CS2022 Overflow

- When we use fixed size Register for arithmetic operands there is the hazard at each micro-ops of overflow.
- If Register R stores an n-bit 2's compliment number, then:

$$-2^{n-1} \leq R \leq 2^{n-1}-1$$

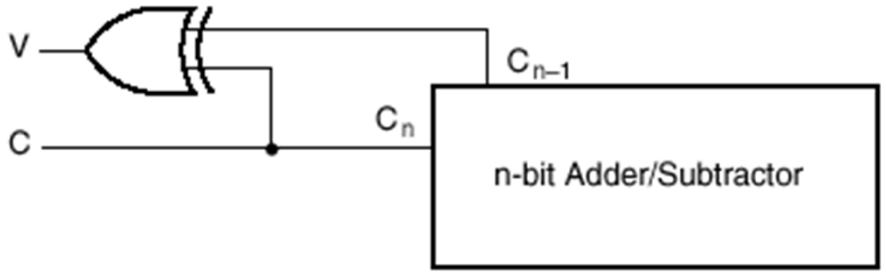
- Note the asymmetry, so even negation can cause overflow, e.g.

$$R \leftarrow -2^{n-1} \textbf{ then } R \leftarrow -R \Rightarrow \text{OVERFLOW}$$

CS2022 Carry and Overflow

- We can detect overflow for all the previous operations simply by recording the status bits C = Carry and V = Overflow (see section 3.10 I Mano and Kime).

$$K_1: C \leftarrow C_n,\ V \leftarrow C_n \oplus C_{n-1}$$

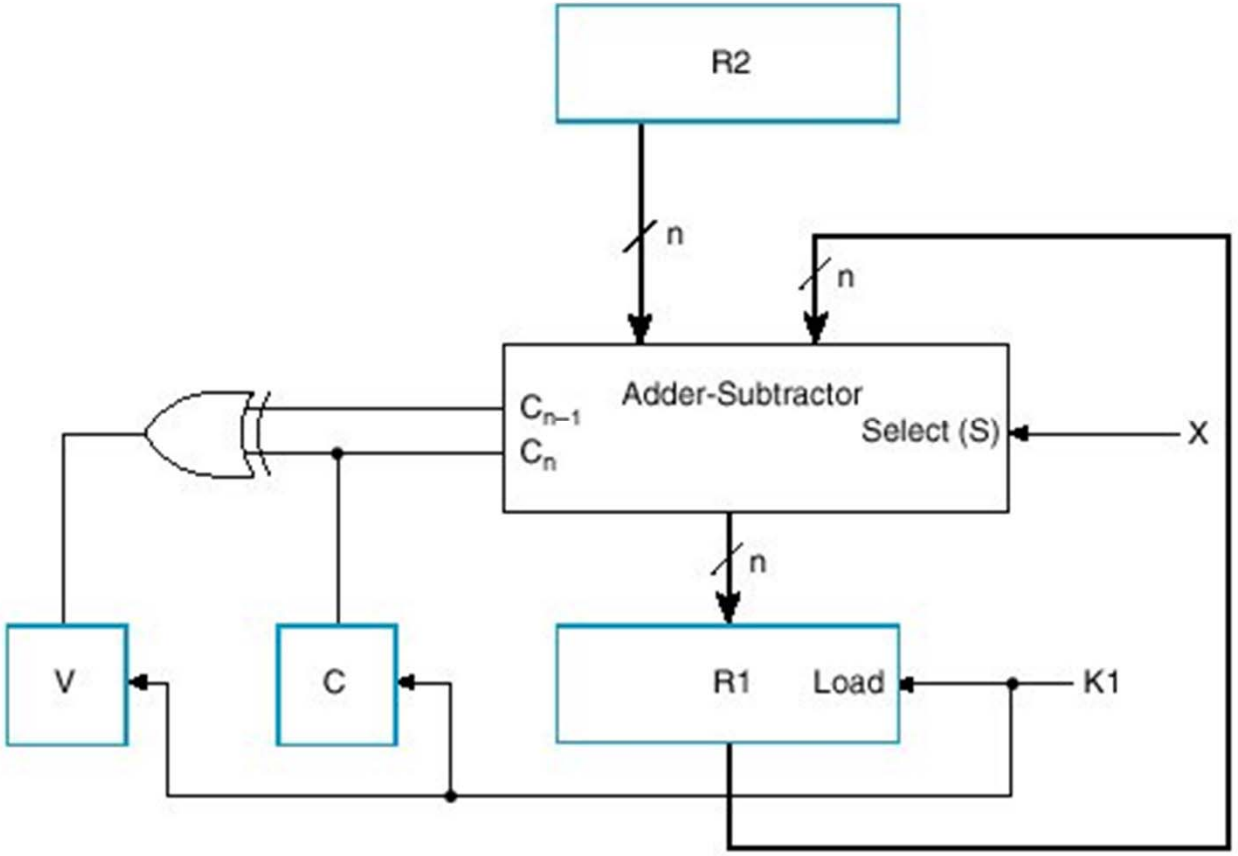


CS2022 Adder-Subtractor

- This then is the basis for the adder-subtractor on the next page which uses control input X to select addition and $\overline{X}$ for subtraction.

$$\overline{X}.K_1: R1 \leftarrow R1 + R2$$
$$X.K_1: R1 \leftarrow R1 + \overline{R2} + 1$$
$$K_1: C \leftarrow C_n,\ V \leftarrow C_n \oplus C_{n-1}$$

CS2022 Add & Subtract Micro operation



CS2022 Logic Micro-operations

- The aim here is to provide an effective set of bit-wise functions. A typical basic set follows:

Symbolic micro-op	Description
$R0 \leftarrow \overline{R1}$	Logical bitwise NOT (1's compliment)
$R0 \leftarrow R1 \wedge R2$	Logical bitwise AND (clears bits)
$R0 \leftarrow R1 \vee R2$	Logical bitwise OR (sets bits)
$R0 \leftarrow R1 \oplus R2$	Logical bitwise XOR (complements bits)

CS2022 ∧ And ∨

- George Bode Prof. Of mathematics in UCC introduced the notation '∧' and '∨' in 1854, and they are used in Register Transfer (RT) notation if it is necessary to distinguish addition from logical OR. For example:

$$K_1 + K_2 : R1 \leftarrow R2 + R3,\ R4 \leftarrow R5 \vee R6$$

Logical OR

addition

Logical OR

- The OR micro-ops will always use ∨.

CS2022 Shift Micro-operations

- These provide lateral bitwise shift which are essential for many basic arithmetic algorithms e.g. multiplication division , square root…
- The minimal set is:

$$R \leftarrow sr\ R \equiv R_i \leftarrow R_{i+1}\ i=0, n-2, R_{n-1} \leftarrow 0$$
$$R \leftarrow sl\ R \equiv R_i \leftarrow R_{i-1}\ i=1, n-1, R_0 \leftarrow 0$$

CS2022 Logical Shifts

- These are logical shifts and from them you can develop variants which handle the end bits differently, e.g. arithmetic shift, rotates…
- Please see table 9-5 in Mano and Kime.

CS2022 Transfer Micro-operations

- Providing choice of path
- Two approaches are used:
  - Multiplexer-based transfer for speed
  - Bus-based transfers for flexibility and economy

CS2022 Transfer with Multiplexer

- The if-then else control structure, when applied to identity micro-ops, results in the destination register requiring selective access to two different source registers.

If (K_1=1) then R0←R1
else if (K_2=1) then R0←R2

K_1: R0←R1, $\overline{K_1}K_2$:R0←R2

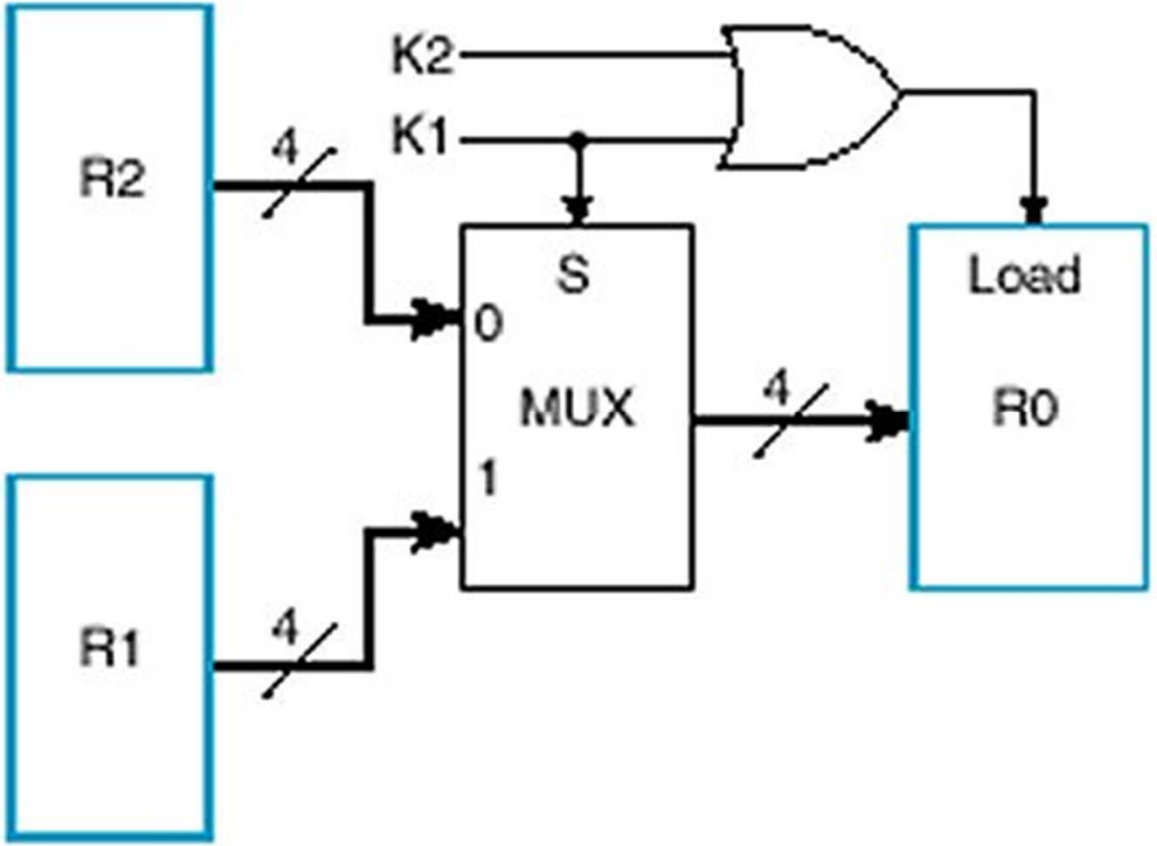
Two Sources - One Destination

- To route two sources to one destination we can use a 2:1 MUX with control input S, data input D0 $\wedge$ D1, and R0 must have a load control $LOAD_{R0}$.
- From the RT description we deduce the following functions for these control inputs.

$$LOAD_{R0} = K_1 + \overline{K}_1 K_2 = K_1 + K_2$$
$$S = K_1 \Rightarrow D1 = R1 \wedge D0 = R2$$

Multiplexer

- From this we derive the circuit schematic, assuming 4-bit Registers.



CS2022 Multiplexer Detail

- Note changes in K1 only affect the MUX – R0 path, so transients are short, allowing fast operation.