

### Unit - 3

#### 1.) Mathematics of Asymmetric Key

\* It is classified into six types they are

\* Prime Numbers

\* Primality Testing

\* Factorization

\* Euler's Theorem

\* Chinese Remainder Theorem

\* Exponentiation and logarithm

#### Prime Number:

\* An integer  $p > 1$

\* If its only divisors of 1 and  $p$

\* It is also known as fundamental theorem of arithmetic



## Primality Testing:

\* A primality test is an algorithm used for determining whether the input is prime.

\* It is used for cryptography.

\* Primality tests do not generally give prime factors, it states whether the input number is prime or not.

\* Some primality tests prove that a number is prime they are

1) The sieve of Eratosthenes

2) Prime factorization

## Euler's theorem:

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

Consider,  $a = 3$  and  $n = 10$

$$3^{\phi(10)} \equiv 1 \pmod{10}$$

$$\phi(10) = 4$$



$$3^4 = 1 \pmod{10}$$

$$81 = 1 \pmod{10}$$

\* The numbers  $a$  and  $n$  must be prime.

### Chinese Remainder Theorem:

\* The Chinese remainder theorem (CRT) is used to solve a set of different congruent equation with one variable but different moduli which are relatively prime.

(\*) It is a one of most useful results of number theory.

Let us consider some linear system of congruences

$$x = a_1 \pmod{n_1}$$

$$x = a_2 \pmod{n_2}$$

...

$$x = a_n \pmod{n_n}$$



~~ACRT~~ CRT States that the equations have a unique solution of the moduli and that are relatively prime

Eg:

$$x = (a_1 M_1 M_1^{-1} + a_2 M_2 M_2^{-1} + \dots + a_n M_n M_n^{-1}) \text{ mod } M$$

Consider a set of Congruent equations

$$x = 2 \pmod{3}$$

$$x = 3 \pmod{5}$$

$$x = 2 \pmod{7}$$

$x = a_1 \pmod{m_1}$	$x = 2 \pmod{3}$
$x = a_2 \pmod{m_2}$	$x = 3 \pmod{5}$
$x = a_3 \pmod{m_3}$	$x = 2 \pmod{7}$

$$x = (a_1 M_1 M_1^{-1} + a_2 M_2 M_2^{-1} + a_3 M_3 M_3^{-1}) \text{ mod } M$$

Given

$$a_1 = 2$$

$$m_1 = 3$$

$$a_2 = 3$$

$$m_2 = 5$$

$$a_3 = 2$$

$$m_3 = 7$$



To find.

$$M_1$$

$$M_1^{-1}$$

$$M_2$$

$$M_2^{-1}$$

$$M$$

$$M_3$$

$$M_3^{-1}$$

To find  $M = m_1 \times m_2 \times m_3$

$$M = 3 \times 5 \times 7$$

$$\boxed{M = 105}$$

To find  $M_1 = \frac{M}{m_1} = \frac{105}{3} = 35$

To find  $M_2 = \frac{M}{m_2} = \frac{105}{5} = 21$

To find  $M_3 = \frac{M}{m_3} = \frac{105}{7} = 15$

To find inverse of M

$$M_1 \times M_1^{-1} = 1 \pmod{m_1}$$

$$35 \times M_1^{-1} = 1 \pmod{3}$$

$$35 \times 2 = 1 \pmod{3}$$

$$M^{-1} = 2$$

$$M_2 \times M_2^{-1} = 1 \text{ Mod } m_2$$

$$21 \times M_2^{-1} = 1 \text{ Mod } 5$$

$$21 \times 1 = 1 \text{ Mod } 5$$

$$M_2^{-1} = 1$$

$$M_3 \times M_3^{-1} = 1 \text{ Mod } m_3$$

$$15 \times M_3^{-1} = 1 \text{ Mod } 7$$

$$15 \times 1 = 1 \text{ Mod } 7$$

$$M_3^{-1} = 1$$

$$= (2 \times 35 \times 2 + 3 \times 21 \times 1 + 2 \times 15 \times 1) \text{ Mod } 105$$

$$= 233 \text{ Mod } 105$$

$$x = 23 //$$



## 2) RSA Algorithm:

\* The Rivest - Shamir - Adleman (RSA) invented this algorithm for converting plaintext to ciphertext.

\* The typical size for  $n$  is 1024 bits or 309 decimal digits.

### Algorithm:

\* pick two large prime numbers  $p$  and  $q$  such that  $p \neq q$ .

\* calculate  $n = p \times q$

\* calculate  $\phi(n) = (p-1) \times (q-1)$

\* pick "e" so that  $\text{GCD}(e, \phi(n)) = 1$

$$1 < e < \phi(n)$$

\* calculate encryption using

$$C = M^e \bmod n$$

Decryption using

$$M = C^d \bmod n$$

Eg:

1)  $p = 17$  (prime) - selected

$q = 11$  (prime) - selected

$e = 7$   $M = 8$

2)  $n = p \times q$

$n = 17 \times 11$

$n = 187$

3)  $\phi(n) = (p-1)(q-1) = 16 \times 10 = 160$

4) Encryption

$C = M^e \mod n$

$1 = [88^7] \mod 187$

$= 11$

5) Decryption  $M = C^d \mod n$

$M = 11^{23} \mod 187$

$= 88$



### 3) Diffie Hellman Key exchange

It is an first published public key algorithm that defines public key cryptography.

The main purpose of this algorithm is to enable two users to securely exchange a key that can be used for encryption of messages.

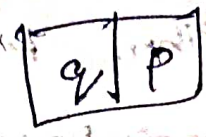
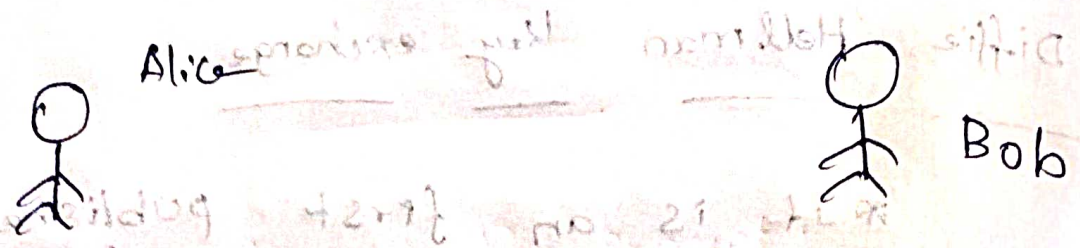
This algorithm is limited to exchange of secret values.

Why this is secure

It uses discrete logarithm algorithm to avoid forced attack.

It is used for the exchange of private key & not public key.





$a \rightarrow$  denotes private value of Alice

$b \rightarrow$  denotes private value of Bob

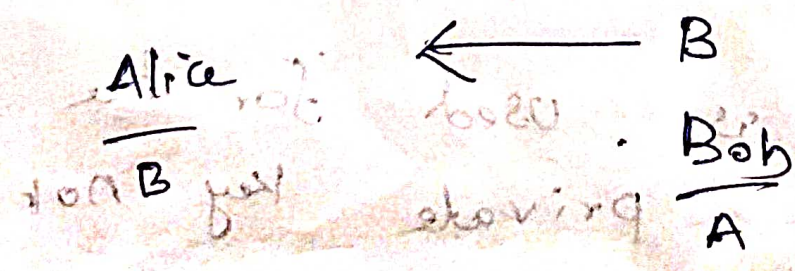
$q \rightarrow$  denotes prime Number

$p \rightarrow$  It denotes primitive root

$$A = p^a \text{ mod } q$$

$$B = p^b \text{ mod } q$$

$A \rightarrow$



$$S = B^a \text{ mod } q$$

$$S = A^b \text{ mod } q$$



Q9:

$$q=13, \quad p=6$$

Alice

Bob

$$b=10$$

private value  $a=3$

$$A = p^a \mod q \quad B = p^b \mod q$$

$$= 6^3 \mod 13 = 6^{10} \mod 13$$

$$= 8 \quad = 4$$

$$S = 8^a \mod q$$

$$S = A^b \mod q$$

$$= 4 \mod 13 \quad S = 8^{10} \mod 13$$

$$= 12$$

$$= 12$$