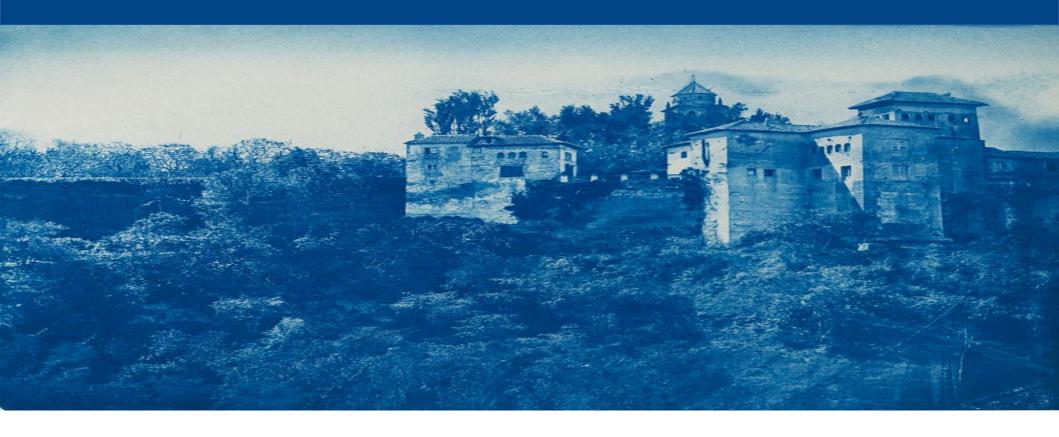
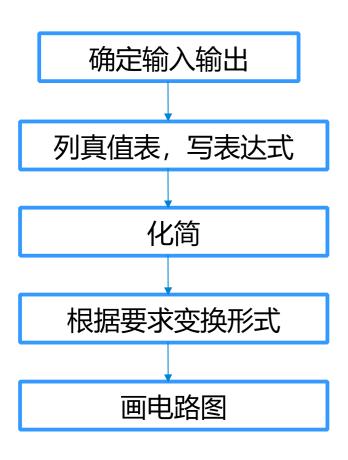
# 数字逻辑第五章



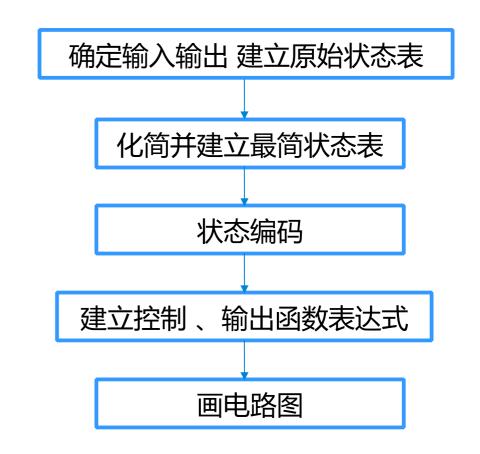
信息科学与工程学院计算机系 杨永全 yangyq@ouc.edu.cn

### 同步时序线路的设计

#### 组合线路设计方法

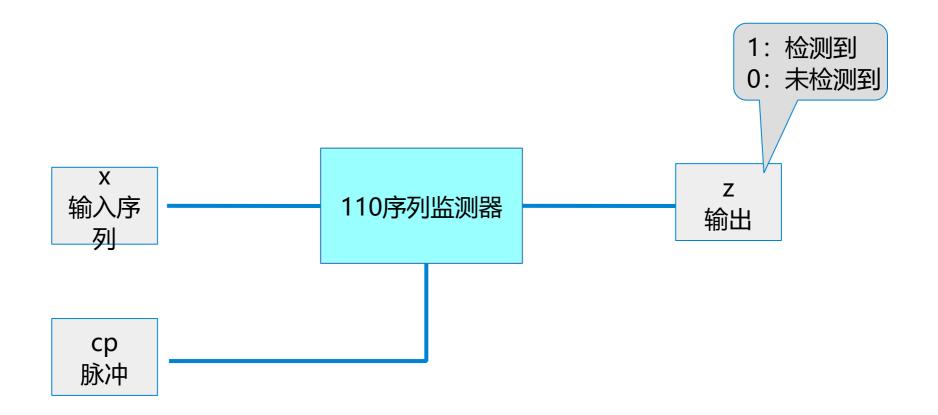


#### 时序线路设计方法



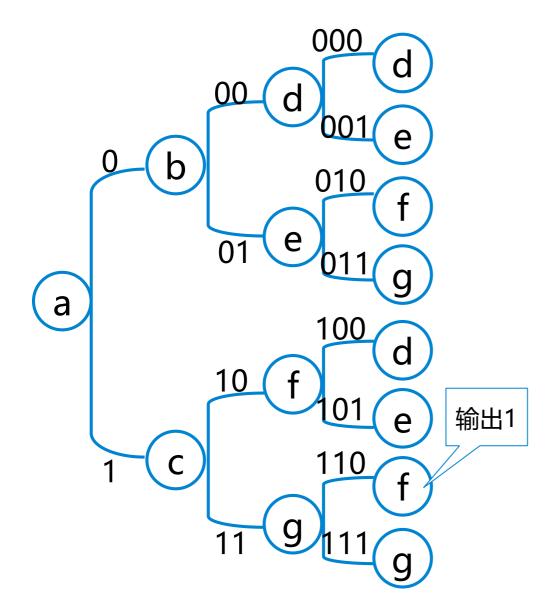
# 引例

- 题目:用与非门和JK触发器设计一个同步时序线路,检测输入为连续的110
- 一、确定输入、输出, 建立原始状态表



# 原始状态表

• 设初始状态为a



#### 原始状态表

S	0	1
a	b	С
b	d	е
С	f	g
d	d	е
е	f	g
f	d	е
g	f	g

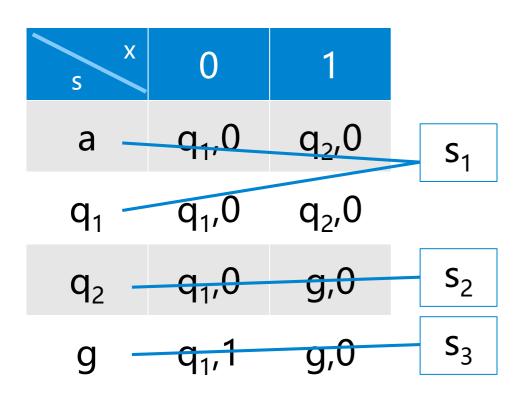
# 化简,建立最简状态表

S	0	1	
a	b,0	c,0	
b —	<del>d,0</del>	e,0	$q_1$
C	f,0	g,0	
d /	d,0	e,0	
e	1,0	g,0	$\rightarrow$ $q_2$
f /	d,0	e,0	
g	f,1	g,0	

S X	0	1
a	q <sub>1</sub> ,0	q <sub>2</sub> ,0
$q_1$	q <sub>1</sub> ,0	q <sub>2</sub> ,0
$q_2$	q <sub>1</sub> ,0	g,0
g	q <sub>1</sub> ,1	g,0

# 化简,建立最简状态表

#### • 得到最终的最简状态表



S	0	1
S <sub>1</sub>	s <sub>1</sub> ,0	s <sub>2</sub> ,0
$S_2$	s <sub>1</sub> ,0	s <sub>3</sub> ,0
<b>S</b> <sub>3</sub>	s <sub>1</sub> ,1	s <sub>3</sub> ,0

# 状态编码

• 三个状态,需要? 位编码

最终得到状态表

	$S_1$	$S_2$	$S_3$
$Y_2$	0	1	1
$Y_1$	0	0	1

S	0	1
00	00,0	10,0
10	00,0	11,0
11	00,1	11,0

# 确定输出及控制函数

根据右侧的状态激励表,可以得到 状态转移表,最终得到控制及输出 函数真值表 激励表 (JK触发器)

X	<b>y</b> <sub>2</sub>	y <sub>1</sub>	y <sub>2</sub> n+	·1 y <sub>1</sub> n+1	J <sub>2</sub> K <sub>2</sub>	J <sub>1</sub> K <sub>1</sub>	Z
0	0	0	0	0	Ф0	0Ф	0
0	1	0	0	0	Ф1	Ф0	0
0	1	1	0	0	Ф1	Ф1	1
1	0	0	1	0	1Ф	Ф0	0
1	1	0	1	1	Ф0	1Ф	0
1	1	1	1	1	Ф0	Ф0	0
0	0	1	Ф	Ф	ФФ	ФФ	Ф
1	0	1	Ф	Ф	ФФ	ФФ	Ф

Q	Qn+1	JK
0	0	0 Ф
0	1	1 Ф
1	0	Ф 1
1	1	Ф 0

### 确定输出及控制函数

根据控制及输出函数真值表,得到 JK、Z和输入之间的关系:

$$-z = \sum(3) + \sum \Phi(1,5)$$

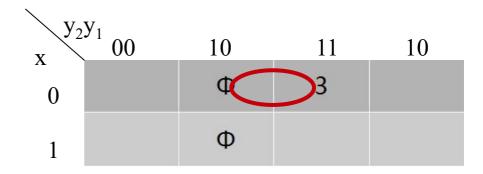
$$-$$
 J1= $\sum$ (6)+ $\sum$ \Phi(1,3,5,7)

- 
$$K1 = \sum(3) + \sum \Phi(0,1,2,4,5,6)$$

$$-$$
 J2= $\sum$ (4)+ $\sum$ \Phi(1,2,3,5,6,7)

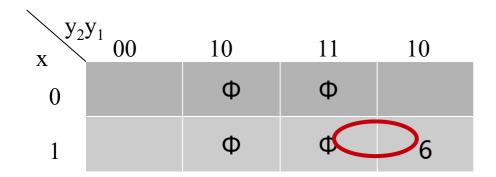
- 
$$K2 = \sum (2,3) + \sum \Phi(0,1,4,5)$$

#### 使用卡诺图化简



$$z = \overline{x} y_1$$

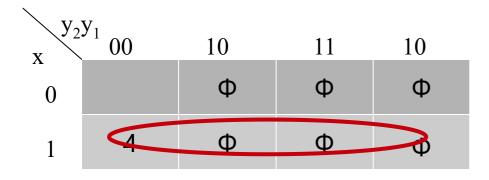
# 通过化简确定输入输出函数

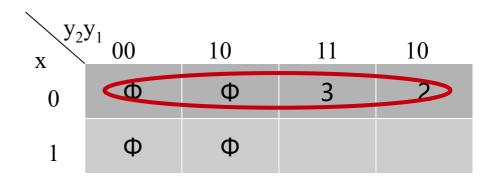


y <sub>2</sub> ;	y <sub>1</sub> 00	10	11	10
0		Ф	3	
1	Ф	Ф		Ф

$$J_1 = x y_2$$

$$K_1 = \overline{x}$$

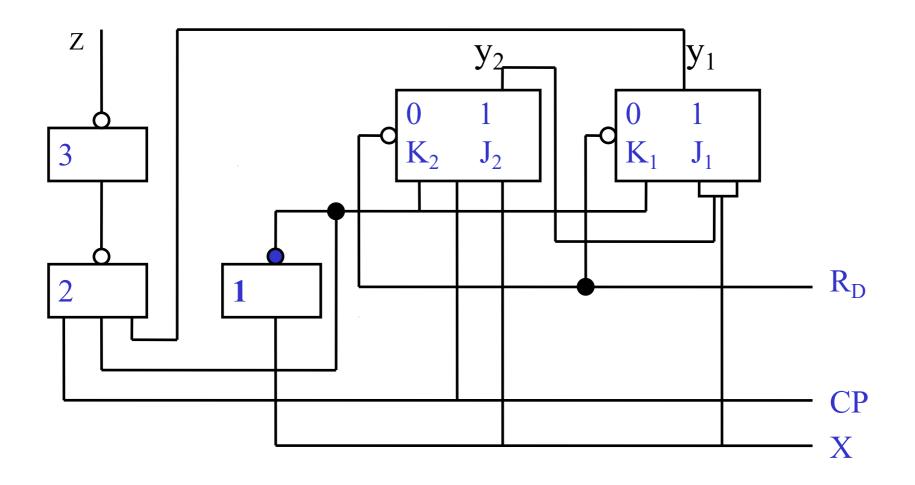




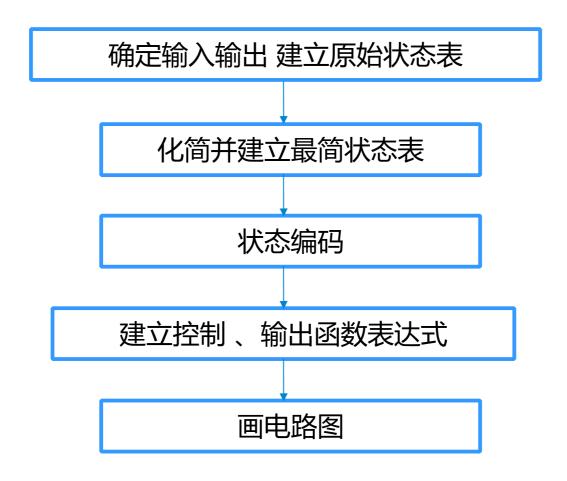
$$J_2 = x$$

$$K_2 = \overline{x}$$

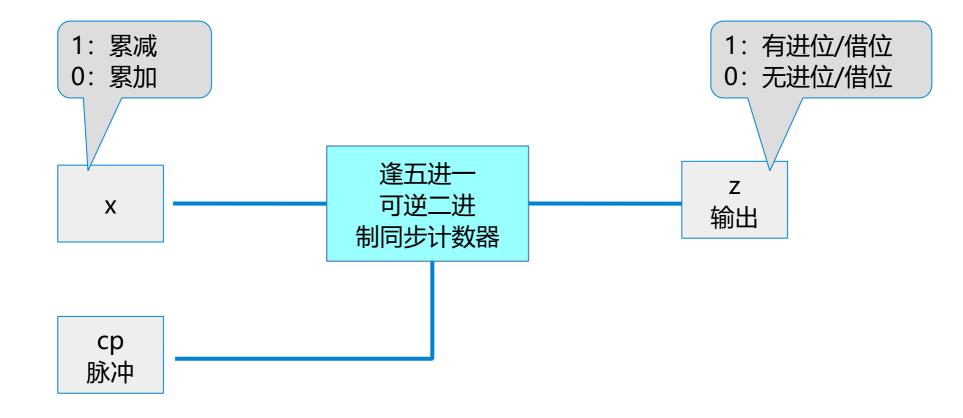
# 画逻辑电路图



# 时序线路设计步骤总结



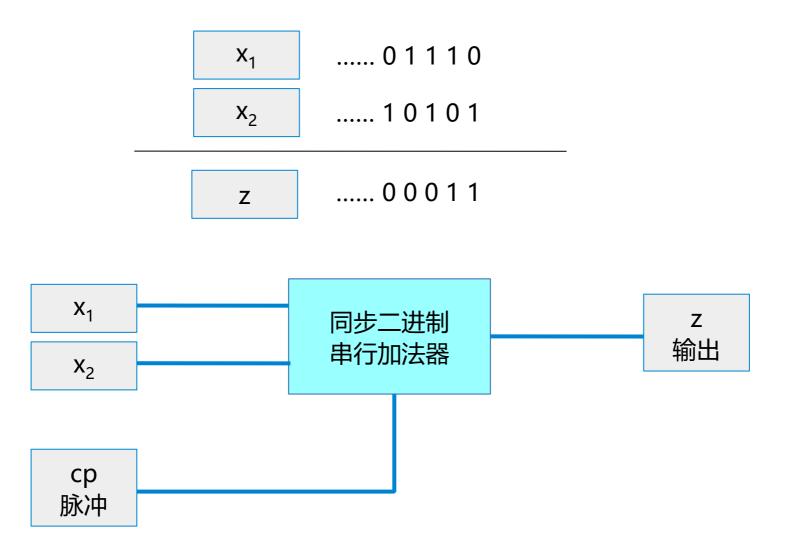
• 例1:建立逢五进一可逆二进制同步计数器



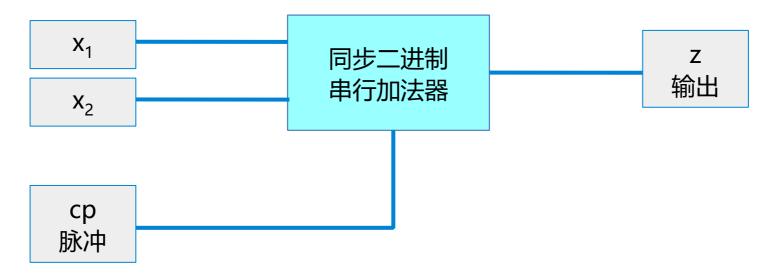
• 例1: 建立逢五进一可逆二进制同步计数器

S	0	1
а	b,0	e,0
b	c,0	a,1
С	d,0	b,0
d	e,0	c,0
е	a,1	d,0

• 例2:同步二进制串行加法器

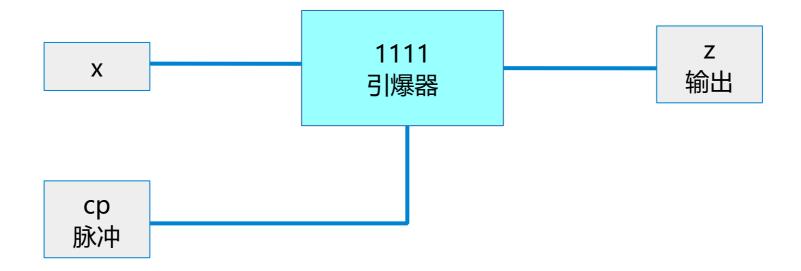


#### • 例2:同步二进制串行加法器



x <sub>1</sub> x <sub>2</sub>	00	01	10	11
а	a,0	a,1	a,1	b,0
b	a,1	b,0	b,0	b,1

• 例3:引爆条件'1111'



• 例3:引爆条件'1111'

S	0	1	S	0	1
a	a,0	b,0	a	a,0	b,0
b	Ф,Ф	c,0	b	a,0	c,0
С	Ф,Ф	d,0	С	a,0	d,0
d	Ф,Ф	Ф,1	d	a,0	Ф,1

考虑输入为0恢复初始状态

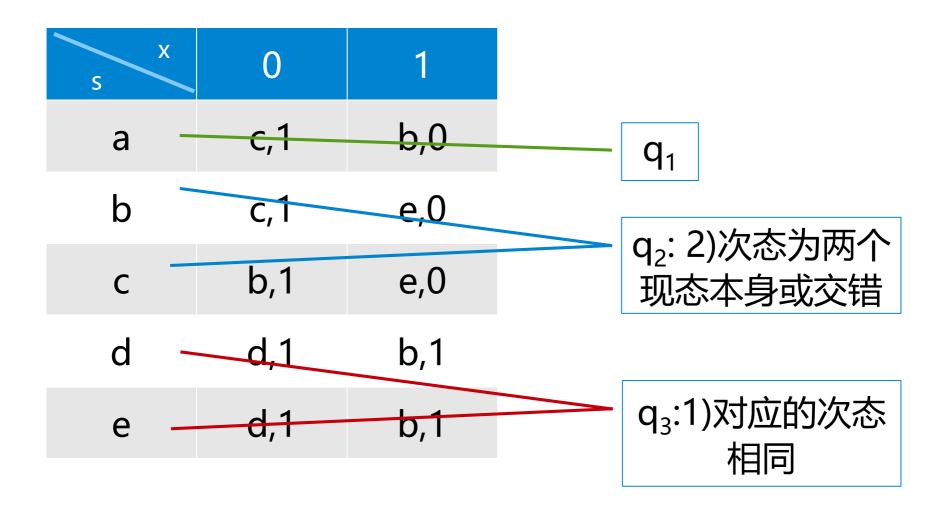
- 一、化简原理: 找出等价状态并将它们合并
- 等价状态: 1、必要条件: 在同样的输入作用下, 有相同的输出
  - 2、同样的输入条件下,相应的次态彼此等价
- 等价次态: 1、对应的次态相同
  - 2、次态为两个现态本身或交错
  - 3、两个次态为状态对封闭链中的一对
  - 4、两个次态的某一后续状态对可以合并

等价状态具有可传递性: AB等价, AC等价 = >BC等价, 则A、B、C为等价类

- 等价类: 彼此等价的状态的集合
- 最大等价类:一个等价类不包含在任何其他等价类中

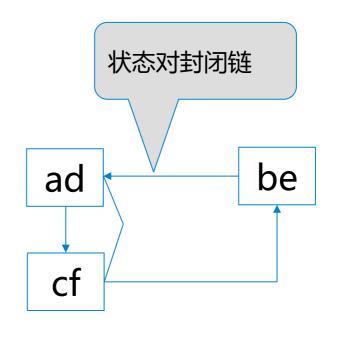
化简原始状态表 = >寻找最大等价类

#### • 例1



#### • 例2

S	0	1
a	e,0	d,0
b	a,1	f,O
С	c,0	a,1
d	b,0	a,0
е	d,1	c,0
f	c,0	d,1
g	h,1	g,1
h	c,1	b,1



$$q_1=(a, d)$$
  
 $q_2=(b, e) q_3=(c, f)$   
 $q_4=(g) q_5=(h)$ 

k次划分法: 先找出输出相同的状态集合,为一次划分,再找第二次输入下输出也相同的集合,依此类推,直到k+1次不能再划分,找到最大等价类集合。

#### • 等价次态:

- a) 对应的次态相同
- b) 次态为两个现态本身或交错
- c) 两个次态为状态对封闭链中的一对
- d) 两个次态的某一后续状态对可以合并

S X	0	1
a	c,0	b,1
b	f,O	a,1
С	d,0	g,0
d	d,1	e,0
е	c,0	e,1
f	d,0	g,0
g	c,1	d,0

• 1、一次划分

$$- q_1 = \{ a_{21}, b_{21}, e_{21} \}$$

$$- q_2 = \{ c_{33}, f_{33} \}$$

$$- q_3 = \{ d_{31}, g_{23} \}$$

• 2、二次划分(把一次划分中下标不同的分出来)

$$- q_1 = \{ a_{21}, b_{21}, e_{21} \}$$

$$- q_2 = \{ c_{33}, f_{33} \} = \{ c_{34}, f_{34} \}$$

$$- q_3 = \{ d_{31} \}$$

$$- q_4 = \{ g_{23} \}$$

• 3、所有下标都相同,划分结束。

S X	0	1
a	e,0	d,0
b	a,1	f,0
С	f,0	a,1
d	b,0	a,0
е	d,1	c,0
f	c,0	d,1
g	h,1	g,1
h	c,1	b,1

• 1、一次划分

$$- q_1 = \{ a_{21}, d_{21} \}$$

$$- q_2 = \{ b_{13}, e_{13} \}$$

$$- q_3 = \{ c_{31}, f_{31} \}$$

$$- q_4 = \{ g_{44}, h_{32} \}$$

• 2、二次划分

$$- q_1 = \{ a_{21}, d_{21} \}$$

$$- q_2 = \{ b_{13}, e_{13} \}$$

$$- q_3 = \{ c_{31}, f_{31} \}$$

$$- q_4 = \{ g_{54} \}$$

$$- q_5 = \{ h_{32} \}$$

• 3、划分结束

- 隐含表法: 是一种直角三角形表格, 表中每一个小格表示一个状态对的等价或 不等价关系。
- 第一步: 做隐含表, 然后顺序比较。

X <sub>1</sub> X <sub>2</sub> S	00	01	11	10	b	X		1		
a	d,0	d,0	f,O	a,0	С	X	af		1	
b	c,1	d,0	e,1	f,0	d	bd af	X	X		
С	c,1	d,0	e,1	a,0	е	Х	df	df	X	
d	d,0	b,0	a,0	f,O	f	√	af X	X	bd	•
е	c,1	f,0	e,1	a,0	'	dg			bg	
f	d,0	d,0	a,0	f,0	g	af	X	X	af	)
g	g,0	g,0	a,0	a,0	h	X	bc af	bc	X	k c
h	b,1	d,0	e,1	a,0		а	b	С	d	•

X						
X	af					
bd af	X	X				
X	df af	df	X			
√	Χ	Χ	bd	X		
dg af	X	X	bg af	X	dg af	
X	bc af	bc	X	bc df	X	X
а	b	С	d	е	f	g

• 第二步:关连比较。继续检查填有隐含条件的那些方格。若检查发现所填的隐含条件肯定不能满足,就在该方格内打 "×"

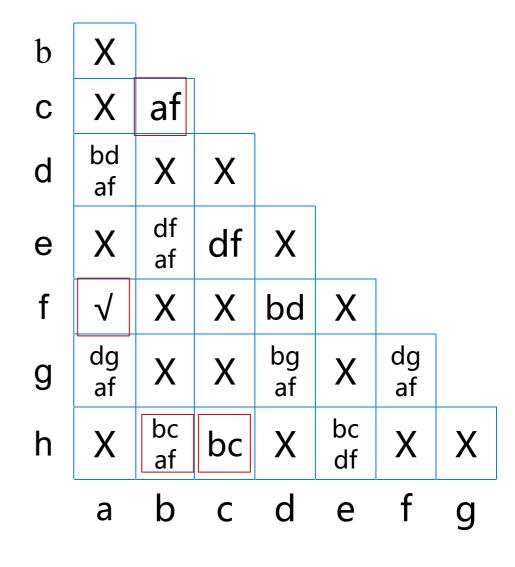
X <sub>1</sub> X <sub>2</sub> S	00	01	11	10	b	X						
a	d,0	d,0	f,0	a,0	С	X	af					
b	c,1	d,0	e,1	f,O	d	bd af	X	X				
С	c,1	d,0	e,1	a,0		X	df	df	X			
d	d,0	b,0	a,0	f,0	e		af				]	
е	c,1	f,0	e,1	a,0	f	<b>√</b>	X	X	bd	X		]
f	d,0	d,0	a,0	f,0	g	dg af	X	X	bg af	X	dg af	
g	g,0	g,0	a,0	a,0	h	X	bc af	bc	X	bc df	X	X
h	b,1	d,0	e,1	a,0		a	b	С	d	е	f	g

• 第三步: 寻找最大等价类 未打 "×"的方格, 都代表一个等价状态对 由此得到全部等价对:

#### 部最大等价类:

[a, f]、[b, c, h]、 [d]、[e]、[g]

• 第四步 状态合并,得最简状态表



#### • 隐含表法总结

- 1、构作隐含表①等价√②不等价×③条件
- 2、顺序比较追踪,找出所有等价状态
- 3、形成最大等价类集合
- 4、构成最简状态表

### 状态编码

• 确定需要几位二进制码

k = [log<sub>2</sub>N] 向上取整

- 次佳编码法:
  - 1)次态相同,现态相邻
  - 2)现态相同,次态相邻
  - 3)输出相同,现态相邻

优先顺序1>2>3

### 状态编码

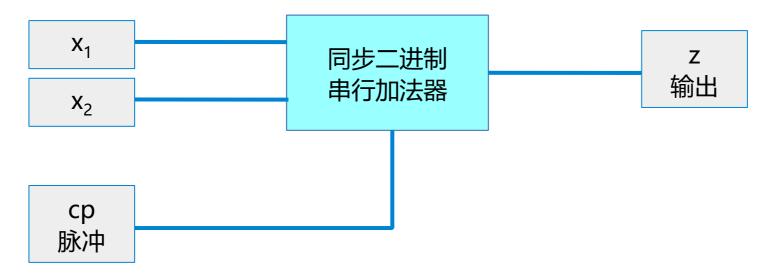
• 例:为下列最简状态表进行编码

S	0	1
a	c,0	d,0
b	c,0	a,0
С	b,0	d,0
d	a,1	b,1

- log<sub>2</sub>4=2 所以需两位二进制码
  - 1)次态相同,现态相邻 ab 相邻, ac 相邻
  - 2)现态相同,次态相邻cd, ca, bd, ab 相邻
  - 3)输出相同,现态相邻 abc 相邻

例1:用JK触发器及与非门设计一个同步二进制串行加法器

•一、确定输入输出,建立原始状态表



x <sub>1</sub> x <sub>2</sub>	00	01	10	11
а	a,0	a,1	a,1	b,0
b	a,1	b,0	b,0	b,1

- 二、化简略
- 三、状态编码 a=0 b=1

x <sub>1</sub> x <sub>2</sub>	00	01	10	11
a	0,0	0,1	0,1	1,0
b	0,1	1,0	1,0	1,1

• 四、列控制、输出函数表达式

Q Qn+1	J	K
0 0	0	Ф
0 1	1	Φ
1 0	Ф	1
1 1	Ф	0

<b>X</b> <sub>1</sub>	<b>X</b> <sub>2</sub>	у	y <sup>n+1</sup>	J	K	Z
0	0	0	0	0	Ф	0
0	0	1	0	Ф	1	1
0	1	0	0	0	Ф	1
0	1	1	1	Ф	0	0
1	0	0	0	0	Ф	1
1	0	1	1	Ф	0	0
1	1	0	1	1	Ф	0
1	1	1	1	Ф	0	1

• 四、列控制、输出函数表达式

$$z = x_1 \overline{x_2} y + \overline{x_1} x_2 \overline{y} + x_1 \overline{x_2} \overline{y} + x_1 x_2 y$$

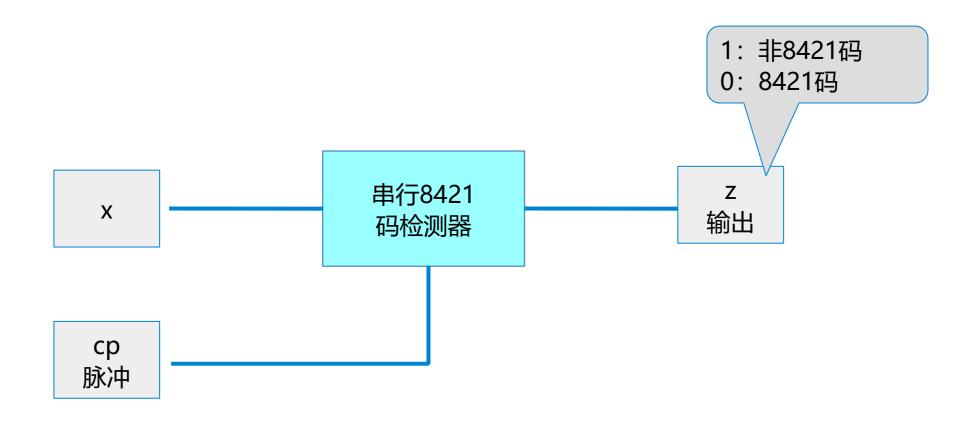
$$J = x_1 y$$

$$K = \overline{x_1} \overline{y}$$

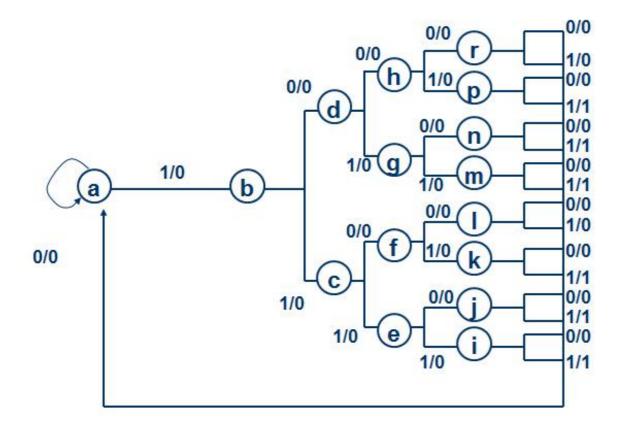
<b>x</b> <sub>1</sub>	<b>X</b> <sub>2</sub>	у	y <sup>n+1</sup>	J	K	Z
0	0	0	0	0	Ф	0
0	0	1	0	Ф	1	1
0	1	0	0	0	Ф	1
0	1	1	1	Ф	0	0
1	0	0	0	0	Ф	1
1	0	1	1	Ф	0	0
1	1	0	1	1	Ф	0
1	1	1	1	Ф	0	1

例2用与非门,与或非门及JK触发器,设计串行8421码检测器。

•一、确定输入输出,建立原始状态表



#### • 一、确定输入输出,建立原始状态表



S	0	1
a	a,0	b,0
b	d,0	c,0
С	f,0	e,0
d	h,0	g,0
e	j,0	i,0
f	1,0	k,0
g	n,0	m,0
h	r,0	p ,0
i	a,0	a,1
j	a,0	a,1
k	a,0	a,1
I	a,0	a,0
m	a,0	a,1
n	a,0	a,1
р	a,0	a,1
r	a,0	a,0

#### • 二、状态表化简(K次划分法)

$$q_1 = (a, b, c, d, e, f, g, h, l, r)$$
  
 $q_2 = (i, j, k, m, n, p)$ 

#### 更新下标:

$$q_1 = (a_{11}, b_{11}, c_{11}, d_{11}, f_{12}, g_{22}, h_{12}, l_{11}, r_{11})$$
  
 $q_2 = (p)$ 

#### 第二次划分:

$$q_1 = (a_{11}, b_{11}, c_{11}, d_{11}, l_{11}, r_{11})$$

$$q_2 = (e_{22}, g_{22})$$

$$q_3 = (f_{12}, h_{12})$$

$$q_4 = (p)$$

S	0	1
a	a,0	b,0
b	d,0	c,0
С	f,0	e,0
d	h,0	g,0
е	j,0	i,0
f	1,0	k,0
g	n,0	m,0
h	r,0	p ,0
i	a,0	a,1
j	a,0	a,1
k	a,0	a,1
I	a,0	a,0
m	a,0	a,1
n	a,0	a,1
р	a,0	a,1
r	a,0	a,0

#### • 更新下标

$$q_1 = (a_{11}, b_{11}, c_{32}, d_{32}, l_{11}, r_{11})$$

$$q_2 = (e_{44}, g_{44})$$

$$q_3 = (f_{14}, h_{14})$$

$$q_4 = (p)$$

#### 第三次划分:

$$q_{1} = (a_{11}, b_{11}, l_{11}, r_{11})$$

$$q_{2} = (c_{32}, d_{32},)$$

$$q_{3} = (e_{44}, g_{44})$$

$$q_{4} = (f_{14}, h_{14})$$

$$q_{5} = (p)$$

S	0	1
a	a,0	b,0
b	d,0	c,0
С	f,0	e,0
d	h,0	g,0
е	j,0	i,0
f	1,0	k,0
g	n,0	m,0
h	r,0	p ,0
i	a,0	a,1
j	a,0	a,1
k	a,0	a,1
I	a,0	a,0
m	a,0	a,1
n	a,0	a,1
р	a,0	a,1
r	a,0	a,0

#### • 更新下标

$$q_{1} = (a_{11}, b_{22}, l_{11}, r_{11})$$

$$q_{2} = (c_{43}, d_{43}, )$$

$$q_{3} = (e_{55}, g_{55})$$

$$q_{4} = (f_{15}, h_{15})$$

$$q_{5} = (p)$$

#### 第四次划分:

$$q_{1} = (a_{11}, l_{11}, r_{11})$$

$$q_{2} = (b_{22})$$

$$q_{3} = (c_{43}, d_{43}, )$$

$$q_{4} = (e_{55}, g_{55})$$

$$q_{5} = (f_{15}, h_{15})$$

$$q_{6} = (p)$$

S	0	1
a	a,0	b,0
b	d,0	c,0
С	f,0	e,0
d	h,0	g,0
е	j,0	i,0
f	1,0	k,0
g	n,0	m,0
h	r,0	p ,0
i	a,0	a,1
j	a,0	a,1
k	a,0	a,1
I	a,0	a,0
m	a,0	a,1
n	a,0	a,1
р	a,0	a,1
r	a,0	a,0

#### • 更新下标

$$q_{1} = (a_{12}, l_{11}, r_{11})$$

$$q_{2} = (b_{33})$$

$$q_{3} = (c_{43}, d_{43}, )$$

$$q_{4} = (e_{55}, g_{55})$$

$$q_{5} = (f_{15}, h_{15})$$

$$q_{6} = (p)$$

#### 第五次划分:

$$q_{1} = (a_{12})$$

$$q_{2} = (b_{33})$$

$$q_{3} = (c_{43}, d_{43}, )$$

$$q_{4} = (e_{55}, g_{55})$$

$$q_{5} = (f_{15}, h_{15})$$

$$q_{6} = (p)$$

$$q_{7} = (l_{11}, r_{11})$$

S	0	1
a	a,0	b,0
b	d,0	c,0
С	f,0	e,0
d	h,0	g,0
e	j,0	i,0
f	1,0	k,0
g	n,0	m,0
h	r,0	p ,0
i	a,0	a,1
j	a,0	a,1
k	a,0	a,1
I	a,0	a,0
m	a,0	a,1
n	a,0	a,1
р	a,0	a,1
r	a,0	a,0

#### • 更新下标

$$q_{1} = (a_{12})$$

$$q_{2} = (b_{33})$$

$$q_{3} = (c_{54}, d_{54},)$$

$$q_{4} = (e_{66}, g_{66})$$

$$q_{5} = (f_{16}, h_{16})$$

$$q_{6} = (p)$$

$$q_{7} = (l_{11}, r_{11})$$

得到最终结果。

S	0	1
a	a,0	b,0
b	d,0	c,0
С	f,0	e,0
d	h,0	g,0
е	j,0	i,0
f	1,0	k,0
g	n,0	m,0
h	r,0	p ,0
i	a,0	a,1
j	a,0	a,1
k	a,0	a,1
I	a,0	a,0
m	a,0	a,1
n	a,0	a,1
р	a,0	a,1
r	a,0	a,0

#### • 更新下标

$$q_{1} = (a_{12})$$

$$q_{2} = (b_{33})$$

$$q_{3} = (c_{54}, d_{54},)$$

$$q_{4} = (e_{66}, g_{66})$$

$$q_{5} = (f_{16}, h_{16})$$

$$q_{6} = (p)$$

$$q_{7} = (l_{11}, r_{11})$$

得到最简状态表。

S X	0	1
<b>q</b> <sub>1</sub>	q <sub>1</sub> ,0	q <sub>2</sub> ,0
$q_2$	q <sub>3</sub> ,0	q <sub>3</sub> ,0
$q_3$	q <sub>5</sub> ,0	q <sub>4</sub> ,0
$q_4$	q <sub>6</sub> ,0	q <sub>6</sub> ,0
<b>q</b> <sub>5</sub>	q <sub>7</sub> ,0	q <sub>6</sub> ,0
$q_6$	$q_{1},0$	q <sub>1</sub> ,1
$q_7$	$q_{1},0$	$q_{1},0$

• 后续步骤不再赘述。