# Computational Complexity Analysis: Least Squares vs. Gradient Descent

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STAT 102B

#### Least Squares Problem

- Given data matrix  $X \in \mathbb{R}^{n \times p}$  and response vector  $y \in \mathbb{R}^n$
- Objective: Find  $\beta \in \mathbb{R}^p$  that minimizes the sum of squared errors (SSE)

$$\min_{\beta} SSE(\beta) = \frac{1}{2n} ||X\beta - y||_2^2$$
 (1)

The closed form solution is given by:

$$\hat{\beta} = (X^{\top}X)^{-1}X^{\top}y \tag{2}$$

 This is exact, but can be computationally intensive for large data sets, especially in terms of number of predictors p

#### Clarification on Notation

In the lecture notes, I used m for the number of observations in order to differentiate from n which is reserved for the dimensionality of the argument of the objective function

$$f(x), x \in \mathbb{R}^n$$

In this slide deck, I use n for the number of observations, in order for the calculations to align with standard calculations on computational in the literature

The notation  $\mathcal{O}(\cdot)$  used in subsequent slides indicates the order of computations required

## Computational Steps for Least Squares

#### Computing the least squares solution requires:

- **Step 1**: Compute  $X^{\top}X \in \mathbb{R}^{p \times p}$ 
  - ▶ To obtain each entry in this  $p \times p$  entry we need to multiply two *n*-dimensional vectors;
    - the vector that contains the *n*-observations for variable *i* with the vector that contains the *n*-observations for variable *j*, with  $i, j = 1, \dots, p$
  - **Each** such vector multiplication requires  $\mathcal{O}(n)$  operations
  - ► Total:  $\mathcal{O}(np^2)$  operations, since  $X^\top X$  has  $p^2$  entries and each entry requires  $\mathcal{O}(n)$  operations
- **Step 2:** Compute  $X^{\top}y \in \mathbb{R}^p$ 
  - To obtain each  $(X^{\top}y)$  we need to multiply again two *n*-dimensional vectors; the vector that contains the *n*observations for each variable, with the *n*-dimensional response vector v
  - Total:  $\mathcal{O}(np)$  operations

# Computational Steps for Least Squares (Continued)

- **Step 3:** Solve the system  $X^{\top}X\beta = X^{\top}y$ 
  - ► Typically requires computing the inverse  $(X^\top X)^{-1}$  ► Matrix inversion requires  $\mathcal{O}(\rho^3)$  operations

  - In practice, we typically solve using Cholesky decomposition, QR decomposition, or eigenvalue decomposition for better numerical stability
  - These methods also require  $\mathcal{O}(p^3)$  operations
- Total computational complexity:

$$\mathcal{O}(np^2 + p^3) \tag{3}$$

#### Computational Burden for Large Data sets

Consider a large-scale problem with:

$$n = 10^6$$
 (samples),  $p = 10^5$  (features) (4)

- $np^2 = 10^6 \cdot (10^5)^2 = 10^{16}$  operations
- $p^3 = (10^5)^3 = 10^{15}$  operations

The dominant term is  $\mathcal{O}(np^2) = \left| 10^{16} \right|$  operations

- Modern CPUs can perform  $\sim 10^9$  operations per second
- This would take approximately  $10^7$  seconds  $\approx 115$  days!
- Memory requirements would be enormous as well

#### Gradient Descent for Least Squares

Instead of computing the closed-form solution, we can iteratively approach the minimum:

- Initialize  $\beta_0$  (often as zeros or small random values)
- Update rule:

$$\beta_{k+1} = \beta^{(t)} - \eta \nabla_{\beta} SSE(\beta)$$
 (5)

where  $\eta > 0$  is the step size

• The gradient of the SSE function is given by:

$$\nabla_{\beta} \frac{1}{2n} \| X\beta - y \|_{2}^{2} = \frac{1}{n} [X^{\top} (X\beta - y)]$$
 (6)

# Computational Steps per Iteration

#### Each iteration of gradient descent requires:

- **Step 1**: Compute  $X\beta \in \mathbb{R}^n$ 
  - Each entry multiplies two p-dimensional vectors and there are n such entries
  - $\triangleright$  Total:  $\mathcal{O}(np)$  operations
- Step 2: Compute residual  $r = X\beta y \in \mathbb{R}^n$ 
  - Simple vector subtraction
  - Total:  $\mathcal{O}(n)$  operations

# Computational Steps per Iteration (Continued)

- **Step 3**: Compute  $X^{\top}r \in \mathbb{R}^p$ 
  - Each entry requires multiplication of two n-dimensional vectors and there are p such entries
  - ► Total:  $\mathcal{O}(np)$  operations
- **Step 4**: Update  $\beta$ 
  - ► Simple vector subtraction
  - Total:  $\mathcal{O}(p)$  operations
- Total computational complexity per iteration:

$$O(np)$$
 (7)

## Computational Advantage for Large Data sets

For our large-scale problem with:

$$n = 10^6$$
 (samples),  $p = 10^5$  (features) (8)

- Cost per iteration:  $np = 10^6 \cdot 10^5 = 10^{11}$  operations
- With a good choice of step size, convergence typically requires  $10^2-10^3$  iterations
- Total cost:  $\sim 10^{13} 10^{14}$  operations

Gradient descent offers a computational advantage of  $\sim 10^2-10^3$  times over the closed-form solution for large-scale problems

### Conditions for Convergence

For gradient descent to find the optimal regression coefficients:

- The matrix  $X^{T}X$  must be **positive definite** 
  - Guarantees the existence of a unique global minimum
  - This is the same condition needed for  $(X^TX)^{-1}$  to exist in the closed-form solution
- The Hessian of the SSE function is  $\frac{1}{n}(X^{T}X)$
- If  $X^{\top}X$  is positive definite, gradient descent with an appropriate step size will converge to the global minimum

### Step Size Selection

• Theoretical bound: For convergence, the step size  $\eta$  must satisfy:

$$0 < \eta < \frac{2}{\lambda_{\max}(X^{\top}X)} \tag{9}$$

where  $\lambda_{\max}(X^{\top}X)$  is the largest eigenvalue of  $X^{\top}X$ 

- Computing  $\lambda_{\max}(X^{\top}X)$  requires  $\mathcal{O}(p^3)$  operations
  - ▶ This defeats the purpose of using gradient descent to avoid  $\mathcal{O}(p^3)$  computations
- Practical solution: Backtracking line search
  - Start with a relatively large step size
  - Decrease it gradually until a sufficient decrease condition is met
  - No need to compute eigenvalues explicitly

### Computational Cost Comparison

Method	Computational Complexity
Closed-form Least Squares	$\mathcal{O}(np^2+p^3)$
Gradient Descent (per iteration)	$\mathcal{O}(np)$
Gradient Descent (total, k iterations)	$\mathcal{O}(k \cdot np)$

For  $n = 10^6$  and  $p = 10^5$ :

- ullet Closed-form:  $\sim 10^{16}$  operations
- $\bullet$  Gradient Descent (1000 iterations):  $\sim 10^{14}$  operations

For large-scale problems, gradient descent offers a significant computational advantage.

# When to Use Fach Method

#### Use closed-form solution when:

- p is small (few features)
- Exact solution is required
- Numerical stability isn't a significant concern
- Available computational resources can handle  $\mathcal{O}(np^2 + p^3)$

#### Use gradient descent when:

- p is large (many features)
- Approximate solution is acceptable
- The problem is "well-conditioned" for faster convergence
- Memory constraints prevent storing or operating on  $X^{\top}X$

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Other Considerations

- Beyond basic gradient descent:

  Momentum methods: Polyak, Nesterov
  - Adaptive learning rates: AdaGrad, ADAM, ADAM-W
  - Stochastic variants: Stochastic gradient
    - Further computational savings with subsampling
    - Ideal for extremely large datasets (n ≫ p)
  - Coordinate descent: Update one coordinate at a time (will be covered after Exam I)
    - Particularly efficient when X is sparse

#### Numerical stability:

- Closed-form solution may suffer from ill-conditioning
- Gradient methods can be more robust with appropriate step sizes