CH08-320201 Algorithms and Data Structures

Prof. Dr. Michael Sedlmair

Jacobs University Spring 2018

Overview today

- 0. Introduction
 - 0.1 Syllabus and Organization
 - 0.2 Goals
 - 0.3 Content
- 1. Foundations
 - 1.1 Definitions
 - 1.2 First example: Insertion Sort

0. Introduction

0.1 Syllabus and Organization

Online resources

Course Website

- https://stevenabreu7.github.io/ads-website/
- infos, slides, homework, etc.
- NOTE: server will change —> stay tuned

Moodle

- https://moodle.jacobs-university.de/course/view.php?id=744
- homework submissions, scores, announcements, forum, etc.
- NOTE: "beta" version
 - you might not have access or even see the course yet
 - we will set up all the details over the next days/weeks
 - —> don't worry / please give us some time

Content

- This course introduces a basic set of data structures and algorithms that form the basis of almost all computer programs. The data structures and algorithms are analyzed in respect to their computational complexity with techniques such as worst case and amortized analysis.
- Topics: Fundamental data structures (lists, stacks, trees, hash tables), fundamental algorithms (sorting, searching, graph traversal).

Prerequisites

• C or C++ or python

```
#Include <SIGIO.h >
int main(void)

{
  int count;
  for (count = 1; count <= 500; count ++)
     printf ("I will not Throw paper dirplanes in class.");
  return 0;
}

MEND 10-3
```

Who "we" are:

• TAs:

- Steven Abreu: s.abreu [@jacobs-university.de]
- Oana Miron: o.miron [@jacobs-university.de]
- Mohit Shrestha: mo.shrestha [@jacobs-university.de]

Lecturer

- Michael Sedlmair: old webpage / new webpage (coming soon)
- Office: Res I, 168.
- Phone: 3051
- E-Mail: m.sedlmair [@jacobs-university.de]
- Office hours: by appointment (usually Mon/Tue)

Lectures

- Times: Tuesdays
 - 08:15am 09:30am **AND**
 - 09:45am 11:00am
- Location: Research II Lecture Hall

Tutorials / Q&As

Tutorials:

- Time: Thursdays 11:15am 12:30am
- Location: Research II Lecture Hall
- What: Discussing last homework & questions current content/homework (try to ask them already here!)

Q&A:

- Time: Weekend (tbd)
- Where: Colleges (tbd)
- What: questions current content/homework

Assignments

Homeworks:

- The homework assignments include theoretical and practical problems that tackle topics from the lectures.
- The homework assignments are handed out on a regular basis (goal: in sync with the lectures).

Submitting your homework

- Solutions that are handed in late lead to reduced credit (-15% per day).
- Exceptions are only made with an official excuse.
- With an official excuse of up to 4 days, the deadline for the respective homework is extended by the same amount of days.
- With an official excuse of more than 4 days, the respective homework will not count.
- Handing in via Moodle.

Exams

- midterm data and location tbd
- final exam data and location tbd
- no quizzes

Grading

- Homework: 35%
- Midterm exam: 25%
- Final exam: 40%

Dates — Lectures

Date	Content		
01.02. (Thu)	Foundations – definition, introductory example		
06.02. (!!! 9:45am !!!)	Foundations – asymptotic analysis		
13.02.	Foundations – divide & conquer as first concept,		
	Foundations – recurrences		
20.02.	Foundations – applying recurrences to divide & conquer		
	Foundations – applying recurrences to divide & conquer		
27.02.	Sorting & Searching – heap sort		
	Sorting & Searching – heap sort		
06.03.	Sorting & Searching – quick sort		
	Sorting & Searching – randomized quick sort		
13.03.	Sorting & Searching – lower bounds		
	Sorting & Searching – sorting in linear time, searching		
20.03.	Fundamental Data Structures – array, stack, queue,		
	Fundamental Data Structures – lists, rooted tree, BST		
27.03	Spring break		

Dates — Lectures

03.04.	Fundamental Data Structures – red-black tree		
	Fundamental Data Structures – red-black tree		
10.04.	Fundamental Data Structures – hash tables		
	Fundamental Data Structures – hash tables		
17.04.	Design Concepts – dynamic programming		
	Design Concepts – greedy algorithms		
24.04.	Graph Algorithms – representations, graph searches		
	Graph Algorithms – minimum spanning tree		
01.05.	Holiday		
08.05.	Graph Algorithms – shortest paths		
	Graph Algorithms – maximum flow		
15.05.	Computational Geometry – line intersections		
	Computational Geometry – convex hull / Voronoi diagrams		

Might slightly change

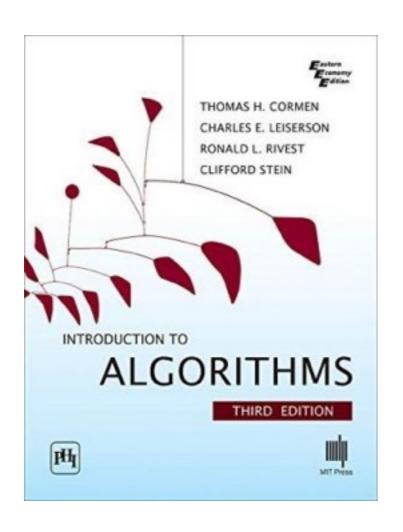
Dates — Assignments

Handout	Deadline	
06.02.	12.02., 23:59	Homework 1 (Foundations)
13.02.	19.02., 23:59	Homework 2 (Foundations)
20.02.	26.02., 23:59	Homework 3 (Foundations)
27.02.	05.03., 23:59	Homework 4 (Sorting & Searching)
06.03.	12.03., 23:59	Homework 5 (Sorting & Searching)
13.03.	19.03., 23:59	Homework 6 (Sorting & Searching)
20.03.	02.04., 23:59	Homework 7 (Fundamental Data Structures)
03.04.	09.04., 23:59	Homework 8 (Fundamental Data Structures)
10.04.	16.04., 23:59	Homework 9 (Fundamental Data Structures)
17.04.	23.04., 23:59	Homework 10 (Design Concepts)
24.04.	07.05., 23:59	Homework 11 (Graph Algorithms)
08.05.	14.05., 23:59	Homework 12 (Graph Algorithms)
not graded		Homework 13 (Computational Geometry)

Might slightly change

Literature

Introduction to Algorithms. Thomas H. Corman, Charles E. Leiserson, Ronald L. Rivest, and Clifford Stein, 3rd edition, MIT Press, 2009.



—> ebook in the library

0.2 Goals

Goals

- The objective of the course is to learn about
 - fundamental algorithms for solving problems efficiently,
 - basic algorithmic concepts,
 - the analysis of algorithms, and
 - fundamental data structures for efficiently storing, accessing, and modifying data.

0.3 Content

Content

- Foundations
- Sorting & Searching
- Fundamental Data Structures
- Design Concepts
- Graph Algorithms
- Computational Geometry

1. Foundations

1.1 Definitions

Definition: Algorithm

- An algorithm is a sequence of computational steps which transforms a set of values (input) to another set of values (desired output).
- It is a tool for solving a well-defined computational problem.
- Step-wise procedure that can be implemented in a computer program.
- Consists of a finite list of well-defined instructions (Turing machine).
- 'Algorithm' stems from 'Algoritmi', the Latin form of al-Khwārizmī, a Persian mathematician, astronomer and geographer.

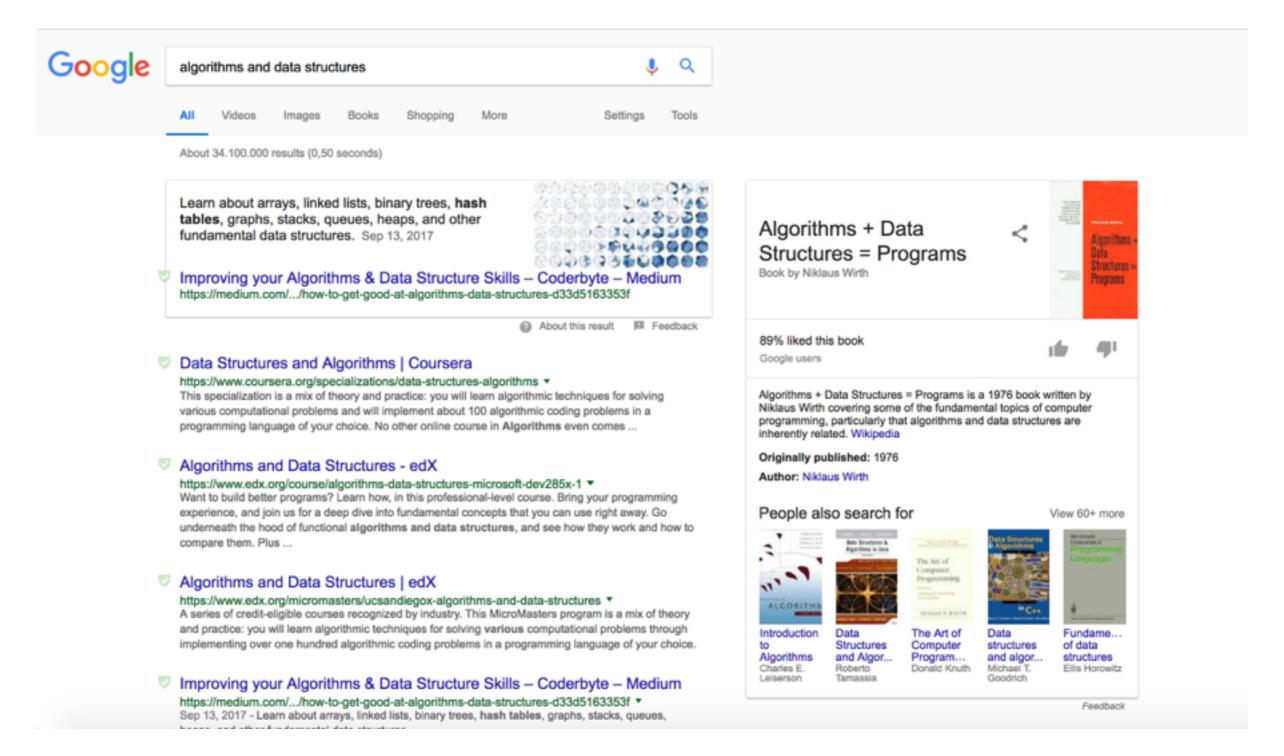
Example: Sorting Problem

- Input: sequence $\langle a_1, a_2, ..., a_n \rangle$ of numbers.
- Output permutation $\langle a_1', a_2', ..., a_n' \rangle$ such that $a_1' \leq a_2' \leq ... \leq a_n'$
- Example (instance of sorting problem):

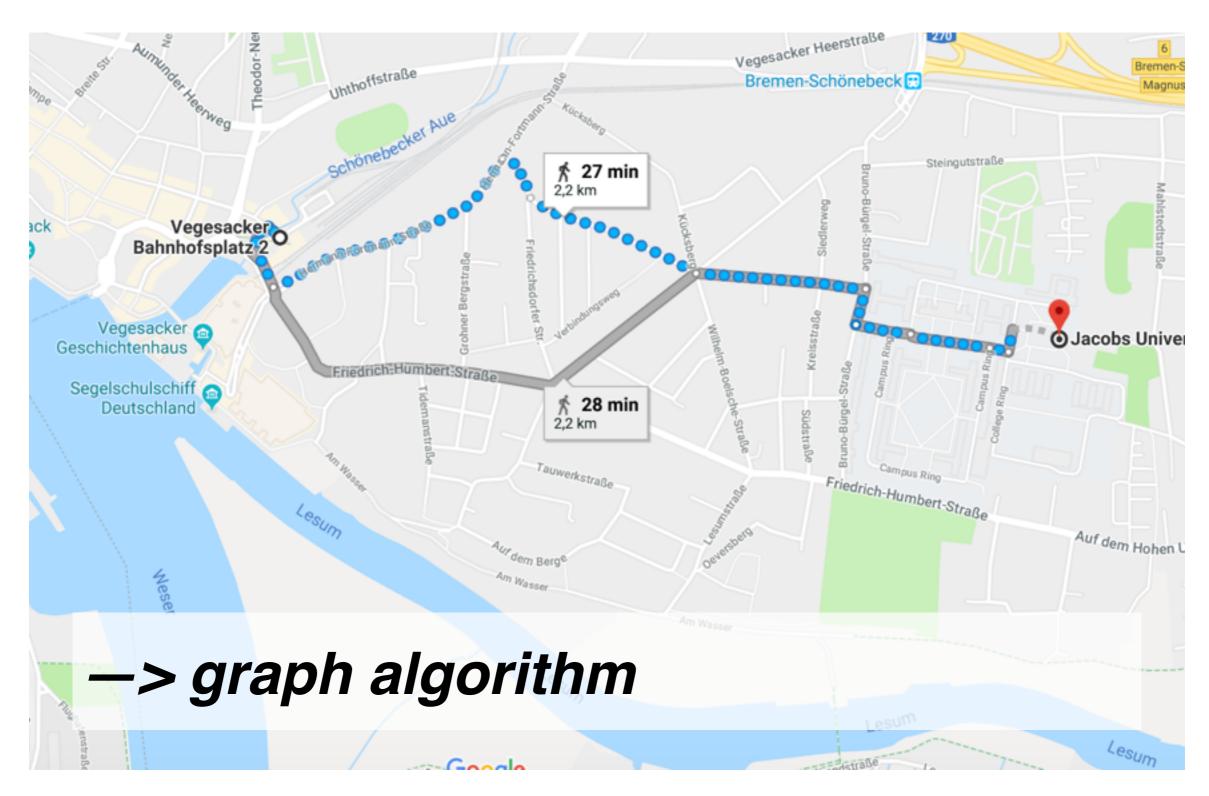
Input: 824936

Output: 234689

Example: Searching

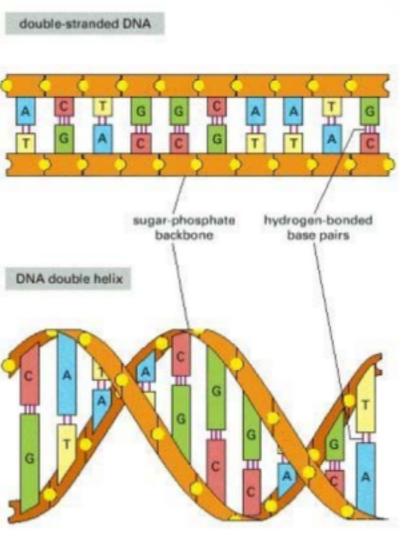


Example: Road map



Example: DNA Sequences





-> string matching

Analysis of algorithms

- The theoretical study of computer-program performance and resource usage.
- Other design goals?
 - correctness
 - functionality
 - robustness
 - reliability
 - user-friendliness
 - programmer time
 - simplicity
 - modularity
 - maintainability
 - extensibility

Performance of Algorithms

- Analysis helps us to understand scalability.
- Performance often draws the line between what is feasible and what is impossible.
- Algorithmic mathematics provides a language for talking about program behavior.
- "Performance is the currency of computing."
- The lessons of program performance generalize to other computing resources.

Definition: Data Structure

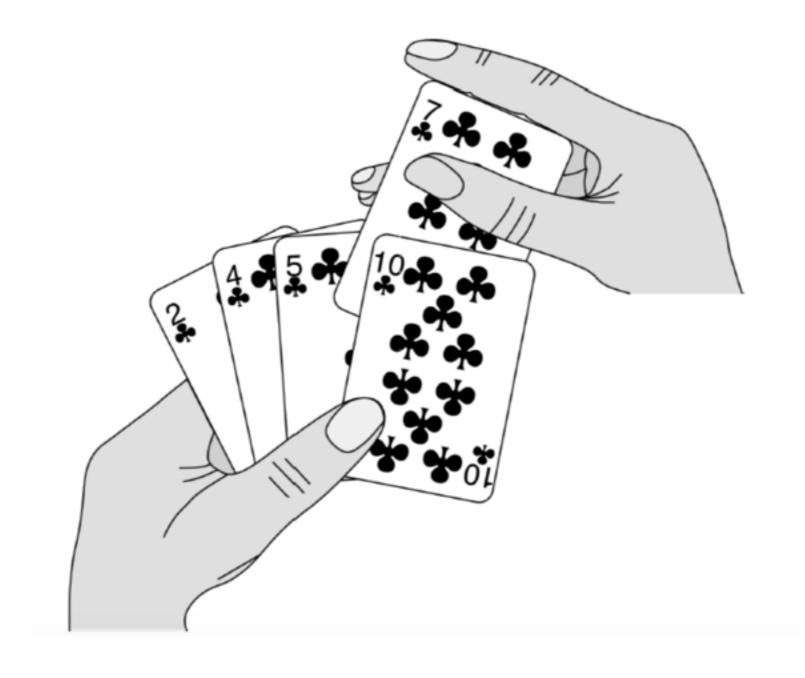
 A data structure is a way to store and organize data in order to facilitate access and modification.

- There is typically no best data structure, but each data structure has its strengths and weaknesses.
- Which data structure to use, depends on the problem that is to be solved.
- Sometimes there is a trade-off between storage (in a data structure) and speed (in accessing a data structure or of an algorithm)

1.2 First example: Insertion Sort

Sorting problem

First algorithm: Insertion Sort

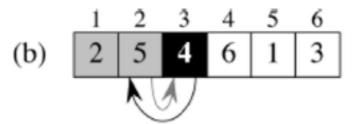


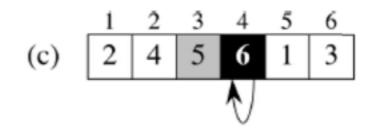
Insertion Sort

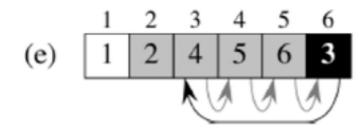
```
INSERTION-SORT (A, n)
 for j = 2 to n
     key = A[j]
     // Insert A[j] into the sorted sequence A[1...j-1].
     i = j - 1
     while i > 0 and A[i] > key
         A[i + 1] = A[i]
          i = i - 1
     A[i+1] = key
```

Example

• Sort A = < 5, 2, 4, 6, 1, 3 >







INSERTION-SORT (A, n)

for
$$j = 2$$
 to n
 $key = A[j]$
// Insert $A[j]$ into the sorted sequence $A[1..j-1]$.
 $i = j-1$
while $i > 0$ and $A[i] > key$
 $A[i+1] = A[i]$
 $i = i-1$
 $A[i+1] = key$

Correctness

```
INSERTION-SORT (A, n)

for j = 2 to n

key = A[j]

// Insert A[j] into the sorted sequence A[1..j-1].

i = j-1

while i > 0 and A[i] > key

A[i+1] = A[i]

i = i-1

A[i+1] = key
```

Loop invariant:

At the start of each iteration of the for loop, the subarray A[1..j-1] consists of elements originally in A[1..j-1], but in sorted order.