

# Public Key Infrastructures

Gianluca Dini

Dept. of Ingegneria dell'Informazione

University of Pisa

Email: [gianluca.dini@unipi.it](mailto:gianluca.dini@unipi.it)

Version: 15/04/2025

1

Public Key Infrastructures

## CERTIFICATES: HOW TO DEFEAT THE MAN-IN-THE-MIDDLE

apr. '25

PKIs

2

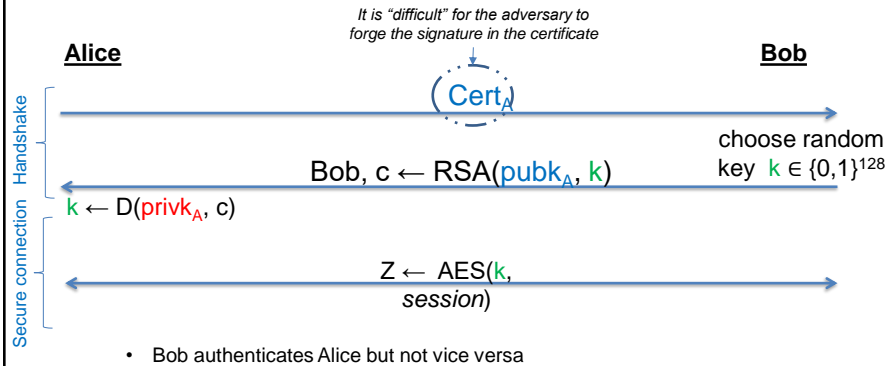
2

# Certificates and Certification Authority

- Certificate format (basic):
  - $Cert_A = Alice, pubK_A, L, \dots, S(priv_{CA}, Alice || pubK_A || L, \dots)$   
with  $L$  = validity period
- A certificate indissolubly binds the identity of a subject ( $Alice$ ) to his/her public key ( $pubK_A$ ); the binding is the digital signature of a Trusted Third Party called **Certification Authority** (CA)
  - In order to verify a certificate you need the CA's public key  $pubK_{CA}$

6

# Key Transport with Certificates



Disclaimer: oy protocol

7

# Diffie-Hellman with certificates

- $Cert_A = Alice, Y_A, L_A, S_{CA}(Alice || Y_A || L_A)$  with
- $Y_A = g^a \bmod p$  Alice's public key, and  $a$  Alice's private key
- $S_{CA}$  digital signature by certification authority CA

```
sequenceDiagram
    participant Alice as Alice  

    participant Bob as Bob  

    Note over Alice: a, Cert_A  

    Note over Bob: b, Cert_B  

    Alice->>Bob: Cert_A  

    Bob->>Alice: Cert_B  

    Note over Alice: K = (Y_B)^a mod p  

    Note over Bob: K = (Y_A)^b mod p  

    Alice->>Bob: Z ← AES(K, session)  

    Note over Alice, Bob: Secure connection
```

*Disclaimer: oy protocol*

apr-25

Public Key Cryptography

8

8

PKIs

# CERTIFICATION AUTHORITIES

apr. '25

PKIs

9

9

## CA's obligations [→]

- CA must be reliable
  - I. CA must verify that owner of  $(\text{privK}_A, \text{pubK}_A)$  pair is really entitled to use that name
    - CA establishes rules/policies to verify that a person has rights to the name
    - Identifying a subject is not easy; depends on country
  - CA must verify that the name (e.g., Alice) goes along with the key  $(\text{privK}_A)$

apr. '25

PKIs

10

10

## CA's obligations [↓]

- CA's certificate must be (immediately) available
  - CA's certificate is released at user registration time
  - CA's certificate is published in newspapers
  - CA's certificate is embedded in a browser installation package (is this secure?)
  - ...

apr. '25

PKIs

11

11

## Trust delegation

- Certification is based on trust delegation/transfer
  1. Bob trusts and delegates CA to verify Alice's identity and attest the authenticity of  $\text{pubK}_A$ ,
  2. Bob trusts the authenticity of CA's  $\text{pubK}_{CA}$   
consequently
- Through certificate  $\text{Cert}_A$  signed by CA, Bob acquires trust (believes) in the authenticity of  $\text{pubK}_A$

apr. '25

PKIs

12

12

## Important to remember

- What a certificate does
  - A certificate defines an indissoluble link between a subject's identifier and public key
- A certificate does not
  - specify the meaning of that link
  - the possible uses of that key
  - make any statement on the trustworthiness of the subject

apr. '25

PKIs

13

13

## Assurance [→]

- How much can I trust that the identifier actually corresponds to the legitimate owner of the key?
- CA Policies
  - Authentication policy
  - Issuance policy
  - These policies are public
    - A child-CA cannot have less restrictive policies
- Assurance is not quantifiable
  - Estimate according to the policy and the application rigor

apr. '25

PKIs

14

14

## Assurance [→]

- Specification, design and implementation contribute to the assurance
- Example: medicine
  - The process
    - A medicine is produced by a known and honorable pharmaceutical manufacturer
    - The medicine is delivered to chemists in a in sealed container
    - When the medicine is sold, the seal is still intact

apr. '25

PKIs

15

15

## Assurance [↓]

- Trust foundations
  - Ministry allows sale if the medicine passes certain tests and complies with certain clinical standards
  - Auditing committees verify that the production process satisfies industrial standards
  - Presence of the safety seal

apr. '25

PKIs

16

16

## In-house or external CA?

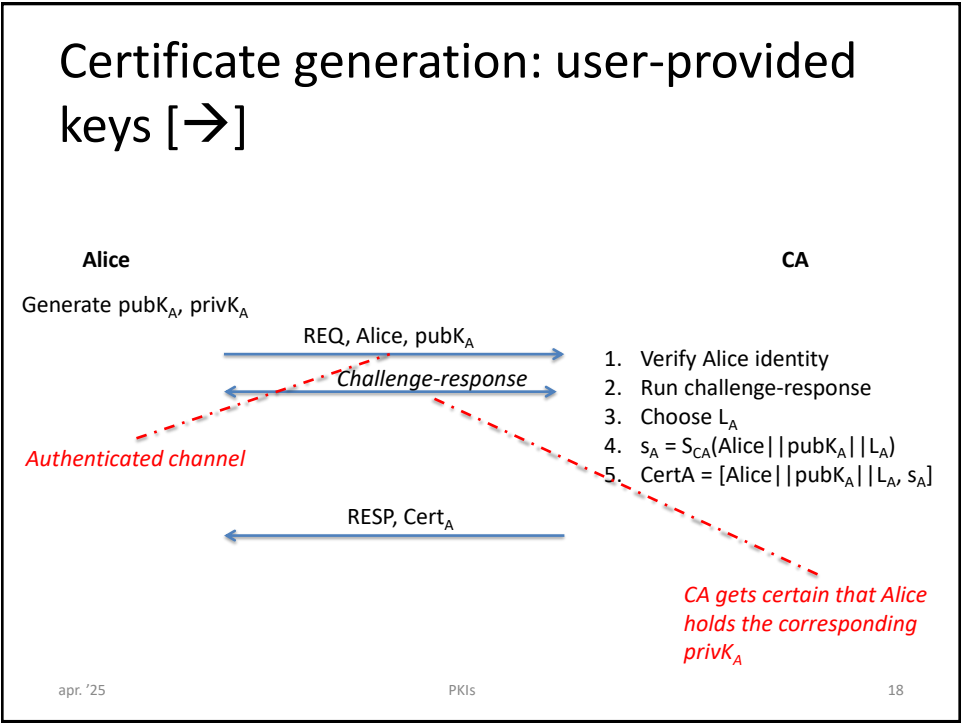
- Implement your own CA or exploit a commercial one?
  - Cost-convenience ratio
    - High quality certification → high costs
    - Low quality certification → high risks
  - In-house
    - Pros – Complete control of the certification process
    - Cons – Cost of the infrastructure; limited scale
  - Commercial
    - Pros – Large scale
    - Cons – Trust delegation; no liability

apr. '25

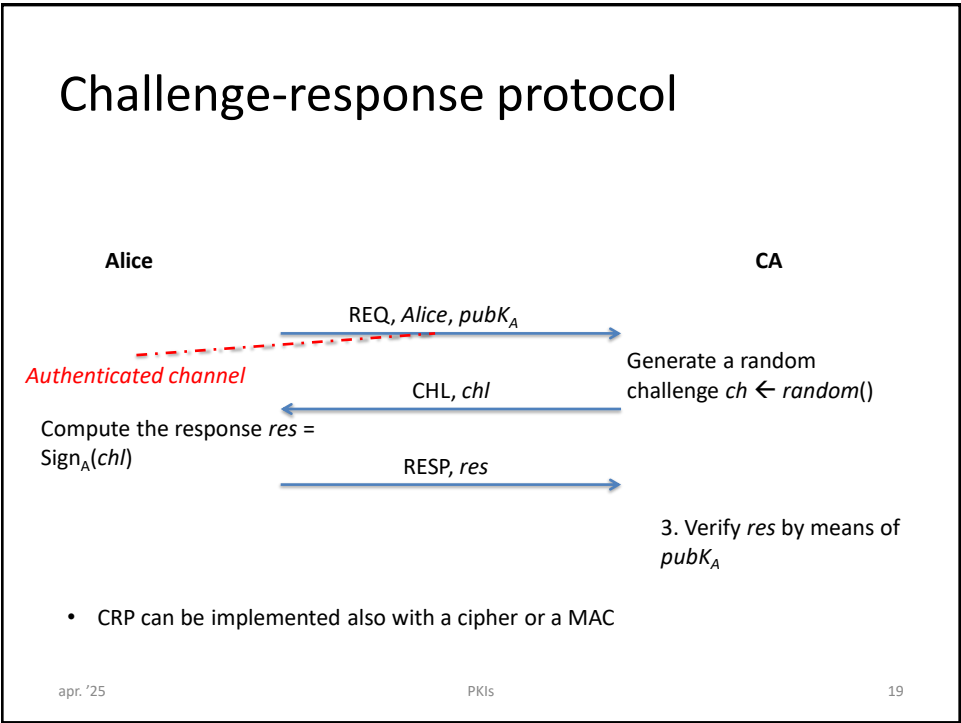
PKIs

17

17



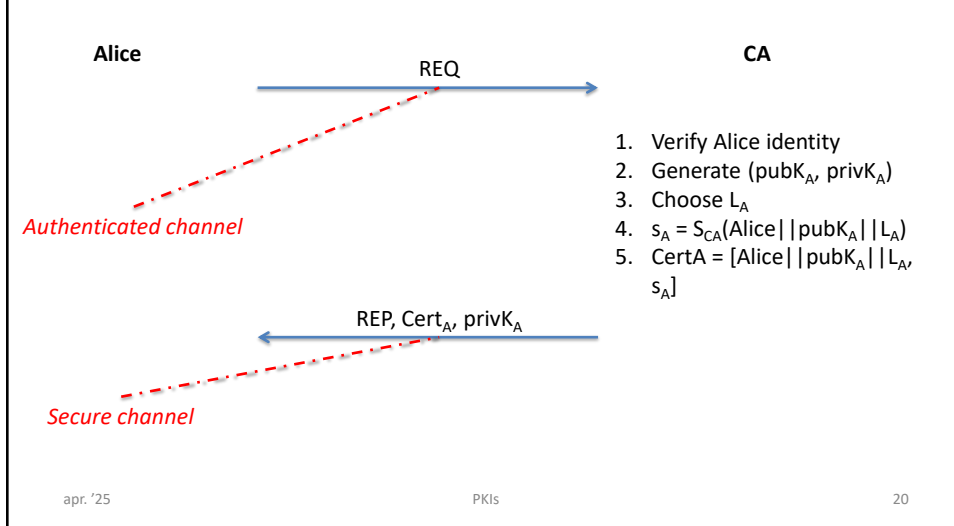
18



19



## Certificate generation: CA-generated keys



20

## On key generation at CA-side

- Fatal crypto flaw in some Taiwan government-certified smartcards makes forgery a snap ([www.arstechnica.com](http://www.arstechnica.com))
- Fatal flaw in the hw RNG
- Smartcards passed two international certifications (FIPS 140-2, Common Criteria)
- Research paper at [AsiaCrypt 2013](#)



apr. '25

PKIs

21

21

## Backup of private key [→]

Problems associated to private key:

- **Public key encryption** → backup of **decryption key**
  - Otherwise, encrypted data may become inaccessible
  - To be **able to decrypt even after key lifetime expiration**
  - Who makes backup matters
    - **Government backs up of citizen's privKs** → **privacy issues**
    - **Company backs up of employee's privKs** → **Encrypted data belong to the company**

apr. '25

PKIs

22

22

## Backup of private key [→]

- **Digital signature** → backup of **signing key**

→ still a private key but used for different purposes

- **Delete the key after key expiration**: **cannot be used anymore after expiration**
- Why to backup a signing key?
  - **Private key backup has adverse impact on non-repudiation**
  - However, **recovery from signing key loss is expensive in large scale apps** as **you must redistribute the pubK** (e.g., Microsoft, Adobe, Visa)
- How to backup
  - **Technological solution: Threshold crypto** (t out of n)

→ another cryptographic scheme: ①

- You want **different key pairs for encryption and signing**

apr. '25

PKIs

23

23 ① Evolution of secret splitting: you have **n shares** but just need **t/m shares** to find your secret.

# Threshold crypto (intuition)

Secret  $\Sigma$

$(t, n)$  secret sharing

$n$  shares

$\sigma_1 \sigma_2 \dots \sigma_n$

$n$  servers

## SECRET SHARING

- The secret  $\Sigma$  (e.g., private key) is split into  $n$  shares
- At least  $t$  shares are necessary to reconstruct the secret
- The system tolerates the compromise of  $t-1$  nodes

PKIs

24

24

PKIs

# EXPIRATION AND REVOCATION OF A CERTIFICATE

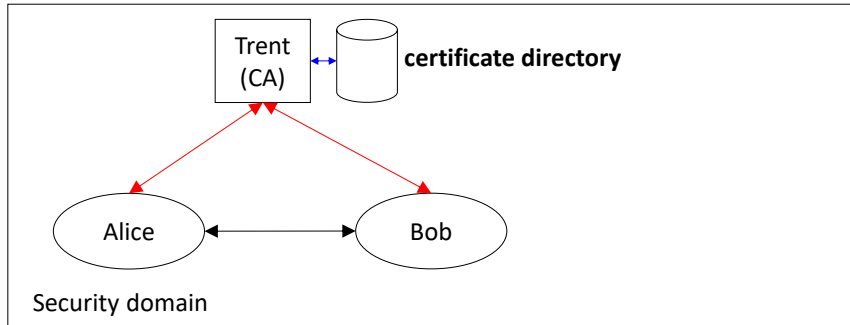
apr. '25

PKIs

25

25

## Single CA Model



- **Security domain** under control of the CA
- **Certificate directory** is a read-only database that stores certificates

apr. '25

PKIs

26

26

## Expired & revoked certificates

*Verifying if a certificate is valid is easy. Doing so to check revocation is not easy.*

- A certificate is **expired** if the validity period is expired
- If the **private key gets compromised before expiration**, then the certificate must be **revoked**
  - Examples: the private key has been revealed; the subject has changed role or left the organization
- **Certificate revocation must be**
  - **Correct:** revocation can be granted only to authorized parties, i.e., the owner or the issuer *Nobody else can do it on your behalf, except your CA for your*
  - **Timely:** revocation must be disseminated to all interested parties as soon as possible *↓ Issuer usually does so under push of authority*

apr. '25

PKIs

27

## How to verify a certificate

- **Bob's verification of Alice's Cert<sub>A</sub>**
  1. Bob obtains CA's public key  $\text{pubK}_{\text{CA}}$  [once at set-up]
  2. Bob verifies validity of CA's public key [once at set-up]
  3. Bob verifies the digital signature in Cert<sub>A</sub> by using  $\text{pubK}_{\text{CA}}$
  4. Bob verifies that Cert<sub>A</sub> is valid
  5. Bob verifies that Cert<sub>A</sub> is not revoked
  6. If all these checks are successful, then Bob accepts  $\text{pubK}_A$  as authentic Alice's key

apr. '25

PKIs

28

28

## Revocation methods *check for revocation*

*→ Periodically browser downloads list of revoked certificates*

- **Offline method: Certificate Revocation List (CRL)**
- **Online method: Online Certificate Status Protocol (OCSP)** *↳ on demand asking*

apr. '25

PKIs

29

29

# CRL

- A CRL is published periodically
- A revoked certificate lies in CRL until expiration
- Δ-CRL for efficiency

*You can find url of this list usually in the certificate itself.*

Timestamp states CRL freshness

serial number, revocation date, revocation reason

Digitally Signed by CA

*You don't put names for privacy reasons: if bro had certificate revoked for being fired/convicted, that's invasion of PKIs privacy.*

apr. '25

30 *if you didn't have timestamp attacker could disseminate old lists and abuse member revoked certificates.*

# OCSP

- Protocol sketch
  - Alice → OCSP: <OCSP RQST, Bob's cert serial nr.>
  - OCSP → Alice: <OCSP RESP, OK|KO><sub>OCSP</sub>
- Protocol Pros
  - Lighter and simpler than CRL protocol
  - Effective if the adversary is not a MIM
- Protocol Cons
  - In the clear → confidentiality issues
  - Exposed to replay attack (nonces are an extension ☹)
  - Browsers silently ignore OCSP if the query times out (→ MIM)

*X* [ ] *X*

apr. '25

PKIs

31

PKIs

THE X.509 STANDARD

apr. '25

PKIs

32

32

X.509 certificate format

Data structure containing 11 fields.

A data structure with several fields

1. Version (vers. 3 nowadays)

2. Serial number

3. Signature algorithm identifier<sup>①</sup>

4. Issuer distinguished name<sup>CA</sup>

5. Validity interval

6. Subject distinguished name<sup>Alice</sup>

7. Subject public key information

8. Issuer unique identifier (v=2,3)

9. Subject unique identifier (v=2,3)

10. Extensions (v=3)

11. Signature

②

① Many digital sig. algorithm, you know which is used for obj.org.

X.509 uses the Abstract Syntax Notation, ASN.1, (RFC 1422)

X.509 has been conceived for X.400 mail standard

X.509 uses Distinguished Names

apr. '25

PKIs

33

33

② We won't look in details. Fields that help you manage problems that may arise with naming system (i.e): an org. you might have same names.

## Distinguished names (X.500)

COUNTRY  
CO=IT

ORGANIZATION  
CO=IT, O=University of Pisa

ORGANIZATIONAL UNIT  
CO=IT, O=University of Pisa, OU=Dipartimento di Ingegneria della Informazione

COMMON NAME  
CO=IT, O=University of Pisa, OU=Dipartimento di Ingegneria della Informazione, CN=Gianluca Dini

How names are constructed:

hierarchical: Common name +  
Organizational Unit + Organization +  
Country

• There are other details beyond our interest

apr. '25

PKIs

34

34

## Example: <https://www.mps.it>

**Certificate name**  
[www.mps.it](https://www.mps.it)  
Consorzio Operativo Gruppo MPS  
Terms of use at [www.verisign.com/rpa](https://www.verisign.com/rpa) (c)00  
Florence  
Italy, IT

**Issuer**  
VeriSign Trust Network  
[www.verisign.com/CPS](https://www.verisign.com/CPS) Incorp.by Ref. LIABILITY LTD.(c)97 VeriSign

**Details**  
**Certificate version:** 3  
**Serial number:** 0x652D0F8ADAB4C7B168A27BBD1C3E9D9D  
**Not valid before:** Mar 2 00:00:00 2004 GMT  
**Not valid after:** Mar 2 23:59:59 2005 GMT  
**Fingerprint: (MD5)** CA CA 88 08 EC D0 8E 49 A6 9A 66 C4 69 31 E0 AE  
**Fingerprint: (SHA-1)** 82 64 CB 69 F0 43 86 43 FF B4 55 D4 25 EF 51 60 65 46 D3 87

*contd*

apr. '25

PKIs

35

35

Foundations of Cybersecurity

16



Example: https://www.mps.it

Public key algorithm: rsaEncryption

Public-Key (1024 bit):

Modulus:

00: E1 80 74 5E E7 E5 54 8B DF 6D 00 95 B5 96 27 AC

10: 66 93 E0 49 B9 6F 5B 73 53 1C BE 1C EB 47 64 B2

20: 12 95 70 E6 CD 50 67 02 88 E3 EE 9D B1 91 49 C8

30: 8D 58 19 4B 86 8F C0 2E 65 E8 F2 D4 82 CC 55 DB

40: 43 BC 66 DA 44 2F 53 B3 48 4B 37 15 F3 AB 67 C1

50: 69 B4 53 23 19 30 1A 19 23 7F 28 E0 E3 C0 6B 18

60: FF 84 C4 AC A9 74 28 DB FF E9 48 CA 75 D5 35 D6

70: 46 FB 7D D4 A7 3F A1 4B 00 60 14 DC D5 00 CF C7

Exponent:

01 00 01

Public key algorithm: sha1WithRSAEncryption

00: 23 A6 FE 90 E3 D9 BB 30 69 CF 43 2C FD 4B CF 67

10: D7 3C 46 22 9A 08 DB 05 1D 45 DC 07 F3 1E 4D 1F

20: 4B 11 23 5B 42 91 14 95 25 88 1F BD 60 E5 6F 84

30: 44 70 7A 95 EC 30 E4 46 4F 37 87 F1 B2 FA 45 04

40: 6F 7C BE 97 25 C7 20 E7 F3 90 55 51 99 3A 72 35

50: 40 F2 E8 E3 36 3A 7D 58 61 9C 91 D6 AC 34 E7 E8

60: 09 27 64 4F 2C 4C C2 D2 A3 32 DB 2B 7E F0 B6 F3

70: 69 96 E4 2B C3 2B 42 ED CA 2C 3C C8 F5 AA E6 71

contd

36

36

Example: https://www.mps.it

I cannot use this certificate to certify something else

Extensions:

X509v3 Basic Constraints: CA:FALSE

X509v3 Key Usage: Digital Signature, Key Encipherment

X509v3 CRL Distribution Points:

URI:http://crl.verisign.com/Class3InternationalServer.crl

X509v3 Certificate Policies:

Policy: 2.16.840.1.113733.1.7.23.3

CPS: https://www.verisign.com/rpa

X509v3 Extended Key Usage: Netscape Server Gated Crypto, Microsoft Server Gated Crypto, TLS Web Server Authentication, TLS Web Client Authentication

Authority Information Access:

OCSP - URI:http://ocsp.verisign.com

Unknown extension object ID 1 3 6 1 5 5 7 1 12:

0\_.][0Y0W0U..image/gif0!0.0...+.....k...j.H.,{..0%.#http://logo.verisign.com/vslogo.gif

37

37

PKIs

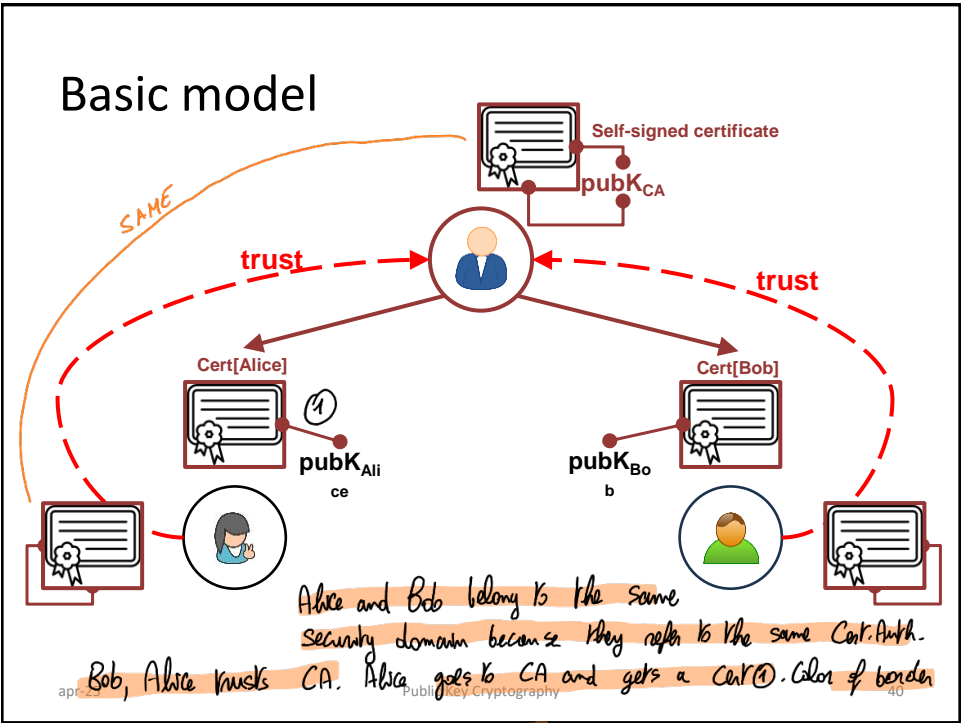
TRUST MODELS

apr. '25

PKIs

38

38



40

## Basic Model

- The Model
  - Every user trusts the root
  - The root releases certificates
- Inconvenient
  - Users have to go to the root in order to get the root self-signed certificate

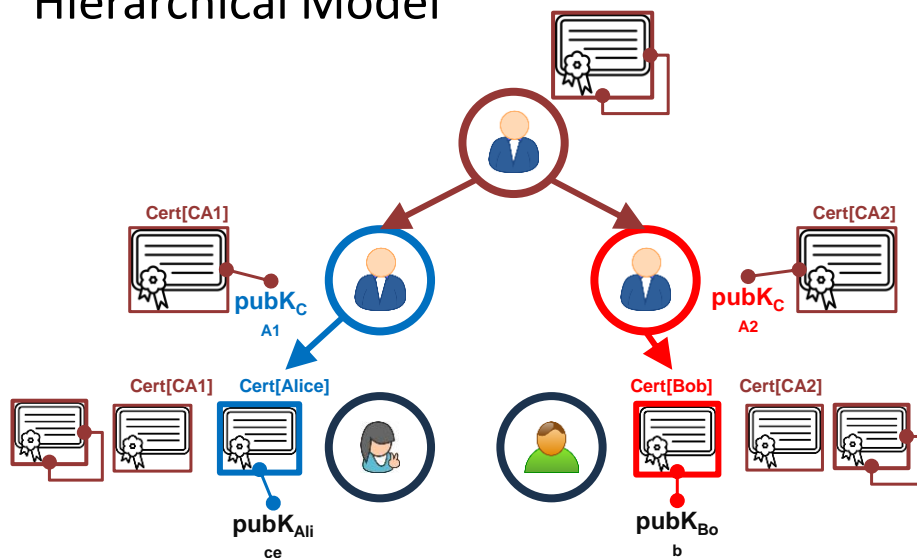
apr. '25

PKIs

41

41

## Hierarchical Model



apr-25

Public Key Cryptography

42

42 Variation of this: hierarchical model: Alice gets certified by the intermediate (subCA CA1, that signs Alice's Cert. Yr can do this because it is delegated by the root CA (root of trust). So Alice and Bob need to know root Cert. and delegates Cert.

So Alice needs to transfer her certificate and the certificate of CA1, that released her cert: Bob might not know it. Longer chains will give longer overhead

## Hierarchical model: constraints on the certification path

- If  $CA_x$  certifies  $CA_y$  (subordinate), the trust that  $CA_x$  has in  $CA_y$  transitively propagates to all CAs reachable from  $CA_y$ . You could use multiple levels ①
  - $CA_x$  may limit this propagation by posing constraints
    1. **on the chain length:** The chain after  $CA_y$  has a limited length [length=0 if you want to stop at previous examples]
    2. **on the set of domains:** CAs in the chain after  $CA_y$  must belong to a predefined set of CAs
- ① By doing this certification, the trust propagates downwards

apr. '25

PKIs

43

43

## Esempio: <https://www.mps.it>

### Certificate name

VeriSign Trust Network

[www.verisign.com/CPS](http://www.verisign.com/CPS) Incorpor. by Ref. LIABILITY LTD.(c)97 VeriSign

### Issuer

VeriSign, Inc.

Class 3 Public Primary Certification Authority

US

### Details

Certificate version: 3

Serial number: 0x254B8A853842CCE358F8C5DDAE226EA4

Not valid before: Apr 17 00:00:00 1997 GMT

Not valid after: Oct 24 23:59:59 2011 GMT

Fingerprint: (MD5) BC 0A 51 FA C0 F4 7F DC 62 1C D8 E1 15 43 4E CC

Fingerprint: (SHA-1) C2 F0 08 7D 01 E6 86 05 3A 4D 63 3E 7E 70 D4 EF 65 C2 CC 4F

apr. '25

PKIs

44

44

### Esempio: https://www.mps.it

**Public key algorithm:** rsaEncryption  
**Public-Key (1024 bit):**  
**Modulus:**  
00: 6F 7B B2 04 AB E7 34 4F 9C 53 A7 02 B2 90 4F 22  
10: F9 3A 3C 5A 8B 51 2B FE CB 42 95 30 70 FE 8A B2  
20: D3 1D C1 B8 5A 49 5C F7 39 4E 4D B7 F3 3B 09 F1  
30: FA E5 28 93 3E 30 F5 63 AA 43 71 27 56 FE A3 BB  
40: CA C4 6C 75 B2 32 C1 07 D9 DD 25 40 F5 5C A9 D4  
50: 15 0A 34 9A ED 42 97 EA BD F1 B2 55 45 73 3C AA  
60: E7 B6 5B 6C 4C F0 AA 3B 36 E6 BC D3 05 D4 BF E1  
70: 2B 65 A2 25 39 18 85 1F 7D 02 19 D6 E8 80 82 D8  
**Exponent:**  
01 00 01  
**Public key algorithm:** sha1WithRSAEncryption  
00: 08 01 EC E4 68 94 03 42 F1 73 F1 23 A2 3A DE E9  
10: F1 DA C6 54 C4 23 3E 86 EA CF 6A 3A 33 AB EA 9C  
20: 04 14 07 36 06 0B F9 88 6F D5 13 EE 29 2B C3 E4  
30: 72 8D 44 ED D1 AC 20 09 2D E1 F6 E1 19 05 38 B0  
40: 3D 0F 9F 7F F8 9E 02 DC 86 02 86 61 4E 26 5F 5E  
50: 9F 92 1E 0C 24 A4 F5 D0 70 13 CF 26 C3 43 3D 49  
60: 1D 9E 82 2E 52 5F BC 3E C6 66 29 01 8E 4E 92 2C  
70: BC 46 75 03 82 AC 73 E9 D9 7E 0B 67 EF 54 52 1A

apr. '25

PKIs

45

45

### Esempio: https://www.mps.it

**Extensions:**  
**X509v3 Basic Constraints:** CA:TRUE, pathlen:0  
**X509v3 Certificate Policies:**  
**Policy:** 2.16.840.1.113733.1.7.1.1  
**CPS:** https://www.verisign.com/CPS  
**X509v3 Extended Key Usage:** TLS Web Server Authentication, TLS Web Client Authentication, Netscape Server Gated Crypto, 2.16.840.1.113733.1.8.1  
**X509v3 Key Usage:** Certificate Sign, CRL Sign  
**Netscape Cert Type:** SSL CA, S/MIME CA  
**X509v3 CRL Distribution Points:**  
**URI:**http://crl.verisign.com/pca3.crl

apr. '25

PKIs

46

Certification Practice Statement

state how it issues certs, and it cannot be broken like the root CA.

46

# Web Model

Model implemented by browsers: assume there are many CAs in the world. When a user receives a certificate from a server, the browser needs self signed cert of the CA. Certificates are in the installation package. When you install browser, you automatically trust every CA. ①

may be hierarchical

Browser

Browser

For this reason, verifying installation package integrity is crucial

apr-25

Public Key Cryptography

47

47

# Examples

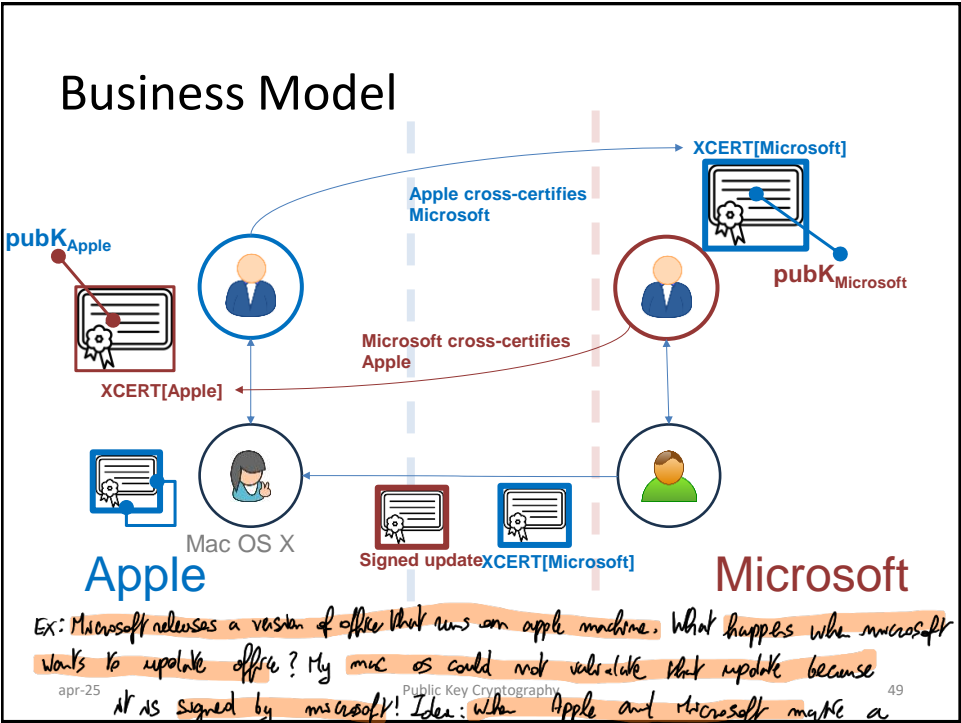
- Firefox
- Chrome
- Edge
- [www.unipi.it](http://www.unipi.it) ← GEANT Vereniging ← USERTrust (root)

apr-25

Public Key Cryptography

48

48



49 Contact to say that they support each other, they also cross certify each other. Apple certifies Microsoft, Microsoft certifies Apple. When you have update, you attach cross signed update

## Incidents

- March 2011 – Comodo
  - 9 fraudulent certs
- Summer 2011 – DigiNotar
  - 500+ fraudulent certs
  - [FOX-IT final report \(long\)](#)
  - [ENISA's resume \(short\)](#)
- January 2013 – Turktrust
  - 100+ fraudulent certs
  - [The TURKTRUST SSL certificate fiasco – what really happened, and what happens next?](#)