Analysis Report

void computeHOGSharedPred<float, float, float, int=8, int=8, int=16, int=16, int=64>(float*, float*, float*, float*, int, int, int, int)

Duration	696.684 μs
Grid Size	[31,1,1]
Block Size	[256,1,1]
Registers/Thread	40
Shared Memory/Block	36 KiB
Shared Memory Requested	96 KiB
Shared Memory Executed	96 KiB
Shared Memory Bank Size	4 B

[0] GeForce GTX 960

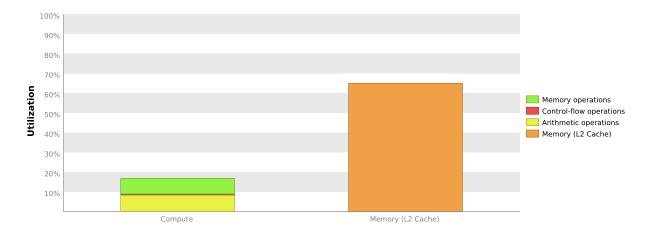
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GPU UUID	GPU-0db32734-f94e-48a7-8b5d-4604317dc554
Compute Capability	5.2
Max. Threads per Block	1024
Max. Shared Memory per Block	48 KiB
Max. Registers per Block	65536
Max. Grid Dimensions	[2147483647, 65535, 65535]
Max. Block Dimensions	[1024, 1024, 64]
Max. Warps per Multiprocessor	64
Max. Blocks per Multiprocessor	32
Single Precision FLOP/s	2.644 TeraFLOP/s
Double Precision FLOP/s	82.624 GigaFLOP/s
Number of Multiprocessors	8
Multiprocessor Clock Rate	1.291 GHz
Concurrent Kernel	true
Max IPC	6
Threads per Warp	32
Global Memory Bandwidth	112.16 GB/s
Global Memory Size	4 GiB
Constant Memory Size	64 KiB
L2 Cache Size	1 MiB
Memcpy Engines	2
PCIe Generation	2
PCIe Link Rate	5 Gbit/s
PCIe Link Width	16

1. Compute, Bandwidth, or Latency Bound

The first step in analyzing an individual kernel is to determine if the performance of the kernel is bounded by computation, memory bandwidth, or instruction/memory latency. The results below indicate that the performance of kernel "void computeHOGSharedPred<f..." is most likely limited by memory bandwidth. You should first examine the information in the "Memory Bandwidth" section to determine how it is limiting performance.

1.1. Kernel Performance Is Bound By Memory Bandwidth

For device "GeForce GTX 960" the kernel's compute utilization is significantly lower than its memory utilization. These utilization levels indicate that the performance of the kernel is most likely being limited by the memory system. For this kernel the limiting factor in the memory system is the bandwidth of the L2 Cache memory.



2. Memory Bandwidth

Memory bandwidth limits the performance of a kernel when one or more memories in the GPU cannot provide data at the rate requested by the kernel. The results below indicate that the kernel is limited by the bandwidth available to the L2 cache.

2.1. Global Memory Alignment and Access Pattern

Memory bandwidth is used most efficiently when each global memory load and store has proper alignment and access pattern.

Optimization: Each entry below points to a global load or store within the kernel with an inefficient alignment or access pattern. For each load or store improve the alignment and access pattern of the memory access.

/home/adas/cuda-workspace/CudaVisionSysDeploy/Release/../src/init/../device/HOG/HOGdescriptor.cuh

	/nome/adas/edda-workspace/edda-visionsysbeproy/recease//sic/infd/device/f100/f100descriptor.edn
Line 142	Global Load L2 Transactions/Access = 31.9, Ideal Transactions/Access = 4 [1999616 L2 transactions for 62720 total executions]
Line 142	Global Load L2 Transactions/Access = 31.9, Ideal Transactions/Access = 4 [1999616 L2 transactions for 62720 total executions]
Line 152	Global Store L2 Transactions/Access = 31.9, Ideal Transactions/Access = 4 [7811 L2 transactions for 245 total executions]
Line 152	Global Store L2 Transactions/Access = 31.9, Ideal Transactions/Access = 4 [7811 L2 transactions for 245 total executions]
Line 152	Global Store L2 Transactions/Access = 31.9, Ideal Transactions/Access = 4 [7811 L2 transactions for 245 total executions]
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	executions]
Line 152	Global Store L2 Transactions/Access = 31.9, Ideal Transactions/Access = 4 [7811 L2 transactions for 245 total executions]
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executions]

2.2. GPU Utilization Is Limited By Memory Bandwidth

The following table shows the memory bandwidth used by this kernel for the various types of memory on the device. The table also shows the utilization of each memory type relative to the maximum throughput supported by the memory. The results show that the kernel's performance is potentially limited by the bandwidth available from one or more of the memories on the device.

Optimization: Try the following optimizations for the memory with high bandwidth utilization.

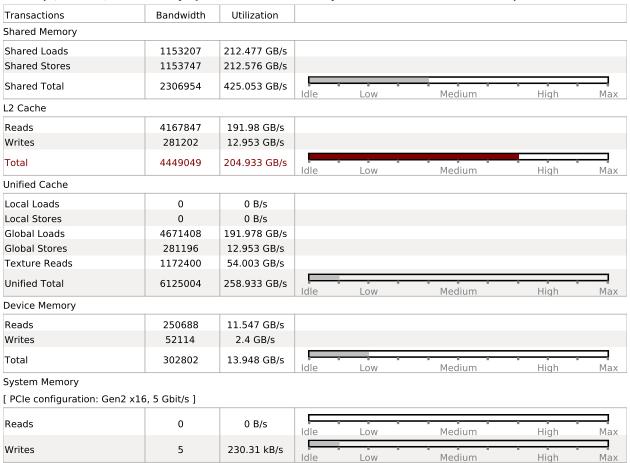
Shared Memory - If possible use 64-bit accesses to shared memory and 8-byte bank mode to achieved 2x throughput.

L2 Cache - Align and block kernel data to maximize L2 cache efficiency.

Unified Cache - Reallocate texture data to shared or global memory. Resolve alignment and access pattern issues for global loads and stores.

Device Memory - Resolve alignment and access pattern issues for global loads and stores.

System Memory (via PCIe) - Make sure performance critical data is placed in device or shared memory.



3. Instruction and Memory Latency

Instruction and memory latency limit the performance of a kernel when the GPU does not have enough work to keep busy. The results below indicate that the GPU does not have enough work because instruction execution is stalling excessively.

3.1. Kernel Profile - PC Sampling

The Kernel Profile - PC Sampling gives the number of samples for each source and assembly line with various stall reasons. Using this information you can pinpoint portions of your kernel that are introducing latencies and the reason for the latency. Samples are taken in round robin order for all active warps at a fixed number of cycles regardless of whether the warp is issuing an instruction or not

Instruction Issued - Warp was issued

Instruction Fetch - The next assembly instruction has not yet been fetched.

Execution Dependency - An input required by the instruction is not yet available. Execution dependency stalls can potentially be reduced by increasing instruction-level parallelism.

Memory Dependency - A load/store cannot be made because the required resources are not available or are fully utilized, or too many requests of a given type are outstanding. Data request stalls can potentially be reduced by optimizing memory alignment and access patterns.

Texture - The texture sub-system is fully utilized or has too many outstanding requests.

Synchronization - The warp is blocked at a __syncthreads() call.

Constant - A constant load is blocked due to a miss in the constants cache.

Pipe Busy - The compute resource(s) required by the instruction is not yet available.

Memory Throttle - Large number of pending memory operations prevent further forward progress. These can be reduced by combining several memory transactions into one.

Not Selected - Warp was ready to issue, but some other warp issued instead. You may be able to sacrifice occupancy without impacting latency hiding and doing so may help improve cache hit rates.

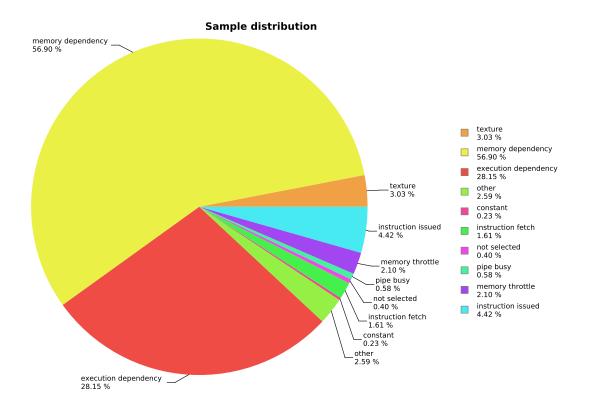
Other - The warp is blocked for a uncommon reason.

Examine portions of the kernel that have high number of samples to know where the maximum time was spent and observe the latency reasons for those samples to identify optimization opportunities.

Cuda Functions	Sample Count	% of Kernel Samples
void computeHOGSharedPred <float, float,="" int="64">(float*, float*, float*,</float,>	52526	100.0
float*, int, int, int)		

Source Files:

/home/adas/cuda-workspace/CudaVisionSysDeploy/Release//src/init//device/HOG/normalizeDescriptor.cuh
/home/adas/cuda-workspace/CudaVisionSysDeploy/Release//src/init//device/HOG/HOGdescriptor.cuh
/usr/local/cuda-7.5/targets/x86_64-linux/include/device_functions.hpp



3.2. GPU Utilization Is Limited By Shared Memory Usage

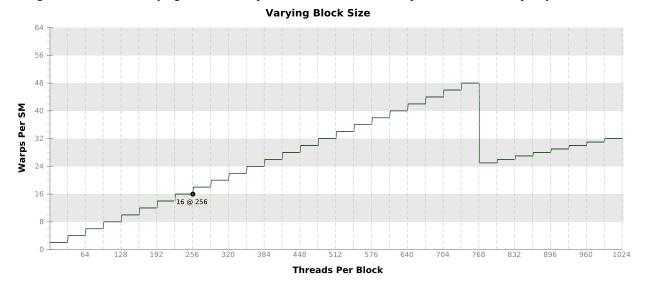
The kernel uses 36 KiB of shared memory for each block. This shared memory usage is likely preventing the kernel from fully utilizing the GPU. Device "GeForce GTX 960" is configured to have 96 KiB of shared memory for each SM. Because the kernel uses 36 KiB of shared memory for each block each SM is limited to simultaneously executing 2 blocks (16 warps). Chart "Varying Shared Memory Usage" below shows how changing shared memory usage will change the number of blocks that can execute on each SM.

Optimization: Reduce shared memory usage to increase the number of blocks that can execute on each SM. You can also increase the number of blocks that can execute on each SM by increasing the amount of shared memory available to your kernel. You do this by setting the preferred cache configuration to "prefer shared".

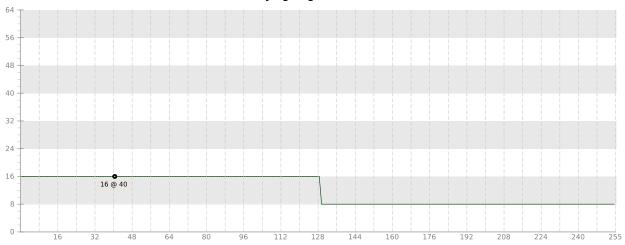
Variable	Achieved	Theoretical	Device Limit	Grid Si	ze: [3	1,1,1]	(31 bl	ocks)	Block	Size: [256,1	,1](2	56 thread
Occupancy Per SM													
Active Blocks		2	32	0	3	6	9 12	2 15	18	21	24	27	30 32
Active Warps	15.08	16	64	0	7	14	21	28	35	42	49	56	664
Active Threads		512	2048	0	256	512	2 76	8 1	024	1280	1536	179	2 2048
Occupancy	23.6%	25%	100%	0%		1	%		50%		75%	, D	100%
Warps													
Threads/Block		256	1024	0	128	256	38	34 5	12	640	768	89	5 1024
Warps/Block		8	32	0	3	6	9 12	2 15	18	21	24	27	30 32
Block Limit		8	32	0	3	6	9 12	2 15	18	21	24	27	30 32
Registers													
Registers/Thread		40	255	0	32	64	9	6 1	L28	160	192	22	4 255
Registers/Block		10240	65536	0		16	k		32k		48k		64k
Block Limit		6	32	0	3	6	9 12	2 15	18	21	24	27	30 32
Shared Memory													
Shared Memory/Block		36864	98304	0			32k			64	k		96k
Block Limit		2	32	0	3	6	9 12	2 15	18	21	24	27	30 32

3.3. Occupancy Charts

The following charts show how varying different components of the kernel will impact theoretical occupancy.

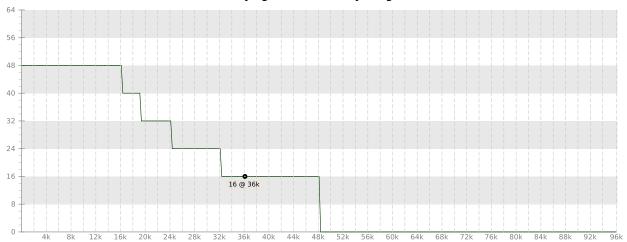


Varying Register Count



Registers Per Thread

Varying Shared Memory Usage



Shared Memory Per Block (bytes)

4. Compute Resources

GPU compute resources limit the performance of a kernel when those resources are insufficient or poorly utilized.

4.1. Kernel Profile - Instruction Execution

The Kernel Profile - Instruction Execution shows the execution count, inactive threads, and predicated threads for each source and assembly line of the kernel. Using this information you can pinpoint portions of your kernel that are making inefficient use of compute resource due to divergence and predication.

Examine portions of the kernel that have high execution counts and inactive or predicated threads to identify optimization opportunities.

Cuda Fuctions:

void computeHOGSharedPred<float, float*, float, int=8, int=8, int=16, int=64>(float*, float*, float*, float*, int, int, int, int)

Maximum instruction execution count in assembly: 62720

Average instruction execution count in assembly: 10397

Instructions executed for the kernel: 5177977

Thread instructions executed for the kernel: 165083799

Non-predicated thread instructions executed for the kernel: 142791205

Warp non-predicated execution efficiency of the kernel: 86.2%

Warp execution efficiency of the kernel: 99.6%

Source files:

/home/adas/cuda-workspace/CudaVisionSysDeploy/Release/../src/init/../device/HOG/normalizeDescriptor.cuh/home/adas/cuda-workspace/CudaVisionSysDeploy/Release/../src/init/../device/HOG/HOGdescriptor.cuh

/usr/local/cuda-7.5/targets/x86 64-linux/include/device functions.hpp

4.2. Function Unit Utilization

Different types of instructions are executed on different function units within each SM. Performance can be limited if a function unit is over-used by the instructions executed by the kernel. The following results show that the kernel's performance is not limited by overuse of any function unit.

Load/Store - Load and store instructions for shared and constant memory.

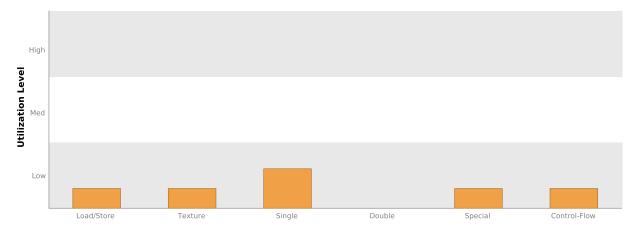
Texture - Load and store instructions for local, global, and texture memory.

Single - Single-precision integer and floating-point arithmetic instructions.

Double - Double-precision floating-point arithmetic instructions.

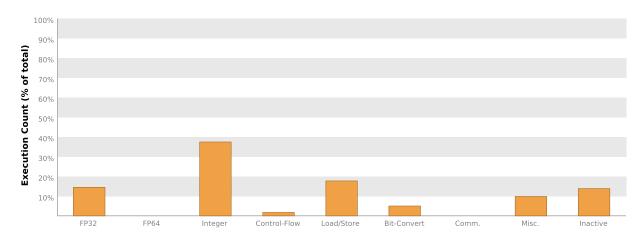
Special - Special arithmetic instructions such as sin, cos, popc, etc.

Control-Flow - Direct and indirect branches, jumps, and calls.



4.3. Instruction Execution Counts

The following chart shows the mix of instructions executed by the kernel. The instructions are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing instructions in that class. The "Inactive" result shows the thread executions that did not execute any instruction because the thread was predicated or inactive due to divergence.



4.4. Floating-Point Operation Counts

The following chart shows the mix of floating-point operations executed by the kernel. The operations are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing operations in that class. The results do not sum to 100% because non-floating-point operations executed by the kernel are not shown in this chart.

