

Adaptive relaying for wireless sensor networks leveraging wake-up receiver

Nour El Hoda Djidi, Antoine Courtay, Matthieu Gautier, Oliver Berder
Univ Rennes, CNRS, IRISA
firstname.lastname@irisa.fr

Abstract—Energy management is a perennial issue in the design of wireless sensor networks that require long-term sustainability. Radio transmissions being the main energy consuming tasks, adaptive transmission power and relaying protocols represent promising solutions to increase network lifetime. In this work, a novel MAC protocol leveraging emerging wake up receivers is proposed for multihop wireless sensor networks. It uses energy-based back-off mechanism to allow nodes to choose the best transmission schemes (relay or not) and corresponding transmission powers. The proposed approach has been implemented on a real hardware platform. Performance evaluation combining analytical models and microbenchmarks demonstrates a potential lifetime gain up to 1.7 when two relays are used.

Index Terms—Wireless sensor networks, Adaptive MAC protocol, Wake-up radio

I. INTRODUCTION

Billions of communicating objects have already been deployed throughout the world in the recent years, and emerging applications will dramatically increase this phenomenon. Energy is the main concern of wireless sensor designers and operators, since they are often placed in harsh environments, and obviously the more sensors you deploy, the less often you want to change the batteries. As radio transmissions represent the main source of energy consumption, the challenge is twofold: (i) reduce the transmission power, adapting the latter to the distance between nodes [1], using channel coding or cooperation between nodes [2], (ii) decrease radio activity thanks to efficient Medium Access Control (MAC) protocols [3], [4] or emerging technologies at radio front-end level.

When the addressee is far from the source node (eventually out of the radio range of the latter) multi-hop routing is commonly used, which consists in using one or several relays to help forwarding information. Since several relays are generally available at each step on the route to the destination, there is however a need of protocol overhead to choose the best relay (regarding criteria as local or global energy, latency ...) and even to choose relaying rather than direct transmission if the destination is directly reachable.

Some recent advances in ultra low power radio receivers made it possible to let the node awake and continuously listen to the wireless channel while consuming few microwatts. These devices, so-called Wake-up Receivers (WuRx), represent a real breakthrough in wireless protocols since the latency is quasi reduced to nothing with very low energy consumption [5]. The main transceiver wakes up only when a specific signal, called Wake-Up Beacon (WUB), is received

by the WuRx. Protocol overhead is also severely decreased, and furthermore it can be used to select the best relay among a certain set, e.g. based on a back-off with energy information.

Among the recent MAC protocols leveraging WuRx, OP-WUM (Opportunistic MAC Protocol Leveraging Wake up Receivers) uses timer-based contention to allow nodes to opportunistically select a receiver among their potential receivers at each packet sending [6]. On the other hand T-ROME and CTP-WUR are a cross-layer multi-hop wake-up routing protocols that combine wake-up and communication radios, especially taking into account the smaller range of wake-up receivers [7], [8]. In the present paper we present the design of a novel cross-layer protocol, called WARP (WuRx-based Adaptive Relaying Protocols) that combines the advantages of opportunistic relay selection and multi-hop schemes, both leveraging WuRx. This protocol is evaluated in terms of energy consumption conducting real experiments. Energy traces are extracted to feed an energy model that highlights the interest of our protocol and verifies its scalability.

The rest of the paper is structured as follows. Section II describes our adaptive protocol and details the relaying policy, mainly based on energy consumption. Experimental validations with real wake-up receivers are proposed in Section III, before conclusions are drawn in Section IV.

II. ADAPTIVE RELAYING MAC PROTOCOL

A. Relaying policy

The network is composed of a transmitting node, the sender, that sends data packet to the destination node and a set of N potential relay nodes. The decision to relay or not a packet is done by the sender considering relaying capability of the relay nodes. The sender also adapts its transmitting power according to the chosen strategy, resulting in a power consumption p_{Sender} : P_{SR} , the power consumption needed by the sender to reach the relay or P_{SD} to reach the destination. In the same manner, the destination power consumption p_{Dest} can be either P_{SD} or P_{RD} , the power consumption needed to link the relay and the destination. This study clearly depends on power profile of the radio transceivers: indeed, depending on the used circuits, the power consumption can vary of a factor 1.5 to 4 between the minimum and the maximum transmitting power levels [1]. If the transmission energy needed for P_{SR} is in the same range than the energy for P_{SD} , relaying will never be an interesting scheme.

We denote η a variable that is equal to 0 if direct transmission is chosen and $\eta = 1$ if relaying is considered. The total energy transmission to send a packet is equal to:

$$E_{\text{Tot}}(p_{\text{Sender}}, \eta, p_{\text{Dest}}) = E_{\text{Sender}}(p_{\text{Sender}}) + E_{\text{Relay}}(\eta) + E_{\text{Dest}}(p_{\text{Dest}}), \quad (1)$$

with E_{Sender} , E_{Relay} and E_{Dest} the energy consumed by the sender, the relay and the destination respectively. In the case of direct transmission, the total energy is $E_{\text{Tot}}^D = E_{\text{Tot}}(P_{SD}, 0, P_{SD})$ while sending a packet requires $E_{\text{Tot}}^R = E_{\text{Tot}}(P_{SR}, 1, P_{RD})$ using relaying.

Different policies can be applied to decide to use one solution or another depending on the metric to optimize e.g. network lifetime, fair packet rate among all senders ... This issue is not fully explored in this paper and a basic policy is proposed to evaluate the feasibility of the protocol. In this study, the goal is to maximize the sender lifetime i.e. until its residual energy is over a hardware failing threshold E_S^{fail} . To this aim, the sender firstly uses all the relays to decrease its energy consumption. Afterwards, it uses direct transmission if he is still alive while all the relays are in a distress state i.e. their residual energy $E_S(i)$, $\forall i \in \{1, \dots, N\}$ is below a critical threshold E_S^{empty} . Consequently, the sender chooses relaying scheme if these two conditions are true:

$$E_{\text{Sender}}(P_{SR}) < E_{\text{Sender}}(P_{SD}), \quad (2)$$

and

$$\exists i \in \{1, \dots, N\} / E_S(i) > E_S^{\text{empty}}. \quad (3)$$

The lifetime T_{Life} of the sender is thus composed of two phases: the first of duration T_{Relay} with relaying transmission and the second lasting T_{Direct} with direct transmission:

$$T_{\text{Life}} = T_{\text{Relay}} + T_{\text{Direct}}. \quad (4)$$

The first term depends on which type of nodes reaches the first its energy lower bound i.e. the sender is in power failure or all relays are in a distress state. The energy storage device of all nodes is assumed to have a finite capacity E_S^{full} . Considering a uniform use of the N relays, the relaying time is computed by:

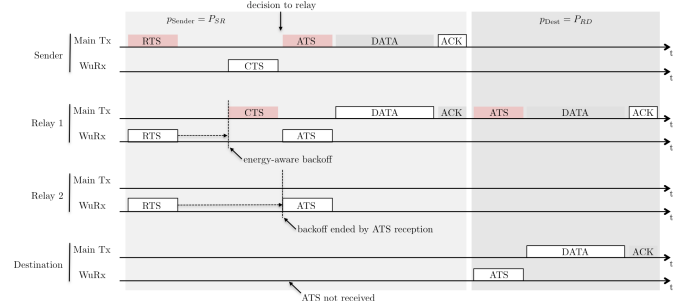
$$T_{\text{Relay}} = \min\left(\frac{E_S^{\text{full}} - E_S^{\text{fail}}}{P_{\text{Sender}}(P_{SR})}, N \frac{E_S^{\text{full}} - E_S^{\text{empty}}}{P_{\text{Relay}}(1)}\right), \quad (5)$$

with $P_{\text{Sender}}(P_{SR})$ and $P_{\text{Relay}}(1)$ the average power consumption of the sender and a relay when relaying is used.

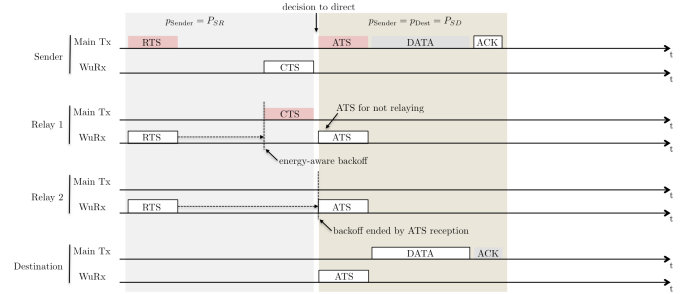
The second phase dedicated to direct transmission is computed by:

$$T_{\text{Direct}} = \frac{E_S^{\text{full}} - T_{\text{Relay}} P_{\text{Sender}}(P_{SR}) - E_S^{\text{fail}}}{P_{\text{Sender}}(P_{SD})}, \quad (6)$$

with $P_{\text{Sender}}(P_{SD})$ the average power consumption of the sender in the case of direct transmission.



(a) MAC protocol while relaying is decided.



(b) MAC protocol while direct transmission is decided.

Fig. 1: Packet transmission using WARP. Red color is used to represent WUB transmission, gray color is used to represent DATA and ACK frames transmission, and white color is used to represent reception. CS are not represented before each transmission. Background color gives the transmission powers p_{Sender} and p_{Dest} of the sender and destination nodes.

B. WARP protocol design

The WARP protocol enables a node to opportunistically select the best communication strategy between direct transmission or choosing among a set of potential relays. In order to implement the previously introduced strategy, the sender needs to know the residual energy $E_S(i)$ of the best relay.

Fig. 1 shows an example of packet transmission using WARP for both relaying (Fig. 1a) and direct transmission (Fig. 1b) schemes. In this example, the sender has two potential relays. When a node wants to send a packet, it sends a Request To Send (RTS) WUB frame. Then, a contention window begins, during which a set of potential receivers answers by sending a Clear To Send (CTS) WUB frame using a backoff, called contention backoff. The backoff is computed from a state metric, i.e. the residual energy of the relay in this study. The more residual energy the node has, the smaller the backoff should be. In this example, the relay 1 computes the shortest backoff, and is thus the first one to answer by sending a CTS WUB.

According to conditions (2) and (3), the sender chooses the best communication scheme. If relaying is decided (Fig. 1a), the sender sends an About To Send (ATS) WUB in order to inform the remaining nodes still competing that the next hop relay was already chosen. In this example, the relay 2 receives the ATS and thus cancels its backoff. Then, the data exchange

takes place, firstly between the sender and the first receiver using the main transceiver with a transmission power P_{SR} . Secondly, the relay forwards the packet to the destination using a transmission power P_{RD} .

If direct transmission is chosen (Fig. 1b), the sender sends an ATS WUB at a transmission power P_{SD} in order to inform both relay nodes and the destination that the packet will be directly transmitted to the destination. Finally, the data exchange takes place between the sender and the destination.

C. WARP energy model

The average power consumption of WARP is derived in this section. Let χ_T be the rate at which the sender transmits packets and we suppose that the relays only forward packets, they do not generate local data.

The energetic costs of transmitting and receiving a single packet (ATS, DATA and ACK) are respectively denoted E_T and E_R and are computed as follows:

$$E_T(p_T) = p_T T_{wub} + p_T T_{data} + P_R T_{ack}, \quad (7)$$

$$E_R(p_T) = P_R T_{data} + p_T T_{ack}, \quad (8)$$

where T_{wub} , T_{data} and T_{ack} are respectively the transmission durations of any WUB (RTS, CTS or ATS), a DATA frame and an ACK frame. p_T and P_R are the power consumed by the main transceiver when sending frames and listening respectively. Hence, the average power consumptions of the sender P_{Sender} , the destination P_{Dest} and the relay P_{Relay} are:

$$P_{Sender}(p_T) = P_W + \chi_T P_{SR} T_{wub} + \chi_T E_T(p_T) + (1 - \chi_T T_{wub} - \chi_T T_{Tx}) P_S, \quad (9)$$

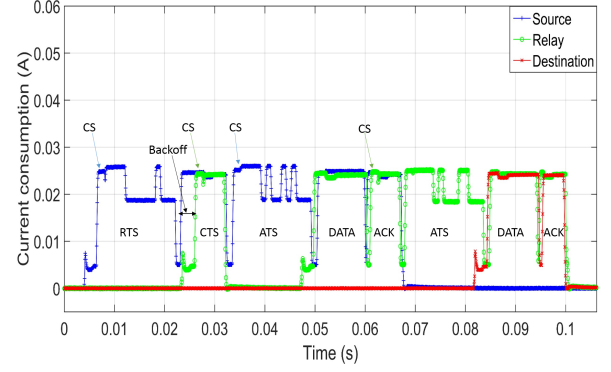
$$P_{Dest}(p_T) = P_W + \chi_T E_R(p_T) + (1 - \chi_T T_{Rx}) P_S, \quad (10)$$

$$P_{Relay}(\eta) = P_W + \chi_T P_{SR} T_{wub} + (1 - \chi_T T_{wub}) P_S + \eta [\chi_T E_R(P_{SR}) + \chi_T E_T(P_{RD}) - (\chi_T T_{Tx} + \chi_T T_{Rx}) P_S], \quad (11)$$

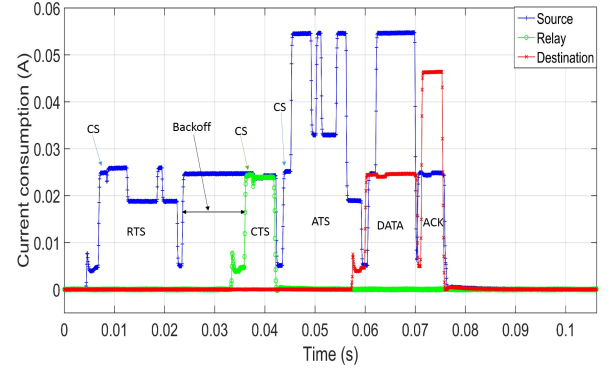
where P_W is the power consumed by the WuRx, P_S is the power consumed by the main transceiver in sleep state, and T_{Tx} and T_{Rx} are the durations during which the main transceiver is active when transmitting and receiving a single packet respectively. They can be computed as follows:

$$T_{Tx} = T_{wub} + T_{data} + T_{ack}, \quad (12)$$

$$T_{Rx} = T_{data} + T_{ack}. \quad (13)$$



(a) Relay backoff sets to 3 ms, relaying is decided.



(b) Relay backoff sets to 12 ms, direct transmission is decided.

Fig. 2: Current consumption using WARP.

III. EXPERIMENTAL VALIDATIONS

A. Experimental setup

Experimentations were realized using the PowWow platform [4] and the WuRx from [5]. PowWow is a modular platform that is equipped with a CC1120 radio chip and a MSP430 low power micro-controller, both from Texas Instruments. The data is sent with a bitrate of 20 kbps. The WuRx receives data with On-Off Keying modulation (OOK) and works in the 868 MHz frequency band with a bitrate of 1 kbps. The sensitivity in these conditions was measured to be -55 dBm. The WuRx embeds a ULP MCU, the PIC12LF1552 from Microchip, that can process address matching. The power consumption of the whole WuRx was measured to be $1.83 \mu W$ in continuous listening mode and $284 \mu W$ when receiving data with the ULP MCU active to process them.

The experimental setup consisted in three nodes: one source, one relay node and a destination. The RTS and CTS WUB were always sent at -11 dBm. The ATS WUB, DATA and ACK frames were sent at -11 dBm when relaying is decided, and were sent at 14 dBm when direct communication is decided.

B. Microbenchmarks

Fig. 2 shows the current consumptions of the sender, the selected relay and the destination for two different backoffs of the relay, 3 ms and 12 ms. Each backoff leads to a different decision, relaying in Fig. 2a and direct transmission in Fig. 2b.

Frame duration and power consumption		Energy per packet per node type and per mode	
T_{wub}	11 ms	Relaying	
T_{CTS}	4 ms	$E_{\text{Sender}}(P_{SR})$	2.57 mJ
T_{ACK}	4 ms	$E_{\text{Relay}}(1)$	2.989 mJ
T_{Data}	7.7 ms	$E_{\text{Dest}}(P_{RD})$	1.023 mJ
P_s	0.12 mW	Direct transmission	
P_R	73.5 mW	$E_{\text{Sender}}(P_{SD})$	4.21 mJ
P_{SR}	73.05 mW	$E_{\text{Relay}}(0)$	0.289 mJ
P_{SD}	141.36 mW	$E_{\text{Dest}}(P_{SD})$	1.29 mJ

TABLE I: Energy, power consumption and time duration extracted from the microbenchmarks.

These measurements were obtained using a Keysight N6705 DC analyzer. The different stages of the proposed MAC protocol, detailed in Section II-B, can be seen in the figure. The main transceiver also offers Channel Sensing (CS) capabilities that are used in this experiment before sending any packet.

These microbenchmarks show the feasibility of the switching mechanism. In Fig. 2a when the WuRx of the relay node receives the RTS WUB, it computes a backoff of 3ms, meaning that the relay node has enough energy to forward the data to the destination, so the source decides to send data to the relay. In the contrary, in Fig. 2b, the back-off is set to 12ms, meaning that the relay has not enough energy to forward data to the destination, so the source decides to directly send data.

From the microbenchmarks, Table I gives measured values in term of energy, power consumption and time duration. These results show that the used transceiver has a P_{SD} value around twice the one of P_{SR} . Even if it is not a favorable case for relaying, $E_{\text{Sender}}(P_{SR})$ is still lower than $E_{\text{Sender}}(P_{SD})$ so relaying will be chosen until either the sender or the relays are in distress states.

C. Analytical evaluation of the sender lifetime

The sender lifetime is analytically evaluated by including experimental measurements given in Table I in the theoretical lifetime (4)-(6) and the energy models of WARP introduced in Section II-C. The lifetime is evaluated in both direct transmission and relaying cases. WARP is also compared to simple direct transmission with no adaptation mechanism in order to evaluate the protocol overhead. The energy storage device is a 0.9 F super-capacitor with a maximum voltage of 5.0 V, and the minimum voltage required to power the node is 2.8 V. Thus E_S^{full} and E_S^{fail} are equal to 11.25 J and 3.528 J respectively. E_S^{empty} is set to 3.534 J so that a relay can still rely twice. The results are shown in Fig. 3.

As $P_{\text{Sender}}(P_{SD})$ is higher than $P_{\text{Relay}}(1)$, the sender will die earlier when using direct transmission than when using relaying. The lifetime increases up to 1.3 times when one relay is used and up to 1.7 times with two relays. However, the sender lifetime will not increase when using more than two relays, in this case it depends on $P_{\text{Sender}}(P_{SR})$ and the sender will die before the relays. At sender death, 40% of the energy has been used per relay for $N=2$. Comparing WARP to direct transmission, one can notice the energy overhead due to the adaptation mechanism.

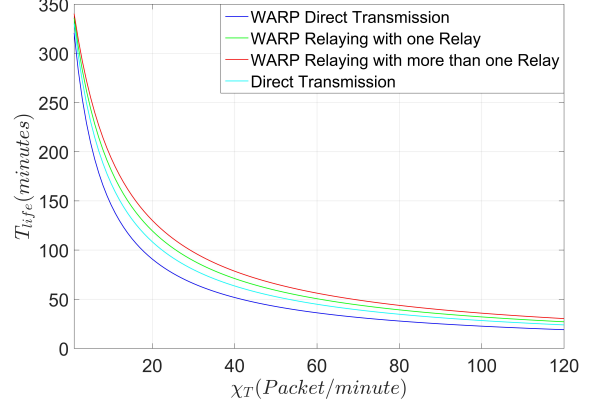


Fig. 3: Sender lifetime as a function of packet rate χ_T .

IV. CONCLUSION

This paper presents WARP, a MAC protocol that adapts the communication schemes, relaying or direct transmission, according to the residual energy of the nodes. The protocol leverages WuRx to opportunistically select the best relay and uses time backoff to send residual energy information to the sender node. The WuRx allows to significantly reduce energy consumption overhead due to the protocol to decide the best strategy. However, the protocol performance clearly depends on the power profile of the radio transceiver. With our setup, a 1.7 lifetime gain compared to direct transmission can be achieved when 2 relays are used. Future works are to use other transceivers and to extend the protocol to multi-hop multi-sender networks and associated decision policies.

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