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1 Data Structure

1.1 Dsu With Rollback [89 lines] bc2588

```
struct dsu_save {
    int v, rnkv, u, rnku;
    dsu_save() {}
    dsu_save(int _v, int _rnkv, int _u, int _rnku)
        : v(_v), rnkv(_rnkv), u(_u), rnku(_rnku) {}
};

struct dsu_with_rollbacks {
    vector<int> p, rnk;
    int comps;
    stack<dsu_save> op;
    dsu_with_rollbacks() {}
    dsu_with_rollbacks(int n) {
        p.resize(n);
        rnk.resize(n);
        for (int i = 0; i < n; i++) {
            p[i] = i;
            rnk[i] = 0;
        }
        comps = n;
    }
    int find_set(int v) { return (v == p[v]) ? v :
        find_set(p[v]); }
    bool unite(int v, int u) {
        v = find_set(v);
        u = find_set(u);
        if (v == u) return false;
        comps--;
        if (rnk[v] > rnk[u]) swap(v, u);
        op.push(dsu_save(v, rnk[v], u, rnk[u]));
```

```
p[v] = u;
    if (rnk[u] == rnk[v]) rnk[u]++;
    return true;
}

void rollback() {
    if (op.empty()) return;
    dsu_save x = op.top();
    op.pop();
    comps++;
    p[x.v] = x.v;
    rnk[x.v] = x.rnkv;
    p[x.u] = x.u;
    rnk[x.u] = x.rnku;
}

};

struct query {
    int v, u;
    bool united;
    query(int _v, int _u) : v(_v), u(_u) {}
};

struct QueryTree {
    vector<vector<query>> t;
    dsu_with_rollbacks dsu;
    int T;
    QueryTree() {}
    QueryTree(int _T, int n) : T(_T) {
        dsu = dsu_with_rollbacks(n);
        t.resize(4 * T + 4);
    }
    void add_to_tree(int v, int l, int r, int ul, int ur,
        query& q) {
        if (ul > ur) return;
        if (l == ul && r == ur) {
            t[v].push_back(q);
            return;
        }
        int mid = (l + r) / 2;
        add_to_tree(2 * v, l, mid, ul, min(ur, mid), q);
        add_to_tree(2 * v + 1, mid + 1, r, max(ul, mid +
            1), ur, q);
    }
    void add_query(query q, int l, int r) {
        add_to_tree(1, 0, T - 1, l, r, q);
    }
    void dfs(int v, int l, int r, vector<int>& ans) {
        for (query& q : t[v]) {
            q.united = dsu.unite(q.v, q.u);
        }
        if (l == r)
            ans[l] = dsu.comps;
        else {
            int mid = (l + r) / 2;
            dfs(2 * v, l, mid, ans);
            dfs(2 * v + 1, mid + 1, r, ans);
        }
        for (query q : t[v]) {
            if (q.united) dsu.rollback();
        }
    }
    vector<int> solve() {
        vector<int> ans(T);
        dfs(1, 0, T - 1, ans);
        return ans;
    }
};
```

1.2 MO with Update [43 lines]

2fbf87

```
//1 indexed
//Complexity:  $O(S \times Q + Q \times \frac{N^2}{S^2})$ 
//S =  $(2*n^2)^{(1/3)}$ 
const int block_size = 2720; // 4310 for 2e5
const int mx = 1e5 + 5;
struct Query {
    int L, R, T, id;
    Query() {}
    Query(int _L, int _R, int _T, int _id) : L(_L),
        R(_R), T(_T), id(_id) {}
    bool operator<(const Query &x) const {
        if (L / block_size == x.L / block_size) {
            if (R / block_size == x.R / block_size) return T <
                x.T;
            return R / block_size < x.R / block_size;
        }
        return L / block_size < x.L / block_size;
    }
} Q[mx];
struct Update {
    int pos;
    int old, cur;
    Update(){};
    Update(int _p, int _o, int _c) : pos(_p), old(_o),
        cur(_c){};
} U[mx];
int ans[mx];
inline void add(int id) {}
inline void remove(int id) {}
inline void update(int id, int L, int R) {}
inline void undo(int id, int L, int R) {}
inline int get() {}
void MO(int nq, int nu) {
    sort(Q + 1, Q + nq + 1);
    int L = 1, R = 0, T = nu;
    for (int i = 1; i <= nq; i++) {
        Query q = Q[i];
        while (T < q.T) update(++T, L, R);
        while (T > q.T) undo(T--, L, R);
        while (L > q.L) add(--L);
        while (R < q.R) add(++R);
        while (L < q.L) remove(L++);
        while (R > q.R) remove(R--);
        ans[q.id] = get();
    }
}
```

1.3 MO [28 lines]

bed3e5

```
const int N = 2e5 + 5;
const int Q = 2e5 + 5;
const int SZ = sqrt(N) + 1;
struct qry {
    int l, r, id, blk;
    bool operator<(const qry& p) const {
        return blk == p.blk ? r < p.r : blk < p.blk;
    }
};
qry query[Q];
ll ans[Q];
void add(int id) {}
void remove(int id) {}
ll get() {}
int n, q;
```

```
void MO() {
    sort(query, query + q);
    int cur_l = 0, cur_r = -1;
    for (int i = 0; i < q; i++) {
        qry q = query[i];
        while (cur_l > q.l) add(--cur_l);
        while (cur_r < q.r) add(++cur_r);
        while (cur_l < q.l) remove(cur_l++);
        while (cur_r > q.r) remove(cur_r--);
        ans[q.id] = get();
    }
}
/* 0 indexed. */
```

1.4 Persistent Segment Tree [53 lines]

359d65

```
const int mxn = 1e5+5; //CHECK here for problem
int root[mxn], leftchild[25*mxn], rightchild[25*mxn],
    value[25*mxn], a[mxn];
int now = 0, n;
int build(int L, int R){
    int node = ++now;
    if(L == R) return node;
    int mid = (L+R)>>1;
    leftchild[node] = build(L, mid);
    rightchild[node] = build(mid+1, R);
    //initialize value[node]
    return node;
}

int update(int nownode, int L, int R, int val){
    int node = ++now;
    if(L == R){
        //value[node] = value[nownode]+1;
        //update value[node]
        return node;
    }
    int mid = (L+R)>>1;
    leftchild[node] = leftchild[nownode];
    rightchild[node] = rightchild[nownode];
    if(mid <= val){ //change condition as required
        leftchild[node] = update(leftchild[nownode], L,
            mid, val);
    }
    else{
        rightchild[node] = update(rightchild[nownode],
            mid+1, R, val);
    }
    //value[node] = value[nownode]+1;
    //update value[node]
    return node;
}

int query(int leftnode, int rightnode, int L, int R, int
    k){
    if(L==R) return L;
    //int leftcnt = value[leftchild[rightnode]]-
        value[leftchild[leftnode]];a
    //change as required
    int mid = (L+R)>>1;
    if(leftcnt >= k){ //change condition as required
        return query(leftchild[leftnode],
            leftchild[rightnode], L, mid, k);
    }
    else{
```

```
        return query(rightchild[leftnode],
            rightchild[rightnode], mid+1, R, k-leftcnt);
    }
}

void persistentsegtree(){
    root[0] = build(0, mxn);
    for(int i=1; i<=n; i++){
        root[i] = update(root[i-1], 0, mxn, a[i]);
    }
}
```

1.5 SQRT Decomposition [96 lines]

a772d3

```
struct sqrtDecomposition {
    static const int sz = 320; //sz = sqrt(N);
    int numberofblocks;

    struct node {
        int L, R;
        bool islazy = false;
        ll lazyval=0;
        //extra data needed for different problems
        void ini(int l, int r) {
            for(int i=l; i<=r; i++) {
                //...initialize as need
            }
            L=l, R=r;
        }
        void semiupdate(int l, int r, ll val) {
            if(l>r) return;
            if(islazy){
                for(int i=L; i<=R; i++){
                    //...distribute lazy to everyone
                }
                islazy = 0;
                lazyval = 0;
            }
            for(int i=l; i<=r; i++){
                //...do it manually
            }
        }
        void fullupdate(ll val){
            if(islazy){
                //...only update lazyval
            }
            else{
                for(int i=L; i<=R; i++){
                    //...everyone are not equal, make them equal
                }
                islazy = 1;
                //update lazyval
            }
        }
        void update(int l, int r, ll val){
            if(l<=L && r>=R) fullupdate(val);
            else semiupdate(max(l, L), min(r, R), val);
        }
        ll semiquery(int l, int r){
            if(l>r) return 0;
            if(islazy){
                for(int i=L; i<=R; i++){
                    //...distribute lazy to everyone
                }
                islazy = 0;
            }
```

```

        lazyval = 0;
    }
    ll ret = 0;
    for(int i=l; i<=r; i++){
        //...take one by one
    }
    return ret;
}
ll fullquery(){
    //return stored value;
}
ll query(int l, int r){
    if(l<=L && r>=R) return fullquery();
    else return semiquery(max(l, L), min(r, R));
}
};
vector<node> blocks;
void init(int n){
    numberofblocks = (n+sz-1)/sz;
    int curL = 1, curR = sz;
    blocks.resize(numberofblocks+5);
    for(int i=1; i<=numberofblocks; i++){
        curR = min(n, curR);
        blocks[i].ini(curL, curR);
        curL += sz;
        curR += sz;
    }
}
void update(int l, int r, ll val){
    int left = (l-1)/sz+1;
    int right = (r-1)/sz+1;
    for(int i=left; i<=right; i++){
        blocks[i].update(l, r, val);
    }
}
ll query(int l, int r){
    int left = (l-1)/sz+1;
    int right = (r-1)/sz+1;
    ll ret = 0;
    for(int i=left; i<=right; i++){
        ret += blocks[i].query(l, r);
    }
    return ret;
}
};

```

1.6 Segment Tree [73 lines]

c1fe4f

```

/*edit:data,combine,build check datatype*/
template<typename T>
struct SegmentTree {
#define lc (C << 1)
#define rc (C << 1 | 1)
#define M ((L+R)>>1)
    struct data {
        T sum;
        data() :sum(0) {};
    };
    vector<data>st;
    vector<bool>isLazy;
    vector<T>lazy;
    int N;
    SegmentTree(int _N) :N(_N) {
        st.resize(4 * N);
        isLazy.resize(4 * N);
        lazy.resize(4 * N);
    }
};

```

```

}
void combine(data& cur, data& l, data& r) {
    cur.sum = l.sum + r.sum;
}
void push(int C, int L, int R) {
    if (!isLazy[C]) return;
    if (L != R) {
        isLazy[lc] = 1;
        isLazy[rc] = 1;
        lazy[lc] += lazy[C];
        lazy[rc] += lazy[C];
    }
    st[C].sum = (R - L + 1) * lazy[C];
    lazy[C] = 0;
    isLazy[C] = false;
}
void build(int C, int L, int R) {
    if (L == R) {
        st[C].sum = 0;
        return;
    }
    build(lc, L, M);
    build(rc, M + 1, R);
    combine(st[C], st[lc], st[rc]);
}
data Query(int i, int j, int C, int L, int R) {
    push(C, L, R);
    if (j < L || i > R || L > R) return data(); //
    default val 0/INF
    if (i <= L && R <= j) return st[C];
    data ret;
    data d1 = Query(i, j, lc, L, M);
    data d2 = Query(i, j, rc, M + 1, R);
    combine(ret, d1, d2);
    return ret;
}
void Update(int i, int j, T val, int C, int L, int R)
{
    push(C, L, R);
    if (j < L || i > R || L > R) return;
    if (i <= L && R <= j) {
        isLazy[C] = 1;
        lazy[C] = val;
        push(C, L, R);
        return;
    }
    Update(i, j, val, lc, L, M);
    Update(i, j, val, rc, M + 1, R);
    combine(st[C], st[lc], st[rc]);
}
void Update(int i, int j, T val) {
    Update(i, j, val, 1, 1, N);
}
T Query(int i, int j) {
    return Query(i, j, 1, 1, N).sum;
}
};

```

1.7 Sqrt Tricks [8 lines]

addf19

1. Size of the block is not always Sqrt, adjust it as necessary. if $O(n/b+b)$ then take $n/b = b$ and calculate b .
2. MO's Algorithm

*it is possible to solve a Mo problem without any remove operation. For L in one block R only increases, for every range we can start L from the last of that block

3. Sqrt Decomposition by time of queries.
 - *keep overall solution and \sqrt{n} updates in a vector and for a query iterate over all of them, when the vector size exceeds \sqrt{n} you can add these updates with overall solution using $O(n)$
4. If sum of N positive numbers are S, there are at most \sqrt{S} distinct values.
5. Randomization
6. Baby step, gaint step

1.8 Treap [152 lines]

a39318

```

struct Treap {
    struct Node {
        int val, priority, cnt; // value, priority, subtree
        size
        Node *l, *r; // left child, right child
        pointer
        Node() {} //rng from template
        Node(int key) : val(key), priority(rng()),
        l(nullptr), r(nullptr) {}
    };
    typedef Node *node;
    node root;
    Treap() : root(0) {}
    int cnt(node t) { return t ? t->cnt : 0; } // return
    subtree size
    void updateCnt(node t) {
        if (t) t->cnt = 1 + cnt(t->l) + cnt(t->r); //
        update subtree size
    }
    void push(node cur) {
        ; // Lazy Propagation
    }

    void combine(node &cur, node l, node r) {
        if (!l) {
            cur = r;
            return;
        }
        if (!r) {
            cur = l;
            return;
        }
        // Merge Operations like in segment tree
    }

    void reset(node &cur) {
        if (!cur) return; // To reset other fields of cur
        except value and cnt
    }

    void operation(node &cur) {
        if (!cur) return;
        reset(cur);
        combine(cur, cur->l, cur);
        combine(cur, cur, cur->r);
    }

    // Split(T,key): split the tree in two tree. Left
    pointer contains all value
    // less than or equal to key. Right pointer contains
    the rest.
};

```

```

void split(node t, node &l, node &r, int key) {
    if (!t)
        return void(l = r = nullptr);
    push(t);
    if (t->val <= key) {
        split(t->r, t->r, r, key), l = t;

    } else {
        split(t->l, l, t->l, key), r = t;
    }
    updateCnt(t);
    operation(t);
}

void splitPos(node t, node &l, node &r, int k, int add
= 0) {
    if (!t) return void(l = r = 0);
    push(t);
    int idx = add + cnt(t->l);
    if (idx <= k)
        splitPos(t->r, t->r, r, k, idx + 1), l = t;
    else
        splitPos(t->l, l, t->l, k, add), r = t;
    updateCnt(t);
    operation(t);
}

// Merge(T1,T2): merges 2 tree into one.The tree with
// root of higher
// priority becomes the new root.
void merge(node &t, node l, node r) {
    push(l);
    push(r);
    if (!l || !r)
        t = l ? l : r;
    else if (l->priority > r->priority)
        merge(l->r, l->r, r), t = l;
    else
        merge(r->l, l, r->l), t = r;
    updateCnt(t);
    operation(t);
}

// insert creates a set.all unique value.
void insert(int val) {
    if (!root) {
        root = new Node(val);
        return;
    }
    node l, r, mid, mid2, rr;
    mid = new Node(val);
    split(root, l, r, val);
    merge(l, l, mid); // these 3 lines will create
    multiset.
    merge(root, l, r);
    /*split(root, l, r, val - 1); // l contains all
    small values.
    merge(l, l, mid); // l contains new val
    too.
    split(r, mid2, rr, val); // rr contains all
    greater values.
    merge(root, l, rr);*/
}

// removes all similar values.
void erase(int val) {
    node l, r, mid;
    /* Removes all similar element*/

```

```

    split(root, l, r, val - 1);
    split(r, mid, r, val);
    merge(root, l, r);
    /*Removes single instance*/
    /*split(root, l, r, val - 1);
    split(r, mid, r, val);
    merge(mid, mid->l, mid->r);
    merge(l, l, mid);
    merge(root, l, r);*/
}

void clear(node cur) {
    if (!cur) return;
    clear(cur->l), clear(cur->r);
    delete cur;
}

void clear() { clear(root); }
void inorder(node t) {
    if (!t) return;
    inorder(t->l);
    cout << t->val << ' ';
    inorder(t->r);
}

void inorder() {
    inorder(root);
    puts("");
}

//1 indexed - xth element after sorting.
int find_by_order(int x) {
    if (!x) return -1;
    x--;
    node l, r, mid;
    splitPos(root, l, r, x - 1);
    splitPos(r, mid, r, 0);
    int ans = -1;
    if (cnt(mid) == 1) ans = mid->val;
    merge(r, mid, r);
    merge(root, l, r);
}

// 1 indexed. index of val in sorted array. -1 if not
// found.
int order_of_key(int val) {
    node l, r, mid;
    split(root, l, r, val - 1);
    split(r, mid, r, val);
    int ans = -1;
    if (cnt(mid) == 1) ans = 1 + cnt(l);
    merge(r, mid, r);
    merge(root, l, r);
    return ans;
}
};

```

1.9 Trie Bit [61 lines]

390174

```

struct Trie {
    struct node {
        int next[2];
        int cnt, fin;
        node() : cnt(0), fin(0) {
            for (int i = 0; i < 2; i++) next[i] = -1;
        }
    };
    vector<node>data;
    Trie() {

```

```

        data.push_back(node());
    }
}

void key_add(int val) {
    int cur = 0;
    for (int i = 30; i >= 0; i--) {
        int id = (val >> i) & 1;
        if (data[cur].next[id] == -1) {
            data[cur].next[id] = data.size();
            data.push_back(node());
        }
        cur = data[cur].next[id];
        data[cur].cnt++;
    }
    data[cur].fin++;
}

int key_search(int val) {
    int cur = 0;
    for (int i = 30; ~i; i--) {
        int id = (val >> i) & 1;
        if (data[cur].next[id] == -1) return 0;
        cur = data[cur].next[id];
    }
    return data[cur].fin;
}

void key_delete(int val) {
    int cur = 0;
    for (int i = 30; ~i; i--) {
        int id = (val >> i) & 1;
        cur = data[cur].next[id];
        data[cur].cnt--;
    }
    data[cur].fin--;
}

bool key_remove(int val) {
    if (key_search(val)) return key_delete(val), 1;
    return 0;
}

int maxXor(int x) {
    int cur = 0;
    int ans = 0;
    for (int i = 30; ~i; i--) {
        int b = (x >> i) & 1;
        if (data[cur].next[!b] + 1 &&
            data[data[cur].next[!b]].cnt > 0) {
            ans += (1LL << i);
            cur = data[cur].next[!b];
        }
        else cur = data[cur].next[b];
    }
    return ans;
}
};

```

2 Dynamic Programming

2.1 Convex Hull Trick [91 lines]

62db2e

```

struct Hull_Static{
    /*all m need to be decreasing order
    if m is in increasing order then negate the
    m(like, add_line(-m,c)),
    remember in query you have to negate the x also*/
    const ll inf=1000000000000000000;
    int min_or_max; //if min then 0 otherwise 1

```



```

int pointer; /*keep track for the best line for
previous query, requires all insert first*/
vector<ll> M, C; //y=m*x+c;
inline void clear(){
    min_or_max=0; //initially with minimum trick
    pointer=0;
    M.clear();
    C.clear();
}
Hull_Static(){clear();}
Hull_Static(int _min_or_max){
    clear();
    this->min_or_max=_min_or_max;
}
bool bad_min(int idx1, int idx2, int idx3){
    return (C[idx3]-C[idx1])*(M[idx1]-M[idx2]) <
           (C[idx2]-C[idx1])*(M[idx1]-M[idx3]);
    return 1.0*(C[idx3]-C[idx1])*(M[idx1]-M[idx2]) <
           1.0*(C[idx2]-C[idx1])*(M[idx1]-M[idx3]); //for
    overflow
}
bool bad_max(int idx1, int idx2, int idx3){
    return (C[idx3]-C[idx1])*(M[idx1]-M[idx2]) >
           (C[idx2]-C[idx1])*(M[idx1]-M[idx3]);
    return 1.0*(C[idx3]-C[idx1])*(M[idx1]-M[idx2]) >
           1.0*(C[idx2]-C[idx1])*(M[idx1]-M[idx3]); //for
    overflow
}
bool bad(int idx1, int idx2, int idx3){
    if(!min_or_max) return bad_min(idx1, idx2, idx3);
    else return bad_max(idx1, idx2, idx3);
}
void add_line(ll m, ll c) /*add line where m is given
in decreasing order
//if (M.size())>0 and M.back()==m) return; //same
gradient, no need to add
above line added from tarango khan, this line cost me
several wa, but some code got ac with this*/
M.push_back(m);
C.push_back(c);
while(M.size())>=3 and bad((int)M.size()-3,
(int)M.size()-2, (int)M.size()-1){
    M.erase(M.end()-2);
    C.erase(C.end()-2);
}
}
ll getval(ll idx, ll x){
    return M[idx]*x+C[idx];
}
ll getminval(ll x){ /*if queries are non-decreasing
order*/
    while(pointer<(int)M.size()-1 and getval(pointer+
1, x)<getval(pointer, x)) pointer++;
    return M[pointer]*x+C[pointer];
}
}
ll getmaxval(ll x){
    while(pointer<(int)M.size()-1 and getval(pointer+
1, x)>getval(pointer, x)) pointer++;
    return M[pointer]*x+C[pointer];
}
}
ll getminvalternary(ll x){
    ll lo=0, hi=(ll)M.size()-1;
    ll ans=inf; /*change with problem*/
    while(lo<=hi){

```

```

        ll mid1=lo+(hi-lo)/3;
        ll mid2=hi-(hi-lo)/3;
        ll val1=getval(mid1, x);
        ll val2=getval(mid2, x);
        if(val1<val2){
            ans=min(ans, val2);
            hi=mid2-1;
        }
        else{
            ans=min(ans, val1);
            lo=mid1+1;
        }
    }
    return ans;
}
ll getmaxvalternary(ll x){
    ll lo=0, hi=(ll)M.size()-1;
    ll ans=-inf; /*change with problem*/
    while(lo<=hi){
        ll mid1=lo+(hi-lo)/3;
        ll mid2=hi-(hi-lo)/3;
        ll val1=getval(mid1, x);
        ll val2=getval(mid2, x);
        if(val1<val2){
            ans=max(ans, val2);
            lo=mid1+1;
        }
        else{
            ans=max(ans, val1);
            hi=mid2-1;
        }
    }
    return ans;
}
};

```

2.2 Divide and Conquer DP [26 lines]

6d8559

```

ll G, L; //total group, cell size
ll dp[8001][801], cum[8001];
ll C[8001]; //value of each cell
inline ll cost(ll l, ll r){
    return cum[r]-cum[l-1]*(r-l+1);
}
void fn(ll g, ll st, ll ed, ll r1, ll r2){
    if(st>ed) return;
    ll mid=(st+ed)/2, pos=-1;
    dp[mid][g]=inf;
    for(int i=r1; i<=r2; i++){
        ll tcost=cost(i, mid)+dp[i-1][g-1];
        if(tcost<dp[mid][g]){
            dp[mid][g]=tcost, pos=i;
        }
    }
    fn(g, st, mid-1, r1, pos);
    fn(g, mid+1, ed, pos, r2);
}
int main(){
    for(int i=1; i<=L; i++)
        cum[i]=cum[i-1]+C[i];
    for(int i=1; i<=L; i++)
        dp[i][1]=cost(1, i);
    for(int i=2; i<=G; i++) fn(i, 1, L, 1, L);
}

```

2.3 Knuth Optimization [32 lines]

911417

/*It is applicable where recurrence is in the form :

$$dp[i][j] = \min_{k < j} \{ dp[i][k] + dp[k][j] \} + C[i][j]$$

condition for applicability is:

$$A[i, j-1] \leq A[i, j] \leq A[i+1, j]$$

Where,

$$A[i][j] - \text{the smallest } k \text{ that gives optimal answer, like-}$$

$$dp[i][j] = dp[i-1][k] + C[k][j]$$

$$C[i][j] - \text{given cost function}$$

also applicable if: $C[i][j]$ satisfies the following 2 conditions:

$$C[a][c] + C[b][d] \leq C[a][d] + C[b][c], a \leq b \leq c \leq d$$

$$C[b][c] \leq C[a][d], a \leq b \leq c \leq d$$

reduces time complexity from $O(n^3)$ to $O(n^2)$ */

```

for(int s=0; s<=k; s++) //s-length(size) of substring
    for(int l=0; l+s<=k; l++){ //l-left point
        int r=l+s; //r-right point
        if(s<2){
            res[l][r]=0; //DP base-nothing to break
            mid[l][r]=1; /*mid is equal to left border*/
            continue;
        }
        int mleft=mid[l][r-1]; /*Knuth's trick: getting
        bounds on m*/
        int mright=mid[l+1][r];
        res[l][r]=inf;
        for(int m=mleft; m<=mright; m++){ /*iterating for m in
        the bounds only*/
            int64 tres=res[l][m]+res[m][r]+(x[r]-x[l]);
            if(res[l][r]>tres){ //relax current solution
                res[l][r]=tres;
                mid[l][r]=m;
            }
        }
    }
}
int64 answer=res[0][k];

```

2.4 LIS O(nlogn) with full path [17 lines]

e7e81f

```

int num[MX], mem[MX], prev[MX], array[MX], res[MX], maxlen;
void LIS(int SZ, int num[]){
    CLR(mem), CLR(prev), CLR(array), CLR(res);
    int i, k;
    maxlen=1;
    array[0]=-inf;
    RFOR(i, 1, SZ+1) array[i]=inf;
    prev[0]=-1, mem[0]=num[0];
    FOR(i, SZ){
        k=lower_bound(array, array+maxlen+1, num[i])-array;
        if(k==1) array[k]=num[i], mem[k]=i, prev[i]=-1;
        else array[k]=num[i], mem[k]=i, prev[i]=mem[k-1];
        if(k>maxlen) maxlen=k;
    }
    k=0;
    for(i=mem[maxlen]; i!=-1; i=prev[i]) res[k++]=num[i];
}

```

2.5 SOS DP [18 lines]

5063f0

/*iterative version

```

for(int mask = 0; mask < (1<<N); ++mask){
    dp[mask][ -1] = A[mask]; //handle base case separately
    (leaf states)
    for(int i = 0; i < N; ++i){
        if(mask & (1<<i))
            dp[mask][i] = dp[mask][i-1] +
            dp[mask^(1<<i)][i-1];
    }
}

```

```

else
    dp[mask][i] = dp[mask][i-1];
}
F[mask] = dp[mask][N-1];
}
//memory optimized, super easy to code.
for(int i = 0; i < (1<<N); ++i)
    F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask <
    (1<<N); ++mask){
    if(mask & (1<<i))
        F[mask] += F[mask^(1<<i)];
}

```

2.6 Sibling DP [26 lines] cfc5ff

```

/*dividing tree into min group such that each group
cost not exceed k*/
ll n,k,dp[mx][mx];
vector<pair<ll,ll>>adj[mx];///must be rooted tree
ll sibling_dp(ll par,ll idx,ll remk){
    if(remk<0)return inf;
    if(adj[par].size()<idx+1)return 0;
    ll u=adj[par][idx].first;
    if(dp[u][remk]!=-1)
        return dp[u][remk];
    ll ret=inf,under=0,sibling=0;
    if(par!=0){//creating new group
        under=1+dfs(u,0,k);
        sibling=dfs(par,idx+1,remk);
        ret=min(ret,under+sibling);
    }
    //divide the current group
    ll temp=remk-adj[par][idx].second;
    for(ll chk=temp;chk>=0;chk--){
        ll siblingk=temp-chk;
        under=0,sibling=0;
        under=dfs(u,0,chk);
        sibling=dfs(par,idx+1,siblingk);
        ret=min(ret,under+sibling);
    }
    return dp[u][remk]=ret;
}

```

3 Flow

3.1 Blossom [58 lines] 1b2a6f

```

// Finds Maximum matching in General Graph
// Complexity O(NM)
// mate[i] = j means i is paired with j
// source: https://codeforces.com/blog/entry
// 92339?#comment=810242
vector<int> Blossom(vector<vector<int>>& graph) {
    //mate contains matched edge.
    int n = graph.size(), timer = -1;
    vector<int> mate(n, -1), label(n), parent(n),
        orig(n), aux(n, -1), q;
    auto lca = [&](int x, int y) {
        for (timer++; ; swap(x, y)) {
            if (x == -1) continue;
            if (aux[x] == timer) return x;
            aux[x] = timer;
            x = (mate[x] == -1 ? -1 : orig[parent[mate[x]]]);
        }
    };
    auto blossom = [&](int v, int w, int a) {

```

```

        while (orig[v] != a) {
            parent[v] = w; w = mate[v];
            if (label[w] == 1) label[w] = 0, q.push_back(w);
            orig[v] = orig[w] = a; v = parent[w];
        }
    };
    auto augment = [&](int v) {
        while (v != -1) {
            int pv = parent[v], nv = mate[pv];
            mate[v] = pv; mate[pv] = v; v = nv;
        }
    };
    auto bfs = [&](int root) {
        fill(label.begin(), label.end(), -1);
        iota(orig.begin(), orig.end(), 0);
        q.clear();
        label[root] = 0; q.push_back(root);
        for (int i = 0; i < (int)q.size(); ++i) {
            int v = q[i];
            for (auto x : graph[v]) {
                if (label[x] == -1) {
                    label[x] = 1; parent[x] = v;
                    if (mate[x] == -1)
                        return augment(x), 1;
                    label[mate[x]] = 0; q.push_back(mate[x]);
                }
                else if (label[x] == 0 && orig[v] != orig[x]) {
                    int a = lca(orig[v], orig[x]);
                    blossom(x, v, a); blossom(v, x, a);
                }
            }
        }
        return 0;
    };
    // Time halves if you start with (any) maximal
    // matching.
    for (int i = 0; i < n; i++)
        if (mate[i] == -1)
            bfs(i);
    return mate;
}

```

3.2 Dinic [72 lines] a786f1

```

/*Complexity: O(V^2 E)
.Call Dinic with total number of nodes.
.Nodes start from 0.
.Capacity is long long data.
.make graph with create edge(u,v,capacity).
.Get max flow with maxFlow(src,des).*/
#define eb emplace_back
struct Dinic {
    struct Edge {
        int u, v;
        ll cap, flow = 0;
        Edge() {}
        Edge(int u, int v, ll cap) :u(u), v(v), cap(cap) {}
    };
    int N;
    vector<Edge> edge;
    vector<vector<int>> adj;
    vector<int> d, pt;
    Dinic(int N) :N(N), edge(0), adj(N), d(N), pt(N) {}
    void addEdge(int u, int v, ll cap) {
        if (u == v) return;
        edge.emplace_back(u, v, cap);

```

```

        adj[u].emplace_back(edge.size() - 1);
        edge.emplace_back(v, u, 0);
        adj[v].emplace_back(edge.size() - 1);
    }
    bool bfs(int s, int t) {
        queue<int> q({s});
        fill(d.begin(), d.end(), N + 1);
        d[s] = 0;
        while (!q.empty()) {
            int u = q.front(); q.pop();
            if (u == t) break;
            for (int k : adj[u]) {
                Edge& e = edge[k];
                if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
                    d[e.v] = d[e.u] + 1;
                    q.emplace(e.v);
                }
            }
        }
        return d[t] != N + 1;
    }
    ll dfs(int u, int T, ll flow = -1) {
        if (u == T || flow == 0) return flow;
        for (int& i = pt[u]; i < adj[u].size(); i++) {
            Edge& e = edge[adj[u][i]];
            Edge& oe = edge[adj[u][i] ^ 1];
            if (d[e.v] == d[e.u] + 1) {
                ll amt = e.cap - e.flow;
                if (flow != -1 && amt > flow) amt = flow;
                if (ll pushed = dfs(e.v, T, amt)) {
                    e.flow += pushed;
                    oe.flow -= pushed;
                    return pushed;
                }
            }
        }
        return 0;
    }
    ll maxFlow(int s, int t) {
        ll total = 0;
        while (bfs(s, t)) {
            fill(pt.begin(), pt.end(), 0);
            while (ll flow = dfs(s, t)) {
                total += flow;
            }
        }
        return total;
    }
};

```

3.3 Flow [6 lines] 6ebca7

Covering Problems:

- > Maximum Independent Set(Bipartite): Largest set of nodes which do not have any edge between them. sol: V-(MaxMatching)
- > Minimum Vertex Cover(Bipartite): -Smallest set of nodes to cover all the edges -sol: MaxMatching
- > Minimum Edge Cover(General graph): -Smallest set of edges to cover all the nodes -sol: V-(MaxMatching) (if edge cover exists, does not exit for isolated nodes)

- > Minimum Path Cover(Vertex disjoint) DAG: -Minimum number of vertex disjoint paths that visit all the nodes -sol: make a bipartite graph using same nodes in two sides, one side is "from" other is "to", add edges from "from" to "to", then ans is V-(MaxMatching)
- > Minimum Path Cover(Vertex Not Disjoint) General graph: -Minimum number of paths that visit all the nodes -sol: consider cycles as nodes then it will become a path cover problem with vertex disjoint on DAG

3.4 HopcroftKarp [67 lines] fac9fc

/ Finds Maximum Matching In a bipartite graph*

.Complexity $O(E\sqrt{V})$

.1-indexed

.No default constructor

.add single edge for (u, v)/*

```
struct HK {
    static const int inf = 1e9;
    int n;
    vector<int> matchL, matchR, dist;
    //matchL contains value of matched node for L part.
    vector<vector<int>> adj;
    HK(int n) : n(n), matchL(n + 1),
    matchR(n + 1), dist(n + 1), adj(n + 1) {}
}
```

```
void addEdge(int u, int v) {
    adj[u].push_back(v);
}
```

```
bool bfs() {
    queue<int> q;
    for (int u = 1; u <= n; u++) {
        if (!matchL[u]) {
            dist[u] = 0;
            q.push(u);
        }
        else dist[u] = inf;
    }
    dist[0] = inf; //unmatched node matches with 0.
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        for (auto v : adj[u]) {
            if (dist[matchR[v]] == inf) {
                dist[matchR[v]] = dist[u] + 1;
                q.push(matchR[v]);
            }
        }
    }
    return dist[0] != inf;
}
```

```
bool dfs(int u) {
    if (!u) return true;
    for (auto v : adj[u]) {
        if (dist[matchR[v]] == dist[u] + 1
            && dfs(matchR[v])) {
            matchL[u] = v;
            matchR[v] = u;
            return true;
        }
    }
    dist[u] = inf;
    return false;
}
```

```
}
int max_match() {
    int matching = 0;
    while (bfs()) {
        for (int u = 1; u <= n; u++) {
            if (!matchL[u])
                if (dfs(u))
                    matching++;
        }
    }
    return matching;
}
};
```

3.5 Hungarian [116 lines] 64902f

/ Complexity: $O(n^3)$ but optimized*

It finds minimum cost maximum matching.

For finding maximum cost maximum matching

add -cost and return -matching()

*1-indexed */*

```
struct Hungarian {
    long long c[N][N], fx[N], fy[N], d[N];
    int l[N], r[N], arg[N], trace[N];
    queue<int> q;
    int start, finish, n;
    const long long inf = 1e18;
    Hungarian() {}
    Hungarian(int n1, int n2) : n(max(n1, n2)) {
        for (int i = 1; i <= n; ++i) {
            fy[i] = l[i] = r[i] = 0;
            for (int j = 1; j <= n; ++j) c[i][j] = inf;
        }
    }
    void add_edge(int u, int v, long long cost) {
        c[u][v] = min(c[u][v], cost);
    }
    inline long long getC(int u, int v) {
        return c[u][v] - fx[u] - fy[v];
    }
    void initBFS() {
        while (!q.empty()) q.pop();
        q.push(start);
        for (int i = 0; i <= n; ++i) trace[i] = 0;
        for (int v = 1; v <= n; ++v) {
            d[v] = getC(start, v);
            arg[v] = start;
        }
        finish = 0;
    }
    void findAugPath() {
        while (!q.empty()) {
            int u = q.front();
            q.pop();
            for (int v = 1; v <= n; ++v) if (!trace[v]) {
                long long w = getC(u, v);
                if (!w) {
                    trace[v] = u;
                    if (!r[v]) {
                        finish = v;
                        return;
                    }
                    q.push(r[v]);
                }
                if (d[v] > w) {

```

```
                    d[v] = w;
                    arg[v] = u;
                }
            }
        }
    }
    void subX_addY() {
        long long delta = inf;
        for (int v = 1; v <= n; ++v) if (trace[v] == 0 &&
            d[v] < delta) {
            delta = d[v];
        }
        // Rotate
        fx[start] += delta;
        for (int v = 1; v <= n; ++v) if (trace[v]) {
            int u = r[v];
            fy[v] -= delta;
            fx[u] += delta;
        }
        else d[v] -= delta;
        for (int v = 1; v <= n; ++v) if (!trace[v] && !d[v]) {
            trace[v] = arg[v];
            if (!r[v]) {
                finish = v;
                return;
            }
            q.push(r[v]);
        }
    }
    void Enlarge() {
        do {
            int u = trace[finish];
            int nxt = l[u];
            l[u] = finish;
            r[finish] = u;
            finish = nxt;
        } while (finish);
    }
    long long maximum_matching() {
        for (int u = 1; u <= n; ++u) {
            fx[u] = c[u][1];
            for (int v = 1; v <= n; ++v) {
                fx[u] = min(fx[u], c[u][v]);
            }
        }
        for (int v = 1; v <= n; ++v) {
            fy[v] = c[1][v] - fx[1];
            for (int u = 1; u <= n; ++u) {
                fy[v] = min(fy[v], c[u][v] - fx[u]);
            }
        }
        for (int u = 1; u <= n; ++u) {
            start = u;
            initBFS();
            while (!finish) {
                findAugPath();
                if (!finish) subX_addY();
            }
            Enlarge();
        }
        long long ans = 0;
    }
}
```

```

for (int i = 1; i <= n; ++i) {
    if (c[i][l[i]] != inf) ans += c[i][l[i]];
    else l[i] = 0;
}
return ans;
}
};

3.6 MCMF [116 lines] 466389
/*Credit: ShahjalalShohag
.Works for both directed, undirected and with negative
cost too
.doesn't work for negative cycles
.for undirected edges just make the directed flag
false
.Complexity: O(min(E^2 * V log V, E log V * flow))*/
using T = long long;
const T inf = 1LL << 61;
struct MCMF {
    struct edge {
        int u, v;
        T cap, cost;
        int id;
        edge(int _u, int _v, T _cap, T _cost, int _id) {
            u = _u;
            v = _v;
            cap = _cap;
            cost = _cost;
            id = _id;
        }
    };
    int n, s, t, mxid;
    T flow, cost;
    vector<vector<int>> g;
    vector<edge> e;
    vector<T> d, potential, flow_through;
    vector<int> par;
    bool neg;
    MCMF() {}
    MCMF(int _n) { // 0-based indexing
        n = _n + 10;
        g.assign(n, vector<int>());
        neg = false;
        mxid = 0;
    }
    void add_edge(int u, int v, T cap, T cost, int id = -1, bool directed = true) {
        if (cost < 0) neg = true;
        g[u].push_back(e.size());
        e.push_back(edge(u, v, cap, cost, id));
        g[v].push_back(e.size());
        e.push_back(edge(v, u, 0, -cost, -1));
        mxid = max(mxid, id);
        if (!directed) add_edge(v, u, cap, cost, -1, true);
    }
    bool dijkstra() {
        par.assign(n, -1);
        d.assign(n, inf);
        priority_queue<pair<T, T>, vector<pair<T, T>>, greater<pair<T, T>>> q;
        d[s] = 0;
        q.push(pair<T, T>(0, s));
        while (!q.empty()) {
            int u = q.top().second;
            T nw = q.top().first;

```

```

q.pop();
if (nw != d[u]) continue;
for (int i = 0; i < (int)g[u].size(); i++) {
    int id = g[u][i];
    int v = e[id].v;
    T cap = e[id].cap;
    T w = e[id].cost + potential[u] - potential[v];
    if (d[u] + w < d[v] && cap > 0) {
        d[v] = d[u] + w;
        par[v] = id;
        q.push(pair<T, T>(d[v], v));
    }
}
}
for (int i = 0; i < n; i++) { // update potential
    if (d[i] < inf) potential[i] += d[i];
}
return d[t] != inf;
}
T send_flow(int v, T cur) {
    if (par[v] == -1) return cur;
    int id = par[v];
    int u = e[id].u;
    T w = e[id].cost;
    T f = send_flow(u, min(cur, e[id].cap));
    cost += f * w;
    e[id].cap -= f;
    e[id ^ 1].cap += f;
    return f;
}
//returns {maxflow, mincost}
pair<T, T> solve(int _s, int _t, T goal = inf) {
    s = _s;
    t = _t;
    flow = 0, cost = 0;
    potential.assign(n, 0);
    if (neg) {
        // run Bellman-Ford to find starting potential
        d.assign(n, inf);
        for (int i = 0, relax = true; i < n && relax; i++) {
            for (int u = 0; u < n; u++) {
                for (int k = 0; k < (int)g[u].size(); k++) {
                    int id = g[u][k];
                    int v = e[id].v;
                    T cap = e[id].cap, w = e[id].cost;
                    if (d[u] > d[v] + w && cap > 0) {
                        d[v] = d[u] + w;
                        relax = true;
                    }
                }
            }
        }
        for (int i = 0; i < n; i++) if (d[i] < inf)
            potential[i] = d[i];
    }
    while (flow < goal && dijkstra()) flow +=
    send_flow(t, goal - flow);
    flow_through.assign(mxid + 10, 0);
    for (int u = 0; u < n; u++) {
        for (auto v : g[u]) {
            if (e[v].id >= 0) flow_through[e[v].id] = e[v ^
1].cap;
        }
    }
}

```

```

}
return make_pair(flow, cost);
}
};

```

4 Game Theory

4.1 Points to be noted [14 lines]

6fe124

>[First Write a Brute Force solution]

>Nim = all xor

>Misere Nim = Nim + corner case: if all piles are 1, reverse(nim)

>Bogus Nim = Nim

>Staircase Nim = Odd indexed pile Nim (Even indexed pile doesn't matter, as one player can give bogus moves to drop all even piles to ground)

>Sprague Grundy: [Every impartial game under the normal play convention is equivalent to a one-heap game of nim]

Every tree = one nim pile = tree root value; tree leaf value = 0; tree node value = mex of all child nodes.

[Careful: one tree node can become multiple new tree roots(multiple elements in one node), then the value of that node = xor of all those root values]

>Hackenbush(Given a rooted tree; cut an edge in one move; subtree under that edge gets removed; last player to cut wins):

Colon: $//G(u) = (G(v_1) + 1) \oplus (G(v_2) + 1) \oplus \dots [v_1, v_2, \dots \text{ are childs of } u]$

For multiple trees ans is their xor

>Hackenbush on graph (instead of tree given an rooted graph):

fusion: All edges in a cycle can be fused to get a tree structure; build a super node, connect some single nodes with that super node, number of single nodes is the number of edges in the cycle.

Sol: [Bridge component tree] mark all bridges, a group of edges that are not bridges, becomes one component and contributes number of edges to the hackenbush. (even number of edges contributes 0, odd number of edges contributes 1)

5 Geometry

5.1 Geometry [384 lines]

6bfd7b

```
namespace Geometry
```

```

{
    #define M_PI(acos(-1.0))
    double eps=1e-8;
    typedef double T; //coordinate point type
    struct pt //Point
    {
        T x,y;
        pt(){}
        pt(T _x,T _y):x(_x),y(_y){}
        pt operator+(pt p){
            return{x+p.x,y+p.y};
        }
        pt operator-(pt p){
            return{x-p.x,y-p.y};
        }
        pt operator*(T d){
            return{x*d,y*d};
        }
    }
}

```



```

pt operator*(pt d){/*I added for General linear
transformation,not sure about that function*/
    return{x*d.x,y*d.y};
}
pt operator/(T d){
    return{x/d,y/d};/*only for floating point*/
}
pt operator/(pt d){/*I added for General linear
transformation,not sure about that function*/
    return{x/d.x,y/d.y};
}
bool operator<(const pt& p)const {
    if(x!=p.x)
        return x<p.x;
    return y<p.y;
}
bool operator==(pt b){
    return x==b.x && y==b.y;
}
bool operator!=(pt b){
    return!(*(this)==b);
}
friend ostream& operator<<(ostream& os,const pt p){
    return os<<"("<p.x<<","<p.y<<")";
}
friend istream& operator>>(istream& is,pt &p){
    is>>p.x>>p.y;
    return is;
}
};
T sq(pt p){
    return p.x*p.x+p.y*p.y;
}
double Abs(pt p){
    return sqrtl(sq(p));
}
pt translate(pt v,pt p){ /*To translate an object by a
vector v*/
    return p+v;
}
pt scale(pt c,double factor,pt p){/*To scale an object
by a certain ratio factor around a center*/
    return c+(p-c)*factor;
}
pt rot(pt p,double a){/*To rotate a point by angle
a*/
    return{p.x*cos(a)-p.y*sin(a),p.x*sin(a)+p.y*
cos(a)};
}
pt perp(pt p){/*To rotate a point 90 degree*/
    return{-p.y,p.x};
}
pt linearTransfo(pt p,pt q,pt r,pt fp,pt fq){/*so far
don't know about that function*/
    return fp+(r-p)*(fq-fp)/(q-p);
}
T dot(pt v,pt w){
    return v.x*w.x+v.y*w.y;
}
bool isPerp(pt v,pt w){
    return dot(v,w)==0;
}
double angle(pt v,pt w){/*Find the smallest angle of
two vector*/

```

```

    double cosTheta=dot(v,w)/Abs(v)/Abs(w);
    return acos(max(-1.0,min(1.0,cosTheta)));
}
T cross(pt v,pt w){
    return v.x*w.y-v.y*w.x;
}
T orient(pt a,pt b,pt c){
    return cross(b-a,c-a); /*if c is left side+ve,c is
right side-ve,on line 0*/
}
bool inAngle(pt a,pt b,pt c,pt p){/*if p is in the
angle*/
    assert(orient(a,b,c)!=0);
    if(orient(a,b,c)<0)
        swap(b,c);
    return orient(a,b,p)>=0 && orient(a,c,p)<=0;
}
double orientedAngle(pt a,pt b,pt c){/*the actual
angle from ab to ac*/
    if(orient(a,b,c)>=0)
        return angle(b-a,c-a);
    else
        return 2*M_PI-angle(b-a,c-a);
}
//line
struct line{
    pt v;
    T c;
    line(){
        line(pt p,pt q){/*From points P and Q*/
            v=(q-p),this->c=cross(v,p);
        }
        line(T a,T b,T c){/*From equation ax+by=c*/
            v=pt(b,-a),this->c=c;
        }
        line(pt v,T c){/*From direction vector v and offset
c*/
            this->v=v,this->c=c;
        }
        double getY(double x){/*self made,not sure if it is
okay*/
            assert(v.x!=0);
            double ret=(double)(c+v.y*x)/v.x;
            return ret;
        }
        double getX(double y){/*self made,not sure if it is
okay*/
            assert(v.y!=0);
            double ret=(double)(c-v.x*y)/-v.y;
            return ret;
        }
        T side(pt p){/*which side a point is*/
            return cross(v,p)-c;
        }
        double dist(pt p){/*point to line dist*/
            return abs(side(p))/Abs(v);
        }
        double sqDist(pt p){/*square dist*/
            return side(p)*side(p)/(double)sq(v);
        }
        line perpThrough(pt p){/*perpendicular line with
point p*/
            return line(p,p+perp(v));
        }
    }

```

```

    bool cmpProj(pt p,pt q){/*compare function to sort
points on a line*/
        return dot(v,p)<dot(v,q);
    }
    line translate(pt t){/*translate with vector t*/
        return line(v,c+cross(v,t));
    }
    line shiftLeft(double dist){/*translate with
distance dist*/
        return line(v,c+dist*Abs(v));
    }
    pt proj(pt p){
        return p-perp(v)*side(p)/sq(v);
    }
    pt refl(pt p){
        return p-perp(v)*2*side(p)/sq(v);
    }
};
bool areParallel(line l1,line l2){
    return(l1.v.x*l2.v.y==l1.v.y*l2.v.x);
}
bool areSame(line l1,line l2){
    return areParallel(l1,l2)and(l1.v.x*l2.c==l2.v.x*
l1.c)and(l1.v.y*l2.c==l2.v.y*l1.c);
}
bool inter(line l1,line l2,pt& out){
    T d=cross(l1.v,l2.v);
    if(d==0)return false;
    out=(l2.v*l1.c-l1.v*l2.c)/d;
    return true;
}
line intBisector(line l1,line l2,bool interior){/*if
change sign then returns the other one*/
    assert(cross(l1.v,l2.v)!=0);
    double sign=interior?1:-1;
    return line(l2.v/Abs(l2.v)+l1.v*sign/Abs(l1.v),
l2.c/Abs(l2.v)+l1.c*sign/Abs(l1.v));
}
//segment
bool inDisk(pt a,pt b,pt p){/*check weather point p is
in diameter AB*/
    return dot(a-p,b-p)<=0;
}
bool onSegment(pt a,pt b,pt p){/*check weather point p
is in segment AB*/
    return orient(a,b,p)==0 and inDisk(a,b,p);
}
bool properInter(pt a,pt b,pt c,pt d,pt& i){
    double oa=orient(c,d,a),
ob=orient(c,d,b),
oc=orient(a,b,c),
od=orient(a,b,d);
//Proper intersection exists iff opposite signs
    if(oa*ob<0 and oc*od<0){
        i=(a*ob-b*oa)/(ob-oa);
        return 1;
    }
    return 0;
}
/*To create sets of points we need a comparison
function*/
struct cmpX{
    bool operator()(pt a,pt b){
        return make_pair(a.x,a.y)<make_pair(b.x,b.y);
    }

```

```

    }
};
set<pt,cmpX>inters(pt a,pt b,pt c,pt d){
    pt out;
    if(properInter(a,b,c,d,out))
        return{out};
    set<pt,cmpX>s;
    if(onSegment(c,d,a))s.insert(a);
    if(onSegment(c,d,b))s.insert(b);
    if(onSegment(a,b,c))s.insert(c);
    if(onSegment(a,b,d))s.insert(d);
    return s;
}
bool LineSegInter(line l,pt a,pt b,pt& out){
    if(l.side(a)*l.side(b)>eps)return 0;
    return inter(l,line(a,b),out);
}
double segPoint(pt a,pt b,pt p){/*returns distance
    from a point p to segment AB*/
    if(a!=b){
        line l(a,b);
        if(l.cmpProj(a,p)and l.cmpProj(p,b))
            return l.dist(p);
    }
    return min(Abs(p-a),Abs(p-b));
}
double segSeg(pt a,pt b,pt c,pt d){/*returns distance
    from a segment AB to segment CD*/
    pt dummy;
    if(properInter(a,b,c,d,dummy))return 0;
    return min(min(segPoint(a,b,c),segPoint(a,b,
        d)),segPoint(c,d,a)),segPoint(c,d,b));
}
/*int latticePoints(pt a,pt b){
    // requires int representation
    return __gcd(abs(a.x-b.x),abs(a.y-b.y))+1;
} // A = i+(b/2)-1; here
A=area,i=pointsinside,b=pointsonline
Polygon*/
bool isConvex(vector<pt>&p){
    bool hasPos=0,hasNeg=0;
    for(int i=0,n=p.size();i<n;i++){
        int o=orient(p[i],p[(i+1)%n],p[(i+2)%n]);
        if(o>0)hasPos=1;
        if(o<0)hasNeg=true;
    }
    return!(hasPos and hasNeg);
}
double areaTriangle(pt a,pt b,pt c){
    return abs(cross(b-a,c-a))/2.0;
}
double areaPolygon(const vector<pt>&p){
    double area=0.0;
    for(int i=0,n=p.size();i<n;i++){
        area+=cross(p[i],p[(i+1)%n]);
    }
    return fabs(area)/2.0;
}
bool pointInPolygon(const vector<pt>&p,pt q){/*returns
    true if pt q is in polygon p*/
    bool c=false;
    for(int i=0,n=p.size();i<n;i++){
        int j=(i+1)%p.size();

```

```

        if((p[i].y<=q.y and q.y<p[j].y or p[j].y<=q.y and
            q.y<p[i].y)and
            q.x<p[i].x+(p[j].x-p[i].x)*(q.y-p[i].y)/
            (p[j].y-p[i].y))
            c=!c;
        }
        return c;
    }
}
ll is_point_in_convex(vector<pt>& p, pt &x) { // O(log
    n)
    ll n = p.size(); /*this function from
    YouKnowWho*/
    if (n < 3) return 1;
    ll a = orient(p[0], p[1], x), b = orient(p[0], p[n
    - 1], x);
    if (a < 0 || b > 0) return 1;
    ll l = 1, r = n - 1;
    while (l + 1 < r) {
        int mid = l + r >> 1;
        if (orient(p[0], p[mid], x) >= 0) l = mid;
        else r = mid;
    }
    ll k = orient(p[l], p[r], x);
    if (k <= 0) return -k;
    if (l == 1 && a == 0) return 0;
    if (r == n - 1 && b == 0) return 0;
    return -1;
}
pt centroidPolygon(vector<pt>&p){/*from rezaul,i don't
    know about that*/
    pt c(0,0);
    double scale=6.0*areaPolygon(p);
    // if(scale<eps)return c;
    for(int i=0,n=p.size();i<n;i++){
        int j=(i+1)%n;
        c=c+(p[i]+p[j])*cross(p[i],p[j]);
    }
    return c/scale;
}
///Circle
pt circumCenter(pt a,pt b,pt c){/*return the center of
    the circle go through point a,b,c*/
    b=b-a,c=c-a;
    assert(cross(b,c)!=0);
    return a+perp(b*sq(c)-c*sq(b))/cross(b,c)/2;
}
bool circle2PtsRad(pt p1,pt p2,double r,pt& c){
    double d2=sq(p1-p2);
    double det=r*r/d2-0.25;
    if(det<0.0)return false;
    double h=sqrt(det);
    c.x=(p1.x+p2.x)*0.5+(p1.y-p2.y)*h;
    c.y=(p1.y+p2.y)*0.5+(p2.x-p1.x)*h;
    return true;
}
int circleLine(pt c,double r,line l,pair<pt,pt>&
    out){/*circle line intersection*/
    double h2=r*r-l.sqDist(c);
    if(h2<0)return 0;/*the line doesn't touch the
    circle;*/
    pt p=l.proj(c);
    pt h=l.v*sqrt(h2)/Abs(l.v);
    out=make_pair(p-h,p+h);
    return 1+(h2>0);
}

```

```

}
int circleCircle(pt c1,double r1,pt c2,double
    r2,pair<pt,pt>& out){/*circle circle intersection*/
    pt d=c2-c1;
    double d2=sq(d);
    if(d2==0){/*concentric circles
        assert(r1!=r2);
        return 0;
    }
    double pd=(d2+r1*r1-r2*r2)/2;
    double h2=r1*r1-pd*pd/d2;//h ~ 2
    if(h2<0)return 0;
    pt p=c1+d*pd/d2,h=perp(d)*sqrt(h2/d2);
    out=make_pair(p-h,p+h);
    return 1+h2>0;
}
int tangents(pt c1,double r1,pt c2,double r2,bool
    inner,vector<pair<pt,pt>>&out){
    if(inner)r2=-r2;/*returns tangent(the line which
    touch a circle in one point)of two circle*/
    pt d=c2-c1;/*the same code can be used to find the
    tangent to a circle passing through a point by
    setting r2 to 0*/
    double dr=r1-r2,d2=sq(d),h2=d2-dr*dr;
    if(d2==0 or h2<0){
        assert(h2!=0);
        return 0;
    }
    for(int sign :{-1,1}){
        pt v=pt(d*dr+perp(d)*sqrt(h2)*sign)/d2;
        out.push_back(make_pair(c1+v*r1,c2+v*r2));
    }
    return 1+(h2>0);
}
//Convex Hull-Monotone Chain
pt H[100000+5];
vector<pt>monotoneChain(vector<pt>&points){
    sort(points.begin(),points.end());
    vector<pt>ret;
    ret.clear();
    int st=0;
    for(int i=0,sz=points.size();i<sz;i++){
        while(st>=2 and
            orient(H[st-2],H[st-1],points[i])<0)st--;
        H[st++]=points[i];
    }
    int taken=st-1;
    for(int i=points.size()-2;i>=0;i--){
        while(st>=taken+2 and
            orient(H[st-2],H[st-1],points[i])<0)st--;
        H[st++]=points[i];
    }
    for(int i=0;i<st;i++)ret.push_back(H[i]);
    return ret;
}
//Convex Hull-Monotone Chain from you_know_who
int sz;
vector<pt> monotoneChain(vector<pt> &v) {
    if(v.size()==1) return v;
    sort(v.begin(), v.end());
    vector<pt> up(2*v.size()+2), down(2*v.size()+2);
    int szup=0, szdw=0;
    for(int i=0;i<v.size();i++) {

```

```

    while(szup>1 && orient(up[szup-2],
up[szup-1], v[i])>=0)
        szup--;
    while(szdw>1 && orient(down[szdw-2],
down[szdw-1], v[i])<=0)
        szdw--;
    up[szup++]=v[i];
    down[szdw++]=v[i];
}
if(szdw>1) szdw--;
reverse(up.begin(), up.begin()+szup);
for(int i=0;i<szup-1;i++) down[szdw++]= up[i];
if(szdw==2 && down[0].x==down[1].x &&
down[0].y==down[1].y)
    szdw--;
sz = szdw;
return down;
}
double cosA(double a,double b,double c){
    double val=b*b+c*c-a*a;
    val/=(2*b*c);
    return acos(val);
}
double triangle(double a,double b,double c){
    double s=(a+b+c)/2;
    return sqrt(1*s*(s-a)*(s-b)*(s-c));
}
}
using namespace Geometry;
5.2 Rotation Matrix [39 lines] f97f03
struct { double x; double y; double z; } Point;
double rMat[4][4];
double inMat[4][1] = {0.0, 0.0, 0.0, 0.0};
double outMat[4][1] = {0.0, 0.0, 0.0, 0.0};
void mulMat() {
    for(int i = 0; i < 4; i++) {
        for(int j = 0; j < 1; j++){
            outMat[i][j] = 0;
            for(int k = 0; k < 4; k++)
                outMat[i][j] += rMat[i][k] * inMat[k][j];
        }
    }
}
void setMat(double ang, double u, double v, double w){
    double L = (u * u + v * v + w * w);
    ang = ang * PI / 180.0; /*converting to radian
value*/
    double u2 = u*u; double v2 = v*v; double w2 = w*w;
    rMat[0][0]=(u2+(v2+w2)*cos(ang))/L;
    rMat[0][1]=(u*v*(1-cos(ang))-w*sqrt(L)*sin(ang))/L;
    rMat[0][2]=(u*w*(1-cos(ang))+v*sqrt(L)*sin(ang))/L;
    rMat[0][3]=0.0;
    rMat[1][0]=(u*v*(1-cos(ang))+w*sqrt(L)*sin(ang))/L;
    rMat[1][1]=(v2+(u2+w2)*cos(ang))/L;
    rMat[1][2]=(v*w*(1-cos(ang))-u*sqrt(L)*sin(ang))/L;
    rMat[1][3]=0.0;
    rMat[2][0]=(u*w*(1-cos(ang))-v*sqrt(L)*sin(ang))/L;
    rMat[2][1]=(v*w*(1-cos(ang))+u*sqrt(L)*sin(ang))/L;
    rMat[2][2]=(w2 + (u2 + v2) * cos(ang)) / L;
    rMat[2][3]=0.0; rMat[3][0]=0.0; rMat[3][1]=0.0;
    rMat[3][2]=0.0; rMat[3][3]=1.0;
}
/*double ang;
double u, v, w; //points = the point to be rotated

```

```

Point point, rotated; //u,v,w=unit vector of line
inMat[0][0] = points.x; inMat[1][0] = points.y;
inMat[2][0] = points.z; inMat[3][0] = 1.0;
setMat(ang, u, v, w); mulMat();
rotated.x = outMat[0][0]; rotated.y = outMat[1][0];
rotated.z = outMat[2][0];*/

```

6 Graph

6.1 2SAT [92 lines]

5289ec

```

struct TwoSat {
    vector<bool>vis;
    vector<vector<int>>>adj, radj;
    vector<int>dfs_t, ord, par;
    int n, intime; //For n node there will be 2*n node in
SAT.
    void init(int N) {
        n = N;
        intime = 0;
        vis.assign(N * 2 + 1, false);
        adj.assign(N * 2 + 1, vector<int>());
        radj.assign(N * 2 + 1, vector<int>());
        dfs_t.resize(N * 2 + 1);
        ord.resize(N * 2 + 1);
        par.resize(N * 2 + 1);
    }
    inline int neg(int x) {
        return x <= n ? x + n : x - n;
    }
    inline void add_implication(int a, int b) {
        if (a < 0) a = n - a;
        if (b < 0) b = n - b;
        adj[a].push_back(b);
        radj[b].push_back(a);
    }
    inline void add_or(int a, int b) {
        add_implication(-a, b);
        add_implication(-b, a);
    }
    inline void add_xor(int a, int b) {
        add_or(a, b);
        add_or(-a, -b);
    }
    inline void add_and(int a, int b) {
        add_or(a, b);
        add_or(a, -b);
        add_or(-a, b);
    }
    inline void force_true(int x) {
        if (x < 0) x = n - x;
        add_implication(neg(x), x);
    }
    inline void add_xnor(int a, int b) {
        add_or(a, -b);
        add_or(-a, b);
    }
    inline void add_nand(int a, int b) {
        add_or(-a, -b);
    }
    inline void add_nor(int a, int b) {
        add_and(-a, -b);
    }
    inline void force_false(int x) {
        if (x < 0) x = n - x;
        add_implication(x, neg(x));
    }
}

```

```

}
inline void topsort(int u) {
    vis[u] = 1;
    for (int v : radj[u]) if (!vis[v]) topsort(v);
    dfs_t[u] = ++intime;
}
inline void dfs(int u, int p) {
    par[u] = p, vis[u] = 1;
    for (int v : adj[u]) if (!vis[v]) dfs(v, p);
}
void build() {
    int i, x;
    for (i = n * 2, intime = 0; i >= 1; i--) {
        if (!vis[i]) topsort(i);
        ord[dfs_t[i]] = i;
    }
    vis.assign(n * 2 + 1, 0);
    for (i = n * 2; i > 0; i--) {
        x = ord[i];
        if (!vis[x]) dfs(x, x);
    }
}
bool satisfy(vector<int>& ret) //ret contains the value
that are true if the graph is satisfiable.
{
    build();
    vis.assign(n * 2 + 1, 0);
    for (int i = 1; i <= n * 2; i++) {
        int x = ord[i];
        if (par[x] == par[neg(x)]) return 0;
        if (!vis[par[x]]) {
            vis[par[x]] = 1;
            vis[par[neg(x)]] = 0;
        }
    }
    for (int i = 1; i <= n; i++) if (vis[par[i]])
        ret.push_back(i);
    return 1;
}
};

```

6.2 BridgeTree [66 lines]

f8e197

```

int N, M, timer, compid;
vector<pair<int, int>> g[mx];
bool used[mx], isBridge[mx];
int comp[mx], tin[mx], minAncestor[mx];
vector<int> Tree[mx]; // Store 2-edge-connected
component tree.(Bridge tree).
void markBridge(int v, int p) {
    tin[v] = minAncestor[v] = ++timer;
    used[v] = 1;
    for (auto& e : g[v]) {
        int to, id;
        tie(to, id) = e;
        if (to == p) continue;
        if (used[to]) minAncestor[v] = min(minAncestor[v],
tin[to]);
        else {
            markBridge(to, v);
            minAncestor[v] = min(minAncestor[v],
minAncestor[to]);
            if (minAncestor[to] > tin[v]) isBridge[id] = true;
            // if (tin[u] <= minAncestor[v]) ap[u] = 1;
        }
    }
}

```

```

    }
}
void markComp(int v, int p) {
    used[v] = 1;
    comp[v] = compid;
    for (auto& e : g[v]) {
        int to, id;
        tie(to, id) = e;
        if (isBridge[id]) continue;
        if (used[to]) continue;
        markComp(to, v);
    }
}
vector<pair<int, int>> edges;
void addEdge(int from, int to, int id) {
    g[from].push_back({ to, id });
    g[to].push_back({ from, id });
    edges[id] = { from, to };
}
void initB() {
    for (int i = 0; i <= compid; ++i) Tree[i].clear();
    for (int i = 1; i <= N; ++i) used[i] = false;
    for (int i = 1; i <= M; ++i) isBridge[i] = false;
    timer = compid = 0;
}
void bridge_tree() {
    initB();
    markBridge(1, -1); //Assuming graph is connected.
    for (int i = 1; i <= N; ++i) used[i] = 0;
    for (int i = 1; i <= N; ++i) {
        if (!used[i]) {
            markComp(i, -1);
            ++compid;
        }
    }
    for (int i = 1; i <= M; ++i) {
        if (isBridge[i]) {
            int u, v;
            tie(u, v) = edges[i];
            // connect two componets using edge.
            Tree[comp[u]].push_back(comp[v]);
            Tree[comp[v]].push_back(comp[u]);
            int x = comp[u];
            int y = comp[v];
        }
    }
}

```

6.3 Centroid Decomposition [49 lines] 3fb5b1

```

ll n, subsize[mx];
vector<ll> adj[mx];
char ans[mx];
bool brk[mx];
void calculatesize(ll u, ll par) {
    subsize[u] = 1;
    for (ll i = 0; i < (ll) adj[u].size(); i++) {
        ll v = adj[u][i];
        if (v == par or brk[v] == true) continue;
        calculatesize(v, u);
        subsize[u] += subsize[v];
    }
}
ll getcentroid(ll u, ll par, ll n) {
    ll ret = u;

```

```

    for (ll i = 0; i < (ll) adj[u].size(); i++) {
        ll v = adj[u][i];
        if (v == par or brk[v] == true) continue;
        if (subsize[v] > (n/2)) {
            ret = getcentroid(v, u, n);
            break;
        }
    }
    return ret;
}
void decompose(ll u, char rank) {
    calculatesize(u, -1);
    ll c = getcentroid(u, -1, subsize[u]);
    brk[c] = true;
    ans[c] = rank;
    for (ll i = 0; i < (ll) adj[c].size(); i++) {
        ll v = adj[c][i];
        if (brk[v] == true) continue;
        decompose(v, rank+1);
    }
}
int main() {
    scanf("%lld", &n);
    for (ll i = 0; i < n-1; i++) {
        ll a, b;
        scanf("%lld %lld", &a, &b);
        adj[a].push_back(b);
        adj[b].push_back(a);
    }
    decompose(1, 'A');
    for (ll i = 1; i <= n; i++) {
        printf("%c", ans[i]);
    }
}

```

6.4 DSU on Tree [56 lines] 391fb6

```

int n;
//extra data you need
vector<int> adj[mxn];
vector<int> *dsu[mxn];
void call(int u, int p = -1) {
    sz[u] = 1;
    for (auto v : adj[u]) {
        if (v != p) {
            dep[v] = dep[u] + 1;
            call(v, u);
            sz[u] += sz[v];
        }
    }
}
void dfs(int u, int p = -1, int isb = 1) {
    int mx = -1, big = -1;
    for (auto v : adj[u]) {
        if (v != p && sz[v] > mx) {
            mx = sz[v];
            big = v;
        }
    }
    for (auto v : adj[u]) {
        if (v != p && v != big) {
            dfs(v, u, 0);
        }
    }
    if (big != -1) {
        dfs(big, u, 1);
    }
}

```

```

    dsu[u] = dsu[big];
}
else {
    dsu[u] = new vector<int>();
}
dsu[u] -> push_back(u);
//calculation
for (auto v : adj[u]) {
    if (v == p || v == big) continue;
    for (auto x : *dsu[v]) {
        dsu[u] -> push_back(x);
        //calculation
    }
}
//calculate ans for node u
if (isb == 0) {
    for (auto x : *dsu[u]) {
        //reverse calculation
    }
}
}
int main() {
    //input graph
    dep[1] = 1;
    call(1);
    dfs(1);
}

```

6.5 Heavy Light Decomposition [73 lines] d0e24f

```

/*Heavy Light Decomposition
Build Complexity O(n)
Query Complexity O(lg^2 n)
Call init() with number of nodes
It's probably for the best to not do "using namespace
hld"*/
namespace hld {
    //N is the maximum number of nodes
    *par, lev, size corresponds to
    parent, depth, subtree-size*/
    //head[u] is the starting node of the chain u is in
    //in[u] to out[u] keeps the subtree indices
    const int N = 100000 + 7;
    vector<int> g[N];
    int par[N], lev[N], head[N], size[N], in[N], out[N];
    int cur_pos, n;
    //returns the size of subtree rooted at u
    *maintains the child with the largest subtree at the
    front of g[u]*/
    //WARNING: Don't change anything here specially with
    size[] if Jon Snow
    int dfs(int u, int p) {
        size[u] = 1, par[u] = p;
        lev[u] = lev[p] + 1;
        for (auto &v : g[u]) {
            if (v == p) continue;
            size[u] += dfs(v, u);
            if (size[v] > size[g[u].front()]) {
                swap(v, g[u].front());
            }
        }
        return size[u];
    }
}
//decomposed the tree in an array

```


//note that there is no physical array here

```
void decompose(int u, int p){
    in[u]++;cur_pos;
    for(auto &v : g[u]){
        if(v==p)continue;
        head[v]=(v==g[u].front()? head[u] : v);
        decompose(v,u);
    }
    out[u]=cur_pos;
}
```

//initializes the structure with _n nodes

```
void init(int _n,int root=1){
    n=_n;
    cur_pos=0;
    dfs(root,0);
    head[root]=root;
    decompose(root,0);
}
```

//checks whether p is an ancestor of u

```
bool isances(int p,int u){
    return in[p]<=in[u]and out[u]<=out[p];
}
```

//Returns the maximum node value in the path u-v

```
ll query(int u,int v){
    ll ret=-INF;
    while(!isances(head[u],v)){
        ret=max(ret,seg.query(1,1,n,in[head[u]],in[u]));
        u=par[head[u]];
    }
    swap(u,v);
    while(!isances(head[u],v)){
        ret=max(ret,seg.query(1,1,n,in[head[u]],in[u]));
        u=par[head[u]];
    }
    if(in[v]<in[u])swap(u,v);
    ret=max(ret,seg.query(1,1,n,in[u],in[v]));
    return ret;
}
```

//Adds val to subtree of u

```
void update(int u,ll val){
    seg.update(1,1,n,in[u],out[u],val);
}
```

```
};
```

6.6 K'th Shortest path [40 lines]

9f3788

```
int m,n,deg[MM],source,sink,K,val[MM][12];
```

```
struct edge{
    int v,w;
}adj[MM][500];
struct info{
    int v,w,k;
    bool operator<(const info &b)const{
        return w>b.w;
    }
};
```

```
priority_queue<info,vector<info>>Q;
void kthBestShortestPath(){
```

```
    int i,j;
    info u,v;
    for(i=0;i<n;i++){
        for(j=0;j<K;j++)val[i][j]=inf;
        u.v=source,u.k=0,u.w=0;
        Q.push(u);
        while(!Q.empty()){
            u=Q.top();
```

```
Q.pop();
for(i=0;i<deg[u.v];i++){
    v.v=adj[u.v][i].v;
    int cost=adj[u.v][i].w+u.w;
    for(v.k=u.k;v.k<K;v.k++){
        if(cost==inf)break;
        if(val[v.v][v.k]>cost){
            swap(cost,val[v.v][v.k]);
            v.w=val[v.v][v.k];
            Q.push(v);
            break;
        }
    }
    for(v.k++;v.k<K;v.k++){
        if(cost==inf)break;
        if(val[v.v][v.k]>cost)swap(cost, val[v.v][v.k]);
    }
}
```

6.7 LCA [46 lines]

9de12b

```
const int Lg = 22;
vector<int>adj[mx];
int level[mx];
int dp[Lg][mx];
void dfs(int u) {
    for (int i = 1;i < Lg;i++){
        dp[i][u] = dp[i - 1][dp[i - 1][u]];
    }
    for (int v : adj[u]) {
        if (dp[0][u] == v)continue;
        level[v] = level[u] + 1;
        dp[0][v] = u;
        dfs(v);
    }
}
```

```
int lca(int u, int v) {
    if (level[v] < level[u])swap(u, v);
    int diff = level[v] - level[u];
    for (int i = 0;i < Lg;i++){
        if (diff & (1 << i))
            v = dp[i][v];
    }
    for (int i = Lg - 1;i >= 0;i--){
        if (dp[i][u] != dp[i][v])
            u = dp[i][u], v = dp[i][v];
    }
    return u == v ? u : dp[0][u];
}
```

```
int kth(int u, int k) {
    for (int i = Lg - 1;i >= 0;i--){
        if (k & (1 << i))
            u = dp[i][u];
    }
    return u;
}
```

//kth node from u to v. 0th is u.

```
int go(int u, int v, int k) {
    int l = lca(u, v);
    int d = level[u] + level[v] - (level[l] << 1);
    assert(k <= d);
    if (level[l] + k <= level[u]) return kth(u, k);
    k -= level[u] - level[l];
    return kth(v, level[v] - level[l] - k);
}
```

/*

LCA(u,v) with root r:
lca(u,v)~lca(u,r)~lca(v,r)

Distance between u,v:
level(u) + level(v) - 2*level(lca(u,v))
*/

6.8 SACK [50 lines]

249f90

```
int sz[maxn];
void getsz(int v,int p){
    sz[v]=1;
    for(auto u : g[v])
        if(u!=p){
            getsz(u,v);
            sz[v]+=sz[u];
        }
}
```

//SACK $O(n \log^2 n)$

```
map<int,int>*cnt[maxn];
void dfs(int v,int p){
    int mx=-1,bigChild=-1;
    for(auto u : g[v])
        if(u!=p){
            dfs(u,v);
            if(sz[u]>mx)mx=sz[u],bigChild=u;
        }
    if(bigChild!=-1)cnt[v]=cnt[bigChild];
    else cnt[v]=new map<int,int>();
    (*cnt[v])[col[v]]++;
    for(auto u : g[v])
        if(u!=p && u!=bigChild){
            for(auto x : *cnt[u])
                (*cnt[v])[x.first]+=x.second;
        }
}
```

//SACK- $O(n \log n)$

```
vector<int>*vec[maxn];
int cnt[maxn];
void dfs(int v,int p,bool keep){
    int mx=-1,bigChild=-1;
    for(auto u : g[v])
        if(u!=p&&sz[u]>mx)mx=sz[u],bigChild=u;
    for(auto u : g[v])
        if(u!=p && u!=bigChild)dfs(u,v,0);
    if(bigChild!=-1)
        dfs(bigChild,v,1),vec[v]=vec[bigChild];
    else vec[v]=new vector<int>();
    vec[v]->push_back(v);cnt[col[v]]++;
    for(auto u : g[v])
        if(u!=p && u!=bigChild)
            for(auto x : *vec[u]){
                cnt[col[x]]++;
                vec[v]->push_back(x);
            }
}
```

/*in this step*vec[v]contains all of the subtree of vertex v.*/

```
if(keep==0)
    for(auto u:*vec[v])cnt[col[u]]--;
```

```
}
```

6.9 SCC [43 lines]

4da431

/*components: number of SCC.

sz: size of each SCC.

comp: component number of each node.

Create reverse graph.

Run find_scc() to find SCC.

Might need to create condensation graph by
 create_condensed().
 Think about indeg/outdeg
 for multiple test cases- clear
 adj/radj/comp/vis/sz/topo/condensed.*/
 vector<int>adj[mx], radj[mx];

```
int comp[mx], vis[mx], sz[mx], components;
vector<int>topo;
void dfs(int u) {
    vis[u] = 1;
    for (int v : adj[u])
        if (!vis[v]) dfs(v);
    topo.push_back(u);
}
void dfs2(int u, int val) {
    comp[u] = val;
    sz[val]++;
    for (int v : radj[u])
        if (comp[v] == -1)
            dfs2(v, val);
}
void find_scc(int n) {
    memset(vis, 0, sizeof vis);
    memset(comp, -1, sizeof comp);
    for (int i = 1; i <= n; i++)
        if (!vis[i])
            dfs(i);
    reverse(topo.begin(), topo.end());
    for (int u : topo)
        if (comp[u] == -1)
            dfs2(u, ++components);
}
vector<int>condensed[mx];
void create_condensed(int n) {
    for (int i = 1; i <= n; i++)
        for (int v : adj[i])
            if (comp[i] != comp[v])
                condensed[comp[i]].push_back(comp[v]);
}
```

7 Math

7.1 Big Sum [13 lines] 8d9520

```
ll bigsum(ll a, ll b, ll m) {
    if (b == 0) return 0;
    ll sum; a %= m;
    if (b & 1) {
        sum = bigsum((a * a) % m, (b - 1) / 2, m);
        sum = (sum + (a * sum) % m) % m;
        sum = (1 + (a * sum) % m) % m;
    } else {
        sum = bigsum((a * a) % m, b / 2, m);
        sum = (sum + (a * sum) % m) % m;
    }
    return sum;
}
```

7.2 CRT [52 lines] 59a568

```
ll ext_gcd(ll A, ll B, ll* X, ll* Y) {
    ll x2, y2, x1, y1, x, y, r2, r1, q, r;
    x2 = 1; y2 = 0;
    x1 = 0; y1 = 1;
    for (r2 = A, r1 = B; r1 != 0; r2 = r1, r1 = r, x2 =
        x1, y2 = y1, x1 = x, y1 = y) {
        q = r2 / r1;
```

```
        r = r2 % r1;
        x = x2 - (q * x1);
        y = y2 - (q * y1);
    }
    *X = x2; *Y = y2;
    return r2;
}
/*-----BlackBox-----*/
class ChineseRemainderTheorem {
    typedef long long vlong;
    typedef pair<vlong, vlong> pll;
    /** CRT Equations stored as pairs of vector. See
        addEquation()*/
    vector<pll> equations;
    public:
    void clear() {
        equations.clear();
    }
    /** Add equation of the form  $x = r \pmod m$ */
    void addEquation(vlong r, vlong m) {
        equations.push_back({ r, m });
    }
    pll solve() {
        if (equations.size() == 0) return { -1, -1 }; /// No
            equations to solve
        vlong a1 = equations[0].first;
        vlong m1 = equations[0].second;
        a1 %= m1;
        /** Initially  $x = a_0 \pmod{m_0}$ */
        /** Merge the solution with remaining equations */
        for (int i = 1; i < equations.size(); i++) {
            vlong a2 = equations[i].first;
            vlong m2 = equations[i].second;
            vlong g = __gcd(m1, m2);
            if (a1 % g != a2 % g) return { -1, -1 }; ///
                Conflict in equations
            /** Merge the two equations*/
            vlong p, q;
            ext_gcd(m1 / g, m2 / g, &p, &q);
            vlong mod = m1 / g * m2;
            vlong x = ((__int128)a1 * (m2 / g) % mod * q % mod
                + (__int128)a2 * (m1 / g) % mod * p % mod) % mod;
            /** Merged equation*/
            a1 = x;
            if (a1 < 0) a1 += mod;
            m1 = mod;
        }
        return { a1, m1 };
    }
};
```

7.3 FFT [85 lines] 4ca8f0

```
template<typename float_t>
struct mycomplex {
    float_t x, y;
    mycomplex<float_t>(float_t _x = 0, float_t _y = 0) :
        x(_x), y(_y) {}
    float_t real() const { return x; }
    float_t imag() const { return y; }
    void real(float_t _x) { x = _x; }
    void imag(float_t _y) { y = _y; }
    mycomplex<float_t>& operator+=(const
        mycomplex<float_t> &other) { x += other.x; y +=
        other.y; return *this; }
```

```
mycomplex<float_t>& operator-=(const
    mycomplex<float_t> &other) { x -= other.x; y -=
    other.y; return *this; }
mycomplex<float_t> operator+(const mycomplex<float_t>
    &other) const { return mycomplex<float_t>(*this) +=
    other; }
mycomplex<float_t> operator-(const mycomplex<float_t>
    &other) const { return mycomplex<float_t>(*this) -=
    other; }
mycomplex<float_t> operator*(const mycomplex<float_t>
    &other) const {
    return {x * other.x - y * other.y, x * other.y +
    other.x * y};
}
mycomplex<float_t> operator*(float_t mult) const {
    return {x * mult, y * mult};
}
friend mycomplex<float_t> conj(const
    mycomplex<float_t> &c) {
    return {c.x, -c.y};
}
friend ostream& operator<<(ostream &stream, const
    mycomplex<float_t> &c) {
    return stream << '(' << c.x << ", " << c.y << ')';
}
};
using cd = mycomplex<double>;
void fft(vector<cd> &a, bool invert) {
    int n = a.size();
    for (int i = 1, j = 0; i < n; i++) {
        int bit = n >> 1;
        for (; j & bit; bit >>= 1)
            j ^= bit;
        j ^= bit;
        if (i < j)
            swap(a[i], a[j]);
    }
    for (int len = 2; len <= n; len <= 1) {
        double ang = 2 * PI / len * (invert ? -1 : 1);
        cd wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i += len) {
            cd w(1);
            for (int j = 0; j < len / 2; j++) {
                cd u = a[i+j], v = a[i+j+len/2] * w;
                a[i+j] = u + v;
                a[i+j+len/2] = u - v;
                w = w*wlen;
            }
        }
    }
    if (invert) {
        for (cd &x : a) {
            double z = n;
            x = x/z;
        }
        //  $x /= n$ ;
    }
}
void multiply (const vector<bool> &a, const
    vector<bool> &b, vector<bool> &res) { //change all
    the bool to your type needed
    vector<cd> fa (a.begin(), a.end()), fb (b.begin(),
    b.end());
```

```

size_t n = 1;
while (n < max (a.size(), b.size())) n <= 1;
n <= 1;
fa.resize (n), fb.resize (n);
fft (fa, false), fft (fb, false);
for (size_t i=0; i<n; ++i)
    fa[i] =fa[i] * fb[i];
fft (fa, true);
res.resize (n);
for (size_t i=0; i<n; ++i)
    res[i] = round(fa[i].real());
while(res.back()==0) res.pop_back();
}

void pow(const vector<bool> &a, vector<bool> &res, long
        long int k){
    vector<bool> po=a;
    res.resize(1);
    res[0] = 1;
    while(k){
        if(k&1){
            multiply(po, res, res);
        }
        multiply(po, po, po);
        k/=2;
    }
}

```

7.4 GaussElimination [39 lines] aa53e0

```

template<typename ld>
int gauss(vector<vector<ld>>& a, vector<ld>& ans) {
    const ld EPS = 1e-9;
    int n = a.size();///number of equations
    int m = a[0].size() - 1;///number of variables
    vector<int>where(m, -1);///indicates which row
    contains the solution
    int row, col;
    for (col = 0, row = 0; col < m && row < n; ++col) {
        int sel = row;///which row contains the maximum
        value/
        for (int i = row + 1; i < n; i++)
            if (abs(a[i][col]) > abs(a[sel][col]))
                sel = i;
        if (abs(a[sel][col]) < EPS) continue;///it's
        basically 0.
        a[sel].swap(a[row]);///taking the max row up
        where[col] = row;
        ld t = a[row][col];
        for (int i = col; i <= m; i++) a[row][i] /= t;
        for (int i = 0; i < n; i++) {
            if (i != row) {
                ld c = a[i][col];
                for (int j = col; j <= m; j++)
                    a[i][j] -= a[row][j] * c;
            }
        }
        row++;
    }
    ans.assign(m, 0);
    for (int i = 0; i < m; i++)
        if (where[i] != -1)
            ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i = 0; i < n; i++) {
        ld sum = 0;
        for (int j = 0; j < m; j++)
            sum += ans[j] * a[i][j];
    }
}

```

```

        if (abs(sum - a[i][m]) > EPS) ///L.H.S!=R.H.S
            ans.clear();///No solution
        }
        return row;
    }
}

```

7.5 GaussMod2 [44 lines] e8fae4

```

template<typename T>
struct Gauss {
    int bits = 60;
    vector<T>table;
    Gauss() {
        table = vector<T>(bits, 0);
    }
    ///call with constructor to define bit size.
    Gauss(int _bits) {
        bits = _bits;
        table = vector<T>(bits, 0);
    }
    int basis()///return rank/size of basis
    {
        int ans = 0;
        for (int i = 0; i < bits; i++)
            if (table[i])
                ans++;
        return ans;
    }
    bool can(T x)///can x be obtained from the basis
    {
        for (int i = bits - 1; i >= 0; i--) x = min(x, x ^
            table[i]);
        return x == 0;
    }
    void add(T x) {
        for (int i = bits - 1; i >= 0 && x; i--) {
            if (table[i] == 0) {
                table[i] = x;
                x = 0;
            }
            else x = min(x, x ^ table[i]);
        }
    }
    T getBest() {
        T x = 0;
        for (int i = bits - 1; i >= 0; i--)
            x = max(x, x ^ table[i]);
        return x;
    }
    void Merge(Gauss& other) {
        for (int i = bits - 1; i >= 0; i--)
            add(other.table[i]);
    }
};

```

7.6 Karatsuba Idea [5 lines] 6944e1

Three subproblems:
 $a = xH \ yH$
 $d = xL \ yL$
 $e = (xH + xL)(yH + yL) - a - d$
 Then $xy = a \cdot r_n + e \cdot r_n/2 + d$

7.7 Linear Diophantine [12 lines] 86858c

*/*x'=x+(k*B/g),y'=y-(k*A/g);infinite soln
 if A=B=0,C must equal 0 and any x,y is solution;
 if A/B=0,(x,y)=(C/A,k)|(k,C/B)*/*

```

bool LDE(int A,int B,int C,int*x,int*y){
    int g=gcd(A,B);
    if(C%g!=0)return false;
    int a=A/g,b=B/g,c=C/g;
    extended_gcd(a,b,x,y);///ax+by=1
    if(g<0)a*=-1;b*=-1;c*=-1;///Ensure gcd(a,b)=1
    *x*=c;*y*=c;///ax+by=c
    return true;///Solution Exists
}

```

7.8 Matrix [100 lines] a33f18

```

template<typename T>
struct Matrix {
    T MOD = 1e9 + 7;///change if necessary
    T add(T a, T b) const {
        T res = a + b;
        if (res >= MOD) return res - MOD;
        return res;
    }
    T sub(T a, T b) const {
        T res = a - b;
        if (res < 0) return res + MOD;
        return res;
    }
    T mul(T a, T b) const {
        T res = a * b;
        if (res >= MOD) return res % MOD;
        return res;
    }
    int R, C;
    vector<vector<T>>mat;
    Matrix(int _R = 0, int _C = 0) {
        R = _R, C = _C;
        mat.resize(R);
        for (auto& v : mat) v.assign(C, 0);
    }
    void print() {
        for (int i = 0; i < R; i++)
            for (int j = 0; j < C; j++)
                cout << mat[i][j] << " \n"[j == C - 1];
    }
    void createIdentity() {
        for (int i = 0; i < R; i++)
            for (int j = 0; j < C; j++)
                mat[i][j] = (i == j);
    }
    Matrix operator+(const Matrix& o) const {
        Matrix res(R, C);
        for (int i = 0; i < R; i++)
            for (int j = 0; j < C; j++)
                res[i][j] = add(mat[i][j] + o.mat[i][j]);
    }
    Matrix operator-(const Matrix& o) const {
        Matrix res(R, C);
        for (int i = 0; i < R; i++)
            for (int j = 0; j < C; j++)
                res[i][j] = sub(mat[i][j] + o.mat[i][j]);
    }
    Matrix operator*(const Matrix& o) const {
        Matrix res(R, o.C);
        for (int i = 0; i < R; i++)
            for (int j = 0; j < o.C; j++)
                for (int k = 0; k < C; k++)

```

```

        res.mat[i][j] = add(res.mat[i][j],
mul(mat[i][k], o.mat[k][j]));
return res;
}
Matrix pow(long long x) {
    Matrix res(R, C);
    res.createIdentity();
    Matrix<T> o = *this;
    while (x) {
        if (x & 1) res = res * o;
        o = o * o;
        x >>= 1;
    }
    return res;
}
Matrix inverse()//Only square matrix && non-zero
determinant
{
    Matrix res(R, R + R);
    for (int i = 0; i < R; i++) {
        for (int j = 0; j < R; j++)
            res.mat[i][j] = mat[i][j];
        res.mat[i][R + i] = 1;
    }
    for (int i = 0; i < R; i++) {
        //find row 'r' with highest value at [r][i]
        int tr = i;
        for (int j = i + 1; j < R; j++)
            if (abs(res.mat[j][i]) > abs(res.mat[tr][i]))
                tr = j;
        //swap the row
        res.mat[tr].swap(res.mat[i]);
        //make 1 at [i][i]
        T val = res.mat[i][i];
        for (int j = 0; j < R + R; j++) res.mat[i][j] /=
val;
        //eliminate [r][i] from every row except i.
        for (int j = 0; j < R; j++) {
            if (j == i) continue;
            for (int k = R + R - 1; k >= i; k--) {
                res.mat[j][k] -= res.mat[i][k] * res.mat[j][i]
/ res.mat[i][i];
            }
        }
    }
    Matrix ans(R, R);
    for (int i = 0; i < R; i++)
        for (int j = 0; j < R; j++)
            ans.mat[i][j] = res.mat[i][R + j];
    return ans;
}
};

```

7.9 Miller-Rabin-Pollard-Rho [68 lines] 3e3e5f

```

ll powmod(ll a, ll p, ll m) {//(a^p % m)
    ll result = 1;
    a %= m;
    while (p) {
        if (p & 1)
            result = (vll)result * a % m;
        a = (vll)a * a % m;
        p >>= 1;
    }
    return result;
}

```

```

bool check_composite(ll n, ll a, ll d, int s) {
    ll x = powmod(a, d, n);
    if (x == 1 || x == n - 1)
        return false;
    for (int r = 1; r < s; r++) {
        x = (vll)x * x % n;
        if (x == n - 1)
            return false;
    }
    return true;
}
bool MillerRabin(ll n) {
    if (n < 2) return false;
    int r = 0;
    ll d = n - 1;
    while ((d & 1) == 0) {
        d >>= 1;
        r++;
    }
    for (int a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
37}) {
        if (n == a) return true;
        if (check_composite(n, a, d, r))
            return false;
    }
    return true;
}
ll mult(ll a, ll b, ll mod) {
    return (vll)a * b % mod;
}
ll f(ll x, ll c, ll mod) {
    return (mult(x, x, mod) + c) % mod;
}
ll rho(ll n) {
    if (n % 2 == 0) return 2;
    ll x = myrand() % n + 1, y = x, c = myrand() % n + 1,
g = 1;
    while (g == 1) {
        x = f(x, c, n);
        y = f(y, c, n);
        y = f(y, c, n);
        g = __gcd(abs(x - y), n);
    }
    return g;
}
set<ll> prime;
void prime_factorization(ll n) {
    if (n == 1) return;
    if (MillerRabin(n)) {
        prime.insert(n);
        return;
    }
    ll x = n;
    while (x == n) x = rho(n);
    prime_factorization(x);
    prime_factorization(n / x);
}
//call prime_factorization(n) for prime factors.
//call MillerRabin(n) to check if prime.

```

7.10 Mod Inverse [5 lines]

772679

```

int modInv(int a, int m) {
    int x, y; //if g==1 Inverse doesn't exist
    int g = gcdExt(a, m, x, y);
}

```

```

        return (x % m + m) % m;
}

```

7.11 NTT [96 lines]

6faca3

```

ll power(ll a, ll p, ll mod) {
    if (p == 0) return 1;
    ll ans = power(a, p/2, mod);
    ans = (ans * ans) % mod;
    if (p % 2) ans = (ans * a) % mod;
    return ans;
}
int primitive_root(int p) {
    vector<int> factor;
    int phi = p - 1, n = phi;
    for (int i = 2; i * i <= n; i++) {
        if (n % i) continue;
        factor.push_back(i);
        while (n % i == 0) n /= i;
    }
    if (n > 1) factor.push_back(n);
    for (int res = 2; res <= p; res++) {
        bool ok = true;
        for (int i = 0; i < factor.size() && ok; i++)
            ok &= power(res, phi / factor[i], p) != 1;
        if (ok) return res;
    }
    return -1;
}
int nttdata(int mod, int &root, int &inv, int &pw) {
    int c = 0, n = mod - 1;
    while (n % 2 == 0) c++, n /= 2;
    pw = (mod - 1) / n;
    int g = primitive_root(mod);
    root = power(g, n, mod);
    inv = power(root, mod - 2, mod);
    return c;
}
const int M = 786433;
struct NTT {
    int N;
    vector<int> perm;
    int mod, root, inv, pw;
    NTT() {}
    NTT(int mod, int root, int inv, int pw) : mod(mod),
root(root), inv(inv), pw(pw) {}
    void precalculate() {
        perm.resize(N);
        perm[0] = 0;
        for (int k = 1; k < N; k <= 1) {
            for (int i = 0; i < k; i++) {
                perm[i] <= 1;
                perm[i + k] = 1 + perm[i];
            }
        }
    }
    void fft(vector<ll> &v, bool invert = false) {
        if (v.size() != perm.size()) {
            N = v.size();
            assert(N && (N & (N - 1)) == 0);
            precalculate();
        }
        for (int i = 0; i < N; i++)
            if (i < perm[i])
                swap(v[i], v[perm[i]]);
    }
}

```



```
for (int len = 2; len <= N; len <=<=1) {
    ll factor = invert ? inv: root;
    for (int i=len; i<pw; i<=<=1)
        factor = (factor * factor) % mod;
    for (int i=0; i<N; i+=len) {
        ll w = 1;
        for (int j=0; j<len/2; j++) {
            ll x = v[i+j], y = (w*v[i+j+len/2])%mod;
            v[i+j] = (x+y)%mod;
            v[i+j+len/2] = (x-y+mod)%mod;
            w = (w*factor)%mod;
        }
    }
    if (invert) {
        ll n1 = power(N, mod-2, mod);
        for (ll &x: v) x = (x*n1)%mod;
    }
}
vector<ll> multiply(vector<ll> a, vector<ll> &b) {
    while (a.size() && a.back() == 0) a.pop_back();
    while (b.size() && b.back() == 0) b.pop_back();
    int n = 1;
    while (n < a.size() + b.size()) n<=<=1;
    a.resize(n);
    b.resize(n);
    fft(a);
    fft(b);
    for (int i=0; i<n; i++) a[i] = (a[i] * b[i])%M;
    fft(a, true);
    while (a.size() && a.back() == 0) a.pop_back();
    return a;
}
// int mod=786433, root, inv, pw;
// nttdata(mod, root, inv, pw);
// NTT nn = NTT(mod, root, inv, pw);
};
```

7.12 No of Digits in n! in base B [7 lines] 86bfaf

```
ll NoOfDigitInNFactInBaseB(ll N,ll B){
    ll i;
    double ans=0;
    for(i=1;i<=N;i++)ans+=log(i);
    ans=ans/log(B),ans=ans+1;
    return (ll)ans;
}
```

7.13 SOD Upto N [16 lines] d8aa2c

```
ll SOD_Upto_N(ll N){
    ll i,j,ans=0;//upto N in Sqrt(N)
    for(i=1;i*i<=N;i++){
        j=N/i;
        ans+=((j*(j+1))/2)-(((i-1)*i)/2);
        ans+=((j-i)*i);
    }
    return ans;
}
ll SODUptoN(ll N){
    ll res=0,u=sqrt(N);
    for(ll i=1;i<=u;i++)
        res+=(N/i)-i;
    res*=2,res+=u;
    return res;
}
```

7.14 Sieve Phi Mobius [26 lines] 353c39

```
const int N = 1e7;
vector<int>pr;
int mu[N + 1], phi[N + 1], lp[N + 1];
void sieve() {
    phi[1] = 1, mu[1] = 1;
    for (int i = 2; i <= N; i++) {
        if (lp[i] == 0) {
            lp[i] = i;
            phi[i] = i - 1;
            pr.push_back(i);
        }
        for (int j = 0; j < pr.size() && i * pr[j] <= N; j++) {
            lp[i * pr[j]] = pr[j];
            if (i % pr[j] == 0) {
                phi[i * pr[j]] = phi[i] * pr[j];
                break;
            }
            else
                phi[i * pr[j]] = phi[i] * phi[pr[j]];
        }
    }
    for (int i = 2; i <= N; i++) {
        if (lp[i / lp[i]] == lp[i]) mu[i] = 0;
        else mu[i] = -1 * mu[i / lp[i]];
    }
}
```

8 Misc

8.1 Bit hacks [12 lines] dd22ef

x & -x is the least bit in x.
iterate over all the subsets of the mask
for (int s=m; ; s=(s-1)&m) {
 ... you can use s ...
 if (s==0) break;
}
c = x&-x, r = x+c; (((r^x) >> 2)/c) | r is the
next number after x with the same number of bits set.
__builtin_popcount(x) //number of ones in binary
__builtin_popcountll(x) // for long long
__builtin_clz(x) // number of leading zeros
__builtin_ctz(x) // number of trailing zeros, they
also have long long version

8.2 Bitset C++ [13 lines] a6a7a4

```
bitset<17>BS;
BS[1] = BS[7] = 1;
cout<<BS._Find_first()<<endl; // prints 1
bs._Find_next(idx) // this function returns first set bit
after index idx.for example:

bitset<17>BS;
BS[1] = BS[7] = 1;
cout<<BS._Find_next(1)<<','<<BS._Find_next(3)<<endl; //
prints 7,7
So this code will print all of the set bits of BS:

for(int i=BS._Find_first();i< BS.size();i =
BS._Find_next(i))
    cout<<i<<endl;
//Note that there isn't any set bit after idx,
BS._Find_next(idx) will return BS.size(); same as
calling BS._Find_first() when bitset is clear;
```

8.3 Template [33 lines] 7aea62

```
// #pragma GCC optimize("O3,unroll-loops")
// #pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;

template <typename A, typename B> ostream&
operator<<(ostream& os, const pair<A, B>& p) {
    return os << '(' << p.first << ", " << p.second <<
    ')'; }
template <typename T_container, typename T = typename
enable_if<!is_same<T_container, string>::value,
typename T_container::value_type>::type> ostream&
operator<<(ostream& os, const T_container& v) { os
<< '{'; string sep; for (const T& x : v) os << sep
<< x, sep = ", "; return os << '}'; }
void dbg_out() { cerr << endl; }
template <typename Head, typename... Tail> void
dbg_out(Head H, Tail... T) { cerr << " " << H;
dbg_out(T...); }

#ifdef SMIE
#define debug(args...) cerr << "(" << #args << "):",
dbg_out(args)
#else
#define debug(args...)
#endif

template <typename T> inline T gcd(T a, T b) { T c;while
(b) { c = b;b = a % b;a = c; }return a; } // better
than __gcd
ll powmod(ll a, ll b, ll MOD) { ll res = 1;a %=
MOD;assert(b >= 0);for (; b >>= 1) { if (b &
1)res = res * a % MOD;a = a * a % MOD; }return res;
}
template <typename T>using orderedSet = tree<T,
null_type, less_equal<T>, rb_tree_tag,
tree_order_statistics_node_update>;
//order_of_key(k) - number of element strictly less than
k
//find_by_order(k) - k'th element in set.(0
indexed)(iterator)
```

```
mt19937
rng(chrono::steady_clock::now().time_since_epoch()
.count());
//uniform_int_distribution<int>(0, i)(rng)
int main(int argc, char* argv[]) {
    ios_base::sync_with_stdio(false);//DON'T CC++
    cin.tie(NULL);//DON'T use for interactive
    int seed = atoi(argv[1]);
}
```

8.4 build [2 lines] 801989

```
#!/bin/bash
>&2 echo -e "Making [$2]\t: $1.cpp" && g++ -std=gnu++17
-Wshadow -Wall -Wextra -Wno-unused-result -O2 -g
-fsanitize=undefined -fsanitize=address $2 "$1.cpp"
-o "$1"
```

8.5 debug [3 lines] **859f78**

```
#!/bin/bash
build "$1" -DSMIE && >&2 echo -e "Running\t\t:
    $1\n-----" && "./$1"
```

8.6 stress [15 lines] **62e61a**

```
#!/bin/bash
build $1 $2 && build $1_gen $2 && build $1_brute $2 &&
for((i = 1; ; ++i)); do
    echo -e "\nTest Case "$i
    ./$1_gen $i > inp
    ./$1 < inp > out1
    ./$1_brute < inp > out2
    diff -w out1 out2 || break
done
echo -e "=====\nINPUT\n-----"
cat inp
echo -e "\nOUTPUT\n-----"
cat out1
echo -e "\nEXPECTED\n-----"
cat out2
```

8.7 vimrc [14 lines] **ffdf4e**

```
filetype plugin indent on
set rnu wfw hls is ar aw wrap mouse=a
```

```
let mapleader=' '
im jk <esc>
tno jk <c-w>N
no <leader>d "_d
im {<cr> {<cr>}<esc>O
nn ff :let @+ = expand("%:~p")<cr>
nn cd :cd %:~h<cr>
```

```
au BufNewFile *.cpp -r ./template.cpp | 14
```

```
ca hash w !cpp -dD -P -fpreprocessed \\\ tr -d
    '[:space:]' \\\ md5sum \\\ cut -c-6
```

9 String**9.1 Aho-Corasick [124 lines]** **2d8d6c**

```
const int NODE=3000500;///Maximum Nodes
const int LGN=30;      ///Maximum Number of Tries
const int MXCHR=53;    ///Maximum Characters
const int MXP=5005;    ///
struct node {
    int val;
    int child[MXCHR];
    vector<int>graph;
    void clear(){
        CLR(child,0);
        val=0;
        graph.clear();
    }
}Trie[NODE+10];
int maxNodeId,fail[NODE+10],par[NODE+10];
int nodeSt[NODE+10],nodeEd[NODE+10];
vlong csum[NODE+10],pLoc[MXP];
void resetTrie(){
    maxNodeId=0;
}
int getNode(){
    int curNodeId=++maxNodeId;
    Trie[curNodeId].clear();
```

```
    return curNodeId;
}
inline void upd(vlong pos){
    csum[pos]++;
}
inline vlong qry(vlong pos){
    vlong res=csum[pos];
    return res;
}
struct AhoCorasick {
    int root,size,euler;
    void clear(){
        root=getNode();
        size=euler=0;
    }
    inline int getname(char ch){
        if(ch=='-')return 52;
        else if(ch>='A' && ch<='Z')return 26+(ch-'A');
        else return(ch-'a');
    }
    void addToTrie(string &s,int id){
        ///Add string s to the Trie in general way
        int len=SZ(s),cur=root;
        FOR(i,0,len-1){
            int c=getname(s[i]);
            if(Trie[cur].child[c]==0){
                int curNodeId=getNode();
                Trie[curNodeId].val=c;
                Trie[cur].child[c]=curNodeId;
            }
            cur=Trie[cur].child[c];
        }
        pLoc[id]=cur;
        size++;
    }
    void calcFailFunction(){
        queue<int>Q;
        Q.push(root);
        while(!Q.empty()){
            int s=Q.front();
            Q.pop();
            ///Add all the children to the queue:
            FOR(i,0,MXCHR-1){
                int t=Trie[s].child[i];
                if(t!=0){
                    Q.push(t);
                    par[t]=s;
                }
            }
            if(s==root){/*Handle special case when s is
            root*/
                fail[s]=par[s]=root;
                continue;
            }
            ///Find fall back of s:
            int p=par[s],f=fail[p];
            int val=Trie[s].val;
            /*Fall back till you found a node who has got val as a
            child*/
            while(f!=root && Trie[f].child[val]==0){
                f=fail[f];
            }
            fail[s]=(Trie[f].child[val]==0)? root :
            Trie[f].child[val];
```

```
///Self fall back not allowed
        if(s==fail[s]){
            fail[s]=root;
        }
        Trie[fail[s]].graph.push_back(s);
    }
}
void dfs(int pos){
    ++euler;
    nodeSt[pos]=euler;
    for(auto x: Trie[pos].graph){
        dfs(x);
    }
    nodeEd[pos]=euler;
}
///Returns the next state
int goTo(int state,int c){
    if(Trie[state].child[c]!=0){/*No need to fall
    back*/
        return Trie[state].child[c];
    }
}
///Fall back now:
int f=fail[state];
while(f!=root && Trie[f].child[c]==0){
    f=fail[f];
}
int res=(Trie[f].child[c]==0)?
root:Trie[f].child[c];
return res;
}
/*Iterate through the whole text and find all the
matchings*/
void findmatching(string &s){
    int cur=root,idx=0;
    int len=SZ(s);
    while(idx<len){
        int c=getname(s[idx]);
        cur=goTo(cur,c);
        upd(nodeSt[cur]);
        idx++;
    }
}
}acorasick;
```

9.2 Double Hashing [50 lines] **1a70c1**

```
struct SimpleHash {
    int len;
    long long base, mod;
    vector<int> P, H, R;
    SimpleHash() {}
    SimpleHash(string str, long long b, long long m) {
        base = b, mod = m, len = str.size();
        P.resize(len + 4, 1), H.resize(len + 3, 0),
        R.resize(len + 3, 0);
        for (int i = 1; i <= len + 3; i++)
            P[i] = (P[i - 1] * base) % mod;
        for (int i = 1; i <= len; i++)
            H[i] = (H[i - 1] * base + str[i - 1] + 1007)
            % mod;
        for (int i = len; i >= 1; i--)
            R[i] = (R[i + 1] * base + str[i - 1] + 1007)
            % mod;
    }
}
```

```

inline int range_hash(int l, int r) {
    int hashval = H[r + 1] - ((long long)P[r - 1 +
1] * H[1] % mod);
    return (hashval < 0 ? hashval + mod : hashval);
}
inline int reverse_hash(int l, int r) {
    int hashval = R[l + 1] - ((long long)P[r - 1 +
1] * R[r + 2] % mod);
    return (hashval < 0 ? hashval + mod : hashval);
}
};
struct DoubleHash {
    SimpleHash sh1, sh2;
    DoubleHash() {}
    DoubleHash(string str) {
        sh1 = SimpleHash(str, 1949313259, 2091573227);
        sh2 = SimpleHash(str, 1997293877, 2117566807);
    }
    long long concate(DoubleHash& B, int l1, int r1,
int l2, int r2) {
        int len1 = r1 - l1 + 1, len2 = r2 - l2 + 1;
        long long x1 = sh1.range_hash(l1, r1),
x2 = B.sh1.range_hash(l2, r2);
        x1 = (x1 * B.sh1.P[len2]) % 2091573227;
        long long newx1 = (x1 + x2) % 2091573227;
        x1 = sh2.range_hash(l1, r1);
        x2 = B.sh2.range_hash(l2, r2);
        x1 = (x1 * B.sh2.P[len2]) % 2117566807;
        long long newx2 = (x1 + x2) % 2117566807;
        return (newx1 << 32) ^ newx2;
    }
    inline long long range_hash(int l, int r) {
        return ((long long)sh1.range_hash(l, r) << 32) ^
sh2.range_hash(l, r);
    }
    inline long long reverse_hash(int l, int r) {
        return ((long long)sh1.reverse_hash(l, r) << 32) ^
sh2.reverse_hash(l, r);
    }
};

```

9.3 KMP [23 lines]

99c570

```

char P[maxn], T[maxn];
int b[maxn], n, m;
void kmpPreprocess() {
    int i=0, j=-1;
    b[0]=-1;
    while(i<m) {
        while(j>=0 and P[i]!=P[j])
            j=b[j];
        j=b[j];
        i++; j++;
        b[i]=j;
    }
}
void kmpSearch() {
    int i=0, j=0;
    while(i<n) {
        while(j>=0 and T[i]!=P[j])
            j=b[j];
        j=b[j];
        i++; j++;
        if(j==m) {
            //pattern found at index i-j
        }
    }
}

```

9.4 Manacher [16 lines]

2b3cab

```

vector<int> manacher_odd(string s) {
    int n = s.size();
    s = "$" + s + "~";
    vector<int> p(n + 2);
    int l = 1, r = 1;
    for(int i = 1; i <= n; i++) {
        p[i] = max(0, min(r - i, p[l + (r - i)]));
        while(s[i - p[i]] == s[i + p[i]]) {
            p[i]++;
        }
        if(i + p[i] > r) {
            l = i - p[i], r = i + p[i];
        }
    }
    return vector<int>(begin(p) + 1, end(p) - 1);
}

```

9.5 Palindromic Tree [30 lines]

9ebc05

```

struct PalindromicTree {
    int n, idx, t;
    vector<vector<int>> tree;
    vector<int> len, link;
    string s; // 1-indexed
    PalindromicTree(string str) {
        s = "$" + str;
        n = s.size();
        len.assign(n + 5, 0);
        link.assign(n + 5, 0);
        tree.assign(n + 5, vector<int>(26, 0));
    }
    void extend(int p) {
        while(s[p - len[t] - 1] != s[p]) t = link[t];
        int x = link[t], c = s[p] - 'a';
        while(s[p - len[x] - 1] != s[p]) x = link[x];
        if(!tree[t][c]) {
            tree[t][c] = ++idx;
            len[idx] = len[t] + 2;
            link[idx] = len[idx] == 1 ? 2 : tree[x][c];
        }
        t = tree[t][c];
    }
    void build() {
        len[1] = -1, link[1] = 1;
        len[2] = 0, link[2] = 1;
        idx = 2;
        for(int i = 1; i < n; i++) extend(i);
    }
};

```

9.6 Suffix Array [78 lines]

0f4693

```

struct SuffixArray {
    vector<int> p, c, rank, lcp;
    vector<vector<int>> st;
    SuffixArray(string const& s) {
        build_suffix(s + char(1));
        p.erase(p.begin());
        build_rank(p.size());
        build_lcp(s); // doesn't work for
        build_sparse_table(lcp.size()); // single char
    }
    void build_suffix(string const& s) {
        int n = s.size();
        const int MX_ASCII = 256;
        vector<int> cnt(max(MX_ASCII, n), 0);
    }
}

```

```

p.resize(n); c.resize(n);
for (int i = 0; i < n; i++) cnt[s[i]]++;
for (int i = 1; i < MX_ASCII; i++) cnt[i] += cnt[i - 1];
for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;
c[p[0]] = 0;
int classes = 1;
for (int i = 1; i < n; i++) {
    if (s[p[i]] != s[p[i - 1]]) classes++;
    c[p[i]] = classes - 1;
}
vector<int> pn(n), cn(n);
for (int h = 0; (1 << h) < n; ++h) {
    for (int i = 0; i < n; i++) {
        pn[i] = p[i] - (1 << h);
        if (pn[i] < 0) pn[i] += n;
    }
    fill(cnt.begin(), cnt.begin() + classes, 0);
    for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;
    for (int i = 1; i < classes; i++) cnt[i] += cnt[i - 1];
    for (int i = n - 1; i >= 0; i--) p[--cnt[c[pn[i]]]] = pn[i];
    cn[p[0]] = 0; classes = 1;
    for (int i = 1; i < n; i++) {
        pair<int, int> cur = {c[p[i]], c[(p[i] + (1 << h)) % n]};
        pair<int, int> prev = {c[p[i - 1]], c[(p[i - 1] + (1 << h)) % n]};
        if (cur != prev) ++classes;
        cn[p[i]] = classes - 1;
    }
    c.swap(cn);
}
}
void build_rank(int n) {
    rank.resize(n, 0);
    for (int i = 0; i < n; i++) rank[p[i]] = i;
}
void build_lcp(string const& s) {
    int n = s.size(), k = 0;
    lcp.resize(n - 1, 0);
    for (int i = 0; i < n; i++) {
        if (rank[i] == n - 1) {
            k = 0;
            continue;
        }
        int j = p[rank[i] + 1];
        while (i + k < n && j + k < n && s[i + k] == s[j + k])
            k++;
        lcp[rank[i]] = k;
        if (k) k--;
    }
}
void build_sparse_table(int n) {
    int lim = __lg(n);
    st.resize(lim + 1, vector<int>(n)); st[0] = lcp;
    for (int k = 1; k <= lim; k++)
        for (int i = 0; i + (1 << k) <= n; i++)
            st[k][i] = min(st[k - 1][i], st[k - 1][i + (1 << (k - 1))]);
}
int get_lcp(int i) { return lcp[i]; }
int get_lcp(int i, int j) {
    if (j < i) swap(i, j);
}

```

```

j--; /*for lcp from i to j we don't need last lcp*/
int K = __lg(j - i + 1);
return min(st[K][i], st[K][j - (1 << K) + 1]);
}
};

9.7 Suffix Automata [118 lines] 179441
const int MXCHR = 26;
/*
take an object of suffixAutomata
call extend(c) for each character c in string
call Process() to initiate the important values
*/
struct suffixAutomata {
/*
len -> largest string length of the corresponding
endpos-equivalent class
link -> longest suffix that is another endpos-equivalent
class
firstpos -> end position of the first occurrence of the
largest string of that node
**/
struct state {
int link, len;
int next[MXCHR];
state() {}
state(int l) {
len = l;
link = -1;
for (int i = 0; i < MXCHR; i++) next[i] = -1;
}
};
vector<state> node;
int sz, last;
vector<int> cnt, distinct, firstPos, occur, SA;
vector<vector<int>> adj; // suffix links tree
// cnt and SA for counting

sort the nodes.
int L;

suffixAutomata() {
node.push_back(state(0));
firstPos.push_back(-1);
occur.push_back(0);
last = 0;
sz = 0;
L = 0;
}
int getID(char c) {
return c - 'a'; // change according to problem
}
void extend(char c) {
int idx = ++sz, p = last, id = getID(c);
L++;
node.push_back(state(node[last].len + 1));
firstPos.push_back(node[idx].len - 1);
occur.push_back(1);

while (p != -1 && node[p].next[id] == -1) {
node[p].next[id] = idx;
p = node[p].link;
}
if (p == -1)
node[idx].link = 0;
else {

```

```

int q = node[p].next[id];
if (node[p].len + 1 == node[q].len)
node[idx].link = q;
else {
int clone = ++sz;
state x = node[q];
x.len = node[p].len + 1;
node.push_back(x);
firstPos.push_back(firstPos[q]);
occur.push_back(0);
while (p != -1 && node[p].next[id] == q) {
node[p].next[id] = clone;
p = node[p].link;
}
node[idx].link = node[q].link = clone;
}
last = idx;
}
void Process() {
cnt.resize(sz + 1);
distinct.resize(sz + 1);
SA.resize(sz + 1);
adj.resize(sz + 1);
for (int i = 0; i <= sz; i++) cnt[node[i].len]++;
for (int i = 1; i <= L; i++) cnt[i] += cnt[i - 1];
for (int i = 0; i <= sz; i++) SA[--cnt[node[i].len]] = i;
for (int i = sz; i > 0; i--) {
int idx = SA[i];
occur[node[idx].link] += occur[idx];
adj[node[idx].link].push_back(idx);
distinct[idx] = 1;
for (int j = 0; j < MXCHR; j++) {
if (node[idx].next[j] != -1)
distinct[idx] += distinct[node[idx].next[j]];
}
} // counts distinct substrings and occurrence of
each state
for (int i = 0; i < MXCHR; i++)
if (node[0].next[i] != -1) distinct[0] +=
distinct[node[0].next[i]];
}
pair<int, int> lcs(string &str) {
int mxlen = 0, bestpos = -1, pos = 0, len = 0;
int u = 0; // LCS of two string. returns start
position and length
for (char c : str) {
int v = getID(c);
while (u && node[u].next[v] == -1) {
u = node[u].link;
len = node[u].len;
}
if (node[u].next[v] != -1) {
len++;
u = node[u].next[v];
}
if (len > mxlen) {
mxlen = len;
bestpos = pos;
}
pos++;
}
return {bestpos - mxlen + 1, mxlen};
}

```

```

}
state &operator[] (int index) { return node[index]; }
};

9.8 Trie [28 lines] 408ef5
const int maxn=100005;
struct Trie{
int next[27][maxn];
int endmark[maxn],sz;
bool created[maxn];
void insertTrie(string& s){
int v=0;
for(int i=0;i<(int)s.size();i++){
int c=s[i]-'a';
if(!created[next[c][v]]){
next[c][v]=++sz;
created[sz]=true;
}
v=next[c][v];
}
endmark[v]++;
}
bool searchTrie(string& s){
int v=0;
for(int i=0;i<(int)s.size();i++){
int c=s[i]-'a';
if(!created[next[c][v]])
return false;
v=next[c][v];
}
return(endmark[v]>0);
}
};

9.9 Z-Algorithm [19 lines] e04285
void compute_z_function(const char*S,int N){
int L=0,R=0;
for(int i=1;i<N;++i){
if(i>R){
L=R=i;
while(R<N && S[R-L]==S[R])++R;
Z[i]=R-L,--R;
}
else{
int k=i-L;
if(Z[k]<R-i+1)Z[i]=Z[k];
else{
L=i;
while(R<N && S[R-k]==S[R])++R;
Z[i]=R-L,--R;
}
}
}
}
}

```


Theoretical Computer Science Cheat Sheet	
Definitions	Series
$f(n) = O(g(n))$	iff \exists positive c, n_0 such that $0 \leq f(n) \leq cg(n) \forall n \geq n_0$.
$f(n) = \Omega(g(n))$	iff \exists positive c, n_0 such that $f(n) \geq cg(n) \geq 0 \forall n \geq n_0$.
$f(n) = \Theta(g(n))$	iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$.
$f(n) = o(g(n))$	iff $\lim_{n \rightarrow \infty} f(n)/g(n) = 0$.
$\lim_{n \rightarrow \infty} a_n = a$	iff $\forall \epsilon > 0, \exists n_0$ such that $ a_n - a < \epsilon, \forall n \geq n_0$.
$\sup S$	least $b \in \mathbb{R}$ such that $b \geq s, \forall s \in S$.
$\inf S$	greatest $b \in \mathbb{R}$ such that $b \leq s, \forall s \in S$.
$\liminf_{n \rightarrow \infty} a_n$	$\lim_{n \rightarrow \infty} \inf \{a_i \mid i \geq n, i \in \mathbb{N}\}$.
$\limsup_{n \rightarrow \infty} a_n$	$\lim_{n \rightarrow \infty} \sup \{a_i \mid i \geq n, i \in \mathbb{N}\}$.
$\binom{n}{k}$	Combinations: Size k subsets of a size n set.
$\begin{bmatrix} n \\ k \end{bmatrix}$	Stirling numbers (1st kind): Arrangements of an n element set into k cycles.
$\{n\}_k$	Stirling numbers (2nd kind): Partitions of an n element set into k non-empty sets.
$\langle n \rangle_k$	1st order Eulerian numbers: Permutations $\pi_1 \pi_2 \dots \pi_n$ on $\{1, 2, \dots, n\}$ with k ascents.
$\langle\langle n \rangle\rangle_k$	2nd order Eulerian numbers.
C_n	Catalan Numbers: Binary trees with $n + 1$ vertices.

14. $\begin{bmatrix} n \\ 1 \end{bmatrix} = (n-1)!$,	15. $\begin{bmatrix} n \\ 2 \end{bmatrix} = (n-1)!H_{n-1}$,	16. $\begin{bmatrix} n \\ n \end{bmatrix} = 1$,	17. $\begin{bmatrix} n \\ k \end{bmatrix} \geq \begin{Bmatrix} n \\ k \end{Bmatrix}$,
18. $\begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}$,	19. $\left\{ \begin{matrix} n \\ n-1 \end{matrix} \right\} = \begin{bmatrix} n \\ n-1 \end{bmatrix} = \binom{n}{2}$,	20. $\sum_{k=0}^n \begin{bmatrix} n \\ k \end{bmatrix} = n!$,	21. $C_n = \frac{1}{n+1} \binom{2n}{n}$,
22. $\left\langle \begin{matrix} n \\ 0 \end{matrix} \right\rangle = \left\langle \begin{matrix} n \\ n-1 \end{matrix} \right\rangle = 1$,	23. $\left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle = \left\langle \begin{matrix} n \\ n-1-k \end{matrix} \right\rangle$,	24. $\left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle = (k+1) \left\langle \begin{matrix} n-1 \\ k \end{matrix} \right\rangle + (n-k) \left\langle \begin{matrix} n-1 \\ k-1 \end{matrix} \right\rangle$,	
25. $\left\langle \begin{matrix} 0 \\ k \end{matrix} \right\rangle = \begin{cases} 1 & \text{if } k=0, \\ 0 & \text{otherwise} \end{cases}$,	26. $\left\langle \begin{matrix} n \\ 1 \end{matrix} \right\rangle = 2^n - n - 1$,	27. $\left\langle \begin{matrix} n \\ 2 \end{matrix} \right\rangle = 3^n - (n+1)2^n + \binom{n+1}{2}$,	
28. $x^n = \sum_{k=0}^n \left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle \binom{x+k}{n}$,	29. $\left\langle \begin{matrix} n \\ m \end{matrix} \right\rangle = \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k$,	30. $m! \left\{ \begin{matrix} n \\ m \end{matrix} \right\} = \sum_{k=0}^n \left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle \binom{k}{n-m}$,	
31. $\left\langle \begin{matrix} n \\ m \end{matrix} \right\rangle = \sum_{k=0}^n \left\{ \begin{matrix} n \\ k \end{matrix} \right\} \binom{n-k}{m} (-1)^{n-k-m} k!$,	32. $\left\langle \begin{matrix} n \\ 0 \end{matrix} \right\rangle = 1$,	33. $\left\langle \begin{matrix} n \\ n \end{matrix} \right\rangle = 0$ for $n \neq 0$,	
34. $\left\langle\langle \begin{matrix} n \\ k \end{matrix} \rangle\rangle = (k+1) \left\langle\langle \begin{matrix} n-1 \\ k \end{matrix} \rangle\rangle + (2n-1-k) \left\langle\langle \begin{matrix} n-1 \\ k-1 \end{matrix} \rangle\rangle$,		35. $\sum_{k=0}^n \left\langle\langle \begin{matrix} n \\ k \end{matrix} \rangle\rangle = \frac{(2n)^{\underline{n}}}{2^n}$,	
36. $\left\{ \begin{matrix} x \\ x-n \end{matrix} \right\} = \sum_{k=0}^n \left\langle\langle \begin{matrix} n \\ k \end{matrix} \rangle\rangle \binom{x+n-1-k}{2n}$,		37. $\left\{ \begin{matrix} n+1 \\ m+1 \end{matrix} \right\} = \sum_k \binom{n}{k} \left\{ \begin{matrix} k \\ m \end{matrix} \right\} = \sum_{k=0}^n \left\{ \begin{matrix} k \\ m \end{matrix} \right\} (m+1)^{n-k}$,	

1. $\binom{n}{k} = \frac{n!}{(n-k)!k!}$,	2. $\sum_{k=0}^n \binom{n}{k} = 2^n$,	3. $\binom{n}{k} = \binom{n}{n-k}$,
4. $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$,	5. $\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$,	
6. $\binom{n}{m} \binom{m}{k} = \binom{n}{k} \binom{n-k}{m-k}$,	7. $\sum_{k=0}^n \binom{r+k}{k} = \binom{r+n+1}{n}$,	
8. $\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}$,	9. $\sum_{k=0}^n \binom{r}{k} \binom{s}{n-k} = \binom{r+s}{n}$,	
10. $\binom{n}{k} = (-1)^k \binom{k-n-1}{k}$,	11. $\left\{ \begin{matrix} n \\ 1 \end{matrix} \right\} = \left\{ \begin{matrix} n \\ n \end{matrix} \right\} = 1$,	
12. $\left\{ \begin{matrix} n \\ 2 \end{matrix} \right\} = 2^{n-1} - 1$,	13. $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = k \left\{ \begin{matrix} n-1 \\ k \end{matrix} \right\} + \left\{ \begin{matrix} n-1 \\ k-1 \end{matrix} \right\}$,	

Harmonic series:	
$H_n = \sum_{i=1}^n \frac{1}{i}$,	$\sum_{i=1}^n iH_i = \frac{n(n+1)}{2} H_n - \frac{n(n-1)}{4}$.
$\sum_{i=1}^n H_i = (n+1)H_n - n$,	$\sum_{i=1}^n \binom{i}{m} H_i = \binom{n+1}{m+1} \left(H_{n+1} - \frac{1}{m+1} \right)$.

Geometric series:	
$\sum_{i=0}^n c^i = \frac{c^{n+1} - 1}{c - 1}, \quad c \neq 1,$	$\sum_{i=0}^{\infty} c^i = \frac{1}{1-c}, \quad c < 1,$
$\sum_{i=0}^n ic^i = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^2}, \quad c \neq 1,$	$\sum_{i=0}^{\infty} ic^i = \frac{c}{(1-c)^2}, \quad c < 1.$

coprime $(k+1)$ -tuple together with n . It is a generalization of Euler's totient, $\phi(n) = J_1(n)$.

$$J_k(n) = n^k \prod_{p|n} \left(1 - \frac{1}{p^k}\right)$$

$$\sum_{d|n} J_k(d) = n^k$$

$$\sum_{d|n} \varphi(d) = n$$

$$\varphi(n) = \sum_{d|n} \mu(d) \cdot \frac{n}{d} = n \sum_{d|n} \frac{\mu(d)}{d}$$

- $a | b \implies \varphi(a) | \varphi(b)$
- $n | \varphi(a^n - 1)$ for $a, n > 1$
- $\varphi(mn) = \varphi(m)\varphi(n) \cdot \frac{d}{\varphi(d)}$ where $d = \gcd(m, n)$

Note the special cases

- $\varphi(2m) = \begin{cases} 2\varphi(m) & \text{if } m \text{ is even} \\ \varphi(m) & \text{if } m \text{ is odd} \end{cases}$
- $\varphi(n^m) = n^{m-1}\varphi(n)$
- $\varphi(\text{lcm}(m, n)) \cdot \varphi(\gcd(m, n)) = \varphi(m) \cdot \varphi(n)$
Compare this to the formula
- $\text{lcm}(m, n) \cdot \gcd(m, n) = m \cdot n$
(See [least common multiple](#).)
- $\varphi(n)$ is even for $n \geq 3$. Moreover, if n has r distinct odd prime factors, $2^r | \varphi(n)$

```

    ///n / x yields the same value for i <= x <= la.
}
return 0;
}

```

Mobius Function and Inversion

Notes

- For any positive integer n , define $\mu(n)$ as the sum of the primitive n th roots of unity. It has values in $\{-1, 0, 1\}$ depending on the factorization of n into prime factors:
 - ✓ $\mu(n) = 1$ if n is a square-free positive integer with an even number of prime factors.
 - ✓ $\mu(n) = -1$ if n is a square-free positive integer with an odd number of prime factors.
 - ✓ $\mu(n) = 0$ if n has a squared prime factor.
- Here, a root of unity, occasionally called a de Moivre number, is any complex number that gives 1 when raised to some positive integer power n .
- An n th root of unity, where n is a positive integer (i.e. $n = 1, 2, 3, \dots$), is a number z (maybe complex) satisfying the equation $z^n = 1$.
- An n th root of unity is said to be primitive if it is not a k th root of unity for some smaller k , that is if

$$z^n = 1 \quad \text{and} \quad z^k \neq 1 \text{ for } k = 1, 2, 3, \dots, n-1.$$

- It is a multiplicative function.

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & \text{if } n = 1, \\ 0 & \text{if } n > 1. \end{cases}$$

$$\varphi(n) = \sum_{d|n} d \mu\left(\frac{n}{d}\right)$$

$$\sum_{d|n} \frac{\mu^2(d)}{\varphi(d)} = \frac{n}{\varphi(n)}$$

$$\sum_{\substack{1 \leq k \leq n \\ \gcd(k, n) = 1}} k = \frac{1}{2} n \varphi(n) \quad \text{for } n > 1$$

$$\frac{\varphi(n)}{n} = \frac{\varphi(\text{rad}(n))}{\text{rad}(n)}$$

where,

$$\text{rad}(n) = \prod_{\substack{p|n \\ p \text{ prime}}} p$$

$$\left\lfloor \frac{n}{\varphi(n)} \right\rfloor \text{ is periodic. } 1, 2, 1, 2, 1, 3, 1, 2, 1, 2, 1, 3, \dots$$

$$\sum_{\substack{1 \leq k \leq n \\ \gcd(k, n) = 1}} \gcd(k-1, n) = \varphi(n) d(n)$$

- where $d(n)$ is number of divisors.
- same equation for $\gcd(ak-1, n)$ where a and n are coprime.
- for every n there is at least one other integer $m \neq n$ such that $\phi(m) = \phi(n)$.

Divisor function

$$\sigma_x(n) = \sum_{d|n} d^x$$

- It is multiplicative i.e. if $\gcd(a, b) = 1 \implies \sigma_x(ab) = \sigma_x(a)\sigma_x(b)$.

$$\begin{aligned} \sum_{i=1}^n [\gcd(i, n) = k] &= \varphi\left(\frac{n}{k}\right) \\ \sum_{k=1}^n \gcd(k, n) &= \sum_{d|n} d \cdot \varphi\left(\frac{n}{d}\right) \\ \sum_{k=1}^n \frac{1}{\gcd(k, n)} &= \sum_{d|n} \frac{1}{d} \cdot \varphi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \varphi(d) \\ \sum_{k=1}^n \frac{k}{\gcd(k, n)} &= \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \varphi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \varphi(d) \\ \sum_{k=1}^n \frac{n}{\gcd(k, n)} &= 2 \cdot \sum_{k=1}^n \frac{k}{\gcd(k, n)} - 1, \text{ for } n > 1 \end{aligned}$$

- Given several integers, with integer x appears c_x times, and some fixed integer m . It is asked that how many integers that are co-prime to m , so,

$$\sum_{i=1}^n c_i [\gcd(i, m) = 1] = \sum_{d|m} \mu(d) \sum_{i=1}^{\lfloor n/d \rfloor} c_{id}$$

The classic version states that if g and f are arithmetic functions satisfying

$$g(n) = \sum_{d|n} f(d) \quad \text{for every integer } n \geq 1$$

then

$$f(n) = \sum_{d|n} \mu(d) g\left(\frac{n}{d}\right) \quad \text{for every integer } n \geq 1$$

- $\sum_{d|n} \mu(d) = [n = 1]$
- $\sum_{i=1}^n \sum_{j=1}^n [\gcd(i, j) = 1] = \sum_{d=1}^n \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^2$
- $\sum_{i=1}^n \sum_{j=1}^n \gcd(i, j) = \sum_{d=1}^n \varphi(d) \left\lfloor \frac{n}{d} \right\rfloor^2$
if $F(n) = \prod_{d|n} f(d)$, then $f(n) = \prod_{d|n} F\left(\frac{n}{d}\right)^{\mu(d)}$

mobius function

```

int mob[N];
void mobius()

```

- Any three consecutive Fibonacci numbers are pairwise coprime, which means that, for every n , $\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = \gcd(F_{n+1}, F_{n+2}) = 1$.
- If p is a prime,

$$\begin{cases} p = 5 & \implies p | F_p, \\ p \equiv \pm 1 \pmod{5} & \implies p | F_{p-1} \\ p \equiv \pm 2 \pmod{5} & \implies p | F_{p+1} \end{cases}$$

- The only nontrivial square Fibonacci number is 144. Attila Pethő proved in 2001 that there is only a finite number of perfect power Fibonacci numbers. In 2006, Y. Bugeaud, M. Mignotte, and S. Siksek proved that 8 and 144 are the only such non-trivial perfect powers.
- If the members of the Fibonacci sequence are taken mod n , the resulting sequence is periodic with period at most $6n$.

Sum of floors

$$\sum_{i=1}^n \left\lfloor \frac{n}{i} \right\rfloor = ?$$

```

int32_t main()
{
    BeatMeScanf;
    int i, j, k, n, m;
    cin >> n;
    ///complexity O(sqrt(n))
    for (int i = 1, last; i <= n; i = last + 1) {
        last = n / (n / i);
        debug(i, last, n/i);
    }
}

```

```

{
    for(int i=1; i<N; i++) mob[i]=3;
    mob[1]=1;
    for(int i=2; i<N; i++){
        if(mob[i]==3){
            mob[i]=-1;
            for(int j=2*i; j<N; j+=i) mob[j]=(mob[j]==3?-1:mob[j]*(-1));
            if(i<=(N-1)/i) for(int j=i*i; j<N; j+=i*i) mob[j]=0;
        }
    }
}

```

GCD and LCM

$$\gcd(a, 0) = a$$

$$\gcd(a, b) = \gcd(b, a \bmod b)$$

- Every common divisor of a and b is a divisor of $\gcd(a, b)$.
- If a divides the product $b \cdot c$, and $\gcd(a, b) = d$, then a/d divides c .
- If m is any integer, then $\gcd(a + m \cdot b, b) = \gcd(a, b)$
- The gcd is a multiplicative function in the following sense: if a_1 and a_2 are relatively prime, then $\gcd(a_1 \cdot a_2, b) = \gcd(a_1, b) \cdot \gcd(a_2, b)$.
- $\gcd(a, b) \cdot \text{lcm}(a, b) = |a \cdot b|$
- $\gcd(a, \text{lcm}(b, c)) = \text{lcm}(\gcd(a, b), \gcd(a, c))$
- $\text{lcm}(a, \gcd(b, c)) = \gcd(\text{lcm}(a, b), \text{lcm}(a, c))$.
- For non-negative integers a and b , where a and b are not both zero,

$$\gcd(n^a - 1, n^b - 1) = n^{\gcd(a, b)} - 1.$$

$$r_8(n) = 16 \sum_{d|n} (-1)^{n+d} d^3$$

Gauss Circle Theorem

- The Gauss circle problem is the problem of determining how many integer lattice points there are in a circle centered at the origin and with radius r .
- Since the equation of this circle is given in Cartesian coordinates by $x^2 + y^2 = r^2$, the question is equivalently asking how many pairs of integers m and n there are such that $m^2 + n^2 \leq r^2$
- If the answer for a given r is denoted by $N(r)$ then
$$N(r) = 1 + 4 \sum_{i=0}^{\infty} \left(\left\lfloor \frac{r^2}{4i+1} \right\rfloor - \left\lfloor \frac{r^2}{4i+3} \right\rfloor \right)$$
- A much simpler sum appears if the sum of squares function $r_2(n)$ is defined as the number of ways of writing the number n as the sum of two squares. Then

$$N(r) = \sum_{n=0}^{r^2} r_2(n).$$

3. Combinatorics

Notes

- $\sum_{0 \leq k \leq n} \binom{n-k}{k} = \text{Fib}_{n+1}$
- $\binom{n}{k} = \binom{n}{n-k}$
- $\binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$
- $k \binom{n}{k} = n \binom{n-1}{k-1}$

- $\binom{n}{k} = \frac{n!}{k!(n-k)!}$
- $\sum_{i=0}^n \binom{n}{i} = 2^n$
- $\sum_{i \geq 0} \binom{n}{2i} = 2^{n-1}$
- $\sum_{i \geq 0} \binom{n}{2i+1} = 2^{n-1}$
- $\sum_{i=0}^k (-1)^i \binom{n}{i} = (-1)^k \binom{n-1}{k}$
- $\sum_{i=0}^k \binom{n+i}{i} = \binom{n+k+1}{k}$
- $\sum_{i=0}^k \binom{n+i}{n} = \binom{n+k+1}{k}$
- $\sum_{i=0}^k \binom{i}{n} = \binom{k+1}{n+1}$
- $1 \cdot \binom{n}{1} + 2 \cdot \binom{n}{2} + 3 \cdot \binom{n}{3} + \dots + n \cdot \binom{n}{n} = n \cdot 2^{n-1}$
- $1^2 \cdot \binom{n}{1} + 2^2 \cdot \binom{n}{2} + 3^2 \cdot \binom{n}{3} + \dots + n^2 \cdot \binom{n}{n} = (n+1) \cdot 2^{n-2}$
- Vandermonde's Identity: $\sum_{k=0}^r \binom{m}{k} \binom{n}{r-k} = \binom{m+n}{r}$
- Hockey-Stick

$$n, r \in \mathbb{N}, n > r, \sum_{i=r}^n \binom{i}{r} = \binom{n+1}{r+1}$$

Identity:

$$\sum_{i=0}^k \binom{k}{i}^2 = \binom{2k}{k}$$

- $\sum_{k=0}^n \binom{n}{k} \binom{n}{n-k} = \binom{2n}{n}$
- $\sum_{k=q}^n \binom{n}{k} \binom{k}{q} = 2^{n-q} \binom{n}{q}$
- $\sum_{i=0}^n 3^i \binom{n}{i} = 4^n$
- $\sum_{i=0}^n \binom{2n}{i} = 2^{2n-1} + \frac{1}{2} \binom{2n}{n}$
- $\sum_{i=1}^n \binom{n}{i} \binom{n-1}{i-1} = \binom{2n-1}{n-1}$

- $\sum_{i=0}^n \binom{2n}{i}^2 = \frac{1}{2} \left\{ \binom{4n}{2n} + \binom{2n}{n}^2 \right\}$
- An integer $n \geq 2$ is prime if and only if all the intermediate binomial coefficients $\binom{n}{1}, \binom{n}{2}, \dots, \binom{n}{n-1}$ are divisible by n .
- $\binom{n+k}{k}$ divides $\frac{\text{lcm}(n, n+1, \dots, n+k)}{n}$
- Kummer's theorem states that for given integers $n \geq m \geq 0$ and a prime number p , the largest power of p dividing $\binom{n}{m}$ is equal to the number of carries when m is added to $n - m$ in base p .
- Number of different binary sequences of length n such that no two 0's are adjacent = Fib_{n+1}
- Combination with repetition: Let's say we choose k elements from an n -element set, the order doesn't matter and each element can be chosen more than once. In that case, the number of different combinations is: $\binom{n+k-1}{k}$
- Number of ways to divide n different persons in n/k equal groups i.e. each having size k is $\binom{n-1}{k-1}$
- The number non-negative solution of the equation $x_1 + x_2 + x_3 + \dots + x_k = n$ is $\binom{n+k-1}{n}$
- Number of binary sequence of length n and with k '1' is $\binom{n}{k}$
- The number of ordered pairs (a, b) of binary sequences of length n , such that the distance between them is k , can be

$$\binom{n}{k} \cdot 2^n$$

calculated as follows:
The distance between a and b is the number of components that differs in a and b — for example, the distance between $(0, 0, 1, 0)$ and $(1, 0, 1, 1)$ is 2).

Catalan numbers

- $C_n = \frac{1}{n+1} \binom{2n}{n}$
- $C_0 = 1, C_1 = 1$ and $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$
- 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786
- Number of correct bracket sequence consisting of n opening and n closing brackets.
- The number of ways to completely parenthesize $n+1$ factors.
- The number of triangulations of a convex polygon with $n+2$ sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).
- The number of ways to connect the $2n$ points on a circle to form n disjoint i.e. non-intersecting chords.
- The number of monotonic lattice paths from point $(0,0)$ to point (n,n) in a square lattice of size $n \times n$, which do not pass above the main diagonal (i.e. connecting $(0,0)$ to (n,n)).
- Number of permutations of length n that can be stack sorted (i.e. it can be shown that the

rearrangement is stack sorted if and only if there is no such index $i < j < k$, such that $a_k < a_i < a_j$).

- ✓ The number of **non-crossing partitions** of a set of n elements.
- ✓ The number of rooted full binary trees with $n+1$ leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.
- ✓ The number of Dyck words of length $2n$. A Dyck word is a string consisting of n X's and n Y's such that no initial segment of the string has more Y's than X's. For example, the following are the Dyck words of length 6: XXXYYY XYXXYY XYXYXY XYYXXY XXYXXY.
- ✓ The number of different ways a convex polygon with $n+2$ sides can be cut into triangles by connecting vertices with straight lines (a form of Polygon triangulation)
- ✓ Number of permutations of $\{1, \dots, n\}$ that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing subsequence. For $n = 3$, these permutations are 132, 213, 231, 312 and 321. For $n = 4$, they are 1432, 2143, 2413, 2431, 3142, 3214, 3241, 3412, 3421, 4132, 4213, 4231, 4312 and 4321
- ✓ Number of ways to tile a staircase shape of height n with n rectangles.

- ✓ $N(n, k) = \frac{1}{n} \binom{n}{k} \binom{n}{k-1}$
- ✓ The number of expressions containing n pairs of parentheses, which are correctly matched and which contain k distinct nestings. For instance, $N(4, 2) = 6$ as with four pairs of parentheses six sequences can be created which each contain two times the sub-pattern '()':

$()(())() (())()() (())(()) (())()() (())()()$

- ✓ The number of paths from $(0, 0)$ to $(2n, 0)$, with steps only northeast and southeast, not straying below the x -axis, with k peaks. And sum of all number of peaks is Catalan number.

Stirling numbers of the first kind

- ✓ The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).
- ✓ $S(n, k)$ counts the number of permutations of n elements with k disjoint cycles.
- ✓ $S(n, k) = (n-1) * S(n-1, k) + S(n-1, k-1)$, where $S(0, 0) = 1, S(n, 0) = S(0, n) = 0$
- ✓ $\sum_{k=0}^n S(n, k) = n!$

Stirling numbers of the second kind

- ✓ Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.
- ✓ $S(n, k) = k * S(n-1, k) + S(n-1, k-1)$,
where $S(0, 0) = 1, S(n, 0) = S(0, n) = 0$
- ✓ $S(n, 2) = 2^{n-1} - 1$
- ✓ $S(n, k) * k!$ = number of ways to color n nodes using colors from 1 to k such that each color is used at least once.

Bell number

- ✓ Counts the number of partitions of a set.
- ✓ $B_{n+1} = \sum_{k=0}^n \binom{n}{k} * B_k$
- ✓ $B_n = \sum_{k=0}^n S(n, k)$, where $S(n, k)$ is stirling number of second kind.
- ✓ The number of multiplicative partitions of a squarefree number with i prime factors is the i -th Bell number, B_i .
- ✓ If a deck of n cards is shuffled by repeatedly removing the top card and reinserting it anywhere in the deck (including its original position at the top of the deck), with exactly n repetitions of this operation, then there are n^n different shuffles that can be performed. Of these, the number that return the deck to its original sorted order is exactly B_n . Thus, the probability that the deck is in its original order after shuffling it in this way is B_n/n^n .

Lucas Theorem

- ✓ If p is prime the $\binom{p^a}{k} \equiv 0 \pmod{p}$
- ✓ For non-negative integers m and n and a prime p , the following congruence relation holds:

$$\binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p},$$

where,

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$, when $m < n$.

Derangement

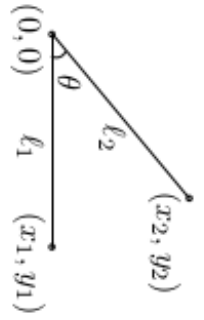
- ✓ A derangement is a permutation of the elements of a set, such that no element appears in its original position.
- ✓ $d(n) = (n-1) * (d(n-1) + d(n-2))$,
where $d(0) = 1, d(1) = 0$
- ✓ $d(n) = \left\lfloor \frac{n!}{e} \right\rfloor, n \geq 1$

4. Burnside Lemma

The task is to count the number of different necklaces from n beads, each of which can be painted in one of the k colors. When comparing two necklaces, they can be rotated, but not reversed (i.e. a cyclic shift is permitted).

Solution:

Identities Cont.		Trees
38. $\binom{n+1}{m+1} = \sum_k \binom{n}{k} \binom{k}{m} = \sum_{k=0}^n \binom{k}{m} n^{\overline{n-k}} = n! \sum_{k=0}^n \frac{1}{k!} \binom{k}{m}$,	39. $\left[\begin{matrix} x \\ x-n \end{matrix} \right] = \sum_{k=0}^n \left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle \binom{x+k}{2n}$,	Every tree with n vertices has $n-1$ edges.
40. $\left\{ \begin{matrix} n \\ m \end{matrix} \right\} = \sum_k \binom{n}{k} \left\{ \begin{matrix} k+1 \\ m+1 \end{matrix} \right\} (-1)^{n-k}$,	41. $\left[\begin{matrix} n \\ m \end{matrix} \right] = \sum_k \left[\begin{matrix} n+1 \\ k+1 \end{matrix} \right] \binom{k}{m} (-1)^{m-k}$,	Kraft inequality: If the depths of the leaves of a binary tree are d_1, \dots, d_n :
42. $\left\{ \begin{matrix} m+n+1 \\ m \end{matrix} \right\} = \sum_{k=0}^m k \left\{ \begin{matrix} n+k \\ k \end{matrix} \right\}$,	43. $\left[\begin{matrix} m+n+1 \\ m \end{matrix} \right] = \sum_{k=0}^m k(n+k) \left[\begin{matrix} n+k \\ k \end{matrix} \right]$,	$\sum_{i=1}^n 2^{-d_i} \leq 1$,
44. $\binom{n}{m} = \sum_k \left\{ \begin{matrix} n+1 \\ k+1 \end{matrix} \right\} \left[\begin{matrix} k \\ m \end{matrix} \right] (-1)^{m-k}$,	45. $(n-m)! \binom{n}{m} = \sum_k \left[\begin{matrix} n+1 \\ k+1 \end{matrix} \right] \left\{ \begin{matrix} k \\ m \end{matrix} \right\} (-1)^{m-k}$, for $n \geq m$,	and equality holds only if every internal node has 2 sons.
46. $\left\{ \begin{matrix} n \\ n-m \end{matrix} \right\} = \sum_k \binom{m-n}{m+k} \binom{m+n}{n+k} \left[\begin{matrix} m+k \\ k \end{matrix} \right]$,	47. $\left[\begin{matrix} n \\ n-m \end{matrix} \right] = \sum_k \binom{m-n}{m+k} \binom{m+n}{n+k} \left\{ \begin{matrix} m+k \\ k \end{matrix} \right\}$,	
48. $\left\{ \begin{matrix} n \\ \ell+m \end{matrix} \right\} \binom{\ell+m}{\ell} = \sum_k \left\{ \begin{matrix} k \\ \ell \end{matrix} \right\} \left\{ \begin{matrix} n-k \\ m \end{matrix} \right\} \binom{n}{k}$,	49. $\left[\begin{matrix} n \\ \ell+m \end{matrix} \right] \binom{\ell+m}{\ell} = \sum_k \left[\begin{matrix} k \\ \ell \end{matrix} \right] \left[\begin{matrix} n-k \\ m \end{matrix} \right] \binom{n}{k}$.	

Geometry
Projective coordinates: triples (x, y, z) , not all x, y and z zero.
$(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0$.
Cartesian Projective
$(x, y) \quad (x, y, 1)$
$y = mx + b \quad (m, -1, b)$
$x = c \quad (1, 0, -c)$
Distance formula, L_p and L_∞ metric:
$\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2}$,
$\left[x_1 - x_0 ^p + y_1 - y_0 ^p \right]^{1/p}$,
$\lim_{p \rightarrow \infty} \left[x_1 - x_0 ^p + y_1 - y_0 ^p \right]^{1/p}$.
Area of triangle $(x_0, y_0), (x_1, y_1)$ and (x_2, y_2) :
$\frac{1}{2} \text{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}$.
Angle formed by three points:

$\cos \theta = \frac{(x_1, y_1) \cdot (x_2, y_2)}{\ell_1 \ell_2}$.
Line through two points (x_0, y_0) and (x_1, y_1) :
$\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0$.
Area of circle, volume of sphere:
$A = \pi r^2, \quad V = \frac{4}{3} \pi r^3$.

$$\sum_{d|n} d = O(n \log \log n).$$

The number of divisors of n is at most around 100 for $n < 5e4$, 500 for $n < 1e7$, 2000 for $n < 1e10$, 200 000 for $n < 1e19$.

Combinatorial (5)

5.1 Permutations

5.1.1 Factorial

n	1	2	3	4	5	6	7	8	9	10
$n!$	1	2	6	24	120	720	5040	40320	362880	3628800
n	11	12	13	14	15	16	17			
$n!$	4.0e7	4.8e8	6.2e9	8.7e10	1.3e12	2.1e13	3.6e14			
n	20	25	30	40	50	100	150	171		
$n!$	2e18	2e25	3e32	8e47	3e64	9e157	6e262	>DBL_MAX		

IntPerm.h

Description: Permutation \rightarrow integer conversion. (Not order preserving.)

Time: $O(n)$

6 lines

```
int permToInt(vi& v) {
    int use = 0, i = 0, r = 0;
    trav(x,v) r=r * ++i + __builtin_popcount(use & ~(1 << x)),
        use |= 1 << x; // (note: minus, not ~!)
    return r;
} // hash-cpp-all = elb8eaea02324af14a3da94f409019b8
```

5.1.2 Cycles

Let $g_S(n)$ be the number of n -permutations whose cycle lengths all belong to the set S . Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp \left(\sum_{n \in S} \frac{x^n}{n} \right)$$

5.1.3 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1)+D(n-2)) = nD(n-1)+(-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

5.1.4 Burnside's lemma

Given a group G of symmetries and a set X , the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g ($g.x = x$).

Sums of powers:

$$\sum_{i=1}^n i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\begin{aligned} \sum_{i=m}^{\infty} f(i) &= \int_m^{\infty} f(x) dx - \sum_{k=1}^{\infty} \frac{B_k}{k!} f^{(k-1)}(m) \\ &\approx \int_m^{\infty} f(x) dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m)) \end{aligned}$$

5.3.2 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n, k) = c(n-1, k-1) + (n-1)c(n-1, k), \quad c(0, 0) = 1$$

$$\sum_{k=0}^n c(n, k) x^k = x(x+1) \dots (x+n-1)$$

$$c(8, k) =$$

$$8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1$$

$$c(n, 2) =$$

$$0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$$

5.3.3 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j :s s.t.

$\pi(j) > \pi(j+1)$, $k+1$ j :s s.t. $\pi(j) \geq j$, k j :s s.t.

$\pi(j) > j$.

$$E(n, k) = (n-k)E(n-1, k-1) + (k+1)E(n-1, k)$$

$$E(n, 0) = E(n, n-1) = 1$$

$$E(n, k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$$

5.3.4 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n, k) = S(n-1, k-1) + kS(n-1, k)$$

$$S(n, 1) = S(n, n) = 1$$

$$S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$$

5.3.5 Bell numbers

Total number of partitions of n distinct elements.

$B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$ For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

5.3.6 Labeled unrooted trees

on n vertices: n^{n-2}

on k existing trees of size n_i : $n_1 n_2 \dots n_k n^{k-2}$

with degrees d_i : $(n-2)! / ((d_1-1)! \dots (d_n-1)!)$

5.3.7 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \quad C_{n+1} = \frac{2(2n+1)}{n+2} C_n, \quad C_{n+1} = \sum C_i C_{n-i}$$

$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with $n+1$ leaves (0 or 2 children).
- ordered trees with $n+1$ vertices.
- ways a convex polygon with $n+2$ sides can be cut into triangles by connecting vertices with straight lines.
- permutations of $[n]$ with no 3-term increasing subseq.