

Double cats as pseudomonads

↪ app. double profunctors

Plan

1. Review 1D

- categories : monads
- profunctors : bimodules
- functors / cofunctors : certain special cases

2. Extend definitions to 2D

- PsDblCat (recover notions known in the literature)
- PsDblCof (introduced M. 2023)
- PsDblProf (WR jw. N. Gambino)

3. Grothendieck construction

profunctors

$A^{\text{op}} \times B \rightarrow \text{set}$

double
profunctors

???

4. Strictification

- double functors
- double cofunctors
- double profunctors

(GrardS, Paré 1999)
(Campbell 2019)

(M. 2023)

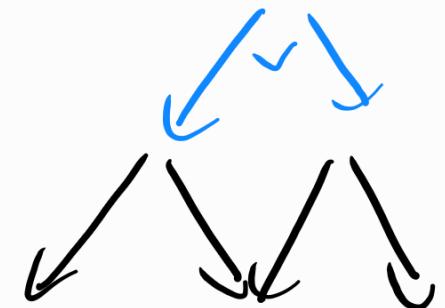
(M. 2024.)

Part 1 Review 1D case

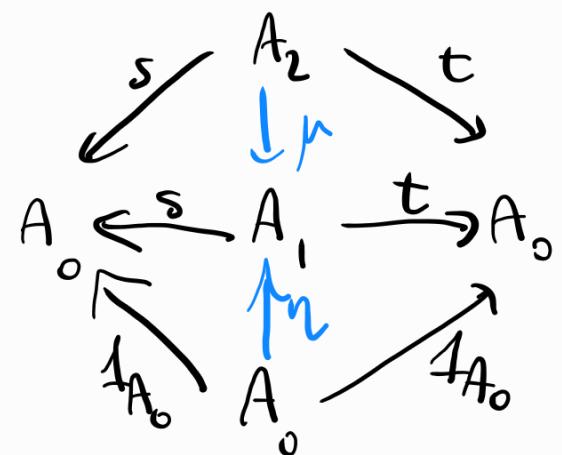
Bénabou 1967 , 1973

Span

Composition
is via pullback



monads
are categories



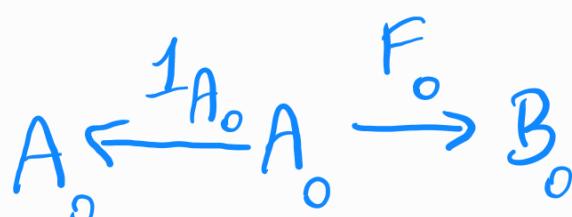
Street 1972

Formal Theory of Monads

Morphisms
of monads



cofunctors



functors

Co-morphisms
of monads



2-cells



equalities

Lack-Street 2002

KL & EM

(co)completions

More general 2-cells which correspond to natural transformations
(& 2-cells
bw cofunctors)

Garner Shulman 2016

Joyal Gambino 2017

Bimod (K)

Bimodules
of monads



profunctors
between cats

Morphisms of
bimodules



Natural transf.
between profunctors

Bimodules

$$A_0 \leftarrow P_i \rightarrow B_0$$

$P_i(x, y)$

heteromorphisms

$$\begin{array}{ccc} & P_i & \\ A_0 & \uparrow P & \rightarrow B_0 \\ A & \downarrow & \\ A_0 & \xrightarrow{P_i} & \end{array}$$

$$P_i \times A_1 \xrightarrow{g} P_i$$

$$\begin{array}{ccc} A_0 & \xrightarrow{P_i} & B_0 \\ \downarrow P_i & \searrow \text{A} & \downarrow B \\ B_0 & & \end{array}$$

$$B_1 \times P_i \xrightarrow{\text{A}} P_i$$

$$A^{\text{op}} \times B \xrightarrow{P} \text{Set}$$

algebras for the distributive law between $B^0(-) \& (-) \circ A_0$,

on $\text{Span}(\text{Set})(A_0, B_0)$

Composition of bimodules

$$A \xrightarrow{P} B \xrightarrow{Q} C$$

... is given by the reflexive coequaliser

$$\begin{array}{ccc} Q \circ B \circ P & \xrightarrow{\quad \beta^{Q \circ 1} \quad} & Q \circ P \\ 1 \circ \alpha^P & \longrightarrow & \longrightarrow Q \circ_B P \end{array}$$

Cofunctors and Functors

\mathfrak{F} free on $A_0 = A_0 \xrightarrow{F_0} B_0$

\mathfrak{A} free on $A_0 = A_0 \xrightarrow{F_0} B_0$

Part 2.

Tricategory

Span (Cat)



$$R'\phi = R$$

$$L'\phi = L$$

Gambino, Lobbia 2021

Formal Theory of Pseudomonads pt 1.

pseudomonads
in Span (Cat) \longleftrightarrow (loosely pseudo)
double cat's

$$(hg)f \begin{array}{c} \cong \\ \downarrow \end{array} f h(gf) \quad f \perp \begin{array}{c} \cong \\ \downarrow \end{array} \begin{array}{c} \cong \\ \downarrow \end{array} f \quad f \begin{array}{c} \cong \\ \downarrow \end{array} \begin{array}{c} \cong \\ \downarrow \end{array} f$$

(α -morphisms)

$$A_0 \xleftarrow{1_{A_0}} A_0 \xrightarrow{F} B_0$$

(loosely pseudo)

double functors

2-cells



(tightly natural,
loosely pseudonat.)

double n.t w/

identity components.

3-cells



equalities

M. 2023

Formal Theory of
Pseudomonads pt 2.

More general



loose ps. nat

2-cells

transformations

More general



double

3-cells

modifications

Morphisms of pseudomonads give a notion of (loosely pseudo) double cofunctors.

$$A_0 \xrightarrow[\text{functor}]{} B_0$$

$$\begin{array}{ccc}
 X & \xrightarrow{x} & X' \\
 F_x f \downarrow & & \downarrow F_{x'} f' \\
 Y & \xrightarrow{y} & Y'
 \end{array}
 \quad
 \begin{array}{ccc}
 F_0 x & \xrightarrow{\beta} & F_0 x' \\
 F_0 f \downarrow & & \downarrow f' \\
 Y & \xrightarrow{y} & Y'
 \end{array}$$

respecting loose composition + identities up to coherent inv. double cells.

General 2-cells/3-cells give a notion of double pseudonatural transformation / double comodification

$$\begin{array}{ccccc}
 & & \xrightarrow{\quad F_x f \quad} & & \xrightarrow{\quad \phi_x \quad} \\
 \begin{matrix} x \\ \downarrow \phi_x \\ x' \end{matrix} & \xrightarrow{x} & \begin{matrix} y \\ \downarrow \phi_y \\ y' \end{matrix} & \begin{matrix} \xrightarrow{\quad \phi_{y'} \quad} \\ \downarrow \phi_y \\ y' \end{matrix} & \begin{matrix} \xrightarrow{\quad \phi_{x'} \quad} \\ \downarrow \phi_{x',f} \\ y' \end{matrix} \\
 & \xrightarrow{x'} & & &
 \end{array}$$

Th^m (M. 2023)

These data assemble into a tricategory PDbGof .

N.B. New even for bicategories.

What about double profunctors ?

Problem :

$$\text{Cat}/B \xrightarrow{F^*} \text{Cat}/A$$

does not preserve

coequalisers, but

Propⁿ (M. 2024)

F^* does preserve
codescent objects.

WIP j.w N. Gambino :

If \mathcal{K} is a Gray-cat
whose homs have
codescent objects separately
preserved by composition
then there is a tricategory

$\text{Bimod}(\mathcal{K})$.

objects : pseudomonads,

hom : pseudoalgebras for
pseudo-distributive law,

Composition : codescent
objects.

Part 3 Grothendieck construction.

$$A_0 \leftarrow P_i \rightarrow B_0$$

(loose) heteromorphisms,

hetero-doublecells

$$\begin{array}{ccc} X & \xrightarrow{P} & Y \\ f \downarrow & \swarrow \hat{P} & \downarrow g \\ X' & \xrightarrow{P'} & Y' \end{array}$$

$$P_i \rightarrow A_0 \times B_0$$

$A_0 \times B_0$ \rightsquigarrow Prof
normal
lax

P_i has actions on either side by A & B :

$$P_i \times A_1 \xrightarrow[r]{A_0} P_i \quad \left\{ \begin{array}{l} \text{functors} \\ \end{array} \right.$$

$$B_1 \times P_i \xrightarrow[l]{B_0} P_i$$

$$\left(\begin{matrix} X \\ f \\ \downarrow \\ X' \end{matrix}, Y \right) \mapsto \begin{matrix} P(X, Y) \\ P(f, Y) \uparrow \\ \text{functor} \\ P(X', Y) \end{matrix}$$

$$\left(\begin{matrix} X \\ \downarrow g \\ Y' \end{matrix}, Y \right) \mapsto \begin{matrix} P(X, Y) \\ P(X, g) \downarrow \\ \text{functor} \\ P(X, Y') \end{matrix}$$

The 2-cells of spans r & l
underlie pseudoalgebras for

$$(-) \circ A, \quad B \circ (-).$$

$$\begin{array}{ccc}
P_i \times A_1 \times A_1 & \xrightarrow{r \times 1} & P_i \times A_1 \\
\downarrow \mu & \cong \chi & \downarrow r \\
P_i \times A_1 & \xrightarrow{r} & P_i
\end{array}
\qquad
\begin{array}{ccc}
P_i & \xrightarrow{1 \times l} & P_i \times A_1 \\
& \searrow \cong \psi & \downarrow r \\
& 1_{P_i} & P_i
\end{array}$$

+ pseudoalgebra axioms

similar for l .

Thus we have pseudofunctors

$$\begin{array}{ccc} A^{\text{op}} & \xrightarrow[\text{loose}]{} & P(-, Y) \\ & \curvearrowright \text{Cat} & \end{array} , \quad \begin{array}{ccc} B & \xrightarrow[\text{loose}]{} & P(X, -) \\ & \curvearrowright \text{Cat} & \end{array}$$

The bimodularity condition

is mediated by an invertible natural transformation.

$$\begin{array}{ccccc} & & B_1 \times P_1 \times A_1 & & \\ & \swarrow l \times 1 & & \searrow 1 \times r & \\ P_1 \times A_1 & \cong & & & B_1 \times P_1 \\ \searrow r & & & & \swarrow l \\ & P_1 & & & \end{array}$$

Th^M (M. 2024)

These data comprise a double functor

$$\left(A^{\text{op}} \times B \right)^T \rightsquigarrow \underline{\text{Prof}}$$

that is (loosely) normal lax,

and (tightly) pseudo.

N.B., This formulation is hard to guess without the pseudomonad perspective on double categories!

Composition of 1D profunctors

$$\underbrace{\left\{ (X \xrightarrow{P} Y, Y \xrightarrow{q} Z) \right\}}$$

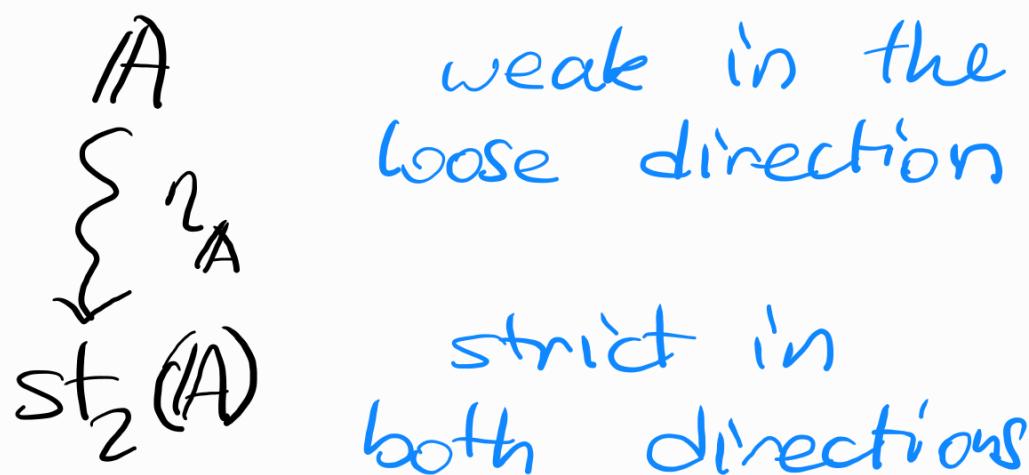
$$\left\langle (p, q) \sim (p', q') \mid \begin{array}{c} Y \\ \downarrow y \\ Y' \end{array} \quad (y p, q) = (p', q' y) \right\rangle$$

Composition of double profunctors

$$\begin{array}{ccccc} X & \xrightarrow{y P} & Y' & \xrightarrow{q} & Z \\ \parallel & & \cong & & \parallel \\ X & \xrightarrow{P} & Y & \xrightarrow{q_y} & Z \end{array}$$

Part 4 : Strictification

(Grandis,
Paré
1999)



$$\begin{array}{ccc} DblCat & \xleftarrow{\perp} & PsDblCat \\ (Campbell 2019) & \xrightarrow{st_2(-)} & Inc \end{array}$$

Th^M (Grandis / Paré
 M_o , Gambino / M.) Let

A, B be strict double categories & $A \xrightarrow{P} B$
be a double profunctor,

1. $(A_{\text{loose}})_1$ free on a graph, P a double pseudofunctor

$\Rightarrow P \cong P'$ a strict double functor (Grandis
Paré)

2. $(B_{\text{loose}})_1$ "

" , P a double pseudocfunctor

$\Rightarrow P \cong P'$ a strict

double cofunctor

(M. 2023)

Th^m (M. 2024)

$(A_{\text{loose}})_1$, $(B_{\text{loose}})_1$ are both free on graphs,

$A \xrightarrow{P} B$ a double profunctor

$\Rightarrow P \cong P'$ where

$$\begin{array}{ccc} P_i \times A & \xrightarrow{r} & P_i \\ A_0 & & \end{array} \quad \left. \begin{array}{c} \\ \end{array} \right\} \text{strict algebras}$$
$$\begin{array}{ccc} B_i \times P_i & \xrightarrow{l} & P \\ B_0 & & \end{array}$$

... but still have a natural
iso. mediating the usual

$$\begin{array}{c}
 B_1 \times P_1 \times A_1 \\
 B_0 \quad A_0
 \end{array}
 \begin{array}{c}
 \swarrow \quad \searrow \\
 B_1 \times_{B_0} P_1 \qquad P_1 \times_{A_0} A_1 \\
 \searrow \qquad \swarrow \\
 P_1
 \end{array}$$

bimodularity
 condition
 \cong

$$\begin{array}{ccccc}
 X & \xrightarrow{x} & X' & \xrightarrow{P} & Y \\
 \parallel & & \cong & & \parallel \\
 X & \xrightarrow{x} & (X' & \xrightarrow{P} & Y') & \xrightarrow{y} & Y
 \end{array}$$

These semi-strict double profunctors correspond to double functors

$$(A^{\text{loose op}} \times B)^T \rightsquigarrow \text{Prof}$$

loose: normal lax

tight : separately strict,
jointly Pseudo.

A strict $\Rightarrow A \xrightarrow{1_A} A$
Strict

$$A \xrightarrow{P} B \xrightarrow{Q} C$$

A, B, C strict

P, Q semi-strict

$\Rightarrow Q \circ P$ semi-strict.