

COFUNCTIONS, MONOIDS, AND SPLIT EPIMORPHISMS

AUSTRALIAN CATEGORY SEMINAR

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§0. MOTIVATION

A category is like a monoid with several objects.

But what is the correct generalisation of monoid homomorphisms?

Obvious answer: A functor between categories.

Thus monoids and homomorphisms form a full subcategory of Cat:

$$\boxed{\text{Mon}} \xleftarrow{\text{f.f.}} \boxed{\text{Cat}}$$

Remark: this inclusion functor has a left adjoint, which maps a small category to the free monoid on its set of morphisms modulo some relations on identities and composites; however no right adjoint exists as initial objects are not preserved.

Less obvious answer: A cofunctor between categories.

The category of small categories and cofunctors also has Mon as a full subcategory:

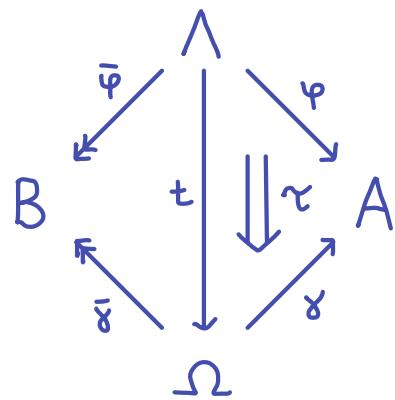
$$\boxed{\text{Mon}} \begin{array}{c} \xrightarrow{\text{f.f.}} \\ \perp \\ \xleftarrow{\Gamma} \end{array} \boxed{\text{Cof}}$$

This inclusion functor has a right adjoint, which maps a small category to its monoid of admissible sections.

Outline of the talk:

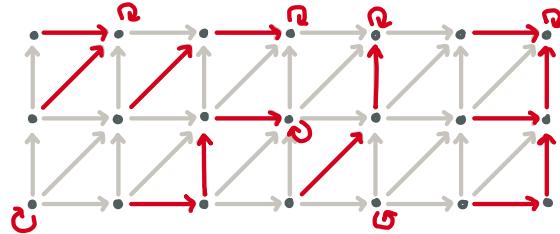
§1. The 2-category of categories and cofunctors

- What is a cofunctor?
- Examples
- Cofunctors as spans of functors
- Morphisms of cofunctors



§2. A right adjoint: the monoid of admissible sections

- An explicit description (vector fields on categories)



- Examples
- Characterisation as a hom category
- An internal perspective

§3. Lenses and split epimorphisms of monoids

- Motivation: generalising semidirect products of groups
- Schreier split epimorphisms between monoids

$$K[f] \xleftarrow[\kappa]{\alpha} A \xleftarrow[f]{\psi} B$$

- Lenses and split opfibrations between monoids

§1. The 2-category of categories and cofunctors

Definition (Higgins, Mackenzie ; Aguiar): A **cofunctor** $\varphi: B \rightarrow A$ between small categories consists of a pair of functions,

$$A_0 \xrightarrow{\varphi_0} B_0 \\ a \longmapsto \varphi_0 a$$

assignment on objects

$$A_0 \times_{B_0} B_1 \xrightarrow{\varphi_1} A_1 \\ (a, \varphi_0 a \xrightarrow{u} b) \longmapsto \varphi_1(a, u): a \rightarrow a'$$

assignment on morphisms

subject to the following axioms:

$$(1) \varphi_0 \text{cod}(\varphi_1(a, u)) = \text{cod}(u)$$

$$(2) \varphi_1(a, 1_{\varphi_0 a}) = 1_a$$

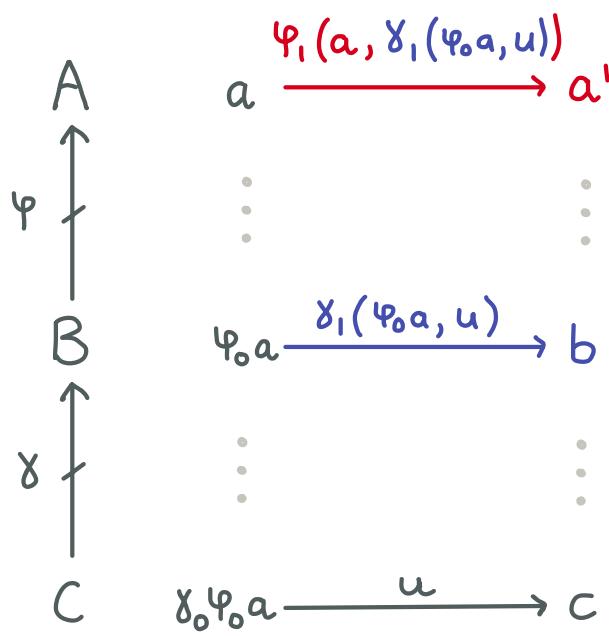
$$(3) \varphi_1(a, v \circ u) = \varphi_1(a', v) \circ \varphi_1(a, u)$$

where $a' = \text{cod } \varphi_1(a, u)$

$$\begin{array}{ccc} A & a & \xrightarrow{\varphi_1(a, u)} a' \\ \uparrow \varphi & \vdots & \vdots \\ B & \varphi_0 a & \xrightarrow{u} b = \varphi_0 a' \end{array}$$

"A cofunctor is a lifting of morphisms against objects, which respects identities and composition"

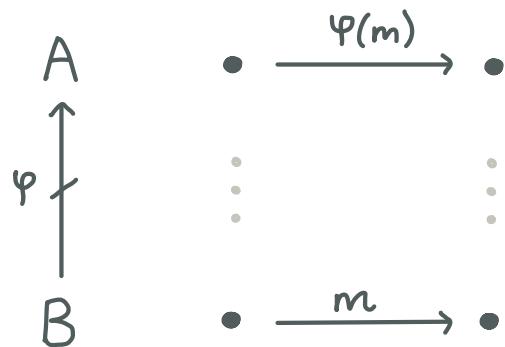
Given cofunctors $\gamma: C \rightarrow B$ and $\varphi: B \rightarrow A$, we can take their composite in the following way:



This yields a category Cof of categories and cofunctors.

Examples:

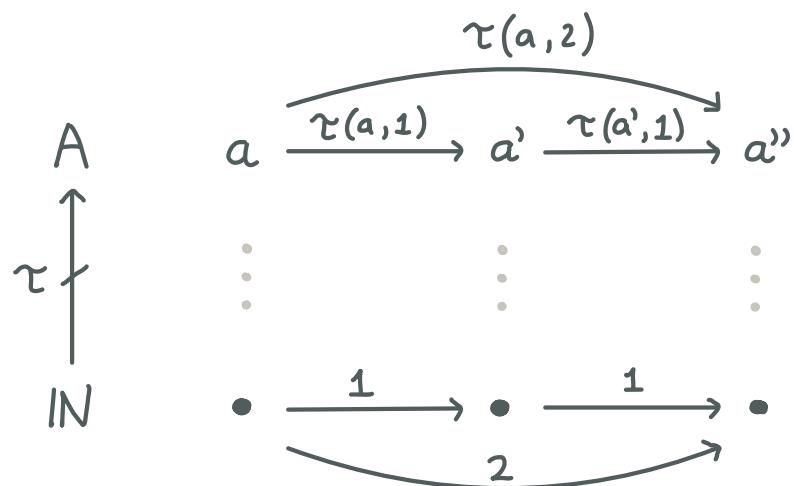
- Every monoid homomorphism $B \xrightarrow{\varphi} A$ yields a cofunctor:



This gives a fully faithful functor $\text{Mon} \longrightarrow \text{Cof}$

- More generally, every bijective-on-objects functor $B \rightarrow A$ yields a cofunctor $B \rightarrow A$.
- Every discrete opfibration $A \rightarrow B$ yields a cofunctor $B \rightarrow A$.

- Every **split opfibration** has an underlying cofunctor given by the splitting.
- More generally, every **delta lens** has an underlying cofunctor.
- Let \mathbb{N} denote the monoid of natural numbers under addition. A cofunctor $\mathbb{N} \xrightarrow{\tau} A$ is the same as a choice of morphism out of every object in A :



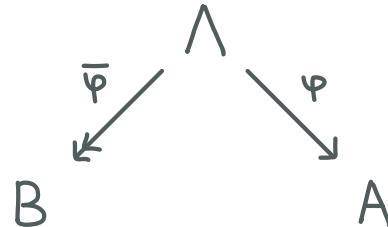
Internally this is the same as a **section** of the domain map:

$$\begin{array}{ccc}
 A_0 & \xrightarrow{\tau} & A_1 \\
 \searrow 1 & & \swarrow d_0 \\
 & A_0 &
 \end{array}$$

- For every category A , there is a unique cofunctor $\mathbb{1} \rightarrow A$ from the category with a single identity arrow. That is, $\mathbb{1}$ is the **initial object** in Cof .

Proposition: Every cofunctor $B \xrightarrow{\varphi} A$ may be represented as a span of functors,

use \longrightarrow for discrete opfibrations



where $\bar{\varphi}$ is a **discrete opfibration** and φ is **identity-on-objects**.

This representation extends to a functor $\text{Cof} \rightarrow \text{Span}_{\text{iso}}(\text{Cat})$, and the pair $(\text{DOpf}^{\text{op}}, \text{Bij})$ is an **orthogonal factorisation system** on Cof .

PROOF (SKETCH): Given $\varphi: B \rightarrow A$, let Λ be the category with the same objects as A and morphisms given by formal pairs $(a, u: \varphi_0 a \rightarrow b)$. Then we have:

$$\begin{array}{ccccc} B & \xleftarrow{\bar{\varphi}} & \Lambda & \xrightarrow{\varphi} & A \\ \varphi_0 a & & a & & a \\ u \downarrow & & \downarrow (a, u) & & \downarrow \varphi_1(a, u) \\ b & & a' & & a' \end{array}$$

where $a' = \text{cod } \varphi_1(a, u)$.

Internally this span of functors is given by the diagram:

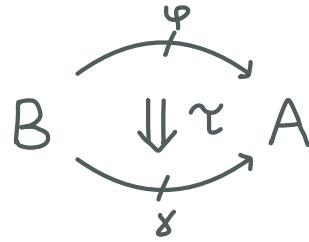
where

$$\Lambda_1 := A_0 \times_{B_0} B_1$$

$$\begin{array}{ccccc} B_0 & \xleftarrow{\varphi_0} & A_0 & \xrightarrow{1} & A_0 \\ d_0 \uparrow & & \uparrow \pi_0 & & \uparrow d_0 \\ B_1 & \xleftarrow{\pi_1} & \Lambda_1 & \xrightarrow{\varphi_1} & A_1 \\ d_1 \downarrow & & \downarrow d_1 \varphi_1 & & \downarrow d_1 \\ B_0 & \xleftarrow{\varphi_0} & A_0 & \xrightarrow{1} & A_0 \end{array}$$

□

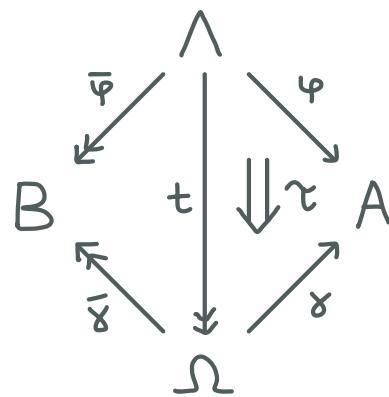
Definition (Aguiar): A **2-cell** between cofunctors,



is an assignment to each object $a \in A$ of a morphism $\tau_a: a \rightarrow a'$ in A such that $\varphi \circ a = \gamma \circ a'$ and for any pair $(a, u: \varphi \circ a \rightarrow b)$ the following diagram in A commutes:

$$\begin{array}{ccc}
 & a \xrightarrow{\varphi_i(a,u)} a'' & \text{where } a' = \text{cod}(\tau_a) \\
 & \tau_a \downarrow \quad \lrcorner \quad \downarrow \tau_{a''} & a'' = \text{cod}(\varphi_i(a,u)) \\
 a' \xrightarrow{\gamma_i(a',u)} a''' & \vdots & a''' = \text{cod}(\tau_{a''}) \\
 & \vdots & = \text{cod}(\gamma_i(a',u)) \\
 \varphi \circ a \xrightarrow{u} b & = & \varphi \circ a'' \\
 \parallel \qquad \qquad \qquad \qquad \parallel & & \\
 \gamma \circ a' \xrightarrow{u} \gamma \circ a''' & &
 \end{array}$$

Proposition: Every 2-cell between cofunctors may be represented as a diagram in Cat :



□

Categories, cofunctors and 2-cells form a **2-category** Cof .

Vertical composition of 2-cells may be understood via pasting in Cat :

$$\begin{array}{ccc}
 \begin{array}{c}
 \begin{array}{ccccc}
 B & \xleftarrow{\bar{\varphi}} & X & \xrightarrow{\varphi} & A \\
 \parallel & & \downarrow t & & \parallel \\
 B & \xleftarrow{\bar{\gamma}} & Y & \xrightarrow{\gamma} & A \\
 \parallel & & \downarrow s & & \parallel \\
 B & \xleftarrow{\bar{\psi}} & Z & \xrightarrow{\psi} & A
 \end{array}
 & = &
 \begin{array}{ccccc}
 B & \xleftarrow{\bar{\varphi}} & X & \xrightarrow{\varphi} & A \\
 \parallel & & \downarrow st & & \parallel \\
 B & \xleftarrow{\bar{\varphi}} & Z & \xrightarrow{\psi} & A
 \end{array}
 \\[10pt]
 \downarrow \varphi \circ \gamma = \sigma \cdot \tau
 \end{array}
 \end{array}$$

Composition of cofunctors corresponds to span composition:

$$\begin{array}{ccccc}
 & \Omega \times_B \Lambda & & & \\
 & \swarrow \pi_0 \quad \vee \quad \searrow \pi_1 & & & \\
 & \Omega & & \Lambda & \\
 & \swarrow \bar{\gamma} \quad \searrow \alpha & & \swarrow \bar{\varphi} \quad \searrow \varphi & \\
 C & & B & & A
 \end{array}$$

Horizontal composition of 2-cells is more involved.

$$\begin{array}{ccc}
 \begin{array}{ccccc}
 W & & X & & A \\
 \swarrow \bar{\psi} \quad \searrow \psi & & \swarrow \bar{\varphi} \quad \searrow \varphi & & \\
 C & \xleftarrow{s} & B & \xleftarrow{t} & A \\
 \nearrow \bar{\kappa} \quad \searrow \kappa & & \nearrow \bar{\chi} \quad \searrow \chi & & \\
 Z & & Y & &
 \end{array}
 & = &
 \begin{array}{ccccc}
 W \times_B X & & & & A \\
 \swarrow \bar{\psi} \pi_W \quad \searrow \psi \pi_X & & \downarrow r & & \\
 C & & Z \times_B Y & & A \\
 \nearrow \bar{\kappa} \pi_Z \quad \searrow \kappa \pi_Y & & & & \\
 & & & &
 \end{array}
 \end{array}$$

To obtain the functor r and natural transformation ρ , first we must use the universal property of discrete opfibrations:

$$\begin{array}{ccc}
 W \times_B X & \xrightarrow{\pi_X} & X \\
 \downarrow \pi_W & \Downarrow \sigma \pi_W & \downarrow \bar{\varphi} \\
 W & \xrightarrow{s} & Z \xrightarrow{\kappa} B
 \end{array} =
 \begin{array}{ccc}
 W \times_B X & \xrightarrow{\pi_X} & X \\
 \downarrow \pi_W & \Downarrow \theta & \downarrow \bar{\varphi} \\
 W & \xrightarrow{\kappa s} & B
 \end{array}$$

Then the functor r is constructed by the universal property of the pullback,

$$\begin{array}{ccccc}
 & & \text{th} & & \\
 & W \times_B X & \dashrightarrow & Z \times_B Y & \xrightarrow{\pi_Y} Y \\
 \pi_W \downarrow & & r & \downarrow \pi_Z & \downarrow \bar{\gamma} \\
 W & \xrightarrow{s} & Z & \xrightarrow{\kappa} B &
 \end{array}$$

while the natural transformation ρ is the horizontal composite:

$$\begin{array}{ccc}
 W \times_B X & \xrightarrow{\pi_X} & X \\
 \Downarrow \theta & \curvearrowright & \Downarrow \tau \\
 h & & \gamma t
 \end{array}$$

Note that for \mathcal{E} with pullbacks, we also have a 2-category $\text{Cof}(\mathcal{E})$, where a 2-cell $\varphi \xrightarrow{\tau} \gamma: B \rightarrow A$ may be specified by:

$$\begin{array}{ccc}
 & A_0 & \xrightarrow{\varphi_0} B_0 \\
 & \downarrow \tau & \uparrow \gamma_0 \\
 A_0 & \xleftarrow{d_0} & A_1 \xrightarrow{d_1} A_0
 \end{array}
 \quad + \text{naturality condition}$$

§2. A right adjoint: the monoid of admissible sections

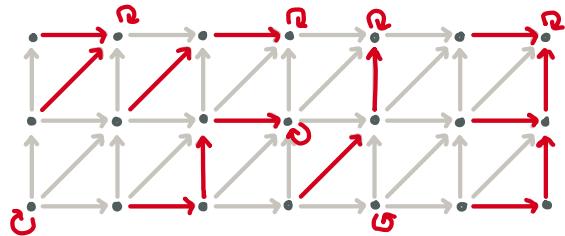
We would like to construct a **right adjoint** to the inclusion.

$$\boxed{\text{Mon}} \begin{array}{c} \xleftarrow{\quad \perp \quad} \\[-1ex] \xleftarrow{\Gamma} \end{array} \begin{array}{c} \xrightarrow{\quad \text{f.f.} \quad} \\[-1ex] \perp \end{array} \boxed{\text{Cof}}$$

How can we construct a monoid from a category?

IDEA: The elements of the monoid are like **vector fields** on the cat.

Example (Spivak, Myers): Consider a full subcategory of $(\mathbb{N}, \leq) \times (\mathbb{N}, \leq)$:



At each object, choose an outgoing arrow to get a family: $(a \xrightarrow{\tau_a} t(a))_{a \in A}$

Definition: The **monoid of admissible sections** $\Gamma(A)$ for a category A is a monoid whose:

- elements are families of morphisms $(a \xrightarrow{\tau_a} t(a))_{a \in A}$
- unit is the family of identities $(a \xrightarrow{1_a} a)_{a \in A}$
- multiplication is given by:

$$(a \xrightarrow{\tau_a} t(a))_{a \in A} * (a \xrightarrow{\sigma_a} s(a))_{a \in A} = (a \xrightarrow{\sigma_{t(a)} \circ \tau_a} st(a))_{a \in A}$$

We will see there are several equivalent characterisations of this monoid.

Note that Γ defines the object part of the right adjoint:

$$\boxed{\text{Set}} \begin{array}{c} \xleftarrow{\quad \perp \quad} \\[-1ex] \xleftarrow{\Gamma} \end{array} \boxed{\text{Mon}} \begin{array}{c} \xleftarrow{\quad \perp \quad} \\[-1ex] \xleftarrow{\Gamma} \end{array} \begin{array}{c} \xrightarrow{\quad \text{f.f.} \quad} \\[-1ex] \perp \end{array} \boxed{\text{Cof}}$$

elements of
the right
adjoint

$$1 \downarrow \Gamma(A)$$

$$\mathbb{N} \downarrow \Gamma(A)$$

$$\mathbb{N} \downarrow \begin{array}{l} \text{admissible} \\ \text{sections} \\ \text{of } A \end{array} \downarrow A$$

Examples

- If A_0 is a discrete category, $\Gamma(A_0) \cong \mathbb{1}$
- If A is a codiscrete category, $\Gamma(A) \cong \text{End}(A_0) = \text{Set}(A_0, A_0)$
- If M is a monoid, then $\Gamma(M) \cong M$
- For any category A , the commutative monoid of endomorphisms of the identity transformation $\text{End}(1_A) = \text{Nat}(1_A, 1_A)$ is a submonoid of the centre $Z(\Gamma(A))$ of $\Gamma(A)$.
- (Garner): The group of extended inner automorphisms of a category $A \in \text{Cof}$ is isomorphic to $\text{Bis}(A)$, the group of bisections, which is the same as the group of invertible elements of $\Gamma(A)$.
- (Aguiar) Recall that a category in Grp is the same as a crossed module of groups. The monoid of admissible sections of such a category is equivalent to Whitehead's monoid of derivations of the crossed module.
- The 2007 paper "External derivations of internal groupoids" by Kasangian, Mantovani, Metere, and Vitale also studies this monoid.
- Every 2-cell $\tau: \varphi \Rightarrow \gamma: B \rightarrow A$ yields an element of $\Gamma(A)$.

Proposition: The monoid of admissible sections is isomorphic to the hom category:

$$\Gamma(A) := \text{Cof}(\mathbb{1}, A)$$

PROOF: Consider a morphism in $\text{Cof}(\mathbb{1}, A)$ given by:

$$\begin{array}{ccccc}
 & & A_0 & & \\
 & \swarrow ! & \downarrow & \searrow i & \\
 1 & & t & \Downarrow \tau & A \\
 & \nwarrow ! & \downarrow & \nearrow i & \\
 & & A_0 & &
 \end{array}$$

Since $\mathbb{1}$ is initial in Cof , the category $\text{Cof}(\mathbb{1}, A)$ is a monoid. The components of the natural transformation τ give exactly a family of morphisms $(a \xrightarrow{\tau_a} t(a))_{a \in A}$, that is, an element of $\Gamma(A)$. Moreover, vertical composition of 2-cells corresponds exactly to composition in $\Gamma(A)$. \square

Corollary: There is a functor $\Gamma: \text{Cof} \longrightarrow \text{Mon}$.

PROOF: Given a cofunctor $B \xrightarrow{\varphi} A$, we get a morphism of monoids by whiskering $\text{Cof}(\mathbb{1}, B) \xrightarrow{\varphi} \text{Cof}(\mathbb{1}, A)$. \square

Note, the statement that Γ is a right adjoint translates to:

$$\text{Cof}(M, A) \cong \text{Mon}(M, \text{Cof}(\mathbb{1}, A))$$

Is this a copower?

Proposition: There is an adjunction of categories:

$$\boxed{\text{Mon}} \quad \begin{array}{c} \xleftarrow{\perp} \\[-1ex] \xleftarrow{\Gamma} \end{array} \quad \boxed{\text{Cof}}$$

PROOF: The unit of the adjunction is a natural isomorphism, and the **counit** can be constructed as the cofunctor $\Gamma(A) \rightarrow A$:

$$\begin{array}{ccc} A & & x \xrightarrow{\tau_x} t(x) \\ \uparrow \varepsilon_A & \vdots & \vdots \\ \Gamma(A) & \bullet \xrightarrow{\quad} \bullet & (a \xrightarrow{\tau_a} t(a))_{a \in A} \end{array}$$

In other words, the counit "**evaluates**" the family at an object $x \in A$ to get a morphism $\tau_x \in A$. One may show the required identities for an adjunction hold. \square

Aside: Recall admissible sections (morphisms in $\text{Cof}(\mathbb{1}, A)$) are the same as cofunctors $\text{IN} \rightarrow A$ (objects in $\text{Cof}(\text{IN}, A)$). What is the relationship between these categories?

Let M be a monoid considered as a one-object category. Then we may construct a category \hat{M} whose objects are morphisms in M (elements of the monoid), and whose hom sets are given by $\hat{M}(f, g) = \{m \in M \mid m \circ f = g \circ m\}$. When $M = \text{Cof}(\mathbb{1}, A)$, then $\hat{M} \cong \text{Cof}(\text{IN}, A)$. Does this construction have a universal property?

Internal perspective

Why is $\Gamma(A)$ called the monoid of admissible **sections**?

For a category \mathcal{E} with pullbacks (think $\mathcal{E} = \text{Set}$) we can define a functor $\Gamma : \text{Cof}(\mathcal{E}) \rightarrow \text{Mon}$ which takes an internal category A to the monoid $\Gamma(A)$ whose underlying set is:

$$\Gamma(A) = \{\sigma : A_0 \rightarrow A_1 \mid d_0 \circ \sigma = 1_{A_0}\} = \{\text{sections to the domain map } d_0\}$$

Given a cofunctor $B \xrightarrow{\varphi} A$, we can obtain a morphism of monoids $\Gamma(B) \rightarrow \Gamma(A)$ in the following way:

$$\begin{array}{ccccc}
 \Gamma(B) & \xrightarrow{\varphi_0} & \Gamma(A) & \xrightarrow{\varphi_1} & \Gamma(A) \\
 B_0 \xrightarrow{\sigma} B_1 & & A_0 \xrightarrow{1 \times \sigma \varphi_0} A_1 & & A_0 \xrightarrow{\varphi_1(1 \times \sigma \varphi_0)} A_1 \\
 \downarrow 1 & & \downarrow 1 & & \downarrow 1 \\
 B_0 & & A_0 & & A_0
 \end{array}$$

This is (more or less) how Aguiar originally defined Γ . However, unlike the previous definition, Γ no longer appears as a right adjoint to the inclusion $\text{Mon}(\mathcal{E}) \rightarrow \text{Cof}(\mathcal{E})$.

Can we fix this?

Assumption: Let \mathcal{E} be a category with finite limits such that for every object $B_0 \in \mathcal{E}$, the functor

$$\mathcal{E} \xrightarrow{B_0 \times (-)} \mathcal{E}/B_0$$

has a chosen right adjoint $\Gamma_{B_0} : \mathcal{E}/B_0 \rightarrow \mathcal{E}$.

Proposition: Under the assumption above, $\text{Mon}(\mathcal{E})$ is a coreflective subcategory of $\text{Cof}(\mathcal{E})$.

$$\boxed{\text{Mon}(\mathcal{E})} \begin{array}{c} \xleftarrow{\quad \perp \quad} \\[-1ex] \xrightleftharpoons[\Gamma]{\text{f.f.}} \end{array} \boxed{\text{Cof}(\mathcal{E})}$$

PROOF: Consider an internal category $B \in \text{Cof}(\mathcal{E})$. Define:

$$\Gamma(B) := \Gamma_{B_0} \left(B_1 \xrightarrow{d_0} B_0 \right) \in \mathcal{E}$$

Since $B_0 \times (-) \dashv \Gamma_{B_0}$, we have a bijection between:

$$Z \xrightarrow{\hat{\sigma}} \Gamma(B) \quad \text{↔} \quad \begin{array}{ccc} B_0 \times Z & \xrightarrow{\sigma} & B_1 \\ \pi_0 \searrow & & \downarrow d_0 \\ & & B_0 \end{array}$$

In particular, when $Z=1$, we have a correspondence between "elements" $1 \rightarrow \Gamma(B)$ and sections of $d_0: B_1 \rightarrow B_0$.

We can show that $\Gamma(B)$ is an internal monoid in \mathcal{E} using the diagrams for B as an internal category in \mathcal{E} . For example, the unit of the monoid comes from the identity map:

$$B_0 \xrightarrow{i} B_1 \quad \text{↔} \quad 1 \xrightarrow{e} \Gamma(B)$$

$$\begin{array}{ccc} B_0 & \xrightarrow{1} & B_0 \\ & \searrow d_0 & \end{array}$$

Multiplication for the monoid (plus the axioms) also follow.

To see functoriality of Γ , consider a cofunctor $(\varphi_0, \varphi_1) : B \rightarrow A$.

$$\begin{array}{ccccc}
 & \langle \pi_0, \varepsilon(\varphi_0 \times 1) \rangle & & & \\
 A_0 \times \Gamma(B) & \xrightarrow{\quad} & \wedge_1 & \xrightarrow{\varphi_1} & A_1 \\
 & \searrow \pi_0 & \downarrow \pi_0 & \nearrow d_0 & \\
 & & A_0 & &
 \end{array}
 \quad \sim \sim \quad
 \begin{array}{ccc}
 \Gamma(B) & \xrightarrow{\Gamma_\varphi} & \Gamma(A)
 \end{array}$$

One may check that the monoid homomorphism axioms hold.

To see that the **unit** of the adjunction is an isomorphism, note that for a monoid M , we have a bijection:

$$\begin{array}{ccc}
 Z & \longrightarrow & \Gamma(M) & \leftarrow \sim \sim & Z & \longrightarrow & M \\
 & & & & \downarrow ! & & \downarrow !
 \end{array}$$

We can easily show that $M \cong \Gamma(M)$ follows.

Finally the **counit** of the adjunction at an internal category B is the internal cofunctor (given as a span of internal functors) below:

$$\begin{array}{ccccc}
 1 & \xleftarrow{!} & B_0 & \xrightarrow{1} & B_0 \\
 \uparrow & & \uparrow \pi_0 & & \uparrow d_0 \\
 \Gamma(B) & \xleftarrow{\pi_1} & B_0 \times \Gamma(B) & \xrightarrow{\varepsilon} & B_1 \\
 \downarrow & & \downarrow d_1, \varepsilon & & \downarrow d_1 \\
 1 & \xleftarrow{!} & B_0 & \xrightarrow{1} & B_0
 \end{array}$$

ε is the counit
of the adjunction
 $B_0 \times (-) \rightarrow \Gamma_{B_0}$

□

SUMMARY (IF SHORT ON TIME)

- In §1 we saw that there is a 2-category Cof of small categories, **cofunctors**, and 2-cells.
- Furthermore, we saw that cofunctors and their 2-cells could be understood via **diagrams in Cat** .
- In §2 we proved that the category of monoids is a **coreflective subcategory** of Cof :

$$\begin{array}{ccc} \boxed{\text{Mon}} & \begin{matrix} \xleftarrow{\quad \perp \quad} \\ \xrightarrow{\quad \text{f.f.} \quad} \end{matrix} & \boxed{\text{Cof}} \\ \Gamma & & \end{array}$$

- We saw that the right adjoint, which takes a category to its **monoid of admissible sections**, has several characterisations including as the hom category $\Gamma(-) = \text{Cof}(\mathbb{1}, -)$.
- Moreover, this adjunction generalises to the **internal setting** with suitable assumptions on \mathcal{E} .

DIRECTIONS FOR FUTURE WORK

- It is possible to construct a **double category** of functors and cofunctors from Span using Pare's **retrocells**. What can we learn from this perspective? What properties of Span allow us to represent cofunctors as spans?
- Replacing \mathcal{E} with a suitable **monoidal category** V , we can generalise the above adjunction to involve monoids and categories in V (this is what Aguiar did) and obtain string diagram proofs.
- Are there **other examples** of $\Gamma(A)$ in the literature? What interesting information about A does it contain? Applications?