# Middle East Technical University Department of Computer Engineering Wireless Systems, Networks and Cybersecurity (WINS) Laboratory



# Secure Implementation of Topological Sort of a Directed Acyclic Graph Algorithm using Homomorphic Encryption

CENG519 Network Security 2021-2022 Spring Term Project Report

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# **Abstract**

In this work, a secure implementation of topological sort of a DAG (Directed Acyclic Graph) is implemented between a server and client using homomorphic encryption. The system is built on a mock client-server architecture where client encrypts a list representation of a DAG using Microsoft SEAL library, and server makes computational operations on encrypted data. The operations that can run on encrypted data are limited, hence the more complicated operations need to be dissembled to more basic operations, making the implementation challenging. The operations that can neither be simplified nor available are instead sent back to client, which returns the outcome to server. This results in a slow solution; but ensures that the data is not decrypted at the server.

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### 1 Introduction

Encryption of data is usually relevant for storing or transportation stages of data. This is partially because traditional cryptography techniques require the data to be unencrypted to be able to process it correctly, which implies that at some stage on a supposedly secure system, the data has to be decrypted and processed in plaintext. Examples to this situation are; sensitive information of companies being processed by 3rd parties, or data analytics extracted from personal data, and other situations where the data is entrusted with the hope that is won't be misused or get stolen [2].

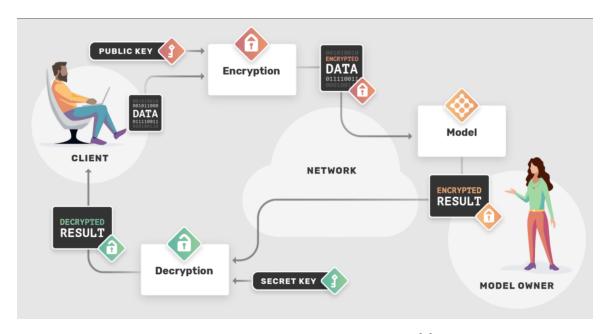


Figure 1 . Homomorphic Encryption [1]

A viable solution to this problem Homomorphic encryption. In principle, the term "homomorphic encryption" generalises encryption techniques that "allow encrypted data to be processed as if it were in plaintext and produce the correct value once decrypted" [3]. However, in order to keep the data safe, the operations that server can run on encrypted data are need to be limited. Currently there are three types of homomorphic encryption scheme that is categorized by the operations they provide [2]:

- 1. Partially homomorphic encryption: only one mathematical operation can be performed an unlimited number of times on the encrypted data.
- 2. Somewhat homomorphic encryption: any number of additions but only one multiplication operation can be performed on the encrypted data.
- 3. Fully homomorphic encryption: any number of additions and multiplications operation can be performed on the encrypted data.

While these limitations help prevent unauthorized decryption, implementing the algorithms that will process the data becomes challenging. The technique demonstrated in this project is fully homomorphic encryption. A major data representation and processing model makes use of Graph Theory; therefore we implement Topological Sorting on Directed Acyclic Graphs using fully homomorphic encryption. The encryption algorithm is provided by Microsoft

# 2 Background and Related Work

#### 2.1 Background

The project requires prerequisites on graph theory, programming with Python and general knowledge of server-client architecture.

#### 2.2 Related Work

While homomorphic encryption is not a new technique, it is still not very wide-spread, resulting in scarce related work.

Some libraries that are currently available are:

- 1. A fully homomorphic encryption library is PALISADE [5]
- 2. awesome-he by jonaschn [6]
- 3. fully-homomorphic-encryption by Google [7]
- 4. TFHE: Fast Fully Homomorphic Encryption over the Torus [8]

Microsoft SEAL is considered "leveled fully homomorphic" by some sources [9]

### 3 Main Contributions

# 3.1 Implementation of Comparison

The algorithm to be implemented is a graph sorting algorithm, and it requires comparisons on data in regular domain. However fully homomorphic encryption provides limited operations: the Microsoft SEAL library which is the backbone of the encryption system does not allow comparison operations. This problem is solved by introducing a mechanism between the server and the client to provide comparison operation.

Since the comparison operation is not an option in the encrypted data on the server side, an interface with the client is implemented. This interface should allow client to be invoked from the server, and it should allow the server to continue the processing from where it left off. In other words, the comparison operation is temporarily handed back to client to make up for the limited availability of mathematical operations on server side.

#### 3.2 Client-Server Interface

The need for such an interface is required because the server can not handle the comparison. This interface should satisfy two fundamental flows besides comparison implementation:

- 1. Server should invoke the client whenever necessary and the client should be able to handle this interrupt.
- 2. After client is processed the data and invoke the server, server should run from where it left of.

Also the data to be compare should be managed in the server. That means instead of sending data, use multiplication with an array that includes only one 1 and the other items are zero.

That results with an array that either all zero or has only one item equals to 1. Therefore index can be detected but data itself can not be accessed.

# 4 Methodology

The code simulates the selected algorithm securely. When it runs, it automatically creates a DAG with selected node size and makes operations on this graph. After resulting with the simulation, it generates results file, plots figures and, draws a visual representation of the graph. The code can be seen in Appendix C.

#### 4.1 Server Implementation

The server -which is represented by a function- should implement the selected algorithm: Kahn's Algorithm for Topological Sorting [10]. As stated on previous sections the limitations should be handled by the server by implementing client-server interface. To satisfy the requirement of client-server interface, following features should be added to server code:

- 1. Server should hold a state. The state is also the main communication mechanism between client and the server.
- 2. Server should yield a state-result pair instead of returning. Yielding provides both invoking the client and the resuming from yield location.

Server is implemented in such a manner that the only encrypted data is the graph itself. The helper variables such as in-degree array and state of the process is held in global scope. Therefore they are not protected. The reason behind that is keeping the encryption operations simple as possible due to scope of the assignment. But all of the operations can also be applied the helper variables.

There are two types of states that server provides:

- 1. State "check if all zero": It is yielded with an encrypted array that has either all of its items are zero or only the first item is one
- 2. State "finished": Sent when the sorting operation is completed.

When the only incoming states which is "not all zero" is set by the client, server knows comparison returns true.

### 4.1.1 Kahn's Algorithm for Topological Sorting

Once all the required mathematical operations are provided either by the encryption library or via client-server communication; Kahn's algorithm for topological sorting can be implemented by executing the following steps [10]:

- 1. Compute in-degree (number of incoming edges) for each of the node.
- 2. Pick all the nodes with in-degree as 0 and add them into a queue: This operation requires comparing in-degree items with 0.
- 3. Pop a node from the queue.
- 4. Decrease in-degree by 1 for all its neighbouring nodes. If in-degree of a neighbouring node is reduced to zero, then add it to the queue: This operations requires finding neighbours which includes comparison operation.
- 5. Repeat Step 3 until the queue is empty.

#### 4.2 Client Implementation

The client is responsible for providing graph and necessary helper variables stated in the server implementation section. Also it should handle the basic requests from the server by using client-server interface. To satisfy the requirement of client-server interface, following features should be added to client code:

- 1. Since the server yields results, the client should iterate the server by calling next on the generator function of the server.
- 2. For every iteration, client should check the state in the response from server and act according to it.

The incoming states is handled as follows:

- 1. State "check if all zero": This state means that server needs a comparison operation. The resulting array will be decrypted by the client and content will be checked if all the items are zero or not. If the all items are not zero, client will set global state to "not all zero" and iterate the server.
- 2. State "finished": This state means sorting operation is completed and the result is achieved.

#### 4.3 Development Environment

The building and running environment is Ubuntu Linux that runs on a Docker container. The related Dockerfile can be seen in Appendix A. Microsoft SEAL encryption library used as the core of the homomorphic encryption. The mock server-client architecture is implemented using Python programming language on Visual Studio Code editor.

#### 5 Results and Discussion

#### 5.1 Results

As we can see in the selected graph in Figure 2, there is a linear relationship between resources consumed and the node count. The pitfalls represent that algorithm is failed to find a topological sort, therefore run time is decreased at that points. Related graphs can be seen in the Appendix A.

#### 5.2 Discussion

The results can be observed from two different perspectives. First of them is the correctness of the results and the other is efficiency of the algorithm. Correctness of the algorithm is verified by network library functions and it is seen that results are not correct sometimes. The pitfall in the graphs shows that algorithm is failed and therefore run time is slow. Efficiency of the algorithm, on the other hand, should be increased. Because the data is transmitted between server and the client multiple times for only one operation, a huge overhead is included in the system. The graph that shows total time required by the node count below shows that.

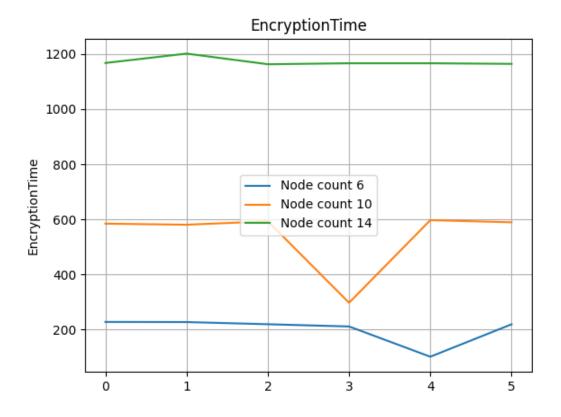


Figure 2 . Execution Time

# 6 Conclusion

The project is developed as a proof-of-concept work. That means only the graph data itself is encrypted, not the other variables. Therefore there are several security vulnerabilities. The helper variables that are not encrypted can be observed and the data becomes a subject to reverse engineering. Also the global state can be tracked to observe the current state of server. The final output is not also encrypted because of simplicity. To overcome such an important issue all of the variables can be encrypted alongside with the graph and necessary functions can be implemented in the client-server interface if can not be implemented in the server. Also such algorithms can be designed to eliminate the extra data transfer between server and the client.

# References

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- [9] (2022) Homomorphic encryption. [Online]. Available: https://en.wikipedia.org/wiki/Homomorphic\_encryption#Fourth-generation\_FHE
- [10] (2022) Kahn's algorithm for topological sorting. [Online]. Available: https://www.geeksforgeeks.org/topological-sorting-indegree-based-solution/

# Appendix A Result Graphs

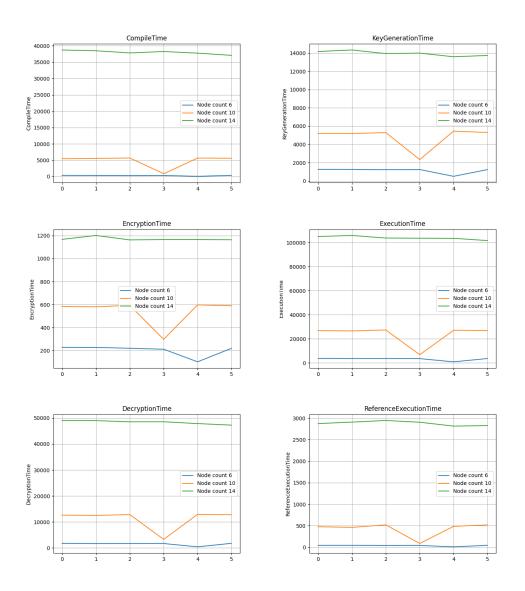


Figure 3 . Results

# Appendix B Docker File

```
#Download base image ubuntu 20.04
  FROM ubuntu:20.04
  # LABEL about the custom image
  LABEL maintainer="ertanon@gmail.com"
  LABEL version="0.1"
  LABEL description="This is a Microsoft EVA development environment image based on
      Ubuntu 20.04"
  # Disable Prompt During Packages Installation
  ARG DEBIAN_FRONTEND=noninteractive
  RUN apt update \
13
      && mkdir development \
      && cd development \
      && apt install -y python3 python3-dev python3-pip \
17
      && cd /usr/local/bin \
      && ln -s /usr/bin/python3 python \
1.8
19
      && pip3 --no-cache-dir install --upgrade pip \
      && apt-get install -y --no-install-recommends apt-utils build-essential sudo git
20
      cmake \
      && apt install -y libboost-all-dev libprotobuf-dev protobuf-compiler \
21
      && apt install -y clang \
      && update-alternatives --install /usr/bin/cc cc /usr/bin/clang 100 \
      && update-alternatives --install /usr/bin/c++ c++ /usr/bin/clang++ 100 \
24
      && pip3 install numpy
25
26
  WORKDIR /development
27
  RUN git clone -b v3.6.4 https://github.com/microsoft/SEAL.git \
29
      && cd SEAL \
      && cmake -S . -B build -DSEAL_THROW_ON_TRANSPARENT_CIPHERTEXT=ON \
      && cmake --build build \
32
      && cmake --install build
33
34
  RUN git clone https://github.com/microsoft/EVA.git \
35
      && cd EVA \
36
      && git submodule update --init \
37
38
      && cmake . \
39
      && make -j \
40
      && pip3 install -e ./python \
      && pip3 install -r examples/requirements.txt \
41
      && cd examples/
43
  RUN pip3 install adhoccomputing \setminus
44
      && pip3 install networkx
45
46
  COPY 519ProjectTemplate /development/519ProjectTemplate
```

# Appendix C Server-Client Code

```
from itertools import tee
  from termios import tcdrain
  from eva import EvaProgram, Input, Output, evaluate
  from eva.ckks import CKKSCompiler
  from eva.seal import generate_keys
  from eva.metric import valuation_mse
  import timeit
  import networkx as nx
  from random import random
10 from os import mkdir
11 import copy
  # Helper packages
  import matplotlib.pyplot as plt
  # Using networkx, generate a random directed acyclic graph
  def generateGraph(n, prob):
      G = nx.gnp_random_graph(n, prob, directed=True)
      nodes = [(u,v) \text{ for } (u,v) \text{ in G.edges}() \text{ if } u < v]
      DAG = nx.DiGraph(nodes)
20
22
      plt.tight_layout()
      nx.draw_networkx(DAG, arrows=True)
23
      plt.savefig(f"results/DAG_{n}.png", format="PNG")
      plt.clf()
26
      return DAG, nodes
27
28
29 # If there is an edge between two vertices its weight is 1 otherwise it is zero
30 # Two dimensional adjacency matrix is represented as a vector
  \# Assume there are n vertices
  # (i,j)th element of the adjacency matrix corresponds to (i*n + j)th element in the
      vector representations
  def serializeGraphZeroOne(GG,vec_size):
      n = len(GG.nodes())
      graphdict = {}
35
      g = []
36
      for row in range(n):
          for column in range(n):
38
39
               if GG.has_edge(row, column):
40
                   weight = 1
41
               else:
                   weight = 0
42
               g.append(weight)
               key = str(row)+'-'+str(column)
44
               graphdict[key] = [weight] # EVA requires str:listoffloat
45
46
      # EVA vector size has to be large, if the vector representation of the graph is
47
      smaller, fill the eva vector with zeros
      for i in range(vec_size - n*n):
48
49
           g.append(0.0)
50
      return g, graphdict
  # To display the generated graph
  def printGraph(graph,n):
      for row in range(n):
           for column in range(n):
               print("{:.5f}".format(graph[row*n+column]), end = '\t')
56
57
          print()
58
  # Eva requires special input, this function prepares the eva input
60 # Eva will then encrypt them
  def prepareInput(n, vec_size):
61
62
      input = {}
63
      GG, nodes = generateGraph(n, 0.5)
```

```
in_degree = [0] * n
65
       for (src, dest) in nodes:
66
            in_degree[dest] += 1
67
       for i in range(vec_size-n):
            in_degree.append(0)
70
71
       # graph is a list
72
       graph, graphdict = serializeGraphZeroOne(GG, vec_size)
73
       input['Graph'] = graph
74
75
       return input, in_degree
   global_state = "start"
   global_in_degree = []
   global_graph_size = 0
   global_vector_size = 0
   def graphanalticprogram(graph):
82
       global global_graph_size
83
       global global_vector_size
84
85
       global global_in_degree
       global global_state
86
87
       # Array for multiplication
88
89
       first_one = []
       for i in range(global_vector_size):
90
           first_one.append(0.0)
91
       first_one[0] = 1
92
93
94
       node_queue = []
95
       for i in range(global_graph_size):
            if global_in_degree[i] == 0:
96
                node_queue.append(i)
97
98
       top_order = []
99
100
       while node_queue:
            # Extract front of queue (or perform dequeue) and add it to topological order
102
            node_index = node_queue.pop(0)
104
            top_order.append(node_index)
105
106
            node_dest_block = graph << (node_index*global_graph_size) # Shifting</pre>
            for k in range(global_graph_size):
                current_dest = node_dest_block << k</pre>
108
                dest_exists_checker = current_dest * first_one # Multiplication
                global_state = "check_if_all_zero"
112
                yield global_state, dest_exists_checker # yielding!!!
114
                if global_state == "not_all_zero": # There is an edge from u to k
                    global_in_degree[k] = global_in_degree[k] - 1
                    if global_in_degree[k] == 0:
116
117
                        node_queue.append(k)
118
119
       yield "finished", top_order
120
# Do not change this
   # the parameter n can be passed in the call from simulate function
   class EvaProgramDriver(EvaProgram):
123
124
       def __init__(self, name, vec_size=4096, n=4):
            self.n = n
            super().__init__(name, vec_size)
       def __enter__(self):
128
129
            super().__enter__()
130
       def __exit__(self, exc_type, exc_value, traceback):
            \verb"super().__exit__(exc_type, exc_value, traceback)"
132
```

```
133
   # Repeat the experiments and show averages with confidence intervals
   # You can modify the input parameters
135
136 # n is the number of nodes in your graph
# If you require additional parameters, add them
   def simulate(n):
138
       global global_state
139
       global global_in_degree
140
       global global_graph_size
141
       global global_vector_size
142
143
       results = {
144
            "Iterations": [],
145
146
147
       m = 256
148
149
       print("Will start simulation for ", n)
150
       config = {}
       config['warn_vec_size'] = 'false'
       config['lazy_relinearize'] = 'true'
       config['rescaler'] = 'always'
       config['balance_reductions'] = 'true'
156
       inputs, global_in_degree = prepareInput(n, m)
       global_graph_size = n
       global_vector_size = m
       graphanaltic = EvaProgramDriver("graphanaltic", vec_size=m,n=n)
160
       with graphanaltic:
161
162
            total_start = timeit.default_timer()
163
            # Graph adjacency list is private information, therefore we will encrpyt it
164
            # In degree and discovered lists are computational helpers and they are not
       encrypted.
            graph = Input('Graph') # Encrypted input
167
            p = graphanalticprogram(graph)
            iter = 0
168
169
            prog = graphanaltic
            compiler = CKKSCompiler(config=config)
171
            while True:
                    global_state, reval = next(p)
                except StopIteration as e:
174
                    print("StopIteration exception occured. Breaking the loop...", e)
                if global_state == "finished":
178
                    print("Sorting is finished")
179
                    print("Topological sort:", reval)
180
                    break
181
182
                Output(f'ReturnedValue{iter}', reval)
183
184
                if global_state == "check_if_all_zero":
185
186
                    result = {}
                    prog.set_output_ranges(30)
187
                    prog.set_input_scales(30)
188
189
190
                    iter_start = timeit.default_timer()
                    compiled_multfunc, params, signature = compiler.compile(prog)
191
                    compiletime = (timeit.default_timer() - iter_start) * 1000.0 #ms
192
                    result["CompileTime"] = compiletime
194
                    iter_start = timeit.default_timer()
public_ctx, secret_ctx = generate_keys(params)
195
196
                    keygenerationtime = (timeit.default_timer() - iter_start) * 1000.0 #
197
       ms
                    result["KeyGenerationTime"] = keygenerationtime
198
```

```
199
200
                    iter_start = timeit.default_timer()
                    encInputs = public_ctx.encrypt(inputs, signature)
201
                    encryptiontime = (timeit.default_timer() - iter_start) * 1000.0 #ms
                    result["EncryptionTime"] = encryptiontime
203
204
                    iter_start = timeit.default_timer()
205
                    encOutputs = public_ctx.execute(compiled_multfunc, encInputs)
206
                    executiontime = (timeit.default_timer() - iter_start) * 1000.0 #ms
207
                    result["ExecutionTime"] = executiontime
208
209
                    iter_start = timeit.default_timer()
210
                    outputs = secret_ctx.decrypt(encOutputs, signature)
211
                    decryptiontime = (timeit.default_timer() - iter_start) * 1000.0 #ms
                    result["DecryptionTime"] = decryptiontime
213
214
                    iter_start = timeit.default_timer()
                    reference = evaluate(compiled_multfunc, inputs)
216
                    referenceexecutiontime = (timeit.default_timer() - iter_start) *
217
       1000.0 #ms
218
                    result["ReferenceExecutionTime"] = referenceexecutiontime
                    if [k for k in reference[f'ReturnedValue{iter}'][:n*n] if k!=0]: # If
220
        not all zeros
                         global_state = "not_all_zero"
221
                    else:
222
                        global_state = "all_zero"
223
224
                    mse = valuation_mse(outputs, reference) # since CKKS does approximate
225
        computations, this is an important measure that depicts the amount of error
                    result["Mse"] = mse
226
                    results ["Iterations"].append(result)
                iter += 1
230
       results["TotalTime"] = (timeit.default_timer() - total_start) * 1000.0 #ms
231
       return results
232
233
234
235
   if __name__ == "__main__":
236
       total results = []
       bm_keys = ["Sim#", "NodeCount", "TotalIteration", "CompileTime", "KeyGenerationTime
237
       ","EncryptionTime","ExecutionTime","DecryptionTime","ReferenceExecutionTime"]
238
       simcnt = 6 #The number of simulation runs, set it to 3 during development
239
       otherwise you will wait for a long time
240
241
       try:
           mkdir("results")
242
       except:
243
           pass
244
245
       resultfile = open("results/results.csv", "w") # Measurement results are collated
246
       in this file for you to plot later on resultfile.write(",".join(bm_keys) + "\n")
247
       resultfile.close()
248
249
250
       print("Simulation campaing started:")
       for nc in range(6, 18, 4): # Node counts for experimenting various graph sizes
251
           n = nc
252
253
           resultfile = open("results/results.csv", "a")
254
           # Dict to hold the data to plot
255
            total_res = {}
256
            total_res["NodeCount"] = n
257
            total_res["CompileTime"] = []
258
            total_res["KeyGenerationTime"] = []
259
           total_res["EncryptionTime"] = []
260
```

```
total_res["ExecutionTime"] = []
261
262
           total_res["DecryptionTime"] = []
           total_res["ReferenceExecutionTime"] = []
263
264
           for i in range(simcnt):
265
                total_res["CompileTime"].append(0)
266
                total_res["KeyGenerationTime"].append(0)
267
                total_res["EncryptionTime"].append(0)
268
                total_res["ExecutionTime"].append(0)
269
                total_res["DecryptionTime"].append(0)
271
                total_res["ReferenceExecutionTime"].append(0)
272
                #Call the simulator
273
                results = simulate(n)
                total_res["TotalTime"] = results["TotalTime"]
275
                iteration = 0
276
               for result in results["Iterations"]:
277
                    iteration += 1
278
279
                    total_res["CompileTime"][i] += result["CompileTime"]
                    total_res["KeyGenerationTime"][i] += result["KeyGenerationTime"]
280
281
                    total_res["EncryptionTime"][i] += result["EncryptionTime"]
                    total_res["ExecutionTime"][i] += result["ExecutionTime"]
282
                    total_res["DecryptionTime"][i] += result["DecryptionTime"]
283
                    total_res["ReferenceExecutionTime"][i] += result["
284
       ReferenceExecutionTime"]
                csv_res = str(i) + "," + str(n) + "," + str(len(results["Iterations"])) +
286
        "," + str(total_res["CompileTime"][i]) + "," + str(total_res["KeyGenerationTime
       "][i]) + "," + str(total_res["EncryptionTime"][i]) + "," + str(total_res["
       ExecutionTime"][i]) + "," + str(total_res["DecryptionTime"][i]) + "," + str(
       total_res["ReferenceExecutionTime"][i]) + "\n"
                resultfile.write(csv_res)
287
           total_results.append(copy.deepcopy(total_res))
289
           resultfile.close()
290
291
       for key in [k for k in bm_keys if k not in ["Sim#", "NodeCount", "TotalIteration",
292
        "Mse"]]:
           plt.cla()
293
294
           for res in total_results:
                if key in res:
295
                    plt.plot(list(range(simcnt)), res[key], label = f"Node count {res['
296
       NodeCount ']}")
           # plt.xlabel("PathLength")
297
           plt.ylabel(key)
298
           plt.title(key)
290
300
           plt.grid()
301
           plt.legend()
           plt.autoscale(enable=True, axis="y", tight=None)
302
           plt.savefig(f"results/{key}.png")
303
           # plt.show(block=False)
304
```