

0.1 Требования к программному обеспечению

Минимальные системные требования: РС с операционной системой Windows версии XP/Vista/7/8/10. Требуется устройства ввода: клавиатура, мышь.

0.2 Средства реализации

Для выполнения работы был выбран язык программирования Python ввиду его простоты.

0.3 Интерфейс

Интерфейс из себя представляет простую консоль, где пользователю ничего делать не нужно.

0.4 Листинг

В данном подразделе приведены листинги программ.

1. алгоритм Кнута - Мориса - Прата
2. алгоритм Бойера - Мура

lstname

```
1  \label{listing:kmp}
2  \caption{
3      — — — — — }
4      def prefix(s):
5          v = [0] * len(s)
6          for i in range(1, len(s)):
7              k = v[i - 1]
8              while k > 0 and s[k] != s[i]:
9                  k = v[k - 1]
10             if s[k] == s[i]:
11                 k = k + 1
12             v[i] = k
13         return v
14     def kmp(s, t):
15         index = -1
16         f = prefix(s)
17         k = 0
18         for i in range(len(t)):
19             while k > 0 and s[k] != t[i]:
20                 k = f[k - 1]
21             if s[k] == t[i]:
22                 k = k + 1
23             if k == len(s):
24                 index = i - len(s) + 1
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25         break
26     return index

```

lstname

```

1  def badCharHeuristic(string, size):
2      badChar = [-1] * 256
3
4      for i in range(size):
5          badChar[ord(string[i])] = i
6
7      return badChar
8
9
10 def search(txt, pat):
11     """
12     A pattern searching function that uses Bad
13     Character
14     Heuristic of Boyer Moore Algorithm
15     """
16     m = len(pat)
17     n = len(txt)
18
19     # create the bad character list by calling
20     # the preprocessing function badCharHeuristic()
21     # for given pattern
22     badChar = badCharHeuristic(pat, m)
23
24     # s is shift of the pattern with respect to text
25     s = 0
26     while s <= n - m:
27         j = m - 1
28
29         # Keep reducing index j of pattern while
30         # characters of pattern and text are matching
31         # at this shift s
32         while j >= 0 and pat[j] == txt[s + j]:
33             j -= 1
34
35         # If the pattern is present at current shift,
36         # then index j will become -1 after the
37         # above loop
38         if j < 0:
39             # print("Pattern occur at shift =
40             # {}".format(s))
41             return s # Return only first entry
42
43     """
44     Shift the pattern so that the next
45     character in text

```

```

42         aligns with the last occurrence of it in
43         pattern.
44         The condition  $s+m < n$  is necessary for
45         the case when
46         pattern occurs at the end of text
47
48          $s += (m - \text{badChar}[\text{ord}(\text{txt}[s + m])])$  if  $s$ 
49          $+ m < n$  else 1)
50         ,,,
51     else:
52         ,,,
53         Shift the pattern so that the bad
54         character in text
55         aligns with the last occurrence of it in
56         pattern. The
57         max function is used to make sure that
58         we get a positive
59         shift. We may get a negative shift if
60         the last occurrence
61         of bad character in pattern is on the
62         right side of the
63         current character.
64         ,,,
65     s += max(1, j - badChar[ord(txt[s + j])])
66
67     return -1

```