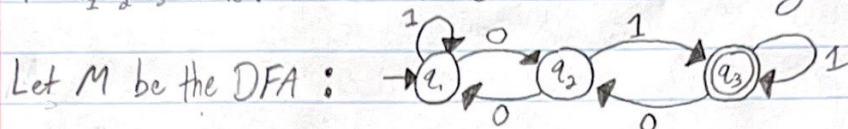


Brian DeFlaminio

1. L is a finite set of binary strings. A language is regular if some finite automaton recognizes it.

where " i " is a binary string

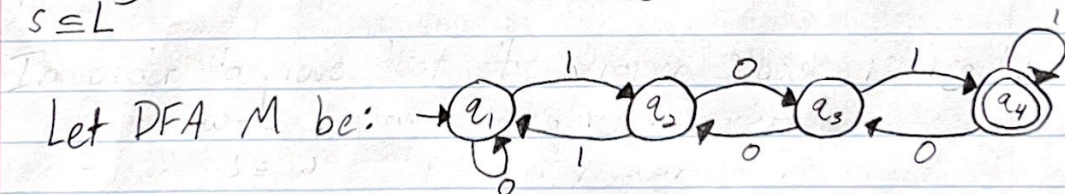
★ Prove that L is regular, for L to be regular $L = \{i \mid M \text{ accepts } i\}$
 $i = i_1 i_2 i_3 \dots i_n$. Let i be the finite binary string "011" where $i \in L$



The input on M is $i = 011$ which leads to the accept state q_3 via the following order of states: Start q_1 , read 0, move to q_2 , read 1, move to q_3 , read 1, move to q_3 , end at q_3 , accept.

We have proved that there exists a DFA, M , that recognizes L , so L must be regular.

2. Prove that the language of all binary strings containing a fixed binary string S of length k as a subsequence is a regular language. Considering that $L = \Sigma^*$ for $\Sigma = \{0, 1\}$ lets arbitrarily define S to be of length $k=4$ and $S = 1011$. Also, $S \subseteq L$



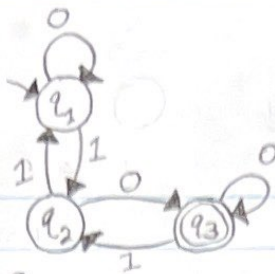
Input $S = 1011$ accepts via the following order of states: start, q_1 , read 1, move to q_2 , read 0, move to q_3 , read 1, move to q_4 , read 1, loop to q_4 , accept.

This process of arbitrarily choosing S ; k could be repeated until all subsets are exhausted and have been given a DFA.

We could then perform a Union operation on every S_1, S_2, S, \dots until we've got the lang of all binary strings which has its regularity preserved through the union operation.

1.6 b, c, d

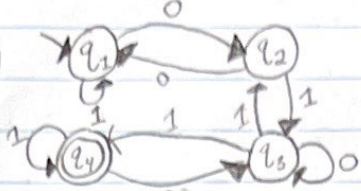
3. b)



(1, 0 accepts)

w begins with a 1 ; ends with 0

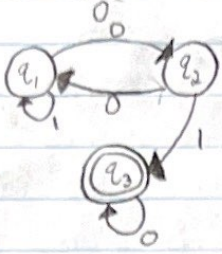
c)



(101011 accepts)

w contains 0101

d)



(010 accepts)

w is at least 3 states ; 3rd symbol is 0

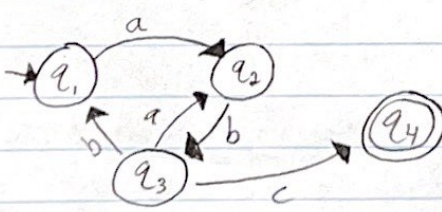
1.20 b, c, d, e

4.

b)	Member	Nonmember
	abab	aabb
	ababab	aaabbb
c)	a	ba
	b	ab
d)	aaa	abb
	aaa, aaaa, aaaaa	baa
e)	aaabaaa	bbbbbb
	babbbab	aaaa

5. Give an NFA where the final character in the input is the first time that character has appeared.

$\Sigma = \{a, b, c\}$



This NFA only accepts strings that end in the letter c.