

# DC

## Unit - 1

### Parallel Computing

\* The goal of PC is to provide performance either in terms of processor power or memory.

\* In PC distribution is frequent

\* It is lower overhead.

\* Assumed to reliable

A short time execution

### Distributed computing

\* The goal of DC is to provide conveniences includes availability, reliability and physical distribution

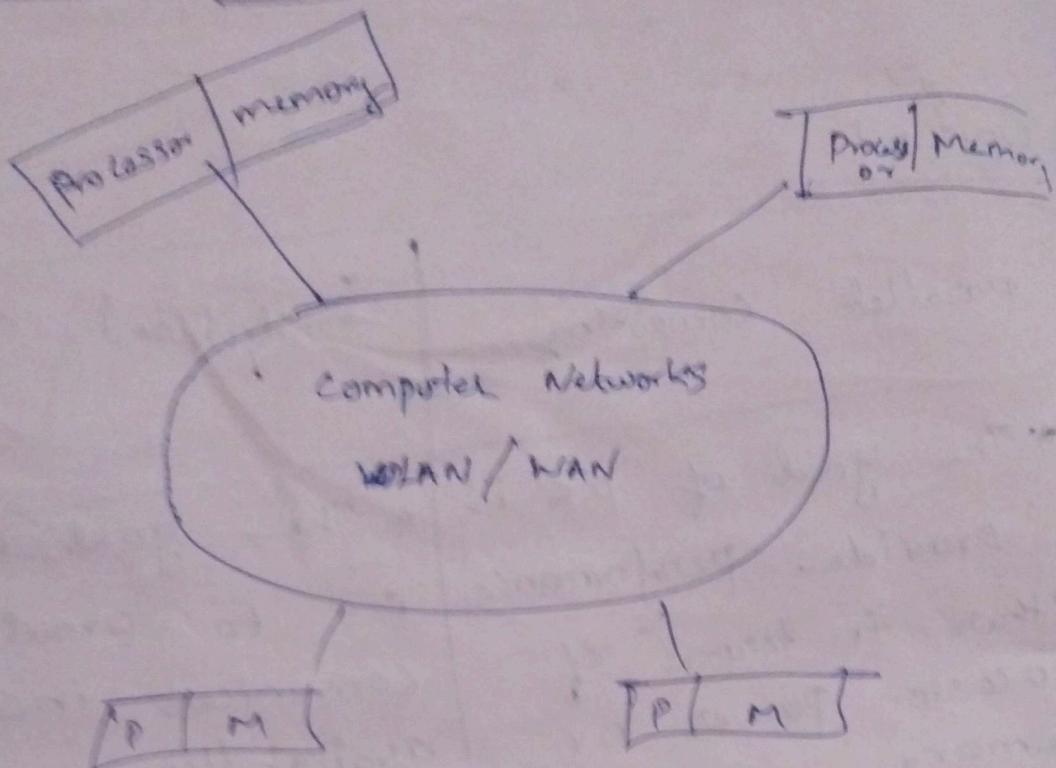
\* In DC distribution is infrequent

\* It is heavier weight

\* Assumed to unreliable

A long time execution

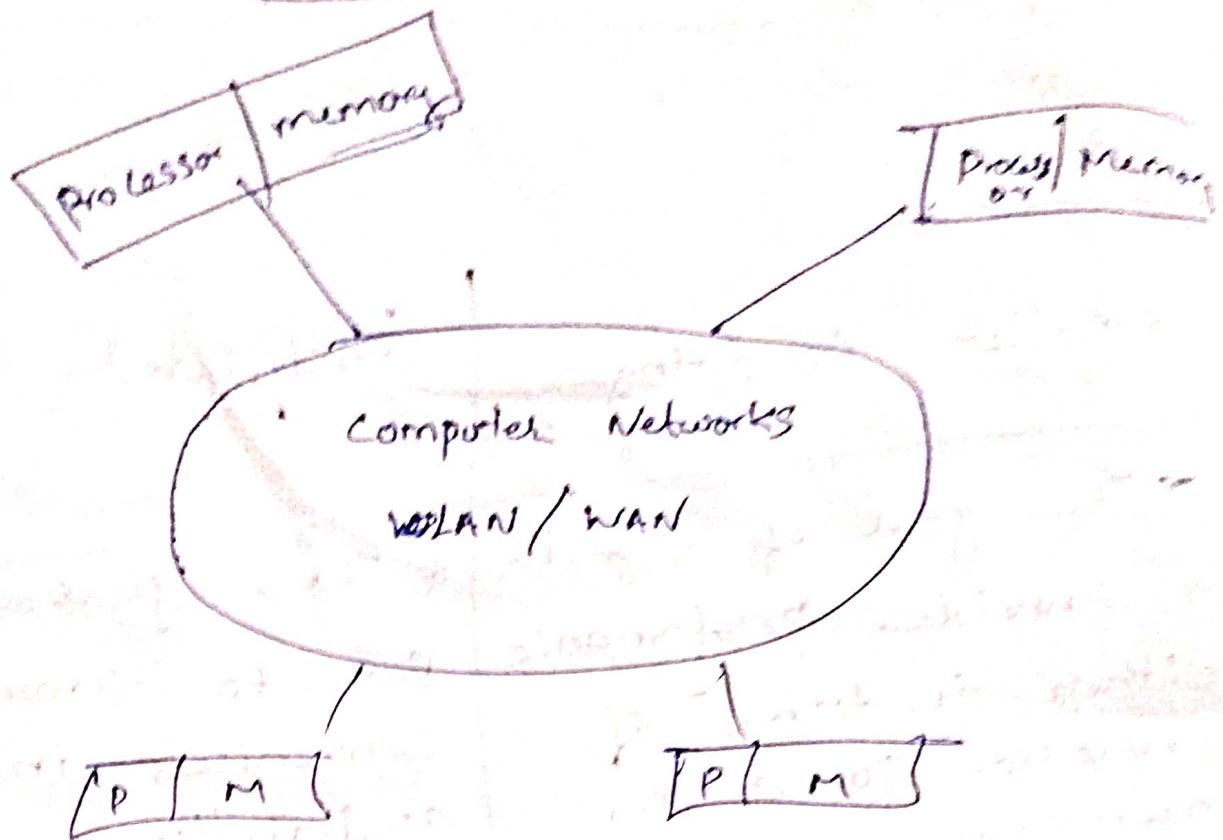
## a. Computer System Components.



## b. Disadvantages of Distributed System

- \* Software: Difficult to develop Software for ds
- \* Network: Saturation, lossy transmission.
- \* Security: Easy to access the Secret data
- \* Absence of global clock.

## 2. Computer System Components.



## 3. Disadvantages of Distributed Systems

\* Software: Difficult to develop Software for a distributed system.

\* Networks: Saturation, lossy transmissions.

\* Security: Easy to access the Secret data.

\* Absence of global clock.

#### 4. Need of DS:

\* Resource Sharing (eg): Printers, disks)

\* Security

\* Reliability

\* consistency of replicated data

\* fault tolerance

#### 5) Message-passing and Shared Memory:

\* Message Passing:

\* Two processes communicate with each other by passing messages.

\* The messages can be both direct and indirect

\* Indirect communications uses mailbox for sending and receiving messages.

\* Message Passing used as a method of communication in microkernels.

\* Message passing System requires synchronization and communication between the two processes.

\* Messages sent by Process can be either fixed or variable in size.

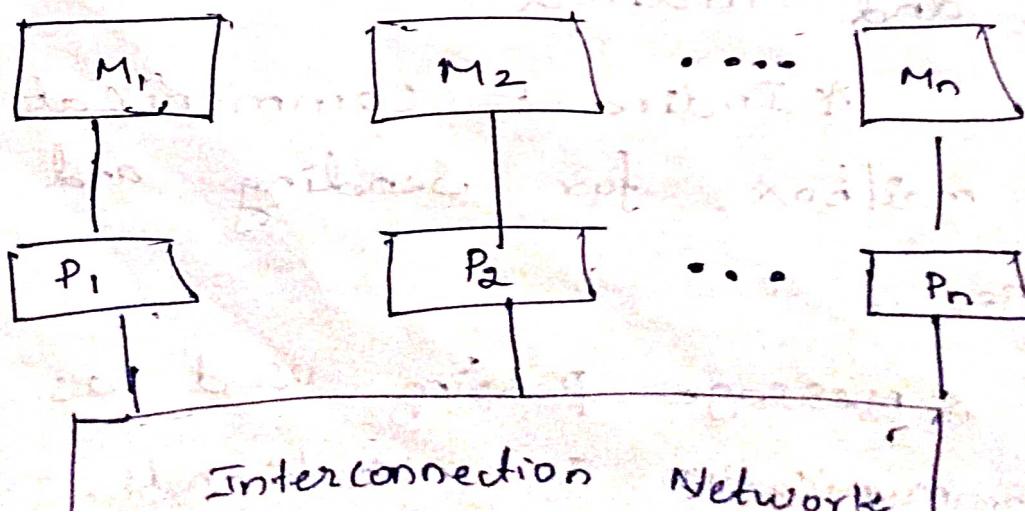
\* The actual function of message passing is normally

- i) send (destination-name, msg)
- ii) receive (source-name, msg)

\* Send Primitives is used for sending a message to destination

\* A process receives information by executing the receiving primitive

Diagram:



- \* This technique can be used in heterogeneous computers

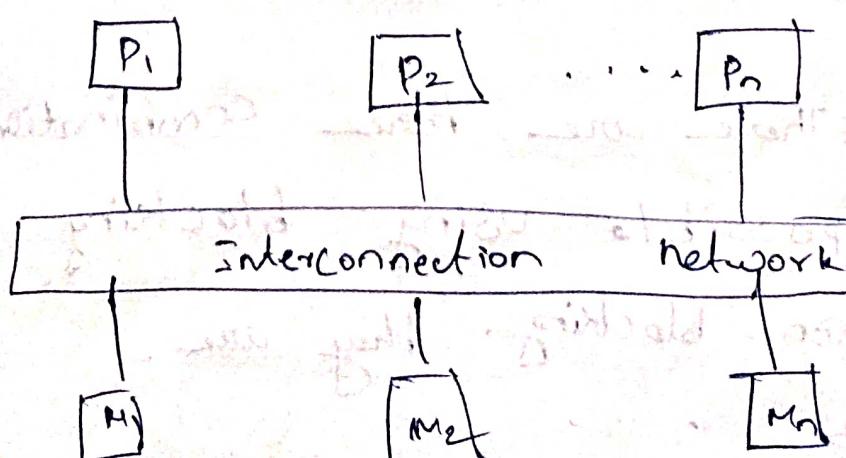
### Shared Memory:

- \* Shared memory uses common shared memory for data sharing
- \* Communication takes place via shared data variables

\* It follows faster communication strategy when compared to message passing technique.

\* It is typically implemented using system calls.

\* In shared memory region, the information can be exchanged by reading and writing data.



## 6. Primitives forms DC (superset of DCE)

\* Message Communication primitives are denoted by `Send()` and `receive()`

\* Message passing primitive commands  
`SEND(msg, dest)`  
`RECEIVE(src, buffer)`

\* Send primitives uses two options for sending data:

\* Buffered

\* unBuffered

\* In buffered option, user data is copied in the kernel buffer. In unbuffered option data gets copied directly from the user buffer to the network.

\* There are three combinations are possible using blocking and non-blocking. They are -

i) Blocking Send, Blocking Receive

ii) Non Blocking Send, blocking Receive

iii) Non Blocking Send, Nonblocking Receive

i) Blocking Send , Blocking Receive:

\* In this method both Sender() and Receiver() are blocked until the message is delivered.

\* This is called Rendezvous.

\* This combination allows for tight synchronization between processes.

ii) Nonblocking Send , Blocking Receive:

\* Sender may continue on, the receiver is blocked until message arrives.

iii) Nonblocking Send , Non-Blocking Receive:

Sending process send message and resume the operation. Receiver receives either a valid message or null message.

## 7. Synchronous Execution:

\* Synchronous execution means

the first task in a program

must finish before processing

the next tasks.

\* Lower and upper bounds

on execution times of processes

can be set.

\* Transmitted messages can

be received within a known

time bound.

\* Drift rates between local  
clocks have a known bound.

\* In synchronous distributed  
systems there is a notation of  
global physical time.

\* Only synchronous distributed  
systems have predictable behavior

in terms of time.

- \* In Synchronous distributed systems it is possible to set timeouts in communication link and failure of process
- \* It is difficult and costly to implement Synchronous distributed systems.

## Asynchronous Execution:

- \* Asynchronous execution means a second task can begin execution before the completion of first task
- \* No bounds on process execution time.
- \* No bound - on message transmission delay
- \* No bounds on drift takes between local clocks

\* No timeouts can be used  
in the case of ~~asynchronous~~ distributed systems.  
\* Asynchronous distributed systems  
are unpredictable in terms of  
timing.

\* Asynchronous are widely used  
for practice and failure detection of  
nodes.

## 8. Design Issues and Challenges:

### Challenges from System Perspective:

\* communication mechanisms.

This task involves designing appropriate mechanism for communication among processes in the network.

\* processes:

Issues involved are

Code migration, thread management at client and server.

## Challenges :

1. Heterogeneity
2. Openness
3. Security
4. Scalability
5. Failure handling
6. Concurrency
7. Transparency.

## Heterogeneity :

modern distributed systems are highly heterogeneous in many dimensions, including available bandwidth, processor speed, disk capacity, security, failure rate. It applies.

1. Computer Network - LAN, WAN, wireless Network, Satellite link
2. Computer hardware - Laptop, Computer, Mobile
3. Operating System : Linux, UNIX, Windows
4. programming language: C, C++, Java, PHP

## Openness:

Openness means that the system can be easily extended and modified.

Openness means ability to plug and play.

& The integration of new components must have the ability to communicate with already exist components in the system.

## Security:

Security becomes even more important in distributed systems.

Authentication, authorization, digital signatures, encryption and

privacy become major issues in the distributed system.

## \*Scalability:

\*A system is said to be Scalable if it can handle the addition of users and resources without buffering or loss of performance.

Scaling techniques

1. Hiding Communications latency.

2. Distribution.

3. Replication.

## \*Failure Handling:

\*Failure in distributed system occurs when there is a fault in program logic or hardware.

Hardware, Software, and networks are not free of failures.

## Techniques for dealing with failure:

1) Detecting failure - finding

2. Masking failure - Hidden

3. Tolerating failure -

4. Recovering failure

5. Redundancy -

## Concurrency:

### Global State of DS.

\* Garbage collection

\* Deadlock

\* Termination

\* Debugging