HW/SW Co-design of Elliptic Curve Cryptography on 8051 B07901010范詠為 B07901013林子軒

Brief Abstract

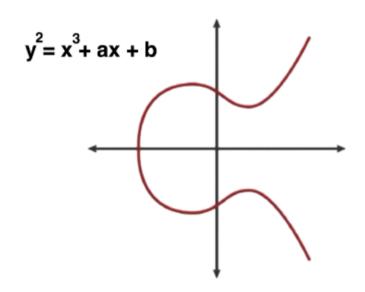
In the field of Information Security, Cryptography is one of the many ways to secure the insecure information channels. In 1985, Neal Koblitz and Victor Miller independently introduced elliptic curve cryptography.

ECC uses the set of points on an elliptic curve along with an addition rule. The unique mathematical structure of the points with the addition rule enables us to perform encryption and decryption of plaintexts. Another reason that supports the feasibility of ECC is the fact that it uses significantly smaller key sizes than the RSA Cryptosystem.

	RSA	ECC
Key Size (security 280)	1024-bits	160-bits
Pros	easy implementation	fast, smaller key size
Cons	slow, longer key size	more complicated

We will try to implement the ECIES (Elliptic Curve Integrated Encryption Scheme) algorithm. Since the performance of 8-bit microcontrollers is often too poor for the implementation of public-key cryptography in software, we will design a hardware accelerator for ECC on an 8051 microcontroller

Algorithm



Algorithm 2 Point doublingINPUT: point $P(x_1, y_1)$.4: $\lambda \leftarrow x_1 \oplus y_1/x_1$ OUTPUT: point 2P.5: $X \leftarrow \lambda^2 \oplus \lambda \oplus a$ 1: if P = -P or $P = \mathcal{O}$ then6: $Y \leftarrow x_1^2 \oplus \lambda \cdot X \oplus X$ 2: return \mathcal{O} 7: return (X, Y)3: else8: end if

Algorithm 3 Point addition

```
INPUT: points P(x_1, y_1), Q(x_2, y_2).
                                                                8:
OUTPUT: point P+Q.
                                                                              \lambda \leftarrow (y_2 \oplus y_1)/(x_2 \oplus x_1)
                                                                9:
 1: if P \neq Q then
                                                                              X \leftarrow \lambda^2 \oplus \lambda \oplus x_1 \oplus x_2 \oplus a
                                                                10:
         if P = -Q then
                                                                              Y \leftarrow \lambda \cdot (x_1 \oplus X) \oplus X \oplus y_1
 2.
                                                               11:
             \mathbf{return}~\mathcal{O}
                                                               12:
                                                                              return (X,Y)
 3:
         else if P = \mathcal{O} then
                                                                         end if
                                                               13:
 4:
         return Q
                                                               14: else
 5:
         else if Q = \mathcal{O} then
                                                                         return 2P
                                                               15:
 6:
             return P
                                                               16: end if
 7:
```

Algorithm 4 Scalar multiple

```
INPUT: point P, integer n.

5: R \leftarrow R + A

OUTPUT: point nP.

6: end if

1: A \leftarrow P

7: n \leftarrow n \gg 1

2: R \leftarrow \mathcal{O}

8: A \leftarrow 2A

3: while n > 0 do

9: end while

4: if n \equiv 1 \mod 2 then

10: return R
```

Algorithm 5 Encryption with the simplified ECIES

```
Input: plaintext x.
                                                            8: k \leftarrow \text{RANDOM}([1, n-1])
OUTPUT: ciphertext (U(x_1, y_1), y).
                                                           9: U(x_1, y_1) \leftarrow kP
                                                          10: V(x_2, y_2) \leftarrow kQ
 1: for char \in x do
                                                          11: for char \in x' do
        char \leftarrow BinaryASCII(char)
        char \leftarrow Padding(char)
                                                                   cipher \leftarrow char \cdot x_2
 3:
                                                          12:
        APPEND(x',char)
                                                                   Append(y, cipher)
 4:
                                                          13:
 5: end for
                                                          14: end for
 6: blocklength \leftarrow \lfloor N/7 \rfloor
                                                          15: return (U, y)
 7: x' \leftarrow \text{Block}(x', blocklength)
```

References

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