

Chapter 2

1. **2.10 Assume that registers x5 and x6 hold the values 0x8000000000000000 and 0xD000000000000000, respectively.**

1) The value of x30 for the following assembly code is 0x_5000000000000000_____

add x30, x5, x6

2) Is the result in x30 the desired result, or has there been overflow? (A)
A. overflow B. no overflow

3) For the contents of registers x5 and x6 as specified above, The value of x30 for the following assembly code is 0x B000000000000000_____

sub x30, x5, x6

4) Is the result in x30 the desired result, or has there been overflow? (B)
A. overflow B. no overflow

5) For the contents of registers x5 and x6 as specified above, The value of x30 for the following assembly code is 0x_D000000000000000_____.

add x30, x5, x6

add x30, x30, x5

6) Is the result in x30 the desired result, or has there been overflow?(A)
A. overflow B. no overflow

2. **2.22 Suppose the program counter (PC) is set to 0x20000000.**

1) What range of addresses can be reached using the RISC-V jump-and-link (jal) instruction? (In other words, what is the set of possible values for the PC after the jump instruction executes?)

The range is [0x_1FF00000_____, 0x_200FFFFE_____](low to high)

2) What range of addresses can be reached using the RISC-V branch if equal (beq) instruction? (In other words, what is the set of possible values for the PC after the branch instruction executes?)

The range is [0x_1FFFF000_____, 0x_20000FFE_____](from low to high)

3. **2.29 Implement the following C code in RISC-V assembly. Hint: Remember that the stack pointer must remain aligned on a multiple of 16.**

```
int fib(int n){
    if (n==0)
        return 0;
    else if (n == 1)
        return 1;
    else
        return fib(n-1) + fib(n-2);
}
```

fib:

```
beq x10, x0, finish //if n==0 return 0
addi x5, x0, 1
beq x10, x5, finish //if n==1 return 1
addi x2, x2, -16
sd x1, 0(x2) //save x1 on stack
sd x10, 8(x2) //save x10 on stack
addi x10, x10, -1 //n-1
jal x1, fib //fib(n-1)
ld x5, 8(x2) //x5 get n
sd x10, 8(x2) //push fib(n-1) onto the stack
addi x10, x5, -2 //n-2
jal x1, fib //fib(n-2)
ld x5, 8(x2) //x5 get fib(n-1)
add x10, x10, x5 //fib(n) = fib(n-1) + fib(n-2)
ld x1, 0(x2) //return saved rd
addi x2, x2, 16 //pop back
```

finish:

```
jalr x0, 0(x1) //Return to caller
```

4. 2.5 Show how the value 0xabcdef12 would be arranged in memory of a little-endian and a big-endian machine. Assume the data are stored starting at address 0 and that the word size is 4 bytes.

a big-endian machine

address	data
3	12
2	ef
1	cd
0	ab

a little-endian machine

address	data
3	ab
2	cd
1	ef
0	12

5. 2.12 Provide the instruction type and assembly language instruction for the following binary value: 0000 0000 0001 0000 1000 0000 1011 0011two

1) the type of the instruction is(A)

A. R B. I C. B D. J

2) the instruction is (A)

A. add x1,x1,x1 B. addi x1,x1,x1 C. beq x1,x2,4 D.jalr x0,0(x2)

chapter3

1. Assume decimal integers 185 and 122 are unsigned 8-bit integers, their bit patterns are A and B, now A and B represent signed 8-bit decimal integers stored in sign-magnitude format, Calculate A + B.
 - 1) The result is _65_(in decimal)
 - 2) (C) is in this calculation.
A. overflow B. underflow C. neither
2. 3.20 Given the bit pattern 0x0C000000, if it is a two's complement integer ,its decimal value is _201326592_____, and an unsigned integer is _201326592_____
3. 3.22 What decimal number does the bit pattern 0x0C000000 represent if it is a floating point number? Use the IEEE 754 standard. B
A. 1.0×2^{-101} B. 1.0×2^{-103} C. 1.05×2^{-101} D. 1.05×2^{-103}
4. 3.23 The binary representation of the decimal number 63.25 assuming the IEEE 754 single precision format is 0x_427D0000_____in hex.