# 人工智能-搜索

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# 搜索

- ■深度优先搜索
- 宽度优先搜索
- A\* 搜索
- α β 搜索

#### 概述

早期的人工智能(AI)重视搜索算法的研究。例如,深度优先搜索(DFS)、宽度优先搜索(BFS)、minmax 搜索、 $\alpha - \beta$  搜索、 $A^*$  算法等。

因为许多 AI 任务可以通过定义状态空间并通过搜索技术来得到有效地解决。

## 状态空间

问题被表达成状态空间(State Space),然后再使用各种搜索技术在状态空间中进行搜索,以得到问题的解答。

本质上,状态空间由节点和边构成。节点表示问题的每一个状态,边表示一个节点到另一个节点 的合法移动。状态空间也定义了一个初始状态和 一个终止状态。

## 状态空间

状态空间可以采取图(包括网格 Grid)和树的形态。

## 几个经典的问题

旅行商(Traveling Salesman)问题

汉诺塔(Towers of Hanoi)问题

8 数码(8-Puzzle)问题

Tic Tac Toe 游戏

# 汉诺塔问题

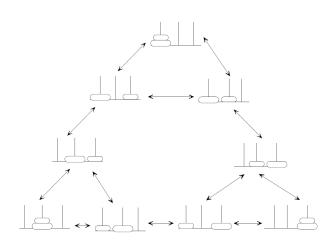


图 1-1: 汉诺塔的状态空间

## 8 数码问题

1	2	3
8		4
7	6	5

1	8	3
2	6	4
7		5

图 1-2:8 数码问题:(a) 初始状态; (b) 目标状态

## 8 数码问题

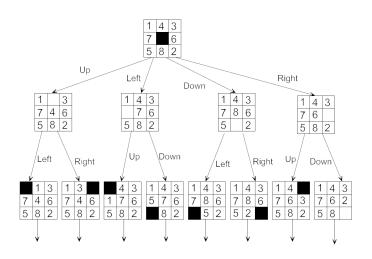


图 1-3:8 数码的状态空间

### Tic Tac Toe 游戏

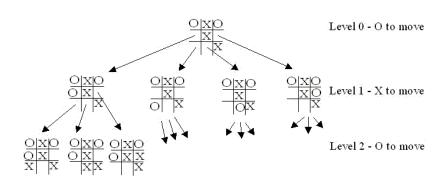


图 1-4: Tic Tac Toe 的状态空间

## 合法动作集

节点在进行状态转移生成子节点时,都要依据合法动作集。

例如,8数码问题的合法动作集是:LEFT、RIGHT、UP、DOWN。

状态空间的分支因子越大, 合法动作集也越大。 有些问题的分支因子非常大, 例如围棋。

## 一般遍历方式(DFS)

使用深度优先搜索对图或树中的节点进行排序:

- Preordering: 以节点首次被访问的次序进行 排序;
- Postordering: 以节点最后被访问的次序进 行排序;
- Reverse Postordering: Postordering 的反序;

## 二叉树遍历方式(DFS)

- Preordering;
- Inordering;
- Postordering;

还有一种是 Level-ordering, 它实际上是宽度优先搜索。

# Level-ordering: 宽度优先搜索

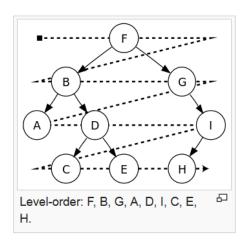


图 1-5: Level-ordering: 宽度优先搜索

# 算法伪代码: 递归实现

Input: A graph G and a vertex v of G Output: All vertices reachable from v labeled as discovered

```
procedure DFS(G, v):
label v as discovered

for all edges from v to w in G.adjacentEdges(v) do
    if vertex w is not labeled as discovered then
    recursively call DFS(G, w)
```

图 1-6: 深度优先搜索的递归实现

# 算法伪代码:非递归实现(Stack)

```
procedure DFS-iterative(G, v):
let S be a stack
S.push(v)
while S is not empty
v = S.pop()
if v is not labeled as discovered:
label v as discovered
for all edges from v to w in G.adjacentEdges(v) do
S.push(w)
```

图 1-7: 深度优先搜索的非递归实现(Stack)

# 示例 Graph

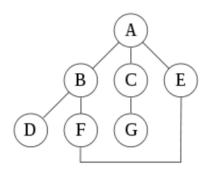


图 1-8: 示例 Graph

递归实现: A, B, D, F, E, C, G; 非递归实现: A, E, F, B, D, C, G;



# 算法伪代码:非递归实现(Open、Closed 表)

```
def dfs (in Start, out State)
open = [Start];
closed = []:
State = failure;
while (open <> []) AND (State <> success)
 begin
  remove the leftmost state from open, call it X:
  if X is the goal, then
   State = success
  else begin
   generate children of X;
   put X on closed
   eliminate the children of X on open or closed
   put remaining children on left end of open
  end else
endwhile
return State:
enddef
```

图 1-9: 深度优先搜索的非递归实现(Open、Closed 表)

## 一个例子

#### Graph/Tree

A to B and C
B to D and E
C to F and G
D to I and J
I to K and L

Start state: A Goal State: E, J

# 执行过程

	OPEN	CLOSED	х	X's Children	State
1.	Α	-			failure
2.			А	B, C	failure
3.	B, C	А			failure
4.	С	А	В	D, E	failure
5.	D, E, C	A, B			failure
6.	E, C	A, B	D	I, J	failure
7.	I, J, E, C	A, B, D			failure
8.	J, E, C	A, B, D	1	K, L	failure
9.	K, L, J, E, C	A, B, D, I			failure
10.	L, J, E, C	A, B, D, I,	К	none	failure
11.	J, E, C	A, B, D, I	L	none	failure
12.	E, C	A, B, D, I	J		success

图 1-10: 例子的执行过程

## 迭代加深的深度优先搜索

Depth First Search with Iterative Deepening, DFID

每次迭代有不同的深度上限值(深度值加1),该 算法可以确保找到最短路径。

## 算法伪代码

图 1-11: 迭代加深的深度优先搜索算法

## 执行过程

Iteration 1: search\_depth =1

	OPEN	CLOSED	х	X's Children	State
1.	А	-			failure
2.			А	B, C	failure
3.	B, C	A			failure
4.	С	A	В		failure
5.	С	A, B			failure
6.		A, B	С		failure
7.		A, B, C			failure

图 1-12: 例子执行过程(DFID)-深度1

## 执行过程

Iteration 2: search\_depth = 2

	OPEN	CLOSED	х	X's Children	State
1.	Α	-			failure
2.			А	B, C	failure
3.	B, C	А			failure
4.	С	А	В	D, E	failure
5.	D, E, C	A, B			failure
6.	E, C	A, B	D		failure
7.	E, C	A, B, D			failure
8.	С	A, B, D	Е		success

图 1-13: 例子执行过程 (DFID) - 深度 2

## DFS 应用

- Finding connected components
- Finding strongly connected components
- Topological sorting
- Generating words in order to plot the Limit Set of a Group
- Planarity testing
- Solving puzzles such as mazes
- Maze generation may use a randomized depth-first search

## 演示

Maze generation

播放视频文件 MAZE\_30x20\_DFS.ogv

## 宽度优先搜索

Breadth First Search, BFS

从一个节点开始,由近到远的方式逐层进行搜索,能够保证找到最短路径。

## 宽度优先搜索-示例

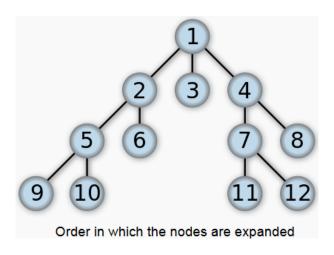


图 1-14: 宽度优先搜索算法的示例

# 非递归实现(Queue)

Input: A graph Graph and a starting vertex root of *G* 

Output: All vertices reachable from root labeled as explored

# 非递归实现(Queue)

```
1 Breadth-First-Search (Graph, root):
       for each node n in Graph:
           n.distance = INFINITY
           n.parent = NIL
       create empty queue Q
 8
       root.distance = 0
10
       O.enqueue (root)
11
12
       while Q is not empty:
13
14
           current = Q.dequeue()
1.5
16
           for each node n that is adjacent to current:
17
               if n.distance == INFINITY:
18
                    n.distance = current.distance + 1
19
                    n.parent = current
20
                    Q.enqueue(n)
```

#### 图 1-15: 宽度优先搜索的非递归实现(Queue)

# 非递归实现(Open、Closed 表)

```
def bfs (in Start, out State)
open = [Start];
closed = []:
State = failure;
while (open <> []) AND (State <> success)
 begin
  remove the leftmost state from open, call it X;
  if X is the goal, then
   State = success
  else begin
   generate children of X;
   put X on closed
   eliminate the children of X on open or closed
   put remaining children on right end of open
  end else
endwhile
return State;
enddef
```

图 1-16: 宽度优先搜索的非递归实现(Open、Closed 表)



#### 示例

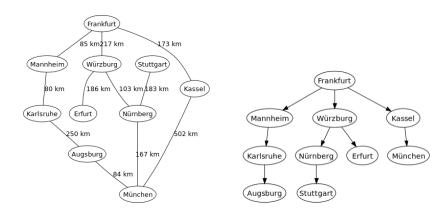


图 1-17: 示例:(a)Germany 城市 Graph; (b)Germany 城市 Tree

#### DFS v.s. BFS

- BFS 能够保证找到最短路径;
- DFS 在搜索时可能会"迷路",必须施加一个搜索深度限制;
- BFS 可能会碰到"组合爆炸"问题,当分支 因子很高时;
- DFS 搜索到的结果可能不是最优的;
- DFS 可能更有效率, 当分支因子很高时;

## 参考文献

- [1] Mark Watson. Practical Artificial Intelligence Programming in Java, 2005.
- [2] Web Page: State Space Representation and Search
- [3] Wikipedia: Depth First Search
- [4] Wikipedia: Tree Traversal
- [5] Wikipedia: Bread First Search