Timus 1146

Tobby and Query

Prof. Edson Alves – UnB/FGA

Problema

Given a 2-dimensional array of positive and negative integers, find the sub-rectangle with the largest sum. The sum of a rectangle is the sum of all the elements in that rectangle. In this problem the sub-rectangle with the largest sum is referred to as the maximal sub-rectangle. A sub-rectangle is any contiguous sub-array of size 1×1 or greater located within the whole array.

As an example, the maximal sub-rectangle of the array:

is in the lower-left-hand corner and has the sum of 15.

1

Entrada e saída

Input

The input consists of an $N\times N$ array of integers. The input begins with a single positive integer N on a line by itself indicating the size of the square two dimensional array. This is followed by N^2 integers separated by white-space (newlines and spaces). These N^2 integers make up the array in row-major order (i.e., all numbers on the first row, left-to-right, then all numbers on the second row, left-to-right, etc.). N may be as large as 100. The numbers in the array will be in the range [-127, 127].

Output

The output is the sum of the maximal sub-rectangle.

Exemplo de entradas e saídas

Sample Input

4

0 -2 -7 0

9 2 -6 2

-4 1 -4 1

-1 8 0 -2

Sample Output

15

- Uma solução de força bruta computaria a soma todas as N^4 submatrizes, sendo que cada soma é feita em $O(N^2)$, de modo que a solução teria complexidade $O(N^6)$
- ullet Contudo, o uso de combinado de somas prefixadas e o algoritmo de Kadane permite identificar a submatriz de soma máxima com complexidade $O(N^3)$
- Para cada par de colunas (i,j), deve ser computado, por meio do algoritmo de Kadane nas somas $p_k(i,j)$, para $1 \le k \le N$, o intervalo de maior soma, onde

$$p_k(i,j) = \sum_{t=i}^{j} a_{kt}$$

 Veja que, dados os limites do problema, mesmo nos casos extremos a soma máxima ainda pode ser armazenada em variáveis inteiras

```
1 #include <hits/stdc++ h>
₃ using namespace std;
4 const int oo { 1000000010 };
6 int kadane(int N, const vector<int>& as)
7 {
      vector\langle int \rangle s(N + 1);
      s[1] = as[1]:
9
10
      for (size_t i = 2; i < as.size(); ++i)</pre>
          s[i] = max(as[i], s[i - 1] + as[i]);
      return *max_element(s.begin() + 1, s.end());
14
15 }
16
17 int solve(int N, const vector<vector<int>>& A)
18 {
19
      vector<vector<int>>> p(N + 1, vector<int>(N + 1, 0));
      int ans = -oo:
20
```

```
for (int i = 1; i \le N; ++i)
22
23
          vector<int> r(N + 1, 0);
24
25
          for (int j = i; j \le N; ++j)
26
               for (int k = 1; k \le N; ++k)
28
                   r[k] += A[k][i];
29
30
               ans = max(ans, kadane(N, r));
31
32
33
34
      return ans;
35
36 }
37
38 int main()
39 {
      ios::sync_with_stdio(false);
```

```
int N;
42
      cin >> N;
43
44
      vector<vector<int>>> A(N + 1, vector<int>(N + 1));
45
46
      for (int i = 1; i \le N; ++i)
47
          for (int j = 1; j \le N; ++j)
48
               cin >> A[i][j];
49
50
      auto ans = solve(N, A);
51
52
      cout << ans << endl:</pre>
53
54
      return 0;
55
56 }
```