

資訊安全導論 HW4

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產生 key

首先生成兩個在 bit 範圍內的隨機質數 p q

用 p q 算出 n 與 ϕ_n

公鑰 e 是隨機的從範圍(2~ ϕ_n)內挑出一個與 ϕ_n 互質的數字

挑法: 1.生成隨機數 2.若為偶數 +1 變奇數 3.判斷是否互質 不互質找下一個奇數

私鑰 d 是公鑰 e 在模 ϕ_n 下的模反元素，因此解此式 $(e*d) \bmod \phi_n == 1$ ，我使用 python 的 pow 函式內建的模反計算(第二個參數設為-1)

```
#生成所有 key
# p q 為大質數
# n 為 p*q phi_n 為(p-1)*(q-1)
# e 為公鑰 d 為私鑰
def key_gen(bit_count):
    p = prime_gen(int(bit_count))
    q = prime_gen(int(bit_count))
    n = p * q
    phi_n = (p-1)*(q-1)
    public_key = random.randint(2,phi_n)
    if public_key % 2 == 0:
        public_key += 1
    while not is_coprime(public_key,phi_n):
        public_key += 2
    private_key = pow(public_key,-1,phi_n)
    print("p = {}".format(p))
    print("q = {}".format(q))
    print("n = {}".format(n))
    print("phi = {}".format(phi_n))
    print("e = {}".format(public_key))
    print("d = {}".format(private_key))
    return
```

質數生成: 1.生成隨機數 2.若為偶數 +1 變奇數 3.判斷是否為質數 不是的話往下一個奇數找

```
#生成一個指定 bit 數範圍內的質數
def prime_gen(bit_count):
    result = random.getrandbits(bit_count)
    if result % 2 == 0:
        result += 1
    while miller_rabin(result) == False:
        result += 2
    return result
```

判斷質數: miller robin 演算法 [米勒-拉賓質數判定法 - 維基百科，自由的百科全書 \(wikipedia.org\)](https://zh.wikipedia.org/wiki/米勒-拉賓質數判定法)

```
#miller rabin 演算法
#用於判斷一個數字是否為質數
def miller_rabin(check_num):
    if check_num == 2 or check_num == 3:
        return True
    if check_num == 1 or check_num < 0 or check_num % 2 == 0:
        return False
    k = 10
    r = 0
    s = check_num-1
    while(s%2==0):
        r += 1
        s //=2
    for i in range(k):
        a = random.randrange(2,check_num-1)
        x = pow(a,s,check_num)
        if(x==1 or x == check_num-1):
            continue
        for j in range(r-1):
            x = pow(x,2,check_num)
            if(x==check_num-1):
                break
        else:
            return False
    return True
```

判斷互質: 1.判斷輸入 ab 大小 大的放第一個參數 2.求 ab 最大公因數 若為 1

則互質

#求 a b 最大公因數(無遞迴版本)，大的數字要放在 a

```
def gcd(a,b):  
    if b <= 1 or b >= a:  
        return -1  
  
    while b != 0:  
        a, b = b, a % b  
    return a
```

#確認 a b 是否互質，也就是確認 a b 的最大公因數是不是 1

#會先判斷 a b 大小，將大的放在第一個參數，以配合 gcd

```
def is_coprime(a, b):  
    if a >= b:  
        return gcd(a, b) == 1  
    else:  
        return gcd(b, a) == 1
```

加密

輸入一訊息，與 n 和公鑰，回傳加密後數字的 base64 編碼

```
#加密函式，輸入待加密訊息(str)與 n 和公鑰
#並將 RSA 加密後的數字透過 base64 編碼，使其能夠顯示
def encrypt(plain_text,n,public_key):
    plain_num = str2num(plain_text)
    cipher_num = RSA_encrypt(plain_num,n,public_key)
    cipher_base64 = base64.b64encode(str(cipher_num).encode('ascii'))
    cipher_text = cipher_base64.decode('ascii')
    return cipher_text
```

RSA 加密，輸入一待加密數字與 n 和公鑰，回傳加密後數字

求: $(\text{plain_num} \wedge \text{公鑰}) \bmod n$

```
#RSA 加密，輸入一數字與 n 和公鑰
#算出加密後的數字
def RSA_encrypt(plain_num,n,public_key):
    cipher_num = fast_power_mod(plain_num,public_key,n)
    return cipher_num
```

fast_power_mod 即為 square and multiply 演算法

將次方數轉為二進制，掃過所有 bit，遇 0 則 square，遇 1 則 square 與

multiply

```
#計算(base^exponent) mod mod_num
#採用 square and multiply algorithm 加速
def fast_power_mod(base, exponent, mod_num):
    bins = bin(exponent)
    result = 1
    for index in range(0,len(bins)):
        result = (result * result) % mod_num
        if bins[index] == '1':
            result = (result * base) % mod_num
    return result
```

解密

輸入加密後的 base64 編碼，與 n 和私鑰，回傳解密後原始訊息

```
#解密函式，輸入加密後訊息(base64 的 str)與 n 和私鑰
#解密出數字後，將數字轉回原始訊息
def decrypt(cipher_text,n,private_key):
    cipher_num = int(base64.b64decode(cipher_text).decode('ascii'))
    plain_num = RSA_decrypt(cipher_num,n,private_key)
    plain_text = num2str(plain_num)
    return plain_text
```

RSA 解密，輸入一加密後數字與 n 和私鑰，回傳原始數字

求: $(\text{cipher_num} \wedge \text{私鑰}) \bmod n$

```
#RSA 解密，輸入一加密數字與 n 和私鑰
#算出解密後的數字
def RSA_decrypt(cipher_num,n,private_key):
    plain_num = fast_power_mod(cipher_num,private_key,n)
    return plain_num
```

CRT 解密

輸入加密後的 base64 編碼，與 p 與 q 和私鑰，回傳解密後原始訊息

```
#CRT 加速之解密函式，輸入加密後訊息(base64 的 str)與 p 與 q 和私鑰
#解密出數字後，將數字轉回原始訊息
def CRT_decrypt(cipher_text,p,q,private_key):
    cipher_num = int(base64.b64decode(cipher_text).decode('ascii'))
    plain_num = RSA_CRT_decrypt(cipher_num,p,q,private_key)
    plain_text = num2str(plain_num)
    return plain_text
```

RSA 加密，輸入一加密後數字與 p 與 q 和私鑰，回傳原始數字

採用中國剩餘定理加速

```
#RSA CRT 解密，輸入一數字與 p 與 q 和私鑰
#算出解密後的數字
def RSA_CRT_decrypt(cipher_num,p,q,private_key):
    dp = private_key % (p-1)
    dq = private_key % (q-1)
    q_inv = pow(q,-1,p)
    m1 = fast_power_mod(cipher_num,dp,p)
    m2 = fast_power_mod(cipher_num,dq,q)
    h = q_inv * (m1-m2) % p
    result = m2 + h * q
    return int(result)
```

結果

Key gen:

```

C:\Users\frakw-notebook\Documents\coding\110-InformationSecurity\HW4\10815057_廖慶彬-python_RSA.py -i
p = 168730963036439668443570793810288765843767799732495019370042993098720677643137140092360362338205439898651
884607243423758756231658060029772151237899590965875652292460214788721137206478209854323070802020696058748015970
29514668067690759066470801404231138669831746031425883995323674114778859182429105530433
q = 1033762905888029785435272360108256708391633246575136148587248466732529732000073134981721265120586244894810
65553101440980347250338863466939685192966231915912527909931874577616975395309825327138874213518728553384906468
876383020305404402189154774196248884532726673069751588313461663032909863430354901421
n = 174427810661736676110827690288391283925469220002721814611063965360674191707681922308670245173629538838540869
1835217743129931254549651388978999241951404072736637570998391987322859785459642228121229331983448659262841620448860
5827779312355864824573643386899015476786924927923566875540390880043899125031620736016295805544703329420958230904
87655016626732449414574034953949557590041384644314545465158193609878387341792591713710385712406643248856977584
2198367523085382545888398938488145695125318932904875860469811495217290543810254072681971692547693718597309424002
829826582943855702499120417371680391072654699630445293
phi = 17442781066173667611082769028839128392546922000272181461106396536067419170768192230867024517362953883835408
69183521774312993125449651388978999241951040727366375709983919873228597854596422281212293319834486592628416204488
6058277793123558648245736433868990154767869249279235668755403908800438991250316207360145051646237510867766603073
784561788333964036741947939686837592611122751872792271686401061049447248527809383881996679578598384962950491646279
13066010326871122797893227378949241270758815001977156825601026733939399135159583544034734217955737900586815416
082271142725024918631591385715401023398957503840170013440
e = 15153146912998936106999404870913373632859500122845080631040948414767546366517041233405686776840980552749925432
90483656996026955204539625920791359312509302469414678017111722487172619426809296765495607535270863484752507738
5936343189850177729149589086626301646266893736979255034872109635790835678059035677211261211481874349023540568
9438720953938934299795714841778646490280252521289250168991161678646706483070929521710487280642464632433657619
58393395500629654267578758440798054067400151565291763186658548865995524012578980908653545282627260144524560132878
38449840951636794511630496587696561924085142487534067
e = 142249084313458659056393238180221855980081142475976395766961757088537509562330037001589034282951781525482373
1889364690970413188750900326258843020638454220069609645451977231556197889237956346301855846357439263799343157
727381656916625329863100437564716921032726301305585486747106675123662012397306239664978365475924363671361543260
17306602166978957703764394930640388947080070987549460072100400381602867137483284831945260532380927494027095269308
760590766667052749406616702461181916544004410370753162677214789810332163838217965744306828548037474944299174603
95070181060708688879383725203864366910222650646852941694

```

加密:

[illegible]

解密:

[illegible]

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