

# SKS (Secure Key Store)

## API and Architecture

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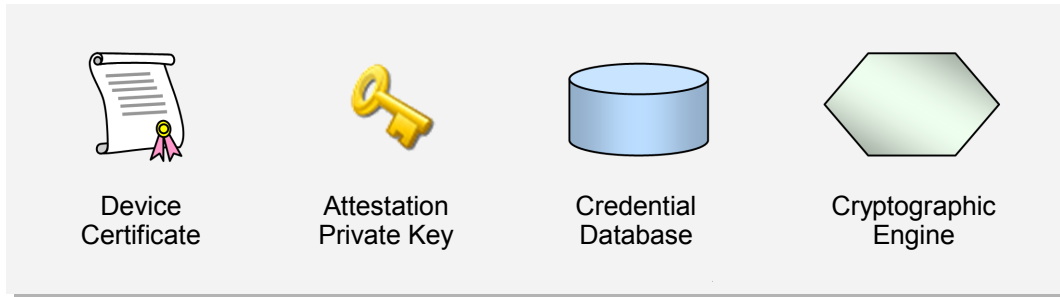
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# Introduction

This document describes the API (Application Programming Interface) and architecture of an electronic component called SKS (Secure Key Store). SKS is essentially an enhanced smart card that is optimized for on-line provisioning of cryptographic keys and associated attributes.

## Architecture

Below is a picture showing the components in the SKS architecture:



All operations inside of a SKS is supposed to be protected from tampering by malicious external entities but the degree of internal protection may vary depending on the environment that the SKS is running in. That is, an SKS housed in a smart card which may be inserted in an arbitrary computer must keep all data within its protected memory, while an SKS that is an integral part of a mobile phone processor may store credential data in the same external Flash where programs are stored, but sealed by an SKS-resident “master key”.

The *Device Certificate* and its associated *Attestation Private Key* form the foundation for the mechanism that secure provisioning of keys, which remains secure even if the surrounding middleware (for *self-contained* SKSes NB) and network are unsecured.

The *Cryptographic Engine* performs in addition to standard cryptographic operations on private and secret keys, the core of the provisioning operations which from an API point-of-view are considerably more complex than the former.

The *Credential Database* holds keys and other data that is related to keys such as protection and extension objects.

## Provisioning API

Although SKS may be regarded as a “component”, it actually comprises of three associated systems: The KeyGen2 protocol, the SKS architecture, and the provisioning API described in this document. *These items are tightly be matched* in order to create a secure and interoperable system. A question that arises is of course how compatible this scheme is with respect to current protocols, APIs, and smart cards. The simple answer is: NOT. The reason why SKS still makes sense is that none of the common protocols, APIs and smart cards support secure on-line provisioning to end-users because *the current generation of smart cards are almost exclusively personalized by more or less proprietary software used by specific card administrators or by automated production facilities*. It is evident that (at least) mobile phones need a scheme that is more consistent with the on-line paradigm since SIM-cards due to operator-bindings do not scale particularly well. “*On the internet anybody can be an operator of something*”.

## “Key User” API

In this document Key User API refers to operations that are used by security applications like SSL-client-certificate authentication, S/MIME, and Kerberos (PKINIT). The Key User API is not a core SKS facility but its implementation is anyway RECOMMENDED, particularly for SKSes that are featured in connectable containers such as smart cards since card middleware have proved to be a major stumbling block for wide-spread adoption of PKI cards for consumers. The described Key User API is fully mappable to the subset of CryptoAPI, PKCS#11, and JCE that most existing PKI-using applications rely on.

The Key User API does not use authenticated sessions like featured in TPM 1.2 because this is a *local security option*, while the provisioning API has its own self-contained (mandatory) authentication scheme.

If another Key User API is used the only requirement is that the things that are provisioned by the Provisioning API, are compatible.

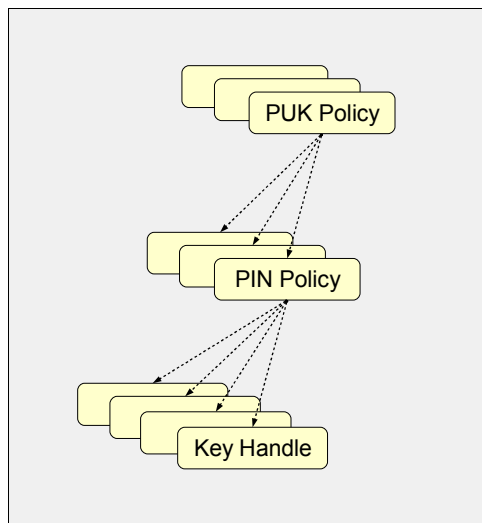
## Objects

The SKS API (as well as its companion protocol KeyGen2), assumes that objects are arranged in a specific fashion in order to work. At the core of the system are the cryptographic keys which are not only used for authentication, signing etc. but also for key life-cycle management and management of key attributes. All provisioned keys, included symmetric dittos, are identified and managed through an X.509 certificate. The reason for this arrangement is that this facilitates *universal object management* as well as supporting the POI (Proof Of Issuance) concept for enabling *secure remote object management by independent issuers*.

Note: unlike smart cards, SKS has no visible file system, only objects.

## Key Protection Objects

Keys may optionally be protected by PIN-codes. Each PIN-protected key have a separate error-counter, but a single PIN policy object may govern multiple keys. A PIN-policy and its associated keys are in turn governed by a PUK (Personal Unlock Key) policy object that can be used to reset error-counters that have passed the limit as defined by the PIN policy.



PIN and PUK policy objects are not directly addressable after provisioning; in order to read PIN and/or PUK policy data, you need to use an associated key handle as input.

The following XML extract shows a matching key generation (provisioning) request in KeyGen2:

```
<CreateObject>
  <PUKPolicy ID="PUK.1" Format="numeric" RetryLimit="3" EncryptedValue="mjRKrcuO ... 1O/e9mgMf3qw">
    <PINPolicy ID="PIN.1" Format="numeric" Grouping="shared" MaxLength="8" MinLength="4"
      PatternRestrictions="three-in-a-row sequence" RetryLimit="3">
      <KeyPair ID="Key.1" KeyUsage="encryption">
        <RSA KeySize="1024"/>
      </KeyPair>
      <KeyPair ID="Key.2" KeyUsage="authentication">
        <RSA KeySize="2048"/>
      </KeyPair>
    </PINPolicy>
  </PUKPolicy>
</CreateObject>
```

This sequence should be interpreted as a request for two RSA keys to be generated, protected by user-defined (within specified policy limits) PINs (the same for both keys), where the PINs are governed by an issuer-defined, and (protocol-wise) secret PUK.

## Key Data Objects

Provisioned keys always have an associated X.509 certificate, while other objects are optional.

## Data Types

The table below shows the data types used by the SKS API. Note that multi-byte integers are stored in big-endian fashion.

Type	Length	Comment
byte	1	Unsigned byte (0 - 0xFF)
bool	1	Byte containing 0x01 (true) or 0x00 (false)
short	2	Unsigned two-byte integer (0 - 0xFFFF)
int	4	Unsigned four-byte integer (0 - 0xFFFFFFFF)
byte[]	2 + length	Array of bytes with a leading "short" holding the length of the data
blob	4 + length	Long array of bytes with a leading "int" holding the length of the data

If an array is followed by a number in brackets ("byte[32]") it means that the array MUST be exactly of that length.

## Return Values

All methods return a single-byte status code. In case the status is  $\neq 0$  there is an error and any expected succeeding values MUST NOT be read as they are not supposed to be available. Instead there is a second return value containing an UTF-8 encoded description in English to be used for logging and debugging purposes as shown below:

Name	Type	Comment
Status	byte	Non-zero (error) value
ErrorMessage	byte[]	A human-readable error description

## Error Codes

The following table shows the standard SKS error-codes:

Name	Value	Comment
ERROR_AUTHENTICATION	1	This error is returned when there is something wrong with a supplied PIN-code. For more detailed information, see HHH and FFF
ERROR_STORAGE	2	There is no persistent storage available for the operation
ERROR_MAC	3	MAC does not match supplied data
ERROR_CRYPTO	4	Various cryptographic errors
ERROR_NO_SESSION	5	Session not found
ERROR_SESSION_VERIFY	6	The final step in the provisioning session failed to verify
ERROR_NO_KEY	7	Key not found
ERROR_ALGORITHM	8	Unknown or not fitting algorithm

## Method List

This section provides a list of all the SKS methods. The number in parenthesis holds the decimal value used to identify the method in a call. Method calls are formatted as strings of bytes where the first byte is the method ID and the succeeding bytes the applicable argument data.

## createProvisioningSession (1)

### Input

Name	Type	Comment
ServerSessionID	byte[32]	Server nonce value
ClientSessionID	byte[32]	Client nonce value
IssuerURI	byte[]	UTF-8 encoded URI identifying the issuer. In <a href="#">KeyGen2</a> this is the URL to which the result of this method is POSTed. The string MUST NOT exceed 1024 bytes
ServerPublicKey	byte[]	RSA server key (in X.509 DER format), for encrypting the session key (SK). The size of the key MUST match RSA capabilities
Updatable	bool	True if the session is supporting post provisioning updates
ClientOperationLimit	short	Constraint for thwarting cryptographic attacks on SK by limiting the number of externally visible SKS-generated signed and/or encrypted data objects
SessionLifeTime	int	Validity of the provisioning session in seconds

### Output

Name	Type	Comment
Status	byte	See <a href="#">Return Values</a>
EncryptedSessionKey	byte[]	Encrypted SK
SessionKeyAttest	byte[]	SK attestation signature
ProvisioningHandle	int	Local handle to created provisioning session

`createProvisioningSession` is the foundation for provisioning keys in an SKS. It performs the following steps in an *atomic* fashion:

- Generates a *random, secret* 32-byte SK (Session Key).
- Internally stores SK, ClientSessionID, ServerSessionID, IssuerURI, Updatable, ClientOperationLimit, current time + SessionLifeTime and returns a handle to the storage location in ProvisioningHandle.
- EncryptedSessionKey = **Encrypt** (ServerPublicKey, SK)
- SessionKeyAttest = **Sign** (AttestationPrivateKey,  
                                  **HMACSHA256** (SK, ClientSessionID ||  
  ServerSessionID ||  
  ServerPublicKey ||  
  IssuerURI ||  
  Updatable ||  
  ClientOperationLimit ||  
  SessionLifeTime))

The purpose of `createProvisioningSession` is creating a shared session key (SK) that is only known by the issuer and the SKS. In addition, the SKS is authenticated to the issuer. SK is used for *authenticating* and *encrypting* data that is exchanged between the issuer and the SKS during subsequent steps in the provisioning session.

If any succeeding operation associated with the provisioning session (through ProvisioningHandle), is regarded as incorrect by the SKS, *the session is immediately terminated and removed from the internal storage*.

### Notes

**Encrypt** uses the [RSAES-PKCS1-v1\\_5](#) algorithm.

**Sign** uses [DIAS](#) or [RSASSA-PKCS1-v1\\_5](#) signatures for RSA keys and [ECDSA](#) for EC keys with SHA256 as the hash function.

`ProvisioningHandle` is guaranteed to be unique and never reused.



## closeProvisioningSession (2)

### Input

Name	Type	Comment
ProvisioningHandle	int	Local handle to a provisioning session
GeneratedKeys	short	<i>Expected result.</i> What the issuer considers “has been ordered”
DeletedKeys	short	
ClonedKeys	short	
ReplacedKeys	short	
ExtensionObjects	short	
MAC	byte[32]	Vouches for the authenticity of the <i>expected result</i> parameters

### Output

Name	Type	Comment
Status	byte	See <a href="#">Return Values</a>
AttestedResponse	byte[32]	Attestation of the string “Success”

`closeProvisioningSession` terminates a provisioning session and returns a proof of successful operation to the issuer. However, success status will only be returned if *all* of the following conditions are valid:

- There is a valid provisioning session associated with `ProvisioningHandle`
- MAC matches the *expected result* parameters
- The *expected result* matches the SKS' internal calculations
- All generated keys are fully provisioned which means that matching public key certificates have been deployed.

When a provisioning session has been successfully closed by this method, it remains stored (until all keys associated with it have been deleted), but will only be a target for future updates if its `Updatable` flag is true.

## abortProvisioningSession (3)

### Input

Name	Type	Comment
ProvisioningHandle	int	Local handle to a provisioning session

### Output

Name	Type	Comment
Status	byte	See <a href="#">Return Values</a>

`abortProvisioningSession` is intended to be used by provisioning middleware if an unrecoverable error occurs in the communication with the issuer, or if a user cancels a session. If there is a matching and still *open* provisioning session, all associated data is removed from the SKS, otherwise an error is returned.

## createPUKPolicy (5)

### Input

Name	Type	Comment
ProvisioningHandle	int	Local handle to a provisioning session
ID	byte[]	PUK ID string with a maximum length of 32 bytes
EncryptedValue	byte[]	PUK value encrypted by the session key (SK)
Format	byte	Format of PUK strings. See format table
RetryLimit	byte	Number of incorrect PUK values ( <i>in a sequence</i> ), forcing the PUK object to permanently lock up. A zero value indicates that there is no limit but that the SKS will introduce an <i>internal</i> 1-10 second delay <i>before</i> acting on an unlock operation in order to thwart exhaustive attacks

### Output

Name	Type	Comment
Status	byte	See <a href="#">Return Values</a>
PUKPolicyHandle	int	Non-zero local handle to created PUK policy object

`createPUKPolicy` creates a PUK policy object that is meant to be referenced by the [createPINPolicy](#) method.

## createPINPolicy (6)

### Input

Name	Type	Comment
ProvisioningHandle	int	Local handle to a provisioning session
ID	byte[]	PIN ID string with a maximum length of 32 bytes
PUKPolicyHandle	int	Handle to a governing PUK policy object or zero
UserDefined	bool	True if PINs belonging to keys governed by the PIN policy are supposed to be set by the user or by the issuer. See KKKKKKKKKKKKKK
Format	byte	Format of PIN strings. See format table
RetryLimit	byte	Number of incorrect PIN values ( <i>in a sequence</i> ), forcing the key to lock up. A zero is not allowed as it would mean that the key would always be locked
Grouping	byte	
PatternRestrictions	byte	
MinLength	byte	
MaxLength	byte	

### Output

Name	Type	Comment
Status	byte	See <a href="#">Return Values</a>
PINPolicyHandle	int	Non-zero local handle to created PIN policy object

`createPINPolicy` creates a PIN policy object that is meant to be referenced by the [createKeyPair](#) method.

If `PUKPolicyHandle` is zero no PUK is associated with the PIN policy object.

A `PUKPolicyHandle` value of `0xFFFFFFFF` presumes that the target SKS supports a “device PUK”, *otherwise an error is returned*. The characteristics of a device PUK is out of scope for the SKS specification.

## createKeyPair (7)

### Input

Name	Type	Comment
ProvisioningHandle	int	Local handle to a provisioning session
ID	byte[]	Key ID string with a maximum length of 32 bytes
PINPolicyHandle	int	Handle to a governing PIN policy object or zero
PINValue	byte[]	PIN value must depending on PINPolicyHandle either be of zero length (the key is not PIN protected), contain a plain-text PIN value defined by the user, or constitute of an encrypted PIN set by the issuer
PrivateKeyBackup	bool	True if the generated private key is to be put in EncryptedPrivateKey for issuer backup
Migratable	bool	True if SKS should permit the key to be exported (in clear text) by the user
Updatable	bool	True if the key is subject to post provisioning updates. Note that this also requires the provisioning session's Updatable flag to be true, otherwise an error is returned
DeleteProtected	bool	True if the key needs a PUK in order to be deleted by a user. Note that this option requires that the key is associated with a PUK, otherwise an error is returned. A conforming SKS may ignore this flag since it does not introduce any vulnerabilities
ImportPrivateKey	bool	TBD
KeyUsage	byte	TBD
AlgorithmData	byte	TDB

### Output

Name	Type	Comment
Status	byte	See <a href="#">Return Values</a>
PublicKey	byte[]	Generated public key in X.509 DER representation
AttestedPublicKey	byte[]	Attestation of the authenticity of the public key and associated attributes
EncryptedPrivateKey	byte[]	<i>Optional.</i> This element will only be created if PrivateKeyBackup is true
KeyHandle	int	Local handle to created key pair object

`createKeyPair` creates an asymmetric key-pair inside of the SKS. To assure the issuer that it actually belongs to the SKS, the public key, together with attributes and the protection objects is signed (attested) by the SKS according to the following algorithm:

## References

KeyGen2

DIAS

RSAES-PKCS1-v1\_5

RSASSA-PKCS1-v1\_5

ECDSA

## Acknowledgments

There is a bunch of organizations, mailing-lists, and individuals that have been instrumental for the creation of SKS. I need to check who would accept to be mentioned :-)