

Problem Analysis Of Stable Match

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Initially all $m \in M$ and $w \in W$ are free

While there is a man m who is free and hasn't proposed to every woman w for which $(m, w) \notin F$

 Choose such a man m

 Let w be the highest-ranked woman in m 's preference list
 to which m has not yet proposed

 If w is free then

(m, w) become engaged

 Else w is currently engaged to m'

 If w prefers m' to m then

m remains free

 Else w prefers m to m'

(m, w) become engaged

m' becomes free

 Endif

 Endif

Endwhile

Return the set S of engaged pairs

Common Problems

- ▶ What data structures are used for input and output ?
- ▶ How to find the unmatched SA efficiently?
- ▶ How to efficiently query the ranking of a SA in a student's preference list?
- ▶ Not to test the code sufficiently

What data structures are used for input and output?

- ▶ Analysis of Input and Output Formats
- SA's name → SA's Appearance No. (Map)
- Student's name → Student's Appearance No. (Map)
- SA's Appearance No. → SA's name (Array)
- Student's Appearance No. → Student's name (Array)
- Apparently, the preference list should be a two-dimensional array. Since the Appearance No. can be easily obtained from Map, it is possible to design the preference list as `int [][]`
- ▶ The output is a list of student names. The i-th SA is match the i-th student. Obviously output is OK using a string array.

How to find the unmatched SA efficiently?

- ▶ Queue or Stack: $O(1)$
- ▶ Initial, all SA are free and add to a queue
- ▶ Each iterator pop a SA from queue, try to match, If he can steal a students from another SA, who has to go back to queue.

How to find a student of the highest rank and not be tried match before for a SA?

- ▶ Simple solution: find from head to tail every time
- ▶ But if a SA was stole a student by another SA, he should find lower rank students from the stole one.
- ▶ We can use a array to store the current preference index of SA

How to efficiently query the ranking of a SA in a student's preference list?

- ▶ Simple solution: using a loop to find the rank of a SA according the SA's Appearance No. in the student's preference list. $O(n)$

- ▶ More efficiently solution:

- 1、 Maintain a reverse list of a student's preference list.

Index: SA's appearance No. \rightarrow value: SA's rank

Actually, we don't need SA's rank \rightarrow SA's appearance No.

- 2、 using map to store SA's appearance No. \rightarrow value: SA's rank

Data Structure List

- SA's name \rightarrow SA's *appearance No.* (Map)
- Student's name \rightarrow Student's *appearance No.* (Map)
- SA's *appearance No.* \rightarrow SA's name (Array)
- Student's *appearance No.* \rightarrow Student's name (Array)
- SA's preference list (int[][])

the first dimension: SA's *appearance No.* \rightarrow SA's preference list;

the second dimension: the rank of student \rightarrow student's *appearance No.*

- Student's preference reverse list (int[][])

the first dimension: student's *appearance No.* \rightarrow student's preference list;

the second dimension: SA's *appearance No.* \rightarrow the rank of SA

- Free SAs (Queue)
- Match status of student \rightarrow SA (Array)
- Match status of SA \rightarrow student (Array)
- When you update above variable , you should be full thought.

Pay Attention

- ▶ The problem of object copy, deep copy and shallow copy . (clone)