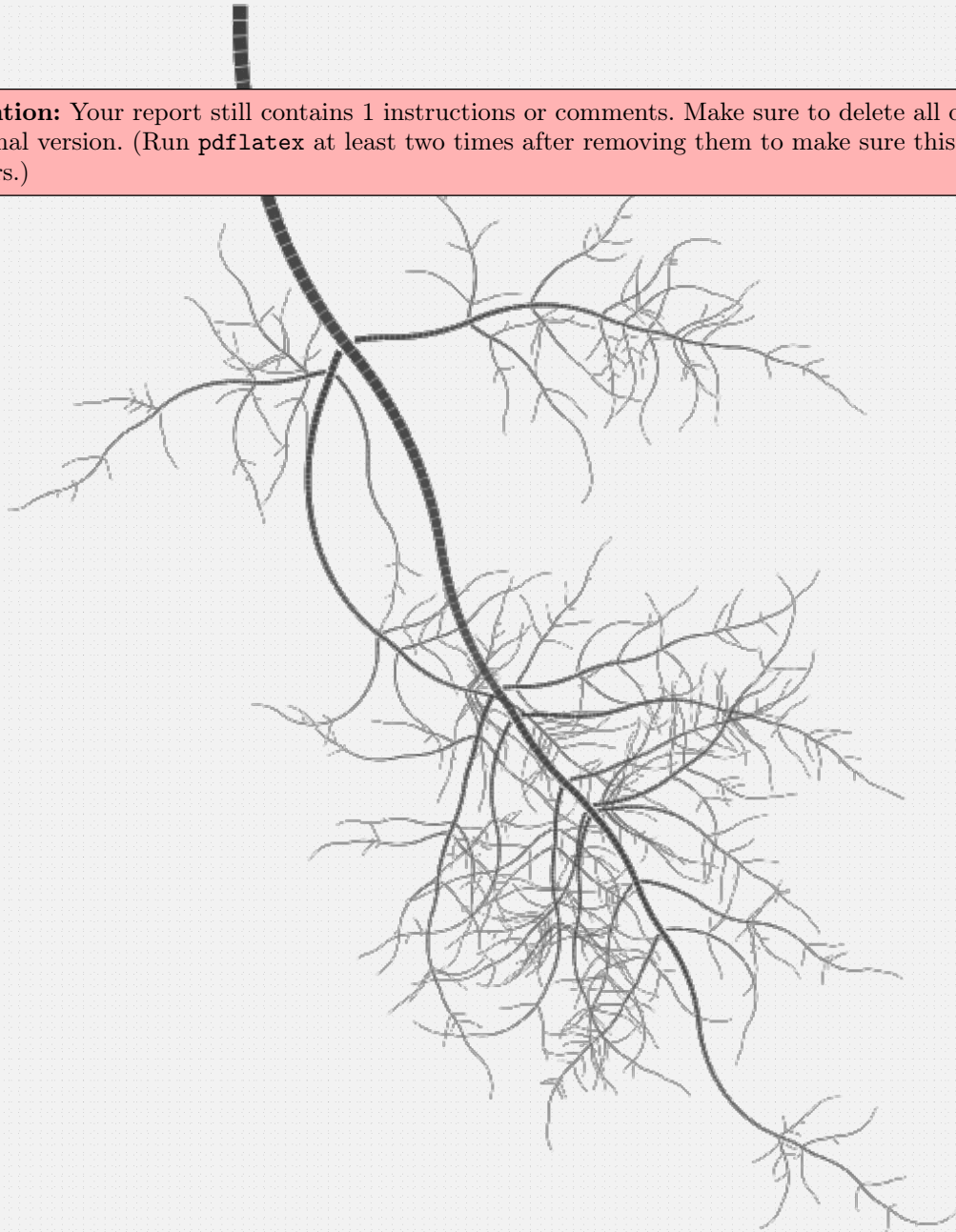


Advanced Programming Languages for A. I.

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FOR A. I.

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In the past decades much research has been done on the characterization and resolution of constraint satisfaction - and constraint optimization problems. This report discusses three challenges ; Sudoku puzzles, Hashiwokakero and a scheduling problem.

1 Sudoku

Sudoku is a well-known puzzle game which needs no introduction. It is typically modelled as a constraint satisfaction problem through the use of `all_different` constraints on rows, columns and blocks. Such global inequalities tend to improve upon the use of binary inequalities. The constraint generating code¹ is fairly trivial and needn't be detailed here.

There are several other ways one could model Sudoku. The widely cited study by Helmut Simonis [10] and subsequent studies provide some ideas. Four ‘dual’ models, two approaches based on a boolean characterisation, a combination of models provided by Laburthe, a model enforcing the singular occurrence of every value in every row, column and block, as well as a model with nothing but channeling constraints were considered². Tests were run on the provided puzzles³ as well as some minimum puzzles provided by Gordon Royle⁴. These are puzzles with a minimal amount of pre-filled cells (17 to be precise [7]), which does not mean that they are harder to solve.

The dual models hold a $N \times N$ array with all the decision variables. Whereas in the classic viewpoint the rows, columns and values of this array correspond to those of the input puzzle, every one of the four dual models changes their roles. The first two switch the role of rows or columns with those of values. In the third model every row and column of the array corresponds to a block and a position. In the fourth dual model every row represents a block, every column a value and every value a position within a block. For each of them it was harder to implement the necessary constraints, usually necessitating the use of auxiliary variables together with appropriate channeling constraints.

In one of his works Laburthe discusses various rules that can be used to resolve Sudoku puzzles, after which he details three models that he associates with the rules [3]. He ends up proposing a model for every level of difficulty of the input puzzle. An attempt was made at implementing his recommendation for ‘difficult’ puzzles. It decreased the average number of backtracks but increased the runtime.

¹ECLiPSe and CHR implementations are available in `/sudoku/model/classic.pl` and in `/sudoku/chr/model/classic.pl`.

²These are implemented with ECLiPSe in the `/sudoku/eclipse/model/` directory, and some CHR versions are in `/sudoku/chr/model/`.

³`/sudoku/benchmarks/benchmarks.pl` provides automatic benchmarking code.

⁴These are available online.

The boolean models include the natural combined model [8] and a more intuitive characterisation resembling an integer programming - or a SAT model [6] (using `occurrences/3` instead of sums, disjunctions and conjunctions). Both of them have $N \times N \times N$ boolean variables b_{rcv} which are true if the cell at row r and column c holds the value v . The natural combined model was cumbersome to implement and performed badly. It was introduced together with an algorithm which was tailored after it, and a constraint for unequality of lists isn't really supported by ECLiPSe⁵.

Note that it is usually not recommended to use a boolean model when integers can be used instead (as pointed out by Rossi [4]).

The last two models were found to be the most performant. The first makes use of the previously mentioned `occurrences/3` constraint to make every value occur just once in every row, column and block.

The second one generates nothing but channeling constraints. It has been demonstrated that this can provide good results despite such constraints being less 'tight' than `all_different` constraints⁶. When Dotú discussed it he was considering QuasiGroups [2]. This was extended⁷ to Sudoku puzzles by making use of three instead of two dual models (since blocks need to be considered as well). The variant in which channeling constraints between all models (one primal, three dual) are generated performed better than the one in which only channeling constraints between the primal and every dual model are applied. These variants are analogous to what Dotú referred to as 'trichanneling' and 'bichanneling'.

1.1 Experiments

Number of backtracks and running time for most of the models are displayed in table 1. Removal of 'big' (`all_different`) constraints in the classic model⁸ led to an increase in runtime which corroborates Demoen's experiences [1].



Figure 1: *Missing(6)* model

An interesting observation is that the two most performant models also have the same number of backtracks. One of them is the model with nothing but channeling constraints, the other uses the `occurrences/3` constraint which enforces arc-consistency. The first will detect when for a given *row-column*, *row-value*, *column-value* or *block-value* combination only one possible *value*, *column*, *row* or *position* remains. It will remove this value, column, ... from the domains of the other primal or dual variables⁹. The second does the same ; it can propagate unequalities when the domain of a variable is reduced to a singleton but also knows when a value can be put in only one particular cell of a row, column or block. It is probably slower because the constraint itself is more generic than reification constraints.

A model combining the classic viewpoint with the fourth dual model was set up. Number of backtracks and runtimes are displayed in table 3.

The `first_fail` heuristic generally outperforms `input_order`. It considers variables with the smallest domain first, rather than considering variables in the order they were given. This increases pruning power as removing a value from the domain will remove a bigger part of the search tree because the branching count increases as you go down the tree (which is not necessarily the case when using `input_order`). This tends to cause the number of backtracks to decrease. Unless the solution lies at the left side of the search tree generated

⁵A custom-made implementation as well as the `~/2` constraint which checks if two terms can be unified were tried. Channeling back to integers with `ic_global:bool_channeling/3` worked better (ironically).

⁶"The reason for this difference is that the primal not-equals constraints detect singleton variables (i.e. those variables with a single value), the channelling constraints detect singleton variables and singleton values (i.e. those values which occur in the domain of a single variable), whilst the primal all-different constraint detects global consistency (which includes singleton variables, singleton values and many other situations)"[5]

⁷The code lies in `sudoku/model/channeling.pl` in which a flag called `extended` can be used to opt for one of two variants.

⁸In his study Demoen gives several *Missing(6)* examples, models in which 6 of the `all_different` constraints are removed. *Missing(7)* models aren't Sudoku, and because of the stark rise in number of backtracks no further experimentation with the removal of 'small' constraints was done. The `eliminate_redundancy` flag can be used to toggle redundancy elimination on and off.

⁹The difference between unequality constraints and channeling is explained in more detail in [5]. Adding unequalities to the implementation slows down the search procedure because the channeling constraints already do this propagation and more.

Code Snippet 1: Channeling constraints for the combined viewpoint model

```
(multifor([Row,Column,Value], 1, N), param(Primal, Dual, K) do
    #=(Primal[Row,Column], Value, Bool), % reification
    block(K, Row, Column, Block), % calls utility function to convert  $R \times C \rightarrow B$ 
    position(K, Row, Column, Position), % calls utility function to convert  $R \times C \rightarrow P$ 
    #=(Dual[Block,Value], Position, Bool) % reification
)
```

by the `input_order` heuristic performance will increase in comparison.

As for the `all_different` constraints, more constraint propagation can be done when making use of `ic_global` since it enforces bound consistency [9] rather than enforcing arc-consistency on the corresponding inequalities. This decreases the number of backtracks for every puzzle but may increase runtime for some of them (such as *sudowiki_nb49*) because the propagation takes time. The global constraints do tend to perform a little better on average (about 11.2 versus 11.5 seconds total runtime).

Combining two viewpoints adds redundancy, leading to more propagation. The number of backtracks ends up decreasing because of this while the runtime may not necessarily decrease due to the larger number of constraints that have to be dealt with. As seen in table 3 the runtimes of the combined model generally laid somewhere in between those of the two original viewpoints.

1.2 CHR

Some of the viewpoints considered previously were implemented in CHR. The runtimes are shown in table 4. Initially the `first_fail` variable heuristic and the `indomain_min` value heuristic were used. These are easy to implement and generally perform quite nicely. Based on previous experiments with ECLiPSe it was concluded that (at least for the given benchmarks) `indomain_max` would be a better choice for the classic model. Using it instead cut runtime by about half, yet increased the runtime of the dual model which shows that indeed, models and heuristics interact. The combined model generally performed at least as well as the worst of the two other models except when runtime was already low.

Because the channeling-only viewpoint performed best in the ECLiPSe experiments an attempt was made at implementing it in CHR. The approach that was used led to complicated code and - ironically - the slowest runtimes seen yet. After a related failed experiment¹⁰ the focus was laid on keeping the code more simple and adding redundant constraints instead.

[Final]

¹⁰At some point a constraint `pos/4` was used to convert from *row-column* combinations to the corresponding blocks and positions (e.g. `pos(2,3,1,6)`). This slowed down the algorithm a lot because it increased the size of the constraint store.

	classic	dual1	dual2	dual3	boolean	laburthe	member	channeling
Total/Average	1981/142	2694/193	3467/248	2178/156	1998/143	14270/1020	1267/91	688/50
Total/Average (minimum)	2628/27	3183/32	3148/32	2612/27	2715/27	59340/594	1909/20	2026/20

Table 1: Runtimes for the various viewpoints (ECLiPSe) - shown in milliseconds. The **first_fail** and **indomain_min** heuristics were used.

	input_order	first_fail	input_order (global)	first_fail (global)
	classic - dual	classic - dual	classic - dual	classic - dual
<i>lambda</i>	495/4712 - 162/1715	149/977 - 72/523	22/3 - 17/0	22/3 - 19/0
<i>hard17</i>	132/873 - 87/389	92/419 - 68/198	20/1 - 27/1	27/1 - 29/1
<i>eastermonster</i>	62/119 - 75/125	58/101 - 84/155	237/51 - 236/49	169/33 - 256/66
<i>tarek_052</i>	72/193 - 68/200	62/130 - 59/137	370/61 - 324/63	224/35 - 180/34
<i>goldennugget</i>	103/520 - 98/441	87/358 - 90/281	687/104 - 517/72	334/76 - 264/44
<i>coloin</i>	368/2209 - 246/2288	61/83 - 55/80	249/88 - 675/194	79/8 - 63/10
<i>extra2</i>	305/4652 - 172/2894	641/7690 - 392/4232	18/0 - 18/0	19/0 - 21/0
<i>extra3</i>	530/4712 - 176/1715	166/977 - 93/523	22/3 - 17/0	22/3 - 19/0
<i>extra4</i>	1445/15116 - 409/6087	234/2097 - 139/1031	23/4 - 19/0	22/3 - 19/0
<i>inkara2012</i>	21/50 - 49/72	91/273 - 75/199	42/3 - 97/3	97/17 - 106/15
<i>clue18</i>	220/1838 - 161/1977	93/439 - 96/493	272/69 - 207/47	12/8 - 62/6
<i>clue17</i>	523/5520 - 329/3509	78/270 - 72/201	17/0 - 18/0	22/0 - 18/0
<i>sudowiki_nb28</i>	324/2851 - 501/4555	322/2221 - 455/3015	1083/413 - 1292/421	638/297 - 875/353
<i>sudowiki_nb49</i>	186/1078 - 127/749	88/655 - 137/672	246/48 - 201/40	224/58 - 275/62
Total/Average (ms)	4782/342 - 2654/190	2216/159 - 1880/135	3302/236 - 3660/262	2013/144 - 2201/157
Total/Average (backtracks)	44443/3175 - 26716/1909	16690/1192 - 11740/838	848/61 - 890/64	542/39 - 591/42

Table 2: Results for the classic and alternative **dual4** viewpoint (ECLiPSe) - shown as milliseconds/backtracks. In both models the **input_order** heuristic outperforms **first_fail** for puzzle *extra2*. This is because the latter tends to direct the search ‘down the wrong alleys’ hence the relatively large number of backtracks. Since the constraints are the same the runtime increases proportionally.

	input_order	first_fail	input_order (global)	first_fail (global)
<i>lambda</i>	160/507	122/188	29/3	36/3
<i>hard17</i>	138/264	116/155	26/1	36/1
<i>eastermonster</i>	79/102	163/218	351/51	395/78
<i>tarek_052</i>	137/170	137/136	625/61	219/39
<i>goldennugget</i>	217/286	50/10	1261/104	30/0
<i>coloin</i>	381/1009	55/19	363/87	67/8
<i>extra2</i>	115/263	120/253	25/0	28/0
<i>extra3</i>	152/507	113/188	32/3	60/3
<i>extra4</i>	266/1111	153/330	40/4	38/3
<i>inkara2012</i>	390/36	201/177	54/3	131/19
<i>clue18</i>	37/730	90/62	351/69	71/8
<i>clue17</i>	182/462	40/1	24/0	24/0
<i>sudowiki_nb28</i>	866/2581	400/713	1667/412	247/60
<i>sudowiki_nb49</i>	250/564	17/252	287/49	198/46
Total/Average (ms)	3369/241	1918/137	5127/37	1575/198
Total/Average (backtracks)	8592/613	2702/193	847/61	268/19

Table 3: Results for the combined (classic + **dual4**) viewpoint (ECLiPSe) - shown as milliseconds/backtracks.

	classic	dual	combined	classic (alt)	dual (alt)	combined (alt)	channeling-only
<i>lambda</i>	0	0	0	0	0	0	
<i>hard17</i>	0	0	0	0	0	0	
<i>eastermonster</i>	0	0	0	0	0	0	
<i>tarek_052</i>	0	0	0	0	0	0	
<i>goldennugget</i>	0	0	0	0	0	0	
<i>coloin</i>	0	0	0	0	0	0	
<i>extra2</i>	0	0	0	0	0	0	
<i>extra3</i>	0	0	0	0	0	0	
<i>extra4</i>	0	0	0	0	0	0	
<i>inkara2012</i>	0	0	0	0	0	0	
<i>clue18</i>	0	0	0	0	0	0	
<i>clue17</i>	0	0	0	0	0	0	
<i>sudowiki_nb28</i>	0	0	0	0	0	0	
<i>sudowiki_nb49</i>	0	0	0	0	0	0	
Total/Average	0	0	0	0	0	0	
Total/Average (minimum)	0	0	0	0	0	0	

Table 4: Results for the CHR Sudoku solver - shown as backtracks/milliseconds. Dual corresponds to dual14 in the ECLiPSe version, dual (bis) with dual13.

2 Hashiwokakero

 **TODO:** Task 2

3 Scheduling Meetings

The last challenge is the scheduling of some meetings, taking into account the preferences of the various persons involved. A constraint optimization problem where the cost is a function of the end time of the last meeting and the number of ‘violations’ (people of lower rank having their meeting after that of people of higher rank). This number of violations is of secondary importance.

Weekend constraints are generated first. If a person doesn’t want to meet on weekends then his or her meeting is not allowed to overlap with the first weekend that follows :

$$((S + StartingDay) \bmod 7) + D < 5$$

In the above constraint S and D represent the start and duration of the person’s meeting. Making direct use of `mod/3` leads to an instantiation error, necessitating the use of an auxiliary variable representing the result of the modulo operation.

Code Snippet 2: Weekend constraints

```
X :: 0..6, % Weekend constraints
_Q * 7 + X #= Start + StartingDay,
X + Duration #< 6
```

Precedence constraints and constraints assuring that no meetings overlap are generated last. The corresponding code is fairly trivial¹¹. The fact that the meeting with the minister should come last is equivalent to adding $N - 1$ precedence constraints with N the total number of persons.

The cost function is defined as $(V_{max} \times E) + V$ where V_{max} is the maximum number of rank violations, E is the end time of the meeting with the minister and V is the actual number of rank violations for a given solution. This ensures that whenever two solutions have a different E , the solution with the smallest E will have the lowest cost (whatever the number of violations V). Yet if two solutions have the same end time E , then it’s the number of violations V that will determine what solution is best.

¹¹All code for this third challenge can be found in `/src/scheduling/scheduling.pl`.

Code Snippet 3: Calculating the number of violations. The cost function is defined as $MaxViolations * (StartTime_{minister} + Duration_{minister}) + Violations$

```

violations(N, Ranks, StartTimes, Violations, MaxViolations) :-
  (for(I, 1, N-1),
   fromto([], InViolations, OutViolations, ViolationList),
   param(StartTimes, N, Ranks) do
     Rank is Ranks[I],
     (for(J, I+1, N-1),
      fromto([], In, Out, List),
      param(Rank, Ranks, StartTimes, I) do
        OtherRank is Ranks[J],
        (Rank < OtherRank -> Out = [(StartTimes[I] #> StartTimes[J])|In] ;
         (Rank > OtherRank -> Out = [(StartTimes[I] #< StartTimes[J])|In] ;
         Out = In))
      ),
      append(InViolations, List, OutViolations)
    ),
    length(ViolationList, MaxViolations),
    Violations #= sum(ViolationList).

```

An additional constraint was used for the cost function, stating that it cannot be smaller than $V_{max} \times D_{tot}$ with D_{tot} the sum of all meeting durations. This makes a difference¹².

Some implied constraints were added to increase performance. In case two persons have a different rank but the same meeting duration and weekend preferences, a corresponding order on their start times can safely be imposed. This mustn't override the precedence constraints.

Table 5 shows the runtime for each benchmark. Two versions are considered ; one ensures that no two meetings overlap by imposing a $(S_1 + D_1 \leq S_2 \text{ or } S_2 + D_2 \leq S_1)$ constraint for every such pair, the other version uses a global version of these same constraints provided by the `ic_edge_finder` library. It's clear that the global version outperforms the other one. The time it takes to propagate the constraints is usually compensated for by the reduction in nodes having to be considered due to the pruning of the search tree.

Instead of making use of implied constraints one can also tinker with the various heuristics provided by the `search/5` procedure. Some of those lend themselves to some benchmarks but not to others.

The `indomain_min` heuristic performed better than `indomain_max` as it is an optimisation problem after all, meaning that selecting the minimum starting time selects solutions with a smaller cost first. A solution with a lower cost will prune the search tree more.

	input_order	input_order + ic_edge_finder
<i>bench1a</i>	207	326
<i>bench1b</i>	135	180
<i>bench1c</i>	314	269
<i>bench2a</i>	3636	2110
<i>bench2b</i>	19269	2043
<i>bench2c</i>	371	366
<i>bench3a</i>	143	138
<i>bench3b</i>	386	389
<i>bench3c</i>	286	309
<i>bench3d</i>	201	333
<i>bench3e</i>	405	371
<i>bench3f</i>	506	413
<i>bench3g</i>	568	479
Total/Average	26422/2033	7721/594

Table 5: Benchmark results for the Scheduling Meetings challenge, shown in milliseconds.

¹²In our tests the total runtime was reduced by a factor of 4 (not when making use of the `ic_edge_finder` version).

Appendix

Reflection

When implementing the Sudoku solver in CHR we took a glimpse at Thom’s implementation. It’s given in his book as a solution to some exercise. Once we understood his approach we had the tendency to implement the viewpoints by making use of the same strategy ; constraints representing variables, constraints representing values and a simple implementation of the `first_fail` heuristic. This worked well, but it wasn’t the most creative thing to do. Further experimentation with other models, heuristics, redundant constraints, ... was our way to (hopefully) compensate for this.

Workload

Some questions were asked online after all the code had been written. There were four of them, all about Sudoku. One on how to enforce equality of lists (we already had the solution but wanted to be sure there was no better alternative). One on looping through a list which is a subscript of an array (we found the appropriate solution ourselves). One on the inner workings of [occurrences/3](#). And a final one on memory usage. Aside from a quick experiment with $\sim=$ /2 no code was rewritten as a result. None of the questions mentioned Sudoku. All the viewpoints were either thought of by ourselves or come from the literature that was cited in this report.

We started working in mid-April and finished a few days before the deadline, each having worked about ??? hours in total. This includes reading (parts of) the educational material, researching, tinkering, programming, debugging and writing the report.

Overview of the Code

Folder	File	Description
/src/sudoku/	utils.pl	Utility functions for Sudoku (CHR & ECLiPSe)
/src/sudoku/benchmarks/	benchmarks.pl	Automatic benchmarking code
/src/sudoku/benchmarks/puzzles/	*	Sudoku benchmarks
/src/sudoku/chr/	solver.pl	Sudoku solver (CHR)
/src/sudoku/chr/model/	*	Sudoku viewpoints (CHR)
/src/sudoku/eclipse/	solver.pl	Sudoku solver (ECLiPSe)
/src/sudoku/eclipse/model/	*	Sudoku viewpoints (ECLiPSe)
/src/hashiwokakero/	jschimpf.pl	Hashiwokakero solution
/src/hashiwokakero/benchmarks/	*	Hashiwokakero benchmarks
/src/scheduling/	scheduling.pl	Scheduling meetings solution

Table 6: Overview of the source code.

References

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