Grafos

Fundamentos

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Faculdade UnB Gama

 \star Os grafos abstraem todas as outras estruturas de dados

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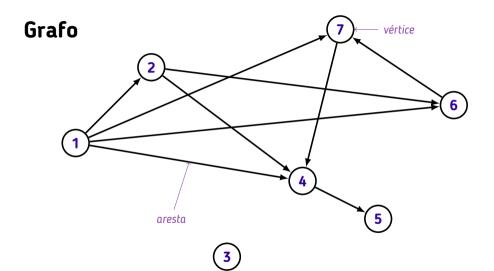
 \star Grafos modelam muitos problemas reais

- \star Os grafos abstraem todas as outras estruturas de dados
- * Grafos modelam muitos problemas reais
- * Travessias em grafos são eficientes e úteis

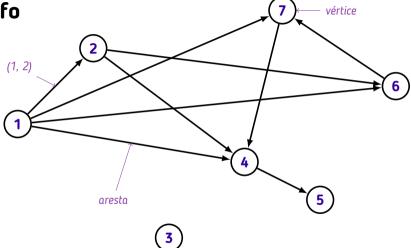
- \star Os grafos abstraem todas as outras estruturas de dados
- * Grafos modelam muitos problemas reais
- * Travessias em grafos são eficientes e úteis
- * Algoritmos clássicos resolvem problemas recorrentes

Grafo

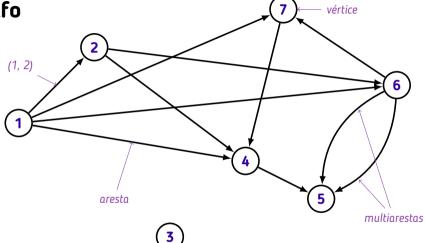
Grafo vértice

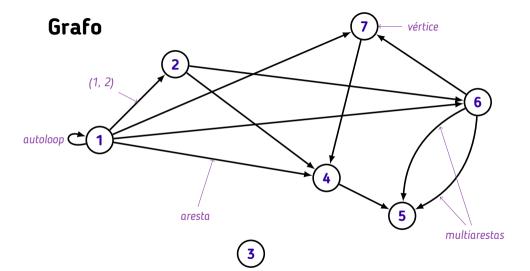


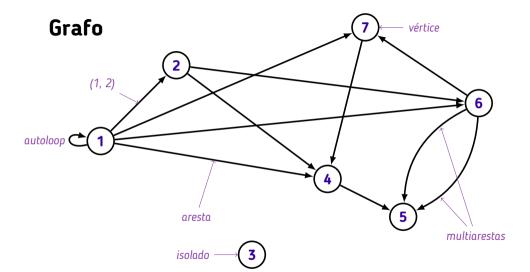
Grafo



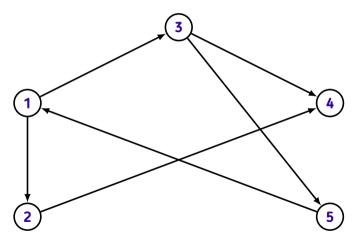
Grafo



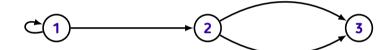




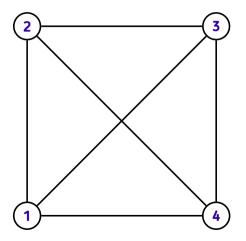
Grafo simples



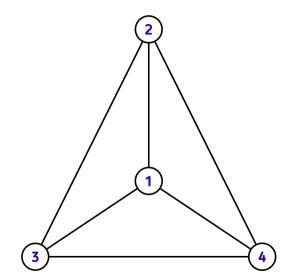
Multigrafo



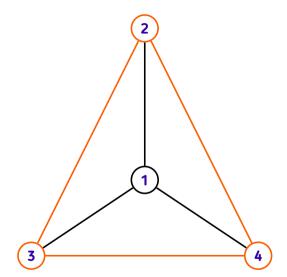
Grafo completo



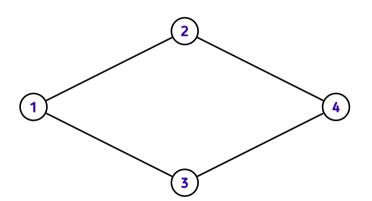
Subgrafo



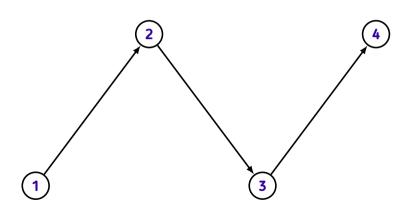
Subgrafo

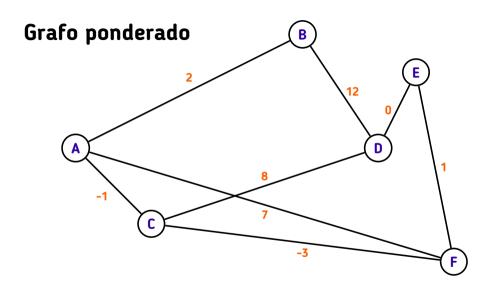


Grafo não-direcionado

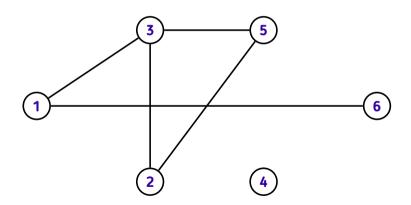


Grafo direcionado

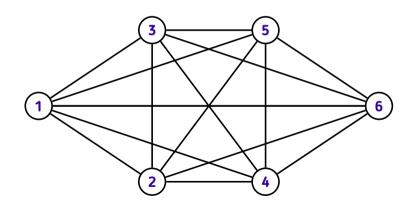




Grafo esparso



Grafo denso

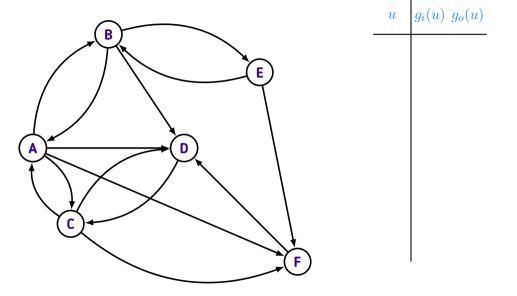


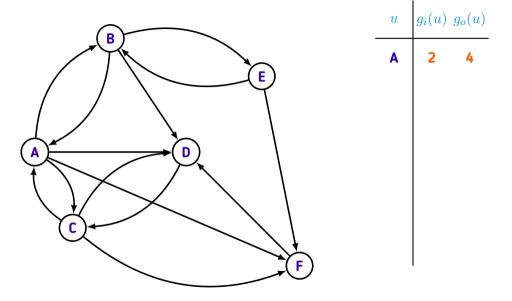
 \star Grau de entrada $g_i(u)$: # arestas que chegam em u

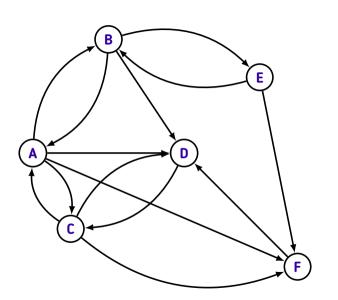
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 \star Grau de saída $g_o(u)$: # arestas que partem de u

- \star Grau de entrada $g_i(u)$: # arestas que chegam em u
- \star Grau de saída $g_o(u)$: # arestas que partem de u
- \star Se G é não-direcionado, então $g_i(u)=g_o(u)$, $orall u\in V$









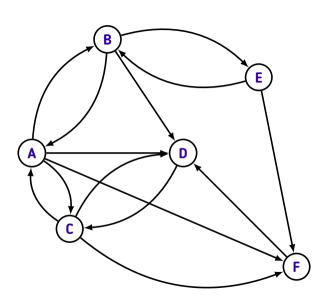
u

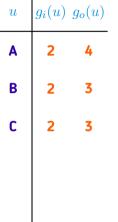
A

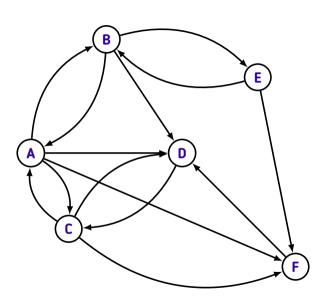
В

2 4

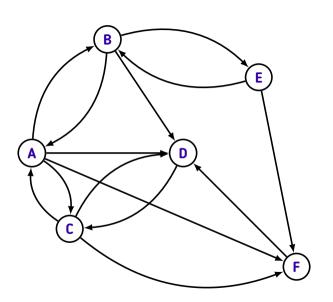
2



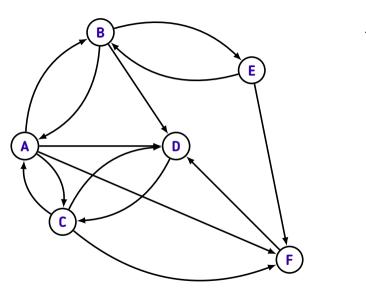




$g_i(u)$	$g_o(u)$
2	4
2	3
2	3
4	1
	2 2 2



u	$g_i(u)$	$g_o(u)$
A	2	4
В	2	3
С	2	3
D	4	1
E	1	2



u	$g_i(u)$	$g_o(u)$
A	2	4
В	2	3
С	2	3
D	4	1

Caminhos

Um caminho é uma sequência não-nula de vértices da forma

$$(u, w_1), (w_1, w_2), (w_2, w_3), \dots, (w_{n-1}, w_n), (w_n, v)$$

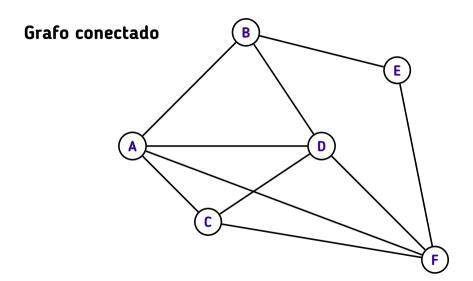
onde \boldsymbol{u} é o ponto de partida e \boldsymbol{v} o ponto de chegada

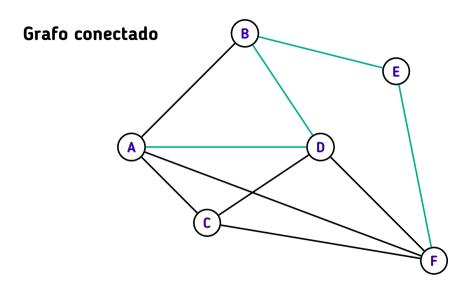
Caminhos

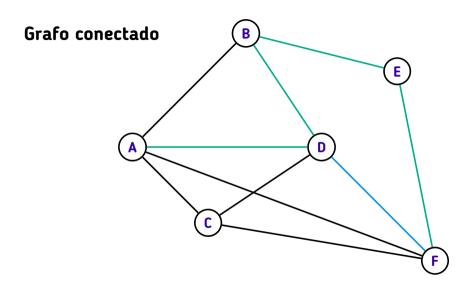
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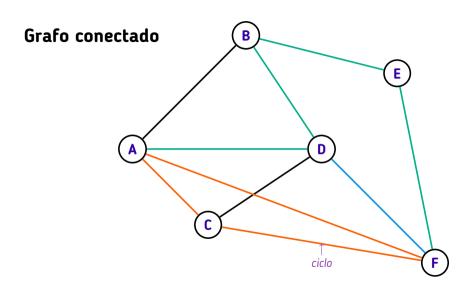
$$(u, w_1), (w_1, w_2), (w_2, w_3), \dots, (w_{n-1}, w_n), (w_n, v)$$

onde u é o ponto de partida e v o ponto de chegada









Referências

- 1. HALIM, Felix; HALIM, Steve. Competitive Programming 3, 2010.
- 2. LAAKSONEN, Antti. Competitive Programmer's Handbook, 2018.
- 3. SKIENA, Steven; REVILLA, Miguel. Programming Challenges, 2003.