Grafos

Ancestrais

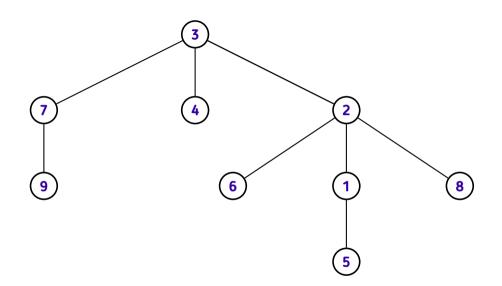
Prof. Edson Alves

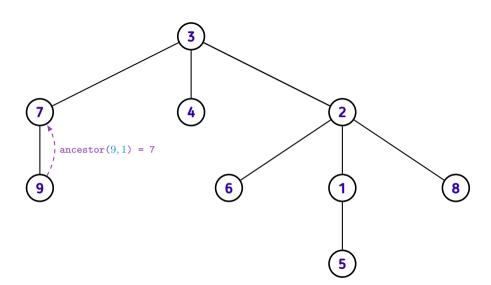
Faculdade UnB Gama

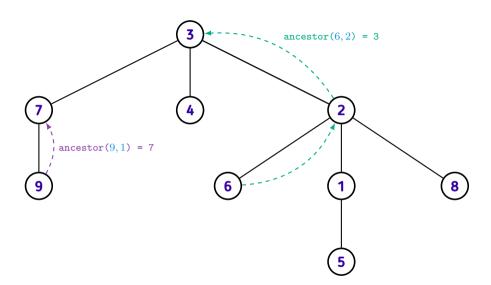
k-ésimo ancestral

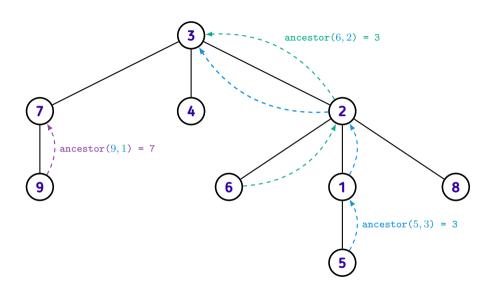
Seja T uma árvore enraizada, u um vértice de T e k um inteiro positivo.

O k-ésimo ancestral de u é o nó v que encerra o caminho que parte de u e segue k níveis, em direção à raiz. Notação: $v = \operatorname{ancestor}(u, k)$.









 \star É possível identificar o k-ésimo ancestral em $O(\log N)$, onde N é o número vértices da árvore, por meio de um algoritmo de programação dinâmica

 \star É possível identificar o k-ésimo ancestral em $O(\log N)$, onde N é o número vértices da árvore, por meio de um algoritmo de programação dinâmica

 \star Seja p(u) o ancestral direto de u na árvore enraizada

 \star É possível identificar o k-ésimo ancestral em $O(\log N)$, onde N é o número vértices da árvore, por meio de um algoritmo de programação dinâmica

 \star Seja p(u) o ancestral direto de u na árvore enraizada

 \star O caso base ocorre com $k=2^0=1$: ancestor(u,1) = p(u)

 \star É possível identificar o k-ésimo ancestral em $O(\log N)$, onde N é o número vértices da árvore, por meio de um algoritmo de programação dinâmica

- \star Seja p(u) o ancestral direto de u na árvore enraizada
- \star O caso base ocorre com $k=2^0=1$: ancestor(u,1) = p(u)
- * A transição é dada por:

$$\mathtt{ancestor}(u,2^k) = \mathtt{ancestor}(\mathtt{ancestor}(u,2^{k-1}),2^{k-1})$$

 \star Seja k um inteiro positivo

- \star Seja k um inteiro positivo
- \star É possível escrever k como a soma de potências distintas de 2:

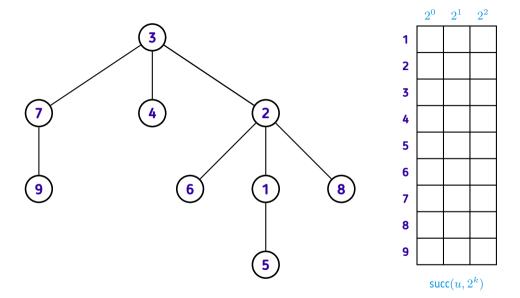
$$k = 2^{\alpha} + 2^{\beta} + \ldots + 2^{\omega}$$

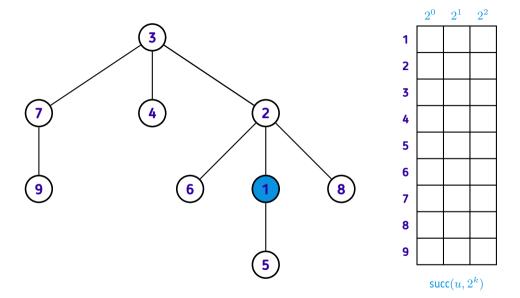
- \star Seja k um inteiro positivo
- \star É possível escrever k como a soma de potências distintas de 2:

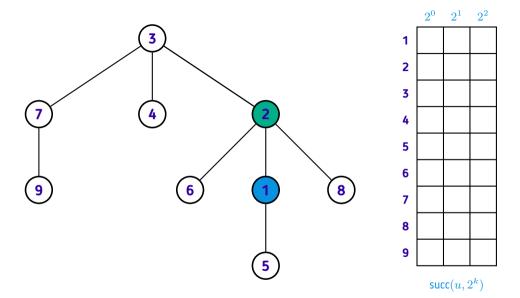
$$k = 2^{\alpha} + 2^{\beta} + \ldots + 2^{\omega}$$

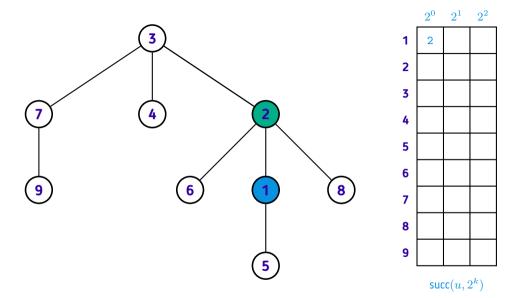
* Deste modo,

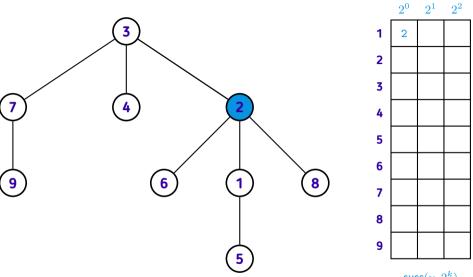
 $ancestor(u, k) = ancestor(ancestor(ancestor(ancestor(u, 2^{\alpha}), 2^{\beta}), \ldots), 2^{\omega})$



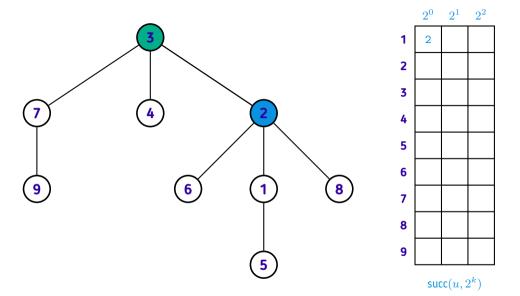


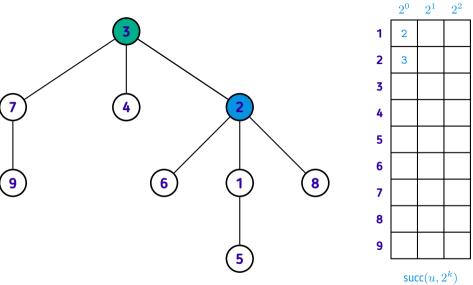


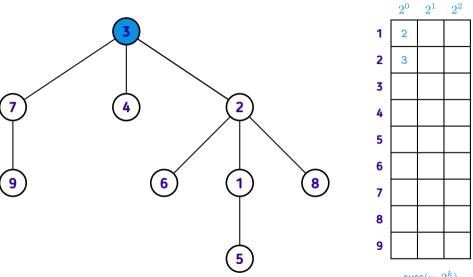




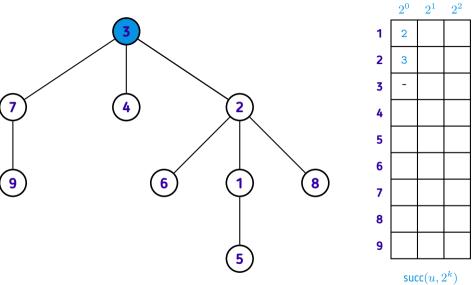
 $\mathsf{succ}(u, 2^k)$

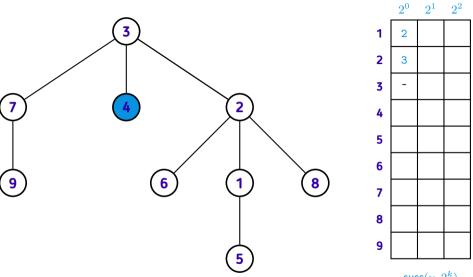




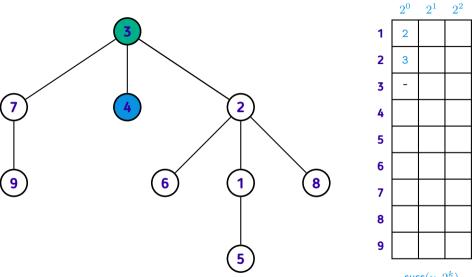


 $\mathsf{succ}(u, 2^k)$

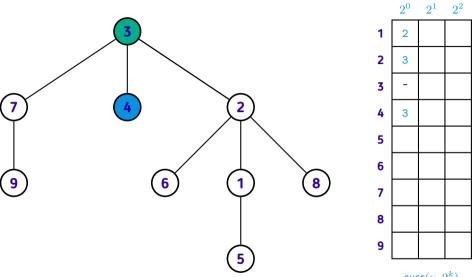




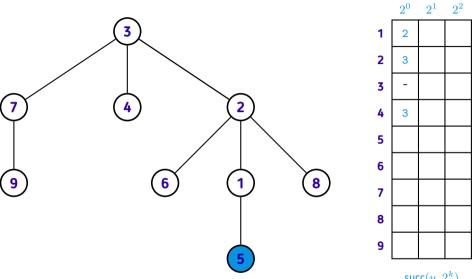
 $\mathsf{succ}(u, 2^k)$



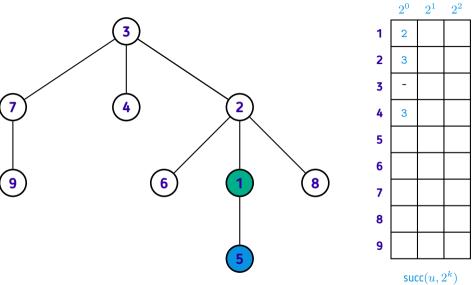
 $\mathsf{succ}(u, 2^k)$

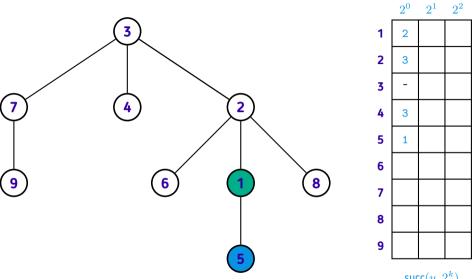


 $succ(u, 2^k)$

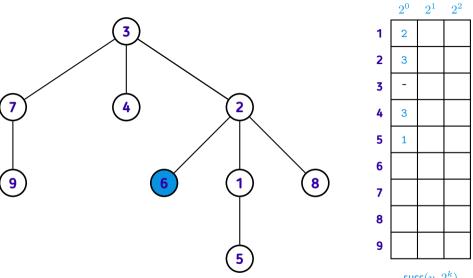


 $succ(u, 2^k)$

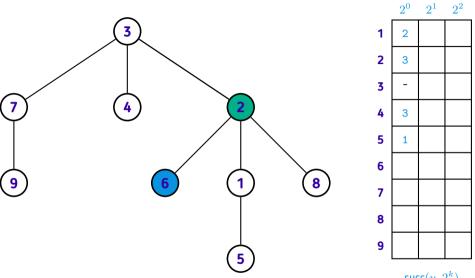




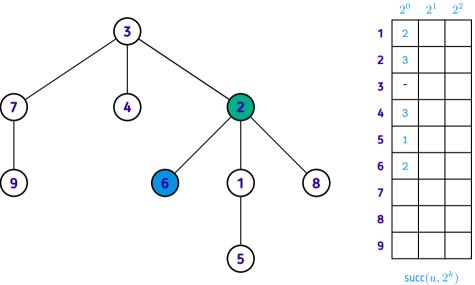
 $succ(u, 2^k)$

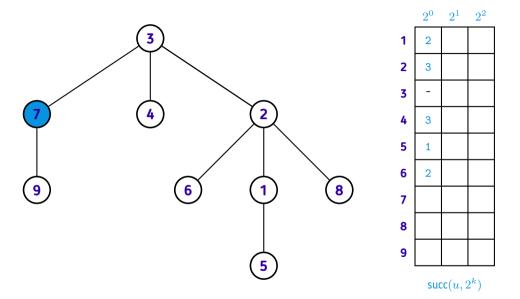


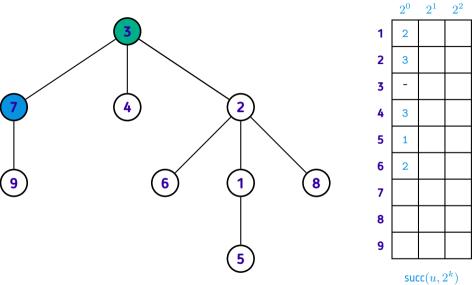
 $succ(u, 2^k)$

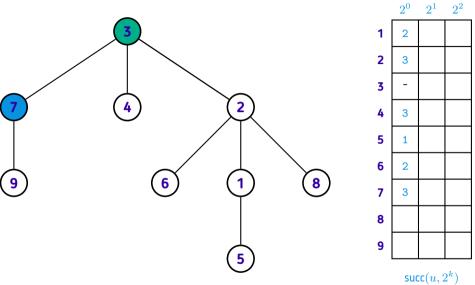


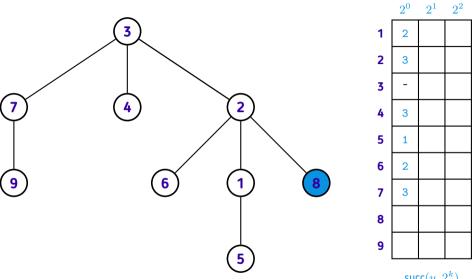
 $succ(u, 2^k)$



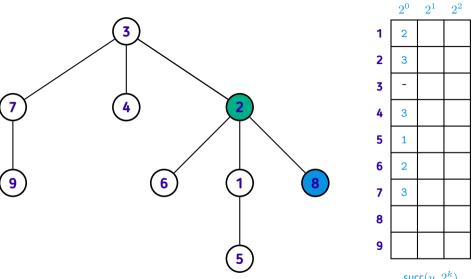




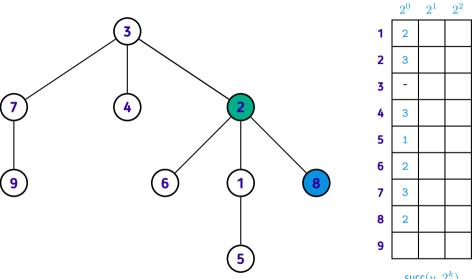




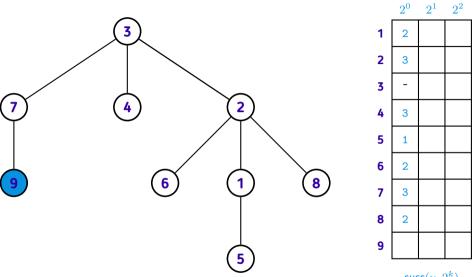
 $succ(u, 2^k)$



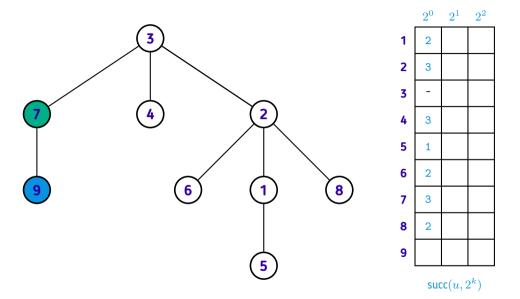
 $succ(u, 2^k)$

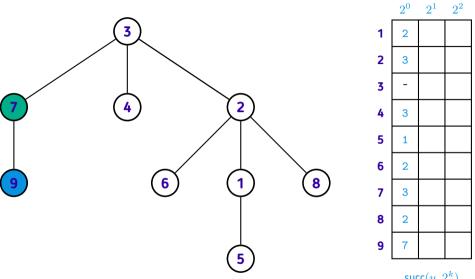


 $succ(u, 2^k)$

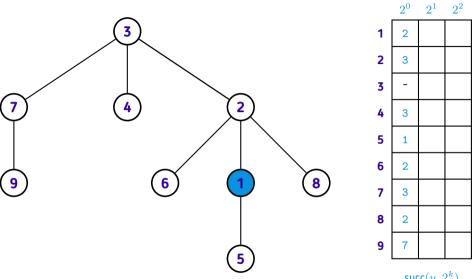


 $succ(u, 2^k)$

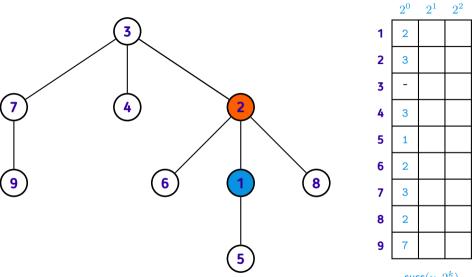




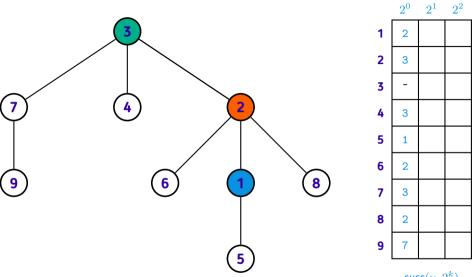
 $succ(u, 2^k)$



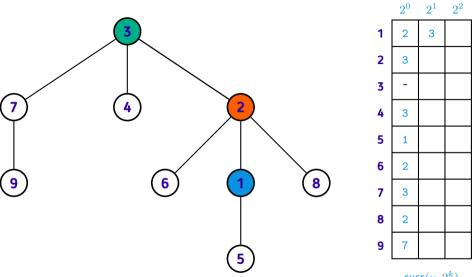
 $succ(u, 2^k)$



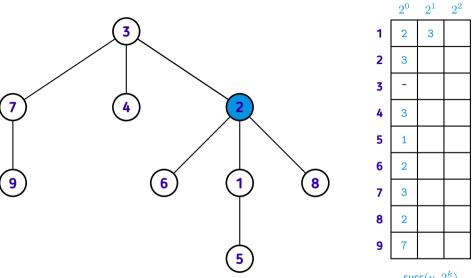
 $succ(u, 2^k)$



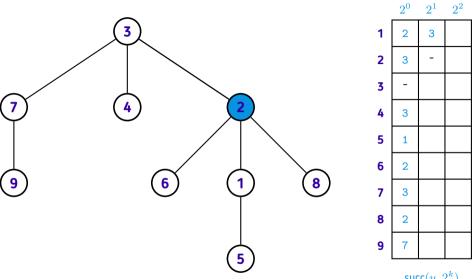
 $succ(u, 2^k)$



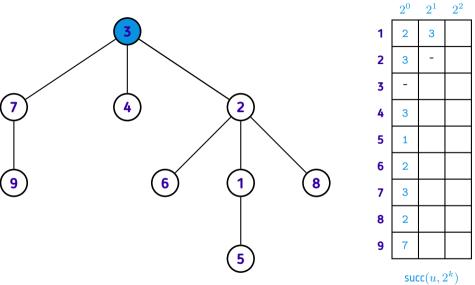
 $succ(u, 2^k)$

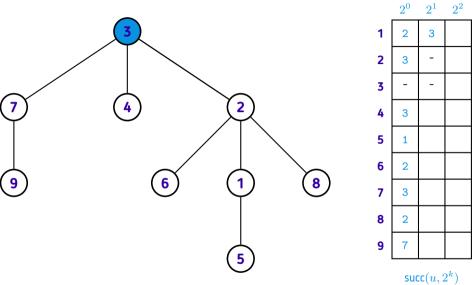


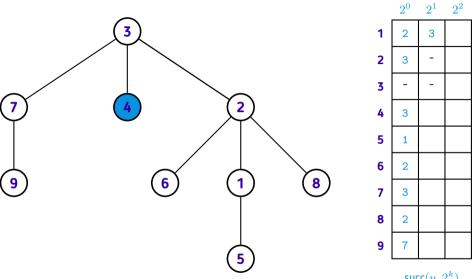
 $succ(u, 2^k)$



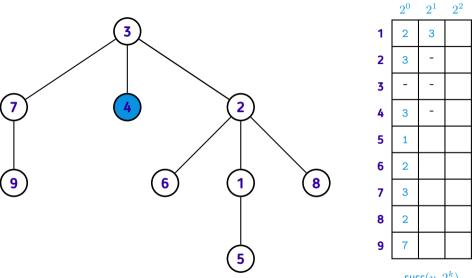
 $succ(u, 2^k)$



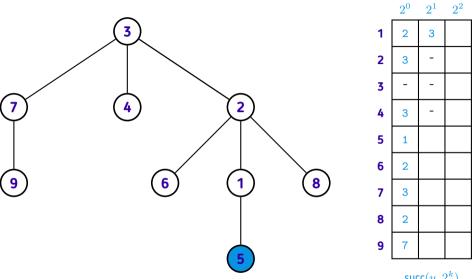




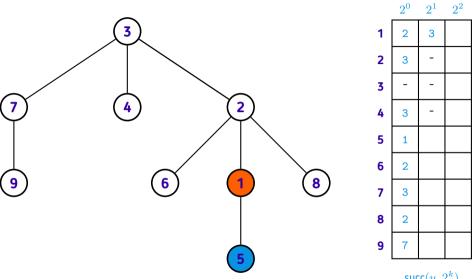
 $succ(u, 2^k)$



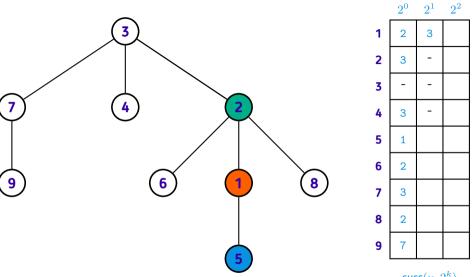
 $succ(u, 2^k)$



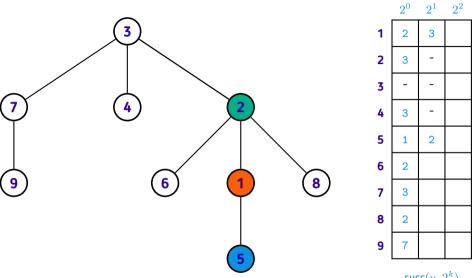
 $succ(u, 2^k)$



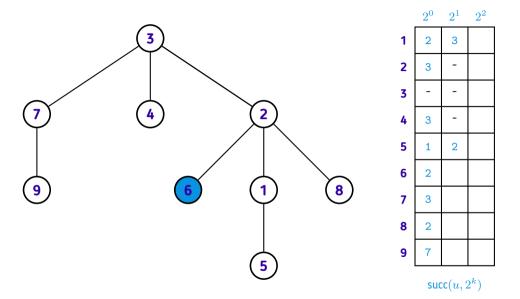
 $succ(u, 2^k)$

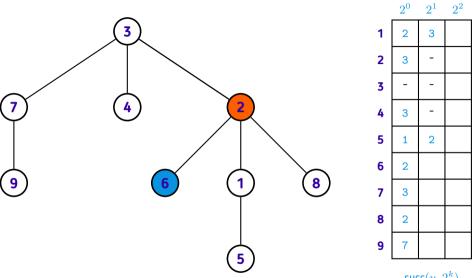


 $succ(u, 2^k)$

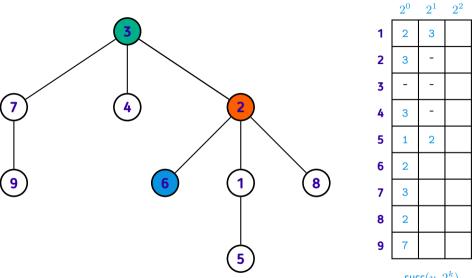


 $succ(u, 2^k)$

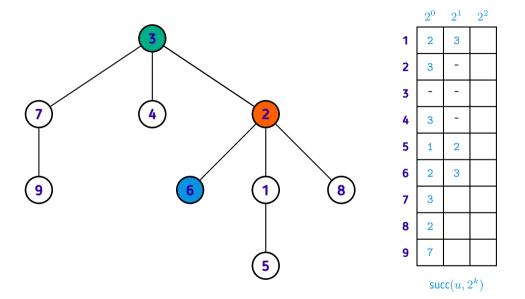


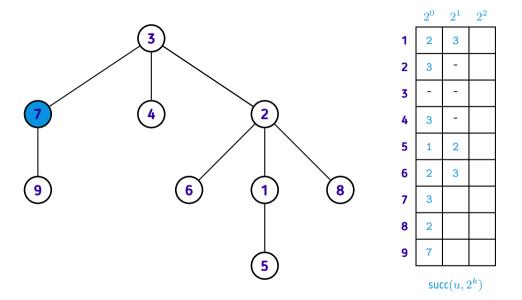


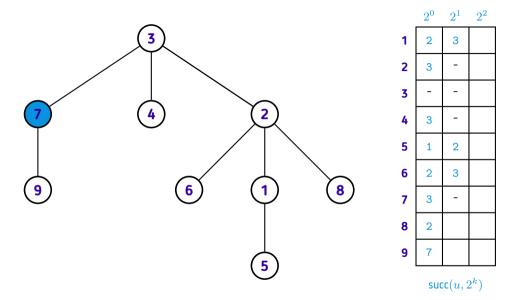
 $succ(u, 2^k)$

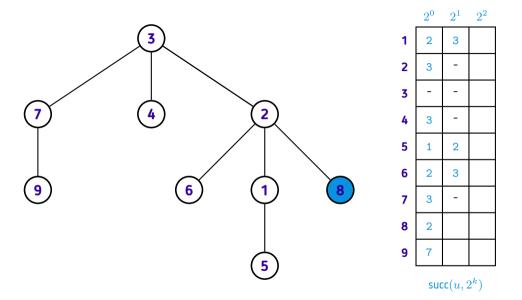


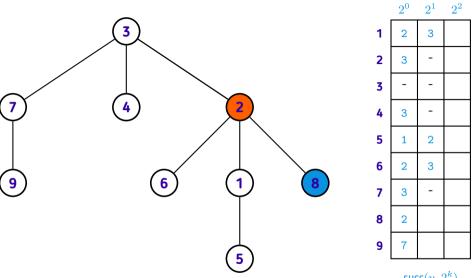
 $succ(u, 2^k)$



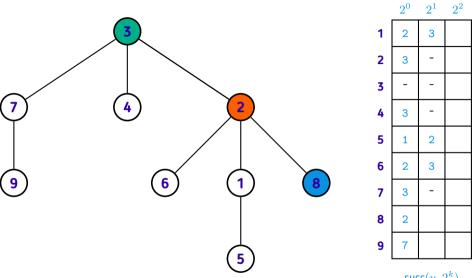




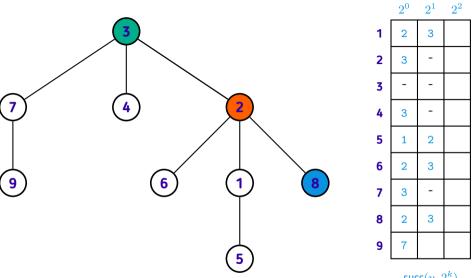




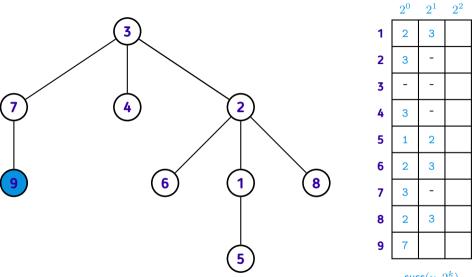
 $succ(u, 2^k)$



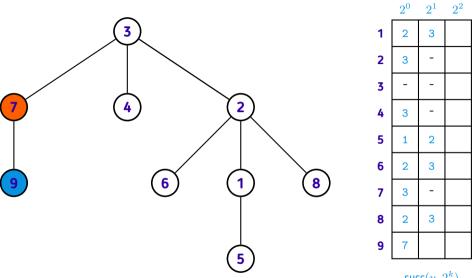
 $succ(u, 2^k)$



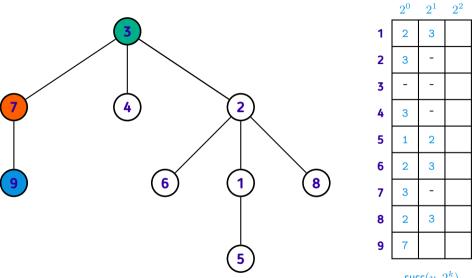
 $succ(u, 2^k)$



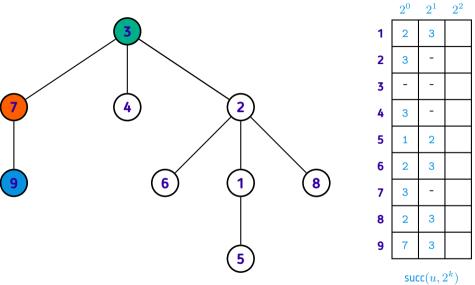
 $succ(u, 2^k)$

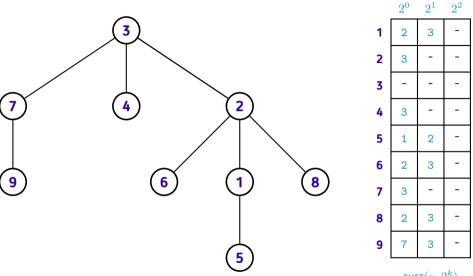


 $succ(u, 2^k)$



 $succ(u, 2^k)$





 $succ(u, 2^k)$

```
int ancestor(int u, int k, int N, const vector<vector<int>>& as)
   if (k >= N)
       return 0;
    int level = 0;
    while (k)
        if (k & 1)
            u = as[u][level];
       k >>= 1;
        ++level;
    return u;
```

```
auto precomp(int root, int N)
{
    int M = 0;
    while ((1 << (M + 1)) <= N)
        ++M:
    vector<vector<int>> as(N + 1, vector<int>(M + 1, 0));
    vector < int > ps(N + 1, 0);
    dfs(root, 0, ps);
    for (int u = 1; u \le N; ++u)
        as[u][0] = ps[u];
    for (int i = 1; i <= M; ++i)
        for (int u = 1; u \le N; ++u)
            as[u][i] = as[as[u][i - 1]][i - 1]:
    return as;
```

Referências

- 1. HALIM, Felix; HALIM, Steve. Competitive Programming 3, 2010.
- 2. LAAKSONEN, Antti. Competitive Programmer's Handbook, 2018.