

CSC 212: Data Structures and Abstractions

Recursion

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Fall 2020



Recursion

Recursion

- Solve a task by reducing it to smaller tasks (**of the same structure**)
- Technically, a recursive function is one that **calls itself**
- General form:
 - ✓ **base case**
 - solution for a **trivial case**
 - it can be used to stop the recursion (prevents “*stack overflow*”)
 - every recursive algorithm needs at least one base case
 - ✓ **recursive call(s)**
 - divide problem into **smaller instance(s)** of the **same structure**

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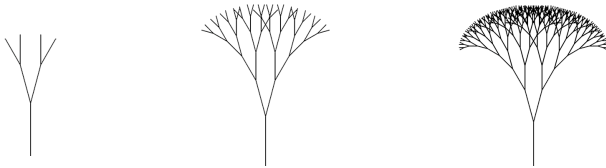
General Form

```
function() {  
    if (this is the base case) {  
        calculate trivial solution  
    } else {  
        break task into subtasks  
        solve each task recursively  
        merge solutions if necessary  
    }  
}
```

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Why recursion?

- Can we live without it?
 - yes, you can write “any program” with arrays, loops, and conditionals
- However ...
 - some formulas are explicitly recursive
 - some problems exhibit a natural recursive solution



<https://courses.cs.washington.edu/courses/cse120/17sp/labs/11/tree.html>

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```
int sum_array(int *A, int n) {  
    // base case  
    if (n == 1) {  
        return A[0];  
    }  
  
    // solve sub-task  
    int sum = sum_array(A, n-1);  
  
    // return sum  
    return A[n-1] + sum;  
}
```

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Recursion call tree

```
int sum_array(int *A, int n) {  
    // base case  
    if (n == 1) {  
        return A[0];  
    }  
  
    // solve sub-task  
    int sum = sum_array(A, n-1);  
  
    // return sum  
    return A[n-1] + sum;  
}
```

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Example: power of a number

$$b^n = \underbrace{b \cdot b \cdot \dots \cdot b}_{n \text{ times}}$$

```
double power(double x, int n) {  
    // base case  
    if (n == 0) {  
        return 1;  
    }  
  
    // recursive call  
    return x * power(x, n-1);  
}
```

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Recursion call tree

```
double power(double x, int n) {  
    // base case  
    if (n == 0) {  
        return 1;  
    }  
    // recursive call  
    return x * power(x, n-1);  
}
```

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Is this faster?

```
double power(double x, int n) {  
    if (n == 0) {  
        return 1;  
    }  
    double half = power(x, n/2);  
    if (n % 2 == 0) {  
        return half * half;  
    } else {  
        return x * half * half;  
    }  
}
```

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Recursion call tree

```
double power(double x, int n) {  
    if (n == 0) {  
        return 1;  
    }  
    double half = power(x, n/2);  
    if (n % 2 == 0) {  
        return half * half;  
    } else {  
        return x * half * half;  
    }  
}
```

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Binary Search

Binary Search

0	1	2	3	4	5	6	7	8	9	10	11	12
1	2	5	10	15	20	22	30	35	40	43	48	51

low

high

k = 48?

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Binary Search

0	1	2	3	4	5	6	7	8	9	10	11	12
1	2	5	10	15	20	22	30	35	40	43	48	51

low

mid

high

k = 48?

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Binary Search

0	1	2	3	4	5	6	7	8	9	10	11	12
1	2	5	10	15	20	22	30	35	40	43	48	51

low

high

k = 48?

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Binary Search

0	1	2	3	4	5	6	7	8	9	10	11	12
1	2	5	10	15	20	22	30	35	40	43	48	51

low

mid

high

k = 48?

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Binary Search

0	1	2	3	4	5	6	7	8	9	10	11	12
1	2	5	10	15	20	22	30	35	40	43	48	51

low mid high

k = 48?

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```
int bsearch(int *A, int lo, int hi, int k) {  
    // base case  
    if (hi < lo) {  
        return NOT_FOUND;  
    }  
    // calculate midpoint index  
    int mid = lo + ((hi-lo)/2);  
    // key found?  
    if (A[mid] == k)  
        return mid;  
    // key in upper subarray?  
    if (A[mid] < k)  
        return bsearch(A, mid+1, hi, k);  
    // key is in lower subarray?  
    return bsearch(A, lo, mid-1, k);  
}
```

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Recursion Tree (binary search)

```
int bsearch(int *A, int lo, int hi, int k) {  
    // base case  
    if (hi < lo) {  
        return NOT_FOUND;  
    }  
    // calculate midpoint index  
    int mid = lo + ((hi-lo)/2);  
    // key found?  
    if (A[mid] == k)  
        return mid;  
    // key in upper subarray?  
    if (A[mid] < k)  
        return bsearch(A, mid+1, hi, k);  
    // key is in lower subarray?  
    return bsearch(A, lo, mid-1, k);  
}
```

Complexity? Best-case, Worst-case, Average-case?

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Example: find peak in
unimodal arrays

Unimodal arrays

- An array is (**strongly**) **unimodal** if it can be split into an increasing part followed by a decreasing part

1	2	5	16	20	18	17	16	15	12	10	8	5
---	---	---	----	----	----	----	----	----	----	----	---	---

How to efficiently find the max?

1	2	5	16	20	18	17	16	15	12	10	8	5
---	---	---	----	----	----	----	----	----	----	----	---	---

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Find the **max** (strongly unimodal)

1	2	5	8	15	20	22	20	15	12	10	8	5
1	2	5	8	15	20	22	20	15	12	10	8	5
1	2	5	8	15	20	22	20	15	12	10	8	5
1	2	5	8	15	20	22	20	15	12	10	8	5
1	2	5	8	15	20	22	20	15	12	10	8	5
1	2	5	8	15	20	22	20	15	12	10	8	5
1	2	5	8	15	20	22	20	15	12	10	8	5
1	2	5	8	15	20	22	20	15	12	10	8	5

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Find the **max** (strongly unimodal)

- Recursion Tree?

- Complexity? Best-case, Worst-case, Average-case?

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Unimodal arrays

- An array is (**weakly**) **unimodal** if it can be split into a nondecreasing part followed by a nonincreasing part

1	2	5	5	15	20	22	22	35	38	13	8	5
---	---	---	---	----	----	----	----	----	----	----	---	---

How to efficiently find the max?

1	2	5	5	15	20	22	22	35	38	13	8	5
---	---	---	---	----	----	----	----	----	----	----	---	---

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Find the **max** (weakly unimodal)

1	2	5	5	15	20	22	22	35	38	13	8	5
---	---	---	---	----	----	----	----	----	----	----	---	---

1	2	5	5	15	20	22	22	35	38	13	8	5
---	---	---	---	----	----	----	----	----	----	----	---	---

1	2	5	5	15	20	22	22	35	38	13	8	5
---	---	---	---	----	----	----	----	----	----	----	---	---

Two recursive calls

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Find the **max** (weakly unimodal)

▸ Recursion Tree?

▸ Complexity?

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