

# Doors

Bruno Carvalho up201606517 and Amadeu Pereira up201605646

Faculdade de Engenharia da Universidade do Porto

**Abstract.** In this paper we describe and explore a simple restraint logic programming solution for the logic problem *Doors*.

**Keywords:** doors · logic puzzle · rooms

## 1 Problem Description

The logic problem *Doors* is a puzzle on a rectangular board whose cells are either empty or contain natural numbers. The board is thought of like a *house*. Each cell is a *room*, and two adjacent cells are separated by a *wall* with one *door*. That door may be either open or closed. If it is open, then the cell can see its adjacent room through the doorway. A man standing in a room can look in all four directions — north, east, west, south — and count the number of *visible rooms*.

The puzzle consists in discovering an assignment of open and closed doors to the walls of the board such that the natural number in each non-empty cell is how many rooms are visible from that cell (including itself). There may be multiple solutions, or none at all.

### 1.1 Examples

4	2	4	2			
5	3	5	3			
4	3		4			
2	3	2	4			

3	5	3	2	2	5	
		4	5			
4	6	4	5		6	
5	7	5	6	1	4	

2	1	4			1	
2	2	5	5	6	4	
1	3	5	4	5	2	
	3	5	5	6	3	

## 2 Approach

A puzzle of size  $n \times m$  is represented internally by three matrices (list of lists): *Board*, of size  $n \times m$ , holding the cell numbers; *Vertical*, of size  $n \times (m-1)$ , holding the vertical walls; and *Horizontal*, of size  $(n-1) \times m$ , holding the horizontal walls.

For each cell  $(R, C)$ ,  $R = 1, 2, \dots, n$ ,  $C = 1, 2, \dots, m$ , indices in *Board*, the left wall has index  $(R, C-1)$  in *Vertical*, the right wall has index  $(R, C)$  in *Vertical*, the top wall has index  $(R-1, C)$  in *Horizontal*, the bottom wall has index  $(R, C)$  in *Horizontal*. Each wall in the board is assigned the number 0 for closed door and 1 for open door.

## 2.1 Restrictions

When solving a puzzle *Board* is fully instantiated, while *Vertical* and *Horizontal* contain domain variables (domain  $\{0, 1\}$ ). Empty cells in the *Board* are represented by a 0, as it is never a valid visible room counter.

1	<i>T</i>	?	?	?	?	?	2
<i>X</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>	<i>G</i>
?	<i>M</i>	?	?	?	?	?	?
?	<i>N</i>	?	?	?	?	?	?

Consider the puzzle above. The horizontal range  $A - G$  consists of 7 rooms and 6 vertical doors: let  $\{b, c, d, e, f, g\}$  be these vertical doors, from left to right.

Focus on room *A*. If  $b = 0$  then *A* sees no rooms to its right. If  $b = 1$  and  $c = 0$  then *A* sees only room *B*. A general formula can be deduced by noticing that closed doors behave as zero elements.

Let  $e_A$  be the total number of rooms *A* sees to its right (east), then

$$e_A = b + b(c + c(d + d(e + e(f + f(g + g \cdot 0)))))) \quad (1)$$

Now, if we analogously define  $w_A$  for west,  $n_A$  for north and  $s_A$  for south, then we find that the number in cell *A* must be  $e_A + w_A + n_A + s_A + 1$ .

Implementing these restrictions in PROLOG is surprisingly simple. We start with a predicate to compute formula:

```
calculate_value([], 0).
calculate_value([H|T], V) :-
    calculate_value(T, V1),
    V #= H + H*V1.
```

Then, for each non-zero cell  $(R, C)$  on the *Board*, we retrieve as a list the four ranges of doors to the right, left, top and bottom of  $(R, C)$ , apply the formula for each list, and finally the restriction:

```
restrict_cell(Board, _, _, [R,C]) :-
    matrixnth1([R,C], Board, 0), !. % empty cell
restrict_cell(Board, Vertical, Horizontal, [R,C]) :-
    matrixnth1([R,C], Board, Value),
    right_total(Vertical, [R,C], Right),
    left_total(Vertical, [R,C], Left),
    top_total(Horizontal, [R,C], Top),
    bot_total(Horizontal, [R,C], Bot),
    Right + Left + Top + Bot + 1 #= Value.
```

## 2.2 Search strategy

The labeling strategy implemented was the `ffc`. This strategy uses "the most constrained heuristic is used: a variable with the smallest domain is selected, breaking ties by (a) selecting the variable that has the most constraints suspended on it and (b) selecting the leftmost one." Even though all of our variables have the same domain we decided to use this labeling strategy since it would provide us the puzzle solutions more efficiently.

## 2.3 Random board generation

It is also implemented a way to generate a random board with a specified number of rows and a specified number of columns. The methodology used here is the opposite of what the game consists. We start by creating the Vertical and the Horizontal matrices and randomly (but with a certain probability) populate them with 0 (door) or 1 (no door). Certainly, the probability of an element to be an 1 is higher than the probability to be an 0.

```
door_prob(0.3).
```

After this step it is pretty straight forward to generate the board. We go through each element of the Board matrix and, using the formula discussed above, calculate each value.

```
calculate_value(Board, Vertical, Horizontal, [R,C]) :-
    matrixnth1([R,C], Board, Value),
    right_total(VERTICAL, [R,C], Right),
    left_total(VERTICAL, [R,C], Left),
    top_total(HORIZONTAL, [R,C], Top),
    bot_total(HORIZONTAL, [R,C], Bot),
    Value is Right + Left + Top + Bot + 1.
```

## 3 Solution presentation

We represent the board using Unicode box drawing characters. Cells are sufficiently sized for all board numbers with `max_width_number/3` and centered with `center_number/2`.

Predicate `write_connector/4` serves to select the appropriate unicode character to connect the cell corners inside and in the edges of the board, and `write_multiple/2` serves to write the same character multiple characters. Using these auxiliary predicates, the board is written in a straightforward fashion, with number rows and horizontal wall rows interleaved.

## 4 Results

```
| ?- doors([[0,4,7,5,0,4],[0,5,7,0,3,0],[3,3,5,2,3,4],[3,2,5,3,4,6]]).
Time: 0.01000s
Resumptions: 7835
Entailments: 1988
Prunings: 2946
Backtracks: 47
Constraints created: 323
```

	4	7	5		4
	5	7		3	
3	3	5	2	3	4
3	2	5	3	4	6

```
yes
```

**Fig. 1.** Example of a solver execution.

In terms of mere efficiency, on square boards we have the following scores:

number of cells	time	constraints created	number of backtracks
9	0.00s	150	0
25	0.00s	626	5
49	0.00s	2058	146
64	0.030s	2963	813
100	0.610s	8767	18734
110	7.10s	52149	383107
150	61.64s	12891362	10509283

**Table 1.** Solver execution statistics on various square board sizes

## 5 Conclusions

We believe our solution is elegant and straightforward, although not terribly efficient. A problem centered around rooms *seeing* each other has a solution whose logic restraints do not involve more than one room at a time.

To extend the simple solution we included detailed execution statistics and a nice board display on the console.

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6 Bruno Carvalho up201606517 and Amadeu Pereira up201605646

## **6 Attachments**

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