

703650 VO Parallel Systems WS2019/2020 13 Dwarfs of HPC

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Overview

- ▶ 13 Dwarfs of HPC
 - abstract application categories
- "Tales from the Proseminar"
 - Daniel's Weird int Problem
 - MPI_Ssend() and MPI_Sendrecv()

Motivation

- MPI API or concepts such as load imbalance are still pretty low-level characteristics of parallel programs
- need to be able to recognize higher-level classes of applications and discuss them

- ▶ this lecture presents the most prominent classes of HPC applications
 - many new applications you encounter will fit into these categories

How and Why are Dwarfs Defined?

- group applications by similarity in computation and data structures
 - first published by Asanovic et al in *The Landscape of Parallel Computing Research: A View from Berkeley*
 - https://www2.eecs.berkeley.edu/Pubs/TechRpts/2006/EECS-2006-183.pdf
- purely algorithmic, implementation-independent
 - enables cross-platform reasoning and cross-application resource sharing (e.g. libraries)
- serve as small, abstract, high-level benchmarks for studying new
 - programming models
 - communication patterns
 - hardware architectures, topologies
 - ..
- used to kick off innovation in all of these aspects



7 Original Dwarfs of HPC

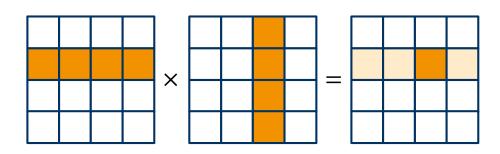
- What are we going do discuss?
 - 1. Dense Linear Algebra
 - 2. Sparse Linear Algebra
 - 3. Spectral Methods
 - 4. N-body Methods
 - ▶ 5. Structured Grids
 - ▶ 6. Unstructured Grids
 - > 7. Monte Carlo Methods

- What have you already heard?
 - matrix mul (first MPI lecture)

- particle-in-cell (this proseminar)
- heat stencil
- Monte Carlo π

Dense Linear Algebra

- data is stored in densely-populated matrices (or vectors)
 - data is stored uncompressed ("as is")
 - data access via strides, often unit stride
- e.g. matrix multiplication, LU decomposition, Gauss-Seidel, ...
- rarely done manually, there are a TON of libraries out there



Dense Linear Algebra: Characteristics

- naïve implementations often memory bound
 - remember: memory wall!
 - caches and prefetching helps
- simple but significant data structure
 - stride often enables/prevents vectorization (SIMD)
 - column- vs. row-major access (often transposed) affects cache efficiency
- still the default measure for performance in HPC
 - e.g. TOP500 uses HPL, a high performance LINPACK benchmark

```
for (int i = 0; i < N; ++i) {
  for (int j = 0; j < N; ++j) {
    double tmp = 0.0;
  for (int k = 0; k < N; ++k) {
    tmp += A[i][k] * B[k][j];
  }
  result[i][j] = tmp;
}</pre>
```

```
vgatherqpd ymm0{k2}, [rax+ymm5*1]
vmulpd ymm0, ymm0, YMMWORD PTR [rdx+rdi]
...
```

Dense Linear Algebra: Optimizations

vectorization

might entail modifications, e.g.
 transposing matrix A or B in matrix mul

loop blocking or tiling

- do not work on single elements but small blocks (e.g. 2x2 or 4x4)
- exploits locality and cache
- also, lots of other HOTs (higher order transformations)

hardware-specific instructions

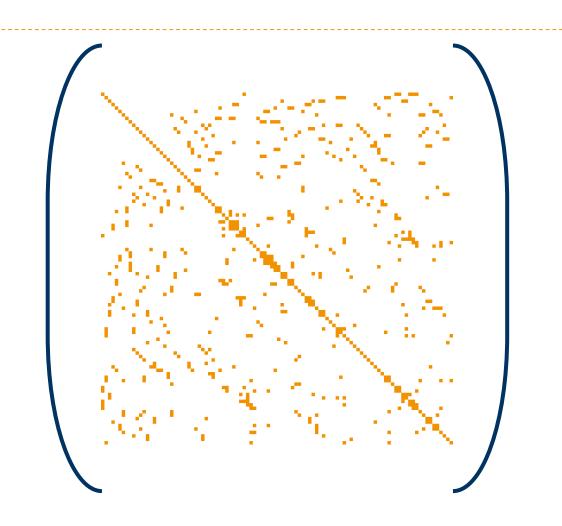
e.g. fused multiply-add (FMA)

```
for (int ii = 0; ii < N; ii += ib) {
 for (int jj = 0; jj < N; jj += jb) {
  for (int k = 0; k < N; ++k) {
  for (int i = ii; i < ii+ib; ++i) {
    for (int j = jj; j < jj+jb; ++j) {
      // ... process single tile
```

Sparse Linear Algebra

- data is stored in sparsely-populated matrices (or vectors)
 - vast majority of data is zero
 - data is stored in compressed format
 - data often accessed indirectly via indices

e.g. conjugate gradient, Google's PageRank, data mining



Sparse Linear Algebra: Characteristics

- computationally or memory limited
 - depends on sparsity of data, data structure representation and algorithm
- different data structures available
 - e.g. coordinate scheme (COO): (i, j, a_{ij})
 - array of structs vs. struct of arrays
 - not necessarily sorted!

```
typedef struct sparseElement {
    int i; int j; double value;
} sparseElement;
sparseElement sparseMatrix[SIZE];
sparseMatrix[0].i = 0;
typedef struct sparseMatrixT {
    int i[SIZE]; int j[SIZE];
    double values[SIZE];
} sparseMatrixT;
sparseMatrixT sparseMatrix;
sparseMatrix.i[0] = 0;
```

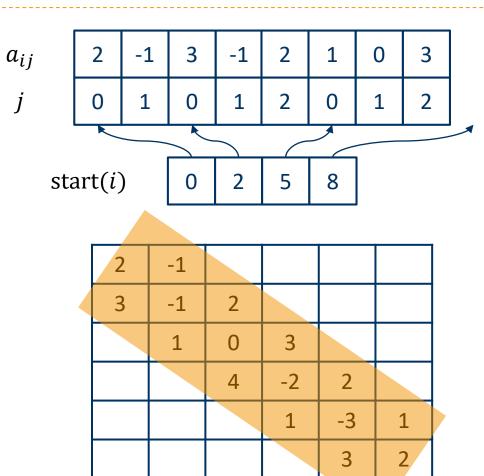
Sparse Linear Algebra: Optimizations

compressed row storage (CRS)

- two arrays of size N
- one holds a_{ij} , the other j
- third array points to start of row i in j
- smaller memory footprint than COO
 - $\triangleright 2N + (m+1) \text{ vs. } 3N$

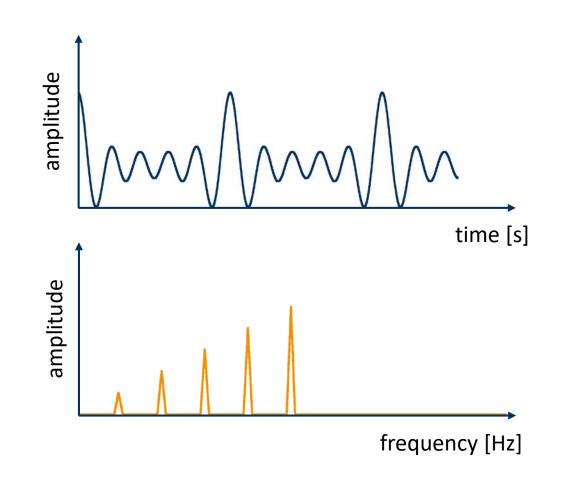
variants

- column-major variant (CCS)
- store small blocks (2x2 or 4x4) instead of single elements, improves SIMDness
- compressed diagonal storage (CDS)
 - ▶ remember: use domain-specific knowledge!



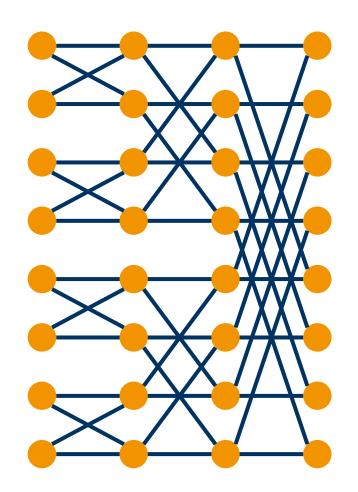
Spectral Methods

- data in frequency domain, not time or space
 - can include multiple stages of computations alternating between local and global communication
- e.g. fast Fourier transform (FFT), audio and video signal processing



Spectral Methods: Characteristics

- usually implemented using butterfly patterns
 - multiple stages of multiply-add
 - often latency limited due to global communication patterns (e.g. all-to-all)
- often resemble structured or unstructured grid methods after transformation to frequency domain



Spectral Methods: Optimizations

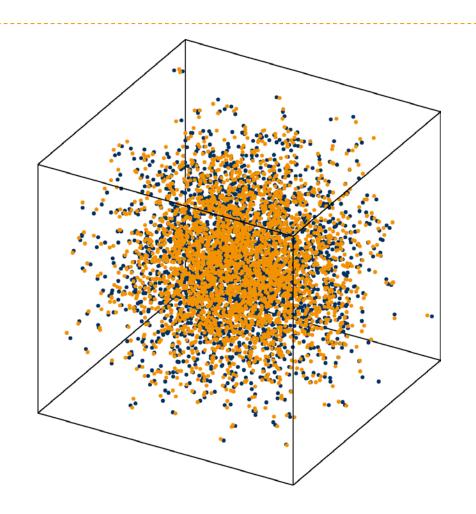
- requires optimization of the transformation to frequency domain
 - relies on transposing data efficiently
- afterwards, consider similar optimizations as for (un)structured grids

- severely restricted scalability to larger HPC systems
 - ongoing research
 - lot's of it in math

N-body Methods

- models interactions between discrete, moving points
 - often requires dynamic data structures
 - varying spatial locality
 - movement affects load balance and data access costs

• e.g. galaxy collision simulations, molecular dynamics, protein folding



N-body Methods: Characteristics

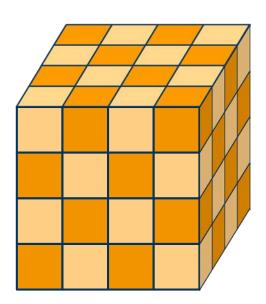
don't want to spoil this week's assignment...

N-body Methods: Optimizations

don't want to spoil this week's assignment...

Structured Grids

- models interactions between discrete, fixed points
 - grid structure described by pattern
 - topological information easily derived
 - usually high spatial locality
 - may be subdivided into finer grid ("adaptive")
- e.g. heat transfer (stencil), computational fluid dynamics (CFD), octrees



Structured Grids: Characteristics

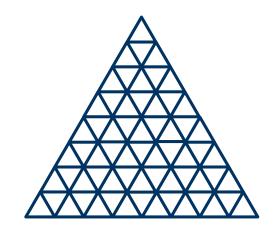
structured nature is a key aspect

- similar characteristics compared to dense linear algebra (e.g. row-major vs. columnmajor)
- memory addresses predictable, facilitates e.g. prefetching
- adaptive grids possible

typically memory bound

- e.g. 7-point stencil in 3D: load 7 data points for computing a single new one
- local communication only (ghost cell exchange)

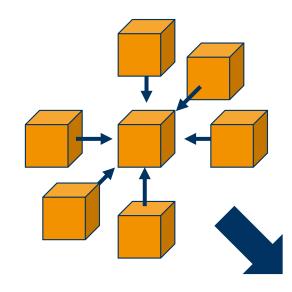
structured grid does not imply rectangular!





Structured Grids: Optimizations

- decomposition, decomposition, decomposition
 - remember in 3D: slabs, poles, cubes
 - highly dependent on type of grid and specific use case
- structured nature of the problem makes analytic prediction possible
 - performance models and simulators for simple stencil kernels (e.g. Kerncraft)

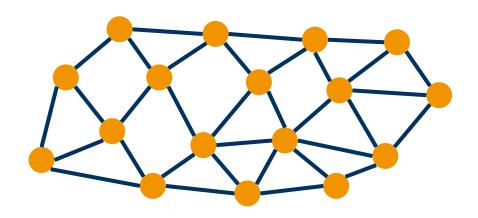


L1 misses: # L2 misses: # L3 misses: # ...

Unstructured Grids

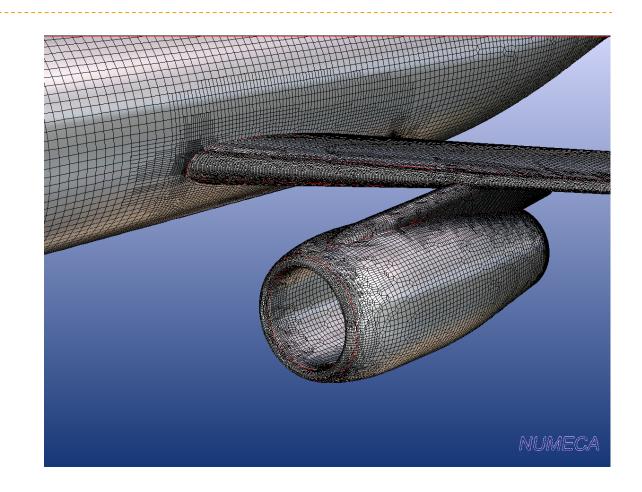
- models interactions between discrete, fixed points
 - grid pattern described explicitly by individual connections
 - irregular geometry and topology
 - usually involves multiple levels of indirection when accessing data

• e.g. computational fluid dynamics (CFD)



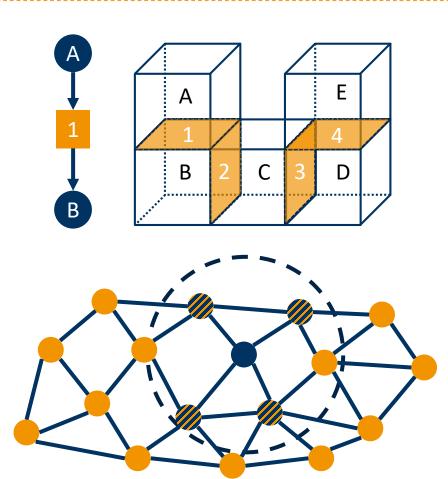
Unstructured Grids: Characteristics

- usually heavily latency bound due to indirect access
 - ... = cells[neighbors[i]]
 - ... = cell.getNeighbor(i)
 - also known as "pointer chasing"
- problems similar to structured grids, e.g.
 - domain decomposition / adaptivity
 - topological information
 - ghost cell exchange
 - but hardly analytically predictable



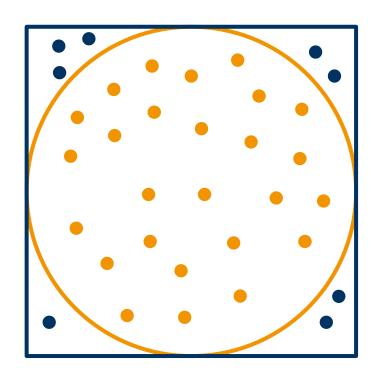
Unstructured Grids: Optimizations

- problem space discretization and topology
 - which types of grid elements?
 - which types of connections?
 - cells, faces, vertices, edges, ...
 - should be efficient to load/store as well
- decomposition, decomposition, decomposition
 - efficient ghost cell exchange required
 - efficient grid navigation e.g. to access neighbors



Monte Carlo Methods

- also known as map-reduce
 - process data independently and merge the results
- models statistical evaluation of repeated random trials
 - communication usually insignificant
 - embarrassingly parallel, multiple copies of sequential method
- e.g. numerical integration, quantum many-body problems, ray tracing



Monte Carlo Methods: Characteristics

- parallelization almost a no-factor
 - similar to multiple sequential programs sharing some resources (e.g. L3 cache, random number generator)
 - relatively inexpensive reduction
- depends heavily on sequential optimizations

Monte Carlo Methods: Optimizations

- not much to do beyond sequential optimization
 - **ILP**
 - prefetching
 - vectorization
 - resource contention (read: fast random number generation)

- consider different hardware
 - **GPUs**
 - FPGAs
 - ...
- try to decrease cost of evaluating a sample
 - increase number of samples if required

Additional Dwarfs

Additional Dwarfs

- ▶ 8. Combinational Logic
- ▶ 9. Graph Traversal
- ▶ 10. Dynamic Programming
- ▶ 11. Backtrack & Branch+Bound
- ▶ 12. Graphical Models
- ▶ 13. Finite State Machine

 focus more (but not exclusively) on integer-heavy applications and machine learning

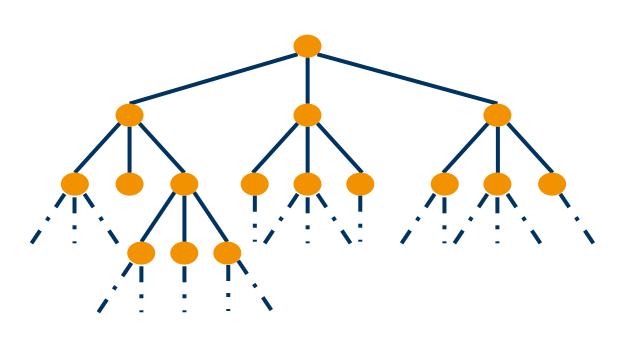
Combinational Logic

- generally involves performing simple operations on large amounts of integer data
 - e.g. computing cyclic redundancy codes (CRC)
- often parallelizable on multiple levels
 - bit-level parallelism
 - block-level parallelism

```
uint8_t compute(uint8_t const msg[], int n) {
    uint8 t rem = 0;
    for (int byte = 0; byte < n; ++byte) {
        rem ^= (msg[byte] << (WIDTH - 8));</pre>
        for (uint8_t bit = 8; bit > 0; --bit) {
            if (rem & TOPBIT) {
                 rem = (rem << 1) ^ POLYNOMIAL;</pre>
             } else {
                 rem = (rem << 1);
    return (rem);
```

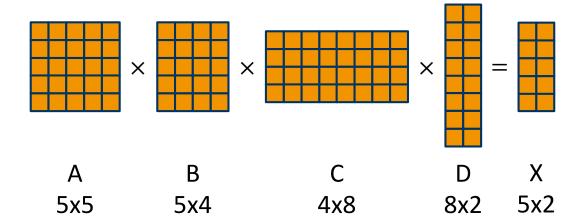
Graph Traversal

- traverse a number of objects in a graph and examine characteristics
 - e.g. searching, sorting, collision detection, decision trees, ...
 - usually heavy on data reads and lookups,
 very little computation and output
- parallelizable over different paths in the graph
 - but indirect accesses are heavily latency-bound (c.f. unstructured grids)



Dynamic Programming

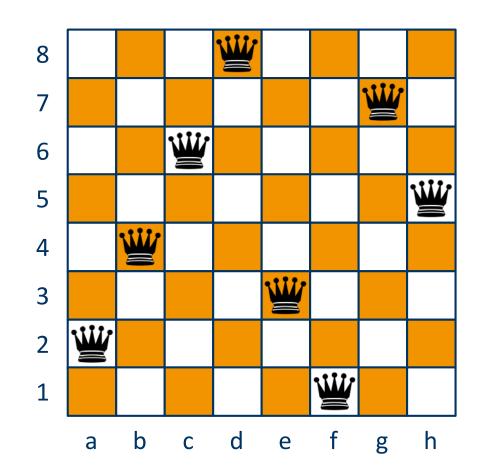
- method of computing solutions by solving simpler, overlapping sub-problems
 - applicable to problems where optimal result is composed of optimal results of sub-problems
 - e.g. matrix-chain-multiplication
- usually based on memoization
 - solve each sub-problem exactly once
 - store and re-use the result



$$A \times (B \times (C \times D)) = X$$
 154 ops
 $(A \times B) \times (C \times D) = X$ 204 ops
 $((A \times B) \times C) \times D = X$ 340 ops

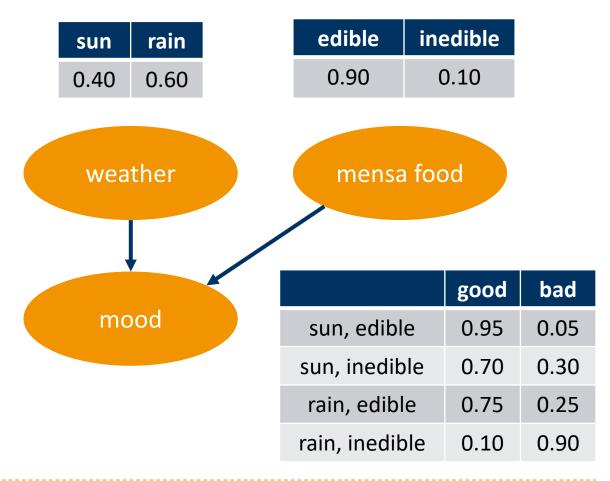
Backtrack & Branch+Bound

- search and optimization problems for very large problem spaces
 - incrementally build solution but discard if determined unsuitable
 - e.g. n-queens problem
- use divide & conquer strategy:
 - break down complex problem into smaller sub-problems until they become solvable
 - solve sub-problems in parallel



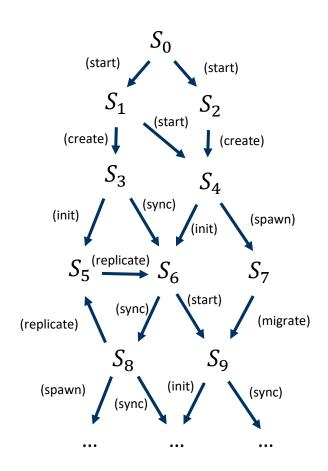
Graphical Models

- also known as probabilistic models
- represent graphs consisting of random variables as nodes and dependencies as edges
 - e.g. Bayesian networks, Hidden Markov models
- ongoing research in math and computer science regarding parallelization and optimization



Finite State Machines

- represent interconnected set of states to be moved among
 - e.g. parsers
- can sometimes be decomposed into multiple state machines that act in parallel
 - ongoing research



Literature Material

- ▶ White Paper by Berkeley University: "The Landscape of Parallel Computing Research: A View from Berkeley" from 2006:
 - https://www2.eecs.berkeley.edu/Pubs/TechRpts/2006/EECS-2006-183.pdf
- ▶ Holds more detailed descriptions and related aspects

Note the Focus on Scientific Problems

- MPI can also be used to implement
 - distributed-memory runtime systems
 - emulate shared memory runtime systems on distributed memory (e.g. PGAS)
 - provide the connecting parallelism layer for shared-memory or sequential systems
 - e.g. use multiple accelerators in separate compute nodes
 - extend shared memory parallelism to distributed memory (MPI+X)
 - **...**
- > still, the majority of codes is of scientific computing nature

Tales from the Proseminar

- sequential 2D heat stencil
 - ▶ 100x100 problem size
 - -O0 (also with -O2) on LCC2
 - gcc 4.8.5 (but also 9.2 and MSVC 2019)
 - switching from long long loop variables to int halves execution time

```
$ /usr/bin/time –f %E ./stencil_long 100 100
0:04.60
$ /usr/bin/time –f %E ./stencil_int 100 100
0:02.31
```

What the heck is going on?

profiled with gprof to find "hot spots", left is int, right is long long

% c	umulative	selt	f
time	seconds	second	ds name
32.92	0.76	0.76	int.c:93
20.06	1.22	0.46	int.c:78
14.8 3	1.56	0.34	int.c:79
10.47	1.80	0.24	<pre>int.c:87</pre>
9.59	2.02	0.22	int.c:86

% C	umulative	self	
time	seconds	second	ls name
36.84	1.69	1.69	long.c 79
31.61	3.15	1.45	long.c 78
16.35	3.90	0.75	long.c 93
5.34	4.15	0.25	long.c:86
3.27	4.30	0.15	long.c:87

```
// get temperature at current position
75
     value_t tc = A[i];
76
     // get temperatures of adjacent cells
     value_t txl = (i % Nx != 0) ? A[i - 1] : tc;
78
     value_t txr = (i % Nx != Nx - 1) ? A[i + 1] : tc;
     // ..... snip .....
     if (Ny > 1)
90
       B[i] = tc + 0.165 * (txl + txr + tyl + tyr + tzl + tzr + (-6 * tc));
91
92
     else
       B[i] = tc + 0.2 * (txl + txr + tzl + tzr + (-4 * tc));
93
     // if ((int)B[i] < (int)A[i])</pre>
```

far-fetched idea, but maybe branch (miss-)predictions? compared with perf stat:

```
2328.650634 task-clock:u (msec) #
                                      0.996 CPUs
            0 context-switches:u #
                                      0.000 K/sec
            0 cpu-migrations:u
                                     0.000 K/sec
          186 page-faults:u
                                      0.080 K/sec
5,739,935,525 cycles:u
                                      2.465 GHz
10,671,532,880 instructions:u
                                      1.86 IPC
1,196,623,336 branches:u
                                  # 513.870 M/sec
    1,030,678 branch-misses:u
                                      0.09%
  2.338387625 seconds time elapsed
```

```
4639.728721 task-clock:u (msec) #
                                     0.998 CPUs
            0 context-switches:u #
                                     0.000 K/sec
            0 cpu-migrations:u
                                 # 0.000 K/sec
          186 page-faults:u
                                     0.040 K/sec
11,444,184,736 cycles:u
                                     2.467 GHz
10,972,976,005 instructions:u
                                     0.96 IPC
                                 # 257.984 M/sec
1,196,977,569 branches:u
    1,030,347 branch-misses:u
                                     0.09%
  4.650327844 seconds time elapsed
```

```
value_t txl = ( i % Nx != 0 ) ? A[i-1] : tc;
      eax, DWORD PTR [rbp-20]
                                              rax, 3
                                       sal
mov
                                             rdx, [rax 8]
cdq
                                       lea
      DWORD PTR [rbp-24]
idiv
                                             rax, QWORD PTR [rbp-32]
                                       mov
      eax, edx
                                       add
                                             rax, rdx
mov
                                       movsd xmm0, QWORD PTR [rax]
test
      eax, eax
je
                                       jmp .L3
      .L2
      eax, DWORD PTR [rbp-20]
mov
cdqe
```

- ► Found instruction information on Agner Fog's blog for Intel Skylake architecture:
 - https://www.agner.org/optimize/instruction tables.pdf

inst.	operands	µорs	latency
idiv	r32	10	26
idiv	r64	57	42-95

if you want to study compiler output: https://godbolt.org/

- root cause of the issues: single loop for multi-dimensional problem space
 - requires % operator to get boundaries (which are strided in linearized space)
 - replace with loop nest and comparison operators (>, <, ==, !=)</p>
- potentially premature optimization and violates step 1 of "Four Steps to Creating an Optimized Parallel Program"
 - ▶ I asked Daniel's team, they claim that wasn't the case
 - argument holds nevertheless

Tales from the Proseminar: MPI_Ssend() & MPI_Sendrecv()

- synchronous send (MPI_Ssend())
 - is NOT the same as non-blocking send
 - is NOT necessarily better than MPI_Send()
 - can be used for debugging and finding deadlocks
 - is not always better than asynchronous send
- MPI_Sendrecv()
 - is blocking on the outside (will return when buffers available)
 - is non-blocking on the inside
 - implemented using MPI_Isend(), MPI_Irecv(), and two MPI_Wait()

Summary

▶ 13 Dwarfs of HPC

- abstract application categories
- facilitate cross-platform reasoning and cross-application component reuse
- 7 older ones, most well-studied physics problems
- ▶ 6 newer ones, partially of theoretical nature, subject to ongoing research
- give a broad perspective on HPC potential and limitations

"Tales from the Proseminar"

- Daniel's Weird int Problem: Don't consider int instructions cheap!
- MPI Ssend() and MPI Sendrecv()

Image Sources

- Dwarfs: http://rpgmke.wikidot.com/moonshae-isles-campaign
- ► CRC: https://barrgroup.com/Embedded-Systems/How-To/CRC-Calculation-C-Code
- ▶ Unstructured Mesh: https://resourcearea.cpu-24-7.com/en/numeca_welcome