

Q5

1.

The differentiate of  $f(n) = \log_2(n)$  is

$$f(n)' = \frac{1}{\ln(2) * n}$$

The differentiate of  $g(n) = 10\sqrt[n]{n}$  is

$$g(n)' = \frac{1}{-10 * n^{11/10}}$$

If we calculate

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{f(n)'}{g(n)'} &= \frac{-10 * n^{11/10}}{\ln(2) * n} \\ &= \frac{-10 * n^{1/10}}{\ln(2)} \end{aligned}$$

We also knew

$$\lim_{n \rightarrow \infty} n n^{1/10} = \infty, \ln(2) \text{ is a constant}$$

Therefore  $\lim_{n \rightarrow \infty} \frac{f(n)'}{g(n)'} = \infty$ , This result indicates

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \text{ (according to L'Hopital's rule)}$$

we can also calculate

$$\lim_{n \rightarrow \infty} \frac{g(n)'}{f(n)'} = \frac{\ln(2) * n}{-10 * n^{11/10}}$$

$$= \frac{\ln(2)}{-10 * n^{1/10}}$$

Recall

$$\lim_{n \rightarrow \infty} n^{1/10} = \infty, \ln(2) \text{ is a constant}$$

Thus

$$\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} = 0$$

From the above two limits and According to the limit asymptotic theorem, we can come to the conclusion that  $g(n) = O(f(n))$ .

2.

For  $g(n) = 2^{n \cdot \log_2 n^2}$  can be simplified by using the log identity ( $2^{\log_2 n^2} = n^2$ ). Based on the log identity, we can reduce  $g(n)$  to  $2^{n \cdot n^2}$ .

For  $f(n) = n^n$ , this can be rewritten to  $f(n) =$

$e^{n \cdot \ln(n)}$  which always grows faster than  $e^n$ , and therefore it will be faster than  $2^n$  as well ( $e > 2$ ).

therefore we have  $f(n) = 2^{n*} n^2 \leq n^n$  for  $n \geq 4$ ,  $c = 1$ . Thus  $f(n) = O(n)$

3.

For

$f(n) = n^{(1 + \sin(\pi * n))}$ , for the sub function  $1 + \sin(\pi * n)$ ,  $n \in \mathbb{N}$  it is worthy to note that this function could has maximum and minimum values

When  $n \bmod 2 \neq 0$  ( $n$  is odd number)

$$1 + \sin(\pi * n) = 0$$

When  $n \equiv \frac{1}{2} + 2 * \pi$

$$1 + \sin(\pi * n) = 2$$

Thus if  $n$  is an odd number

$$f(n) = n^0 = 1$$

$$g(n) = n$$

We know that  $g(n)$  grows faster than  $f(n)$

However if  $n \equiv \frac{1}{2} + 2 * \pi$

$$f(n) = n^2$$

$$g(n) = n$$

We know that  $n^2$  is definitely growing faster than  $n$ , therefore  $f(n)$  grows faster than  $g(n)$ . Thus, neither  $f(n) = O(g(n))$  nor  $g(n) = O(f(n))$ .