

# REPRESENTATION THEOREM FOR ENRICHED CATEGORIES

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# I. MOTIVATION

- Universal properties are useful to make many constructions in mathematics

Ex: (co)limits, adjunctions, Kan extensions, ...

- Encoded by the representability of a certain presheaf

Ex:  $G: I \rightarrow \mathcal{C}$  has a limit  $L \in \mathcal{C} \Leftrightarrow \mathcal{C}(I, L) \cong \mathcal{C}^I(\Delta(I), G)$ .

- Not always convenient to work with.

Ex: in  $\infty$ -category theory, homotopy type theory.

- Looking for an internal characterization.

Ex: Limit of a functor  $G \leftarrow$  terminal object in the category of cones  $\Delta \downarrow G$ .

## II. 1-CATEGORICAL STORY

$\mathcal{C}$ -category

A presheaf  $F: \mathcal{C}^{\text{op}} \rightarrow \text{Set}$  is **representable** if there is:

- an object  $L \in \mathcal{C}$ , and
- a natural isomorphism  $\mathcal{C}(-, L) \cong F$  ( $\in \text{Set}^{\mathcal{C}^{\text{op}}}$ )

Examples:

- **Limit:**  $G: I \rightarrow \mathcal{C}$  - functor  
 $F := \mathcal{C}^I(\Delta(-), G): \mathcal{C}^{\text{op}} \rightarrow \text{Set}, C \mapsto \mathcal{C}^I(\Delta C, G)$   
set of cones  
G has a limit  $\Leftrightarrow F$  is representable.
- **Adjunction:**  $L: \mathcal{C} \rightarrow \mathcal{D}$  - functor  
For  $D \in \mathcal{D}, F_D = \mathcal{D}(L(-), D): \mathcal{C}^{\text{op}} \rightarrow \text{Set}, C \mapsto \mathcal{D}(LC, D)$   
L has a right adjoint  $\Leftrightarrow F_D$  is representable  $\forall D \in \mathcal{D}$ .

Grothendieck construction:  $\mathcal{S}: \text{Set}^{\mathcal{C}^\text{op}} \rightarrow \text{Cat}_{/\mathcal{C}}$

$F: \mathcal{C}^\text{op} \rightarrow \text{Set} \mapsto \mathcal{S}_F \rightarrow \mathcal{C}$

where  $\mathcal{S}_F$  is the **category of elements** of  $F$

- obj: pairs  $(C, x)$  of  $C \in \mathcal{C}, x \in FC$
- mor  $(C, x) \rightarrow (D, y)$ :  $C \xrightarrow{f} D \in \mathcal{C}$  st  $Ff(y) = x$  ( $FD \xrightarrow{Ff} FC$ )

Remark: It corresponds to the category:

$$\bigsqcup_{C \in \mathcal{C}} FC \xrightleftharpoons[s]{t} \bigsqcup_{C, D \in \mathcal{C}} \mathcal{C}(C, D) \times FD \hookleftarrow \dots \in \text{Set}$$

Example: **Limit**:  $G: I \rightarrow \mathcal{C}; F = \mathcal{C}^I(\Delta(-), G)$

$\mathcal{S}_F \cong \Delta \downarrow G$  is the **category of cones** over  $G$

- obj: pairs  $(C, \kappa)$  of  $C \in \mathcal{C}, \kappa: \Delta C \Rightarrow G$
- mor  $(C, \kappa) \rightarrow (D, \mu)$ :  $C \xrightarrow{f} D \in \mathcal{C}$  st  $\Delta C \xrightarrow[\kappa \Rightarrow G \Leftarrow \mu]{\Delta f} \Delta D$

Theorem:  $\mathcal{L}_c : \text{Set}^{\mathcal{C}^\text{op}} \rightarrow \text{Cat}/c$  is fully faithful.

A functor  $P : A \rightarrow \mathcal{C}$  is a **discrete fibration** if for all  $A \in A$  and  $C \xrightarrow{g} PA \in \mathcal{C}$ , there is a unique  $B \xrightarrow{f} A \in A$  st  $Pf = g$

Remark: This is equivalent to saying that

$$\begin{array}{ccc} \text{mor } A & \xrightarrow{P} & \text{mor } \mathcal{C} \\ t \downarrow \lrcorner & & \downarrow t \\ \text{ob } A & \xrightarrow{P} & \text{ob } \mathcal{C} \end{array} \quad \text{is a pullback in Set}$$

Example: Grothendieck construction

$F : \mathcal{C}^\text{op} \rightarrow \text{Set}$ ,  $\mathcal{L}_c F \rightarrow \mathcal{C}$  is a discrete fibration

$$\begin{array}{ccc} \bigsqcup_{C \in \mathcal{C}} \mathcal{C}(C, D) \times FD & \longrightarrow & \text{mor } \mathcal{C} = \bigsqcup_{CD \in E} \mathcal{C}(C, D) \\ t \downarrow \lrcorner & & \downarrow t \\ \bigsqcup_{C \in \mathcal{C}} FC & \longrightarrow & \text{ob } \mathcal{C} = \bigsqcup_{C \in \mathcal{C}} * \end{array}$$

Notation:  $\text{Fib}/\mathcal{C} \subseteq \text{Cat}/\mathcal{C}$  full subcategory spanned by the discrete fibrations

Theorem:  $f_{\mathcal{C}}: \text{Set}^{\mathcal{C}^{\text{op}}} \rightarrow \text{Fib}/\mathcal{C}$  is an equivalence of categories.

Representation theorem: A presheaf  $F: \mathcal{C}^{\text{op}} \rightarrow \text{Set}$  is representable  $\Leftrightarrow f_{\mathcal{C}} F$  has a terminal object.

Example: Limit:  $G: I \rightarrow \mathcal{C}, F = \mathcal{C}^I(\Delta(-), G)$

$G$  has a limit  $\Leftrightarrow F$  is representable

$\Leftrightarrow f_{\mathcal{C}} F = \Delta \downarrow G$  has a terminal object

Remark: An object  $T$  in a category  $\mathcal{C}$  is terminal  $\Leftrightarrow$  the projection  $\mathcal{C}_{/T} \rightarrow \mathcal{C}$  is an isomorphism of categories.

$$\begin{array}{ccc} \mathcal{C}_{/T} & \longrightarrow & \mathcal{C}^2 \\ \downarrow & & \downarrow \\ \mathcal{C} & \xrightarrow{(\text{id}, \{T\})} & \mathcal{C} \times \mathcal{C} \end{array} \quad \epsilon \text{Cat}$$

### III PROBLEMS IN HIGHER DIMENSIONS

A **2-category**  $\mathcal{C}$  is a category enriched in  $\text{Cat}$ , i.e., it consists of:

- a set of objects  $\text{ob}\mathcal{C}$
- a **category** of morphisms  $\mathcal{C}(C,D)$ ,  $\forall C,D \in \mathcal{C}$
- + associative and unital compositions

$$C \xrightarrow{\quad} D + C \xrightarrow{\perp} D$$



Example: **Cat** -2-category of categories, functors, and natural transformations

$\mathcal{C}$ -2-category

A **2-presheaf** on  $\mathcal{C}$  is a 2-functor  $F: \mathcal{C}^\alpha \rightarrow \text{Cat}$ .

Example: **Representable** 2-presheaf: for  $L \in \mathcal{C}$ ,

$$\mathcal{C}(-, L): \mathcal{C}^\alpha \rightarrow \text{Cat}, C \mapsto \mathcal{C}(C, L)$$

A 2-presheaf  $F: \mathcal{C}^{\text{op}} \rightarrow \text{Cat}$  is **representable** if there is:

- an object  $L \in \mathcal{C}$
- a 2-natural isomorphism  $\mathcal{C}(-, L) \cong F(\in \text{2Cat}(\mathcal{C}^{\text{op}}, \text{Cat}))$

Example: **2-limit**:  $G: I \rightarrow \mathcal{C}$  - 2-functor

$F := \mathcal{C}^I(\Delta(-), G): \mathcal{C}^{\text{op}} \rightarrow \text{Cat}, C \mapsto \mathcal{C}^I(\Delta C, G) \leftarrow$   
category of 2-cones & modifications

$G$  has a 2-limit  $\Leftrightarrow F$  is representable

$L \in \mathcal{C}$  - 2-limit of  $G$ ; then  $\forall C \in \mathcal{C}$

$$\mathcal{C}(C, L) \xrightarrow{\sim} \mathcal{C}^I(\Delta C, G) \in \text{Cat}$$

$$C \rightarrow L \rightsquigarrow \Delta C \rightarrow G \quad (= \Delta C \rightarrow \Delta L \rightarrow G)$$

$$C \overbrace{\parallel}^{\text{II}} L \rightsquigarrow \Delta C \overbrace{\parallel}^{\text{III}} G \quad (= \Delta C \overbrace{\parallel}^{\text{III}} \Delta L \rightarrow G)$$

An object  $T$  in a 2-category  $\mathcal{C}$  is **2-terminal** if  $\mathcal{C}(C, T) \cong 1 \quad \forall C \in \mathcal{C}$

Try 1:  $\mathcal{S}_\mathcal{C} : \mathbf{2}\text{Cat}(\mathcal{C}^\text{op}, \mathbf{Cat}) \rightarrow \mathbf{2}\text{Cat}/\mathcal{C}$ ,  $F : \mathcal{C}^\text{op} \xrightarrow{\sim} \mathbf{Cat} \mapsto \mathcal{S}_\mathcal{C} F \rightarrow \mathcal{C}$

where  $\mathcal{S}_\mathcal{C} F$  is the (strict) 2-category of elements of  $F$

- obj: pairs  $(C, x)$  of  $C \in \mathcal{C}$ ,  $x \in FC$
- mor  $(C, x) \rightarrow (D, y)$ :  $C \xrightarrow{f} D \in \mathcal{C}$  st  $Ff(y) = x$
- 2-mor  $(C, x) \xrightarrow{\quad f \quad} (D, y)$ :  $C \xrightarrow{f} D \in \mathcal{C}$  st  $(F\alpha)_y = \text{id}_x$  ( $Ff \xrightarrow{F\alpha} Fg$ )

⚠ Does not see the morphisms of  $FC$

E.g. 2-limit of  $G \not\rightarrowtail$  2-terminal object in the (strict) 2-category of cones  $\Delta \downarrow G = \mathcal{S}_\mathcal{C} \mathcal{C}^I(\Delta(-), G)$ .

- obj: pairs  $(C, \kappa)$  of  $C \in \mathcal{C}$ ,  $\kappa : \Delta C \Rightarrow G$
  - mor  $(C, \kappa) \rightarrow (D, \mu)$ :  $C \xrightarrow{f} D \in \mathcal{C}$  st  $\Delta C \xrightarrow{\Delta f} \Delta D$
  - 2-mor  $(C, \kappa) \xrightarrow{\quad f \quad} (D, \mu)$ :  $C \xrightarrow{f} D \in \mathcal{C}$  st  $\begin{array}{ccc} \kappa & \Rightarrow & \mu \\ \downarrow & & \downarrow \\ G & \xrightarrow{\Delta f} & \Delta D \\ \kappa & \Rightarrow & \mu \\ \downarrow & & \downarrow \\ G & \xrightarrow{\Delta g} & \Delta D \end{array}$
- [Eckmann-M.] counterexamples.

Try 2:  $\int_e^{\text{lax}}: \text{2Cat}(\mathcal{C}^\text{op}, \text{Cat}) \rightarrow \text{2Cat}/\mathcal{C}$ ,  $F: \mathcal{C}^\text{op} \rightarrow \text{Cat} \mapsto \int_e^{\text{lax}} F \rightarrow \mathcal{C}$

where  $\int_e^{\text{lax}} F$  is the **lax** 2-category of elements of  $F$

- mor  $(C, x) \rightarrow (D, y)$ : pairs  $(f, \varphi)$  of  $C \xrightarrow{f} D \in \mathcal{C}$ ,  $F(f)(y) \xrightarrow{\varphi} x \in FC$

Theorem: [Buckley]  $\int_{\mathcal{C}}^{\text{lax}}: \text{2Cat}(\mathcal{C}^\text{op}, \text{Cat}) \rightarrow \text{2Cat}/\mathcal{C}$  is **fully faithful**.

But: 2-limit of  $G \not\leftrightarrow$  2-terminal object in the lax 2-category of cones  $\Delta \downarrow^{\text{lax}} G = \int_e^{\text{lax}} \mathcal{C}^I(\Delta(-), G)$

- mor  $(C, \kappa) \rightarrow (D, \mu)$ : pairs  $(f, \varphi)$  of  $C \xrightarrow{f} D$ ,  $\Delta C \xrightarrow{\Delta f} \Delta D$

[Eckingman - M.] **Counterexamples**

$$\begin{array}{ccc} & \Delta f & \\ \pi \searrow & \varphi & \swarrow \mu \\ G & & D \end{array}$$

Other approach: [Gagna-Harpaz-Lorari]

2-limit of  $G \rightsquigarrow$  "2-final" object in  $\Delta \downarrow^{\text{lax}} G$   
with respect to cartesian morphisms ( $\varphi$  invertible)

## IV SOLUTION: DOUBLE CATEGORIES

A **double category**  $\mathbf{A}$  is an internal category  
to categories

$$\mathbf{A}_0 \begin{array}{c} \xleftarrow{s} \\[-1ex] \xrightleftharpoons[t]{\quad} \\[-1ex] \xrightarrow{c} \end{array} \mathbf{A}_1 \subset \mathbf{A}_1 \times_{\mathbf{A}_0} \mathbf{A}_1 \in \mathbf{Cat}$$

category of objects  
& vertical morphisms

category of horizontal  
morphisms & squares

There is a (horizontal) embedding

$$\text{Int}: \mathbf{2Cat} \rightarrow \mathbf{DblCat}, \quad \mathcal{C} \mapsto \text{ob } \mathcal{C} \begin{array}{c} \xleftarrow{s} \\[-1ex] \xrightleftharpoons[t]{\quad} \\[-1ex] \xrightarrow{c} \end{array} \bigsqcup_{C,D \in \mathcal{C}} \mathcal{C}(C,D)$$

Sees  $\mathcal{C}$  as a double cat. with **trivial** vertical morphisms

The functor Int has a right adjoint

$$\text{Enr}: \mathbf{DblCat} \rightarrow \mathbf{2Cat}, \quad \mathbf{A} \mapsto \begin{cases} \text{ob } (\mathbf{A}_0) \\ \text{Enr } \mathbf{A}(A,B) = A_1 \times_{A_0 \times A_0} \{(A,B)\} \end{cases}$$

Forgets the **vertical** morphisms of  $\mathbf{A}$ .

[Paré]  $\mathcal{S}\mathcal{C}_c : 2\text{Cat}(\mathcal{C}^\alpha, \text{Cat}) \rightarrow \text{DbCat}_{/\text{IntC}}$

$F: \mathcal{C}^{\text{op}} \rightarrow \text{Cat} \mapsto \mathbb{S}_{\mathcal{C}} F \rightarrow \text{Int}\mathcal{C}$  where  $\mathbb{S}_{\mathcal{C}} F$  has :

where  $\mathcal{S}\mathcal{E}F$  is the double category of elements of  $F$

- obj: pairs  $(C, x)$  of  $C \in \mathcal{C}$ ,  $x \in FC$
  - hor mor  $(C, x) \rightarrow (D, y)$ :  $C \xrightarrow{f} D \in \mathcal{C}$  st  $Ff(y) = x$
  - ver mor  $(C, x) \dashrightarrow (C, x')$ :  $x \xrightarrow{\varphi} x' \in FC$

- sq  $(C, x) \xrightarrow{f} (D, y) : C \xrightarrow{\begin{smallmatrix} f \\ \perp \alpha \end{smallmatrix}} D \in \mathcal{C}$  st  $x = Ff(y) \xrightarrow{Ff(\psi)} FP(y)$   
 $\varphi \downarrow \quad \downarrow \psi$   
 $(C, x') \xrightarrow{g} (D, y')$   
 $(F\alpha)_y \downarrow \quad \searrow \varphi \quad \downarrow (F\alpha)_{y'}$   
 $Fg(y) \xrightarrow{Fg(\psi)} Fg(y') = x'$

Remark: It corresponds to the double category:

$$\coprod_{C \in \mathcal{C}} FC \xrightleftharpoons[\tau]{\sigma} \coprod_{C, D \in \mathcal{C}} \mathcal{C}(C, D) \times FD \xleftarrow{\epsilon} \dots \in \text{Cat}$$

Observation:  $S_e = \text{Enr } S \cap S_e$

[Paré] An object  $T$  in a double category  $\mathcal{A}$  is **double terminal** if

$$-\forall A \in \mathcal{A}, \exists ! A \rightarrow T \in \mathcal{A} \quad - \forall \begin{matrix} A \\ \downarrow \\ B \end{matrix} \in \mathcal{A}, \exists ! \begin{matrix} A \xrightarrow{\quad} T \\ \downarrow \\ B \xrightarrow{\quad} T \end{matrix} \in \mathcal{A}$$

Remark: This is equivalent to requiring that the projection  $\mathcal{A}_{/T} \rightarrow \mathcal{A}$  is an isomorphism of double categories.

$$\begin{array}{ccc} \mathcal{A}_{/T} & \longrightarrow & \mathcal{A}^{\text{Int} 2} \\ \downarrow & & \downarrow \\ \mathcal{A} & \xrightarrow{(\text{id}, \{T\})} & \mathcal{A} \times \mathcal{A} \end{array} \quad \in \text{DblCat}$$

bilimits, pseudo-functors

Theorem: [dingman-M.] A 2-presheaf  $F: \mathcal{C}^{\text{op}} \rightarrow \text{Cat}$  is representable  $\Leftrightarrow \int \mathcal{C} F$  has a double terminal object.

Example: **2-limits:**  $G: I \rightarrow \mathcal{C}$ ,  $F = \mathcal{C}^I(\Delta(-), G)$

$G$  has a 2-limit  $\Leftrightarrow$  the double category of cones

$\Delta \Downarrow G = \int \mathcal{C} F$  has a double terminal object

- ver mor  $(C, \kappa) \mapsto (C, \kappa'): \Delta C \xrightarrow{\text{J}} \int \mathcal{C} G$

Theorem: [Clingman-M.] Let  $\mathcal{A}$  be a double category with **tobulators**. Then an object  $T \in \mathcal{A}$  is double terminal  $\Leftrightarrow T \in \text{Enr} \mathcal{A}$  is 2-terminal.

Lemma: Let  $F: \mathcal{C}^{\text{op}} \rightarrow \text{Cat}$  be a 2-presheaf. Suppose that  $\mathcal{C}$  has **tensors by 2** and  $F$  preserves them, then  $\int_{\mathcal{C}} F$  has tobulators.

Cor: Under the hypotheses of the lemma,  $F$  is representable  $\Leftrightarrow \int_{\mathcal{C}} F = \text{Enr} \int_{\mathcal{C}} F$  has a 2-terminal object.

Observation:  $G: I \rightarrow \mathcal{C}$  - 2-functor.

If  $\mathcal{C}$  has tensors by 2, then  $F = \mathcal{C}^I(\Delta(-), G)$  preserves them.

Cor: Let  $G: I \rightarrow \mathcal{C}$  be a 2-functor and suppose that  $\mathcal{C}$  has tensors by 2. Then  $G$  has a 2-limit  $\Leftrightarrow$  the 2-category of cones  $\Delta \downarrow G$  has a 2-terminal object.

## V. CASE OF ENRICHED CATEGORIES

$(\mathcal{V}, \times, 1)$ -cartesian closed category with pullbacks.

Need  $\mathcal{V}$  to be **extensive**, i.e., it has all small coproducts that "behave well"

Consequence:  $\text{Set} \xrightarrow{\perp} \mathcal{V}$ ,  $S \mapsto \bigsqcup 1$   
 $\uparrow$   
 $U = \mathcal{V}(1, -)$

A  $\mathcal{V}$ -enriched category  $\mathcal{C}$  consists of:

- a set of objects
- a hom-object  $\mathcal{C}(C, D) \in \mathcal{V} \quad \forall C, D \in \mathcal{C}$
- + associative and unital compositions

Example:  $\mathcal{D}$  is  $\mathcal{V}$ -enriched with:

- object set  $\text{ob } \mathcal{D}$
- homs  $Y^X$  given by the internal homs  $\forall X, Y \in \mathcal{V}$

Remark: Set  $\hookrightarrow \mathcal{V}$  induces an inclusion  $\text{Cat} \hookrightarrow \mathcal{V}\text{-Cat}$   
 $\mathcal{C}$ - $\mathcal{V}$ -enriched category

A  $\mathcal{V}$ -presheaf on  $\mathcal{C}$  is a  $\mathcal{V}$ -functor  $F: \mathcal{C}^{\text{op}} \rightarrow \mathcal{V}$

Example: Representable  $\mathcal{V}$ -presheaf: for  $L \in \mathcal{C}$ ,

$$\mathcal{C}(-, L): \mathcal{C}^{\text{op}} \rightarrow \mathcal{V}, C \mapsto \mathcal{C}(C, L).$$

A  $\mathcal{V}$ -presheaf  $F: \mathcal{C}^{\text{op}} \rightarrow \mathcal{V}$  is **representable** if there is:

- an object  $L \in \mathcal{C}$
- a  $\mathcal{V}$ -natural isomorphism  $\mathcal{C}(-, L) \cong F$  ( $\in \mathcal{V}\text{-Cat}(\mathcal{C}^{\text{op}}, \mathcal{V})$ )

Example:  $\mathcal{V}$ -enriched limits:  $G: I \rightarrow \mathcal{C}$ - $\mathcal{V}$ -functor

$$F := \mathcal{C}^I(\Delta(-), G): \mathcal{C}^{\text{op}} \rightarrow \mathcal{V}, C \mapsto \mathcal{C}^I(\Delta C, G) \leftarrow \text{object of } \mathcal{V}\text{-cones in } \mathcal{V}$$

$G$  has a  $\mathcal{V}$ -enriched limit  $\Leftrightarrow F$  is representable.

An internal category to  $\mathcal{V}$   $A$  is a diagram

$$A_0 \xleftarrow[\tau]{i} A_1 \subset A_1 \times_{A_0} A_1 \in \mathcal{V} \quad \text{+ relations}$$

There is an embedding

$$\text{Int}: \mathcal{V}\text{-Cat} \rightarrow \text{Cat}(\mathcal{V}), \mathcal{C} \mapsto \text{ob } \mathcal{C} \xleftarrow[\tau]{i} \bigsqcup_{C,D \in \mathcal{C}} \mathcal{C}(C,D)$$

[Cottrell-Fujii-Power] Int has a right adjoint

$$\text{Enr}: \text{Cat}(\mathcal{V}) \rightarrow \mathcal{V}\text{-Cat}, A \mapsto \begin{cases} UA_0 \\ \text{Enr } A(A,B) = A_1 \times_{A_0 \times A_0} \{(A,B)\} \end{cases}$$

Grothendieck construction:

$$f_e: \mathcal{V}\text{-Cat}(\mathcal{C}^\text{op}, \mathcal{V}) \rightarrow \text{Cat}(\mathcal{V}) /_{\text{Int } \mathcal{C}}, F: \mathcal{C}^\text{op} \rightarrow \mathcal{V} \mapsto f_e F \rightarrow \text{Int } \mathcal{C}$$

where  $f_e F$  is the internal category to  $\mathcal{V}$

$$\bigsqcup_{C \in \mathcal{C}} FC \xleftarrow[\tau]{i} \bigsqcup_{C,D \in \mathcal{C}} \mathcal{C}(C,D) \times FD \xleftarrow{c} \dots \in \mathcal{V}$$

Remark: [Beardsley-Wong] have a "lax" version

Theorem: [M-Sarazola-Verdugo]

$\mathcal{S}_e : \mathcal{V}\text{-Cat}(\mathcal{C}^\alpha, \mathcal{V}) \rightarrow \text{Cat}(\mathcal{V}) /_{\text{Int}\mathcal{C}}$  is **fully faithful**

An internal functor  $P : A \rightarrow \text{Int}\mathcal{C}$  in  $\text{Cat}(\mathcal{V})$  is a **discrete fibration** if

$$\begin{array}{ccc} A_1 & \xrightarrow{P} & (\text{Int}\mathcal{C})_1 = \bigsqcup_{C, D \in \mathcal{C}} C(D) \\ t \downarrow \perp & & \downarrow t \\ A_0 & \xrightarrow{P} & (\text{Int}\mathcal{C})_0 = \text{ob}\mathcal{C} \end{array}$$

is a pullback in  $\mathcal{V}$

Example: Grothendieck construction

$F : \mathcal{C}^\alpha \rightarrow \mathcal{V}$ ,  $\mathcal{S}_e F \rightarrow \text{Int}\mathcal{C}$  is a discrete fibration

Notation:  $\text{Fib}/e \subseteq \text{Cat}(\mathcal{V}) /_{\text{Int}\mathcal{C}}$  full subcategory spanned by the discrete fibrations over  $\text{Int}\mathcal{C}$

Theorem: [M-Sarazola-Verdugo]

$\mathcal{S}_e : \mathcal{V}\text{-Cat}(\mathcal{C}^\alpha, \mathcal{V}) \rightarrow \text{Fib}/e$  is an **equivalence** of categories.

An object  $T$  in an internal category  $\mathcal{A}$  to  $\mathcal{V}$  is **terminal** if  $\mathcal{A}_{/T} \rightarrow \mathcal{A}$  is an isomorphism in  $\text{Cat}(\mathcal{V})$ .

$$\begin{array}{ccc} \mathcal{A}_{/T} & \longrightarrow & \mathcal{A}^{\text{Int}2} \\ \downarrow \mathcal{J} & & \downarrow \in \text{Cat}(\mathcal{V}) \\ \mathcal{A} & \xrightarrow{(\text{id}, \{T\})} & \mathcal{A} \times \mathcal{A} \end{array}$$

Theorem: [M.-Sarazola-Verdugo]

A  $\mathcal{V}$ -presheaf  $F: \mathcal{C}^\text{op} \rightarrow \mathcal{V}$  is representable  
 $\iff \int_e F$  has a terminal object.

Example:  $\mathcal{V}$ -limits:  $G: I \rightarrow \mathcal{C}, F = \mathcal{C}^I(\Delta(-), G)$

$G$  has a  $\mathcal{V}$ -limit  $\iff$  the internal category of  $\mathcal{V}$ -cones  $\Delta \downarrow G := \int_e F$  has a terminal object.

Conj: If  $\mathcal{V}$  is "generated" by  $\{X_i\} \subseteq \mathcal{V}$ ,  $\mathcal{C}$  has tensors by the  $X_i$ 's and  $F: \mathcal{C}^\text{op} \rightarrow \mathcal{V}$  preserves them, then  $F$  is representable  $\iff \text{Enr} \int_e \mathcal{V}$  has a  $\mathcal{V}$ -terminal object  $T$  such that homs to  $T$  are isomorphic to 1

## VI. APPLICATIONS

- $\mathcal{V} = \text{Set}$ :  $\text{Set}\text{-Cat} = \text{Cat} = \text{Cat}(\text{Set})$ 
  - ~ retrieve 1-categorical case
- $\mathcal{V} = \text{Cat}$ :  $\text{Cat}\text{-Cat} = 2\text{Cat}$ ,  $\text{Cat}(\text{Cat}) = \text{DblCat}$ 
  - ~ another proof of the representation theorem
- $\mathcal{V} = (n-1)\text{Cat}$ :  $(n-1)\text{Cat}\text{-Cat} = n\text{-Cat}$ ,  $\text{Cat}((n-1)\text{Cat})$ 
  - ~ the Grothendieck construction takes values in **internal categories to  $(n-1)$ -categories**  
 $(\neq n\text{-fold categories})$
  - ~  **$n$ -limits**  $\leftrightarrow$  terminal objects in the internal category to  $(n-1)\text{Cat}$  of  $n$ -cones
  - ~ With Rosekh & Rovelli, we construct limits in  **$(\infty, n)$ -categories** using this characterization.
  - ~  $n$ -limits in an  $n$ -category with **tensors by  $\Sigma^{n-1} 1$**   
 $\leftrightarrow$   $n$ -terminal objects in the  $n$ -category of  $n$ -cones.