

Midterm CS2040C Sem 1 2021-2022

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1 Question 01

- (1) What will be printed if we run that code? (any answer with the first two most significant digits correct will be accepted)

$[N \times \log_2(2N) \times 2 \rightarrow 2048 \times 12 \times 2 = 49152]$

- (2) What is the time complexity of the code above in terms of N ?

$[O(N \log(N))]$

- (3) If N in line 5 is changed to 1000000 (1 Million), do you think your code can run in less than 1s if your computer can do 100 Million basic operations in 1s? Write Y/N in the box:

$[N]$

- (4) If the comment `"/"` in line 8 is removed so that line 8 is part of the code, will the time complexity that you have answered in part 2 changes? Write Y/N in the box:

$[N]$

- (5) If line 10 is changed to `for (int k = 0; k < N; ++k)`, the time complexity of the code above in terms of N changes into $O(5)$

$[N^2 \log(N)]$

2 Question 02

- (1) `#include <1>`

$[algorithm]$

- (2) `for (auto 2name : names)`

$[\&]$

- (3) `std::sort(names.3(), names.end(), [](const std::string &a, const std::string &b) {}`

$[begin]$

- (4) `if (a.length() 4 b.length())`

$[!=]$

- (5) `return a.length() 5 b.length();`

$[<]$

3 Question 03

- (1) Choose the Data Structure to be used

```
std::vector
```

- (2) How do you implement add(v)

[using the vector method

```
push_back(v)
```

with the input being the element. This method implementation has a $O(1)$ time complexity]

- (3) How do you implement count()

[using the vector method size() with the output being the size of the vector or the number of items in the ADT. The time complexity will be $O(1)$]

- (4) How do you implement returnAnyWithEqualProbability()

[let the index be:

```
int index = rand() % size();
```

where this will randomly choose an index within the size of the ADT. Then once the index is found, use the STL vector in built

```
vector.remove(index);
```

Finding the random index is $O(1)$ while removing the element is $O(N)$]

4 Question 04

What time complexity?

[$O(N)$]

5 Question 05

What is the time complexity?

[$O(N)$]

6 Question 06

What is the time complexity?

[$O(N \log(N))$]

7 Question 07

What is the time complexity?

[$O(N^2)$, because we have no tail, we must traverse until we find the end]

8 Question 08

explain how newSort() works and why it can sort properly (HARD)

[Now the Quicksort doesn't sort the whole array entirely, leaving some portions of it unsorted. Thus, with the addition of insertionSort, it will sort the rest of the array in $O(K)$ times, meaning the time complexity is $O((N - K) \log(N - K) + K)$]

9 Question 09

1. What is the time complexity?

[$O(N^2)$, since K is large, we reach once $\text{high-low}+1 \leq K$ very early and have to resort to insertionSort]

2. What is the time complexity?

[$O(N \log(N))$, since K is very small, we are always $\text{high-low}+1 > K$ and we keep recursing, resulting in mostly quickSort all the way]

10 Question 10

What is the expected time complexity of newSort() in terms of K (that can range from 1 to N) and N ? (HARD)

from $O(N(\log(N) - \log(K)))$, we only quick sort a portion of the array. What's that portion you may ask? why that's $N - K$. Thus, we quicksort $O(N(\log(N) - \log(K))) = O(N \log(\frac{N}{K}))$. When partitioning, we have to multiply by $O(N)$. Next, we insertionSort the rest of what's left, which is $\text{length} \leq K$. For each element its correct location must be in the same subarray. So each element is at most K swaps away from correct position. Thus, $O(NK)$ for insertion sort.

Therefore answer is: $O(N \log(\frac{N}{K}) + NK)$

11 Question 11

- (1) Fill In the Blanks

[1]

- (2) Fill In the Blanks

[3]

- (3) Fill In the Blanks

[9]

- (4) Fill In the Blanks

[7]

- (5) Fill In the Blanks

[N]

12 Question 12

```
#include <iostream>
#include <algorithm> // blanks
#include <vector>
#include <utility>
#include <unordered_map>

using namespace std;
typedef pair<int, int> pii;
typedef pair<int, pii> pip;
typedef unordered_map<int, pii> mip; // {value, {firstOccur, count}}

int main() {
    int N, K; cin >> N >> K; // 1 <= N <= 1M, 1 <= K <= N

    mip map;
    for (int i = 0; i < N; ++i) { // O(N) here
        int v; cin >> v;
        if (map.find(v) != map.end()) {
            ++map[v].second;
        } else {
            map[v] = {i, 1};
        }
    }

    vector<pip> A;
    for (auto &e : map) {
        A.push_back({e.first, e.second});
    }

    // sort
    sort(A.begin(), A.end(), [](pip &a, pip &b){ // {value, {firstOccur, count}}
        if (a.second.second != b.second.second) {
            return a.second.second > b.second.second;
        } else {
            return a.second.first < b.second.first;
        }
    });

    for (auto a : A) {
        for (auto i = 0; i < a.second.second; ++i)
            cout << a.first << " ";
    }
    cout << endl;
    return 0;
}
```

[the overall time complexity of my C++ code above is $O(N + N^2 \log(N^2))$, where N^2 is the number of unique numbers in N inputs]