

Boolean Algebra: (Switching Algebra)

Three main operators: $\left\{ \begin{array}{l} \text{NOT: } ()', (\overline{}), \sim () \\ \text{AND: } \bullet \\ \text{OR: } + \end{array} \right.$

Precedence: $\left\{ \begin{array}{l} () \\ \text{NOT} \\ \text{AND} \\ \text{OR} \end{array} \right.$

Axioms: A set of basic definitions
that we assume to be true.

↳ use Axioms to develop theorems.

$$(T3): \quad X + X = X \quad \checkmark$$

Perfect induction: The theorem needs to be
checked to be true
for all input combinations
(truth Table)

X	X + X	
0	0 + 0	$\xrightarrow{A4D} 0 = X$
1	1 + 1	$\xrightarrow{A3D} 1 = X$

$$(T9): \quad X + \underline{X \cdot Y} = X$$

X	Y	$X \cdot Y$	$X + X \cdot Y$
0	0	0	0
0	1	0	0
1	0	0	1
1	1	1	1

For theorems involving more than 3 variables
we can use "finite induction" approach to
prove them:

→ n variables

① $n=2$, use perfect induction or
other approaches to prove the theorem
is valid.

(2) Assume the theorem is valid for n variables, then show it is going to be valid for $n+1$ variables.

Use Boolean Axiom and Theorems above T9, to prove T9.

$$X + X \cdot Y = X$$

$$X \cdot 1 + X \cdot Y$$

used (T1D)

$$= X \cdot (1 + Y)$$

used (T8)

$$= X \cdot 1$$

used (T2)

$$= X$$

$$T8: X \cdot Y + X \cdot Z = X \cdot (Y + Z)$$

Use
Finite induction to prove DeMorgan's
Theorems. (T13)

$$[A \cdot (B' \cdot C' + B \cdot C)]'$$

$$\begin{cases} (X + Y)' = X' \cdot Y' \\ (X \cdot Y)' = X' + Y' \end{cases}$$

$$= A' + (B' \cdot C' + B \cdot C)'$$

$$= A' + (B' \cdot C')' \cdot (B \cdot C)'$$

$$= A' + \underbrace{(B + C) \cdot (B' + C')}_{b}$$

$$\underbrace{B \cdot B'}_{0} + \underbrace{B \cdot C' + C \cdot B'}_{0} + \underbrace{C \cdot C'}_{0}$$

$$= A' + B \cdot C' + C \cdot B'$$

$$F(x, y, z) = x \cdot y + x' \cdot y' + x' \cdot z'$$

Express $F(x, y, z)$ as Sum of minterms:

$$x \cdot y = x \cdot y \cdot \underset{\substack{\uparrow \\ (z+z')}}{1} = x \cdot y \cdot (z+z') = x \cdot y \cdot z + x \cdot y \cdot z'$$

$$x' \cdot y' = x' \cdot y' \cdot \underset{\substack{\uparrow \\ (z+z')}}{1} = x' \cdot y' \cdot (z+z') = x' \cdot y' \cdot z + x' \cdot y' \cdot z'$$

$$X' \cdot Z' = X' \cdot Z' \cdot \underset{\substack{\uparrow \\ (Y+Y')}}{1} = X' \cdot Z' \cdot Y + X' \cdot Z' \cdot Y'$$

$$F(X, Y, Z) = X \cdot Y \cdot Z + X \cdot Y \cdot Z' + X' \cdot Y' \cdot Z + X' \cdot Y' \cdot Z'$$

$+ X' \cdot Z' \cdot Y + \cancel{X' \cdot Z' \cdot Y'}$
 $(X+X=X)$

$\cancel{X+X=X}$
 $X' \cdot Y \cdot Z'$

$$= \sum_{X, Y, Z} (m_7, m_6, m_1, m_0, m_2)$$

① Expand the function into sum of products.

② In each product, if there is a literal missing, "AND" the product with "1" and replace "1" with (missing literal + its complement)

$F(x, y, z) = x \cdot y + x' \cdot y' + y \cdot z$

Express $F(x, y, z)$ as product of maxterms.

$$(A + BC) = (A + B) \cdot (A + C) \quad \leftarrow \text{Theorem}$$

$$\begin{aligned}
 & X.Y + \underbrace{X'.Y'}_A + \underbrace{\frac{Y'}{B}.\frac{Z'}{C}} \\
 = & (\underbrace{X.Y}_{\uparrow C \uparrow B} + \underbrace{X'.Y'}_A + Y') \cdot (\underbrace{X.Y}_{\uparrow C \uparrow B} + \underbrace{X'.Y' + Z'}_A) \\
 = & (\underbrace{X'.Y' + Y' + Y}_{{}=1}) \cdot (\underbrace{X'.Y' + Y' + X}_{\downarrow}) \cdot (\underbrace{X'.Y' + Z' + X}_{\downarrow}) \cdot (\underbrace{X'.Y' + Z' + Y}_{{}=1}) \\
 = & (\underbrace{X + Y' + Y'}_{=1}) \cdot (\underbrace{X + Y' + X'}_{=1}) \cdot (\underbrace{X + Z' + X'}_{=1}) \cdot (\underbrace{X + Z' + Y'}_{=1}) \\
 = & (Z' + Y + X') \cdot (Z' + Y + Y') \\
 = & 1 + Z' = 1
 \end{aligned}$$

$$= (X + Y') \cdot \underbrace{(X + Z' + Y')}_{\text{maxterm}} \cdot \underbrace{(X' + Y + Z')}_{\text{maxterm}}$$

$$(X + Y') = (X + Y' + \underbrace{0}_{\substack{\uparrow \\ Z \cdot Z'}}) = \underbrace{(X + Y')}_A + \underbrace{Z}_B \cdot \underbrace{Z'}_C$$

$$= (X + Y' + Z) \cdot (X + Y' + Z')$$

$$F(x, y, z) = (X + Y' + Z) \cdot (X + Y' + Z') \cdot (X' + Y + Z')$$

$$= \prod_{x, y, z} (M_2, M_3, M_5)$$

$$X.Y + X'.Z + \underbrace{Y.Z}_1 = \underbrace{X.Y + X'.Z}$$

$$X.Y + X'.Z + 1.Y.Z$$

$$X.Y + X'.Z + (X+X').Y.Z$$

$$= \underline{X.Y} + \underline{X'.Z} + \underline{X.Y.Z} + \underline{X'.Y.Z}$$

$$= X.Y [1+Z] + X'.Z [1+Y]$$

$$= X.Y + X'.Z$$