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MIDS 22 AL Move On Experience

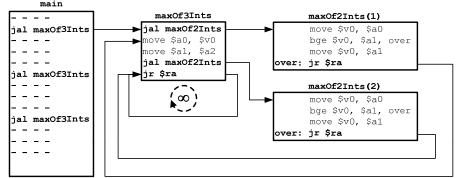
What if maxOf3Ints calls another function?

MIPS32 AL – More On Functions (following C++ lead - will the cow come home?)

```
.data
.word 1, 2, 3, 1, 2
.asciiz "\nmax(1, 2, 3) = "
.asciiz "\nmax(2, 3, 1) = "
.asciiz "\nmax(3, 1, 2) = "
                                                                                                                                           move $a0, $t1
li $v0, 1
syscall
max_of1:
max_of2:
max_of3:
                                                                                                                                           lw $a0, 8($s0)
lw $a1, 12($s0)
lw $a2, 16($s0)
jal maxOf3Ints
                          .text
.globl main
main:
                         la $s0, someInts
                                                                                                                                            move $t1, $v0
la $a0, max_of3
li $v0, 4
syscall
                          lw $a0, 0($s0)
lw $a1, 4($s0)
lw $a2, 8($s0)
jal maxOf3Ints
                                                                                                                                           move $a0, $t1
li $v0, 1
syscall
                          move $t1, $v0
la $a0, max_of1
li $v0, 4
syscall
                                                                                                                                           li $v0, 10
syscall
                          move $a0, $t1
li $v0, 1
syscall
                                                                                                                 maxOf3Ints: jal maxOf2Ints
                                                                                                                                           move $a0, $v0
move $a1, $a2
jal maxOf2Ints
jr $ra
                          lw $a0, 4($s0)
lw $a1, 8($s0)
lw $a2, 12($s0)
jal maxOf3Ints
                                                                                                                move $t1, $v0
la $a0, max_of2
li $v0, 4
syscall
        (continued on the right)
```

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MIPS32 AL – More On Functions (following C++ lead – cow is never seen again)



- In HLL (e.g.: C++), compiler takes care of necessary bookkeeping
- In AL, bookkeeping responsibility rests with programmer
 - Remember: there's *only ONE set of registers* available
 - Specifically: there's only ONE \$ra that must be shared by all function calls

.

MIPS32 AL – More On Functions (how 2 keep up w/ function-call bookkeeping)

- Function calls/returns in most programming languages:
 - ◆ Follow strict LIFO (last-in, first-out) order
- So, *stack* is "best suited function-call bookkeeping device"
 - ◆ Main reason every machine has runtime (system, call, ...) stack
- Some machines have memory stack (special hardware) as part of architecture (*e.g.*, VAX)
 - Hardware supports stack management
 - ◆ Have special instructions (*e.g.*, push and pop)
- MIPS does not support stack management in hardware
 - Stack has to be implemented via software convention
 - Use of stack segment governed by agreed upon social contract
 - ◆ Normal data transfer instructions (e.g., **1w** and **sw**) are used
 - Like (what we did before) using stack segment for local variables

MIPS32 AL – More On Functions (for case in hand, we've known enough to fix)

MIPS32 AL – More On Functions (but a generally-usable convention is needed)

- Should support ...
 - ◆ Any function calling any function any # of times from anywhere

 ✓ Including function calling "itself" (recursion), directly or indirectly
 - Each function should also be able to pass any # of arguments
- Should ensure that ...
 - Every function returns to where it should when it completes
- Should enable/facilitate ...
 - Establishment and preservation of each function's *environment*
 - So that each function can successfully/correctly perform what is intended
 - - Information about (describing state of) function during execution
 - Arguments, values of local variables, which statement being executed, ...
 - - Values of all registers (whose contents are subject to change) that function references



MIPS32 AL – More On Functions (a function-call convention for non-leaf funcs)

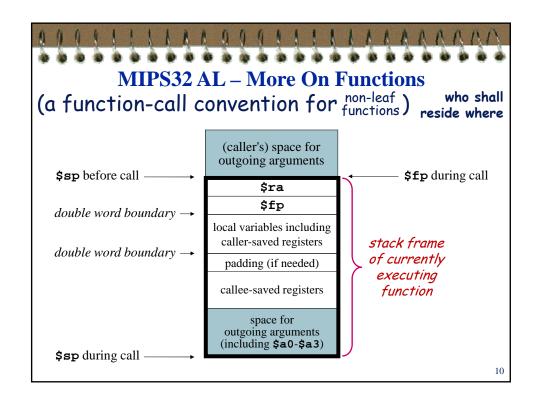
Part of Convention: MIPS Register Usage Convention

Name	Reg. Number	Usage	Preserved Across Call	
\$zero	0	constant 0	(NA)	
\$v0-\$v1	2-3	function results	N	
\$a0-\$a3	4-7	function arguments	N	
\$t0-\$t7	8-15	temporaries (caller-saved)	N	
\$s0-\$s7	16-23	saved (callee-saved)	Y	
\$t8-\$t9	24-25	temporaries (caller-saved)	N	
\$k0-\$k1	26-27	reserved for OS kernel	N	
\$gp	28	global pointer	Y	
\$sp	29	stack pointer	Y	
\$fp	30	frame pointer	Y	
\$ra	31	return address	Y	

MIPS32 AL - More On Functions (a function-call convention for non-leaf funcs) ■ What *callee* (called function) must do • Allocate stack space (create *stack frame* or *activation record*) Subtract frame size from stack pointer (\$sp): can use subu or addiu
 Minimum frame size is 32 bytes (per MIPS software architecture)
 Total bytes must be divisible by 8 (aligned for floating-point numbers) Always allocate: 2 words for \$ra and \$fp, and 4 words for \$a0-\$a3 May-need-space items: callee-saved registers (\$s0-\$s7), caller-saved registers (\$t0-\$t9), local variables, and outgoing arguments ◆ Save values of any callee-saved register that may be used ◆ Always save value of \$ra and \$fp Save values of any \$50-\$57 that callee intends to use ◆ Set \$fp to "old" value of \$sp **body** Perform task (save/restore caller-saved registers accordingly) Place return result in proper location (e.g., \$v0) for caller Restore values (previously saved) of callee-saved registers ◆ Deallocate stack space (restore \$sp to its "old" value) ◆ Transfer control back to caller: jr \$ra

MIPS32 AL – More On Functions (a function-call convention for non-leaf funcs) What caller (calling function) must do Save values of any caller-saved registers (\$t0-\$t9,...) Only those that caller wants preserved Pass arguments to callee Up to 4 arguments via \$a0, \$a1, \$a2 and \$a3 Any additional arguments via stack (pushed onto stack) Transfer control to callee Use jal (we won't be using jalr) Wait for callee's return Retrieve value returned by callee (if any) Restore values (previously saved) of any caller-saved registers

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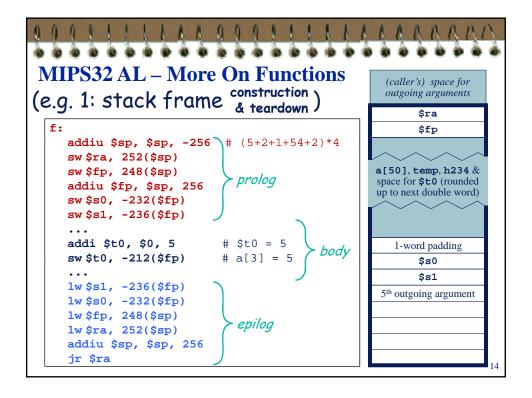


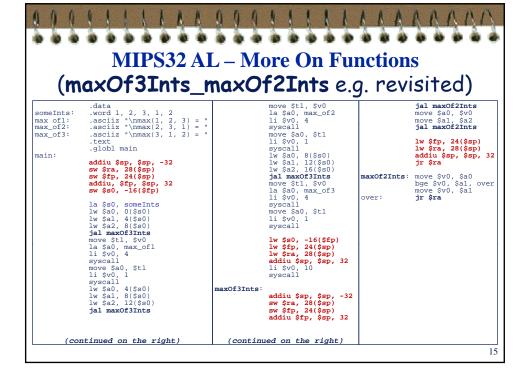
MIPS32 AL – More On Functions (a function-call convention for non-leaf func $_s^n$)

- Additional points to note:
 - ♦ 4 words allocated for \$a0-\$a3 are not used by current function
 - Reserved for use by any function current function may call
 - (Thus, these are slots allocated by caller but are used by callee!)
 - Caller must allocate *outgoing arguments section* based on the function it calls that has the <u>maximum</u> number of arguments
 - ☞ Subject to the *minimum* requirement of 4 arguments
 - So that local variables section starts on double word boundary AND total frame size is multiple of 8:
 - ☞ Insert (if necessary) 1-word padding after callee-saved registers section
 - Round up (if necessary) size of local variables section to next multiple of 8
 - Allocation of stack frame (by callee) done *only once* per function
 Once \$sp has been modified, *don't modify it again* for the rest of function
 - Values of argument registers \$a0-\$a3 and return-value registers
 \$v0-\$v1 are not required to be preserved across function calls
 - Function can thus change their values without saving/restoring them
 - ◆ Program begins with OS calling main, so main is 1st callee
 - Function is callee (gets called by another) 1st, then caller (calls another)

MIPS32 AL - More On Functions (a function-call convention for non-leaf funcs) ■ What *callee* (called function) must do ◆ Allocate stack space (create *stack frame* or *activation record*) Subtract frame size from stack pointer (\$sp): can use subu or addiu Minimum frame size is 32 bytes (per MIPS software architecture) Total bytes must be divisible by 8 (aligned for floating-point numbers) Always allocate: 2 words for \$ra and \$fp, and 4 words for \$a0-\$a3 May-need-space items: callee-saved registers (\$s0-\$s7), caller-saved registers (\$£0-\$£9), local variables, and outgoing arguments How do we know how much space How do we know what until we know whether there are we will use until we any ... or how many ... there are? write code for task? (plan ahead as best we can & make adjustments along the way)

```
MIPS32 AL - More On Functions
      (e.g. 1: stack frame analysis/design)
int f(int arg1, int arg2)
                                            assume we need to use
                                              $s0 and $s1,
                                  2 local
    int a[50], temp;
                                            preserve $t0, and are
                                 variables
                                            short of registers (thus
                                            need memory space)
    a[3] = 5;
                                            for temp and h234
    g(a[1], a[2], a[3], a[4], a[5]);
                                                    non-leaf
                                                   function w/
    int h234 = h(a[2], a[3], a[4]);
                                                  maximum of
                                                   5 outgoing
                              another local variable
                                                   arguments
    return a[3];
```





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just discussed, for non-leaf functions

- Case $1 \Rightarrow$ no need for any stack space

 - ◆ Another *e.g.* (for better comparison with other cases to see next)

```
int g(int x, int y)
{
   return x + y;
}
```

- No need for any stack frame



MIPS32 AL – More On Functions (function-call convention adapted for leaf funcs)

- Case $2a \Rightarrow$ need stack space (<= 32 bytes)
 - For local variables

```
int g(int x, int y)
                           g: addiu $sp, $sp, -32
                                                    # push min-sized stack frame
                              sw $ra, 28($sp)
   int a[3];
                              sw $fp, 24($sp)
   ... // calculations
                              addi $fp, $sp, 32
   return a[0];
                                                    # calculations
                              lw $v0, -32($fp)
                                                    # put result in $v0
                              lw $fp, 24($sp)
                              lw $ra, 28($sp)
                              addiu $sp, $sp, 32
                                                    # pop stack frame
                                                    # return
```

- Minimum frame size
 - \sim Space requirement of 6 (= 3+1+2) words is less than 32 bytes
- ♦ Calculations should not change any callee-saved registers

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MIPS32 AL – More On Functions (function-call convention adapted for leaf funcs)

- Case $2b \Rightarrow$ need stack space (> 32 bytes)
 - For local variables

```
int g(int x, int y)
                           g: addiu $sp, $sp, -72
                                                    # push stack frame
                              sw $ra, 68($sp)
{
   int a[15];
                              sw $fp, 64($sp)
   ... // calculations
                              addi $fp, $sp, 72
   return a[0];
                                                    #calculations
                              lw $v0, -72($fp)
                                                    # put result in $v0
                              lw $fp, 64($sp)
                              lw $ra, 68($sp)
                              addiu $sp, $sp, 72
                                                    # pop stack frame
                                                    # return
```

- ◆ Stack frame of 18 (= 15+1+2) words
 - Local variables section padded so that it's divisible by 8
- ♦ Calculations should not change any callee-saved registers

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MIPS32 AL – More On Functions (function-call convention adapted for leaf func $_s^n$)

- Case $3a \Rightarrow$ need stack space (≤ 32 bytes)
 - For local variables and callee-saved registers

```
int g(int x, int y)
                           g: addiu $sp, $sp, -32
                                                      # push min-size stack frame
                               sw $ra, 28($sp)
                               sw $fp, 24($sp)
   int a[3];
   ... // calculations
                               addi $fp, $sp, 32
   return a[0];
                               sw $s0, -32($fp)
                                                      # calculations ($s0 used)
                               lw $v0, -24($fp)
                                                      # put result in $v0
                               lw $s0, -32($fp)
                               lw $fp, 24($sp)
                               lw $ra, 28($sp)
                               addiu $sp, $sp, 32
                                                      # pop stack frame
                               jr $ra
                                                      # return
```

- Stack frame of 8 (= 3+1+1+1+2) words
 - Local variables section padded so that it's divisible by 8
 - Callee-saved registers section padded so that it's divisible by 8

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MIPS32 AL – More On Functions (function-call convention adapted for leaf funcs)

- Case $3b \Rightarrow$ need stack space (> 32 bytes)
 - For local variables and callee-saved registers

```
int g(int x, int y) {
    int a[25];
    int a[26];
    int a[27];
    int a[27];
    int a[28];
    int a[28];
```

- \bullet Stack frame of 32 (= 25+1+3+1+2) words
 - Local variables section padded so that it's divisible by 8
 - Callee-saved registers section padded so that it's divisible by 8

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Function:				
No stack frame	needed 🗆	leaf	non	
Total bytes (>	теат	leaf		
\$ra and \$fp (8)				
local var	to-DW filler			
section	local var			
Section	\$t0-\$t9			
padding (0 or				
callee-saved	\$s0-\$s7			
reg section				
arguments	add'l arg			
section	\$a0-\$a3 (16)			

2.1