KUET\_Effervescent Team Notebook

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```
if(L == R)
    return node;
  int mid = (L+R)>>1;
 leftchild[node] = build(L, mid);
  rightchild[node] = build(mid+1, R);
  //initialize value[node]
 return node;
int update(int nownode, int L, int R, int val){
  int node = ++now;
  if(L == R){
    //value[node] = value[nownode]+1;
    //update value[node]
   return node;
  int mid = (L+R)>>1;
  leftchild[node] = leftchild[nownode];
  rightchild[node] = rightchild[nownode];
  if (mid <= val) {//change condition as required
    leftchild[node] = update(leftchild[nownode],
    L, mid, val);
  }
  else{
    rightchild[node] =
    update(rightchild[nownode], mid+1, R, val);
  //value[node] = value[nownode]+1;
  //update value[node]
 return node;
int query(int leftnode, int rightnode, int L, int
    R. int k)
  if(L==R) return L;
  //int leftcnt = value[leftchild[rightnode]]-
    value[leftchild[leftnode]];a
  //change as required
  int mid = (L+R) >> 1;
  if(leftcnt >= k){//change condition as required
    return query(leftchild[leftnode],
    leftchild[rightnode], L, mid, k);
 else{
    return query(rightchild[leftnode],
    rightchild[rightnode], mid+1, R, k-leftcnt);
void persistentsegtree(){
 root[0] = build(0, mxn);
 for(int i=1; i<=n; i++){
   root[i] = update(root[i-1], 0, mxn, a[i]);
```

```
1.4 SQRT Decomposition [106 lines]
struct sqrtDecomposition
    static const int sz = 320; //sz = sqrt(N);
    int numberofblocks;
    struct node
 {
    int L, R;
    bool islazy = false;
    11 lazyval=0;
    //extra data needed for different problems
    void ini(int 1, int r)
      for(int i=1; i<=r; i++){
        //...initialize as need
     L=1, R=r;
    void semiupdate(int 1, int r, 11 val){
      if(1>r) return;
      if(islazy){
        for(int i=L; i<=R; i++){
          //...distribute lazy to everyone
        islazv = 0;
        lazyval = 0;
      for(int i=1; i<=r; i++){</pre>
        //...do it manually
    void fullupdate(ll val){
      if(islazy){
        //...only update lazyval
      else{
        for(int i=L; i<=R; i++){
          //...everyone are not equal, make them
    equal
        islazy = 1;
        //update lazyval
    void update(int 1, int r, ll val){
      if(l<=L && r>=R) fullupdate(val);
      else semiupdate(max(1, L), min(r, R), val);
    11 semiquery(int 1, int r){
      if(1>r) return 0;
      if(islazy){
        for(int i=L; i<=R; i++){
          //...distribute lazy to everyone
        islazy = 0;
        lazyval = 0;
```

```
11 \text{ ret} = 0;
      for(int i=1; i<=r; i++){
        //...take one by one
      return ret;
    11 fullquery(){
      //return stored value;
    11 query(int 1, int r){
      if(1<=L && r>=R) return fullquery();
      else return semiquery(max(1, L), min(r, R));
  };
  vector<node> blocks;
  void init(int n){
    numberofblocks = (n+sz-1)/sz;
    int curL = 1, curR = sz;
    blocks.resize(numberofblocks+5);
    for(int i=1; i<=numberofblocks; i++){</pre>
      curR = min(n, curR);
      blocks[i].ini(curL, curR);
      curL += sz;
      curR += sz;
 }
  void update(int 1, int r, 11 val){
    int left = (1-1)/sz+1;
    int right = (r-1)/sz+1;
    for(int i=left; i<=right; i++){</pre>
      blocks[i].update(1, r, val);
 }
  11 query(int 1, int r){
    int left = (1-1)/sz+1;
    int right = (r-1)/sz+1;
    11 \text{ ret} = 0;
    for(int i=left; i<=right; i++){</pre>
      ret += blocks[i].query(l, r);
    return ret;
 }
1.5 Segment Tree [73 lines]
/*edit:data,combine,build check datatype*/
template<typename T>
struct SegmentTree {
```

#define lc ( $C \ll 1$ )

#define rc (C << 1 | 1)

```
void combine(data& cur, data& 1, data& r) {
  st[C].sum = (R - L + 1) * lazy[C];
data Query(int i, int j, int C, int L, int R) {
  if (j < L \mid | i > R \mid | L > R) return data();
  if (i <= L && R <= j) return st[C];
  data d1 = Query(i, j, lc, L, M);
  data d2 = Query(i, j, rc, M + 1, R);
void Update(int i, int j, T val, int C, int L,
  if (j < L \mid | i > R \mid | L > R) return;
```

#define M ((L+R)>>1)

data() :sum(0) {};

st.resize(4 \* N);isLazy.resize(4 \* N);

if (L != R) {

lazv[C] = 0;

if (L == R) {

return;

lazv.resize(4 \* N);

SegmentTree(int \_N) :N(\_N) {

cur.sum = 1.sum + r.sum:

if (!isLazy[C]) return;

lazy[lc] += lazy[C];

lazy[rc] += lazy[C];

isLazy[lc] = 1;

isLazy[rc] = 1;

isLazy[C] = false;

st[C].sum = 0;

build(lc, L, M);

push(C, L, R);

data ret;

return ret;

push(C, L, R);

isLazy[C] = 1;

lazy[C] = val;

push(C, L, R);

int R) {

return;

build(rc, M + 1, R);

// default val O/INF

combine(ret, d1, d2);

if (i <= L && R <= j) {

void push(int C, int L, int R) {

void build(int C, int L, int R) {

combine(st[C], st[lc], st[rc]);

struct data {

vector<data>st; vector<bool>isLazy;

vector<T>lazy;

T sum:

int N;

```
Update(i, j, val, lc, L, M);
    Update(i, j, val, rc, M + 1, R);
    combine(st[C], st[lc], st[rc]);
  void Update(int i, int j, T val) {
    Update(i, j, val, 1, 1, N);
  T Query(int i, int j) {
    return Query(i, j, 1, 1, N).sum;
};
1.6 Trie Bit [61 lines]
struct Trie {
  struct node {
    int next[2];
    int cnt, fin;
    node() :cnt(0), fin(0) {
      for (int i = 0; i < 2; i++) next[i] = -1;
  };
  vector<node>data;
  Trie() {
    data.push_back(node());
  void key_add(int val) {
    int cur = 0;
    for (int i = 30; i \ge 0; i--) {
      int id = (val >> i) & 1;
      if (data[cur].next[id] == -1) {
        data[cur].next[id] = data.size();
        data.push_back(node());
      cur = data[cur].next[id];
      data[cur].cnt++:
    data[cur].fin++;
  int key_search(int val) {
    int cur = 0;
    for (int i = 30; ~i; i--) {
      int id = (val >> i) & 1;
      if (data[cur].next[id] == -1) return 0;
      cur = data[cur].next[id];
    return data[cur].fin;
  void key_delete(int val) {
    int cur = 0:
    for (int i = 30; ~i; i--) {
      int id = (val >> i) & 1;
      cur = data[cur].next[id];
      data[cur].cnt--;
    data[cur].fin--;
  bool key_remove(int val) {
```

```
if (key_search(val)) return key_delete(val),
   return 0;
  int maxXor(int x) {
   int cur = 0;
   int ans = 0;
   for (int i = 30; ~i; i--) {
     int b = (x >> i) & 1;
     if (data[cur].next[!b] + 1 &&
    data[data[cur].next[!b]].cnt > 0) {
       ans += (1LL \ll i);
        cur = data[cur].next[!b];
      else cur = data[cur].next[b];
   return ans;
2 Dynamic Programming
2.1 Convex Hull Trick [91 lines]
struct Hull Static{
/*all m need to be decreasing order
if m is in increasing order then negate the
    m(like, add\_line(-m, c)),
remember in query you have to negate the x also*/
  int min_or_max;//if min then 0 otherwise 1
  int pointer; /*keep track for the best line for
    previous query, requires all insert first*/
  vector<ll>M,C;//y=m*x+c;
  inline void clear(){
   min_or_max=0;//initially with minimum trick
   pointer=0;
   M.clear();
   C.clear():
  Hull_Static(){clear();}
  Hull_Static(int _min_or_max){
    clear():
   this->min_or_max=_min_or_max;
  bool bad_min(int idx1,int idx2,int idx3){
    return(C[idx3]-C[idx1])*(M[idx1]-M[idx2]) <
    (C[idx2]-C[idx1])*(M[idx1]-M[idx3]);
   return 1.0*(C[idx3]-C[idx1])*(M[idx1]-M[idx2])
   1.0*(C[idx2]-C[idx1])*(M[idx1]-M[idx3]);//for
    overflow
  bool bad_max(int idx1,int idx2,int idx3){
   return(C[idx3]-C[idx1])*(M[idx1]-M[idx2]) >
    (C[idx2]-C[idx1])*(M[idx1]-M[idx3]);
   return 1.0*(C[idx3]-C[idx1])*(M[idx1]-M[idx2])
   1.0*(C[idx2]-C[idx1])*(M[idx1]-M[idx3]);//for
    overflow
```

```
11 \text{ mid1}=lo+(hi-lo)/3;
  bool bad(int idx1,int idx2,int idx3){
                                                                 11 \text{ mid} 2=\text{hi}-(\text{hi}-\text{lo})/3;
    if(!min_or_max)return bad_min(idx1,idx2,idx3);
                                                                 11 val1=getval(mid1,x);
    else return bad_max(idx1,idx2,idx3);
                                                                 11 val2=getval(mid2,x);
                                                                 if(val1<val2){</pre>
  void add_line(ll m,ll c){/*add line where m is
                                                                   ans=max(ans,val2);
    given in decreasing order
                                                                   lo=mid1+1;
//if(M.size()>0 \text{ and } M.back()==m)return;//same
    gradient, no need to add
                                                                 else{
above line added from tarango khan, this line cost
                                                                   ans=max(ans,val1);
    me sevaral wa, but some code got ac with this*/
                                                                   hi=mid2-1;
    M.push_back(m);
    C.push_back(c);
                                                              }
    while(M.size()>=3 and bad((int)M.size()-3,
                                                              return ans;
    (int)M.size()-2,(int)M.size()-1)){
      M.erase(M.end()-2);
                                                          };
      C.erase(C.end()-2);
                                                          2.2 Divide and Conquer DP [26 lines]
                                                          11 G,L;///total group, cell size
                                                          ll dp[8001][801], cum[8001];
  ll getval(ll idx,ll x){
                                                          11 C[8001]; ///value of each cell
    return M[idx]*x+C[idx];
                                                          inline ll cost(ll 1,ll r){
                                                            return(cum[r]-cum[l-1])*(r-l+1);
  ll getminval(ll x){/*if queries are
    non-decreasing order*/
                                                          void fn(ll g,ll st,ll ed,ll r1,ll r2){
    while(pointer<(int)M.size()-1 and
                                                            if(st>ed)return;
    getval(pointer+
                                                            11 \text{ mid}=(\text{st+ed})/2, \text{pos}=-1;
    1,x)<getval(pointer,x))pointer++;
                                                            dp[mid][g]=inf;
    return M[pointer] *x+C[pointer];
                                                            for(int i=r1;i<=r2;i++){
                                                              ll tcost=cost(i,mid)+dp[i-1][g-1];
  11 getmaxval(ll x){
                                                              if(tcost<dp[mid][g]){</pre>
    while(pointer<(int)M.size()-1 and
                                                                   dp[mid][g]=tcost,pos=i;
    getval(pointer+
    1,x)>getval(pointer,x))pointer++;
    return M[pointer] *x+C[pointer];
                                                            fn(g,st,mid-1,r1,pos);
                                                            fn(g,mid+1,ed,pos,r2);
  11 getminvalternary(ll x){
    ll lo=0,hi=(ll)M.size()-1;
                                                          int main(){
    11 ans=inf;/*change with problem*/
                                                            for(int i=1;i<=L;i++)
    while(lo<=hi){
                                                               \operatorname{cum}[i] = \operatorname{cum}[i-1] + C[i];
      11 \text{ mid1=lo+(hi-lo)/3};
                                                            for(int i=1;i<=L;i++)
      11 \text{ mid} 2=\text{hi}-(\text{hi}-\text{lo})/3;
                                                              dp[i][1]=cost(1,i);
      ll val1=getval(mid1,x);
                                                            for(int i=2;i<=G;i++)fn(i,1,L,1,L);
      11 val2=getval(mid2,x);
      if(val1<val2){</pre>
        ans=min(ans, val2);
                                                          2.3 Knuth Optimization [32 lines]
        hi=mid2-1:
                                                          /*It is applicable where recurrence is in the form
      }
                                                                                                                     }
      else{
                                                          dp[i][j] = mini < k < j \{ dp[i][k] + dp[k][j] \} + C[i][j]
        ans=min(ans, val1);
                                                          condition for applicability is:
        lo=mid1+1:
                                                          A[i, j-1] \leq A[i, j] \leq A[i+1, j]
                                                          A[i][j]-the smallest k that gives optimal
    return ans;
                                                               answer.like-
                                                          dp[i][j] = dp[i-1][k] + C[k][j]
  11 getmaxvalternary(ll x){
                                                          C[i][j]-given cost function
    ll lo=0, hi=(ll)M.size()-1;
                                                          also applicable if: C[i][j]satisfies the following
    11 ans=-inf;/*change with problem*/
                                                              2 conditions:
    while(lo<=hi){
                                                          C[a][c]+C[b][d] <= C[a][d]+C[b][c], a <= b <= c <= d
```

```
C[b][c] <= C[a][d], a <= b <= c <= d
reduces time complexity from O(n^3) to O(n^2)*/
for(int s=0;s<=k;s++)//s-length(size)of substring
 for(int l=0; l+s <= k; l++) {//l-left point}
    int r=1+s;//r-right point
    if(s<2){
      res[1][r]=0;//DP base-nothing to break
      mid[l][r]=l;/*mid is equal to left border*/
      continue:
    int mleft=mid[l][r-1];/*Knuth's trick: getting
    bounds on m*/
    int mright=mid[l+1][r];
    res[l][\bar{r}]=inf;
    for(int m=mleft;m<=mright;m++){/*iterating for</pre>
    m in the bounds only*/
      int64 tres=res[1] [m]+res[m] [r]+(x[r]-x[1]);
      if(res[1][r]>tres){//relax current solution
        res[l][r]=tres;
        mid[l][r]=m;
int64 answer=res[0][k];
2.4 LIS O(nlogn) with full path [17 lines]
int num[MX],mem[MX],prev[MX],array[MX],res[MX],
    maxlen:
void LIS(int SZ,int num[]){
  CLR(mem), CLR(prev), CLR(array), CLR(res);
  int i,k;
  maxlen=1;
  array[0]=-inf;
  RFOR(i,1,SZ+1) array[i]=inf;
  prev[0] = -1, mem[0] = num[0];
  FOR(i,SZ){
    k=lower_bound(array,array+maxlen+1,num[i])-
    if (k==1) array[k]=num[i],mem[k]=i,prev[i]=-1;
    array[k]=num[i],mem[k]=i,prev[i]=mem[k-1];
    if(k>maxlen) maxlen=k;
  for(i=mem[maxlen];i!=-1;i=prev[i])res[k++]=
    num[i];
2.5 SOS DP [18 lines]
//iterative version
for(int mask = 0; mask < (1<<N); ++mask){</pre>
  dp[mask][-1] = A[mask]; //handle base case
    separately (leaf states)
  for(int i = 0; i < N; ++i){
    if(mask & (1<<i))
      dp[mask][i] = dp[mask][i-1] +
    dp[mask^(1<<i)][i-1];
```

```
dp[mask][i] = dp[mask][i-1];
 F[mask] = dp[mask][N-1];
//memory optimized, super easy to code.
for(int i = 0; i < (1 << N); ++i)
 F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask <
    (1<<N); ++mask){
  if(mask & (1<<i))
    F[mask] += F[mask^(1<<i)];
3 Flow
3.1 Blossom [58 lines]
// Finds Maximum matching in General Graph
// Complexity O(NM)
// mate[i] = j means i is paired with j
// source: https://codeforces.com/blog/entry
    /92339?#comment-810242
vector<int> Blossom(vector<vector<int>>& graph) {
  //mate contains matched edge.
  int n = graph.size(), timer = -1;
  vector<int> mate(n, -1), label(n), parent(n),
    orig(n), aux(n, -1), q;
  auto \bar{l}ca = [\&](int x, int y) {
    for (timer++; ; swap(x, y)) {
      if (x == -1) continue;
      if (aux[x] == timer) return x;
      aux[x] = timer;
      x = (mate[x] == -1 ? -1 :
    orig[parent[mate[x]]]);
    }
  auto blossom = [&](int v, int w, int a) {
    while (orig[v] != a) {
      parent[v] = w; w = mate[v];
      if (label[w] == 1) label[w] = 0,
    q.push_back(w);
      orig[v] = orig[w] = a; v = parent[w];
  };
  auto augment = [&](int v) {
    while (v != -1) {
      int pv = parent[v], nv = mate[pv];
      mate[v] = pv; mate[pv] = v; v = nv;
  };
  auto bfs = [&](int root) {
    fill(label.begin(), label.end(), -1);
    iota(orig.begin(), orig.end(), 0);
    label[root] = 0; q.push_back(root);
    for (int i = 0; i < (int)q.size(); ++i) {
      int v = q[i];
     for (auto x : graph[v]) {
        if (label[x] == -1) {
```

```
label[x] = 1; parent[x] = v;
          if (mate[x] == -1)
            return augment(x), 1;
          label[mate[x]] = 0;
   q.push_back(mate[x]);
        else if (label[x] == 0 && orig[v] !=
   orig[x]) {
          int a = lca(orig[v], orig[x]);
          blossom(x, v, a); blossom(v, x, a);
     }
   }
   return 0;
 // Time halves if you start with (any) maximal
    matching.
 for (int i = 0; i < n; i++)
   if (mate[i] == -1)
      bfs(i);
 return mate;
3.2 Dinic [72 lines]
/*.Complexity: O(V^2 E)
  .Call Dinic with total number of nodes.
  .Nodes start from 0.
  .Capacity is long long data.
  .make graph with create edge(u, v, capacity).
  .Get max flow with maxFlow(src,des).*/
#define eb emplace_back
struct Dinic {
 struct Edge {
   int u, v;
   11 cap, flow = 0;
   Edge() {}
   Edge(int u, int v, ll cap) : u(u), v(v),
    cap(cap) {}
 };
  int N;
 vector<Edge>edge;
 vector<vector<int>>adj;
 vector<int>d, pt;
 Dinic(int N) : N(N), edge(0), adj(N), d(N), pt(N)
  void addEdge(int u, int v, ll cap) {
   if (u == v) return:
   edge.eb(u, v, cap);
    adj[u].eb(edge.size() - 1);
   edge.eb(v, u, 0);
   adj[v].eb(edge.size() - 1);
 bool bfs(int s, int t) {
    queue<int>q({ s });
   fill(d.begin(), d.end(), N + 1);
   d[s] = 0;
   while (!q.empty()) {
```

int u = q.front();q.pop();

if (u == t) break;

```
for (int k : adj[u]) {
        Edge\& e = edge[k];
        if (e.flow < e.cap \&\& d[e.v] > d[e.u] + 1) {
          d[e.v] = d[e.u] + 1;
          q.emplace(e.v);
    return d[t] != N + 1;
  ll dfs(int u, int T, ll flow = -1) {
    if (u == T || flow == 0) return flow;
    for (int& i = pt[u];i < adj[u].size();i++) {</pre>
      Edge& e = edge[adj[u][i]];
      Edge& oe = edge[adj[u][i] ^ 1];
      if (d[e.v] == d[e.u] + 1) {
        11 amt = e.cap - e.flow;
        if (flow !=-1 \&\& amt > flow) amt = flow;
        if (ll pushed = dfs(e.v, T, amt)) {
          e.flow += pushed;
          oe.flow -= pushed;
          return pushed;
    return 0;
  ll maxFlow(int s, int t) {
    11 \text{ total} = 0;
    while (bfs(s, t)) {
      fill(pt.begin(), pt.end(), 0);
      while (ll\ flow = dfs(s, t)) {
        total += flow;
    return total:
};
3.3 Flow [6 lines]
Covering Problems:
> Maximum Independent Set(Bipartite): Largest set
    of nodes which do not have any edge between
    them. sol: V-(MaxMatching)
> Minimum Vertex Cover(Bipartite): -Smallest set
    of nodes to cover all the edges -sol:
    MaxMatching
> Minimum Edge Cover(General graph): -Smallest set
    of edges to cover all the nodes -sol:
    V-(MaxMatching) (if edge cover exists, does
    not exit for isolated nodes)
> Minimum Path Cover(Vertex disjoint) DAG:
    -Minimum number of vertex disjoint paths that
    visit all the nodes -sol: make a bipartite
    graph using same nodes in two sides, one side
    is "from" other is "to", add edges from "from"
```

to "to", then ans is V-(MaxMatching)

```
> Minimum Path Cover(Vertex Not Disjoint) General
    graph: -Minimum number of paths that visit all
    the nodes -sol: consider cycles as nodes then
    it will become a path cover problem with
    vertex disjoint on DAG
3.4 HopCroftKarp [67 lines]
/*. Finds Maximum Matching In a bipartite graph
  . Complexity O(E\sqrt{V})
  .1-indexed
  .No default constructor
  .add single edge for (u, v)*/
struct HK {
  static const int inf = 1e9;
  vector<int>matchL, matchR, dist;
  //matchL contains value of matched node for L
  vector<vector<int>>adj;
  HK(int n) : n(n), matchL(n + 1),
  matchR(n + 1), dist(n + 1), adj(n + 1) {
  void addEdge(int u, int v) {
    adj[u].push_back(v);
                                                         1-indexed */
                                                      struct Hungarian {
  bool bfs() {
    queue<int>q;
   for (int u = 1; u \le n; u++) {
                                                        queue<int> q;
     if (!matchL[u]) {
        dist[u] = 0;
        q.push(u);
                                                        Hungarian() {}
      else dist[u] = inf;
    dist[0] = inf;///unmatched node matches with
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (auto v : adj[u]) {
        if (dist[matchR[v]] == inf) {
          dist[matchR[v]] = dist[u] + 1;
          q.push(matchR[v]);
     }
    return dist[0] != inf;
  bool dfs(int u) {
    if (!u) return true;
   for (auto v : adj[u]) {
                                                          finish = 0;
      if (dist[matchR[v]] == dist[u] + 1
          && dfs(matchR[v])) {
        matchL[u] = v;
        matchR[v] = u;
        return true;
                                                            q.pop();
```

```
dist[u] = inf;
   return false;
 int max_match() {
   int matching = 0;
   while (bfs()) {
     for (int u = 1; u \le n; u++) {
       if (!matchL[u])
          if (dfs(u))
            matching++;
   return matching;
3.5 Hungarian [116 lines]
/* Complexity: O(n^3) but optimized
   It finds minimum cost maximum matching.
  For finding maximum cost maximum matching
  add -cost and return -matching()
 long long c[N][N], fx[N], fy[N], d[N];
 int 1[N], r[N], arg[N], trace[N];
  int start, finish, n;
  const long long inf = 1e18;
 Hungarian(int n1, int n2) : n(max(n1, n2)) {
   for (int i = 1; i <= n; ++i) {
     fy[i] = 1[i] = r[i] = 0;
     for (int j = 1; j \le n; ++j) c[i][j] = inf;
 void add_edge(int u, int v, long long cost) {
   c[u][v] = min(c[u][v], cost);
 inline long long getC(int u, int v) {
   return c[u][v] - fx[u] - fy[v];
 void initBFS() {
   while (!q.empty()) q.pop();
   q.push(start);
   for (int i = 0; i <= n; ++i) trace[i] = 0;
   for (int v = 1; v \le n; ++v) {
     d[v] = getC(start, v);
     arg[v] = start;
 void findAugPath() {
   while (!q.empty()) {
      int u = q.front();
```

```
for (int v = 1; v \le n; ++v) if (!trace[v])
      long long w = getC(u, v);
      if (!w) {
        trace[v] = u;
        if (!r[v]) {
         finish = v;
          return;
        q.push(r[v]);
      if (d[v] > w) {
        d[v] = w;
        arg[v] = u;
void subX_addY() {
  long long delta = inf;
  for (int v = 1; v <= n; ++v) if (trace[v] == 0
  && d[v] < delta) {
    delta = d[v];
  // Rotate
  fx[start] += delta:
  for (int v = 1; v \le n; ++v) if (trace[v]) {
    int u = r[v];
    fy[v] -= delta;
    fx[u] += delta;
  else d[v] -= delta;
  for (int v = 1; v \le n; ++v) if (!trace[v] &&
  !d[v]) {
   trace[v] = arg[v];
    if (!r[v]) {
      finish = v;
      return;
    q.push(r[v]);
void Enlarge() {
  do {
   int u = trace[finish];
    int nxt = l[u]:
    l[u] = finish:
    r[finish] = u;
    finish = nxt;
  } while (finish);
long long maximum_matching() {
  for (int u = 1; u \le n; ++u) {
   fx[u] = c[u][1];
   for (int v = 1; v \le n; ++v) {
      fx[u] = min(fx[u], c[u][v]);
```

```
n = _n + 10;
    for (int v = 1; v \le n; ++v) {
      fy[v] = c[1][v] - fx[1];
      for (int u = 1; u \le n; ++u) {
                                                          mxid = 0:
        fy[v] = min(fy[v], c[u][v] - fx[u]);
    for (int u = 1; u \le n; ++u) {
      start = u;
      initBFS();
      while (!finish) {
        findAugPath();
        if (!finish) subX_addY();
      Enlarge();
                                                          true);
    long long ans = 0;
    for (int i = 1; i \le n; ++i) {
      if (c[i][1[i]] != inf) ans += c[i][1[i]];
      else l[i] = 0;
    }
                                                          d[s] = 0;
    return ans;
};
3.6 MCMF [116 lines]
/*Credit: ShahjalalShohag
                                                            q.pop();
  . Works for both directed, undirected and with
    negative cost too
  .doesn't work for negative cycles
  .for undirected edges just make the directed
    flag false
  . Complexity: O(min(E^2 *V log V, E logV *V))
    flow))*/
using T = long long;
const T inf = 1LL << 61;
struct MCMF {
  struct edge {
    int u, v;
    T cap, cost;
                                                            }
    int id;
    edge(int _u, int _v, T _cap, T _cost, int _id)
                                                          potential
     u = _u;
     v = v;
      cap = _cap;
      cost = _cost;
      id = _id;
  };
  int n, s, t, mxid;
  T flow, cost;
  vector<vector<int>> g;
  vector<edge> e;
  vector<T> d, potential, flow_through;
  vector<int> par;
  bool neg;
                                                          return f;
  MCMF() {}
  MCMF(int _n) { // O-based indexing
```

```
g.assign(n, vector<int>());
  neg = false;
void add_edge(int u, int v, T cap, T cost, int
  id = -1, bool directed = true) {
  if (cost < 0) neg = true;
  g[u].push_back(e.size());
  e.push_back(edge(u, v, cap, cost, id));
  g[v].push_back(e.size());
  e.push_back(edge(v, u, 0, -cost, -1));
  mxid = max(mxid, id);
  if (!directed) add_edge(v, u, cap, cost, -1,
bool dijkstra() {
  par.assign(n, -1);
  d.assign(n, inf);
  priority_queue<pair<T, T>, vector<pair<T,</pre>
  T>>, greater<pair<T, T>> > q;
  q.push(pair<T, T>(0, s));
  while (!q.empty()) {
    int u = q.top().second;
    T nw = q.top().first;
    if (nw != d[u]) continue;
    for (int i = 0; i < (int)g[u].size(); i++) {
      int id = g[u][i];
      int v = e[id].v;
      T cap = e[id].cap;
      T w = e[id].cost + potential[u] -
  potential[v];
      if (d[u] + w < d[v] \&\& cap > 0) {
        d[v] = d[u] + w;
        par[v] = id;
        q.push(pair<T, T>(d[v], v));
  for (int i = 0; i < n; i++) { // update
    if (d[i] < inf) potential[i] += d[i];</pre>
  return d[t] != inf;
T send_flow(int v, T cur) {
  if (par[v] == -1) return cur;
  int id = par[v];
  int u = e[id].u;
  T w = e[id].cost;
  T f = send_flow(u, min(cur, e[id].cap));
  cost += f * w;
  e[id].cap -= f;
  e[id ^1].cap += f;
```

```
//returns {maxflow, mincost}
  pair<T, T> solve(int _s, int _t, T goal = inf) {
   t = _t;
    flow = 0, cost = 0;
    potential.assign(n, 0);
    if (neg) {
      // run Bellman-Ford to find starting
    potential
      d.assign(n, inf);
      for (int i = 0, relax = true; i < n &&
    relax; i++) {
        for (int u = 0; u < n; u++) {
          for (int k = 0; k < (int)g[u].size();
    k++) {
            int id = g[u][k];
            int v = e[id].v;
            T cap = e[id].cap, w = e[id].cost;
            if (d[v] > d[u] + w && cap > 0) {
              d[v] = d[u] + w;
              relax = true;
       }
      for (int i = 0; i < n; i++) if (d[i] < inf)
    potential[i] = d[i];
    while (flow < goal && dijkstra()) flow +=
    send_flow(t, goal - flow);
    flow_through.assign(mxid + 10, 0);
    for (int u = 0; u < n; u++) {
     for (auto v : g[u]) {
        if (e[v].id >= 0) flow_through[e[v].id] =
    e[v ^ 1].cap;
    return make_pair(flow, cost);
};
4 Game Theory
4.1 Points to be noted [14 lines]
>[First Write a Brute Force solution]
```

```
>Nim = all xor
>Misere Nim = Nim + corner case: if all piles are
   1, reverse(nim)
>Bogus Nim = Nim
>Staircase Nim = Odd indexed pile Nim (Even
   indexed pile doesnt matter, as one player can
```

give bogus moves to drop all even piles to

ground)
>Sprague Grundy: [Every impartial game under the
normal play convention is equivalent to a
one-heap game of nim]

```
Every tree = one nim pile = tree root value; tree
    leaf value = 0; tree node value = mex of all
    child nodes.
[Careful: one tree node can become multiple new
    tree roots(multiple elements in one node),
    then the value of that node = xor of all those
    root values]
>Hackenbush(Given a rooted tree; cut an edge in
    one move; subtree under that edge gets
    removed; last player to cut wins):
Colon: //G(u) = (G(v1) + 1) \oplus (G(v2) + 1) \oplus \cdots [v1, v2, \cdots]
    are childs of u]
For multiple trees ans is their xor
>Hackenbush on graph (instead of tree given an
    rooted graph):
fusion: All edges in a cycle can be fused to get a
    tree structure; build a super node, connect
    some single nodes with that super node, number
    of single nodes is the number of edges in the
Sol: [Bridge component tree] mark all bridges, a
    group of edges that are not bridges, becomes
    one component and contributes number of edges
    to the hackenbush. (even number of edges
    contributes 0, odd number of edges contributes
    1)
5 Geometry
5.1 Rotation Matrix [39 lines]
struct { double x; double y; double z; } Point;
double rMat[4][4];
double inMat[4][1] = \{0.0, 0.0, 0.0, 0.0\};
double outMat [4][1] = \{0.0, 0.0, 0.0, 0.0\};
void mulMat() {
 for(int i = 0; i < 4; i++){
   for(int j = 0; j < 1; j++){
      outMat[i][j] = 0;
      for(int k = 0; k < 4; k++)
        outMat[i][j] += rMat[i][k] * inMat[k][j];
   }
 }
void setMat(double ang, double u, double v, double
  double L = (u * u + v * v + w * w);
  ang = ang * PI / 180.0; /*converting to radian
  double u2 = u*u; double v2 = v*v; double w2 =
  rMat[0][0]=(u2+(v2+w2)*cos(ang))/L;
  rMat[0][1] = (u*v*(1-cos(ang))-w*sqrt(L)*
    sin(ang))/L;
  rMat[0][2] = (u*w*(1-cos(ang))+v*sqrt(L)*
    sin(ang))/L;
  rMat[0][3]=0.0;
  rMat[1][0]=(u*v*(1-cos(ang))+w*sqrt(L)*
    sin(ang))/L;
  rMat[1][1] = (v2+(u2+w2)*cos(ang))/L;
```

```
rMat[1][2]=(v*w*(1-cos(ang))-u*sqrt(L)*|
    sin(ang))/L;
  rMat[1][3]=0.0:
 rMat[2][0]=(u*w*(1-cos(ang))-v*sqrt(L)*
    sin(ang))/L;
  rMat[2][1] = (v*w*(1-cos(ang))+u*sqrt(L)*
    sin(ang))/L;
  rMat[2][2]=(w2 + (u2 + v2) * cos(ang)) / L;
 rMat[2][3]=0.0; rMat[3][0]=0.0; rMat[3][1]=0.0;
  rMat[3][2]=0.0; rMat[3][3]=1.0;
/*double ang;
  double u, v, w; //points = the point to be
    rotated
  Point point, rotated; //u,v,w=unit vector of
  inMat[0][0] = points.x; inMat[1][0] = points.y;
  inMat[2][0] = points.z; inMat[3][0] = 1.0;
  setMat(ang, u, v, w); mulMat();
  rotated.x = outMat[0][0]; rotated.y =
    outMat[1][0];
  rotated.z = outMat[2][0];*/
6 Graph
6.1 2SAT [92 lines]
struct TwoSat {
  vector<bool>vis;
  vector<vector<int>>adj, radj;
  vector<int>dfs_t, ord, par;
  int n, intime; //For n node there will be 2*n
    node in SAT.
  void init(int N) {
    n = N;
    intime = 0;
    vis.assign(N * 2 + 1, false);
    adj.assign(N * 2 + 1, vector\langle int \rangle());
    radj.assign(N * 2 + 1, vector\langle int \rangle());
    dfs_t.resize(N * 2 + 1);
    ord.resize(N * 2 + 1);
    par.resize(N * 2 + 1);
  inline int neg(int x) {
    return x \le n ? x + n : x - n;
  inline void add_implication(int a, int b) {
    if (a < 0) a = n - a:
    if (b < 0) b = n - b;
    adj[a].push_back(b);
    radj[b].push_back(a);
  inline void add_or(int a, int b) {
    add_implication(-a, b);
    add_implication(-b, a);
  inline void add_xor(int a, int b) {
    add or(a, b):
    add_or(-a, -b);
```

```
inline void add_and(int a, int b) {
  add_or(a, b);
  add_or(a, -b);
  add_or(-a, b);
inline void force_true(int x) {
  if (x < 0) x = n - x;
  add_implication(neg(x), x);
inline void add_xnor(int a, int b) {
  add or(a. -b):
  add_or(-a, b);
inline void add_nand(int a, int b) {
  add_or(-a, -b);
inline void add_nor(int a, int b) {
  add_and(-a, -b);
inline void force_false(int x) {
  if (x < 0) x = n - x;
  add_implication(x, neg(x));
inline void topsort(int u) {
  vis[u] = 1:
  for (int v : radj[u]) if (!vis[v]) topsort(v);
  dfs_t[u] = ++intime;
inline void dfs(int u, int p) {
  par[u] = p, vis[u] = 1;
  for (int v : adj[u]) if (!vis[v]) dfs(v, p);
void build() {
  int i, x;
 for (i = n * 2, intime = 0; i >= 1; i--) {
    if (!vis[i]) topsort(i);
    ord[dfs_t[i]] = i;
  vis.assign(n * 2 + 1, 0);
  for (i = n * 2; i > 0; i--) {
   x = ord[i];
    if (!vis[x]) dfs(x, x);
bool satisfy(vector<int>& ret)//ret contains the
  value that are true if the graph is
  satisfiable.
  build();
  vis.assign(n * 2 + 1, 0);
  for (int i = 1; i \le n * 2; i++) {
    int x = ord[i];
    if (par[x] == par[neg(x)]) return 0;
    if (!vis[par[x]]) {
      vis[par[x]] = 1;
      vis[par[neg(x)]] = 0;
```

```
}
    for (int i = 1;i <= n;i++) if (vis[par[i]])
    ret.push_back(i);
    return 1;
6.2 BellmanFord [57 lines]
#include <bits/stdc++.h>
using namespace std;
const int mx = 1e5+6;
const int INF = 0x3f3f3f3f;
struct edge{
 int u,v;
 int cost;
vector<edge>e;
vector<int>path(mx);
int dist[mx];
   Time-complexity: O(|V|*|E|)
   Space-complexity: O(|V|)
   To find any negative cycle assign dist[i] = 0
   We can use floyd-warshall algorithm to find
    negative cycle too.
 *Handle LL carefully.
void bellmanford(int s,int n){
  int m = e.size();
  memset(dist,0x3f3f3f3f,sizeof dist);
  dist[s] = 0;
  int x;
  for (int i = 0; i < n; ++i) {
    x = -1:
    for (int j = 0; j < m; ++j)
      if (dist[e[i].u] < INF)</pre>
        if (dist[e[j].v] > dist[e[j].u] +
    e[j].cost) {
          dist[e[j].v] = max(-INF, dist[e[j].u] +
    e[i].cost);
          path[e[j].v] = e[j].u;
          x = e[i].v;
  }
  if (x == -1)
    cout << "No negative cycle from " << s;</pre>
  else {
    for (int i = 0; i < n; ++i)
      y = path[y];
    vector<int> path;
    for (int cur = y; ; cur = path[cur]) {
      path.push_back(cur);
      if (cur == y && path.size() > 1)
        break;
```

```
reverse(path.begin(), path.end());
    cout << "Negative cycle: ";</pre>
    for (size_t i = 0; i < path.size(); ++i)</pre>
      cout << path[i] << ' ';
int main(){
6.3 BridgeTree [66 lines]
int N, M, timer, compid;
vector<pair<int, int>> g[mx];
bool used[mx], isBridge[mx];
int comp[mx], tin[mx], minAncestor[mx];
vector<int> Tree[mx]; // Store 2-edge-connected
    component tree. (Bridge tree).
void markBridge(int v, int p) {
 tin[v] = minAncestor[v] = ++timer;
  used[v] = 1;
 for (auto& e : g[v]) {
    int to, id;
    tie(to, id) = e;
    if (to == p) continue;
    if (used[to]) minAncestor[v] =
    min(minAncestor[v], tin[to]);
    else {
      markBridge(to, v);
      minAncestor[v] = min(minAncestor[v],
    minAncestor[to]);
      if (minAncestor[to] > tin[v]) isBridge[id]
    = true:
      // if (tin[u] \le minAncestor[v]) ap[u] = 1;
 }
void markComp(int v, int p) {
  used[v] = 1;
  comp[v] = compid;
  for (auto\& e : g[v]) {
    int to, id;
    tie(to, id) = e;
    if (isBridge[id]) continue;
    if (used[to]) continue;
    markComp(to, v);
vector<pair<int, int>> edges;
void addEdge(int from, int to, int id) {
  g[from].push_back({ to, id });
  g[to].push_back({ from, id });
  edges[id] = { from, to };
void initB() {
 for (int i = 0; i <= compid; ++i)
    Tree[i].clear();
 for (int i = 1; i <= N; ++i) used[i] = false;
```

```
for (int i = 1; i <= M; ++i) isBridge[i] =
    false;
  timer = compid = 0;
void bridge_tree() {
  initB();
  markBridge(1, -1); //Assuming graph is
    connected.
  for (int i = 1; i <= N; ++i) used[i] = 0;
  for (int i = 1; i <= N; ++i) {
    if (!used[i]) {
      markComp(i, -1);
      ++compid;
  for (int i = 1; i <= M; ++i) {
    if (isBridge[i]) {
      int u, v;
      tie(u, v) = edges[i];
      // connect two componets using edge.
      Tree[comp[u]].push_back(comp[v]);
      Tree[comp[v]].push_back(comp[u]);
      int x = comp[u];
      int y = comp[v];
6.4 Centroid Decomposition [49 lines]
11 n,subsize[mx];
vector<ll>adj[mx];
char ans[mx];
bool brk[mx]:
void calculatesize(ll u,ll par){
  subsize[u]=1;
  for(ll i=0;i<(ll)adj[u].size();i++){
    11 v=adj[u][i];
```

```
ll n,subsize[mx];
vector<ll>adj[mx];
char ans[mx];
bool brk[mx];
void calculatesize(ll u,ll par){
    subsize[u]=1;
    for(ll i=0;i<(ll)adj[u].size();i++){
        ll v=adj[u][i];
        if(v==par or brk[v]==true)continue;
        calculatesize(v,u);
        subsize[u]+=subsize[v];
    }
}
ll getcentroid(ll u,ll par,ll n){
    ll ret=u;
    for(ll i=0;i<(ll)adj[u].size();i++){
        ll v=adj[u][i];
        if(v==par or brk[v]==true)continue;
        if(subsize[v]>(n/2)){
            ret=getcentroid(v,u,n);
            break;
     }
    return ret;
}
void decompose(ll u,char rank){
        calculatesize(u,-1);
```

```
11 c=getcentroid(u,-1,subsize[u]);
                                                         //calculation
  for(ll i=0;i<(ll)adj[c].size();i++){
                                                        for(auto v: adj[u]){
                                                          if (v == p \mid \mid v == big) continue;
    if(brk[v]==true)continue;
                                                          for(auto x: *dsu[v]){
    decompose(v,rank+1);
                                                             dsu[u]->push_back(x);
                                                             //calculation
  for(11 i=0;i< n-1;i++){
                                                        //calculate ans for node u
    scanf("%lld %lld",&a,&b);
                                                        if(isb == 0){
                                                          for(auto x: *dsu[u]){
    adj[a].push_back(b);
                                                             //reverse calculation
    adj[b].push_back(a);
                                                        }
  for(11 i=1;i<=n;i++){
   printf("%c",ans[i]);
                                                      int main()
                                                        scanf("%d", &n);
6.5 DSU on Tree [69 lines]
                                                        for(int i=1; i<n; i++){
                                                           int u, v;
//extra data you need
                                                          scanf("%d %d", &u, &v);
vector<int> adj[mxn];
                                                          adj[u].pb(v);
vector<int> *dsu[mxn];
                                                          adj[v].pb(u);
void call(int u, int p=-1){
  for(auto v: adj[u]){
                                                        dep[1] = 1;
                                                        call(1);
      dep[v] = dep[u]+1;
                                                        dfs(1);
      sz[u] += sz[v];
                                                       6.6 Dijkstra [33 lines]
                                                       #include <bits/stdc++.h>
                                                       #define ff first
void dfs(int u, int p = -1, int isb = 1){
                                                       #define ss second
                                                      using namespace std;
  for(auto v: adj[u]){
                                                       const int mx = 1e5 + 5;
    if(v != p && sz[v]>mx){
                                                      using ll = long long;
                                                      using pll = pair<11, 11>;
                                                      vector<pll>adj[mx];
                                                      int dis[mx];
                                                      bool vis[mx];
  for(auto v: adi[u]){
                                                      //Complexity O(V+ElogV)
    if(v != p \&\& v != big){
                                                      void Dijkstra(int src) {
                                                        priority_queue<pll, vector<pll>, greater<pll>
                                                        pq.push({ 0,src });
                                                        memset(dis, 0x3f3f3f3f, sizeof dis);
                                                        memset(vis, 0, sizeof vis);
    dsu[u] = dsu[big];
                                                        dis[src] = 0;
                                                        while (!pq.empty()) {
                                                          int u = pq.top().ss;
    dsu[u] = new vector<int>():
                                                          pq.pop();
                                                          if (vis[u]) continue;
  dsu[u]->push_back(u);
                                                          vis[u] = true;
```

brk[c]=true;

ans[c]=rank;

int main(){

int n;

sz[u] = 1:

}

if(v != p){

call(v, u);

int mx=-1, big=-1;

mx = sz[v];

dfs(v, u, 0);

dfs(big, u, 1);

big = v;

if(big != -1){

else{

ll a,b;

11 v=adj[c][i];

scanf("%lld",&n);

decompose(1, 'A');

```
for (auto v : adj[u]) {
      if (dis[v.ss] > dis[u] + v.ff) {
        dis[v.ss] = dis[u] + v.ff;
        pq.push({ dis[v.ss],v.ss });
int main() {
6.7 Heavy Light Decomposition [73 lines]
/*Heavy Light Decomposition
Build Complexity O(n)
Query Complexity O(lq^2 n)
Call init()with number of nodes
It's probably for the best to not do"using
    namespace hld"*/
namespace hld {
  //N is the maximum number of nodes
  /*par, lev, size corresponds to
    parent, depth, subtree-size*/
  //head[u] is the starting node of the chain u is
  //in[u] to out[u] keeps the subtree indices
  const int N=100000+7;
  vector<int>g[N];
  int par[N],lev[N],head[N],size[N],in[N],out[N];
  int cur_pos,n;
  //returns the size of subtree rooted at u
  /*maintains the child with the largest subtree
    at the front of q[u]*/
  //WARNING: Don't change anything here specially
    with size[]if Jon Snow
  int dfs(int u,int p){
    size[u]=1,par[u]=p;
   lev[u] = lev[p] + 1;
    for(auto &v : g[u]){
      if (v==p) continue;
      size[u] += dfs(v,u);
      if(size[v]>size[g[u].front()]){
        swap(v,g[u].front());
    return size[u];
  //decomposed the tree in an array
  //note that there is no physical array here
  void decompose(int u,int p){
    in[u]=++cur_pos;
   for(auto &v : g[u]){
      if(v==p)continue;
      head[v]=(v==g[u].front()? head[u]: v);
      decompose(v,u);
```

```
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```

```
Q.pop();
  //initializes the structure with _n nodes
                                                           for(i=0;i<deg[u.v];i++){
                                                             v.v=adj[u.v][i].v;
                                                             int cost=adj[u.v][i].w+u.w;
                                                             for(v.k=u.k;v.k<K;v.k++){
                                                               if(cost==inf)break;
                                                               if(val[v.v][v.k]>cost){
                                                                 swap(cost,val[v.v][v.k]);
                                                                 v.w=val[v.v][v.k];
  //checks whether p is an ancestor of u
                                                                 Q.push(v);
                                                                 break;
    return in[p] <= in[u] and out[u] <= out[p];
  //Returns the maximum node value in the path u-v
                                                             for(v.k++;v.k<K;v.k++){
                                                               if(cost==inf)break;
                                                               if(val[v.v][v.k]>cost)swap(cost,
                                                           val[v.v][v.k]);
      ret=max(ret,seg.query(1,1,n,in[head[u]],
                                                           }
                                                       6.9 LCA [46 lines]
                                                       const int Lg = 22;
     ret=max(ret,seg.query(1,1,n,in[head[u]],
                                                       vector<int>adj[mx];
                                                       int level[mx];
                                                       int dp[Lg][mx];
                                                       void dfs(int u) {
                                                        for (int i = 1; i < Lg; i++)
    ret=max(ret,seg.query(1,1,n,in[u],in[v]));
                                                           dp[i][u] = dp[i - 1][dp[i - 1][u]];
                                                        for (int v : adj[u]) {
                                                           if (dp[0][u] == v)continue;
                                                           level[v] = level[u] + 1;
                                                           dp[0][v] = u;
                                                           dfs(v);
                                                        }
                                                       int lca(int u, int v) {
                                                         if (level[v] < level[u])swap(u, v);</pre>
int m,n,deg[MM],source,sink,K,val[MM][12];
                                                         int diff = level[v] - level[u];
                                                        for (int i = 0; i < Lg; i++)
                                                           if (diff & (1 << i))
                                                             v = dp[i][v];
                                                        for (int i = Lg - 1; i >= 0; i--)
                                                           if (dp[i][u] != dp[i][v])
                                                             u = dp[i][u], v = dp[i][v];
                                                        return u == v ? u : dp[0][u];
                                                       int kth(int u, int k) {
                                                        for (int i = Lg - 1; i >= 0; i--)
                                                           if (k \& (1 < \bar{i}))
                                                             u = dp[i][u];
                                                        return u;
                                                       //kth node from u to v. Oth is u.
                                                      int go(int u, int v, int k) {
                                                        int l = lca(u, v);
                                                        int d = level[u] + level[v] - (level[1] << 1);</pre>
```

u=Q.top();

out[u]=cur\_pos;

head[root]=root;

decompose(root,0);

11 query(int u,int v){

u=par[head[u]];

u=par[head[u]];

11 ret=-INF;

in[u])):

swap(u,v);

return ret;

};

};

struct edge{ int v,w;

}adj[MM] [500];

int v,w,k;

return w>b.w;

for(i=0;i<n;i++)

while(!Q.empty()){

struct info{

int i,j;

info u,v;

Q.push(u);

bool isances(int p,int u){

while(!isances(head[u],v)){

while(!isances(head[u],v)){

if(in[v] < in[u])swap(u,v);

//Adds val to subtree of u

void update(int u,ll val){

6.8 K'th Shortest path [40 lines]

seg.update(1,1,n,in[u],out[u],val);

bool operator<(const info &b)const{</pre>

priority\_queue<info, vector<info>>Q;

for(j=0;j<K;j++)val[i][j]=inf;

void kthBestShortestPath(){

u.v=source,u.k=0,u.w=0;

 $n=_n;$ 

cur\_pos=0; dfs(root,0);

void init(int \_n,int root=1){

```
if (level[1] + k <= level[u]) return kth(u, k);</pre>
  k -= level[u] - level[l];
 return kth(v, level[v] - level[l] - k);
   LCA(u,v) with root r:
   lca(u,v)^{l}ca(u,r)^{l}ca(v,r)
   Distance between u, v:
   level(u) + level(v) - 2*level(lca(u,v))
6.10 SACK [50 lines]
int sz[maxn]:
void getsz(int v,int p){
  sz[v]=1;
 for(auto u : g[v])
    if(u!=p){
      getsz(u,v);
      sz[v] += sz[u];
//SACK O(nlog^2n)
map<int,int>*cnt[maxn];
void dfs(int v,int p){
  int mx=-1,bigChild=-1;
  for(auto u : g[v])
  if(u!=p){
     dfs(u,v);
     if(sz[u]>mx)mx=sz[u],bigChild=u;
  if (bigChild!=-1) cnt [v]=cnt [bigChild];
  else cnt[v]=new map<int,int>();
 (*cnt[v])[col[v]]++;
  for(auto u : g[v])
  if(u!=p && u!=bigChild){
    for(auto x :*cnt[u])
      (*cnt[v])[x.first]+=x.second;
//SACK-O(nlogn)
vector<int>*vec[maxn];
int cnt[maxn];
void dfs(int v,int p,bool keep){
  int mx=-1,bigChild=-1;
  for(auto u : g[v])
  if(u!=p&&sz[u]>mx)mx=sz[u],bigChild=u;
  for(auto u : g[v])
  if (u!=p \&\& u!=bigChild)dfs(u,v,0);
  if(bigChild!=-1)
    dfs(bigChild, v, 1), vec[v] = vec[bigChild];
  else vec[v]=new vector<int>();
  vec[v]->push_back(v);cnt[col[v]]++;
  for(auto u : g[v])
  if(u!=p && u!=bigChild)
     for(auto x :*vec[u]){
       cnt[col[x]]++;
```

assert(k <= d);</pre>

```
vec[v]->push_back(x);
/*in this step*vec[v]contains all of the subtree
    of vertex v.*/
  if(keep==0)
    for(auto u:*vec[v])cnt[col[u]]--;
6.11 SCC [43 lines]
/*components: number of SCC.
sz: size of each SCC.
comp: component number of each node.
Create reverse graph.
Run find_scc() to find SCC.
Might need to create condensation graph by
    create condensed().
Think about indeg/outdeg
for multiple test cases- clear
    adj/radj/comp/vis/sz/topo/condensed.*/
vector<int>adj[mx], radj[mx];
int comp[mx], vis[mx], sz[mx], components;
vector<int>topo;
void dfs(int u) {
  vis[u] = 1;
  for (int v : adj[u])
    if (!vis[v]) dfs(v);
  topo.push_back(u);
void dfs2(int u, int val) {
  comp[u] = val:
  sz[val]++;
  for (int v : radj[u])
    if (comp[v] == -1)
      dfs2(v, val);
void find_scc(int n) {
  memset(vis, 0, sizeof vis);
  memset(comp, -1, sizeof comp);
  for (int i = 1; i \le n; i++)
    if (!vis[i])
      dfs(i);
  reverse(topo.begin(), topo.end());
  for (int u : topo)
    if (comp[u] == -1)
      dfs2(u, ++components);
vector<int>condensed[mx];
void create_condensed(int n) {
  for (int i = 1; i \le n; i++)
    for (int v : adj[i])
      if (comp[i] != comp[v])
        condensed[comp[i]].push_back(comp[v]);
}
7 Math
7.1 Big Sum [13 lines]
ll bigsum(ll a, ll b, ll m) {
  if (b == 0) return 0;
```

```
ll sum; a %= m;
  if (b & 1) {
    sum = bigsum((a * a) % m, (b - 1) / 2, m);
    sum = (sum + (a * sum) % m) % m;
    sum = (1 + (a * sum) % m) % m;
 } else {
    sum = bigsum((a * a) % m, b / 2, m);
    sum = (sum + (a * sum) % m) % m;
 return sum;
7.2 CRT [52 lines]
11 ext_gcd(l1 A, l1 B, l1* X, l1* Y) {
 ll x2, y2, x1, y1, x, y, r2, r1, q, r;
 x2 = 1; y2 = 0;
 x1 = 0; y1 = 1;
 for (r2 = A, r1 = B; r1 != 0; r2 = r1, r1 = r,
   x2 = x1, y2 = y1, x1 = x, y1 = y) {
    q = r2 / r1;
   r = r2 \% r1;
   x = x2 - (q * x1);
   y = y2 - (q * y1);
  *X = x2; *Y = y2;
 return r2;
/*----*/
class ChineseRemainderTheorem {
 typedef long long vlong;
 typedef pair<vlong, vlong> pll;
  /** CRT Equations stored as pairs of vector. See
    addEgation()*/
 vector<pll> equations;
 public:
 void clear() {
    equations.clear();
  /** Add equation of the form x = r \pmod{m}*/
  void addEquation(vlong r, vlong m) {
    equations.push_back({ r, m });
 pll solve() {
    if (equations.size() == 0) return \{-1,-1\};
    /// No equations to solve
    vlong a1 = equations[0].first;
    vlong m1 = equations[0].second;
    a1 \%= m1;
    /** Initially x = a_0 \pmod{m_0}*/
    /** Merge the solution with remaining
    equations */
    for (int i = 1; i < equations.size(); i++) {</pre>
      vlong a2 = equations[i].first;
      vlong m2 = equations[i].second;
      vlong g = \_gcd(m1, m2);
     if (a1 % g != a2 % g) return { -1,-1 }; ///
    Conflict in equations
      /** Merge the two equations*/
      vlong p, q;
```

```
ext_gcd(m1 / g, m2 / g, &p, &q);
      vlong mod = m1 / g * m2;
      vlong x = ((_int128)a1 * (m2 / g) \% mod * q
    % mod + (__int128)a2 * (m1 / g) % mod * p %
    mod) % mod;
      /** Merged equation*/
      a1 = x;
      if (a1 < 0) a1 += mod;
      m1 = mod:
    return { a1, m1 };
7.3 FFT [85 lines]
template<typename float_t>
struct mycomplex {
  float_t x, y;
  mycomplex<float_t>(float_t _x = 0, float_t _y =
    0) : x(_x), y(_y) {}
  float_t real() const { return x; }
  float_t imag() const { return y; }
  void real(float_t _x) { x = _x; }
  void imag(float_t _y) { y = _y; }
  mycomplex<float_t>& operator+=(const
    mycomplex<float_t> &other) { x += other.x; y
    += other.y; return *this; }
  mycomplex<float_t>& operator-=(const
    mycomplex<float_t> &other) { x -= other.x; y
    -= other.y; return *this; }
  mycomplex<float_t> operator+(const
    mycomplex<float_t> &other) const { return
    mycomplex<float_t>(*this) += other; }
  mycomplex<float_t> operator-(const
    mycomplex<float_t> &other) const { return
    mycomplex<float_t>(*this) -= other; }
  mycomplex<float_t> operator*(const
    mycomplex<float_t> &other) const {
    return {x * other.x - y * other.y, x * other.y
    + other.x * y};
  mycomplex<float_t> operator*(float_t mult) const
   return {x * mult, y * mult};
  friend mycomplex<float_t> conj(const
    mycomplex<float_t> &c) {
    return \{c.x, -c.y\};
  friend ostream& operator << (ostream & stream,
    const mycomplex<float_t> &c) {
    return stream << '(' << c.x << ", " << c.v <<
    ')';
using cd = mycomplex<double>;
```

```
void fft(vector<cd> & a, bool invert) {
  int n = a.size();
  for (int i = 1, j = 0; i < n; i++) {
    int bit = n \gg 1;
   for (; j & bit; bit >>= 1)
      j ^= bit;
    j ^= bit;
    if (i < j)
      swap(a[i], a[j]);
  for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 : 1);
    cd wlen(cos(ang), sin(ang));
    for (int i = 0; i < n; i += len) {
      cd w(1);
      for (int j = 0; j < len / 2; j++) {
        cd u = a[i+j], v = a[i+j+len/2] * w;
        a[i+j] = u + v;
        a[i+j+len/2] = u - v;
        w = w*wlen;
     }
   }
  if (invert) {
   for (cd \& x : a){
     double z = n:
     z=1/z;
     x = x*z;
   // x /= n;
void multiply (const vector<bool> & a, const
    vector<bool> & b, vector<bool> & res)
    {//change all the bool to your type needed
  vector<cd> fa (a.begin(), a.end()), fb
    (b.begin(), b.end());
  size_t n = 1;
  while (n < max (a.size(), b.size())) n <<= 1;
  fa.resize (n), fb.resize (n);
  fft (fa, false), fft (fb, false);
  for (size_t i=0; i<n; ++i)
   fa[i] =fa[i] * fb[i];
  fft (fa, true);
  res.resize (n);
  for (size_t i=0; i<n; ++i)
    res[i] = round(fa[i].real());
  while(res.back()==0) res.pop_back();
void pow(const vector<bool> &a, vector<bool>
    &res, long long int k){
  vector<bool> po=a;
  res.resize(1);
  res[0] = 1;
  while(k){
    if(k&1){
      multiply(po, res, res);
```

```
multiply(po, po, po);
   k/=2;
7.4 GaussElimination [39 lines]
template<typename ld>
int gauss(vector<vector<ld>>& a, vector<ld>& ans)
  const ld EPS = 1e-9;
 int n = a.size();///number of equations
 int m = a[0].size() - 1;///number of variables
 vector<int>where(m, -1);///indicates which row
    contains the solution
 int row, col;
 for (col = 0, row = 0; col < m && row < n; ++col)
    int sel = row;///which row contains the
    maximum value/
    for (int i = row + 1; i < n; i++)
      if (abs(a[i][col]) > abs(a[sel][col]))
    if (abs(a[sel][col]) < EPS) continue; ///it's
    basically 0.
    a[sel].swap(a[row]);///taking the max row up
    where[col] = row;
    ld t = a[row][col];
    for (int i = col; i <= m; i++) a[row][i] /= t;
    for (int i = 0; i < n; i++) {
      if (i != row) {
       ld c = a[i][col];
       for (int j = col; j <= m; j++)
          a[i][j] -= a[row][j] * c;
     }
    }
    row++;
 ans.assign(m, 0);
 for (int i = 0; i < m; i++)
    if (where[i] != -1)
      ans[i] = a[where[i]][m] / a[where[i]][i];
 for (int i = 0; i < n; i++) {
    ld sum = 0;
    for (int j = 0; j < m; j++)
      sum += ans[j] * a[i][j];
    if (abs(sum - a[i][m]) > EPS) ///L.H.S!=R.H.S
      ans.clear();//No solution
 return row;
7.5 GaussMod2 [44 lines]
template<typename T>
struct Gauss {
 int bits = 60;
 vector<T>table:
 Gauss() {
   table = vector<T>(bits, 0);
```

```
//call with constructor to define bit size.
  Gauss(int _bits) {
    bits = _bits;
    table = vector<T>(bits, 0);
  int basis()//return rank/size of basis
    int ans = 0;
    for (int i = 0;i < bits;i++)
      if (table[i])
        ans++;
    return ans;
  bool can(T x)//can x be obtained from the basis
    for (int i = bits - 1; i >= 0; i--) x = min(x, x)
    ^ table[i]);
    return x == 0;
  void add(T x) {
    for (int i = bits - 1; i >= 0 && x; i--) {
      if (table[i] == 0) {
        table[i] = x;
        x = 0;
      else x = min(x, x \hat table[i]);
  T getBest() {
    \bar{T} x = 0;
    for (int i = bits - 1; i >= 0; i--)
      x = max(x, x ^ table[i]);
    return x;
  void Merge(Gauss& other) {
    for (int i = bits - 1; i \ge 0; i--)
    add(other.table[i]);
};
7.6 Karatsuba Idea [5 lines]
Three subproblems:
a = xH yH
d = xL yL
e = (xH + xL)(yH + yL) - a - d
Then xy = a rn + e rn/2 + d
```

```
7.7 Linear Diophatine [12 lines]
/*x'=x+(k*B/q), y'=y-(k*A/q); infinite soln
if A=B=0, C must equal O and any x,y is solution;
if A/B=0, (x,y)=(C/A,k)/(k,C/B)*/
bool LDE(int A,int B,int C,int*x,int*y){
  int g=gcd(A,B);
  if(C%g!=0)return false;
  int a=A/g, b=B/g, c=C/g;
  extended_gcd(a,b,x,y); //ax+by=1
```

```
if(g<0)a*=-1;b*=-1;c*=-1;//Ensure\ qcd(a,b)=1
 *x*=c; *y*=c; //ax+by=c
  return true://Solution Exists
7.8 Matrix [100 lines]
template<typename T>
struct Matrix {
  T MOD = 1e9 + 7; /// change if necessary
  T add(T a, T b) const {
    T res = a + b;
    if (res >= MOD) return res - MOD;
    return res;
  T sub(T a, T b) const {
    T res = a - b;
    if (res < 0) return res + MOD;
    return res;
  T mul(T a, T b) const {
   T res = a * b;
    if (res >= MOD) return res % MOD;
    return res;
  int R, C;
  vector<vector<T>>mat:
  Matrix(int _R = 0, int _C = 0) {
    R = R, C = C;
   mat.resize(R);
    for (auto& v : mat) v.assign(C, 0);
  void print() {
   for (int i = 0; i < R; i++)
      for (int j = 0; j < C; j++)
        cout << mat[i][j] << " \n"[j == C - 1];
  void createIdentity() {
   for (int i = 0; i < R; i++)
      for (int j = 0; j < C; j++)
        mat[i][j] = (i == j);
  Matrix operator+(const Matrix& o) const {
    Matrix res(R, C);
   for (int i = 0; i < R; i++)
      for (int j = 0; j < C; j++)
        res[i][j] = add(mat[i][j] + o.mat[i][j]);
  Matrix operator-(const Matrix& o) const {
    Matrix res(R, C);
   for (int i = 0; i < R; i++)
      for (int j = 0; j < C; j++)
        res[i][j] = sub(mat[i][j] + o.mat[i][j]);
  Matrix operator*(const Matrix& o) const {
    Matrix res(R, o.C);
   for (int i = 0; i < R; i++)
     for (int j = 0; j < o.C; j++)
        for (int k = 0; k < C; k++)
```

```
res.mat[i][j] = add(res.mat[i][j],
    mul(mat[i][k], o.mat[k][j]));
    return res;
 Matrix pow(long long x) {
    Matrix res(R, C);
    res.createIdentity();
    Matrix<T> o = *this;
    while (x) {
     if (x \& 1) res = res * o;
     0 = 0 * 0:
     x >>= 1;
   return res;
 Matrix inverse()///Only square matrix &
    non-zero determinant
    Matrix res(R, R + R);
    for (int i = 0; i < R; i++) {
     for (int j = 0; j < R; j++)
        res.mat[i][j] = mat[i][j];
     res.mat[i][R + i] = 1;
    for (int i = 0; i < R; i++) {
      ///find row 'r' with highest value at [r][i]
      int tr = i;
     for (int j = i + 1; j < R; j++)
        if (abs(res.mat[j][i]) >
    abs(res.mat[tr][i]))
          tr = j;
      ///swap the row
     res.mat[tr].swap(res.mat[i]);
      ///make 1 at [i][i]
     T val = res.mat[i][i];
     for (int j = 0; j < R + R; j++) res.mat[i][j]
      ///eliminate [r][i] from every row except i.
     for (int j = 0; j < R; j++) {
       if (j == i) continue;
        for (int k = R + R - 1; k >= i; k--) {
          res.mat[j][k] -= res.mat[i][k] *
    res.mat[j][i] / res.mat[i][i];
    Matrix ans(R, R);
    for (int i = 0; i < R; i++)
     for (int j = 0; j < R; j++)
        ans.mat[i][j] = res.mat[i][R + j];
    return ans;
7.9 Miller-Rabin-Pollard-Rho [68 lines]
ll powmod(ll a, ll p, ll m) {///(a^p \% m)}
 11 result = 1:
 a \%= m;
 while (p) {
```

```
if (p & 1)
      result = (vll)result * a % m;
    a = (vll)a * a % m;
   p >>= 1;
 return result;
bool check_composite(ll n, ll a, ll d, int s) {
  ll x = powmod(a, d, n);
  if (x == 1 | | x == n - 1)
   return false:
 for (int r = 1; r < s; r++) {
   x = (vll)x * x % n;
   if (x == n - 1)
      return false;
 return true;
bool MillerRabin(ll n) {
  if (n < 2) return false;
  int r = 0;
 11 d = n - 1;
  while ((d \& 1) == 0) {
   d >>= 1;
   r++:
  for (int a : {2, 3, 5, 7, 11, 13, 17, 19, 23,
    29, 31, 37}) {
    if (n == a) return true;
    if (check_composite(n, a, d, r))
      return false;
 return true;
11 mult(11 a, 11 b, 11 mod) {
 return (vll)a * b % mod;
ll f(ll x, ll c, ll mod) {
  return (mult(x, x, mod) + c) % mod;
ll rho(ll n) {
  if (n \% 2 == 0) return 2;
  11 x = myrand() \% n + 1, y = x, c = myrand() \% n
   + 1, g = 1;
  while (g == 1) {
   x = f(x, c, n);
   y = f(y, c, n);
   y = f(y, c, n);
    g = \_gcd(abs(x - y), n);
 return g;
set<ll>prime;
void prime_factorization(ll n) {
  if (n == 1) return;
 if (MillerRabin(n)) {
    prime.insert(n);
```

```
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```

```
NTT(int mod, int root, int inv, int pw) :
                                                          mod(mod), root(root), inv(inv), pw(pw) {}
                                                        void precalculate() {
                                                          perm.resize(N);
                                                          perm[0] = 0;
                                                          for (int k=1; k<N; k<<=1) {
                                                            for (int i=0; i<k; i++) {
                                                              perm[i] <<= 1;
//call prime_factorization(n) for prime factors.
                                                              perm[i+k] = 1 + perm[i];
                                                        void fft(vector<ll> &v, bool invert = false) {
                                                          if (v.size() != perm.size()) {
                                                            N = v.size();
                                                            assert(N && (N&(N-1)) == 0);
                                                            precalculate();
                                                          for (int i=0; i<N; i++)
                                                            if (i < perm[i])
                                                              swap(v[i], v[perm[i]]);
                                                          for (int len = 2; len <= N; len <<=1) {
                                                            11 factor = invert ? inv: root;
                                                            for (int i=len; i<pw; i<<=1)
                                                              factor = (factor * factor) % mod;
                                                            for (int i=0; i<N; i+=len) {
                                                              11 w = 1;
                                                              for (int j=0; j<len/2; j++) {
                                                                11 x = v[i+j], y = (w*v[i+j+len/2]) mod;
                                                                v[i+j] = (x+y) \% mod;
                                                                v[i+j+len/2] = (x-y+mod)\%mod;
                                                                w = (w*factor)%mod;
                                                          if (invert) {
                                                            ll n1 = power(N, mod-2, mod);
                                                            for (11 &x: v) x = (x*n1) \text{mod};
                                                        }
                                                        vector<11> multiply(vector<11> a, vector<11> &b)
                                                          while (a.size() && a.back() == 0)
int nttdata(int mod, int &root, int &inv, int &pw)
                                                          a.pop_back();
                                                          while (b.size() \&\& b.back() == 0)
                                                          b.pop_back();
                                                          int n = 1;
                                                          while (n < a.size() + b.size()) n <<=1;
                                                          a.resize(n);
                                                          b.resize(n);
                                                          fft(a);
                                                          for (int i=0; i<n; i++) a[i] = (a[i] *
                                                          b[i])%M;
                                                          fft(a, true);
                                                          while (a.size() && a.back() == 0)
                                                          a.pop_back();
                                                          return a;
```

return;

while (x == n) x = rho(n);

prime\_factorization(n / x);

//call MillerRabin(n) to check if prime.

int g = gcdExt(a, m, x, y);

return (x % m + m) % m;

ll power(ll a, ll p, ll mod) {

ll ans = power(a, p/2, mod);

if (p==0) return 1;

ans = (ans \* ans) % mod;

int primitive\_root(int p) {

int phi = p-1, n = phi;

factor.push\_back(i);

while (n%i==0) n/=i;

for (int i=2; i\*i<=n; i++) {

continue;

for (int res =2; res<=p; res++) {

vector<int> factor;

bool ok = true:

if (ok) return res;

int c = 0, n = mod-1;

pw = (mod-1)/n;

const int M = 786433:

vector<int> perm;

int mod, root, inv, pw;

while (n%2==0) c++, n/=2;

root = power(g, n, mod);

int g = primitive\_root(mod);

inv = power(root, mod-2, mod);

int x, y; //if q==1 Inverse doesn't exist

ans = (ans \* a) % mod;

factor.push\_back(n);

for (int i=0; i<factor.size() && ok; i++)

ok &= power(res, phi/factor[i], p) != 1;

prime\_factorization(x);

7.10 Mod Inverse [5 lines]

7.11 NTT [96 lines]

if(p%2)

return ans;

if (n%i)

if (n>1)

return -1;

return c;

struct NTT {

int N;

 $NTT()\{\}$ 

int modInv(int a, int m) {

11 x = n;

```
//
         int mod=786433, root, inv, pw;
 //
         nttdata(mod, root, inv, pw);
 //
         NTT \ nn = NTT(mod, root, inv, pw);
};
7.12 No of Digits in n! in base B [7 lines]
11 NoOfDigitInNFactInBaseB(11 N,11 B){
  ll i;
  double ans=0;
  for(i=1;i<=N;i++)ans+=log(i);
  ans=ans/log(B),ans=ans+1;
  return(11)ans;
7.13 SOD Upto N [16 lines]
11 SOD_UpTo_N(11 N){
  ll i,j,ans=0;///upto N in Sqrt(N)
  for(i=1;i*i<=N;i++){
    j=N/i;
    ans+=((j*(j+1))/2)-(((i-1)*i)/2);
    ans+=((j-i)*i);
  return ans;
11 SODUptoN(11 N){
 11 res=0,u=sqrt(N);
  for(ll i=1;i<=u;i++)
    res+=(N/i)-i;
  res*=2,res+=u;
 return res;
7.14 Sieve-Phi-Mobius [26 lines]
const int N = 1e7;
vector<int>pr;
int mu[N + 1], phi[N + 1], lp[N + 1];
void sieve() {
    phi[1] = 1, mu[1] = 1;
    for (int i = 2; i <= N; i++) {
        if (lp[i] == 0) {
            lp[i] = i;
            phi[i] = i - 1;
            pr.push_back(i);
        for (int j = 0; j < pr.size() && i * pr[j]
    <= N; j++) {
            lp[i * pr[j]] = pr[j];
            if (i % pr[j] == 0) {
                phi[i * pr[j]] = phi[i] * pr[j];
                break:
            else
                phi[i * pr[j]] = phi[i] *
    phi[pr[j]];
    for (int i = 2; i \le N; i++) {
        if (lp[i / lp[i]] == lp[i]) mu[i] = 0;
```

```
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```

```
8.1 Bit hacks [12 lines]
# x & -x is the least bit in x.
# iterate over all the subsets of the mask
for (int s=m; ; s=(s-1)\&m) {
 ... you can use s ...
 if (s==0) break;
# c = x\&-x, r = x+c; (((r^x) >> 2)/c) | r is the
next number after x with the same number of bits
# __builtin_popcount(x) //number of ones in binary
  __builtin_popcountll(x) // for long long
# __builtin_clz(x) // number of leading zeros
  __builtin_ctz(x) // number of trailing zeros,
      they also have long long version
8.2 Bitset C++ [13 lines]
bitset<17>BS;
BS[1] = BS[7] = 1:
cout<<BS._Find_first()<<endl; // prints 1</pre>
bs._Find_next(idx). This function returns first
    set bit after index idx.for example:
bitset<17>BS;
BS[1] = BS[7] = 1;
cout<<BS._Find_next(1)<<','<<BS._Find_next(3)</pre>
    <<endl; // prints
    7,7
So this code will print all of the set bits of BS:
for(int i=BS._Find_first();i< BS.size();i =</pre>
    BS._Find_next(i))
    cout << i << endl;
//Note that there isn't any set bit after idx,
    BS._Find_next(idx) will return BS.size(); same
    as calling BS._Find_first() when bitset is
    clear:
8.3 Template [30 lines]
// #pragma GCC optimize("03,unroll-loops")
// #pragma GCC
    target("avx2, bmi, bmi2, lzcnt, popcnt")
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <typename A, typename B> ostream&
    operator << (ostream& os, const pair <A, B>& p) {
    return os << '(' << p.first << ", " <<
    p.second << ')'; }
```

else mu[i] = -1 \* mu[i / lp[i]];

}

8 Misc

```
template <typename T_container, typename T =
    typename enable_if<!is_same<T_container,
    string>::value, typename
    T_container::value_type>::type> ostream&
    operator << (ostream & os, const T_container & v)
    { os << '{'; string sep; for (const T& x : v)
    os << sep << x, sep = ", "; return os << '}';
void dbg_out() { cerr << endl; }</pre>
template <typename Head, typename... Tail> void
    dbg_out(Head H, Tail... T) { cerr << " " <<</pre>
    H; dbg_out(T...); }
#ifdef SMIE
#define debug(args...) cerr << "(" << #args <<
    "):", dbg_out(args)
#else
#define debug(args...)
#endif
template <typename T> inline T gcd(T a, T b) { T
    c; while (b) { c = b; b = a \% b; a = c; }return
    a; } // better than __gcd
ll powmod(ll a, ll b, ll MOD) { ll res = 1;a %=
    MOD; assert(b >= 0); for (; b; b >>= 1) { if (b
    & 1)res = res * a % MOD; a = a * a % MOD;
    }return res: }
template <typename T>using orderedSet = tree<T,
    null_type, less_equal<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
//order_of_key(k) - number of element strictly
    less than k
//find_by_order(k) - k'th element in set. (0)
    indexed)(iterator)
using ll = long long;
mt19937
rng(chrono::steady_clock::now().time_since_epoch()
//uniform_int_distribution<int>(0, i)(rnq)
int main() {
  ios_base::sync_with_stdio(false);//DON'T CC++
  cin.tie(NULL);//DON'T use for interactive
8.4 Vimrc [26 lines]
filetype plugin indent on
set hls is ar ai cul wrap lbr nu rnu et magic sc
    aw sb spr gd so=2 tm=400 mouse=a sw=4 sts=4
    ts=4 ls=2 bs=indent,eol,start
au TerminalOpen * setlocal nonu nornu
au BufNewFile *.cpp -r ./template.cpp | 14
let mapleader = " "
imap jk <esc>
imap <c-h> <left>
noremap <c-c> "+y
nnoremap <leader>f :let @+ = expand('\%:p')<cr>
noremap <leader>d "_d
```

```
nnoremap <leader>ca :%y+<cr>
nnoremap <leader>cd :cd %:h<cr>
nnoremap <leader>b :ls<cr>:b
nnoremap <leader>vim :vs ~/.vimrc<cr>
nnoremap <leader>r %x<c-o>x
inoremap {<cr>> {<cr>>}<esc>0
inoremap ( <c-]>()<left>
inoremap (( (
vnoremap ( <esc>`>a)<esc>`<i(<esc>
inorea aLL <esc>yT(f.abegin(), <c-r>".end()
inorea raLL <esc>yT(f.arbegin(), <c-r>".rend()
8.5 build [3 lines]
#!/bin/bash
>&2 echo -e "Compiling $1.cpp with -std=gnu++17
g++ -std=gnu++17 -Wshadow -Wall -Wextra
    -Wno-unused-result -02 -g -fsanitize=undefined
    -fsanitize=address $2 "$1.cpp" -o "$1"
8.6 debug [6 lines]
#!/bin/bash
bar="----"
build $1 -DSMIE && >&2 echo -e "\tRunning
    $1\n$bar" && ts=$(date +\%s\%N) &&
\time -f "$bar\nMemory Usage : %M KB" "./$1" &&
>&2 echo -e "Execution Time : $((($(date + %s%N) -
    $ts)/1000000)) ms"
8.7 stress-test [16 lines]
#!/bin/bash
build $1 $2 && build $1_gen $2 && build $1_brute
    $2 &&
for((i = 1; ++i)); do
    echo -e "\nTest Case "$i
    ./$1_gen $i > inp
    echo -e "=======\nINPUT\n----"
    cat inp
    ./$1 < inp > out1
    echo -e "\nOUTPUT\n----"
    cat out1
    ./$1_brute < inp > out2
    echo -e "\nEXPECTED\n----"
    cat out2
    echo ""
    diff -w out1 out2 || break
done
9 String
9.1 Aho-Corasick [124 lines]
const int NODE=3000500;//Maximum Nodes
const int LGN=30;
                     ///Maximum Number of Tries
const int MXCHR=53; ///Maximum Characters
const int MXP=5005;
                     ///
struct node {
  int val;
  int child[MXCHR];
```

```
int maxNodeId,fail[NODE+10],par[NODE+10];
int nodeSt[NODE+10],nodeEd[NODE+10];
vlong csum[NODE+10],pLoc[MXP];
 int curNodeId=++maxNodeId;
 Trie[curNodeId].clear():
inline void upd(vlong pos){
inline vlong qry(vlong pos){
  inline int getname(char ch){
    else if(ch \ge A' \&\& ch \le Z')return 26+(ch - A');
  void addToTrie(string &s,int id){
  //Add string s to the Trie in general way
    int len=SZ(s), cur=root;
      int c=getname(s[i]);
     if(Trie[cur].child[c]==0){
        int curNodeId=getNode();
        Trie[curNodeId].val=c;
        Trie[cur].child[c]=curNodeId;
      cur=Trie[cur].child[c];
  void calcFailFunction(){
    //Add all the children to the queue:
```

vector<int>graph;

CLR(child,0);

graph.clear();

void clear(){

val=0;

}Trie[NODE+10];

void resetTrie(){

return curNodeId;

vlong res=csum[pos];

int root, size, euler;

root=getNode(); size=euler=0;

FOR(i,0,len-1){

pLoc[id]=cur;

queue<int>Q;

Q.pop();

Q.push(root);

while(!Q.empty()){

int s=Q.front();

size++:

if(ch=='-')return 52;

else return(ch-'a');

struct AhoCorasick {

void clear(){

maxNodeId=0;

csum[pos]++;

return res;

int getNode(){

```
FOR(i,0,MXCHR-1){
       int t=Trie[s].child[i];
       if(t!=0){
         Q.push(t);
         par[t]=s;
     if(s==root){/*Handle special case when s is
       fail[s]=par[s]=root;
       continue;
//Find fall back of s:
     int p=par[s],f=fail[p];;
     int val=Trie[s].val;
/*Fall back till you found a node who has got val
   as a child*/
     while(f!=root && Trie[f].child[val]==0){
       f=fail[f]:
    fail[s]=(Trie[f].child[val]==0)? root :
   Trie[f].child[val];
//Self fall back not allowed
     if(s==fail[s]){
       fail[s]=root;
     Trie[fail[s]].graph.push_back(s);
 void dfs(int pos){
   ++euler;
   nodeSt[pos] = euler;
   for(auto x: Trie[pos].graph){
     dfs(x);
   nodeEd[pos] = euler;
//Returns the next state
 int goTo(int state,int c){
   if(Trie[state].child[c]!=0){/*No need to fall
     return Trie[state].child[c];
 //Fall back now:
   int f=fail[state];
   while(f!=root && Trie[f].child[c]==0){
     f=fail[f];
   int res=(Trie[f].child[c]==0)?
   root:Trie[f].child[c];
   return res;
/*Iterate through the whole text and find all the
   matchings*/
 void findmatching(string &s){
   int cur=root,idx=0;
   int len=SZ(s);
   while(idx<len){
```

```
}acorasick;
9.2 Double Hasing [50 lines]
struct SimpleHash {
  int len:
  long long base, mod;
  vector<int> P, H, R;
  SimpleHash() {}
  SimpleHash(const char* str, long long b, long
    long m) {
    base = b, mod = m, len = strlen(str);
    P.resize(len + 4, 1), H.resize(len + 3, 0),
    R.resize(len + 3, 0);
    for (int i = 1; i <= len + 3; i++)
      P[i] = (P[i - 1] * base) \% mod;
    for (int i = 1; i <= len; i++)
      H[i] = (H[i - 1] * base + str[i - 1] + 1007)
    % mod:
    for (int i = len; i >= 1; i--)
      R[i] = (R[i + 1] * base + str[i - 1] + 1007)
    % mod;
  inline int range_hash(int 1, int r) {
    int hashval = H[r + 1] - ((long long)P[r - 1]
    + 1] * H[1] % mod);
    return (hashval < 0 ? hashval + mod :
   hashval);
 }
  inline int reverse_hash(int 1, int r) {
    int hashval = R[1 + 1] - ((long long)P[r - 1]
    + 1] * R[r + 2] % mod);
    return (hashval < 0 ? hashval + mod :
    hashval):
 }
struct DoubleHash {
  SimpleHash sh1, sh2;
  DoubleHash() {}
  DoubleHash(const char* str) {
    sh1 = SimpleHash(str, 1949313259, 2091573227);
    sh2 = SimpleHash(str, 1997293877, 2117566807);
  long long concate(DoubleHash& B , int 11 , int
    r1 , int 12 , int r2) {
    int len1 = r1 - 11+1, len2 = r2 - 12+1;
    long long x1 = sh1.range_hash(l1, r1) ,
         x2 = B.sh1.range_hash(12, r2);
    x1 = (x1 * B.sh1.P[len2]) \% 2091573227;
    long long newx1 = (x1 + x2) \% 2091573227;
    x1 = sh2.range_hash(l1, r1);
    x2 = B.sh2.range_hash(12, r2);
```

int c=getname(s[idx]);

cur=goTo(cur,c);

upd(nodeSt[cur]);

idx++;

```
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```

```
x1 = (x1 * B.sh2.P[len2]) \% 2117566807;
                                                        link[idx]=len[idx]==1?2:tree[x][c];
  long long newx2 = (x1 + x2) \% 2117566807;
                                                      t=tree[t][c];
inline long long range_hash(int 1, int r) {
                                                    void build(){
  return ((long long)sh1.range_hash(1, r) << 32)
                                                      len[1]=-1, link[1]=1;
                                                      len[2]=0, link[2]=1;
                                                      idx=t=2;
inline long long reverse_hash(int 1, int r) {
                                                      for(int i=1;i<n;i++) extend(i);</pre>
  return ((long long)sh1.reverse_hash(l, r) <<
                                                  };
                                                   9.5 Suffix Array [78 lines]
                                                   struct SuffixArray {
                                                   vector<int> p, c, rank, lcp;
                                                   vector<vector<int>> st;
                                                   SuffixArray(string const& s) {
                                                    build_suffix(s + char(1));
                                                    p.erase(p.begin());
                                                    build_rank(p.size());
                                                    build_lcp(s);
                                                    build_sparse_table(lcp.size());
                                                   void build_suffix(string const& s) {
                                                    int n = s.size();
                                                    const int MX_ASCII = 256;
                                                    vector<int> cnt(max(MX_ASCII, n), 0);
                                                    p.resize(n); c.resize(n);
                                                    for (int i = 0; i < n; i++) cnt[s[i]]++;
                                                    for (int i=1; i<MX_ASCII; i++) cnt[i]+=cnt[i-1];
                                                    for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;
                                                    c[p[0]] = 0;
                                                    int classes = 1;
                                                    for (int i = 1; i < n; i++) {
                                                      if (s[p[i]] != s[p[i-1]]) classes++;
                                                      c[p[i]] = classes - 1;
                                                    vector<int> pn(n), cn(n);
                                                    for (int h = 0; (1 << h) < n; ++h) {
                                                      for (int i = 0; i < n; i++) {
                                                        pn[i] = p[i] - (1 << h);
                                                        if (pn[i] < 0) pn[i] += n;
                                                      fill(cnt.begin(), cnt.begin() + classes, 0);
                                                      for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;
                                                      for (int i=1; i < classes; i++)</pre>
                                                      cnt[i]+=cnt[i-1]:
                                                      for (int i=n-1;i>=0;i--)
                                                      p[--cnt[c[pn[i]]]=pn[i];
                                                      cn[p[0]] = 0; classes = 1;
                                                      for (int i = 1; i < n; i++) {
                                                        << h)) % n]};
                                                        pair<int, int> prev = {c[p[i-1]], c[(p[i-1])
                                                      + (1 << h)) % n]};
                                                        if (cur != prev) ++classes;
                                                        cn[p[i]] = classes - 1;
                                                      c.swap(cn);
```

return (newx1 << 32) ^ newx2;

32) ^ sh2.reverse\_hash(1, r);

while(j>=0 and P[i]!=P[i])

while( $j \ge 0$  and T[i]!=P[j])

//pattern found at index i-j

tree.assign(n+5, vector<int>(26,0));

while (s[p-len[t]-1]!=s[p]) t=link[t];

while (s[p-len[x]-1]!=s[p]) x=link[x];

sh2.range\_hash(1, r);

};

9.3 KMP [23 lines]

int b[maxn],n,m;

int i=0, j=-1;

j=b[j];

i++;j++;

b[i]=j;

void kmpSearch(){

j=b[j];

i++;j++;

9.4 Palindromic Tree [30 lines]

vector<vector<int>> tree;

struct PalindromicTree{

vector<int> len,link;

string s; // 1-indexed PalindromicTree(string str){

len.assign(n+5,0);

void extend(int p){

if(!tree[t][c]){

tree[t][c]=++idx;

len[idx]=len[t]+2;

link.assign(n+5,0);

int x=link[t],c=s[p]-'a';

 $if(i==m){}$ 

int n,idx,t;

s="\$"+str:

n=s.size();

int i=0, j=0;

while(i<n){

while(i<m){

b[0] = -1:

char P[maxn],T[maxn];

void kmpPreprocess(){

```
void build_rank(int n) {
  rank.resize(n, 0);
  for (int i = 0; i < n; i++) rank[p[i]] = i;
void build_lcp(string const& s) {
  int n = s.size(), k = 0;
  lcp.resize(n - 1, 0);
  for (int i = 0; i < n; i++) {
    if (rank[i] == n - 1) {
      k = 0;
      continue;
    int j = p[rank[i] + 1];
    while (i + k < n \&\& j + k < n \&\& s[i+k] ==
    s[i+k]
      k++;
    lcp[rank[i]] = k;
    if (k) k--;
  }
void build_sparse_table(int n) {
  int lim = __lg(n);
  st.resize(lim + 1, vector<int>(n)); st[0] = lcp;
  for (int k = 1; k \le \lim_{k \to \infty} k + 1)
    for (int i = 0; i + (1 << k) <= n; i++)
      st[k][i] = min(st[k-1][i], st[k-1][i+1][i]
    (1 << (k - 1))];
int get_lcp(int i) { return lcp[i]; }
int get_lcp(int i, int j) {
  if (j < i) swap(i, j);
  j--; /*for lcp from i to j we don't need last
    lcp*/
  int K = -\lg(j - i + 1);
  return min(st[K][i], st[K][j - (1 << K) + 1]);
};
9.6 Suffix Automata [119 lines]
const int MXCHR = 26;
take an object of suffixAutomata
call extend(c) for each character c in string
```

```
call Process() to initiate the important values
*/
struct suffixAutomata {
     * len -> largest string length of the
    corresponding endpos-equivalent class
     * link -> longest suffix that is another
    endpos-equivalent\ class
     * firstpos -> end position of the first
    occurrence of the largest string of
     *that node
```

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struct state {
    int link, len;
    int next[MXCHR];
    state() {}
    state(int 1) {
        len = 1;
       link = -1;
       for (int i = 0; i < MXCHR; i++)
next[i] = -1;
   }
};
vector<state> node;
int sz, last;
vector<int> cnt, distinct, firstPos, occur,
vector<vector<int>> adj; // suffix links tree
// cnt and SA for counting sort the nodes.
int L;
suffixAutomata() {
    node.push_back(state(0));
    firstPos.push_back(-1);
    occur.push_back(0);
    last = 0:
    sz = 0:
    L = 0;
int getID(char c) {
    return c - 'a'; // change according to
void extend(char c) {
    int idx = ++sz, p = last, id = getID(c);
    L++;
    node.push_back(state(node[last].len + 1));
    firstPos.push_back(node[idx].len - 1);
    occur.push_back(1);
    while (p != -1 \&\& node[p].next[id] == -1)
{
        node[p].next[id] = idx;
        p = node[p].link;
    if (p == -1)
        node[idx].link = 0;
    else {
        int q = node[p].next[id];
        if (node[p].len + 1 == node[q].len)
            node[idx].link = q;
        else {
            int clone = ++sz;
            state x = node[q];
            x.len = node[p].len + 1;
            node.push_back(x);
            firstPos.push_back(firstPos[q]);
            occur.push_back(0);
```

```
while (p != -1 && node[p].next[id]
== q) {
                node[p].next[id] = clone;
                p = node[p].link;
           node[idx].link = node[q].link =
clone;
   last = idx;
void Process() {
    cnt.resize(sz + 1);
    distinct.resize(sz + 1);
    SA.resize(sz + 1);
    adj.resize(sz + 1);
    for (int i = 0; i \le sz; i++)
cnt[node[i].len]++;
    for (int i = 1; i <= L; i++) cnt[i] +=
cnt[i - 1];
    for (int i = 0; i <= sz; i++)
SA[--cnt[node[i].len]] = i;
   for (int i = sz; i > 0; i--) {
        int idx = SA[i];
        occur[node[idx].link] += occur[idx];
        adj[node[idx].link].push_back(idx);
        distinct[idx] = 1;
        for (int j = 0; j < MXCHR; j++) {
            if (node[idx].next[j] != -1)
                distinct[idx] +=
distinct[node[idx].next[j]];
   } // counts distinct substrings and
occurance of each state
    for (int i = 0; i < MXCHR; i++)
        if (node[0].next[i] != -1) distinct[0]
+= distinct[node[0].next[i]];
pair<int, int> lcs(string &str) {
    int mxlen = 0, bestpos = -1, pos = 0, len
    int u = 0; // LCS of two string. returns
start position and length
   for (char c : str) {
        int v = getID(c);
        while (u && node[u].next[v] == -1) {
            u = node[u].link;
            len = node[u].len;
        if (node[u].next[v] != -1) {
            len++;
            u = node[u].next[v];
        if (len > mxlen) {
            mxlen = len;
            bestpos = pos;
        pos++;
```

```
return {bestpos - mxlen + 1, mxlen};
    state &operator[](int index) { return
    node[index]; }
};
9.7 Trie [28 lines]
const int maxn=100005;
struct Trie{
  int next[27][maxn];
  int endmark[maxn],sz;
  bool created[maxn];
  void insertTrie(string& s){
    int v=0;
    for(int i=0;i<(int)s.size();i++){</pre>
      int c=s[i]-'a';
      if(!created[next[c][v]]){
        next[c][v]=++sz;
        created[sz]=true;
      v=next[c][v];
    endmark[v]++;
  bool searchTrie(string& s){
    int v=0;
    for(int i=0;i<(int)s.size();i++){</pre>
      int c=s[i]-'a';
      if(!created[next[c][v]])
        return false;
      v=next[c][v];
    return(endmark[v]>0);
};
9.8 Z-Algorithm [19 lines]
void compute_z_function(const char*S,int N){
  int L=0,R=0;
  for(int i=1;i<N;++i){</pre>
    if(i>R){
      L=R=i;
      while (R < N \&\& S[R-L] == S[R]) ++R;
      Z[i]=R-L,--R;
    else{
      int k=i-L;
      if(Z[k]<R-i+1)Z[i]=Z[k];
      else{
        L=i;
        while (R < N \&\& S[R-k] == S[R]) ++ R;
        Z[i]=R-L,--R;
 }
```