

Chapter 4: Network Layer

Chapter goals:

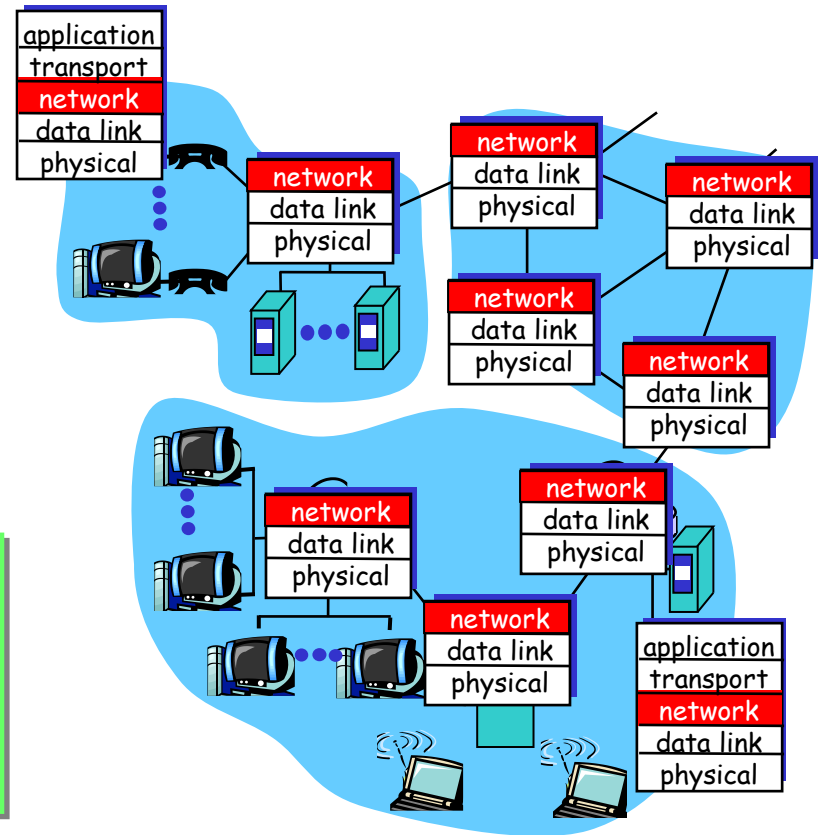
- ❑ understand principles behind network layer services:
 - network layer service models
 - forwarding versus routing
 - how a router works
 - routing (path selection)
 - dealing with scale
 - advanced topics: IPv6, mobility
- ❑ instantiation, implementation in the Internet

Chapter 4: Network Layer

- ❑ 4.1 Introduction
- ❑ 4.2 Virtual circuit and datagram networks
- ❑ 4.3 What's inside a router
- ❑ 4.4 IP: Internet Protocol
 - Datagram format
 - IPv4 addressing
 - NAT
 - ICMP
 - IPv6
- ❑ 4.5 Routing algorithms
 - Link state
 - Distance Vector
 - Hierarchical routing
- ❑ 4.6 Routing in the Internet
 - RIP
 - OSPF
 - BGP
- ❑ 4.7 Broadcast and multicast routing

Network layer

- ❑ on sending side encapsulates segments into datagrams
 - ❑ on rcving side, delivers segments to transport layer
 - ❑ network layer protocols in every host, router
- ❑ Router examines header fields in all IP datagrams passing through it



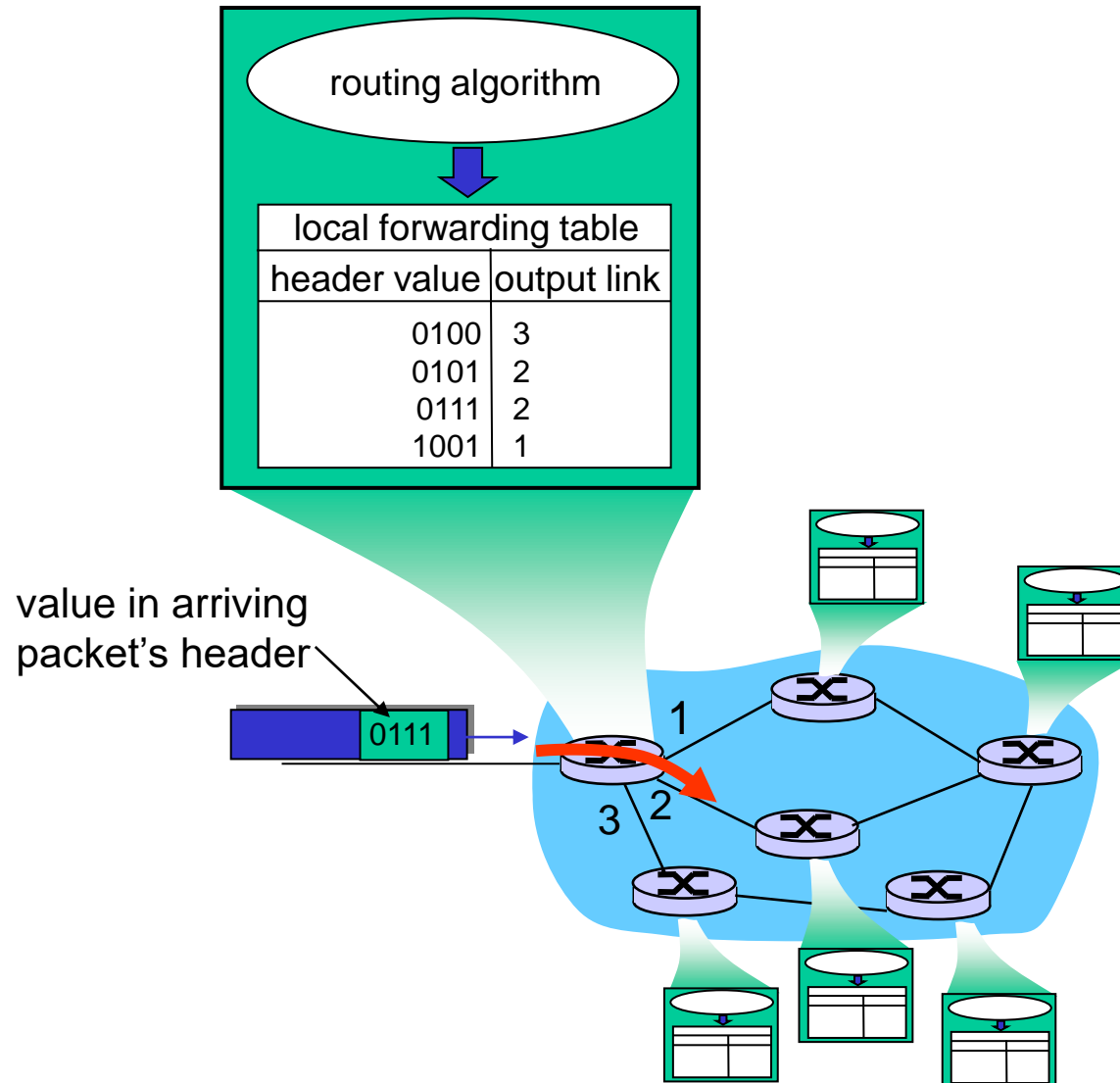
Key Network-Layer Functions

- ❑ *forwarding*: move packets from router's input to appropriate router output
- ❑ *routing*: determine route taken by packets from source to dest.
 - *Routing algorithms*

analogy:

- ❑ *routing*: process of planning trip from source to dest
- ❑ *forwarding*: process of correct left turns, right turns, exits, etc.

Interplay between routing and forwarding



Connection setup

- ❑ important function in some network architectures:
 - ATM, frame relay, X.25

- ❑ Before datagrams flow, two hosts and intervening routers establish virtual connection
 - Routers get involved

- ❑ Network and transport layer cnctn service:
 - **Network:** between two hosts
 - **Transport:** between two processes

Network service model

Q: What *service model* for “channel” transporting datagrams from sender to rcvr?

Example services for individual datagrams:

- ❑ guaranteed delivery
- ❑ Guaranteed delivery with less than 40 msec delay

Example services for a flow of datagrams:

- ❑ In-order datagram delivery
- ❑ Guaranteed minimum bandwidth to flow
- ❑ Restrictions on changes in inter-packet spacing

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Network layer connection and connection-less service

- ❑ Datagram network provides network-layer connectionless service
- ❑ VC network provides network-layer connection service
- ❑ Analogous to the transport-layer services, but:
 - **Service:** host-to-host
 - **No choice:** network provides one or the other
 - **Implementation:** in the core

Virtual circuits

“source-to-dest path behaves much like telephone circuit”

- performance-wise
- network actions along source-to-dest path

- ❑ call setup, teardown for each call *before* data can flow
- ❑ each packet carries VC identifier (not destination host address)
- ❑ every router on source-dest path maintains “state” for each passing connection
- ❑ link, router resources (bandwidth, buffers) may be *allocated* to VC

VC implementation

A VC consists of:

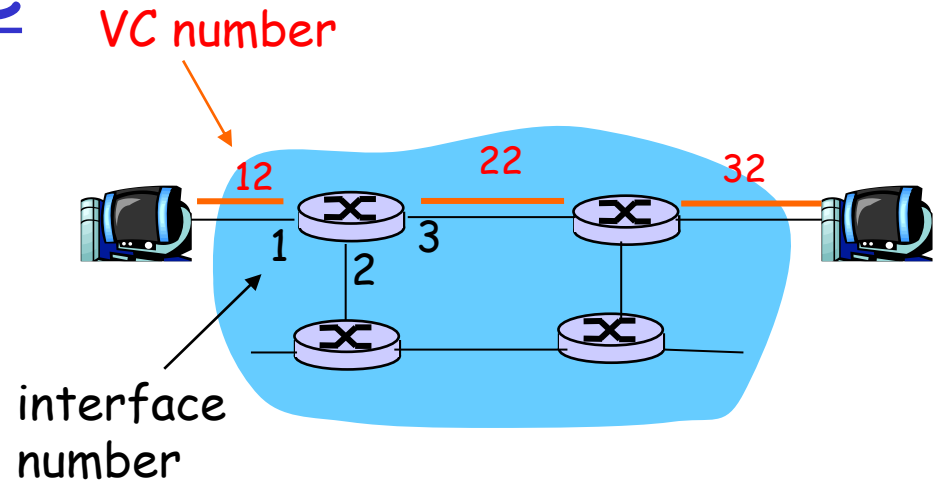
1. Path from source to destination
2. VC numbers, one number for each link along path
3. Entries in forwarding tables in routers along path

Example next slide

- ❑ Packet belonging to VC carries a VC number.
- ❑ VC number must be changed on each link.
 - New VC number comes from forwarding table

Forwarding table

Forwarding table in northwest router:

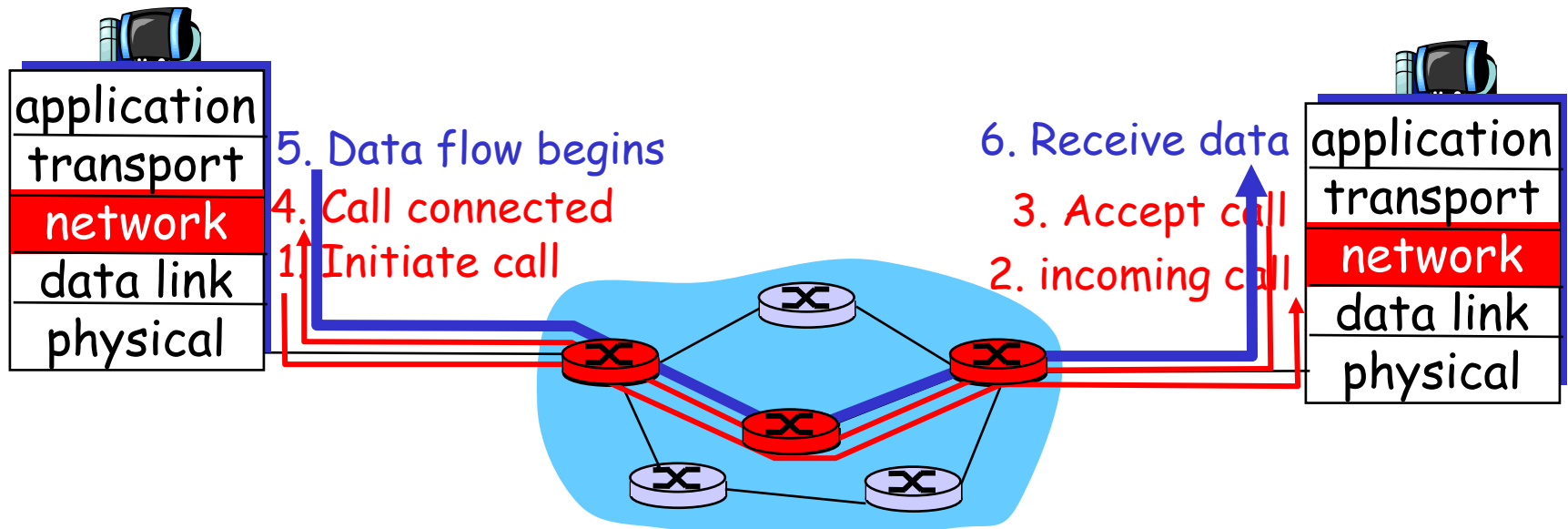


Incoming interface	Incoming VC #	Outgoing interface	Outgoing VC #
1	12	3	22
2	63	1	18
3	7	2	17
1	97	3	87
...

Routers maintain connection state information!

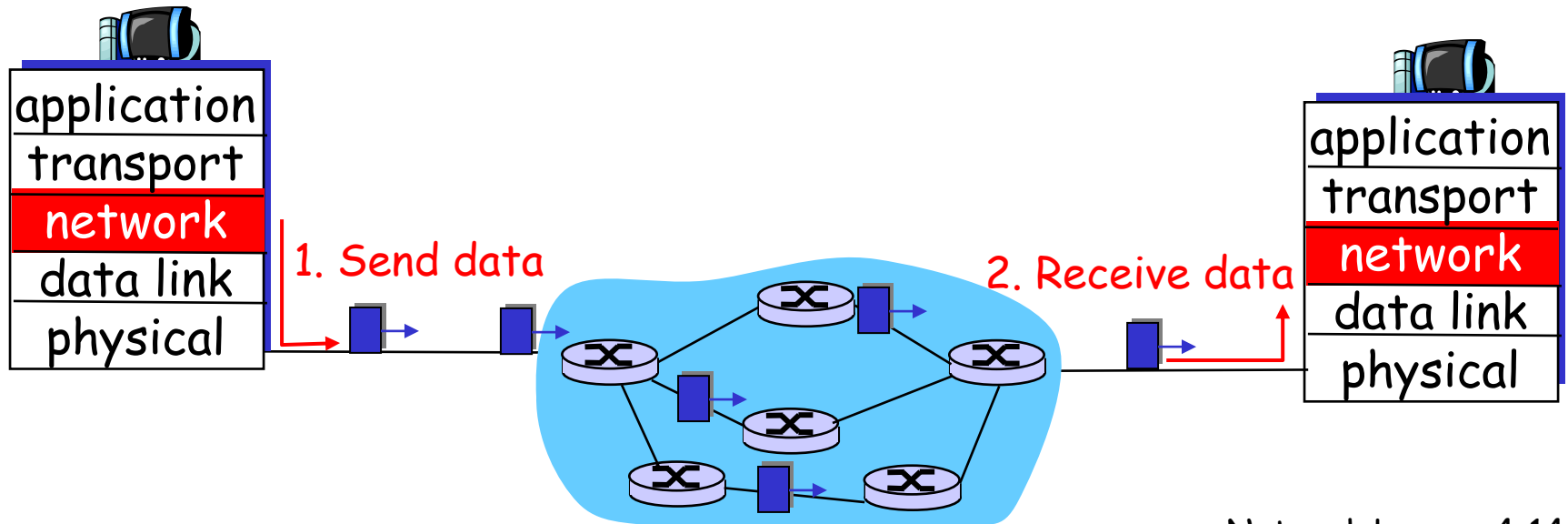
Virtual circuits: signaling protocols

- ❑ used to setup, maintain teardown VC
- ❑ used in ATM, frame-relay, X.25
- ❑ not used in today's Internet



Datagram networks

- ❑ no call setup at network layer
- ❑ routers: no state about end-to-end connections
 - no network-level concept of “connection”
- ❑ packets forwarded using destination host address
 - packets between same source-dest pair may take different paths



Forwarding table

4 billion
possible entries

<u>Destination Address Range</u>	<u>Link Interface</u>
11001000 00010111 00010000 00000000 through 11001000 00010111 00010111 11111111	0
11001000 00010111 00011000 00000000 through 11001000 00010111 00011000 11111111	1
11001000 00010111 00011001 00000000 through 11001000 00010111 00011111 11111111	2
otherwise	3

Longest prefix matching

<u>Prefix Match</u>	<u>Link Interface</u>
11001000 00010111 00010	0
11001000 00010111 00011000	1
11001000 00010111 00011	2
otherwise	3

Examples

DA: 11001000 00010111 00010110 10100001

Which interface?

DA: 11001000 00010111 00011000 10101010

Which interface?

Datagram or VC network: why?

Internet

- ❑ data exchange among computers
 - “elastic” service, no strict timing req.
- ❑ “smart” end systems (computers)
 - can adapt, perform control, error recovery
 - simple inside network, complexity at “edge”
- ❑ many link types
 - different characteristics
 - uniform service difficult

ATM

- ❑ evolved from telephony
- ❑ human conversation:
 - strict timing, reliability requirements
 - need for guaranteed service
- ❑ “dumb” end systems
 - telephones
 - complexity inside network