FIRST CHARACTER
12257	106 224 024		CALL	HEXGET	CONVERT IT TO BINARY
12262	043		RTC		QUIT IF NOT HEX
12263	012		SRC		PUT IT IN LEFT HALF OF BYTE
12264	012		SRC		
12265	012		SRC		
12266	012		SRC		
12267	310		LBA		SAVE IT
12270	306		LAL		BUMP THE MEMORY POINTER
12271	004 001		AD	1	
12273	360		LLA		
12274	307		LAM		GET THE SECOND CHARACTER
12275	106 224 024		CALL	HEXGET	CONVERT IT TO BINARY
12300	043		RTC		QUIT IF NOT HEX
12301	261		ORB		MERGE THE TWO HALVES
12302	310		LBA		LEAVE RESULT IN B REGISTER
12303	306		LAL		BUMP THE MEMORY POINTER AGAIN
12304	004 001		AD	1	
12306	360		LLA		
12307	007		RET		
12400			TP		

12400	001	HEXPKT:	DC	1	OUTPUT FILE IS LOGICAL FILE ONE
12401	117 025		DA	HEXWBF	WRITE FROM WRITE BUFFER
12403	000		DC	0	
12404	000 000	HEXADR:	DA	0	CURRENT CORE ADDRESS
12406		HEXBPF:	SKIP	70	
12514	123	HEXWBP:	DC	HEXWBF+4	WRITE BUFFER POINTER
12515	000 000		DC	0.0	ROOM FOR PARITY CHECKS
12517	000 000 000 000	HEXWBF:	DC	0.0.0.0	ROOM FOR H AND L
12523			SKIP	128	ROOM FOR THE DATA
12723		HEXWBE:	EQU	\$	
13000			SET	013000	ROOM FOR THE BOOT BLOCK
13000		BOOTS:	SKIP	01000	
14000		BOOTE:	EQU	\$	

DONE

14000

SET 0140000

: OPERATING SYSTEM ROUTINE ENTRY POINT TABLE

14000	104 052 030	SNFR\$*	JMP	SNFRX	
14003	104 230 030	SSFR\$*	JMP	SSFRX	
14006	104 376 030	SBFW\$*	JMP	SBFWX	
14011	104 005 031	SNFW\$*	JMP	SNFWX	
14014	104 072 031	SSFW\$*	JMP	SSFWX	
14017	104 368 031	PEOF\$*	JMP	PEOFX	
14022	104 375 031	PBOF\$*	JMP	PBOFX	
14025	104 004 032	BSP\$*	JMP	BSPX	
14030	104 016 032	CPDN\$*	JMP	CPDNX	
14033	104 030 032	CPFN\$*	JMP	CPFNX	
14036	104 075 034	TRW\$*	JMP	TRWX	
14041	104 114 034	TFNR\$*	JMP	TFNRX	
14044	104 127 034	TFNW\$*	JMP	TFNWX	
14047	104 104 032	ERR\$*	JMP	ERRX	
					: SERIAL NUMERIC FILE READ
14052	106 355 031	SNFRX:	CALL	RTCI	INITIALIZE THE RE-TRY COUNT
14055	106 152 032	SNFRS:	CALL	GETPKT	GET THE PACKET PARAMETERS
14060	106 027 033		CALL	RBK\$	START READING THE RECORD
14063	106 363 032		CALL	READ\$	GET THE RECORD TYPE
14066	330		LDA		SAVE IT
44067	106 363 032		CALL	READ\$	GET THE RECORD TYPE COMPLEMENTED
14072	054 377		XR	0377	UN-COMPLEMENT IT
14074	273		OPD		MAKE SURE THEY MATCH
14075	110 217 030		JFZ	SNFRR	TRY AGAIN IF THEY DON'T
14100	074 201		OP	0201	SEE IF IT IS A FILE MARKER
14102	150 062 032		JTZ	FEACT	QUIT IF IT IS
14105	074 347		OP	0347	SEE IF IT IS A SYMBOLIC RECORD
14107	150 074 032		JTZ	TEACT	TYPE ERROR IF IT IS
14112	074 303		OP	0303	MAKE SURE IT IS A NUMERIC RECORD
14114	110 217 030		JFZ	SNFRR	
14117	106 363 032		CALL	READ\$	GET THE PARITY CHECKS
14122	330		LDA		
14123	370		LMA		STORE PARITY IN FIRST BYTE OF BUFFER
14124	306		LAL		
14125	004 001		AD	1	
14127	360		LLA		
14130	305		LAH		
14131	014 000		AC	0	
14133	350		LHA		
14134	106 363 032		CALL	READ\$	
14137	140 217 030		JTC	SNFRR	TRY AGAIN IF RECORD OVER ALREADY
14142	340		LEA		
14143	370		LMA		STORE PARITY IN SECOND BYTE OF BUFFER
14144	306		LAL		
14145	004 001		AD	1	
14147	360		LLA		

14150	305		LAH		
14151	014 000		AC	0	
14153	350		LHA		
14154	106 363 032	SNFRL:	CALL READ\$		READ THE REST OF THE RECORD
14157	140 205 030		JTC SNFRE		QUIT IT AT END OF RECORD
14162	370		LMA		STORE THE BYTE OF DATA
14163	320		LCA		SAVE IT
14164	253		XRD		ACCUMULATE THE PARITIES
14165	330		LDA		
14166	302		LAC		
14167	254		XRE		
14170	012		SRC		
14171	340		LEA		
14172	306		LAL		BUMP THE MEMORY POINTER
14173	004 001		AD 1		
14175	360		LLA		
14176	305		LAH		
14177	014 000		AC 0		
14201	350		LHA		
14202	104 154 030		JMP SNFRL		DO THE NEXT BYTE
14205	303	SNFRE:	LAD		CHECK THE PARITY TOTALS
14206	264		ORE		
14207	110 217 030		JFZ SNFRR		TRY AGAIN IF THEY AREN'T BOTH ZERO
14212	106 016 033		CALL WAIT\$		ELSE WAIT FOR THE OPERATION TO BE COMPLETED
14215	250		XRA		CLEAR THE CARRY TGGLE
14216	007		RET		AND RETURN
14217	106 324 031	SNFRR:	CALL DECRTC		BACK UP AND TRY AGAIN
14222	120 055 030		JFS SNFRS		UNLESS RTC IS NEGATIVE
14225	104 102 032		JMP PEACT		IN WHICH CASE, PARITY ERROR EXIT
 : SERIAL SYMBOLIC FILE READ					
14230	106 355 031	SSFRX:	CALL RTCI		INITIALIZE THE RE-TRY COUNT
14233	106 152 032	SSFRS:	CALL GETPKT		GET PACKET PARAMETERS
14236	106 027 033		CALL RBK\$		START THE READ
14241	106 363 032		CALL READ\$		GET THE RECORD TYPE
14244	330		LDA		SAVE IT
14245	106 363 032		CALL READ\$		GET THE RECORD TYPE COMPLIMENTED
14250	054 377		XR 0377		UN-COMPLEMENT IT
14252	273		CPD		THEY MUST MATCH
14253	110 365 030		JFZ SSFRR		
14256	074 201		CP 0201		QUIT IF IT IS AN EOF RECORD
14260	150 062 032		JTZ FEACT		
14263	074 303		CP 0303		TYPE ERROR IF IT IS A NUMERIC RECORD
14265	150 074 032		JTZ TEACT		
14270	074 347		CP 0347		MAKE SURE IT IS A SYMBOLIC RECORD
14272	110 365 030		JFZ SSFRR		
14275	106 363 032		CALL READ\$		INITIALIZE THE PARITY ACCUMULATORS
14300	330		LDA		
14301	106 363 032		CALL READ\$		
14304	340		LEA		
14305	140 365 030	SSFRL:	JTC SSFRR		TRY AGAIN IF THE RECORD IS OVER ALREADY
14310	106 363 032		CALL READ\$		READ THE REST OF THE RECORD
14313	140 350 030		JTC SSFRE		QUIT IF THE RECORD IS ENDED

14316	260		ORA	CHECK THE VERTICAL PARITY
14317	130 365 030		JFP SSFRR	TRY AGAIN IF IT IS FALSE
14322	320		LCA	SAVE THE BYTE
14323	044 177		ND 0177	STRIP THE VERTICAL PARITY
14325	370		LMA	STORE THE BYTE
14326	302		LAC	ACCUMULATE THE PARITIES
14327	253		XRD	
14330	330		LDA	
14331	302		LAC	
14332	254		XRE	
14333	012		SRC	
14334	340		LEA	
14335	306		LAL	BUMP THE MEMORY POINTER
14336	004 001		AD 1	
14340	360		LLA	
14341	305		LAH	
14342	014 000		AC 0	
14344	350		LHA	
14345	104 310 030		JMP SSFRL	DO THE NEXT CHARACTER
14350	006 015	SSFRE:	LA 015	TERMINATE STRING WITH AN 015
14352	370		LMA	
14353	303		LAD	CHECK THE PARITY SUMS
14354	264		ORE	
14355	110 365 030		JFZ SSFRR	TRY AGAIN IF BOTH AREN'T ZERO
14360	106 016 033		CALL WAIT\$	ELSE WAIT FOR THE OPERATION TO COMPLETE
14363	250		XRA	CLEAR THE CARRY TOGGLE
14364	007		RET	AND RETURN
14365	106 324 031	SSFRR:	CALL DECRTC	BACK UP AND TRY AGAIN
14370	120 233 030		JFS SSFRS	UNLESS RTC IS NEGATIVE
14373	104 102 032		JMP PEACT	IN WHICH CASE, PARITY ERROR EXIT
 : SERIAL BLOCK FILE WRITE				
14376	106 152 032	SBFWX:	CALL GETPKT	
14401	342		LEC	
14402	104 026 031		JMP SBFWE	PUT THE LENGTH IN THE E REGISTER
 : SERIAL NUMERIC FILE WRITE				
14405	106 152 032	SNFWX:	CALL GETPKT	GET THE PACKET PARAMETERS
14410	106 277 032		CALL SAVHL	SAVE THE BUFFER STARTING ADDRESS
14413	036 000		LD 0	INITIALIZE THE PARITY ACCUMULATORS
14415	046 000		LE 0	
14417	307	SNFWPG:	LAM	GENERATE THE PARITY TOTALS
14420	106 213 031		CALL PARGEN	
14423	110 017 031		JFZ SNFWPG	
14426	106 034 033	SBFWE:	CALL WBK\$	START UP THE WRITE
14431	036 303		LD 0303	WRITE OUT RECORD TYPE NUMERIC
14433	106 002 033		CALL WRITE\$	
14436	036 074		LD 074	WRITE OUT ITS COMPLEMENT
14440	106 002 033		CALL WRITE\$	
14443	337	SNFWL:	LDM	WRITE OUT THE REST OF THE RECORD
14444	106 002 033		CALL WRITE\$	
14447	306		LAL	BUMP THE MEMORY POINTER

14450	024	001	AD	1	
14452	360		LLA		
14453	305		LAH		
14454	014	000	AC	0	
14456	350		LHA		
14457	304		LAE		DECREMENT THE BUFFER LENGTH COUNT
14460	024	001	SU	1	
14462	340		LEA		
14463	110	043 031	JFZ	SSFWL	
14466	106	016 033	CALL	WAIT\$	WAIT FOR THE OPERATION TO BE COMPLETE
14471	007		RET		

: SERIAL SYMBOLIC FILE WRITE

14472	106	152 032	SSFWX:	CALL	GETPKT	GET THE PACKET PARAMETERS
14475	106	277 032		CALL	SAVHL	SAVE THE START OF BUFFER ADDRESS
14500	036	000		LD	0	INITIALIZE THE PARITY ACCUMULATORS
14502	046	000		LE	0	
14504	307		SSFWPG:	LAM		GENERATE THE PARITY TOTALS
14505	074	015		CP	015	CHECK FOR END OF BUFFER
14507	150	131 031		JTZ	SSFWPS	
14512	260			ORA		
14513	170	120 031		JTP	SSFWPT	GENERATE THE VERTICAL PARITY BIT
14516	054	200		XR	0200	
14520	370		SSFWPT:	LMA		
14521	026	002		LC	2	WRITE OUT CORRECTLY PARITIED CHAR
14523	106	213 031		CALL	PARGEN	FAKE OUT PARGEN LENGTH COUNTER
14526	104	104 031		JMP	SSFWPG	
14531	106	232 031	SSFWPS:	CALL	PARSTO	
14534	106	034 033		CALL	WBK\$	START UP THE WRITE
14537	036	347		LD	0347	PUT OUT RECORD TYPE SYMBOLIC
14541	106	002 033		CALL	WRITE\$	
14544	036	030		LD	0330	PUT OUT THE TYPE COMPLEMENTED
14546	106	002 033		CALL	WRITE\$	
14551	046	002		LE	2	DONT CHECK FOR 015 IN 1ST TWO PARITY BYTE
14553	304		SSFWL:	LAE		
14554	024	001		SU	1	DECREMENT FUDGE COUNTER
14556	340			LEA		
14557	307			LAM		
14560	120	174 031		JFS	SSFHW	GET CHARACTER FROM BUFFER
14563	074	015		CP	015	EREG NOT NEG SO DONT CHECK FOR 015
14565	110	174 031		JFZ	SSFHW	CHECK FOR END OF STRING
14570	106	016 033		CALL	WAIT\$	NOT END OF STRING SO WRITE IT OUT
14573	007			RET		ITS A 015 SO END OF STRING
14574	330		SSFWH:	LDA		SO RETURN.
14575	106	002 033		CALL	WRITE\$	WRITE THE BUFFERED CHARACTER
14600	306			LAL		
14601	004	001		AD	1	BUMP THE MEMORY POINTER
14603	360			LLA		
14604	305			LAH		
14605	014	000		AC	0	
14607	350			JMP	SSFWL	DO THE NEXT CHARACTER
14610	104	153 031				

14613	310	PARGEN:	LBA	SAVE THE BYTE
14614	253		XRD	
14615	330		LDA	
14616	301		LAB	
14617	254		XRE	
14620	012		SRC	
14621	340		LEA	
14622	106 353 036	CALL	INCHL	
14625	302		LAC	
14626	024 001		SU 1	DECREMENT THE BUFFER LENGTH COUNT
14630	320		LCA	
14631	013		RFZ	
14632	306	PARSTO:	LAL	
14633	066 333 056 032		HL HLSAV+1	DO NEXT BYTE IF NOT ZERO CALCULATE NUMBER OF SHIFT MOD 8
14637	227		SUM	
14640	044 007		ND 7	
14642	320		LCA	
14643	302	PSLOOP:	LAC	SHIFT CIRCULATING PARITY BACK THAT MANY
14644	024 001		SU 1	
14646	320		LCA	
14647	106 260 031		JTS PSTORE	
14652	304		LAE	
14653	002		SLC	
14654	340		LEA	
14655	104 243 031		JMP PSLOOP	
14660	106 316 032	PSTORE:	CALL RESHL	STORE THE CIRC. PARITY
14663	106 364 036		CALL DECHL	
14666	374		LME	
14667	106 364 036		CALL DECHL	STORE THE XOR PARITY
14672	373		LMD	
14673	066 254 056 032		HL PKTADR	GET THE PACKET PARAMETERS AGAIN
14677	347		LEM	
14700	066 255 056 032		HL PKTADR+1	
14704	337		LDM	
14705	106 152 032		CALL GETPKT	
14710	302		LAC	
14711	004 002		AD 2	INIT THE BUFFER LENGTH COMPENSATE FOR THE TWO PARITY ACCUMS
14713	340		LEA	PUT LENGTH IN E-REGISTER
14714	106 364 036		CALL DECHL	BACK UP BUFFER POINTER TO PARITY ACCUMS
14717	106 364 036		CALL DECHL	
14722	250		XRA	RETURN WITH ZERO CONDITION TRUE
14723	007		RET	
: BACK UP AND DECREMENT THE RE-TRY COUNT				
14724	106 041 033	DECRTC:	CALL BKSP\$	BACK UP ONE RECORD
14727	106 016 033	DECRTC:	CALL WAIT\$	WAIT FOR IT
14732	066 365 056 031		HL RTC	DECREMENT THE RE-TRY COUNT
14736	307		LAM	
14737	024 001		SU 1	
14741	370		LMA	
14742	066 254 056 032		HL PKTADR	RESTORE THE PACKET ADDRESS
14746	347		LEM	
14747	066 255 056 032		HL PKTADR+1	

14753 337 LDM
 14754 007 RET

: INITIATE THE RE-TRY COUNT

14755 066 365 056 031 RTCI: HL RTC SET THE RE-TRY COUNT TO THREE
 14761 006 003 LA 3
 14763 370 LMA
 14764 007 RET

14765 000 RTC: DC 0 RE-TRY COUNT STORAGE

: POSITION TO THE END OF THE FILE

14766 106 152 032 PEOFX: CALL GETPKT
 14771 106 310 033 CALL PEF\$
 14774 007 RET

: POSITION TO THE BEGINNING OF THE FILE (AFTER FILE NUMBER RE

14775 106 152 032 PBOFX: CALL GETPKT
 15000 106 056 033 CALL PBF\$
 15003 007 RET

: BACKSPACE ONE RECORD

15004 106 152 032 BSPX: CALL GETPKT
 15007 106 041 033 CALL BKSP\$
 15012 106 016 033 CALL WAIT\$
 15015 007 RET

: CHANGE PHYSICAL DEVICE NUMBER

15016 320 CPDNX: LCA SAVE THE PFN
 15017 353 LHD
 15020 364 LLE
 15021 046 257 036 032 DE LFT INIT THE LFT INDEX
 15025 104 037 032 JMP CLFT THE REST IF LIKE CPFN\$

: CHANGE PHYSICAL FILE NUMBER

15030 320 CPFNX: LCA SAVE THE PFN
 15031 353 LHD
 15032 364 LLE
 15033 046 260 036 032 DE LFT+1 INIT THE LFT INDEX
 15037 307 CLFT: LAM GET THE LOGICAL FILE NUMBER
 15040 260 ORA
 15041 160 247 032 JTS GDFNER CHECK IT'S RANGE
 15044 074 010 CP 8
 15046 120 247 032 JFS GDFNER
 15051 002 SLC
 15052 204 ADE
 15053 360 LLA
 15054 303 LAD

15055	014 000	AC	Ø
15057	350	LHA	
15060	372	LMC	CHANGE THE PFN
15061	007	RET	

: END ACTION RETURN POINTS

15062	106 041 033	FEACT:	CALL BKSP\$	BACK UP TO THE END OF FILE
15065	106 016 033		CALL WAIT\$	WAIT FOR IT
15070	064 001		OR 1	SET THE CARRY TOGGLE
15072	012		SRC	
15073	007		RET	
15074	106 016 033	TEACT:	CALL WAIT\$	WAIT FOR RECORD TO FINISH
15077	064 001		OR 1	TYPE ERROR RETURNS NON-ZERO
15101	007		RET	
15102	006 004	PEACT:	LA 4	INTERNAL ERROR D IF PARITY ERROR

: INTERNAL ERROR HANDLER

15104	066 150 056 032	ERRX:	HL ERRS	
15110	004 100		AD 'A'-1	
15112	370		LMA	
15113	066 125 056 032		HL ERRMSG	
15117	106 151 036		CALL DSPLY\$	
15122	104 064 000		JMP BOOT\$	
15125	011 000 013 013	ERRMSG:	DC \$11.0.013.11. 'INTERNAL ERROR'	
15150	040 015	ERRS:	DC ' ',015	

: GET THE DEVICE NUMBER IN THE B REGISTER
 : THE PHYSICAL FILE NUMBER IF 'PFN'
 : THE LENGTH IN THE C REGISTER
 : THE BUFFER STARTING ADDRESS IN HL

15152	066 254 056 032	GETPKT:	HL PKTADR	SAVE THE PACKET ADDRESS
15156	374		LME	
15157	066 255 056 032		HL PKTADR+1	
15163	373		LMD	
15164	353		LHD	GET THE LOGICAL FILE NUMBER
15165	364		LLE	
15166	307		LAM	
15167	260		ORA	CATCH LOGICAL FILE NUMBER OUT OF RANGE
15170	160 247 032		JTS GDFNER	
15173	074 010		CP 8	
15175	120 247 032		JFS GDFNER	
15200	002		SLC	INDEX INTO THE LOGICAL FILE TABLE
15201	004 257		AD LFT	
15203	360		LLA	
15204	006 032		LA LFT>8	
15206	014 000		AC Ø	
15210	350		LHA	
15211	317		LBM	GET THE DEVICE NUMBER IN THE B REGISTER
15212	106 353 036		CALL INOHL	
15215	327		LCM	GET THE PHYSICAL FILE NR IN THE C REG
15216	066 256 056 032		HL PFN	SAVE IT IN CORE

15222	372	LMC	
15223	353	LHD	GET THE BUFFER STARTING ADDRESS
15224	364	LLE	
15225	106 353 036	CALL INCHL	
15230	347	LEM	
15231	106 353 036	CALL INCHL	
15234	337	LDM	
15235	106 353 036	CALL INCHL	
15240	327	LCM	GET THE LENGTH
15241	353	LHD	PUT THE BSA IN HL
15242	364	LLE	
15243	106 334 032	CALL ADR\$	SELECT THE PROPER PHYSICAL DEVICE
15246	007	RET	
15247	006 010	GDFNER: LA 8	LOGICAL FILE NUMBER OUT OF RANGE NETS
15251	104 047 030	JMP ERR\$	YOU AN INTERNAL ERROR NUMBER EIGHT
15254	000 000	PKTADR: DA Ø	CURRENT PACKET ADDRESS STORAGE
15256	000	PFN: DC Ø	CURRENT PHYSICAL FILE NUMBER STORAGE

: OPERATING SYSTEM LOGICAL FILE TABLE

15257	000 000	LFT: DC Ø.Ø	LFØ IS A NULL DEVICE
15261	001 000	DC 1.Ø	LF1 FOR DECK 1
15263	002 000	DC 2.Ø	LF2 FOR DECK 2
15265	001 001	DC 1.1	LF3 IS CTOS CATALOG
15267	002 000	DC 2.Ø	LF4 IS CTOS DATA SOURCE FILE
15271	002 001	DC 2.1	LF5 IS CTOS DATA OBJECT FILE
15273	000 000	DC Ø.Ø	
15275	001 040	DC 1.32	LF7 IS ASM OBJECT SCRATCH FILE

: UTILITY ROUTINES

15277	305	SAVHL: LAH	
15300	316	LBL	
15301	066 332 056 032	HL HLSAV	
15305	370	LMA	
15306	066 333 056 032	HL HLSAV+1	
15312	371	LMB	
15313	361	LLB	
15314	350	LHA	
15315	007	RET	
15316	066 332 056 032	RESHL: HL HLSAV	
15322	307	LAM	
15323	066 333 056 032	HL HLSAV+1	
15327	367	LLM	
15330	350	LHA	
15331	007	RET	

15332	000 000	HLSAV: DA Ø	
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: CASSETTE MECHANISM DRIVER

15334	006 360	ADR\$: LA Ø360	ADDRESS THE CASSETTE MECHANISM
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15336	121	EX	ADR	
15337	301	LAB		SELECT THE PROPER DECK
15340	074 001	OP	1	
15342	150 357 032	JTZ	DEK1AD	
15345	074 002	OP	2	
15347	150 361 032	JTZ	DEK2AD	
15352	006 001	LA	1	BAD DEVICE NUMBER IS ERROR 'A'
15354	104 047 030	JMP	ERR\$	
15357	155	DEK1AD:	EX DECK1	SELECT DECK ONE
15360	007		RET	
15361	157	DEK2AD:	EX DECK2	SELECT DECK TWO
15362	007		RET	
 : READ A CHARACTER INTO THE A REGISTER				
15363	123	READ\$:		
15363	123	DEKRED:	EX STATUS	WAIT FOR IRG OR READ READY
15364	101		IN	
15365	044 024		ND	024
15367	150 363 032		JTZ	DEKRED
15372	002		SLC	
15373	002		SLC	
15374	002		SLC	
15375	002		SLC	
15376	043		RTC	
15377	125		EX DATA	QUIT IF INTER-RECORD GAP ELSE GET THE CHARACTER
15400	101		IN	
15401	007		RET	
 : WRITE A CHARACTER FROM THE D REGISTER				
15402	123	WRITE\$:		
15402	123	DEKWRT:	EX STATUS	
15403	101		IN	
15404	044 011		ND	011
15406	150 002 033		JTZ	DEKWRT
15411	012		SRC	
15412	043		RTC	
15413	303		LAD	ERROR IF DECK READY WRITE THE DATA
15414	127		EX WRITE	
15415	007		RET	
15416		WAIT\$:		
15416	026 001	DEKHAT:	LC 1	WAIT FOR DECK READY
15420	123	WAIT:	EX STATUS	
15421	101		IN	
15422	242		NDC	
15423	150 026 033		JTZ	WAIT
15426	007		RET	
 : FIRE UP BLOCK READ				
15427		RBK\$:		
15427	106 016 033	DEKRBK:	CALL DEKHAT	WAIT FOR THE DECK TO BE READY

15432	161	EX RBK	FIRE UP THE READ BLOCK
15433	007	RET	
; FIRE UP BLOCK WRITE			
15434		WBK\$:	
15434	106 016 033	DEKWBK: CALL DEKWAT	WAIT FOR THE DECK TO BE READY
15437	163	EX WBK	
15440	007	RET	FIRE UP BLOCK WRITE
; BACKSPACE ONE RECORD			
15441		BKSP\$:	
15441	106 016 033	DEKBSP: CALL DEKWAT	
15444	167	EX BSP	
15445	007	RET	
; REWIND THE TAPE			
15446		REWIND\$:	
15446	106 016 033	DEKREW: CALL DEKWAT	
15451	175	EX REWND	
15452	106 016 033	CALL DEKWAT	
15455	007	RET	
; POSITION TO THE BEGINNING OF THE FILE			
15456		PBF\$:	
15456	106 016 033	DEKPBF: CALL DEKWAT	WAIT FOR ANY PREVIOUS OPERATIONS
15461	066 256 056 032	HL PFN	GET THE DESIRED FILE NUMBER
15465	317	LBM	
15466	173	EX SB	
15467	104 131 033	JMP BWAIT	START SEARCHING BACKWARDS
15472	106 333 033	BBACK: CALL DEKSTP	
15475	106 355 033	CALL DEKFNS	STOP THE TAPE
15500	104 111 033	JMP FSKIP	SEARCH FOR A FILE MARKER
15503	106 355 031	FNEXT: CALL RTCI	
15506	106 375 033	CALL DEKFNN	INITIALIZE THE RE-TRY COUNT
15511	303	FSKIP: LAD	SEARCH FOR NEXT FILE MARKER
15512	271	OPB	SEE IF WE ARE THERE YET
15513	160 103 033	JTS FNEXT	
15516	150 274 033	JT2 DEKTHE	
15521	106 333 033	CALL DEKSTP	STILL FURTHER TO GO
15524	006 007	LA 7	WE ARE THERE
15526	104 047 030	JMP ERR\$	ELSE STOP THE TAPE
15531	106 355 031	BWAIT: CALL RTCI	ERROR EXIT SEVEN
15534	026 006	LC 6	
15536	106 020 033	CALL WAIT	INITIALIZE THE RE-TRY COUNT
15541	044 002	ND 2	WAIT FOR READ READY OR LEADER
15543	110 256 033	JFZ BSTOP	
15546	365	LLH	CATCH LEADER
15547	354	LHE	PUSH THE CHARACTER ONTO THE STACK
15550	343	LED	

15551	330	LDA		
15552	106 363 032	CALL	DEKRED	GET THE NEXT RECORD CHARACTER
15555	100 146 033	JFC	BREAD	
15560	304	LAE		GET THE SECOND RECORD CHARACTER
15561	054 377	XR	0377	UN-COMPLEMENT IT
15563	273	CPD		SEE IF IT MATCHES THE FIRST
15564	110 256 033	JFZ	BSTOP	
15567	074 303	CP	0303	IGNORE NUMERIC RECORDS
15571	150 131 033	JTZ	BWAIT	
15574	074 347	CP	0347	IGNORE SYMBOLIC RECORDS
15576	150 131 033	JTZ	BWAIT	
15601	074 201	CP	0201	ELSE IT MUST BE A FILE MARKER
15603	110 256 033	JFZ	BSTOP	
15606	306	LAL		GET THE FILE NUMBER COMPLEMENTED
15607	054 377	XR	0377	
15611	275	OPH		IT MUST MATCH THE FILE NUMBER
15612	110 256 033	JFZ	BSTOP	
15615	046 000	LE	0	FLIP OVER THE FILE NUMBER
15617	036 010	LD	8	
15621	305	LAH		
15622	012	SRC		
15623	350	LHA		
15624	304	LAE		
15625	210	ACA		
15626	340	LEA		
15627	303	LAD		
15630	024 001	SU	1	
15632	330	LDA		
15633	110 221 033	JFZ	FLIP	COMPARE IT TO THE DESIRED FILE NUMBER
15636	304	LAE		
15637	271	CPB		
15640	160 072 033	JTS	BBACK	WE MUST GO IN THE OTHER DIRECTION
15643	110 131 033	JFZ	BWAIT	WE AREN'T BACK FAR ENOUGH
15646	106 333 033	CALL	DEKSTP	ELSE STOP THE TAPE
15651	161	EX	RBK	POSITION TO AFTER THE FILE LABEL
15652	106 016 033	CALL	DEKWAT	WAIT FOR IT
15655	007	RET		AND QUIT
15656	106 333 033	BSTOP:	CALL	DEKSTP
15661	161	EX	RBK	STOP THE TAPE
15662	106 341 033	CALL	DOKRTC	TRY THAT RECORD AGAIN
15665	160 070 034	JTS	DEKBAD	DECREMENT THE RE-TRY COUNT
15670	173	EX	SB	QUIT IT TOO MANY RE-TRIES
15671	104 146 033	JMP	BREAD	RE-INITIATE BACKWARD MOTION
15674	106 333 033	DEKTHE:	CALL	DEKSTP
15677	167	EX	BSP	STOP THE TAPE
15700	106 016 033	CALL	DEKWAT	APPROACH THE GAP FROM FORWARD DIRECTION
15703	161	EX	RBK	
15704	106 016 033	CALL	DEKWAT	
15707	007	RET		AND QUIT
		: POSITION TO THE END OF THE FILE		
15710		PEF\$:		
15710	106 355 033	DEKPEF:	CALL DEKFNS	SEARCH FOR THE NEXT FILE MARKER

15713	106 333 033	CALL DEKSTP	STOP THE TAPE
15716	167	EX BSP	POSITION IT TO AFTER THE LAST RECORD
15717	106 016 033	CALL DEKWAT	IN A FORWARD DIRECTION
15722	167	EX BSP	
15723	106 016 033	CALL DEKWAT	
15726	161	EX RBK	
15727	106 016 033	CALL DEKWAT	
15732	007	RET	
; STOP THE TAPE AND RE-SELECT THE PROPER DECK			
15733	177	DEKSTP: EX TSTOP	STOP THE TAPE
15734	106 016 033	CALL DEKWAT	WAIT FOR IT TO STOP
15737	007	RET	
; BACK UP THE TAPE AND DECREMENT THE RE-TRY COUNT			
15740	167	DEKRTO: EX BSP	
15741	106 016 033	DEKRTO: CALL DEKWAT	
15744	066 365 056 031	HL RTC	
15750	307	LAM	
15751	024 001	SU 1	
15753	370	LMA	
15754	007	RET	
; SEARCH FORWARD FOR A FILE MARKER			
15755	106 355 031	DEKFNS: CALL RTCI	INITIATE THE RE-TRY COUNT
15760	106 016 033	DEKFNA: CALL DEKWAT	WAIT FOR THE DECK TO BE READY
15763	171	EX SF	START FORWARD MOTION
15764	104 375 033	JMP DEKFNN	INSPECT THE NEXT RECORD
15767	106 363 032	DEKFNW: CALL DEKRED	WAIT FOR BLOCK TO BE OVER
15772	100 367 033	JFC DEKFNW	
15775	026 004	DEKFNN: LC 4	WAIT FOR DATA
15777	106 020 033	CALL WAIT	
16002	106 363 032	CALL DEKRED	GET THE RECORD TYPE
16005	330	LDA	SAVE THE CHARACTER
16006	106 363 032	CALL DEKRED	GET THE RECORD TYPE COMPLEMENTED
16011	054 377	XR 0377	UN-COMPLEMENT IT
16013	273	CPD	THEY MUST MATCH
16014	110 057 034	JFZ DEKFNE	
16017	074 303	CP 0303	IGNORE NON-FILE MARKERS
16021	150 367 033	JTZ DEKFNW	
16024	074 347	CP 0347	
16026	150 367 033	JTZ DEKFNW	
16031	074 201	OP 0201	ELSE IT MUST BE A FILE MARKER
16033	110 057 034	JFZ DEKFNE	
16036	106 363 032	CALL DEKRED	GET THE FILE NUMBER
16041	330	LDA	SAVE IT
16042	106 363 032	CALL DEKRED	GET THE FILE NUMBER COMPLEMENTED
16045	054 377	XR 0377	UN-COMPLEMENT IT
16047	273	CPD	THEY MUST MATCH
16050	110 057 034	JFZ DEKFNE	
16053	106 363 032	CALL DEKRED	THIS MUST BE THE END OF THE RECORD

16056	043		RTC	
16057	106 333 033	DEKFNE:	CALL DEKSTP	STOP THE TAPE
16062	106 340 033		CALL DEKRTC	BACK UP AND COUNT TRY
16065	120 360 033		JFS DEKFNA	TRY AGAIN IF NOT ALREADY TOO MANY
16070	006 002	DEKBAD:	LA 2	ELSE UNLOADABLE RECORD
16072	104 047 030		JMP ERR\$	
 : SPECIAL TAPE ROUTINES				
16075	106 152 032	TRWX:	CALL GETPKT	REWIND THE TAPE
16100	301		LAB	
16101	074 002		OP 2	
16103	006 001		LA 1	
16105	110 047 030		JFZ ERR\$	ONLY REWIND THE FRONT DECK
16110	106 046 033		CALL REWND\$	
16113	007		RET	
16114	106 152 032	TFNFX:	CALL GETPKT	READ A FILE NUMBER
16117	106 355 033	TRWFNR:	CALL DEKFNS	FIND A FILE MARKER
16122	106 333 033		CALL DEKSTP	STOP THE TAPE AFTER IT
16125	323		LCD	PUT THE FILE NUMBER IN THE C REGISTER
16126	007		RET	
16127	106 152 032	TFNWX:	CALL GETPKT	WRITE A FILE NUMBER
16132	106 016 033		CALL DEKWAT	
16135	163		EX WBK	
16136	036 201		LD 0201	
16140	106 002 033		CALL DEKWRT	
16143	036 176		LD 0176	
16145	106 002 033		CALL DEKWRT	
16150	066 256 056 032		HL PFN	
16154	337		LDM	
16155	106 002 033		CALL DEKWRT	
16160	303		LAD	
16161	054 377		XR 0377	
16163	330		LDA	
16164	106 002 033		CALL DEKWRT	
16167	106 016 033		CALL DEKWAT	
16172	007		RET	TERMINATE THE WRITE OPERATION

DONE

16200 SET 016200
 ITSEY BITSEY DEBUG

16200	066 273 056 035	DEBUG\$# HL CURADR
16204	347	LEM
16205	066 274	LL CURADR+1
16207	337	LDM
16210	026 005	LC 5
16212	066 306	LL DSPADR+4
16214	106 226 035	CALL CONBIN
16217	066 273	LL CURADR
16221	347	LEM
16222	066 274	LL CURADR+1
16224	357	LHM
16225	364	LLE
16226	347	LEM
16227	036 000	LD Ø
16231	026 003	LC 3
16233	066 313 056 035	HL DSPDAT+2
16237	106 226 035	CALL CONBIN
16242	066 275	LL DISP
16244	106 151 036	CALL DSPLY\$
16247	066 315	LL INBUF
16251	026 010	LC 8
16253	106 000 036	CALL KEYIN\$
16256	066 315	LL INBUF
16260	106 155 035	CALL CONOCT
16263	074 015	OP 015
16265	152 355 034	CTZ NEWADR
16270	074 111	OP 'I'
16272	152 364 034	CTZ INCADR
16275	074 104	OP 'D'
16277	152 002 035	CTZ DECADR
16302	074 115	OP 'M'
16304	152 020 035	CTZ MODIFY
16307	074 056	OP ','
16311	152 031 035	CTZ ENTER
16314	074 114	OP 'L'
16316	152 121 035	CTZ LSAVE
16321	074 110	OP 'H'
16323	152 144 035	CTZ HSAVE
16326	074 107	OP 'G'
16330	152 111 035	CTZ GET
16333	074 106	OP 'F'
16335	152 101 035	CTZ FRONT
16340	074 117	OP 'O'
16342	150 042 012	JTZ OS\$
16345	074 103	OP 'C'
16347	150 042 035	JTZ GOTO
16352	104 200 034	JMP DEBUG\$
16355	066 273	NEWADR: LL CURADR
16357	374	LME

16360	066 274		LL	CURADR+1
16362	373		LMD	
16363	007		RET	
16364	066 273	INCADR:	LL	CURADR
16366	307		LAM	
16367	004 001		AD	1
16371	370		LMA	
16372	066 274		LL	CURADR+1
16374	307		LAM	
16375	014 000		AC	0
16377	370		LMA	
16400	250		XRA	
16401	007		RET	
16402	066 273	DECADR:	LL	CURADR
16404	307		LAM	
16405	024 001		SU	1
16407	370		LMA	
16410	066 274		LL	CURADR+1
16412	307		LAM	
16413	034 000		SB	0
16415	370		LMA	
16416	250		XRA	
16417	007		RET	
16420	066 273	MODIFY:	LL	CURADR
16422	337		LDM	
16423	066 274		LL	CURADR+1
16425	357		LHM	
16426	363		LLD	
16427	374		LME	
16430	007		RET	
16431	106 020 035	ENTER:	CALL	MODIFY
16434	056 035		LH	CURADR>8
16436	106 364 034		CALL	INCADR
16441	007		RET	
16442	066 051	GOTO:	LL	BRANCH+1
16444	374		LME	
16445	066 052		LL	BRANCH+2
16447	373		LMD	
16450	106 000 000	BRANCH:	CALL	0
16453	300	LSAVI:	LAA	
16454	300	HSAVI:	LAA	
16455	066 370 056 035		HL	ASAVE
16461	370		LMA	
16462	066 371		LL	BSAVE
16464	371		LMB	
16465	066 372		LL	CSAVE
16467	372		LMC	
16470	066 373		LL	DSAVE
16472	373		LMD	
16473	066 374		LL	ESAVE
16475	374		LME	
16476	104 200 034		JMP	DEBUG\$
16501	016 001	FRONT:	LB	1

16503 106 112 000 CALL LOAD2\$
 16506 104 115 035 JMP GETLOD

16511 314 GET: LBE
 16512 106 100 000 CALL LOAD\$
 16515 053 GETLOD: RTZ
 16516 151 EX BEEP
 16517 250 XRA
 16520 007 RET

16521 066 053 056 035 LSAVE: HL LSAVI
 16525 016 306 LB 0306
 16527 304 HLSAVM: LAE
 16530 044 007 ND 7
 16532 002 SLC
 16533 002 SLC
 16534 002 SLC
 16535 074 070 OP 070
 16537 023 RFS
 16540 261 ORB
 16541 370 LMA
 16542 250 XRA
 16543 007 RET

16544 066 054 056 035 HSAVE: HL HSAVI
 16550 016 305 LB 0305
 16552 104 127 035 JMP HLSAVM

: CONVERT OCTAL TO BINARY

16555 036 000 CONOCT: LD 0
 16557 343 LED
 16560 317 CONLOP: LBM
 16561 106 353 036 CALL INCHL
 16564 301 LAB
 16565 074 070 CP '8'
 16567 023 RFS
 16570 074 060 CP '0'
 16572 063 RTS
 16573 044 007 ND 7
 16575 320 LOA
 16576 303 LAD
 16577 044 037 ND 037
 16601 002 SLC
 16602 002 SLC
 16603 002 SLC
 16604 330 LDA
 16605 304 LAE
 16606 002 SLC
 16607 002 SLC
 16610 002 SLC
 16611 340 LEA
 16612 044 007 ND 7
 16614 263 ORD

16615	330	LDA
16616	304	LAE
16617	044 370	ND 0370
16621	262	ORC
16622	340	LEA
16623	104 160 035	JMP CONLOP

: CONVERT BINARY TO OCTAL (RIGHT TO LEFT)

16626	304	CONBIN: LAE
16627	044 007	ND 7
16631	004 060	AD '0'
16633	370	LMA
16634	106 364 036	CALL DECHL
16637	304	LAE
16640	012	SRC
16641	012	SRC
16642	012	SRC
16643	044 037	ND 037
16645	340	LEA
16646	303	LAD
16647	012	SRC
16650	012	SRC
16651	012	SRC
16652	330	LDA
16653	044 340	ND 0340
16655	264	ORE
16656	340	LEA
16657	303	LAD
16660	044 037	ND 037
16662	330	LDA
16663	302	LAC
16664	024 001	SU 1
16666	320	LCA
16667	110 226 035	JFZ CONBIN
16672	007	RET

: STORAGE

16673	000 000	CURADR: DA 0
16675	011 000 013 013	DISP: DC 011.0.013.11.021
16702	040 040 040 040	DSPADR: DC
16711	040 040 040 015	DSPDAT: DC .015
16715	040 040 040 040	INBUF: DC
16770	000	SET 016770
16770	000	ASAVE: DC 0
16771	000	BSAVE: DC 0
16772	000	CSAVE: DC 0
16773	000	DSAVE: DC 0
16774	000	ESAVE: DC 0
16775	001 002	DC 1.2

DONE

17000

SET 017000

KEYBOARD ENTRY ROUTINE

- ACCEPTS A STRING OF CHARACTERS FROM THE KEYBOARD AND PUTS THEM IN MEMORY STARTING WITH THE ADDRESS GIVEN IN THE H AND L REGISTERS AND AT A DISPLAY POSITION DESCRIBED BY THE D (HORZ) AND E (VERT) REGISTERS. THE MAXIMUM NUMBER OF CHARACTERS ACCEPTED IS TAKEN FROM THE C REGISTER UPON ENTRY.
- OVERFLOW OFF THE END OF A DISPLAY LINE IS NOT PERMITTED AND IF EITHER THE MAXIMUM COUNT OR DISPLAY BOUNDARY IS EXCEEDED, SUCCESSIVE CHARACTERS WILL GO IN THE LAST POSITION OVER AND OVER. AN 015 WILL TERMINATE INPUT REQUEST.
- THE CURSOR IS TURNED ON UPON ENTRY AND OFF UPON EXIT.

17000	006 341	KEYIN\$#	LA 0341	ADDRESS THE KEYBOARD
17002	121		EX ADR	
17003	312		LBC	
17004	006 020		LA 020	
17006	131		EX COM1	
17007	106 326 036	KILOOP:	CALL CHAIT	MAKE SURE THE DISPLAY IS READY
17012	123	KWLOOP:	EX STATUS	GET A CHARACTER FROM THE KEYBOARD
17013	101		IN	
17014	044 002		ND 2	
17016	150 012 036		JTZ KWLOOP	
17021	125		EX DATA	
17022	101		IN	
17023	074 010		OP 010	CATCH BACKSPACE
17025	150 105 036		JTZ KBSP	
17030	074 030		OP 030	CATCH DELETE
17032	150 113 036		JTZ KDEL	
17035	074 100		OP 0100	REVERSE THE SHIFT KEY FUNCTION
17037	160 044 036		JTS KSTORE	
17042	054 040		XR 040	
17044	370	KSTORE:	LMA	STORE THE CHARACTER
17045	074 015		CP 015	CATCH THE ENTER KEY
17047	150 102 036		JTZ KEND	
17052	127		EX WRITE	ELSE DISPLAY THE CHARACTER
17053	303		LAD	CATCH CURSOR AT SCREEN BOUNDARY
17054	074 117		OP 79	
17056	120 007 036		JFS KILOOP	
17061	301		LAB	DECREMENT THE CHARACTER COUNT
17062	024 001		SU 1	
17064	160 007 036		JTS KILOOP	ALREADY ABOVE THE MAXIMUM
17067	310		LBA	
17070	303		LAD	BUMP THE CURSOR POSITION FOR REAL
17071	004 001		AD 1	
17073	330		LDA	
17074	106 353 036	CALL	INCHL	BUMP THE MEMORY LOCATION
17077	104 007 036	JMP	KILOOP	DO THE NEXT CHARACTER
17102	250	KEND:	XRA	
17103	131		EX COM1	TURN OFF THE CURSOR
17104	007		RET	

17105 106 124 036	KBSP:	CALL KBSPR	BACKSPACE ONE CHARACTER
17110 104 007 036		JMP KILOOP	
17113 106 124 036	KDEL:	CALL KBSPR	BACKSPACE TO THE BEGINNING OF THE ENTRY
17116 110 113 036		JFZ KDEL	
17121 104 007 036		JMP KILOOP	
17124 301	KBSPR:	LAB	INCREMENT THE CHARACTER COUNTER
17125 272		CPC	UNLESS AT THE BEGINNING OF THE ENTRY
17126 053		RTZ	
17127 004 001		AD 1	
17131 310		LBA	
17132 303		LAD	DECREMENT THE SCREEN POSITION
17133 024 001		SU 1	
17135 330		LDA	
17136 106 364 036		CALL DECHL	DECREMENT THE MEMORY POINTER
17141 106 326 036		CALL CHAIT	MAKE SURE THE DISPLAY IS READY
17144 006 040		LA 0440	ERASE THE CHARACTER
17146 127		EX WRITE	
17147 260		ORA	RETURN WITH ZERO CONDITION FALSE
17150 007		RET	

ORT DISPLAY ROUTINE

- DISPLAYS A STRING OF CHARACTERS WHICH ARE IN MEMORY STARTING WITH THE ADDRESS GIVEN IN THE H AND L REGISTERS AND AT THE POSITION DESCRIBED BY THE D (HORZ) AND E (VERT) REGISTERS.
- OVERFLOW OFF THE END OF A LINE IS NOT PERMITTED.
- SPECIAL CONTROL CHARACTERS TERMINATE THE LINE AND ALLOW MOVEMENT OF THE CURSOR, ERASURE OF THE SCREEN OR LINE, AND ROLL-UP OF THE ENTIRE SCREEN.

- ENTRY VALUES: D - HORIZONTAL CURSOR POSITION (0 TO 79)
E - VERTICAL CURSOR POSITION (0 TO 11)
HL - FIRST CHARACTER LOCATION IN STRING
DE - CURSOR POSITION AFTER LAST CHAR
HL - MEMORY LOCATION AFTER TERM CHAR
- CONTROL CHARACTERS: 003 - END OF THE STRING.
011 - A NEW HORIZONTAL POSITION FOLLOWS
013 - A NEW VERTICAL POSITION FOLLOWS
015 - END OF LINE (DOES CR/LF)
021 - ERASE TO THE END OF THE FRAME
022 - ERASE TO THE END OF THE LINE
023 - ROLL UP THE SCREEN ONE LINE

17151 006 341	DSPLY\$#	LA 0341	ADDRESS THE DISPLAY
17153 121		EX ADR	
17154 250		XRA	TURN OFF THE CURSOR
17155 131	DOCOM:	EX COM1	DO THE CONTROL COMMAND
17156 106 326 036	DLOOP:	CALL CHAIT	MAKE SURE THE DISPLAY IS READY
17161 317		LBM	GET A CHARACTER FROM THE STRING
17162 106 353 036		CALL INCHL	BUMP THE STRING POINTER
17165 301		LAB	CHECK FOR CONTROL CHARACTERS
17166 044 177		ND 0177	STRIP ANY PARITY

17170	074 003	OP	3		
17172	150 265 036	JTZ	ENDOS	END OF STRING	
17175	074 011	OP	011		
17177	150 300 036	JTZ	PHORZ	POSITION HORIZONTALLY	
17202	074 013	OP	013		
17204	150 274 036	JTZ	PVERT	POSITION VERTICALLY	
17207	074 015	OP	015		
17211	150 245 036	JTZ	ENDOL	END OF LINE	
17214	074 021	OP	021		
17216	150 307 036	JTZ	EEOF	ERASE TO THE END OF THE FRAME	
17221	074 022	OP	022		
17223	150 314 036	JTZ	EEOL	ERASE TO THE END OF THE LINE	
17226	074 023	OP	023		
17230	150 321 036	JTZ	ROLLUP	ROLL UP THE SCREEN	
17233	127	EX	WRITE	PUT OUT THE CHARACTER	
17234	303	LAD		BUMP THE CURSOR POSITION	
17235	074 117	OP	79	UNLESS AT THE END OF THE LINE	
17237	014 000	AC	0		
17241	330	LDA			
17242	104 156 036	JMP	DLOOP		
17245	036 000	ENDOL:	LD	• 0	RETURN CURSOR TO START OF NEXT LINE
17247	304		LAE		BUMP THE LINE COUNTER
17250	004 001		AD	1	
17252	340		LEA		
17253	074 014		CP	12	
17255	160 265 036		JTS	ENDOS	THERE IS ROOM FOR THE NEXT LINE
17260	046 013		LE	11	ELSE KEEP THE LINE COUNTER AT ELEVEN
17262	006 010		LA	010	AND ROLL THE SCREEN UP ONE LINE
17264	131		EX	COM1	
17265	106 326 036	ENDOS:	CALL	CHAIT	MAKE SURE THE DISPLAY IS READY
17270	006 020		LA	020	TURN ON THE CURSOR
17272	131		EX	COM1	
17273	007		RET		RETURN
17274	347	PVERT:	LEM		SET THE VERTICAL POSITION
17275	104 301 036		JMP	NCHAR	
17300	337	PHORZ:	LDM		
17301	106 353 036	NCHAR:	CALL	INCHL	SET THE HORIZONTAL POSITION
17304	104 156 036		JMP	DLOOP	BUMP THE STRING POINTER TO THE NXT CHAR
17307	006 004	EEOF:	LA	4	
17311	104 155 036		JMP	DOCOM	
17314	006 002	EEOL:	LA	2	
17316	104 155 036		JMP	DOCOM	
17321	006 010	ROLLUP:	LA	010	
17323	104 155 036		JMP	DOCOM	
17326	123	CHAIT:	EX	STATUS	WAIT FOR THE DISPLAY TO BE READY
17327	101		IN		
17330	012		SRC		
17331	100 326 036		JFC	CHAIT	
17334	303		LAD		
17335	260		ORA		
17336	063		RTS		MAKE SURE CURSOR IS IN CORRECT POSITION PREVENT CURSOR POSITIONS OUT OF RANGE

17337	074	120	CP	80	
17341	023		RFS		
17342	133		EX	COM2	
17343	304		LAF		
17344	260		ORA		
17345	063		RTS		
17346	074	014	CP	12	
17350	023		RFS		
17351	135		EX	COM3	
17352	007		RET		
17353	306		INCHL*	LAL	BUMP MEMORY POINTER UP
17354	004	001		AD	1
17356	360			LLA	
17357	305			LAH	
17360	014	000		AC	0
17362	350			LHA	
17363	007			RET	
17364	306		DEOHL*	LAL	BUMP MEMORY POINTER DOWN
17365	024	001		SU	1
17367	360			LLA	
17370	305			LAH	
17371	034	000		SB	0
17373	350			LHA	
17374	007			RET	

DONE

LIBRARY CATALOG

17404		SET	017404	
17404	037	CAT#*	DC 037	STARTING ADDRESS FOR LOADER
17405	010		DC 010	
17406	340		DC 0340	STARTING ADDRESS COMPLEMENTED
17407	367		DC 0367	
17410		CAT#*	RPT 14	SPACE FOR 14 ENTRIES
17410	040 040 040 040		DC	
17420	040 040 040 040		DC	
17430	040 040 040 040		DC	
17440	040 040 040 040		DC	
17450	040 040 040 040		DC	
17460	040 040 040 040		DC	
17470	040 040 040 040		DC	
17500	040 040 040 040		DC	
17510	040 040 040 040		DC	
17520	040 040 040 040		DC	
17530	040 040 040 040		DC	
17540	040 040 040 040		DC	
17550	040 040 040 040		DC	
17560	040 040 040 040		DC	
17570	052		DC *	
17571	000	ALPFN#	DC 0	AUTO-LOAD PHYSICAL FILE NUMBER
END OF PHYSICAL FILE 1				
17572	040 040 040 040	SYMBOL#	DC	ITEM SYMBOL STORAGE
LOAD AND EXECUTE				
17601	106 100 000	MAUTO\$*	CALL LOAD\$	LOAD THE GIVEN FILE
17604	100 075 000	MAUTO:	JFC RUN\$	EXECUTE IT IF GOOD LOAD
17607	104 004 000		JMP BOOT\$	ELSE RE-LOAD THE OPERATING SYSTEM
17612	106 112 000	MAUT2\$*	CALL LOADE\$	LOAD DECK TWO FILE
17615	104 204 037		JMP MAUTO	
SYMBOLIC FILE LOADER				
17620	353	MLOAD\$*	LHD	GET PACKET ADDRESS
17621	364		LLE	
17622	046 172 036 037		DE SYMBOL	PUT THE NAME IN THE LOOKUP ITEM
17626	026 006		LC 6	
17630	106 345 037		CALL BLKTFR	
17633	046 010 036 037		DE CAT	LOOK IT UP IN THE LIBRARY CATALOG
17637	106 264 037		CALL LOOKUP	
17642	074 005		CP 5	SEE IF IT IS THE THE CATALOG
17644	013		RFZ	ZERO FLAG FALSE IF IT ISN'T
17645	306		LAL	CALCULATE THE FILE NUMBER
17646	044 370		ND 0370	
17650	024 010		SU CAT	

17652	012	SRC	
17653	012	SRC	
17654	012	SRC	
17655	004 002	AD 2	FIRST ENTRY IS PHYSICAL FILE TWO
17657	310	LBA	
17660	106 100 000	CALL LOAD\$	
17663	007	RET	
: SYMBOL LOOKUP ROUTINE			
17664	353	LOOKUP* LHD	CHECK FIRST ENTRY IN TABLE
17665	364	LLE	
17666	104 334 037	JMP LSTART	
17671	066 172 056 037	LOOKPU: HL SYMBOL	GET THE ITEM STARTING ADDRESS
17675	327	LSLOOP: LCM	GET THE NEXT ITEM CHARACTER
17676	106 365 037	CALL INCSPW	GET THE NEXT TABLE ADDRESS
17701	307	LAM	GET THE NEXT TABLE CHARACTER
17702	272	OPC	
17703	110 322 037	JFZ LDIFP	THEY DON'T MATCH
17706	306	LAL	SEE IF AT THE END OF THE ENTRY
17707	044 007	ND 7	
17711	074 005	CP 5	
17713	053	RTZ	
17714	106 365 037	CALL INCSPW	THE ITEM HAS BEEN FOUND IF SO
17717	104 275 037	JMP LSLOOP	GET THE NEXT ITEM ADDRESS
17722	306	LDIFF: LAL	AND TRY THE NEXT CHARACTER
17723	044 370	ND 0370	BUMP THE TABLE POINTER TO NEXT ENTRY
17725	004 010	AD 8	
17727	360	LLA	
17730	305	LAH	
17731	014 000	AC 0	
17733	350	LHA	
17734	307	LSTART: LAM	GET THE TABLE FIRST CHARACTER
17735	074 101	CP 'A'	END OF TABLE IF IT IS NOT ALPHA
17737	063	RTS	
17740	335	LDH	SAVE THE TABLE ADDRESS
17741	346	LEL	
17742	104 271 037	JMP LOOKPU	AND TRY NEXT TABLE ENTRY
: BLOCK TRANSFER FROM HL TO DE C CHARACTERS			
17745	317	BLKTFR* LBM	GET A SOURCE CHARACTER
17746	106 365 037	CALL INCSPW	GET NEXT DESTINATION LOCATION
17751	371	LMB	PUT IT IN A DESTINATION LOCATION
17752	106 365 037	CALL INCSPW	GET NEXT SOURCE ADDRESS
17755	302	LAC	DECREMENT THE COUNT
17756	024 001	SU 1	
17760	320	LCA	
17761	110 345 037	JFZ BLKTFR	DO NEXT CHAR IF NOT ZERO
17764	007	RET	
: INCREMENT HL AND SWAP IT WITH DE			
17765	306	INCSPW* LAL	

17766	004	001	AD	1
17770	364		LLE	
17771	340		LEA	
17772	305		LAH	
17773	014	000	AC	0
17775	353		LHD	
17776	330		LDA	
17777	007		RET	

DONE