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Listing 1: GEOMETRY

Listing 2: Geometry

```
// C++ routines for computational geometry.
double INF = 1e100:
double EPS = 1e-12;
struct PT {
 double x, y;
  PT() {}
 PT(double x, double y) : x(x), y(y) {}
  PT(const PT \&p) : x(p.x), y(p.y) {}
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
 PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
 PT operator * (double c) const { return PT(x*c, y*c ); }
 PT operator / (double c) const { return PT(x/c, y/c ); }
bool operator <(const PT &p) const { return (x < p.x ? true : (x == p.x && y < p.y)); }
bool operator ==(const PT &p) const { return fabs(x-p.x) < EPS && fabs(y-p.y) < EPS; }
bool operator !=(const PT &p) const { return !((*this) == p); }
};
double dot(PT p, PT q)
                          { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
 return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
 return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
 return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
 double r = dot(b-a,b-a);
 if (fabs(r) < EPS) return a;
 r = dot(c-a, b-a)/r;
 if (r < 0) return a;
 if (r > 1) return b;
 return a + (b-a)*r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                         double a, double b, double c, double d)
```

```
return fabs(a*x+b*v+c*z-d)/sqrt(a*a+b*b+c*c):
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
 return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
 return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
 if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
     return false:
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
 if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
 return true:
//Returns of list of intersection points between segments s1, and s2
//If they do not intersect, the result is an empty vector
//If they intersect at exactly 1 point, the result contains that point
//If they overlap for non-0 distance, the left and right points of that intersection
// are returned
vector<PT> intersect(PT a, PT b, PT c, PT d) {
    vector<PT> res:
    if(fabs(cross(b-a, c-b)) < EPS && fabs(cross(b-a, d-b)) < EPS &&
        fabs(cross(d-c, a-d)) < EPS && fabs(cross(d-c, b-d)) < EPS) {
        PT min_s1 = min(a, b), max_s1 = max(a, b);
        PT min s2 = min(c, d), max s2 = max(c, d);
        if(min_s1 < min_s2) {</pre>
            if(max_s1 < min_s2) return res;//return false;</pre>
        }
        else if (min_s2 < min_s1 && max_s2 < min_s1) return res; //return false;
        PT start = max(min_s1, min_s2), end = min(max_s1, max_s2);
        if(start==end) res.push_back(start);//overlap is one point
            res.push_back(min(start, end));
            res.push_back(max(start, end));
        return res; //return true; //remove overlap code block
    double x1 = (b.x-a.x), y1 = (b.y-a.y), x2 = (d.x - c.x), y2 = (d.y - c.y);
    double u1 = (-y1*(a.x - c.x) + x1 * (a.y - c.y))/(-x2 * y1 + x1 * y2);
    double u2 = (x2 * (a.y - c.y) - y2 * (a.x - c.x)) / (-x2 * y1 + x1 * y2);
    if(u1 >= 0 \&\& u1 <= 1 \&\& u2 >= 0 \&\& u2 <= 1)
        res.push_back(PT((a.x + u2*x1),(a.y+u2*y1))); //return true;
    return res;//return false;
```

```
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
 b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS):
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b=(a+b)/2;
 c = (a + c)/2;
 return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for (int i = 0; i < p.size(); i++){
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
      p[j].y \le q.y && q.y < p[i].y) &&
      q.x < p[i].x + (p[i].x - p[i].x) * (q.y - p[i].y) / (p[i].y - p[i].y))
      c = !c;
 }
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
 for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true:
    return false;
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
 vector<PT> ret;
 b = b-a:
  a = a-c;
  double A = dot(b, b):
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C:
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
 if (D > EPS)
   ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
```

```
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
 vector<PT> ret:
 double d = sqrt(dist2(a, b));
 if (d > r+R || d+min(r, R) < max(r, R)) return ret;</pre>
 double x = (d*d-R*R+r*r)/(2*d):
 double y = sqrt(r*r-x*x);
 PT v = (b-a)/d;
 ret.push_back(a+v*x + RotateCCW90(v)*y);
 if (y > 0)
   ret.push_back(a+v*x - RotateCCW90(v)*y);
 return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
 double area = 0;
 for(int i = 0; i < p.size(); i++) {</pre>
   int j = (i+1) % p.size();
   area += p[i].x*p[j].y - p[j].x*p[i].y;
 return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
 PT c(0,0);
 double scale = 6.0 * ComputeSignedArea(p);
 for (int i = 0; i < p.size(); i++){</pre>
   int j = (i+1) % p.size();
   c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
 return c / scale;
```

Listing 3: Convex Hull

```
11 dist(pair<int, int> a, pair<int, int> b) {
    int dx1 = b.first - a.first;
    int dy1 = b.second - a.second;
    return 1LL * dx1 * dx1 + 1LL * dy1 * dy1;
}

11 ccw(pair<int, int> a, pair<int, int> b, pair<int, int> c) {
    int dx1 = b.first - a.first;
    int dy1 = b.second - a.second;
    int dx2 = c.first - a.first;
    int dy2 = c.second - a.second;
    return 1LL * dx1 * dy2 - 1LL * dy1 * dx2;
}

vector<pair<int, int>> getPoly(vector<pair<int, int> > &poly) {
    int n = poly.size();
```

```
pair<int, int> least = poly[0];
int pos = 0:
for(int i = 0; i < n; i++) {</pre>
    if(poly[i] < least) {</pre>
        least = poly[i];
        pos = i;
    }
}
swap(poly[0], poly[pos]);
sort(poly.begin() + 1, poly.end(), [&](const pair<int, int> &a, const pair<int, int>
     \hookrightarrow &b) {
    auto x = ccw(poly[0], a, b);
    if(x != 0) return x > 0;
    return dist(poly[0], a) < dist(poly[0], b);</pre>
});
vector<pair<int, int> > hull;
for(auto &i : poly) {
    while(hull.size() >= 2 && ccw(hull[hull.size()-2], hull.back(), i) <= 0){</pre>
         hull.pop_back();
    hull.push_back(i);
}
return hull;
```

Listing 4: Pick's Theorem

```
struct Point{
    long long x, y;
   Point(){}
    Point(long long x, long long y) : x(x), y(y){}
};
/// twice the area of polygon
long long double_area(Point poly[], int n){
    long long res = 0;
    for (int i = 0, j = n - 1; i < n; j = i++){
        res += ((poly[j].x + poly[i].x) * (poly[j].y - poly[i].y));
    return abs(res):
}
/// number of lattice points strictly on polygon border
long long on_border(Point poly[], int n){
   long long res = 0;
   for (int i = 0, j = n - 1; i < n; j = i++){
        res += __gcd(abs(poly[i].x - poly[j].x), abs(poly[i].y - poly[j].y));
   }
    return res;
}
/// number of lattice points strictly inside polygon
long long interior(Point poly[], int n){
    long long res = 2 + double_area(poly, n) - on_border(poly, n);
    return res / 2;
```

Listing 5: Closest Pair of Points

```
//returns pair of indexes into sorted array of points
pair<int, int> closestPair(vector<pair<11, 11>> &pts) {
   int n = pts.size();
   sort(pts.begin(), pts.end());
   set<pair<pli>ll, ll>, int>> s;
   11 best_dist = 1e18;
   int j = 0;
   pair<int, int> pos;
   for (int i = 0; i < n; ++i) {
        11 d = ceil(sqrt(best dist)):
        while (pts[i].first - pts[j].first > best_dist) {
            s.erase({{pts[j].second, pts[j].first}, j});
           j += 1;
       }
        auto it1 = s.lower_bound({{pts[i].second - d, pts[i].first}, -1});
        auto it2 = s.upper_bound({{pts[i].second + d, pts[i].first}, i});
        for (auto it = it1; it != it2; ++it) {
            11 dx = pts[i].first - it->first.second;
            11 dy = pts[i].second - it->first.first;
            ll dist = dx * dx + dy * dy;
            if(best_dist > dist) {
               best_dist = dist;
               pos.first = it->second;
               pos.second = i;
            }
        s.insert({{pts[i].second, pts[i].first}, i});
   }
   return pos;
```

Listing 6: **GRAPHS**

Listing 7: Bipartite Matching

```
// Implementation of Hopcroft-Karp algorithm
struct BipartiteMatcher {
    vector<vector<int>> G;
    vector<int> L, R, Viz;
    vector<bool> visitedA, visitedB;
   BipartiteMatcher(int n, int m) :
        G(n), L(n, -1), R(m, -1), Viz(n), visitedA(n), visitedB(m) {}
    void AddEdge(int a, int b) {
        G[a].push_back(b);
   }
   bool Match(int node) {
        if(Viz[node])
            return false;
        Viz[node] = true;
        for(auto vec : G[node]) {
            if(R[vec] == -1 || Match(R[vec])) {
                L[node] = vec;
                R[vec] = node:
                return true;
        }
```

```
return false:
   }
    int Solve() {
        bool ok = true:
        while(ok) {
            ok = false;
            fill(Viz.begin(), Viz.end(), 0);
            for(int i = 0; i < L.size(); ++i)</pre>
                if(L[i] == -1)
                     ok |= Match(i);
        }
        int ret = 0;
        for(int i = 0; i < L.size(); ++i)</pre>
            ret += (L[i] != -1);
        return ret;
   }
    void dfs(int node) {
        visitedA[node] = true;
        for(auto to : G[node]) {
            if(!visitedB[to] && R[to] != -1) {
                visitedB[to] = true;
                dfs(R[to]);
        }
    //returns all the nodes in the min vertex cover
    //O(time\ to\ find\ max\ matching)\ +\ O(n)
    //first vector: all nodes in the cover on the left side
    //second vector: all nodes in the cover on the right side
    pair<vector<int>, vector<int>> getMinVertexCover() {
        Solve();
        for(int i = 0; i < L.size(); ++i) {</pre>
            if(L[i] == -1) {
                 dfs(i):
        }
        pair<vector<int>, vector<int>> cover;
        for(int i = 0; i < L.size(); ++i) {</pre>
            if(!visitedA[i]) cover.first.push_back(i);
        for(int i = 0; i < R.size(); ++i) {</pre>
            if(visitedB[i]) cover.second.push_back(i);
        }
        return cover;
    }
};
```

Listing 8: Bridges

```
//Doesn't handle multiple edges
const int Max = 3e5+2;
vector<int> adj[Max], bridgeTree[Max];
bool visited[Max]:
int timeIn[Max], currTime = 0, minTime[Max];
vector<pair<int, int> > bridges;
vector<int> p(Max,-1);
int find(int x) {return p[x] < 0 ? x : p[x] = find(p[x]);}
void merge(int x, int y) {
```

```
if((x=find(x)) == (y=find(y))) return;
    if(p[y] < p[x]) swap(x,y);
   p[x] += p[y];
   p[y] = x;
void dfs(int node, int prev) {
    visited[node] = true;
    timeIn[node] = minTime[node] = ++currTime;
    for(int to : adi[node]) {
        if(to != prev) {
            minTime[node] = min(minTime[node], timeIn[to]);
        if(visited[to]) continue;
        dfs(to, node):
        minTime[node] = min(minTime[node], minTime[to]);
        if(minTime[to] > timeIn[node]) {
            bridges.push_back({node, to});
        } else {
            merge(node, to);
int main() {
    ios::sync_with_stdio(false);
    cin.tie(0);
    cout.tie(0);
    int n, m;
    cin >> n >> m;
   for(int i = 1; i <= n; ++i) {
        minTime[i] = timeIn[i] = Max;
        visited[i] = false;
   }
    int u. v:
   for(int i = 0; i < m; ++i) {</pre>
        cin >> u >> v;
        adj[u].push_back(v);
        adj[v].push_back(u);
    currTime = 0:
   dfs(1,1);
    for(auto &p : bridges) {
        u = find(p.first);
        v = find(p.second);
        bridgeTree[u].push_back(v);
        bridgeTree[v].push_back(u);
   }
    return 0;
```

Listing 9: Centroid Decomposition

```
const int Max = 2e5+2:
vector<int> adj[Max];
int sizes[Max], parent[Max];
bool removed[Max]:
void dfs2(int node, int par) {
    sizes[node] = 1;
```

```
for(int to : adj[node]) {
        if(to != par && !removed[to]) {
            dfs2(to, node);
            sizes[node] += sizes[to];
   }
int findCentroid(int node) {
   dfs2(node, node):
   bool found = true;
   int sizeCap = sizes[node]/2;
    int par = node;
   while(found) {
       found = false:
       for(int to : adj[node]) {
            if(to != par && !removed[to] && sizes[to] > sizeCap) {
                found = true:
                par = node;
                node = to;
                break;
   }
    return node;
void dfs1(int node, int par) {
   removed[node] = true;
   parent[node] = par;
   for(int to : adj[node]) {
       if(!removed[to]) {
            dfs1(findCentroid(to), node);
   }
//dfs1(findCentroid(1), 0);
```

Listing 10: Count Paths of Each Length

```
const int Max = 1e6+10:
int n, sizes[Max];
vector<int> adj[Max], cntPathLength[Max];
11 prod[Max], cntTotalPathLengths[Max] = {0}, cntTotalPathLengthsNaive[Max] = {0};
bool removed[Max];
void dfs2(int node, int par, int root, int currDist) {
   if(cntPathLength[root].size() <= currDist) {</pre>
       cntPathLength[root].push_back(0);
   }
   cntPathLength[root][currDist]++;
   sizes[node] = 1;
   for(int to : adj[node]) {
       if(to != par && !removed[to]) {
            dfs2(to, node, root, currDist+1);
            sizes[node] += sizes[to]:
   }
```

```
int findCentroid(int node) {
    dfs2(node, node, node, 1);
    bool found = true;
    int sizeCap = sizes[node]/2;
    int par = node;
    while(found) {
        found = false:
        for(int to : adj[node]) {
            if(to != par && !removed[to] && sizes[to] > sizeCap) {
                found = true:
                par = node;
                node = to;
                break;
       }
   }
   return node;
void dfs1(int node, int par) {
    removed[node] = true;
    int maxLength = 1;
   for(int to : adj[node]) {
        if(to != par && !removed[to]) {
            cntPathLength[to].clear();
            cntPathLength[to].push_back(0);
            dfs2(to, to, to, 1);
            maxLength = max(maxLength, (int)cntPathLength[to].size());
    vector<int> temp(maxLength, 0);
    temp[0]++;
    for(int to : adj[node]) {
        if(to != par && !removed[to]) {
            vector<int> prod;
            multiply(temp, cntPathLength[to], prod);
            for(int i = 0; i < prod.size(); ++i) {</pre>
                cntTotalPathLengths[i] += prod[i];
            for(int i = 0; i < cntPathLength[to].size(); ++i) {</pre>
                temp[i] += cntPathLength[to][i];
       }
   }
   for(int to : adj[node]) {
        if(to != par && !removed[to]) {
            dfs1(findCentroid(to), node);
        }
   }
```

Listing 11: Cut Vertices

```
const int Max = 1e5+2;
vector<int> adj[Max];
bool visited[Max], cutNode[Max];
int timeIn[Max], currTime = 0, minTime[Max];
void dfs(int node, int prev) {
```

```
visited[node] = true:
    timeIn[node] = minTime[node] = ++currTime;
    int numChildren = 0;
    for(int to : adj[node]) {
       if(to != prev) {
            minTime[node] = min(minTime[node], timeIn[to]);
       if(visited[to]) continue;
       numChildren++:
       if(node == prev && numChildren > 1) {
            cutNode[node] = true;
       }
       dfs(to, node);
       minTime[node] = min(minTime[node], minTime[to]);
       if(node != prev && minTime[to] >= timeIn[node]) {
            cutNode[node] = true;
   }
int main() {
   ios::sync_with_stdio(false);
   cin.tie(0);
   cout.tie(0);
   int n, m;
   cin >> n >> m;
   for(int i = 0; i < Max; ++i) {</pre>
       minTime[i] = timeIn[i] = Max;
       visited[i] = cutNode[i] = false;
   }
   int u, v;
   for(int i = 0; i < m; ++i) {</pre>
       cin >> u >> v;
       adj[u].push_back(v);
       adj[v].push_back(u);
   }
   currTime = 0;
   for(int i = 0; i < n; ++i) {
       if(!visited[i]) {
            dfs(i,i);
   }
   for(int i = 0; i < n; ++i) {</pre>
       if(cutNode[i]) {
   }
   return 0;
```

Listing 12: Floyd Warshall

```
for (int k = 0; k < n; k++){
    for (int i = 0; i < n; i++){
        for (int j = 0; j < n; j++){
            if (w[i][j] > w[i][k] + w[k][j]){
                 w[i][j] = w[i][k] + w[k][j];
                 prev[i][j] = k;
        }
    }
}
```

```
}
}
```

Listing 13: DSU on Tree

```
const int Max = 1e5+3:
int color[Max], Time = 1, timeIn[Max], timeOut[Max], ver[Max], size[Max], cnt[Max],

    heavyChild[Max], Depth[Max] = {0}, answer[Max];
vector<int> adi[Max]:
void dfs(int node, int prev) {
    timeIn[node] = Time;
    ver[Time] = node;
    Time++;
    size[node] = 1;
    int largest = heavyChild[node] = -1;
    Depth[node] = 1 + Depth[prev];
    for(int to : adj[node]) {
        if(to == prev) continue;
        dfs(to, node);
        size[node] += size[to];
        if(size[to] > largest) {
            largest = size[to];
            heavyChild[node] = to;
        }
    timeOut[node] = Time;
void dfs1(int node, int prev, bool keep = true) {
    for(int to : adj[node]) {
        if(to == prev || to == heavyChild[node]) continue;
        dfs1(to, node, false);
    if(heavyChild[node] != -1) {
        dfs1(heavyChild[node], node, true);
    cnt[color[node]]++;
    for(int to : adj[node]) {
        if(to == prev || to == heavyChild[node]) continue;
        for(int i = timeIn[to]; i < timeOut[to]; ++i) {</pre>
            cnt[color[ver[i]]]++;
        }
    }
    if(!keep) {
        for(int i = timeIn[node]; i < timeOut[node]; ++i) {</pre>
            cnt[color[ver[i]]]--;
int main() {
    ios::sync_with_stdio(false);
    cin.tie(0):
```

cout.tie(0);

int n:

cin >> n:

dfs(1, 1);
dfs1(1, 1):

for(int i = 1; i <= n; ++i) {

```
cout << answer[i] << ' ';
}
cout << '\n';
return 0;
}</pre>
```

Listing 14: Lowest Common Ancestor (LCA)

```
const int Max = 1e5+3, Log = 20+1;
vector<int> adj[Max];
int memo[Max][Log];
int depth[Max];
void dfs(int node, int prev, int Depth = 0) {
    depth[node] = Depth;
    memo[node][0] = prev;
    for(int i = 1; i < Log; ++i) {</pre>
        memo[node][i] = memo[memo[node][i-1]][i-1];
   }
   for(int to : adj[node]) {
        if(to == prev) continue;
        dfs(to, node, Depth+1);
   }
}
int LCA(int x, int y) {
    if(depth[x] < depth[y]) swap(x,y);</pre>
    int diff = depth[x] - depth[y];
    for(int k = Log-1; k >= 0; --k) {
        if(diff&(1<<k)) {</pre>
            x = memo[x][k];
   }
    for(int k = Log-1; k \ge 0; --k) {
        if(memo[x][k] != memo[y][k]) {
            x = memo[x][k];
            y = memo[y][k];
        }
   }
    if(x != y) x = memo[x][0];
    return x;
//dfs(1,1);
```

Listing 15: Strongly Connected Components (SCC)

```
int n,m;
vector<vector<int> > adj, adjInv;
vector<int> scc;
int sccID;
vector<bool> visited;

void dfs1(int curr, stack<int> &seen) {
    visited[curr] = true;
    for(int x : adj[curr]) {
        if(!visited[x]) {
            dfs1(x, seen);
        }
    }
}
```

```
seen.push(curr);
void dfs2(int curr) {
    visited[curr] = true;
    scc[curr] = sccID;
    for(int x : adjInv[curr]) {
        if(!visited[x]) {
            dfs2(x);
    }
}
void calcSCC() {
    visited.resize(n+1,false);
    stack<int> seen;
    for(int i = 1; i <= n; ++i) {
        if(!visited[i]) {
            dfs1(i, seen);
        }
    visited.clear();
    visited.resize(n+1,false);
    sccID = 0;
    while(!seen.empty()) {
        while(!seen.empty() && visited[seen.top()]) seen.pop();
        if(!seen.empty()) {
            dfs2(seen.top());
            sccID++;
        }
    }
}
```

Listing 16: MAX FLOW

Listing 17: Dinic's Max Flow

```
struct maxflow {
    struct edge {
        11 a, b, cap, flow;
   11 n, s, t;
    vector<ll> d, ptr, q;
    vector<edge> e;
    vector<vector<ll>>> g;
    maxflow(int _n, int _s, int _t) : n(_n), s(_s), t(_t) {
        d.resize(n);
        ptr.resize(n);
        q.resize(n);
        g.resize(n);
   }
    void addedge(ll a, ll b, ll cap) {
        edge e1 = { a, b, cap, 0 };
        edge e2 = \{ b, a, 0, 0 \};
        g[a].push_back((11) e.size());
```

```
e.push_back(e1);
        g[b].push_back((11) e.size());
        e.push_back(e2);
    }
    bool bfs() {
        11 ah=0, at=0:
        q[qt++] = s;
        d.assign(d.size(), -1);
        d[s] = 0:
        while(qh < qt && d[t] == -1) {
            ll v = q[qh++];
            for(size_t i=0; i<g[v].size(); ++i) {</pre>
                11 id = g[v][i],
                    to = e[id].b;
                if(d[to] == -1 && e[id].flow < e[id].cap) {</pre>
                     q[qt++] = to;
                     d[to] = d[v] + 1;
            }
        }
        return d[t] != -1;
    }
    11 dfs(ll v, ll flow) {
        if(!flow) return 0;
        if(v == t) return flow;
        for(; ptr[v]<(l1)g[v].size(); ++ptr[v]) {</pre>
            11 id = g[v][ptr[v]];
            11 to = e[id].b;
            if(d[to] != d[v] + 1) continue;
            ll pushed = dfs(to, min (flow, e[id].cap - e[id].flow));
            if(pushed) {
                e[id].flow += pushed;
                e[id^1].flow -= pushed;
                return pushed;
            }
        }
        return 0;
    }
    11 getflow() {
        11 \text{ flow} = 0;
        for(;;) {
            if(!bfs()) break;
            ptr.assign(ptr.size(), 0);
            while(ll pushed = dfs(s,inf)) {
                flow += pushed;
        return flow;
    }
};
```

Listing 18: Hungarian Algorithm

```
// this is one-indexed
// jobs X workers cost matrix
// cost[i][j] is cost of job i done by worker j
// #jobs must be <= #workers</pre>
```

```
// Default finds min cost; to find max cost set all costs[i][j] to -costs[i][j]
11 HungarianMatch(const vector<vector<11>>& a) {
    ll n = a.size()-1;
    ll m = a[0].size()-1;
    vector<ll> u(n+1), v(m+1), p(m+1), way(m+1);
    for(ll i = 1; i <= n; ++i) {
        p[0] = i;
        11 i0 = 0;
        vector<ll> minv(m+1, inf);
        vector<char> used(m+1, false);
        do {
            used[j0] = true;
            11 i0 = p[j0], delta = inf, j1;
            for(ll j = 1; j <= m; ++j)</pre>
                if(!used[j]) {
                    11 cur = a[i0][j] - u[i0] - v[j];
                     if(cur < minv[j])</pre>
                         minv[j] = cur, way[j] = j0;
                     if(minv[j] < delta)</pre>
                        delta = minv[j], j1 = j;
            for(11 j = 0; j <= m; ++j)
                if(used[j])
                     u[p[j]] += delta, v[j] -= delta;
                else
                    minv[j] -= delta;
            j0 = j1;
        } while(p[j0] != 0);
        do {
            11 j1 = way[j0];
            p[j0] = p[j1];
            j0 = j1;
        } while(j0);
    // For each N, it contains the M it selected
    vector < ll > ans(n+1):
    for(ll \ j = 1; \ j <= m \ ; \ ++j)
        ans[p[j]] = j;
    return -v[0];
```

Listing 19: Min Cost Max Flow

```
struct mincostmaxflow {
    struct edge {
        ll a, b, cap, cost, flow;
        size_t back;
    };

    vector<edge> e;
    vector<vector<1l>> g;
    ll n, s, t;
    ll k = inf; // The maximum amount of flow allowed

mincostmaxflow(int _n, int _s, int _t) : n(_n), s(_s), t(_t) {
        g.resize(n);
}
```

```
}
void addedge(ll a, ll b, ll cap, ll cost) {
    edge e1 = {a,b,cap,cost,0,g[b].size()};
    edge e2 = {b,a,0,-cost,0,g[a].size()};
    g[a].push_back((ll) e.size());
    e.push back(e1);
    g[b].push_back((11) e.size());
    e.push_back(e2);
}
// Returns {flow, cost}
pair<ll,ll> getflow() {
    11 flow = 0, cost = 0;
    while(flow < k) {</pre>
        vector<ll> id(n, 0);
        vector<ll> d(n, inf);
        vector<11> q(n);
        vector<ll> p(n);
        vector<size_t> p_edge(n);
        11 qh=0, qt=0;
        q[qt++] = s;
        d[s] = 0;
        while(qh != qt) {
            ll v = q[qh++];
            id[v] = 2;
            if(qh == n) qh = 0;
            for(size_t i=0; i<g[v].size(); ++i) {</pre>
                edge& r = e[g[v][i]];
                if(r.flow < r.cap && d[v] + r.cost < d[r.b]) {
                    d[r.b] = d[v] + r.cost;
                    if(id[r.b] == 0) {
                        q[qt++] = r.b;
                        if(qt == n) qt = 0;
                    else if(id[r.b] == 2) {
                        if(--qh == -1) qh = n-1;
                        q[qh] = r.b;
                    }
                    id[r.b] = 1;
                    p[r.b] = v:
                    p_edge[r.b] = i;
            }
        if(d[t] == inf) break;
        11 addflow = k - flow;
        for(ll v=t; v!=s; v=p[v]) {
            11 pv = p[v]; size_t pr = p_edge[v];
            addflow = min(addflow, e[g[pv][pr]].cap - e[g[pv][pr]].flow);
        for(11 v=t; v!=s; v=p[v]) {
            ll pv = p[v]; size_t pr = p_edge[v], r = e[g[pv][pr]].back;
            e[g[pv][pr]].flow += addflow;
            e[g[v][r]].flow -= addflow;
            cost += e[g[pv][pr]].cost * addflow;
        flow += addflow;
    }
    return {flow,cost};
}
```

|};

Listing 20: STRINGS

Listing 21: KMP String Matching

```
struct KMP Match {
    vector<int> T;
    string pat;
    KMP_Match() {};
    KMP_Match(string pattern) : pat(pattern) {this->buildTable(pat);};
    void buildTable(string pattern) {
        pat = pattern;
        T.clear();
        T.resize(pat.length()+1);
        int i = 0, j = -1;
        T[i] = j;
        while(i < pat.size()) {</pre>
            while(j >= 0 && pat[i] != pat[j]) j = T[j];
            i++, j++;
            T[i] = j;
        }
    vector<int> find(string txt, bool all = true) {
        int m = 0, i = 0;
        vector<int> matches;
        while(m + i < txt.length()) {</pre>
            if(pat[i] == txt[m+i]) {
                if(i == pat.length()-1) {
                     matches.push_back(m);
                     if(!all) return matches;
                     m = m + i - T[i];
                     i = T[i]:
                }
                i++;
            } else {
                if(T[i] != -1) {
                     m = m + i - T[i];
                    i = T[i]:
                } else {
                     i = 0;
                     m++;
        return matches;
};
int fail[1000005];
// Checks if two arrays are rotationally equvalent
bool KMPints(vector<ll> a, vector<ll> b){
    for(int i = 0; i < a.size(); i++) {</pre>
        b.push_back(b[i]);
    }
    int p = 0;
    for(int i = 1; i < a.size(); i++) {</pre>
```

```
while(p && a[i] != a[p]) p = fail[p];
    if(a[i] == a[p]) p++;
    fail[i + 1] = p;
}
p = 0;
for(auto &i : b) {
    while(p && i != a[p]) p = fail[p];
    if(i == a[p]) p++;
    if(p == a.size()) return true;
}
return false;
```

Listing 22: Manachers (Palindromes)

```
// Returns the longest palindrome
string cntPals(string &s) {
    string T = "";
    for(int i = 0; i < s.size(); ++i) {</pre>
        T += "#" + s.substr(i,1);
   T += '#':
    vector<int> P(T.size(),0);
    int center = 0, boundary = 0, maxLen = 0, resCenter = 0, cnt = 0;
   for(int i = 1; i < T.size()-1; ++i) {</pre>
        int iMirror = 2 * center - i;
        P[i] = (boundary > i ? min(boundary - i, P[iMirror]) : 0);
        while(i-1-P[i] >= 0 && i+1+P[i] <= (int)T.size()-1 && T[i+1+P[i]] == T[i-1-P[i]])
            P[i]++;
        if(i + P[i] > boundary) {
            center = i;
            boundary = i+P[i];
        }
        if(P[i] > maxLen) {
            maxLen = P[i]:
            resCenter = i;
        cnt += (P[i]+1)/2;
   }
    return s.substr((resCenter - maxLen)/2, maxLen);
    //return cnt;//number of palindromes
```

Listing 23: Suffix Array

```
const int MAXN = 1 << 21;
string s;
int N, gap;
int sa[MAXN], pos[MAXN], lcp[MAXN], tmp[MAXN];

bool sufCmp(int i, int j) {
    if(pos[i] != pos[j]) return pos[i] < pos[j];
    i += gap;
    j += gap;
    return (i < N && j < N) ? pos[i] < pos[j] : i > j;
}

void buildSA() {
    N = s.length();
```

```
for(int i = 0; i < N; ++i) {</pre>
        sa[i] = i:
        pos[i] = s[i];
    for(gap = 1;; gap *= 2) {
        sort(sa, sa + N, sufCmp);
        for(int i = 0: i < N-1: ++i)
            tmp[i+1] = tmp[i] + sufCmp(sa[i], sa[i+1]);
        for(int i = 0; i < N; ++i) pos[sa[i]] = tmp[i];</pre>
        if(tmp[N-1] == N-1) break;
   }
void buildLCP() {
    N = s.size();
    for(int i = 0, k = 0; i < N; ++i) {
        if(pos[i] != 0 ) {
            for(int j = sa[pos[i]-1]; s[i+k] == s[j+k];) k++;
            lcp[pos[i]] = k;
            if(k) k--;
        }
    }
}
```

Listing 24: Trie

```
const int K = 26;//character size
struct node {
    int next[K];
    bool leaf = 0;
    char pch;
    int p = -1;
    int id;
    node(int p = -1, char ch = '#'):p(p),pch(ch) {
        fill(next,next+K,-1);
};
vector<node> t(1);//adj list
void add_string(string s, int id) {
    int c = 0;
    for(char ch: s){
        int v = ch-'a';
        if(t[c].next[v] == -1) {
            t[c].next[v] = t.size();
            t.emplace_back(c,ch);
        c = t[c].next[v];
    t[c].leaf = 1;
    t[c].id = id;
void remove_string(string s) {
    int c = 0:
    for(char ch: s){
        int v = ch-'a';
        if(t[c].next[v] == -1) {
```

```
return;
}
c = t[c].next[v];
}
t[c].leaf = 0;

int find_string(string s) {
    int c = 0;
    for(char ch: s){
        int v = ch-'a';
        if(t[c].next[v] == -1) {
            return -1;
        }
        c = t[c].next[v];
}
if(!t[c].leaf) return -1;
return t[c].id;
}
```

Listing 25: Z Algorithm

```
string s;
//z[i] is the length of the longest substring
//starting from s[i] which is also a prefix of s
int z[100010];//change size here
void zAlg() {
   int n = s.size();
   int L = 0, R = 0;
   for (int i = 1; i < n; i++) {
       if (i > R) {
           L = R = i:
            while (R < n \&\& s[R-L] == s[R]) R++;
            z[i] = R-L; R--;
       } else {
            int k = i-L;
            if (z[k] < R-i+1) z[i] = z[k];
            else {
                while (R < n \&\& s[R-L] == s[R]) R++;
                z[i] = R-L; R--;
       }
   }
```

Listing 26: DATA STRUCTURES

Listing 27: Longest Increasing Subsequence (LIS)

```
// Returns an array with the indexes of the LIS
template <class T>
vector<int> LIS(vector<T>& v) {
   if(v.size() == 0) return {};
   vector<int> p(v.size(), -1);
   vector<int> t(v.size(), 0);
```

```
int lis = 1:
    for(int i = 1; i < v.size(); i++) {</pre>
        if(v[i] <= v[t[0]]) {</pre>
            t[0] = i:
       }
        else if(v[i] > v[t[lis - 1]]) {
            p[i] = t[lis - 1];
            t[lis++] = i;
       }
        else {
            int 1 = -1;
            int r = lis - 1;
            while (r - 1 > 1) {
                int m = 1 + (r - 1) / 2;
                if(v[t[m]] >= v[i]) r = m;
                else 1 = m;
            p[i] = t[r - 1];
            t[r] = i;
       }
   }
    vector<int> ans;
    for(int i = t[lis - 1]; i >= 0; i = p[i]) {
        ans.push_back(i);
    reverse(ans.begin(), ans.end());
   return ans;
vector<int> LIS(vector<int> &arr) {//longest non-decreasing sequence
    vector<int> longest(arr.size(),0);
   multiset<int> seq;
    for(int i = 0; i < arr.size(); ++i) {</pre>
        seq.insert(arr[i]);
        auto it = seq.upper_bound(arr[i]);
        if(it != seq.end()) seq.erase(it);
        longest[i] = seq.size();
   return longest;
vector<int> LSIS(vector<int> &arr) {//longest strictly increasing sequence
    vector<int> longest(arr.size(),0);
    multiset<int> seq;
    for(int i = 0; i < arr.size(); ++i) {</pre>
        seq.insert(arr[i]);
        auto it = seq.lower_bound(arr[i]);
        if(it != seq.end()) seq.erase(it);
        longest[i] = seq.size();
   }
   return longest;
```

Listing 28: Count Rectangles

```
//given a 2D boolean matrix, calculate cnt[i][j]
//cnt[i][j] = the number of times an (i * j) rectangle appears in the matrix
```

```
//such that all cells in the rectangle are false
vector<vector<int>> getNumRectangles(vector<vector<bool>> &grid) {
    vector<vector<int>> cnt:
    const int rows = grid.size();
    if(rows == 0) return cnt;
    const int cols = grid[0].size();
    if(cols == 0) return cnt;
    cnt.resize(rows+1, vector<int>(cols+1, 0));
    vector<vector<int>> arr(rows+2, vector<int>(cols+1, 0));
    for(int i = 1; i <= rows; ++i) {
        for(int j = 1; j <= cols; ++j) {</pre>
            arr[i][j] = 1 + arr[i][j-1];
            if(grid[i-1][j-1]) arr[i][j] = 0;
        }
    }
    for(int j = 1; j <= cols; ++j) {</pre>
        arr[rows+1][j] = 0;
        stack<pair<int, int> > st;
        st.push({0,0});
        for(int i = 1; i <= rows+1; ++i) {</pre>
             pair<int, int> curr = {i, arr[i][j]};
            while(arr[i][j] < st.top().second) {</pre>
                 curr = st.top();
                 st.pop();
                 cnt[i-curr.first][curr.second]++;
                 cnt[i-curr.first][max(arr[i][j], st.top().second)]--;
            st.push({curr.first, arr[i][j]});
        }
    }
    for(int j = 1; j <= cols; ++j) {</pre>
        for(int i = rows-1; i >= 1; --i) {
             cnt[i][j] += cnt[i+1][j];
        for(int i = rows-1; i >= 1; --i) {
             cnt[i][j] += cnt[i+1][j];
    }
    for(int i = 1; i <= rows; ++i) {</pre>
        for(int j = cols-1; j >= 1; --j) {
            cnt[i][j] += cnt[i][j+1];
        }
    }
    return cnt;
```

Listing 29: Disjoint Set

```
vector<int> p(1000001,-1);//change size here if needed
int find(int x) {return p[x] < 0 ? x : p[x] = find(p[x]);}
void merge(int x, int y) {
    if((x=find(x)) == (y=find(y))) return;
    if(p[y] < p[x]) swap(x,y);
    p[x] += p[y];
    p[y] = x;
}</pre>
```

Listing 30: Indexed Set

Listing 31: Largest Range Where Element Is Max/Min

```
#include <bits/stdc++.h>
using namespace std;
int main() {
   int n:
   cin >> n:
    vector<int> arr(n), lefts(n), rights(n);
    for(int &x : arr) cin >> x;
    stack<int> positions;
    for(int i = 0; i < n; ++i) {
        while(!positions.empty() && arr[positions.top()] >= arr[i]) positions.pop();
        if(positions.empty()) {
            lefts[i] = -1;
       } else {
            lefts[i] = positions.top();
        positions.push(i);
    while(!positions.empty()) positions.pop();
    for(int i = n-1; i >= 0; --i) {
        while(!positions.empty() && arr[positions.top()] >= arr[i]) positions.pop();
        if(positions.empty()) {
            rights[i] = n;
        } else {
            rights[i] = positions.top();
        positions.push(i);
   }
    return 0:
```

Listing 32: Online Convex Hull

```
const ll is_query = -(1LL<<62);
struct Line {
    ll m, b;
    mutable function<const Line*()> succ;
    bool operator<(const Line& rhs) const {
        if (rhs.b!= is_query) return m < rhs.m;
        const Line* s = succ();
        if (!s) return 0;
        ll x = rhs.m;
        return b - s->b < (s->m - m) * x;
    }
};
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum bool bad(iterator y) {
        auto z = next(y);
    }
}
```

```
if (y == begin()) {
            if (z == end()) return 0;
            return y->m == z->m && y->b <= z->b;
        auto x = prev(y);
        if (z == end()) return y->m == x->m && y->b <= x->b;
        return (x-b-y-b) * (long double)(z-m-y-m) >= (y-b-z-b) * (long double)
             \hookrightarrow double)(v->m - x->m);
   }
    void insert line(ll m. ll b) {
        auto y = insert({ m, b });
        y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
        if (bad(y)) { erase(y); return; }
        while (next(y) != end() && bad(next(y))) erase(next(y));
        while (y != begin() && bad(prev(y))) erase(prev(y));
    }
    11 getVal(11 x) {
        auto 1 = *lower_bound((Line) { x, is_query });
        return 1.m * x + 1.b;
    }
};
```

Listing 33: Queue Max

```
struct queueMax {
   stack<pair<ll, ll>> s1, s2;
   int size() {
       return s1.size() + s2.size();
   }
   bool isEmpty() {
       return (size() == 0);
   }
   void clear() {
       while(!s1.empty()) {
            s1.pop();
       }
       while(!s2.empty()) {
            s2.pop();
       }
   }
   11 getMax() {
       if(isEmpty()) {
           return -1e18;
       if(!s1.empty() && !s2.empty()) {
            return max(s1.top().second, s2.top().second);
       if(!s1.empty()) {
           return s1.top().second;
       }
       return s2.top().second;
   }
   void push(ll val) {
       if(s2.empty()) {
            s2.push({val, val});
       } else {
            s2.push({val, max(val, s2.top().second)});
   }
   void pop() {
```

```
if(s1.empty()) {
    while(!s2.empty()) {
        if(s1.empty()) {
            s1.push({s2.top().first, s2.top().first});
        } else {
            s1.push({s2.top().first, max(s2.top().first, s1.top().second)});
        }
        s2.pop();
    }
}
s1.pop();
}
```

Listing 34: Two Sat

```
struct twosat {
    vector<vector<int> > adj, adjInv;
    vector<int> scc;
    int sccID;
    vector<bool> visited, assignment;
    twosat(int nodes) {
        n = 2*nodes:
        adj.resize(n);
        adjInv.resize(n);
        scc.resize(n);
        assignment.resize(n/2);
    }
    //X AND Y = (X OR X) AND (Y OR Y)
    //X NAND Y = (!X OR !Y)
    //X NOR Y = (!X OR !X) AND (!Y OR !Y)
    //X XOR Y = (X OR Y) AND (!X OR !Y)
    //X XNOR Y = (!Y OR X) AND (!X OR Y)
    void add(int i, bool statusI, int j, bool statusJ) {
        const int from1 = i+(!statusI)*(n/2);
        const int to1 = j+statusJ*(n/2);
        adj[from1].push_back(to1);
        adjInv[to1].push_back(from1);
        const int from 2 = j + (!statusJ) * (n/2);
        const int to2 = i+statusI*(n/2);
        adj[from2].push_back(to2);
        adjInv[to2].push_back(from2);
    }
    void dfs1(int curr, stack<int> &seen) {
        visited[curr] = true;
        for(int x : adj[curr]) {
            if(!visited[x]) {
                dfs1(x, seen);
        seen.push(curr);
    void dfs2(int curr) {
        visited[curr] = true;
        scc[curr] = sccID;
```

```
for(int x : adjInv[curr]) {
            if(!visited[x]) {
                dfs2(x);
        }
    }
    bool solve() {
        visited.resize(n+1,false);
        stack<int> seen:
        for(int i = 0; i < n; ++i) {
            if(!visited[i]) {
                dfs1(i, seen);
        }
        visited.clear();
        visited.resize(n+1,false);
        sccID = 0;
        while(!seen.empty()) {
            while(!seen.empty() && visited[seen.top()]) seen.pop();
            if(!seen.empty()) {
                dfs2(seen.top());
                sccID++:
        }
        for(int i = 0; i < n/2; ++i) {
            if(scc[i] == scc[i+n/2]) {
                return false;
            assignment[i] = scc[i] < scc[i+n/2];</pre>
        }
        return true;
    }
};
```

```
left++;
    while(right > q.r) {
        remove(right);
        right--;
    answer[q.index] = answerToQuery;
for(int i = 0; i < q; ++i) cout << answer[i] << '\n';</pre>
```

int main() {

int a:

cin >> q;

add(0);

vector<query> queries(q); for(int i = 0; i < q; ++i) {

answer[i] = 0;

answerToQuery = 0; for(auto &q : queries) {

queries[i].index = i;

while(left > q.1) {

add(left);

while(right < q.r) {</pre>

add(right);

while(left < q.1) {</pre> remove(left);

left--:

right++;

cin >> queries[i].1 >> queries[i].r;

sort(queries.begin(), queries.end(), cmp);

int left = 0, right = 0; //store inclusive ranges, start at [0,0]

Listing 35: RANGE DATA STRUCTURES

Listing 36: Mo's Algorithm

```
#include <bits/stdc++.h>
using namespace std;
const int Max = 1e6+2;
int block, answer[Max], answerToQuery;
struct query {
    int 1, r, index;
bool cmp(query x, query y) {
    if(x.1/block == y.1/block) return x.r < y.r;</pre>
    return x.1 < y.1;</pre>
void add(int pos) {
void remove(int pos) {
```

Listing 37: Fenwick Tree

```
struct fenwickTree {
    vector<ll> bit;
    int n;
    fenwickTree() {
        n = (int)1e5+3;
        bit.assign(n,0);
    11 sum(int r) {
        11 ret = 0;
        for(; r >= 0; r = (r&(r+1))-1)
            ret += bit[r];
        return ret;
    void add(int idx, ll d) {
        for(; idx < n; idx = idx | (idx+1))</pre>
            bit[idx] += d;
    11 sum(int 1, int r) {
        return sum(r) - sum(l-1):
}ft;
```

#define LOGSZ 17

```
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ):
// add v to value at x
void set(int x. int v) {
 while(x \le N) {
   tree[x] += v;
   x += (x \& -x):
 }
// get cumulative sum up to and including x
int get(int x) {
 int res = 0;
 while(x) {
   res += tree[x]:
   x -= (x \& -x);
 return res;
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
 int idx = 0, mask = N;
 while(mask && idx < N) {</pre>
   int t = idx + mask;
   if(x \ge tree[t]) {
     idx = t;
     x -= tree[t];
   mask >>= 1;
 return idx;
```

Listing 38: Fenwick Tree 2D

```
struct fenwickTree2D {
   vector<vector<ll>>> bit;
   int n, m;
   //2D matrix, with n rows, and m columns
   fenwickTree2D(int _n, int _m) {
       n = _n;
       m = _m;
       bit.resize(n+1, vector<ll>(m+1,0));
   //returns sum of rows [0..i], and [0..j] inclusive
   11 sum(int i, int j) {
       11 ret = 0;
       for(; i \ge 0; i = (i&(i+1))-1)
           for(int jj = j; jj >= 0; jj = (jj&(jj+1))-1)
               ret += bit[i][jj];
       return ret;
   }
   //adds value to location i, j
   void add(int i, int j, ll d) {
       for(; i <= n; i = i | (i+1))
```

```
for(int jj = j; jj <= m; jj = jj | (jj+1))</pre>
                 bit[i][j] += d;
};
```

Listing 39: Search Buckets

```
// Author: Neal Wu
// search buckets provides two operations on an array:
// 1) set array[i] = x
// 2) count how many i in [start, end) satisfy array[i] < value
// Both operations take sqrt(N log N) time. Amazingly, because of the cache efficiency
     \hookrightarrow this is faster than the
// (log N)^2 algorithm until N = 2-5 million.
template<typename T>
struct search_buckets {
    // values are just the values in order. buckets are sorted in segments of

→ BUCKET_SIZE (last segment may be smaller)
    int N, BUCKET_SIZE;
    vector<T> values, buckets;
    search_buckets(const vector<T> &initial = {}) {
         init(initial);
    }
    int get_bucket_end(int bucket_start) const {
         return min(bucket_start + BUCKET_SIZE, N);
    void init(const vector<T> &initial) {
         values = buckets = initial;
        N = values.size();
         BUCKET_SIZE = 3 * sqrt(N * log(N + 1)) + 1;
        for (int start = 0; start < N; start += BUCKET_SIZE)</pre>
             sort(buckets.begin() + start, buckets.begin() + get_bucket_end(start));
    }
    int bucket_less_than(int bucket_start, T value) const {
         auto begin = buckets.begin() + bucket_start;
         auto end = buckets.begin() + get_bucket_end(bucket_start);
         return lower_bound(begin, end, value) - begin;
    int less_than(int start, int end, T value) const {
         int count = 0;
         int bucket_start = start - start % BUCKET_SIZE;
         int bucket_end = min(get_bucket_end(bucket_start), end);
         if (start - bucket_start < bucket_end - start) {</pre>
             while (start > bucket_start)
                 count -= values[--start] < value;</pre>
        } else {
             while (start < bucket_end)</pre>
                 count += values[start++] < value;</pre>
         if (start == end) return count;
         bucket_start = end - end % BUCKET_SIZE;
         bucket end = get bucket end(bucket start);
         if (end - bucket_start < bucket_end - end) {</pre>
            while (end > bucket_start)
                count += values[--end] < value;</pre>
```

```
} else {
            while (end < bucket end)
                count -= values[end++] < value;</pre>
        while (start < end && get_bucket_end(start) <= end) {</pre>
            count += bucket_less_than(start, value);
            start = get_bucket_end(start);
        }
        return count;
    }
    int prefix_less_than(int n, T value) const {
        return less_than(0, n, value);
    void modify(int index, T value) {
        int bucket_start = index - index % BUCKET_SIZE;
        int old_pos = bucket_start + bucket_less_than(bucket_start, values[index]);
        int new_pos = bucket_start + bucket_less_than(bucket_start, value);
        if (old_pos < new_pos) {</pre>
            copy(buckets.begin() + old_pos + 1, buckets.begin() + new_pos,

    buckets.begin() + old_pos);
            new_pos--;
        } else {
            copy_backward(buckets.begin() + new_pos, buckets.begin() + old_pos,

    buckets.begin() + old_pos + 1);
        buckets[new_pos] = value;
        values[index] = value;
   }
};
```

Listing 40: Seg Tree

```
struct SegmentTree {
    vector<11> treeSum, treeMax, lazy;
   ll n, root, size;
   SegmentTree(int currSize) : n(currSize), root(1) {
       11 x = (11)(ceil(log2(currSize)));
        size = 2*(11)pow(2, x);
        treeSum.resize(size, 0):
        treeMax.resize(size, 0);
        lazy.resize(size, 0);
   }
   SegmentTree(const vector<ll> &arr) : n(arr.size()), root(1) {
       11 x = (11)(ceil(log2(n)));
        size = 2*(11)pow(2, x);
        treeSum.resize(size);
       treeMax.resize(size);
       lazy.resize(size, 0);
       build(arr, root, 0, n-1);
   }
   void build(const vector<ll> &arr, int node, int start, int end) {
       if(start == end) treeMax[node] = treeSum[node] = arr[start];
       else {
            11 \text{ mid} = (\text{start+end})/2:
            build(arr, 2*node, start, mid):
            build(arr, 2*node+1, mid+1, end);
            treeSum[node] = treeSum[2*node] + treeSum[2*node+1];
            treeMax[node] = max(treeMax[2*node], treeMax[2*node+1]);
```

```
}
    void pendingUpdate(int node, int start, int end) {
        if(lazy[node]) {
            treeSum[node] += (end-start+1) * lazy[node];
            treeMax[node] += lazy[node];
            if(start != end) {
                lazy[2*node] += lazy[node];
                lazy[2*node+1] += lazy[node];
            lazy[node] = 0;
        }
    void updateRange(int 1, int r, 11 diff) {updateRange(root, 0, n-1, 1, r, diff);}
    void updateRange(int node, int start, int end, int 1, int r, 11 diff) {
        pendingUpdate(node, start, end);
        if(start > end || start > r || end < 1) return;</pre>
        if(start >= 1 && end <= r) {
            treeSum[node] += (end-start+1) * diff;
            treeMax[node] += diff;
            if(start != end) {
                lazv[2*node] += diff:
                lazy[2*node+1] += diff;
            return;
        11 mid = (start + end) / 2;
        updateRange(2*node, start, mid, 1, r, diff);
        updateRange(2*node+1, mid+1, end, 1, r, diff);
        treeSum[node] = treeSum[2*node] + treeSum[2*node+1];
        treeMax[node] = max(treeMax[2*node], treeMax[2*node+1]);
    11 querySum(int 1, int r) {return querySum(root, 0, n-1, 1, r);}
    11 querySum(int node, int start, int end, int 1, int r) {
        if(r < start || end < 1) return 0;</pre>
        pendingUpdate(node, start, end);
        if(1 <= start && end <= r) return treeSum[node];</pre>
        11 \text{ mid} = (\text{start+end})/2:
        return querySum(2*node, start, mid, 1, r) + querySum(2*node+1, mid+1, end, 1, r);
    11 queryMax(int 1, int r) {return queryMax(root, 0, n-1, 1, r);}
    11 queryMax(int node, int start, int end, int 1, int r) {
        if(r < start || end < 1) return -1e18;</pre>
        pendingUpdate(node, start, end);
        if(1 <= start && end <= r) return treeMax[node];</pre>
        11 \text{ mid} = (\text{start+end})/2;
        return max(queryMax(2*node, start, mid, 1, r), queryMax(2*node+1, mid+1, end, 1,
             \hookrightarrow r));
};
```

Listing 41: Sparse Table

```
struct sparseTable {
    vector<vector<ll> > memo;
    vector<int> logTwo;
    int maxPow:
    sparseTable(const vector<11> &arr) {
        int n = arr.size();
        logTwo.resize(n+1,0):
```

```
for(int i = 2; i <= n; ++i) logTwo[i] = 1 + logTwo[i/2];
    maxPow = logTwo[n]+1;
    memo.resize(maxPow, vector<1l>(n));
    for(int j = 0; j < maxPow; ++j) {
        for(int i = 0; i < n; ++i) {
            if(i+(1<<j)-1<n) {
                memo[j][i] = (j?min(memo[j-1][i], memo[j-1][i+(1<<(j-1))]):arr[i]);
        } else break;
      }
    }
}
ll query(int l, int r) {
    int j = logTwo[r-1+1];
    return min(memo[j][l], memo[j][r-(1<<j)+1]);
};</pre>
```

Listing 42: Sparse Table 2D

```
//***zero based indexing
int table[11][1010][11][1010],logTwo[1010];
struct sparseTable2D {
    int n, m;
    sparseTable2D(const vector<vector<int>> &Matrix) {
        n = Matrix.size();
        m = Matrix[0].size();
         logTwo[0] = -1;
        for(int i = 1; i <= max(n,m); ++i) {
             logTwo[i] = 1 + logTwo[i/2];
        for(int ir=0; ir<n; ++ir) {</pre>
             for(int ic=0; ic<m; ++ic)</pre>
                 table[0][ir][0][ic] = Matrix[ir][ic];
             for(int jc=1; jc<=logTwo[m]; ++jc)</pre>
                 for(int ic=0; ic+(1<<(jc-1))<m; ++ic)</pre>
                      table[0][ir][jc][ic] = min(table[0][ir][jc-1][ic],

    table[0][ir][jc-1][ic+(1<<(jc-1))]);
</pre>
         for(int jr=1; jr<=logTwo[n]; ++jr)</pre>
             for(int ir=0; ir<n; ++ir)</pre>
                 for(int jc=0; jc<=logTwo[m]; ++jc)</pre>
                      for(int ic=0; ic<m; ++ic)</pre>
                          table[jr][ir][jc][ic] = min(table[jr-1][ir][jc][ic],

    table[jr-1][ir+(1<<(jr-1))][jc][ic]);
</pre>
    }
    int queryMin(int x1, int y1, int x2, int y2){
        int kx = logTwo[x2-x1+1];
        int ky = logTwo[y2-y1+1];
         int min_R1 = min(table[kx][x1][ky][y1], table[kx][x1][ky][y2+1-(1<<ky)]);</pre>
        int min_R2 = min(table[kx][x2+1-(1<<kx)][ky][y1],
              \hookrightarrow table[kx][x2+1-(1<<kx)][ky][y2+1-(1<<ky)]);
        return min(min_R1, min_R2);
    }
};
```

Listing 43: MATH

Listing 44: Chinese Remainder Theorem

```
inline 11 LCM(11 a, 11 b) {
    return a / \_gcd(a, b) * b;
inline ll normalize(ll x, ll mod) {
    x \% = mod:
    if (x < 0) x += mod;
    return x;
struct gcd_type { 11 x, y, d; };
gcd_type ex_gcd(ll a, ll b) {
    if (b == 0) return {1, 0, a};
    gcd_type pom = ex_gcd(b, a % b);
    return {pom.y, pom.x - a / b * pom.y, pom.d};
// Chinese remainder theorem: find z such that
// z \% mods[i] = vals[i] for all i. Note that the solution is
// unique modulo M = lcm_i \pmod{[i]}. Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
pair<11, 11> crt(const vector<int> &vals, const vector<int> &mods) {
    for(int i = 0; i < vals.size(); i++) {</pre>
        normalize(vals[i], mods[i]);
    ll ans = vals[0]:
    11 Lcm = mods[0];
    for(int i = 1; i < vals.size(); i++) {</pre>
        auto pom = ex_gcd(Lcm, mods[i]);
        int x1 = pom.x;
        int d = pom.d;
        if((vals[i] - ans) % d != 0) {
            return {0,-1};
        ans = normalize(ans + x1 * (vals[i] - ans) / d % (mods[i] / d) * Lcm, Lcm *
             \hookrightarrow mods[i] / d);
        Lcm = LCM(Lcm, mods[i]); // you can save time by replacing above lcm * n[i] / d
             \hookrightarrow by lcm = lcm * n[i] / d
    }
    return {ans, Lcm};
```

Listing 45: Combinatorics, FastPow, ModInv

```
//returns x such that x*a == 1 (mod b)
//x only exists if gcd(a, b) == 1
ll modInverse(ll a, ll b){
    return 1<a ? b - modInverse(b%a,a)*b/a : 1;
}

//returns a^pw % mod
ll fastPow(ll a, ll pw, ll mod) {
    ll res = 1;
    a %= mod;
    while(pw > 0) {
        if(pw & 1) res = (res*a)%mod;
        a = (a*a)%mod;
}
```

```
pw >>= 1;
   }
    return res;
struct NchooseK {
    int range, mod:
    vector<ll> fact,ifact;
    NchooseK(int currMod) {
        this->mod = currMod;
        range = (int)1e6+3;
        //range = currMod-1;//TODO: uncomment this to use chooseLucas
        fact.resize(range+1);
        ifact.resize(range+1);
        calcFacts();
   }
    void calcFacts() {
        fact[0] = 1:
        for(ll i = 1; i <= range; ++i) {</pre>
            fact[i] = (1LL*fact[i-1]*i)%this->mod;
        ifact[range] = fastPow(fact[range], this->mod-2, this->mod);
        for(int i = range-1; i >= 0; --i) {
            ifact[i] = (1LL*ifact[i+1]*(i+1))%this->mod;
        }
   }
    //helper function for calcChoose
    11 modFact(ll n, ll &e) const {
        if(n <= 1) return 1;</pre>
        11 res = modFact(n/this->mod, e);
        e += n/this->mod;
        if ((n/this->mod)\%2 == 1) return

    res*(fact[n%this->mod]*(this->mod-1)%this->mod)%this->mod;

        return res*fact[n%this->mod]%this->mod;
   }
    //n choose k with n \ge mod
    //***use prime moduli***
   11 calcChoose(ll n, ll k) const {
        11 \ e1 = 0, \ e2 = 0, \ e3 = 0;
        11 a1 = modFact(n, e1);
        11 a2 = modFact(k, e2);
        11 a3 = modFact(n-k, e3);
        if (e1-e2-e3 > 0) return 0;
        return (a1*fastPow (a2*a3%this->mod, this->mod-2, this->mod)%this->mod);
   }
   //classic n choose k
    //***use prime moduli***
   11 choose(int n, int k) const {
        if(k < 0 | | k > n | | n < 0) return 0;
        return ((1LL*fact[n]*ifact[k])%this->mod * 1LL*ifact[n-k])%this->mod;
   }
   //lucas theorem to calculate n choose k in O(\log(k))
    //***use prime moduli***
    //***can only use with: prime moduli < 1e6***
    11 chooseLucas(11 n, 11 k) const {
```

Listing 46: Commonly Used Primes

```
10<sup>3</sup> + {-9,-3,9,13}
10<sup>6</sup> + {-17,3,33}
10<sup>9</sup> + {7,9,21,33,87}
```

Listing 47: Count Primes

```
const int MAX =1000005:
bool prime[MAX];
int prec[MAX];
vector<int> P:
11 rec(ll N, int K) {
    if (N <= 1 || K < 0) return 0;
    if (N \leq P[K]) return N-1;
    if (N < MAX && ll(P[K])*P[K] > N) return N-1 - prec[N] + prec[P[K]];
    const int LIM = 250:
    static int memo[LIM*LIM][LIM];
    bool ok = N < LIM*LIM:</pre>
    if (ok && memo[N][K]) return memo[N][K];
    ll ret = N/P[K] - rec(N/P[K], K-1) + rec(N, K-1);
    if (ok) memo[N][K] = ret;
    return ret;
// init_count_primes();
// count_primes(x);
// Time complexity: Around O(N ^ 0.75)
// Constants to configure:
// - MAX is the maximum value of sqrt(N) + 2
// increase MAX to increase time efficiency
11 count_primes(11 N) {
    if (N < MAX) return prec[N];</pre>
    int K = prec[(int)sqrt(N) + 1];
    return N-1 - rec(N, K) + prec[P[K]];
void init_count_primes() {
    prime[2] = true;
```

ll expo(ll x, ll n, ll m){ ll res = 1:

```
for (int i = 3: i < MAX: i += 2) prime[i] = true:</pre>
for (int i = 3; i*i < MAX; i += 2)
    if (prime[i])
        for (int j = i*i; j < MAX; j += i+i)</pre>
             prime[j] = false;
for(int i = 0; i < MAX; ++i) if (prime[i]) P.push_back(i);</pre>
for(int i = 1; i < MAX; ++i) prec[i] = prec[i-1] + prime[i];</pre>
```

Listing 48: Discrete Log

```
while (n) {
        if (n & 1) res = (res * x) % m;
        x = (x * x) % m;
        n >>= 1;
   return (res % m);
ll extended_gcd(ll a, ll b, ll& x, ll& y){
    /// Bezout's identity, ax + by = gcd(a,b)
    if (!b){
       y = 0, x = 1;
        return a;
   }
   11 g = extended_gcd(b, a % b, y, x);
   y = ((a / b) * x);
   return g;
ll inverse_modulo(ll a, ll m){
   /// inverse exists if and only if a and m are co-prime
   11 x, y, inv;
    extended_gcd(a, m, x, y);
    inv = (x + m) \% m:
    return inv;
    * returns smallest x such that (g^x) \% p = h, -1 if none exists
    * p must be a PRIME
    * function returns x, the discrete log of h with respect to g modulo p
ll discrete_log(ll g, ll h, ll p){
    if (h \ge p) return -1;
    if ((g \% p) == 0){
        if (h == 1) return 0; /// return -1 if strictly positive lleger solution is
            \hookrightarrow required
        else return -1;
   }
   unordered_map <11, 11> mp;
   ll i, q, r, m = ceil(sqrt(p)); /// may be change to sqrtl() ?
   ll d = 1, inv = expo(inverse_modulo(g, p), m, p);
    for (i = 0; i \le m; i++)
        if (!mp[d]) mp[d] = i + 1;
        d *= g;
        if (d >= p) d \%= p;
```

```
for (q = 0; q \le m; q++){
    r = mp[d];
    if (r) return ((m * q) + (--r));
    d *= inv:
    if (d >= p) d %= p;
return -1;
```

Listing 49: Fast Fourier Transform (FFT)

```
/* emaxx implementation */
/* Multiplication with arbitrary modulos
      use ntt if mod is prime and can be written as 2**k * c + 1
      if not, use Chinese Reminder Theorem
      or transform A(x) = A1(x) + A2(x)*c decompose into A(x)/c and A(x)/c
                    B(x) = B1(x) + B2(x)*c
          where c ~= sart(mod)
          A * B = A1*B1 + c*(A1*B2 + A2*B1) * c**2(A2*B2)
          with all values < sqrt(mod) subpolynomials have coefficientes < mod * N after
      \hookrightarrow fft multiply decreasing changes of rounding error
const double PI=acos(-1);
typedef complex<double> base;
inline void fft (vector<base> & a. bool invert) {
    int n=(int) a.size();
    for (int i=1, j=0; i<n; ++i) {
         int bit=n>>1;
        for (;j>=bit;bit>>=1)
            j-=bit;
         j+=bit;
        if(i<j)</pre>
             swap(a[i],a[j]);
    for (int len=2; len<=n; len<<=1) {</pre>
         double ang = 2*PI/len * (invert ? -1 : 1);
         base wlen(cos(ang), sin(ang));
        for (int i=0; i<n; i+=len) {</pre>
             base w(1);
            for (int j=0; j<len/2; ++j) {</pre>
                 base u=a[i+j], v=a[i+j+len/2]*w;
                 a[i+j]=u+v;
                 a[i+j+len/2]=u-v;
                 w*=wlen;
        }
    }
    if (invert)
        for(int i=0;i<n;++i)</pre>
            a[i]/=n;
// a, b => coefs to multiply, res => resulting coefs
// a[0]. b[0]. res[0] = coef x^0
inline void multiply (const vector<int> & a, const vector<int> & b, vector<int> & res)
     \hookrightarrow {//change res to ll if needed
    if(a.size() * b.size() <= 256) {</pre>
```

struct Fraction {

```
res.resize(a.size() + b.size(), 0);
    for(int i = 0; i < (int)a.size(); i++)</pre>
        for(int j = 0; j < (int)b.size(); j++)</pre>
            res[i + j] += 1LL * a[i] * b[j];
    return;
}
vector<base> fa (a.begin(), a.end()), fb (b.begin(), b.end());
while (n<max(a.size(),b.size())) n<<=1;</pre>
n<<=1:
fa.resize(n),fb.resize(n);
fft (fa,false); fft(fb,false);
for (size_t i=0; i<n; ++i)</pre>
    fa[i]*=fb[i];
fft (fa, true);
res.resize (n);
for(size_t i=0; i<n; ++i)</pre>
    res[i]=(int)(fa[i].real()>0 ? fa[i].real()+0.5 : fa[i].real()-0.5);
```

Listing 50: Fibonacci Sequence

```
unordered_map<11,11> table;
int fib(int n) {//**0(log(n))**
    if(n<2) return 1;
    if(table.find(n) != table.end()) return table[n];
    table[n] = (fib((n+1) / 2)*fib(n/2) + fib((n-1) / 2)*fib((n-2) / 2)) % mod;
    return table[n];
}</pre>
```

Listing 51: Fractions

```
int a, b;
Fraction() {
    a = 0;
    b = 1:
}
Fraction(int _a, int _b) {
    a = _a;
    b = _b;
    norm();
}
Fraction(int x) {
    a = x;
    b = 1;
Fraction operator + (const Fraction& other) const {
    return Fraction(a * other.b + b * other.a, b * other.b);
Fraction operator - (const Fraction& other) const {
    return Fraction(a * other.b - b * other.a, b * other.b);
Fraction operator * (const Fraction& other) const {
    return Fraction(a * other.a, b * other.b);
Fraction operator / (const Fraction& other) const {
    assert(other.a != 0);
    return Fraction(a * other.b, b * other.a);
```

```
bool operator < (const Fraction& other) const {</pre>
        return a*other.b < b*other.a;</pre>
    bool operator <= (const int& other) const {</pre>
        return (*this < other || *this == other);</pre>
    bool operator > (const Fraction& other) const {
        return a*other.b > b*other.a;
    bool operator == (const Fraction& other) const {
        return a*other.b == b*other.a;
    bool operator != (const Fraction& other) const {
        return !(*this == other);
    bool operator == (int& other) const {
        return a == b*other;
    bool operator != (int& other) const {
        return !(*this == other);
    void norm() {
        if (b < 0) {
            a = -a;
            b = -b;
        if (a == 0) b = 1;
            int g = __gcd(llabs(a), llabs(b));
            a /= g;
            b /= g;
istream& operator >> (istream& cin, Fraction& p) {
    cin >> p.a;
    p.b = 1;
    return cin;
ostream& operator << (ostream& cout, Fraction& p) {</pre>
    cout << p.a << '/' << p.b;
    return cout;
Fraction abs(Fraction f) {
    f.a = abs(f.a):
    f.b = abs(f.b);
    return f;
```

Listing 52: Gaussian Elimination

```
int Rank;
Fraction Det;
void gauss(vector<vector<Fraction>> &a) {
   int n = a.size();
   int m = a[0].size();//possible RTE here
   Det = 1.0, Rank = 0;
   vector<int> where(max(n, m)+1, -1);
   for(int col = 0, row = 0; col < m && row < n; ++col) {</pre>
```

```
int sel = row:
    for(int i = row+1; i < n; ++i)</pre>
        if(abs(a[i][col]) > abs(a[sel][col])) sel = i;
    if(abs(a[sel][col]) == 0) { Det = 0.0; continue; }
    for(int j = 0; j <= m; ++j) swap(a[sel][j], a[row][j]);</pre>
    if(row != sel) Det = Fraction(0,1)-Det;
    Det = Det * a[row][col]:
    where[col] = row;
    Fraction s = (Fraction(1,1) / a[row][col]);
    for(int j = 0; j <= m; ++j) a[row][j] = a[row][j] * s;
    for(int i = 0; i < n; ++i) if (i != row && abs(a[i][col]) != 0) {
        Fraction t = a[i][col];
        for(int j = 0; j <= m; ++j) a[i][j] = a[i][j] - (a[row][j] * t);
    }
    ++row, ++Rank;
}
```

Listing 53: Matrix Multiplication

```
const 11 \mod = 1e9+7;
vector<vector<ll> > mult(vector<vector<ll> > a, vector<vector<ll> > b) {
    if(a.size() == 0) return {};
   if(a[0].size() == 0) return {};
   if(b.size() == 0) return {};
   if(b[0].size() == 0) return {};
    if(a[0].size() != b.size()) return {};
    int resultRow = a.size(), resultCol = b[0].size(), n = a[0].size();
    vector<vector<ll> > product(resultRow, vector<ll>(resultCol));
   for(int i = 0; i < resultRow; ++i) {</pre>
       for(int j = 0; j < resultCol; ++j) {</pre>
            product[i][j] = 0;
            for(int k = 0; k < n; ++k) {
                product[i][j] += a[i][k] * b[k][j];
                product[i][j] %= mod;
           }
       }
   }
    return product;
vector<vector<ll> > power(vector<vector<ll> > matrix, 11 b) {
   if (b <= 1) return matrix;</pre>
   vector<vector<ll> > temp = power(matrix, b/2);
   if (b % 2 == 0) return mult(temp, temp);
    return mult(mult(temp, temp), matrix);
```

Listing 54: Pollard Rho

```
11 mulMod(l1 x, l1 y, l1 p) {
    if (y == 0) return 0;
    if (x < 100011100011100011111 / y) return x * y % p;</pre>
    ll mid = mulMod((x+x)\%p, y>>111, p);
    if (y & 1) return (mid + x) % p;
    else return mid:
11 powMod(l1 x, l1 k, l1 m) {
    if (k == 0) return 1;
```

```
if ((k & 1)) return mulMod(x,powMod(x, k-1, m), m);
    else return powMod(mulMod(x,x,m), k/2, m);
bool suspect(ll a, ll s, ll d, ll n) {
    11 x = powMod(a, d, n);
    if (x == 1) return true;
    for (int r = 0: r < s: ++r) {
        if (x == n - 1) return true;
        x = mulMod(x, x, n);
    return false;
// {2,7,61,-1}
                                    is for n < 4759123141 (= 2^32)
// {2,3,5,7,11,13,17,19,23,-1} is for n < 10^{15} (at least)
bool isPrime(ll n) {
    if (n \le 1 \mid | (n > 2 \&\& n \% 2 == 0)) return false;
    11 test[] = {2,3,5,7,11,13,17,19,23,-1};
    11 d = n - 1, s = 0;
    while (d \% 2 == 0) ++s, d /= 2;
    for (int i = 0; test[i] < n && test[i] != -1; ++i)</pre>
        if (!suspect(test[i], s, d, n)) return false;
    return true:
// Killer prime: 555555555711 (fail when not used mulMod)
11 pollard_rho(ll n, ll seed) { // always call factorize
 11 x. y:
 x = y = rand() \% (n - 1) + 1;
  int head = 1, tail = 2;
  while (true) {
   x = mulMod(x, x, n);
    x = (x + seed) \% n;
    if (x == y) return n;
    11 d = \_gcd(max(x - y, y - x), n);
    if (1 < d && d < n) return d:
    if (++head == tail) y = x, tail <<= 1;
void factorize(11 n, vector<11> &divisor) { // pollard-rho factorization
 if (n == 1) return:
 if (isPrime(n)) divisor.push_back(n);
  else {
    11 d = n;
    while (d \ge n) d = pollard_rho(n, rand() % (n - 1) + 1);
    factorize(n / d, divisor);
    factorize(d, divisor);
```

Listing 55: Mobius Prime Sieve

```
const int Max = 1e6;
int seive[Max]. mob[Max]:
//returns 0 iff there exists a prime p s.t. n\%(p^2)=0
//returns -1 iff n has an odd number of distinct prime factors
//returns 1 iff n has an even number of distinct prime factors
int mobius(int n) {
   int &temp = mob[n];
```

```
if(temp != -2) return temp;
    int factors = 0, counter = 1, prev = 0;
   while(n > 1) {
       if(prev == seive[n]) counter++;
        else counter = 1;
       if(counter == 2) {
            return temp = 0:
       prev = seive[n];
       n /= seive[n]:
       factors++;
   }
   if(counter == 2) {
       return temp = 0;
   return temp = (factors%2==0?1:-1);
void doSeive() {
   for (int i = 0; i < Max; ++i) {</pre>
       mob[i] = -2;
       seive[i] = i:
   seive[0] = seive[1] = -1;
   for (int i = 2; i*i < Max; ++i) {</pre>
       if (seive[i] == i) {
            for (int j = 2 * i; j < Max; j += i) seive[j] = min(seive[j], i);</pre>
   }
```

Listing 56: Rabin-Miller Primality Test

```
11 power(11 x, 11 y, 11 p) {
    11 res = 1;
    x = x \% p;
    while (y > 0) {
        if (y & 1) res = (res*x) % p;
        y = y >> 1;
        x = (x*x) \% p;
   }
    return res;
bool miillerTest(ll d, ll n) {
    11 a = 2 + rand() \% (n - 4);
    11 x = power(a, d, n);
    if (x == 1 \mid | x == n-1) return true;
    while (d != n-1) {
        x = (x * x) % n;
        if (x == 1) return false;
        if (x == n-1) return true;
   }
    return false;
// It returns false if n is composite and returns true if n
// is probably prime. k is an input parameter that determines
// accuracy level. Higher value of k indicates more accuracy.
```

```
bool isPrime(ll n, ll k) {
    if (n \le 1 \mid | n == 4) return false:
    if (n <= 3) return true;</pre>
    11 d = n - 1:
    while (d \% 2 == 0) d /= 2;
    for (11 i = 0; i < k; i++)
         if (!miillerTest(d, n))
              return false;
    return true;
```

Listing 57: Sum Floor Arithmetic Series

```
//computes:
//[p/q] + [2p/q] + [3p/q] + ... + [np/q]
//(p, q, n are natural numbers)
//[x] = floor(x)
11 cnt(11 p, 11 q, 11 n) {
    11 t = \_gcd(p, q);
    p = p/t;
    q = q/t;
    11 s = 0;
    11 z = 1;
    while((q > 0) \&\& (n > 0)) {
        //(point A)
        t = p/q;
        s += z*t*n*(n+1)/2;
        p -= q*t;
        //(point B)
        t = n/q;
        s += z*p*t*(n+1) - z*t*(p*q*t + p + q - 1)/2;
        n -= q*t;
        //(point C)
        t = n*p/q;
        s += z*t*n;
        n = t:
        swap(p,q);
        z = -z;
    return s;
```

Listing 58: Sum of Kth Powers

```
#define MAX 1000010
#define MOD 1000000007
//Faulhaber'the sum of the k-th powers of the first n positive integers
//1^k + 2^k + 3^k + 4^k + \dots + n^k
//0(k*log(k))
//Usage: lgr::lagrange(n, k)
namespace lgr{
    short factor[MAX]:
   int P[MAX], S[MAX], ar[MAX], inv[MAX];
    inline int expo(int a, int b){
```

```
int res = 1:
    while (b){
         if (b & 1) res = (long long)res * a % MOD;
        a = (long long)a * a % MOD;
    }
    return res;
}
int lagrange(long long n, int k){
    if (!k) return (n % MOD);
    int i, j, x, res = 0;
    if (!inv[0]){
        for (i = 2, x = 1; i < MAX; i++) x = (long long)x * i % MOD;
        inv[MAX - 1] = expo(x, MOD - 2);
        for (i = MAX - 2; i >= 0; i--) inv[i] = ((long long)inv[i + 1] * (i + 1)) %
    }
    k++;
    for (i = 0; i <= k; i++) factor[i] = 0;</pre>
    for (i = 4; i <= k; i += 2) factor[i] = 2;
    for (i = 3; (i * i) \le k; i += 2){
         if (!factor[i]){
             for (j = (i * i), x = i << 1; j <= k; j += x){
                 factor[j] = i;
        }
    }
    for (ar[1] = 1, ar[0] = 0, i = 2; i \le k; i++){
        if (!factor[i]) ar[i] = expo(i, k - 1);
         else ar[i] = ((long long)ar[factor[i]] * ar[i / factor[i]]) % MOD;
    for (i = 1; i <= k; i++){</pre>
        ar[i] += ar[i - 1];
         if (ar[i] >= MOD) ar[i] -= MOD;
    }
    if (n <= k) return ar[n];</pre>
    P[0] = 1, S[k] = 1;
    for (i = 1; i \le k; i++) P[i] = ((long long)P[i - 1] * ((n - i + 1) % MOD)) %
    for (i = k - 1; i \ge 0; i--) S[i] = ((long long)S[i + 1] * ((n - i - 1) % MOD))
         \hookrightarrow % MOD;
    for (i = 0; i \le k; i++){
        x = (long long)ar[i] * P[i] % MOD * S[i] % MOD * inv[k - i] % MOD * inv[i] %
              \hookrightarrow MOD;
        if ((k - i) & 1){
             res -= x;
             if (res < 0) res += MOD;</pre>
        } else{
             if (res >= MOD) res -= MOD;
    }
    return (res % MOD):
}
```

Listing 59: Euler's Totient

```
// Euler's totient function counts the positive integers
```

Listing 60: MISCELLANEOUS

Listing 61: Random Numbers

```
//Source: http://codeforces.com/blog/entry/61587
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
inline int getRand(int 1, int r) {
    uniform_int_distribution<int> uid(1, r);
    return uid(rng);
}
```

Listing 62: Josephus Problem

```
/// Josephus problem, n people numbered from 1 to n stand in a circle.
/// Counting starts from 1 and every k'th people dies
/// Returns the position of the m'th killed people
/// For example if n = 10 and k = 3, then the people killed are 3, 6, 9, 2, 7, 1, 8, 5,
     \hookrightarrow 10, 4 respectively
int josephus(int n, int k, int m){
    for (m = n - m, i = m + 1; i \le n; i++){
        m += k;
        if (m >= i) m %= i;
    return m + 1;
/// O(k log(n))
long long josephus2(long long n, long long k, long long m){ /// hash = 583016
    if (k <= 1) return n - m;</pre>
    long long i = m;
    while (i < n){
        long long r = (i - m + k - 2) / (k - 1);
        if ((i + r) > n) r = n - i;
        else if (!r) r = 1;
        i += r;
```

Random numbers and permutations

```
m = (m + (r * k)) % i;
}
return m + 1;
}
```

Listing 63: Farey Sequence

```
//Farey sequence, grows n^2/3 with n
def farey_function(n, descending=False):
    """Print the nth Farey sequence, either ascending or descending."""
    a, b, c, d = 0, 1, 1, n
    if descending:
        a, c = 1, n-1
    print "%d/%d" % (a,b)
    while (c <= n and not descending) or (a > 0 and descending):
        k = int((n + b) / d)
        a, b, c, d = c, d, (k*c-a), (k*d-b)
print "%d/%d" % (a,b)
```

Listing 64: Python Test Cases

```
# List of N random integers from interval [a,b]:
print( [ randint(a, b) for i in range(N) ] )
# List of N random real numbers from interval [a,b]:
print( [ uniform(a, b) for i in range(N) ] )
# Random permutation of numbers 1 to N:
seq = list(range(1, N+1))
shuffle(seq)
print(seq)
# Random matrices
# Matrix of size NxN with random integers from interval [a,b]:
for r in range(N):
   print( [ randint(a, b) for c in range(N) ] )
# With zeros on diagonal:
for r in range(N):
   print([randint(a, b) if r != c else 0 for c in range(N)])
# Symmetric matrix:
matrix = [ [0] * N for r in range(N) ]
for r in range(N):
   for c in range(r+1):
        matrix[r][c] = matrix[c][r] = randint(a, b)
print(matrix)
# If you want to output matrix with spaces between the numbers:
for row in matrix:
   print( ' '.join( map(str, row) ) )
# Random tree on N vertices:
# Output is the list of edges. Vertices are numbered from O - (N-1).
print( [ (randint(0, i), i+1) for i in range(N-1) ] )
# As Michal SForiek pointed out, the above code will generate rather shallow trees, with
    \hookrightarrow expected depth only O(\log(N)).
```

```
# To generate the deep trees use:
alpha = 3 # affects the depth of the tree. Smaller value generates deeper trees.
# If alpha == 0, then the code will generate a path, i.e. the deepest possible tree.
print( [ (randint(max(0, i-alpha), i), i+1) for i in range(N-1) ] )
\# Random graph on \mathbb N vertices (may be unconnected):
# Output is the list of edges. Vertices are numbered from 0 - (N-1).
print([(i,j) for i in range(N) for j in range(i) if randint(0,1)])
# Random connected graph on N vertices:
# Simply union the edges of random tree and random graph:
print(set((randint(0, i), i+1) for i in range(N-1)) | set((i,j) for i in range(N) for j
    \hookrightarrow in range(i) if randint(0,1)))
# Random string of length N:
# Letters of your choice:
print( ''.join( choice('ABCabc123') for i in range(N) ) )
# Upper-case / lower-case letters:
from string import *
print( ''.join( choice(ascii_uppercase) for i in range(N) ) ) # upper-case
print( ''.join( choice(ascii_lowercase) for i in range(N) ) ) # lower-case
# matching a regex
import re
letters_re = re.compile(r'[A-Za-z0-9]') # regex of letters we are interested in
print( ''.join(chr(c) for c in range(256) if letters_re.match(chr(c))) )
```

Listing 65: The only thing that matters

```
inoremap kj <esc>
set ts=4 sw=4 sts=0
set expandtab
syntax on
```