Networking Introduction

OSI vs TCP/IP

OSI		TCP/IP		
Application				
Presentation	Application Protocols	Application		
Session				
Transport	Ports UDP TCP	Transport		
Network	IP	Internet		
Data Link	MAC	Network Access		
Physical		Network Access		

Encapsulation Decapsulation 01A3041132AB 01A30411ED14 **Application Application** DATA DATA 192.168.10.15 192.168.10.12 **Ports** 49791 49791 TCP or UDP 80 80 **Transport** Transport **IP Addresses** 192.168.10.12 49791 192.168.10.12 192.168.10.15 80 192.168.10.15 Internet 80 Internet **MAC Addresses** 49791 01A3041132AB 01A3041132AB 49791 Network Network 80 192.168.10.15 01A30411ED14 192.168.10.15 01A30411ED14 Access **Access**

Addressing

- Physical Media Access Control (MAC Address) of NIC
- Logical IP Address of NIC
- Network Services Port numbers

Physical Address (MAC Address)

- MAC Address basically made up of two parts
 - Vendor Address
 - Random Address
- 013E1FBB23A1
 - 013E1FBB23A1 First part is the Vendor Address
 - 013E1FBB23A1 Last part is the Random Address
- FFFFFFFFFF
 - Special MAC address for Broadcasts
- This is how computers communicate with each other NIC to NIC

Logical Addresses (IP)

- IP addresses can be used to:
 - Provide a unique address for a host
 - Provide a way to group hosts on a subnet
 - Facilitate the transmission of network packets (data) between
 - Local hosts (on the same subnet)
 - Remote hosts (on different subnets separated by routers)
 - Provide private addresses
 - Provide public addresses
 - Allow private addresses to be translated to public and back again

IP Addresses and Binary

- 131.107.1.4
 - Dotted decimal notation
 - Each decimal number represents an octet (8 binary digits)
 - 0 = 00000000 (Minimum value)
 - 255 = 11111111 (Maximum value)
- 131 = 10000011
- 107 = 01101011
- 1 = 00000001
- 4 = 00000100
- 131.107.1.4 = 10000011 01101011 00000001 00000100

Binary

- Consists of 0 or 1 (off / on)
- Counting in binary:
 - 00000000 bin = 0 dec
 - 00000001 bin = 1 dec
 - 00000010 bin = 2 dec
 - 00000011 bin = 3 dec
 - 00000100 bin = 4 dec
 - 00000101 bin = 5 Dec
 -
 - 11111111 bin = 255 dec

For each 1 in binary add the corresponding decimal values:

Binary to Decimal conversion								
1	1	1	1	1	1	1	1	
128	64	32	16	8	4	2	1	
1111111 bin => 128+64+32+16+8+4+2+1 = 255 dec								

Binary to Decimal conversion								
1 0 1 0 1 1								
128	64	32	16	8	4	2	1	
10101011 bin => 128+32+8+2+1 = 171 dec								

Binary to Decimal conversion								
0	0	0	1	0	0	1	0	
128	64	32	16	8	4	2	1	
10010 bin => 16+2 = 18 dec								

IP Address (Network vs Host)

- IP Addresses describe two things Network section & Host section
 - Network section: is like the name of a street
 - Host section: is like the house number
- Examples:
 - 192.168.0.1
 - 192.168.0.1 Network section, 192.168.0.1 Host section
 - 131.107.1.4
 - 131.107.1.4 Network section, 131.107.1.4 Host section
 - 10.34.2.1
 - 10.34.2.1 Network section, 10.34.2.1 Host section
- Subnet Mask
 - Shows which section is the network and which is the host

Subnet Mask

- Determines which part of the IP address is the:
 - Network section
 - Host section
- Helps to determine if a destination host is on:
 - Local subnet
 - Remote subnet.
- Traditional Class Subnet masks

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Dotted Decimal Binary
Class A = 255.0.0.0 = 11111111 00000000 00000000 00000000 = /8
Class B = 255.255.0.0 = 11111111 1111111 1111111 00000000 = /16
Class C = 255.255.255.0 = 11111111 1111111 1111111 00000000 = /24
Class D = no subnet mask specified (Used for multicast)
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- Traditional masks wasted a lot of address
 - Solution = Classless inter-domain routing (CIDR)
 - CIDR allows for bitwise manipulation of the subnet mask

Original IP Classes (assignable addresses)

Class A

- First octet 1-126 (Binary 00000000 01111111) [0 & 127 by convention are not used)
- First octet is the network section
- Last three octets are the Host section

Class B

- First octet 128-191 (Binary 10000000 10111111)
- First two octets are the network section
- Last two octets are the Host section

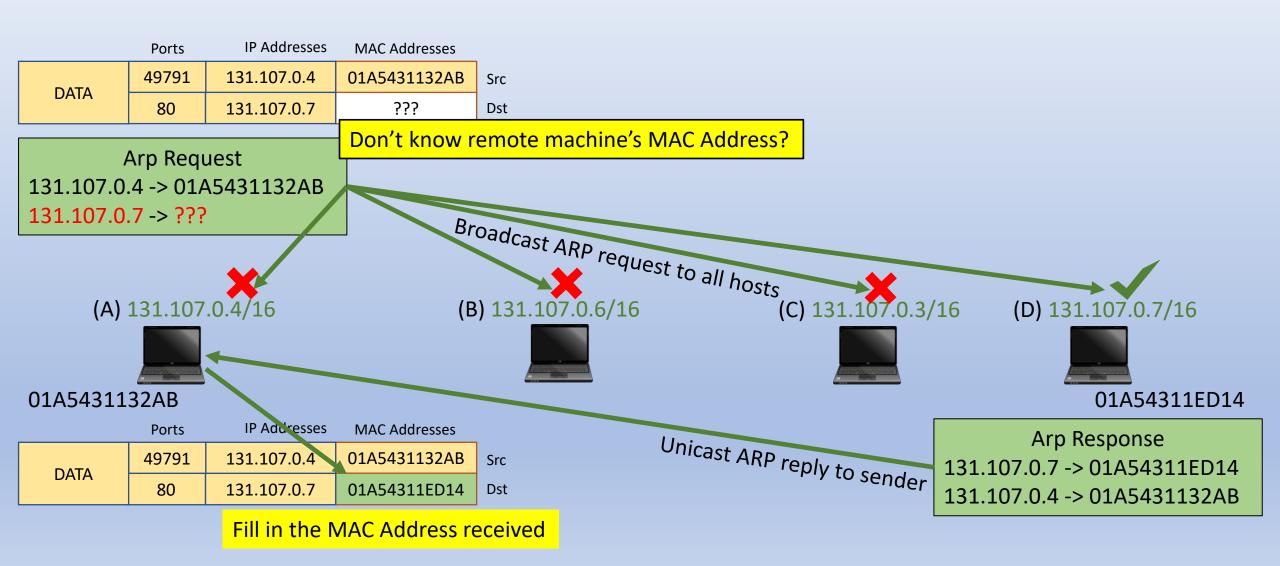
Class C

- First octet 192-223 (Binary 11000000 11011111)
- First three octets are the network section
- Last octet is the Host section

ARP – Address Resolution Protocol

- ARP resolves an IP address into its corresponding MAC address
- ARP request is sent as a broadcast packet
- Asks who owns an IP address
 - IP owner answers with its MAC address with a unicast packet
 - ARP resolution information is cached in memory (for both the sender and receiver)
 - ARP cache entries last about 10 minutes in memory
- Routers do not pass broadcasts
 - Cannot ARP for an IP on a remote subnet
 - To communicate with a remote host, route packet through a router
 - Set Default Gateway IP so hosts can be used to find the default router
- Use arp -a to show the contents of the ARP cache
- Use arp -d to delete the contents of the ARP Cache

ARP in action (Host A to Host D)

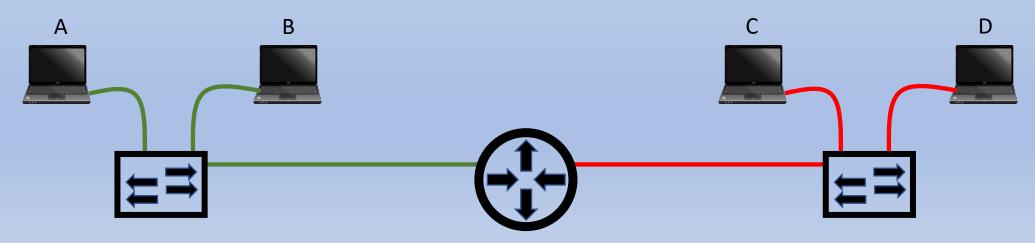


Default Gateway

- Typically set on each host
 - Allows the host to route traffic out of the local subnet
 - Points to the default router
 - Is used when the host needs to:
 - Send traffic to a remote network
 - There is no specific route for the intended destination
- Only used for remote destinations
- Sender sends an ARP request for the Default Gateway's IP address
- Sender receives the MAC address of the router's NIC

Determine if destination is local or remote

- Sending host
 - ANDs its IP address with the senders Subnet Mask
 - ANDs the destination IP address with the senders Subnet Mask
- Compare the results of the previous ANDing calculations
 - if the results are the same then the destination IP is local (Computers A & B)
 - if the results are different then destination IP is remote (Computers B & C)



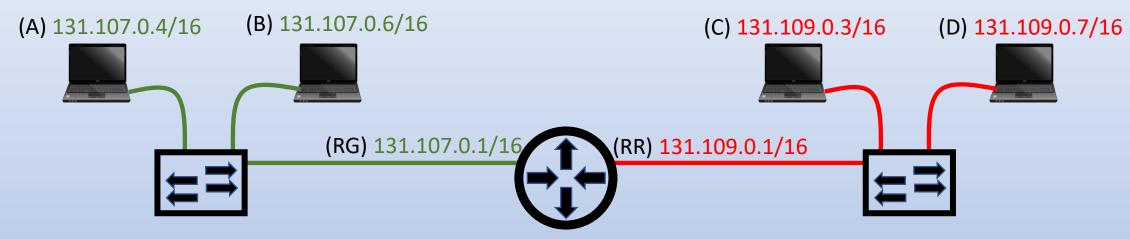
Determine how to route traffic

- Each time a network transmission happens
 - Sender need to determine if destination is on the local or remote network
- If the destination IP address is on the local network
 - ARP for the destination IP address
 - This informs the sender of the remote host's MAC address.
 - Send the network packet to the MAC address of the destination host
- If the destination IP address is on a remote network
 - ARP for the IP address of the Default Gateway
 - This informs the sender of the default gateway's MAC address
 - Send the network packet to the MAC address of the default router

Boolean AND calculation

BOOLEAN AND							
0	AND	0	=	0			
0	AND	1	=	0			
1	AND	0	=	0			
1	AND	1	=	1			

Determining if hosts are local or remote



- A -> B (Local)
 - 131.107.0.4 AND 255.255.0.0 = 131.107.0.0
 - 131.107.0.6 AND 255.255.0.0 = 131.107.0.0
- B -> C (Remote)
 - 131.107.0.6 AND 255.255.0.0 = 131.107.0.0
 - 131.109.0.3 AND 255.255.0.0 = **131.109.0.0**

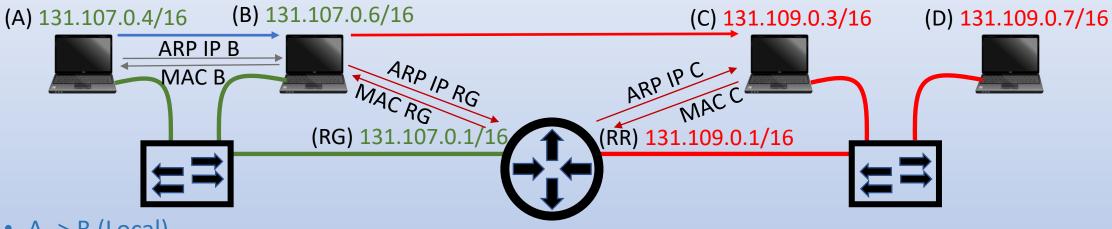
AND calculation (A -> B)

- 131.107.0.4 AND 255.255.0.0
 - 10000011 01101011 00000000 00000100 (131.107.0.4)
 - 11111111 1111111 00000000 00000000 (255.255.0.0)
 - 10000011 01101011 00000000 00000000 (131.107.0.0 AND Result)
- 131.107.0.6 AND 255.255.0.0
 - 10000011 01101011 00000000 00000110 (131.107.0.6)
 - 11111111 1111111 00000000 00000000 (255.255.0.0)
 - 10000011 01101011 00000000 00000000 (131.107.0.0 AND Result)
- Results are the same this means the two addresses are local

AND calculation (B -> C)

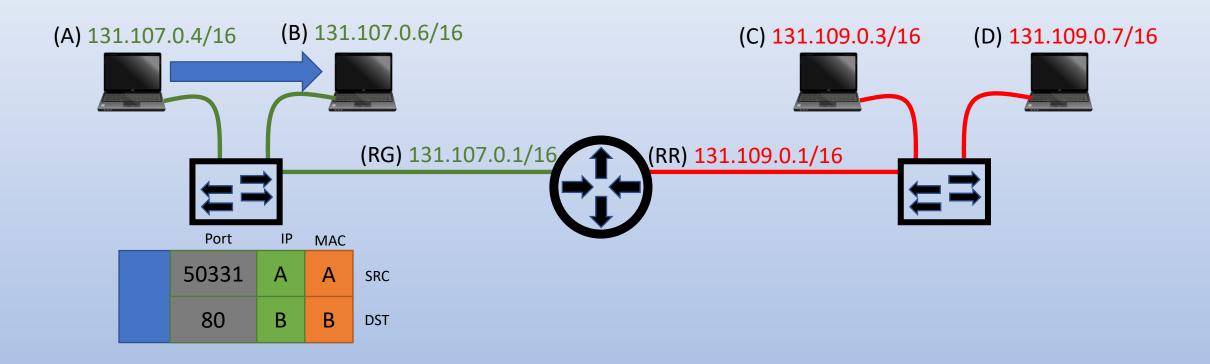
- 131.107.0.6 AND 255.255.0.0
 - 10000011 01101011 00000000 00000110 (131.107.0.6)
 - 11111111 1111111 00000000 00000000 (255.255.0.0)
 - 10000011 01101011 00000000 00000000 (131.107.0.0 AND Result)
- 131.109.0.3 AND 255.255.0.0
 - 10000011 01101101 00000000 00000011 (131.109.0.3)
 - 11111111 1111111 00000000 00000000 (255.255.0.0)
 - 10000011 01101101 00000000 00000000 (131.109.0.0 AND Result)
- Results are different this means the two addresses are remote

ARP Requests

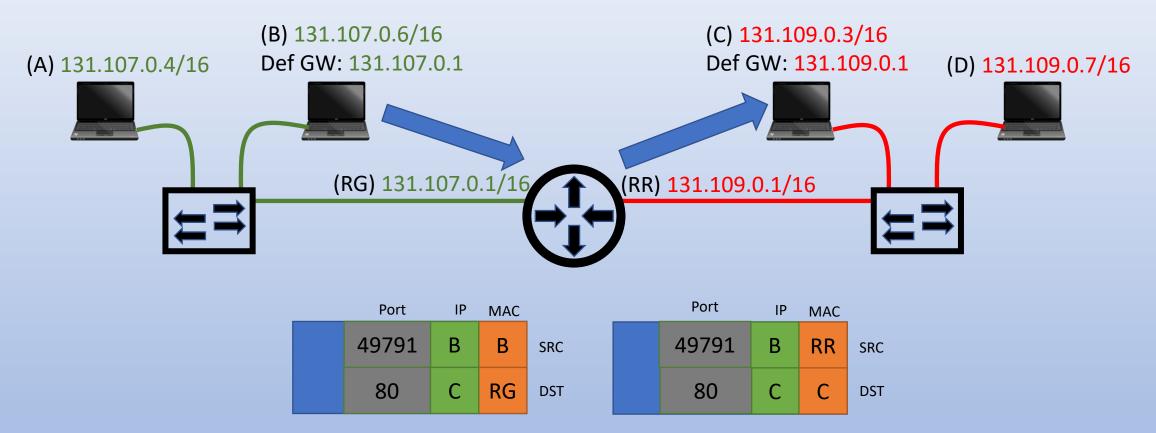


- A -> B (Local)
 - From A, ARP for destination IP B
 - Reply with MAC address from B
- B -> C (Remote)
 - From B, ARP for the router RG IP address
 - Reply with MAC address from RG
 - Send the network packet to the routers MAC address
 - From Router, ARP for destination C IP address
 - Reply with MAC address from C

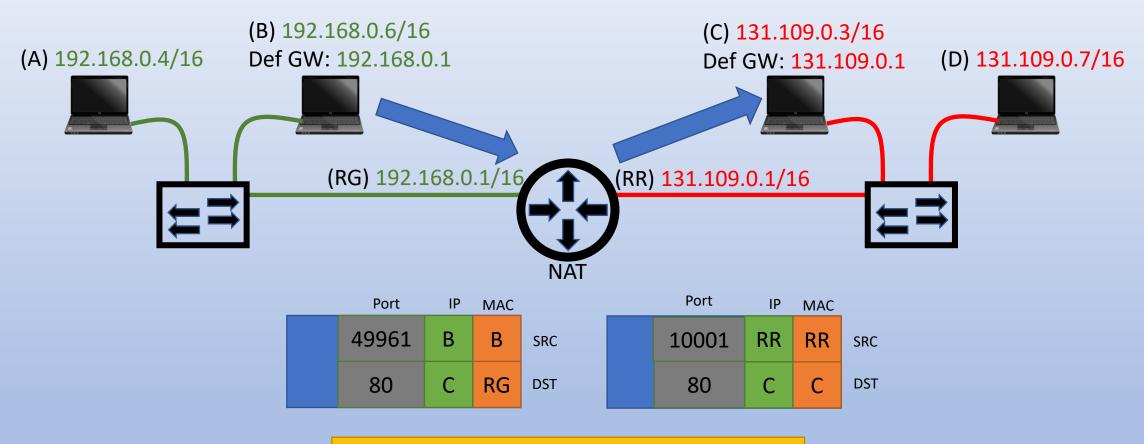
Routing local (A -> B [Webserver])



Routing remote (B -> C [Webserver])



Routing and NAT (B -> C [Webserver])



NAT TABLE

192.168.0.6:**49961** ---> 131.109.0.1:**10001**

Subnetting Introduction

Why do we subnet

- To convert a single network into smaller networks
- To reduce wasting of IP addresses
- To control network traffic congestion
- To manage the growth of a network
- To make network security more granular

How do we subnet

- You need to answer several questions before you are ready to subnet
 - What network address are we starting with
 - How many smaller networks do we need (Subnets)
 - How many IP addresses do we need on each subnet (minimum: 1 per NIC)
 - What growth do we expect to happen within the subnets
 - Is it possible to achieve the result with the given conditions

What network are we starting with

- Class A, B, C each have default subnet masks /8, /16, /24
- Some IP addresses on each network have special purposes
- We will use 131.12.0.0/16 as an example:
 - Network ID: 131.12.0.0/16
 - Broadcast ID: 131.12.255.255/16
 - The first address in each subnet is its Network ID
 - The last address in each subnet is its Broadcast ID
 - The Network ID and Broadcast ID are not used for addressing hosts
 - So, the usable (valid) addresses are: 131.12.0.1 -> 131.12.255.254 / 16

Subnetting example - 131.12.0.0/16

- Responses from an organisation to our questions:
- How many subnets are required
 - 28
- How many IP addresses are required per subnet
 - 1300
- What growth do we expect on these subnets
 - 5% growth in the number of subnets
 - 10% growth in IP addresses
- Therefore, max subnets: 28 + 5% = 28 + 1.4 (~2) = 30
- Therefore, max IP addresses: 1300 + 10% = 1300 + 130 = 1430

Example continued - Calculate subnets

- Subnets Required: 30
- IP addresses / subnet required : 1430
- Starting Network ID: 131.12.0.0/16
- /16 means
 - That the first 16 bits are network
 - This leaves the last 16 bits for us to slice up into smaller subnets
- How many bits are required to make at least 30 subnets
- 5 Bits
 - 00000 -> 11111 = is 32 possible numbers (we need 30)
- That leaves eleven bits for hosts (16 − 5)
 - 0000000000 -> 1111111111 = 2048 possible numbers 2 (one for Net ID and one for BC ID)
 - 2046 IP addresses on each of the 32 subnets
- The subnetting is possible!

How is the subnetting performed

- The new subnet mask is the old mask + the network bits required
 - Old mask = 16 bits
 - Subnet bits required for subnetting = 5 bits
 - New subnetting mask = 21 bits
- Old 131.12.0.0/16 in binary looks like this:
 - 10000011 00001100 00000000 00000000
 - <u>11111111 1111111 00000000 00000000 (16 bits)</u>
- New 131.12.0.0/21 in binary looks like this:
 - 10000011 00001100 00000000000000000 (Subnet portion Host portion)
 - <u>11111111 1111111 11111</u>000 00000000 (21 bits)

How to create the first subnet

- 10000011 00001100 <mark>00000</mark>000 00000000
- The first subnet would look like this in binary:
 - 10000011 00001100 00000 000 00000000 Subnet ID
 - 10000011 00001100 00000000000001 First valid address
 - 10000011 00001100 00000111 11111110 Last valid address
 - 10000011 00001100 00000111 11111111 Broadcast ID
- In Decimal
 - 131.12.0.0 Subnet ID
 - 131.12.0.1 First valid address
 - 131.12.7.254 Last valid address
 - 131.12.7.255 Broadcast ID

How to create the second subnet

- 10000011 00001100 <mark>00001</mark>000 00000000
- The second subnet would look like this in binary:
 - 10000011 00001100 00001000 00000000 Subnet ID
 - 10000011 00001100 00001000 00000001 First valid address
 - 10000011 00001100 00001111 11111110 Last valid address
 - 10000011 00001100 00001111 11111111 Broadcast ID
- In Decimal
 - 131.12.8.0 Subnet ID
 - 131.12.8.1 First valid address
 - 131.12.15.254 Last valid address
 - 131.12.15.255 Broadcast ID

How to create the third subnet

- 10000011 00001100 <mark>00010</mark>000 00000000
- The third subnet would look like this in binary:
 - 10000011 00001100 00010 000 00000000 Subnet ID
 - 10000011 00001100 00010000 00000001 First valid address
 - 10000011 00001100 00010111 11111110 Last valid address
 - 10000011 00001100 00010 111 1111111 Broadcast ID
- In Decimal
 - 131.12.16.0 Subnet ID
 - 131.12.16.1 First valid address
 - 131.12.23.254 Last valid address
 - 131.12.23.255 Broadcast ID

How to create the fourth subnet

- 10000011 00001100 <mark>00011</mark>000 00000000
- The fourth subnet would look like this in binary:
 - 10000011 00001100 00011 000 00000000 Subnet ID
 - 10000011 00001100 00011000 00000001 First valid address
 - 10000011 00001100 00011111 11111110 Last valid address
 - 10000011 00001100 00011111 11111111 Broadcast ID
- In Decimal
 - 131.12.24.0 Subnet ID
 - 131.12.24.1 First valid address
 - 131.12.31.254 Last valid address
 - 131.12.31.255 Broadcast ID

Supernetting Introduction

Why Supernetting

- Optimising Route tables
- Creating larger networks by combining networks
- It is the opposite to subnetting

How to Supernet

- Combined networks are only in sets of 2,4,8,16,32 ...
- Supernetting only can combine subsequent networks together
- The host section of the supernet must:
 - Start with all binary 0's
 - End with all binary 1's

Supernet Example

• Situation:

- Need 1000 hosts on a single network
- Class C networks are all we have to work with
- Each class C network only gives us 256 2 = 254 host addresses

What is required

- Determine the number of bits to accommodate 1000 hosts?
 - Ten bits minimum are required to provide 1000 hosts
 - Using groups of class C we would need 4 subnets (1024 2 = 1022 hosts)

Supernet Solution

- Using these 4 networks
 - 191.9.8.0/24, 191.9.9.0/24, 191.9.10.0/24, 191.9.11.0/24
- 10111111 00001001 000010<mark>00</mark> 00000000 /24
- 10111111 00001001 000010<mark>01</mark> 00000000 /24
- 10111111 00001001 000010<mark>10</mark> 00000000 /24
- 10111111 00001001 000010<mark>11</mark> 00000000 /24

Bits 23 and 24 start at 00 and end with 11 this is a requirement

- •
- 10111111 00001001 000010<mark>00 00000000 /22</mark>
- 191.9.8.0/22 [1 subnet, 1022 hosts (1024 2)]

By removing bits 23 and 24 from the subnet mask we create 1 subnet from 4 subnets

Variable Length Subnet Masks (VLSM)

Why VLSM

- The subnets we need are rarely the same size
- We need some subnets that only require two host addresses (WANs)
- VLSM is a more efficient use of network address space
- It is much more practical to use VLSM

VLSM Example

- 1st network requires 100 hosts (Sydney)
- 2nd network requires 20 hosts (Philippines)
- 3rd network requires 15 hosts (Perth)
- 4th network is a WAN links that require 2 hosts (Per -> Syd)
- 5th network is a WAN links that require 2 hosts (Syd -> Php)
- Network to subnet is 189.10.10.0/24



Perth

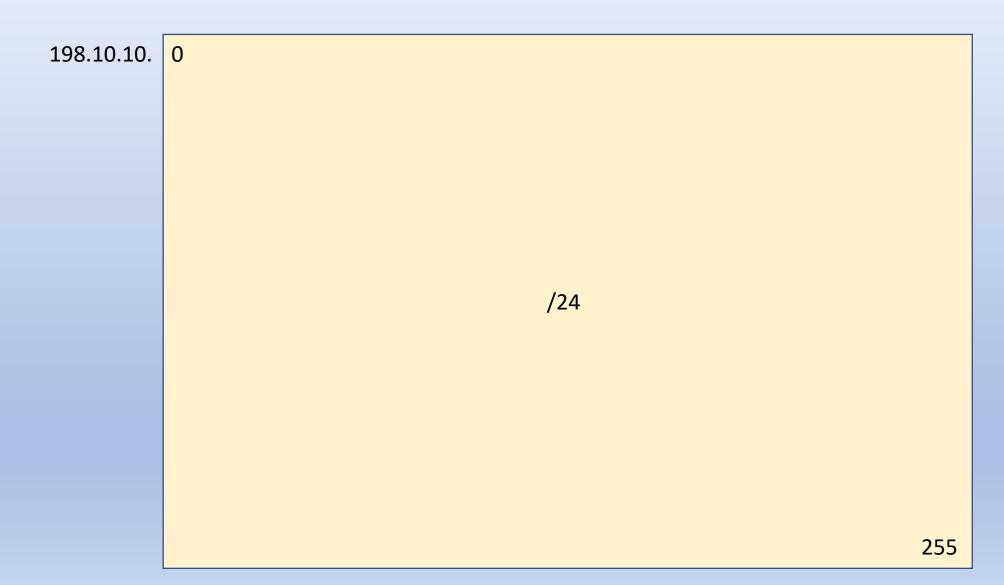


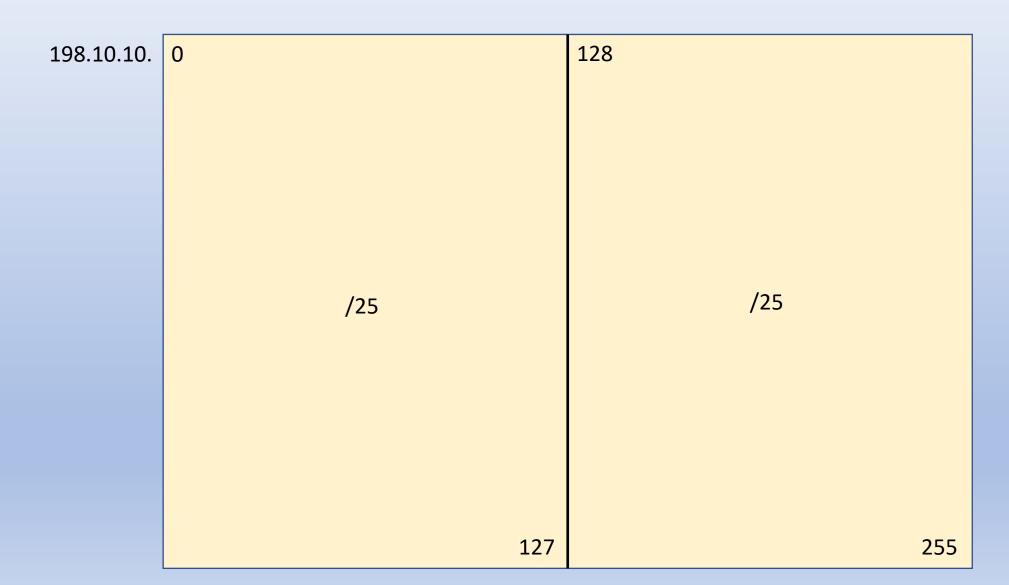
Sydney

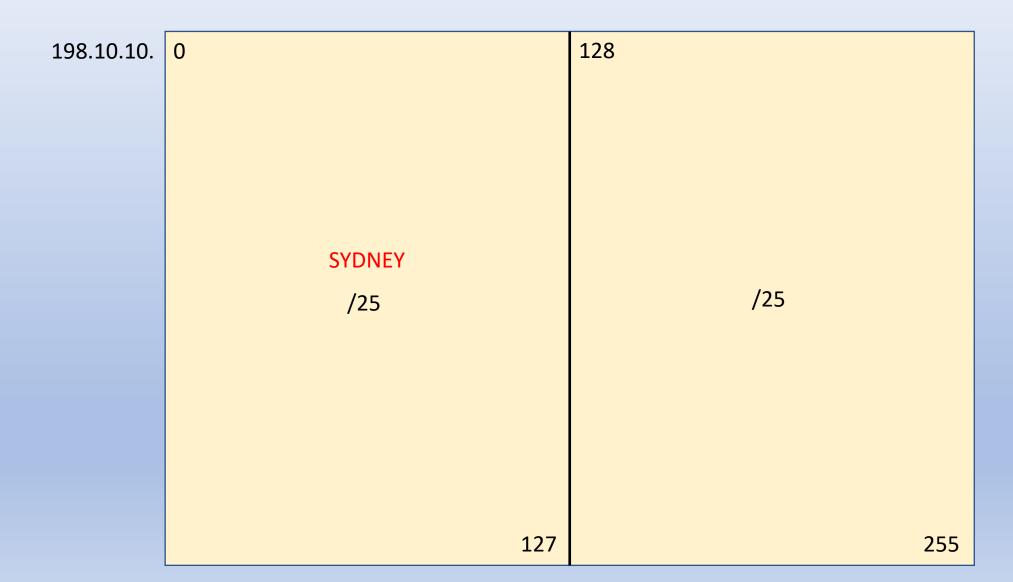


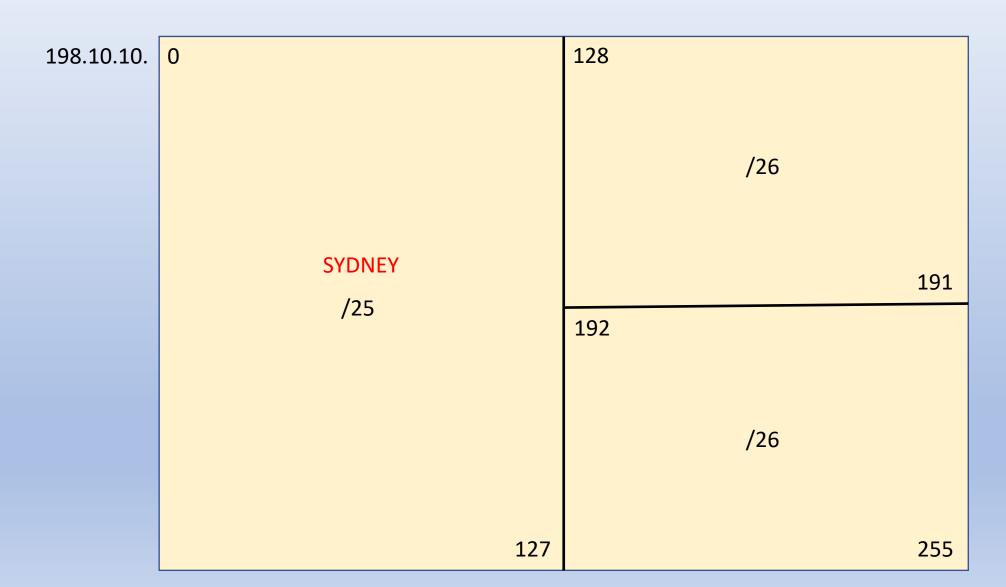
Philippines

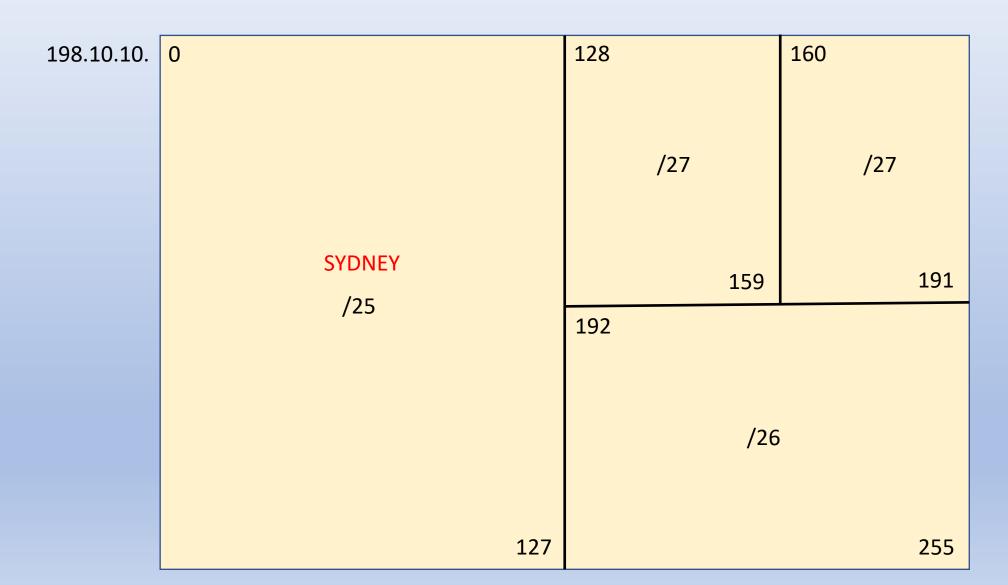


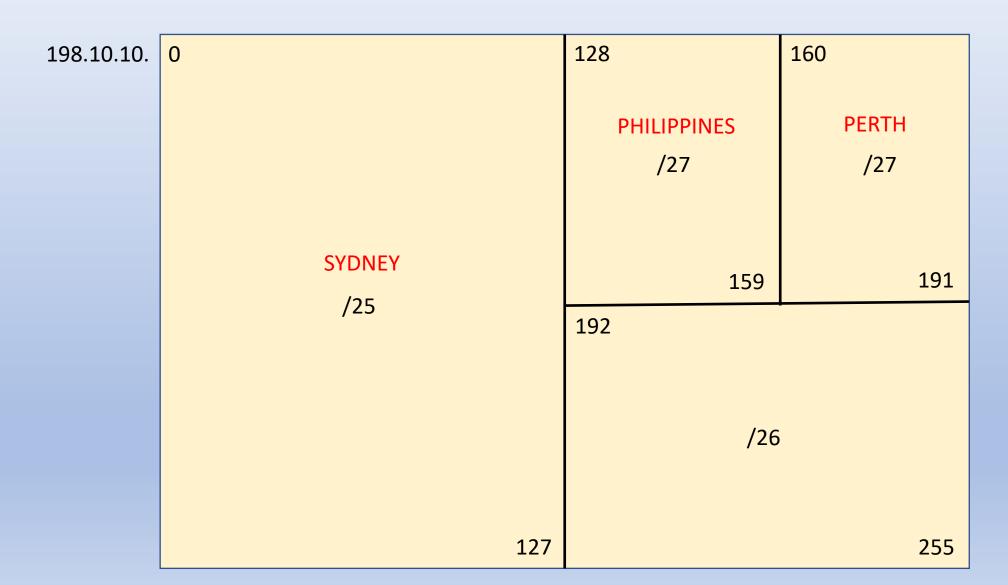








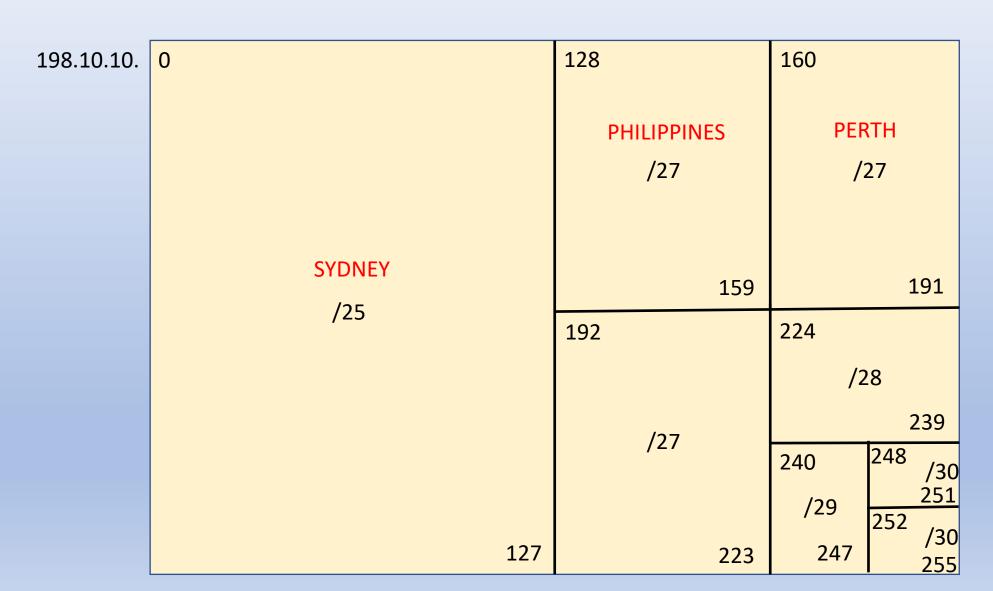


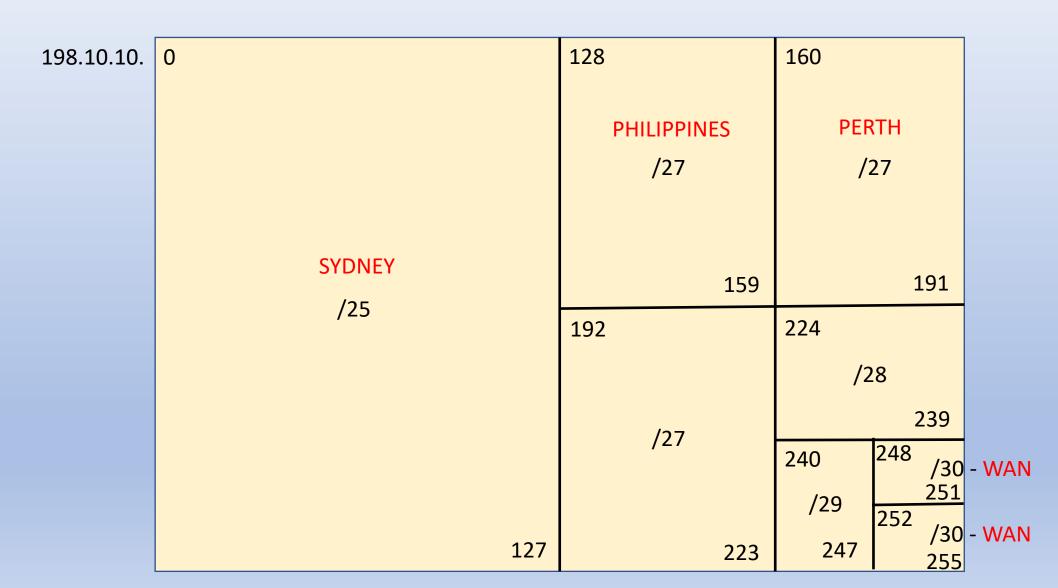


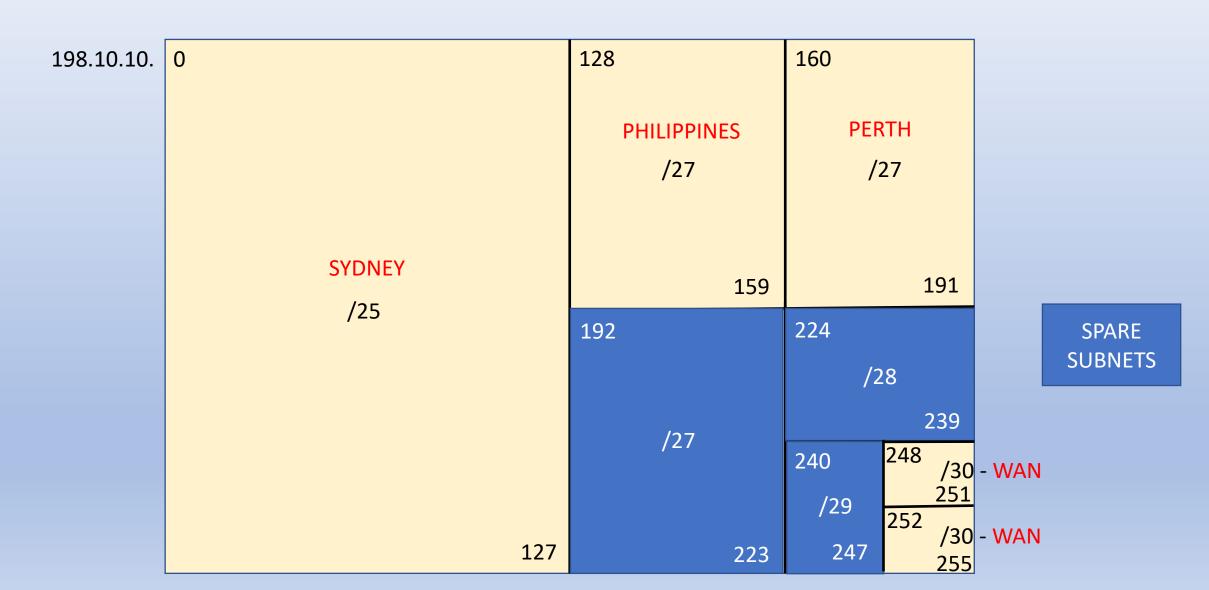
198.10.10.	0	128	160
		DI III IDDINIEC	DEDTIL
		PHILIPPINES	PERTH
		/27	/27
	SYDNEY		
		159	191
	/25	192	224
		/27	/27
	127	223	255

198.10.10.	0	128	160	
		PHILIPPINES /27	PERTH /27	
	SYDNEY	159	191	
	/25	192	224	
			/28	
		/27	239	
			240	
			/28	
	127	223	255	

198.10.10.	0	128	160	
		PHILIPPINES	PERTH	
		/27	/2	27
	SYDNEY	159		191
	/25	192	224	
			/2	8
		/27	239	
		/2/	240	248
			/29	/29
	127	223	247	255







VLSM Solution

- Sydney
 - 198.10.10.0/25 (198.10.10.1 198.10.10.126) [100/126 hosts]
- Philippines
 - 198.10.10.128/27 (198.10.10.129 198.10.10.158) [20/30 hosts]
- Perth
 - 198.10.10.160/27 (198.10.10.161 198.10.10.190) [15/30 hosts]
- WAN
 - 198.10.10.248/30 (198.10.10.249 198.10.10.250) [2/2 hosts]
- WAN
 - 198.10.10.252/30 (198.10.10.253 198.10.10.254) [2/2 hosts]
- Spare
 - Three subnets spare for future network growth