## **Grafos**

**Fundamentos** 

**Prof. Edson Alves** 

Faculdade UnB Gama

 $\star$  Os grafos abstraem todas as outras estruturas de dados

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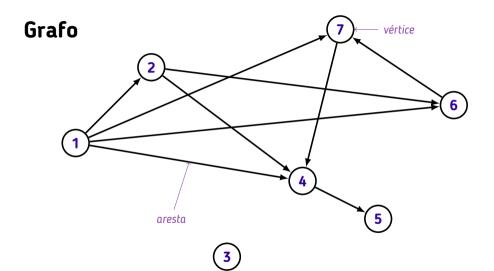
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- $\star$  Os grafos abstraem todas as outras estruturas de dados
- \* Grafos modelam muitos problemas reais
- \* Travessias em grafos são eficientes e úteis

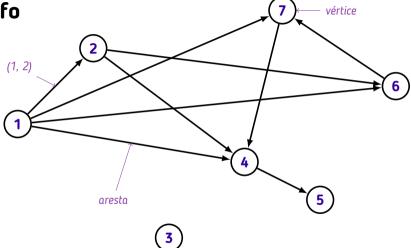
- $\star$  Os grafos abstraem todas as outras estruturas de dados
- \* Grafos modelam muitos problemas reais
- \* Travessias em grafos são eficientes e úteis
- \* Algoritmos clássicos resolvem problemas recorrentes

# Grafo

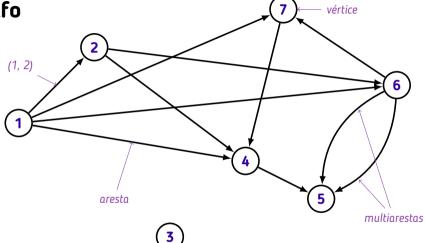
Grafo vértice

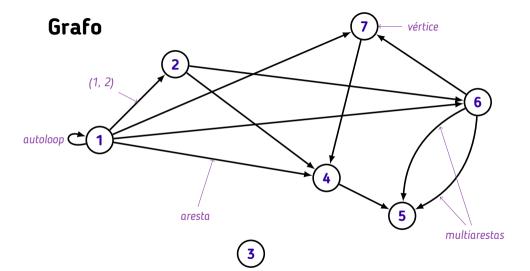


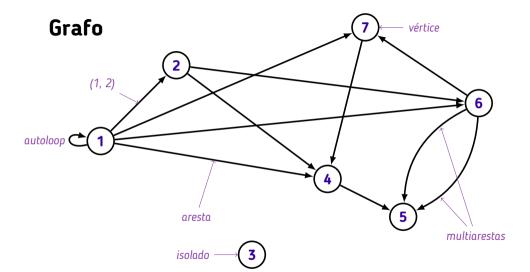
## Grafo



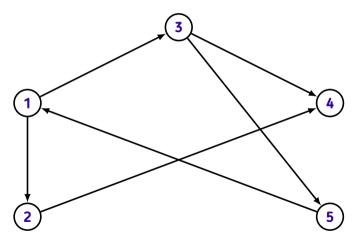
## Grafo



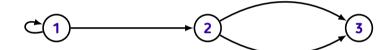




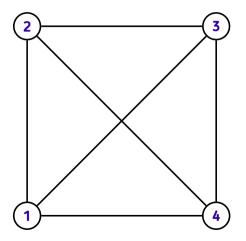
## **Grafo simples**



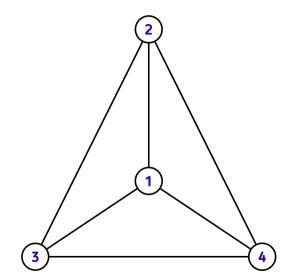
## Multigrafo



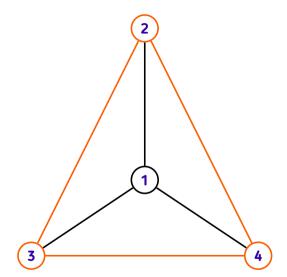
## Grafo completo



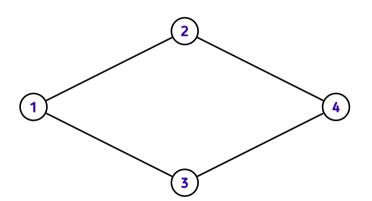
## Subgrafo



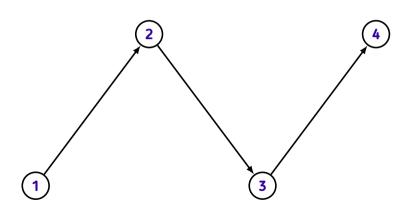
## Subgrafo

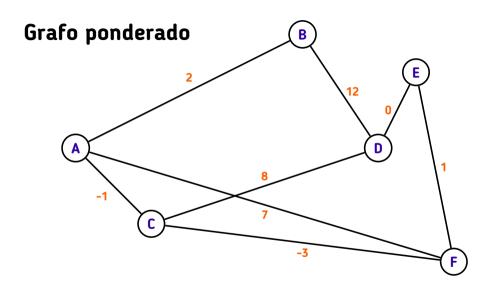


### Grafo não-direcionado

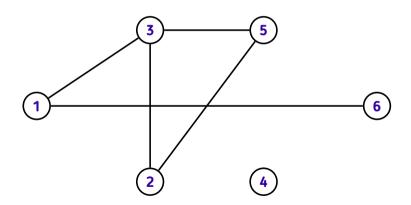


## Grafo direcionado

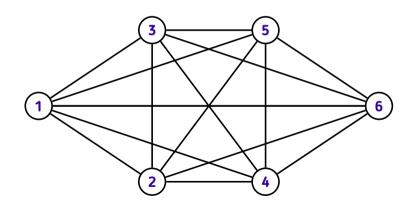




## Grafo esparso



## Grafo denso

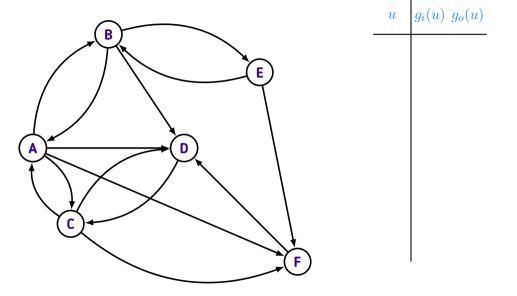


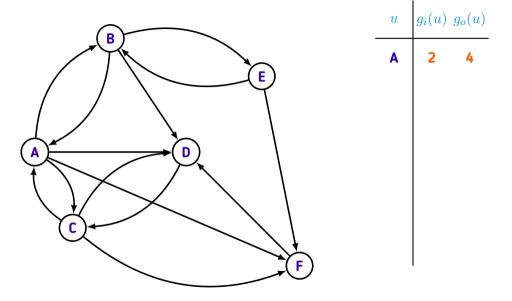
 $\star$  Grau de entrada  $g_i(u)$ : # arestas que chegam em u

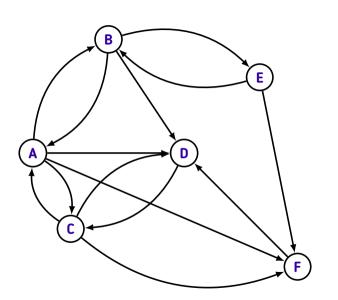
 $\star$  Grau de entrada  $g_i(u)$ : # arestas que chegam em u

 $\star$  Grau de saída  $g_o(u)$ : # arestas que partem de u

- $\star$  Grau de entrada  $g_i(u)$ : # arestas que chegam em u
- $\star$  Grau de saída  $g_o(u)$ : # arestas que partem de u
- $\star$  Se G é não-direcionado, então  $g_i(u)=g_o(u)$ ,  $orall u\in V$









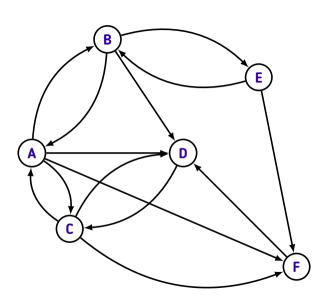
u

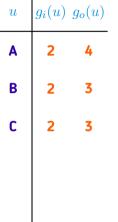
A

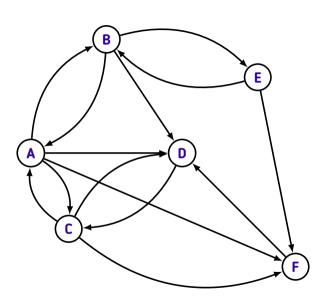
В

2 4

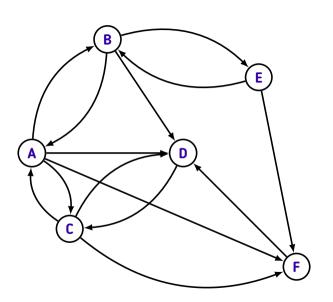
2



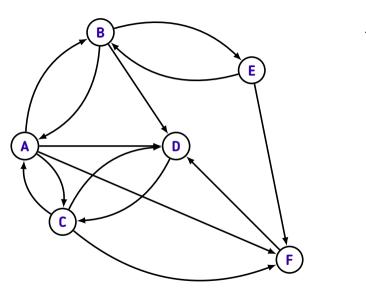




$g_i(u)$	$g_o(u)$
2	4
2	3
2	3
4	1
	2 2 2



u	$g_i(u)$	$g_o(u)$
A	2	4
В	2	3
С	2	3
D	4	1
E	1	2



u	$g_i(u)$	$g_o(u)$
A	2	4
В	2	3
С	2	3
D	4	1

#### **Caminhos**

Um caminho é uma sequência não-nula de arestas da forma

$$(u, w_1), (w_1, w_2), (w_2, w_3), \dots, (w_{n-1}, w_n), (w_n, v)$$

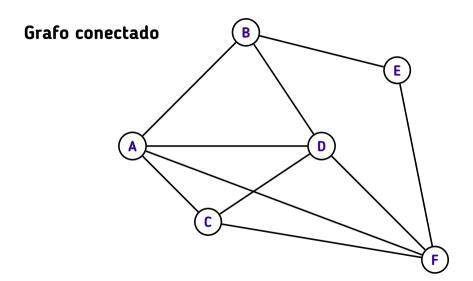
onde  $\boldsymbol{u}$  é o ponto de partida e  $\boldsymbol{v}$  o ponto de chegada

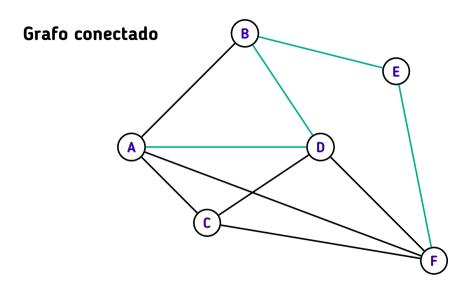
#### **Caminhos**

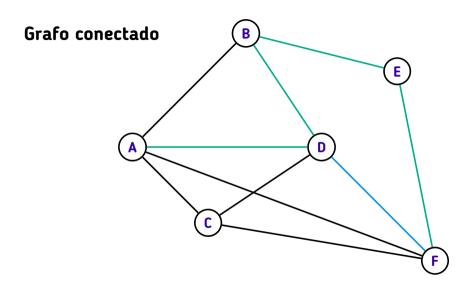
Um caminho é uma sequência não-nula de arestas da forma

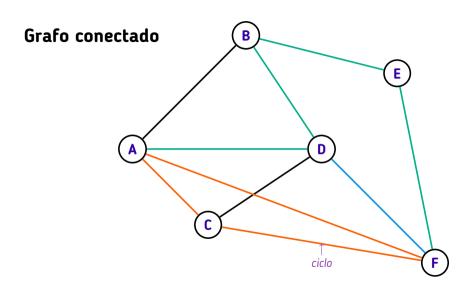
$$(u, \widehat{w_1}), (\widehat{w_1}, w_2), (w_2, \widehat{w_3}), \dots, (w_{n-1}, \widehat{w_n}), (\widehat{w_n}, v)$$

onde u é o ponto de partida e v o ponto de chegada









#### Referências

- 1. HALIM, Felix; HALIM, Steve. Competitive Programming 3, 2010.
- 2. LAAKSONEN, Antti. Competitive Programmer's Handbook, 2018.
- 3. SKIENA, Steven; REVILLA, Miguel. Programming Challenges, 2003.