

3a)

```
void f1(int n)
{
    int i=2;
    while(i < n){
        /* do something that takes O(1) time */
        i = i*i;
    }
}
```

$$n = 32$$

$$(1) \quad i = 2 \quad n = 32$$

$$(2) \quad i = 4 \quad n = 32$$

$$(3) \quad i = 16 \quad n = 32$$

$$(4) \quad i = 32 \quad n = 32$$

The loop starts at 2 and squares itself/doubles everytime the loop iterates. If the initial value of $n = 32$, then it takes 4 iterations for i to reach 32.

$$\therefore \boxed{O(\log i)}$$

b.

```

void f2(int n)
{
    for(int i=1; i <= n; i++){
        if( (i % (int)sqrt(n)) == 0){
            for(int k=0; k < pow(i,3); k++) {
                /* do something that takes O(1) time */
            }
        }
    }
}

```

$$T(n) = \sum_{i=1}^n [\theta(1) + O(\sum_{k=0}^{i^3} \theta(1))] = \theta(n) + \sum_{i=1}^n i^3 = \theta(n^4)$$

$$T(n) = \theta(n) + \sum_{i=1}^n i^3 = \theta(n^4)$$

c.

```

for(int i=1; i <= n; i++){
    for(int k=1; k <= n; k++){
        if( A[k] == i){
            for(int m=1; m <= n; m=m+m){
                // do something that takes O(1) time
                // Assume the contents of the A[] array are not changed
            }
        }
    }
}

```

$$T(n) = \sum_{i=1}^n \sum_{k=1}^n (\theta(1) + O(\sum_{m=1}^{\log n} \theta(1))) = \theta(n^2) + \sum_{i=1}^n \sum_{k=1}^n \log n = \theta(n^2) + \theta(n^2 \log n) = \theta(n^2)$$

$$T(n) = \theta(n^2) + \sum_{i=1}^n \sum_{k=1}^n \log n = \theta(n^2) + \theta(n^2 \log n) = \theta(n^2)$$

$$= \theta(n^2) + \theta(n^2 \log n) = \theta(n^2)$$

```

d, int f (int n)
{
    int *a = new int [10];
    int size = 10;
    for (int i = 0; i < n; i++)  $\Theta(n)$ 
    {
        if (i == size)
        {
            int newsize = 3*size/2;
            int *b = new int [newsize];
            for (int j = 0; j < size; j++) b[j] = a[j];  $\Theta(n)$ 
            delete [] a;  $\Theta(n)$ 
            a = b;
            size = newsize;
        }
        a[i] = i*i;
    }
}

```

$$\begin{aligned}
 \Theta(\text{size}) &= \Theta(3^{k/2}) \rightarrow \text{newsize} = 3^{\text{size}/2} \\
 T(n) &= \sum_{k=0}^{\log_3 n} [\Theta(3^{k/2}) + (n + \log_3 n)] \\
 &= \Theta(3^{\log_3 n / 2}) = \Theta(n^{1/2}) \\
 &= \Theta(\sqrt{n})
 \end{aligned}$$

The first loop runs $n+1$ times. Assuming it loop always executes, the second loop won't run 10 times. Size updates $(3^{\text{size}/2})$ which is $\Theta(3^{\log_3 n})$