

ALGORITHMS

DATA STRUCTURE

Timing an algorithm

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ThS.T

```
long startTime = System.currentTimeMillis();           // record the starting time
/* (run the algorithm) */
long endTime = System.currentTimeMillis();             // record the ending time
long elapsed = endTime - startTime;                   // compute the elapsed time
```

Code Fragment 4.1: Typical approach for timing an algorithm in Java.

The “Big-Oh” Notation

- ▶ Read in book page 168
- ▶ Let $f(n)$ and $g(n)$ be functions mapping positive integers to positive real numbers. We say that $f(n)$ is $O(g(n))$ if there is a real constant $c > 0$ and an integer constant $n_0 \geq 1$ such that $f(n) \leq c \cdot g(n)$, for $n \geq n_0$
- ▶ The big-Oh notation allows us to ignore constant factors and lower-order terms and focus on the main components of a function that affect its growth
- ▶ Example: $5n^4 + 3n^3 + 2n^2 + 4n + 1$ is $O(n^4)$

SEARCH ALGORITHMS

ALGORITHMS

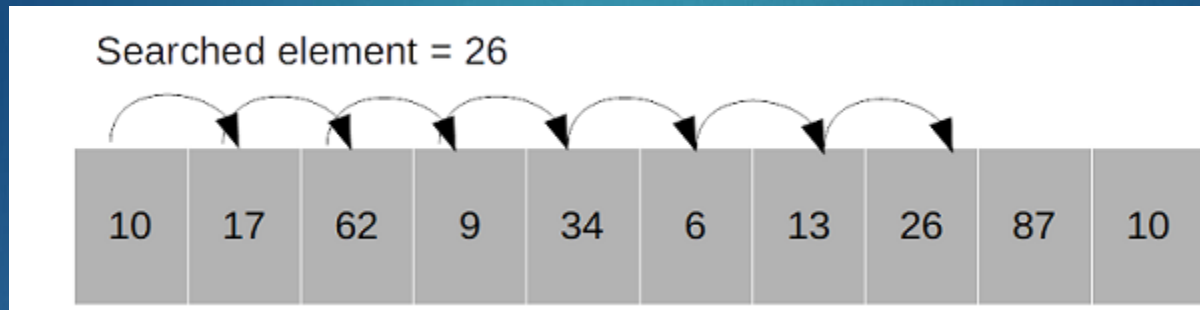
Linear Search

SEARCH ALGORITHMS

LINEAR SEARCH

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RULE

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- ▶ LinearSearch(A , size, target)
 - ▶ For i=0 to size -1
 - ▶ If(A[i] == target) return i
 - ▶ Else return -1

EXAMPLE

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45

45	34	64	55	67	12	57
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BEST
CASE

1	34	45	55	67	12	57
---	----	----	----	----	----	----

1	34	9	55	67	12	45
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WORST
CASE

RUNNING TIME

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- ▶ Best case: 1 comparable time
- ▶ Worst case: n comparable times
- ▶ AVG case: $(n+1) / 2$ comparable times
- ▶ $O(n)$

IMPLEMENT LINEAR SEARCH

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```
▶ public boolean linearSearching(int[] array, int  
    target){  
▶ for(...){  
▶ if(array[i]== target){  
▶ //TODO  
▶ }  
}
```

USING NON-RECURSIVE TO PROGRAM LINEAR SEARCH

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```
▶ public boolean linearSearching(int[] array, int  
    target){  
    for(...){  
        If(array[i]== target){  
            //TODO  
        }  
    }  
}
```

USING RECURSIVE TO PROGRAM LINEAR SEARCH

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```
▶ public boolean linearSearching(int[] array, int target, int  
    cursor)  
    {  
        if(cursor == array.length - 1){  
            // STOP CONDITION  
        }else{  
            // RECURSIVE CONDITION  
            // if array[cursor] == target  
            // if array[cursor] != target  
        }  
    }
```

Binary Search

SEARCH ALGORITHM

HOW IT WORK ?

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- Compare x with the middle element.
- If x matches with middle element, we return the mid index.
- Else If x is greater than the mid element, then x can only lie in right half subarray after the mid element. So we recur for right half.
- Else (x is smaller) recur for the left half.

EXAMPLE

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45

1	3	45	67	76	86	91	134
0	1	2	3	4	5	6	7

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Step1 : size of array $\Rightarrow N = 8$; high = $N - 1 = 7$, low = 0, mid = $(\text{low} + \text{high}) / 2 = 7 / 2 = 3$

Array[3] = 67 > 45 \Rightarrow high = mid - 1 = 3 - 1 = 2

Step2 : high = 2, low = 0, mid = $(\text{low} + \text{high}) / 2 = 2 / 2 = 1$

Array[1] = 3 < 45 \Rightarrow low = mid + 1 = 1 + 1 = 2

Step3: high = 2, low = 2 \Rightarrow mid = $(\text{low} + \text{high}) / 2 = 4 / 2 = 2$

Array[2] = 45 = 45 \Rightarrow stop

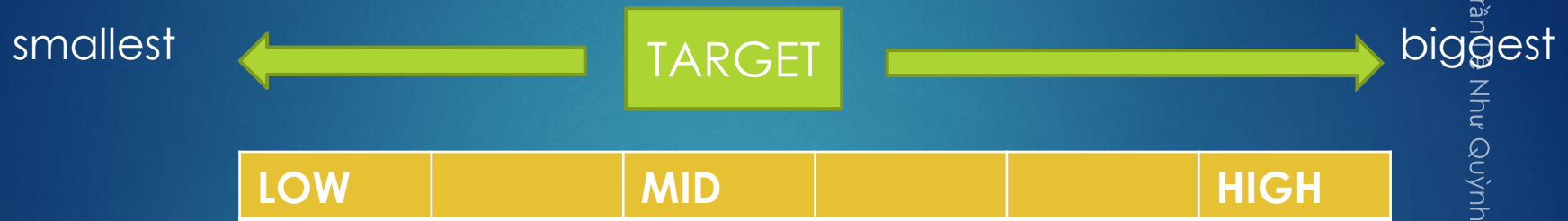
RUNNING TIME

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- ▶ Best: $\text{ceil}(\log_2(n)) + 1$
- ▶ Worst: $\text{floor}(\log_2(n)) + 1$
- ▶ AVG: $\text{approx.} \log_2(n) + 1$
- ▶ $O(\log(n))$

USING RECURSIVE TO IMPLEMENT BINARY SEARCH



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- Just using for sorted array
- Low = begin of array, high = end of array, $mid = (high + low) / 2$
- STOP CONDITION
 - Target equals with value of element at middle
- CONTINUOUS CONDITION
 - Target larger than mid \rightarrow Recursive with low = mid + 1, high not change
 - Target smaller than mid \rightarrow Recursive with low not change, high = mid - 1

USING RECURSIVE TO IMPLEMENT BINARY SEARCH

```
/**
 * Returns true if the target value is found in the indicated portion of the data array.
 * This search only considers the array portion from data[low] to data[high] inclusive.
 */
public static boolean binarySearch(int[ ] data, int target, int low, int high) {
    if (low > high)
        return false; // interval empty; no match
    else {
        int mid = (low + high) / 2;
        if (target == data[mid])
            return true; // found a match
        else if (target < data[mid])
            return binarySearch(data, target, low, mid - 1); // recur left of the middle
        else
            return binarySearch(data, target, mid + 1, high); // recur right of the middle
    }
}
```

NO USING RECURSIVE TO IMPLEMENT BINARY SEARCH

```
► BinarySearch(A[0..N-1], value) {  
    low = 0  
    high = N - 1  
    while (low <= high) {  
        mid = (high + low) / 2  
        if (A[mid] = value) return value;  
        else if (A[mid] > value) high = mid - 1  
        else  
            low = mid + 1  
    }  
    return -1111; // no element in array equals value  
► }
```

Exercise

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- ▶ Coding Linear search
- ▶ Coding Binary search

Challenge 1

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- ▶ Using Binary search vs linear search, run by step
 - ▶ Finding target number = 7
 - ▶ Finding target number = 5

4	6	7	5	8	9	10	12	15
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COMPARE
BINARY SEARCH
AND LINEAR SEARCH
IN THIS CASE

Challenge 2

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- ▶ Using Binary search vs linear search, run by step
 - ▶ Finding target number = 120
 - ▶ Finding target number = 56

23	45	56	67	78	92	101	120	135
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COMPARE
BINARY SEARCH
AND LINEAR SEARCH
IN THIS CASE