

# HW2

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## 1

**1.1**     $L_1 = \{ V_1, \Sigma_1, R_1, S_1, \}$   
           $V_1 = \{ S, a, b \}$   
           $\Sigma_1 = \{ a, b \}$   
           $R_1 = \{ S \rightarrow SbSbSaS | SbSaSbS | SaSbSbS | e \}$   
           $S_1 = \{ S \}$

**1.2**     $L_2 = \{ V_2, \Sigma_2, R_2, S_2, \}$   
           $V_2 = \{ T, a, b \}$   
           $\Sigma_2 = \{ a, b \}$   
           $R_2 = \{ T \rightarrow aTb | aaTb | e \}$   
           $S_2 = \{ T \}$

1.3  $M = \{ K, \Sigma, \Gamma, \Delta, s, F \}$

$$K = \{ p, q \}$$

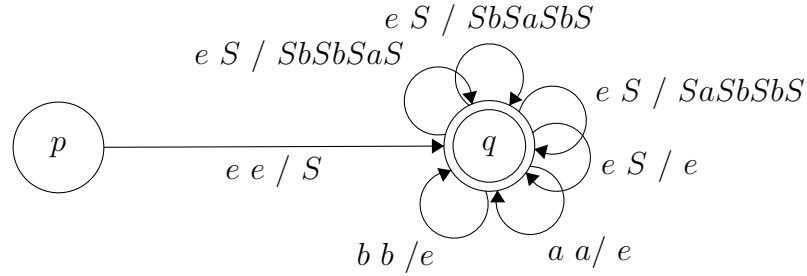
$$\Sigma = \{ S, a, b \}$$

$$\Gamma = \{ S, a, b \}$$

$$\Delta = \{ (p, e, e)(q, S), \\ (p, e, S)(q, SbSbSaS), \\ (q, e, S)(q, SbSaSbS), \\ (q, e, S)(q, SaSbSbS), \\ (q, e, S)(q, e), \\ (q, e, a)(q, e), \\ (q, e, b)(q, e) \}$$

$$s = \{ p \}$$

$$F = \{ q \}$$



1.4  $L_1$  and  $L_2$  are both context free grammars, with disjoint sets of non terminals. Let  $K$  be a new symbol and

$$L_3 = \{ (V_1 \cup V_2 \cup K), (\Sigma_1 \cup \Sigma_2), (R_1 \cup R_2 \cup (K \rightarrow S, K \rightarrow T)), K \}$$

$$V_3 = \{ K, S, T, a, b \}$$

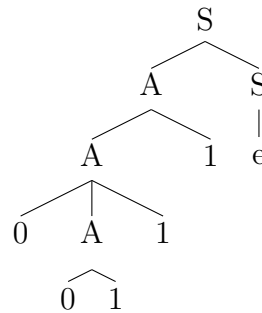
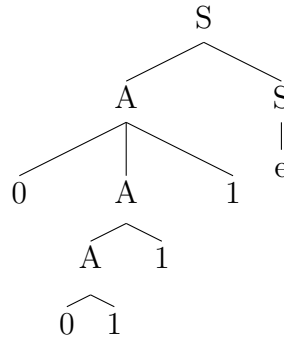
$$\Sigma_3 = \{ a, b \}$$

$$R_3 = \{ K \rightarrow S | T, T \rightarrow aTb | aaTb | e, S \rightarrow SbSbSaS | SbSaSbS | SaSbSbS | e \}$$

$$S_3 = \{ K \}$$

## 2

- 2.1 If a word  $w$  has more than one parse tree that represents it then language is ambiguous. And as showed below we can draw 2 different parse trees for word "00111" with given CFL.



- 2.2  $\{S \rightarrow AS|e, A \rightarrow 0C1, C \rightarrow 0C1 \mid B, B \rightarrow B1 \mid e\}$

- 2.3  $S \rightarrow AS, A \rightarrow 0C1, C \rightarrow 0C1, C \rightarrow B, B \rightarrow B1, B \rightarrow e, S \rightarrow e$

