



PRACTICAL FILE OF OPERATING SYSTEM

BTech: III Year

Department of Computer Science & Information Technology

Name of the Student : Burhanuddin Sahil

Branch & section : CSIT-1

Roll No. : 0827CI201055

Year : 2022

**Department of Computer Science & Information Technology
AITR, Indore,**

ACROPOLIS INSTITUTE OF TECHNOLOGY & RESEARCH, INDORE

Department of Computer Science & Information Technology

Certificate

This is to certify that the experimental work entered in this journal as per the BTech III year syllabus prescribed by the RGPV was done by **Prof.Vandana Kate** in 5th semester in the Laboratory of this institute during the academic year July 2022 - Dec 2022

Signature of Head

Signature of the Faculty

ACROPOLIS INSTITUTE OF TECHNOLOGY & RESEARCH, INDORE

GENERAL INSTRUCTIONS FOR LABORATORY CLASSES

DO'S

- Without Prior permission do not enter into the Laboratory.
- While entering into the LAB students should wear their ID cards.
- The Students should come with proper uniform.
- Students should maintain silence inside the laboratory.
- After completing the laboratory exercise, make sure to shutdown the system properly.

DONT'S

- Students bringing the bags inside the laboratory.
- Students using the computers in an improper way.
- Students scribbling on the desk and mishandling the chairs.
- Students using mobile phones inside the laboratory.
- Students making noise inside the laboratory.

SYLLABUS

CSIT-502 – Operating System

Branch: Computer Science Information Technology V Semester

Course: CSIT 502 Operating System

Unit I

Introduction to Operating Systems, Evaluation of OS, Types of operating Systems, system protection, Operating system services, Operating System structure, System Calls and System Boots, Operating System design and implementation, Spooling and Buffering.

Unit II

Basic concepts of CPU scheduling, Scheduling criteria, Scheduling algorithms, algorithm evaluation, multiple processor scheduling. Process concept, operations on processes, threads, inter process communication, precedence graphs, critical section problem, semaphores, classical problems of synchronization

Unit III

Deadlock problem, deadlock characterization, deadlock prevention, deadlock avoidance, deadlock detection, recovery from deadlock, Methods for deadlock handling. Concepts of memory management, logical and physical address space, swapping, Fixed and Dynamic Partitions, Best-Fit, First-Fit and Worst Fit Allocation, paging, segmentation, and paging combined with segmentation.

Unit IV

Concepts of virtual memory, Cache Memory Organization, demand paging, page replacement algorithms, allocation of frames, thrashing, demand segmentation, Role of Operating System in Security, Security Breaches, System Protection, and Password Management.

Unit V

Disk scheduling, file concepts, File manager, File organization, access methods, allocation methods, free space managements, directory systems, file protection, file organization & access mechanism, file sharing implement issue, File Management in Linux, introduction to distributed systems.

HARDWARE REOUIREMENTS:

Processors - 2.0 GHz or Higher
RAM - 256 MB or Higher
Hard Disk - 20 GB or Higher

SOFTWARE REQUIREMENTS:

Linux: Ubuntu / OpenSUSE / Fedora / Red Hat / Debian / Mint OS
WINDOWS: XP/7
Linux could be loaded in individual PCs.

RATIONALE:

The purpose of this subject is to cover the underlying concepts Operating System .This syllabus provides a comprehensive introduction of Operating System, Process Management, Memory Management, File Management and I/O management.

PREREQUISITE:

The students should have general idea about Operating System Concept, types of Operating System and their functionality.

Lab Plan

Operating System

CSIT-502

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4.	Program to implement Banker's algorithm	17-23
5.	Program to implement FIFO page replacement algorithm.	24-28
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Name of Department :- CSIT Name of Laboratory :- Operating System

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4.		Program to implement Banker's algorithm	17-23		
5.		Program to implement FIFO page replacement algorithm.	24-28		
6.		Program to implement LRU page replacement algorithm	29-32		
7.		Program to implement Disk Scheduling (FCFS) algorithm	33-36		
8.		Program to implement Disk Scheduling (SSTF) algorithm	37-40		
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Experiment-1

FCFS SCHEDULING

Name of Student : Burhanuddin Sahil		Class : CSIT-1
Enrollment No : 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c++ program to implement the FCFS SCHEDULING.

FACILITIES REQUIRED

a) Facilities Required Doing The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	WINDOWS XP/7	

b) Concept of FCFS:

- Jobs are executed on first come, first serve basis.
- Easy to understand and implement.
- Poor in performance as average wait time is high.

Process	Arrival Time	Execute Time	Service Time
P0	0	5	0
P1	1	3	5
P2	2	8	8
P3	3	6	16



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Set the waiting of the first process as '0' and its burst time as its turn around time

Step 5: for each process in the Ready Q calculate

(a) Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n-1)

(b) Turn around time for Process(n)= waiting time of Process(n)+ Burst time for process(n)

Step 6: Calculate

(a) Average waiting time = Total waiting Time / Number of process

(b) Average Turnaround time = Total Turnaround Time / Number of process

Step 7: Stop the process

d) Program :

```
#include <bits/stdc++.h>
using namespace std;

class fcfs
{
    public :

    void completion_time( int burst_time [] ,int n ,int cmp_time[])
    {
        cmp_time[0] = burst_time[0];

        for(int i=1 ; i<n; i++)
        {
            cmp_time[i] = cmp_time[i-1]+burst_time[i];
        }
    }

    void turnaround_time(int cmp_time[] ,int n ,int arr_time[] , int turn_time[] )
    {
        for(int k=0; k<n ; k++)
        {
            turn_time[k] = cmp_time[k] - arr_time[k];
        }
    }

    void waiting_time(int burst_time[] ,int n ,int waiting_time[] , int turn_time[] )
    {
        for(int k=0; k<n ; k++)
        {
            waiting_time[k] = turn_time[k] - burst_time[k];
        }
    }

};
```

```

int main()
{
    fcfs ob;

    int bt[4] = {3,4,6,7};
    int ct[4];
    int tt[4];
    int wt[4];
    int at[4] = {3,1,0,2};
    sort(at ,at+4) ;

    ob.completion_time(bt , 4 , ct);
    ob.turnaround_time(ct , 4 , at, tt);
    ob.waiting_time(bt , 4 , wt , tt);

    for(int i = 0 ; i< 4 ;i++)
    {
        cout<<i<<" "<<at[i]<<" "<<bt[i]<<" "<<ct[i]<<" "<<tt[i]<<" "<<wt[i]<<endl;
    }
    return 0;
}

```

e) Output:

```
C:\Users\burha\OneDrive\Documents\operating system\FCFS\First_Come_First_Serve_Algo.exe
0 0 3 3 0
1 1 4 7 6 2
2 2 6 13 11 5
3 3 7 20 17 10
-----
Process exited after 0.07587 seconds with return value 0
Press any key to continue . . .
```

f) Result:

Average Waiting Time

Average Turnaround Time

Experiment-2

SJF Scheduling

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c++ program to implement SJF CPU Scheduling Algorithm.

FACILITIES REQUIRED

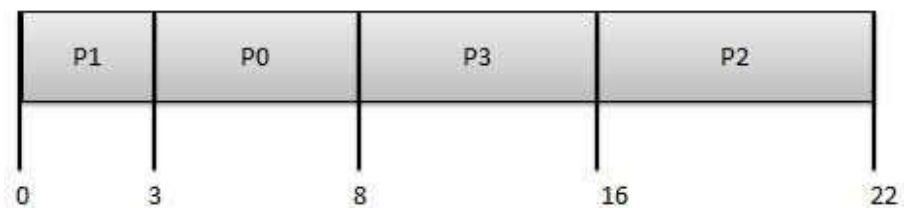
a) Facilities Required Doing The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	WINDOWS XP/7	

b) Concept of SJF:

- Best approach to minimize waiting time.
- Processer should know in advance how much time process will take.

Process	Arrival Time	Execute Time	Service Time
P0	0	5	0
P1	1	3	3
P2	2	8	8
P3	3	6	16



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Start the Ready Q according the shortest Burst time by sorting according to lowest to highest burst time.

Step 5: Set the waiting time of the first process as '0' and its turnaround time as its burst time.

Step 6: For each process in the ready queue, calculate

(c) Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n-1)

(d) Turnaround time for Process(n)= waiting time of Process(n)+ Burst time for process(n)

Step 7: Calculate

(c) Average waiting time = Total waiting Time / Number of process

(d) Average Turnaround time = Total Turnaround Time / Number of process

Step 8: Stop the process

d) Program:

```
#include <bits/stdc++.h>
using namespace std;
struct data
{
    int pid;
    int bt;
    int at;
    int wt;
    int tat;
};
bool cmp( data x , data y)
{
    return ( x.bt<y.bt);
}

int main()
{
    data a[30];

    int i,k,n;
    float wtavg , tatavg;

    cout<<"Enter the Number of process"<<endl;
    cin>>n;
    for(i=0;i<n;i++)
    {
        a[i].pid =i;
        a[i].at=0;
        cout<<"Enter The Burst time for the process "<<i<<endl;
        cin>>a[i].bt;
    }

    sort( a , a+n , cmp);
    a[0].wt =0;

    a[0].tat = a[0].bt;
```

```
    for(i=1;i<n;i++)
    {
        a[i].wt = a[i-1].wt + a[i-1].bt;
        a[i].tat = a[i-1].tat + a[i].bt;
    }
    cout<<endl<<"P.id\tBt\tWt\tTat"<<endl;
    for(i=0;i<n;i++)
    {
        cout<<a[i].pid<<"\t"<<a[i].bt<<"\t"<<a[i].wt<<"\t"<<a[i].tat<<endl;
    }
```


e) Output:

```
C:\Users\burha\OneDrive\Documents\operating system\SJF(With Same(0) Arrival Time)\SJF(With Same(0) Arrival Time).exe
Enter the Number of process
3
Enter The Burst time for the process 0
3
Enter The Burst time for the process 1
2
Enter The Burst time for the process 2
1
P.id  Bt   Wt   Tat
2     1   0    1
1     2   1    3
0     3   3    6

-----
Process exited after 4.089 seconds with return value 0
Press any key to continue . . .
```

f) Result:

Average Waiting Time

Average Turnaround Time

Experiment-3

ROUND ROBIN Scheduling

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c program to implement Round Robin scheduling.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

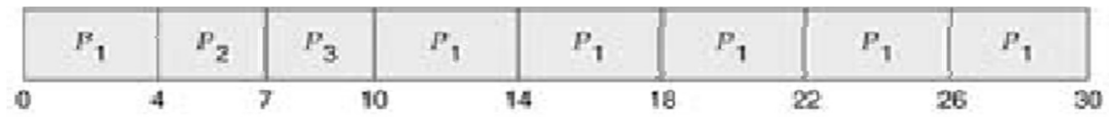
S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of Round Robin Scheduling:

This Algorithm is designed especially for time-sharing systems. A small unit of time, called time slices or **quantum** is defined. All runnable processes are kept in a circular queue. The CPU scheduler goes around this queue, allocating the CPU to each process for a time interval of one quantum. New processes are added to the tail of the queue. The CPU scheduler picks the first process from the queue, sets a timer to interrupt after one quantum, and dispatches the process. If the process is still running at the end of the quantum, the CPU is preempted and the process is added to the tail of the queue. If the process finishes before the end of the quantum, the process itself releases the CPU voluntarily. Every time a process is granted the CPU, a **context switch** occurs, this adds overhead to the process execution time.

	Burst
Process	Time
P_1	24
P_2	3

P_2	3
Average	



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue and time quantum (or) time slice

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Calculate the no. of time slices for each process where

$$\text{No. of time slice for process}(n) = \text{burst time process}(n) / \text{time slice}$$

Step 5: If the burst time is less than the time slice then the no. of time slices =1.

Step 6: Consider the ready queue is a circular Q, calculate

(a) Waiting time for process(n) = waiting time of process(n-1)+ burst time of process(n-1) + the time difference in getting the CPU from process(n-1)

(b) Turn around time for process(n) = waiting time of process(n) + burst time of process(n)+ the time difference in getting CPU from process(n).

Step 7: Calculate

(g) Average waiting time = Total waiting Time / Number of process

(h) Average Turnaround time = Total Turnaround Time / Number of process

Step 8: Stop the process

d) Program:

```
#include <bits/stdc++.h>
using namespace std;

void RoundRobin() {
    int n, qt;
    cout << "enter number of processes" << endl;
    cin >> n;

    int p[n];
    int Bt[n];
    for (int i = 0; i < n; i++) {
        p[i] = i + 1;
    }

    for (int i = 0; i < n; i++) {
        cout << "enter burst of processes" << endl;
        cin >> Bt[i];
    }

    cout << "enter quantum time" << endl;
    cin >> qt;
    cout << "Ready Queue :- " << endl;
    int tot = 0;
    for (int i = 0; i < n; i++) {
        tot = tot + Bt[i];
    }
    int x = 0;


    while (x <= tot) {

        for (int i = 0; i < n; i++) {
            if (Bt[i] > qt) {
                cout << p[i] << " ";
                x = x + qt;
                Bt[i] = Bt[i] - qt;
            } else if (Bt[i] <= qt and Bt[i] != 0) {
                x = x + qt;
                cout << p[i] << " ";
            }
        }
    }
}
```

```
        Bt[i] = 0;
    }
}
}
}
int main() {

    RoundRobin();
    return 0;
}
```

e) Output:

 C:\Users\burha\OneDrive\Documents\operating system\Round Robin\Round Robin.exe

```
enter number of processes
4
enter burst of processes
6
enter burst of processes
3
enter burst of processes
12
enter burst of processes
4
enter quantum time
2
Ready Queue :-
1 2 3 4 1 2 3 4 1 3 3 3 3
-----
Process exited after 18.47 seconds with return value 0
Press any key to continue . . .
```

f) Result:

Average Waiting Time

Average Turnaround Time

Experiment-4

BANKER ALGORITHM

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c program to implement deadlock avoidance & Prevention by using Banker's Algorithm.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of BANKER'S Algorithm:

The Banker's algorithm is a resource allocation and deadlock avoidance algorithm that tests for safety by simulating the allocation of predetermined maximum possible amounts of all resources, and then makes an "s-state" check to test for possible deadlock conditions for all other pending activities, before deciding whether allocation should be allowed to continue.

- Always keep so many resources that satisfy the needs of at least one client
- Multiple instances.
- Each process must a priori claim maximum use.
- When a process requests a resource it may have to wait.
- When a process gets all its resources it must return them in a finite amount of time.

c) Algorithm:

1. Start the program.
2. Get the values of resources and processes.
3. Get the avail value.
4. After allocation find the need value.
5. Check whether it's possible to allocate.
6. If it is possible then the system is in safe state.
7. Else system is not in safety state.
8. If the new request comes then check that the system is in safety.
9. Or not if we allow the request.
10. Stop the program.

d) Program:

```
#include<iostream>

using namespace std;

int alloc[50][50];

int maxi[50][50];

int need[50][50];

int avail[50];

int check_safety(int j,int nr)
{
    for(int i=0;i<nr;i++)
    {
        if(need[j][i]>avail[i])
            return 0;
    }
    return 1;
}

int check(bool a[],int n)
{
    for(int i=0;i<n;i++)
    {
        if(a[i]==false)
```

```

        return 0;

    }

    return 1;

}

int main()

{

    int np=100;

    int nr=100;

    cout<<"\nEnter the no of processes : ";

    cin>>np;

    cout<<"\nEnter the no of resources : ";

    cin>>nr;

    cout<<"\nEnter the allocation data : \n";

    for(int i=0;i<np;i++)

    for(int j=0;j<nr;j++)

    cin>>alloc[i][j];

    cout<<"\nEnter the requirement data : \n";

    for(int i=0;i<np;i++)

    for(int j=0;j<nr;j++)

    cin>>maxi[i][j];

```

```

for(int i=0;i<np;i++)

for(int j=0;j<nr;j++)

need[i][j]=maxi[i][j]-alloc[i][j];

cout<<"\nEnter the availability matrix : \n";

for(int i=0;i<nr;i++)

cin>>avail[i];

int ex_it=nr;

int flg;

bool completed[np];

while(10)

{

    for(int i=0;i<np;i++)

    {

        if(!completed[i] && check_safety(i,nr))

        {

            for(int j=0;j<nr;j++)

            avail[j]+=alloc[i][j];

        }

        completed[i]=true;

```

```

    }

    flg=check(completed,np);

    ex_it--;

    if(flg==1 || ex_it==0)

        break;
}

cout<<"\nThe final availability matrix \n";

for(int i=0;i<nr;i++)

    cout<<avail[i]<<" ";

cout<<"\n ----- Result ----- \n";

if(flg==1)

    cout<<"There is no deadlock";

else

    cout<<"Sorry there is a possibility of deadlock";

return 0;

}

```

e) Output:

C:\Users\burha\OneDrive\Documents\operating system\New folder\Bankers problem.exe

```
Enter the no of processes : 3
Enter the no of resources : 2
Enter the allocation data :
1 2 5 2 3 5
Enter the requirement data :
0 2 3 2 4 7
Enter the availability matrix :
2 4 0 5 7 1
The final availability matrix
11 13
----- Result -----
There is no deadlock
-----
Process exited after 45.89 seconds with return value 0
Press any key to continue . . .
```

f) Result:

The Sequence is :-

Experiment-5

FIFO PAGE REPLACEMENT

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement page replacement algorithm FIFO.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Fifo Page Replacement:

- Treats page frames allocated to a process as a circular buffer:
- When the buffer is full, the oldest page is replaced. Hence first-in, first-out: A frequently used page is often the oldest, so it will be repeatedly paged out by FIFO. Simple to implement: requires only a pointer that circles through the page frames of the process.

reference string

7 0 1 2 0 3 0 4 2 3 0 3 2 1 2 0 1 7 0 1

7	7	7	2	2	2	4	4	4	0	0	0	7	7	7
	0	0	0	3	3	3	2	2	2	1	1	1	0	0
		1	1	1	0	0	0	3	3	3	2	2	2	1

page frames

- FIFO Replacement manifests Belady's Anomaly:

more frames \Rightarrow more page faults

- Reference string: 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5()

3 Frames:-9 page fault

4 Frames: - 10 page fault

C) Algorithm:

Step 1: Create a queue to hold all pages in memory

Step 2: When the page is required replace the page at the head of the queue

Step 3: Now the new page is inserted at the tail of the queue

d) Program:

```
#include <bits/stdc++.h>
using namespace std;

int PageFaults(int pages[], int n, int frames)
{
    unordered_set <int> set;

    queue <int> indexes;

    int countPageFaults = 0;

    for (int i=0; i < n; i++)
    {
        if (set.size() < frames)
        {
            if (set.find(pages[i]) == set.end())
            {
                set.insert(pages[i]);

                countPageFaults++;

                indexes.push(pages[i]);
            }
        }
        else
        {
            if (set.find(pages[i]) == set.end())
            {
                int val = indexes.front();
                indexes.pop();
```

```

        set.erase(val);

        set.insert(pages[i]);
        indexes.push(pages[i]);

        countPageFaults++;
    }
}

return countPageFaults;
}

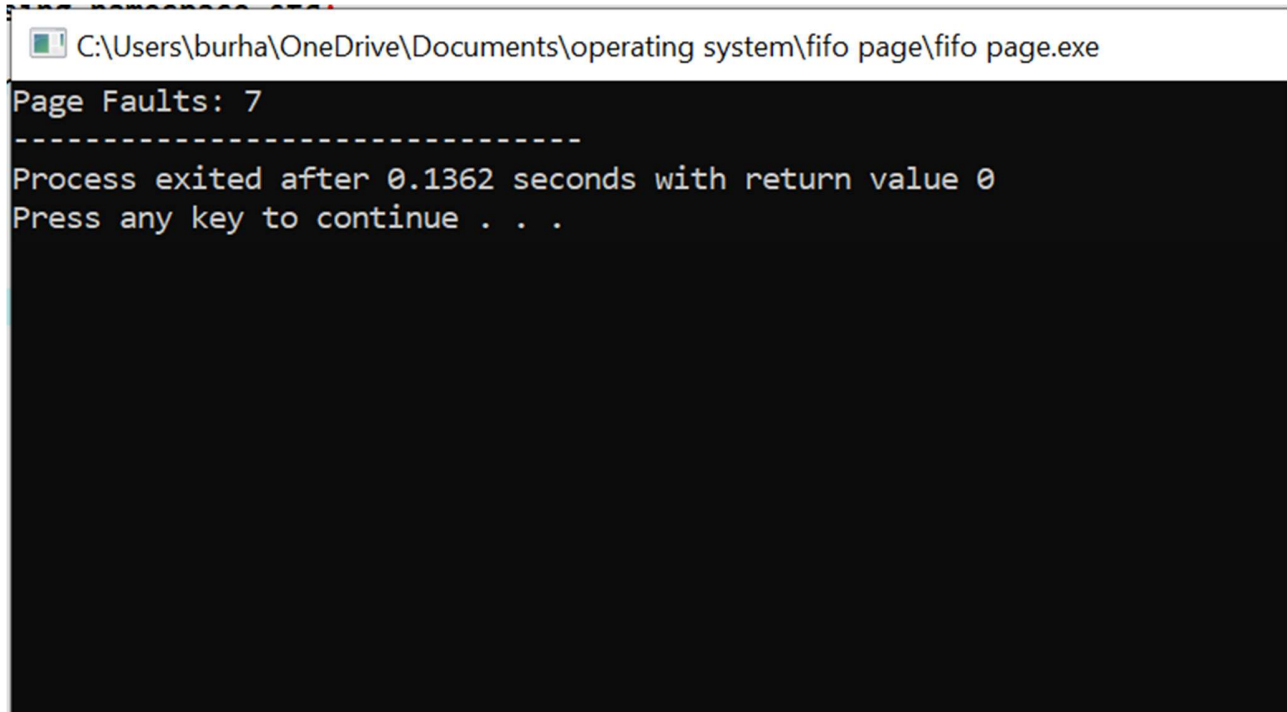
int main()
{
    int pages[] = {7,0,1,2,0,3,0,4,2,3,0,3,2};
    int n = sizeof(pages)/sizeof(pages[0]);
    int frames = 4;

    cout << "Page Faults: " << PageFaults(pages, n, frames);

    return 0;
}

```

e) Output:



```
C:\Users\burha\OneDrive\Documents\operating system\fifo page\fifo page.exe
Page Faults: 7
-----
Process exited after 0.1362 seconds with return value 0
Press any key to continue . . .
```

f) Result:

No. of page faults.....

Experiment-6

LRU PAGE REPLACEMENT

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement page replacement algorithm LRU.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept of LRU Algorithm:

Pages that have been heavily used in the last few instructions will probably be heavily used again in the next few. Conversely, pages that have not been used for ages will probably remain unused for a long time. when a page fault occurs, throw out the page that has been unused for the longest time. This strategy is called LRU (Least Recently Used) paging.

Page reference stream:

1 2 3 2 1 5 2 1 6 2 5 6 3 1 3 6 1 2 4 3

1 1 1 1 3 2 1 5 2 1 6 2 5 6 6 1 3 6 1 2

2 2 3 2 1 5 2 1 6 2 5 6 3 1 3 6 1 2 4

3 2 1 5 2 1 6 2 5 6 3 1 3 6 1 2 4 3

* * * * * * * * * *

LRU

Total 11 page faults

c) Algorithm:

Step 1: Create a queue to hold all pages in memory

Step 2: When the page is required replace the page at the head of the queue

Step 3: Now the new page is inserted at the tail of the queue

Step 4: Create a stack

Step 5: When the page fault occurs replace page present at the bottom of the stack

d) Program:

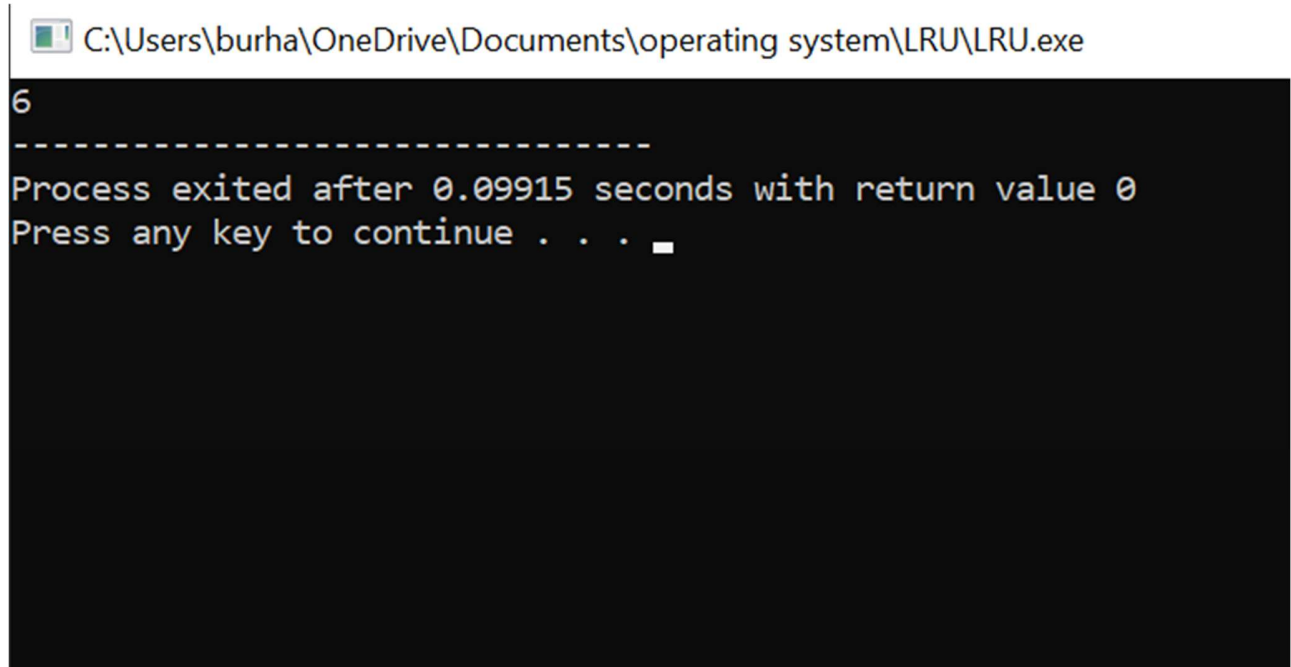
```
#include <iostream>
#include <bits/stdc++.h>

using namespace std;

int pageFault(int pages[], int n, int memcapacity){
    int pagefault = 0;
    vector <int> v;
    int i;
    for(int i=0;i<=n;i++){
        auto it= find(v.begin(), v.end(), pages[i]);
        if(it==v.end()){
            if(v.size()==memcapacity){
                v.erase(v.begin());
            }
            v.push_back(pages[i]);
            pagefault++;
        }
        else{
            v.erase(it);
            v.push_back(pages[i]);
        }
    }
    return pagefault;
}

int main() {
    int pages[] = {7,0,1,2,0,3,0,4,2,3,0,3,2};
    int n=13;
    int memcapacity = 4;
    cout<<pageFault(pages, n, memcapacity);
    return 0;
}
```

e) Output:



```
C:\Users\burha\OneDrive\Documents\operating system\LRU\LRU.exe
6
-----
Process exited after 0.09915 seconds with return value 0
Press any key to continue . . .
```

f) Result:

No. of pages faults

Experiment-7

FCFS Disk Scheduling Algorithm

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement FCFS Disk Scheduling Algorithm

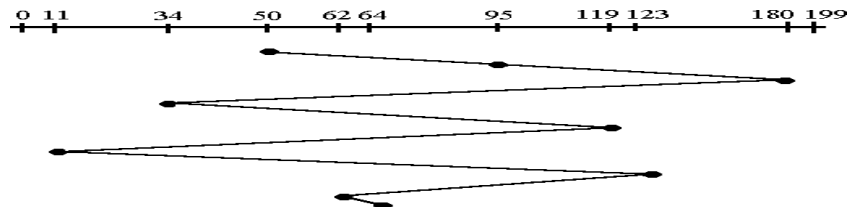
FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept of FCFS Disk Scheduling Algorithm:

All incoming requests are placed at the end of the queue. Whatever number that is next in the queue will be the next number served. Using this algorithm doesn't provide the best results. To determine the number of head movements you would simply find the number of tracks it took to move from one request to the next. For this case it went from 50 to 95 to 180 and so on. From 50 to 95 it moved 45 tracks. If you tally up the total number of tracks you will find how many tracks it had to go through before finishing the entire request. In this example, it had a total head movement of 640 tracks. The disadvantage of this algorithm is noted by the oscillation from track 50 to track 180 and then back to track 11 to 123 then to 64. As you will soon see, this is the worse algorithm that one can use.



c) Algorithm:

Step 1: Create a queue to hold all requests in disk

Step 2: Move the head to the request in FIFO order (Serve the request first that came first)

Step 3: Calculate the total head movement required to serve all request.

d) Program:

```
#include<bits/stdc++.h>
using namespace std;

class Diskscheduling
{
    public :
    float FCFS ( int arr[] , int n , int init)
    {
        float ans = abs (init - arr[0]);
        for(int i=1;i<n;i++)
        {
            ans = ans + abs(arr[i]-arr[i-1]);
        }

        return (ans/n);
    }
};

int main()
{
    int n;
    cout<<"Enter the No. of disk"<<endl;
    cin>>n;
    int disk[n];
    cout<<"Enter the disk sequence"<<endl;
    for(int i=0;i<n;i++)
    {
        cin>>disk[i];
    }

    cout<<"Enter the Position of Head"<<endl;
    int init;
    cin>>init;

    Diskscheduling obj;
    cout<<endl<<"Average :- "<<obj.FCFS(disk,n,init);

}
```

e) Output:

```
C:\Users\burha\OneDrive\Documents\operating system\Disk Scheduling\Disk Scheduling.exe
Enter the No. of disk
9
Enter the disk sequence
55
58
39
18
90
160
150
38
184
Enter the Position of Head
100

Average :- 55.3333
-----
Process exited after 43.83 seconds with return value 0
Press any key to continue . . .
```

f) Result:

Total Head Movement Required Serving All Requests

Experiment-8

SSTF Disk Scheduling Algorithm

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch:2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement SSTF Disk Scheduling Algorithm

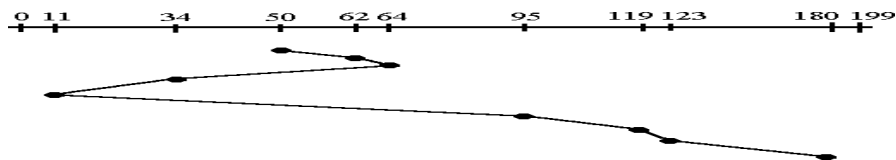
FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept of SSTF Disk Scheduling Algorithm:

In this case request is serviced according to next shortest distance. Starting at 50, the next shortest distance would be 62 instead of 34 since it is only 12 tracks away from 62 and 16 tracks away from 34. The process would continue until all the process are taken care of. For example the next case would be to move from 62 to 64 instead of 34 since there are only 2 tracks between them and not 18 if it were to go the other way. Although this seems to be a better service being that it moved a total of 236 tracks, this is not an optimal one. There is a great chance that starvation would take place. The reason for this is if there were a lot of requests close to each other the other requests will never be handled since the distance will always be greater.



c) Algorithm:

Step 1: Create a queue to hold all requests in disk

Step 2: Calculate the shortest seek time every time before moving head from current head position

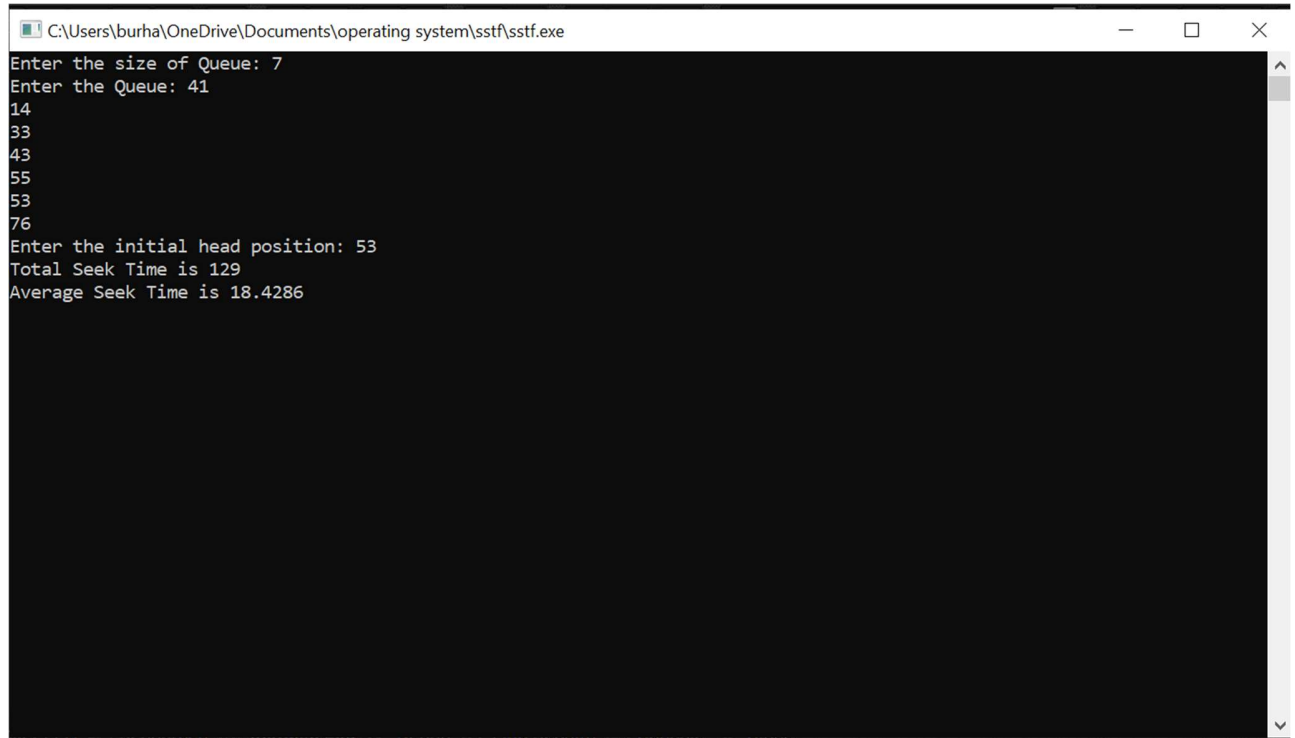
Step 3: Calculate the total head movement required to serve all request.

d) Program:

```
#include<iostream>
#include<conio.h>
#include<math.h>
using namespace std;
int main()
{
    int queue[100],t[100];
    int head,seek=0,n,i,j,temp;
    float avg;
    cout<<"Enter the size of Queue: ";
    cin>>n;
    cout<<"Enter the Queue: ";
    for(i=0;i<n;i++)
        cin>>queue[i];
    cout<<"Enter the initial head position: ";
    cin>>head;

    for(i=1;i<n;i++)
        t[i]=abs(head-queue[i]);
    for(i=0;i<n;i++){
        for(j=i+1;j<n;j++){
            if(t[i]>t[j]){
                temp=t[i];
                t[i]=t[j];
                t[j]=temp;
                temp=queue[i];
                queue[i]=queue[j];
                queue[j]=temp;
            }
        }
    }
    for(i=1;i<n-1;i++){
        seek=seek+abs(head-queue[i]);
        head=queue[i];
    }
    cout<<"Total Seek Time is "<<seek<<endl;
    avg=seek/(float)n;
    cout<<"Average Seek Time is "<<avg;
    return 0;
}
```

e) Output:



```
C:\Users\burha\OneDrive\Documents\operating system\sstf\sstf.exe
Enter the size of Queue: 7
Enter the Queue: 41
14
33
43
55
53
76
Enter the initial head position: 53
Total Seek Time is 129
Average Seek Time is 18.4286
```

f) Result:

Total Head Movement Required Serving All Requests

Experiment-9

SRTF Scheduling

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c program to implement SRTF scheduling.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of SRTF Scheduling:

- Associate with each process the length of its next CPU burst. Use these lengths to schedule the process with the shortest time.
- Two schemes:
 1. non pre-emptive – once CPU given to the process it cannot be preempted until completes its CPU burst.
 2. Preemptive – if a new process arrives with CPU burst length less than remaining time of current executing process, preempt. This scheme is known as the Shortest-Remaining-Time-First (SRTF).

Example of Preemptive SJF

Process	Arrival Time	Burst Time
P1	0.0	7
P2	2.0	4
P3	4.0	1
P4	5.0	4
SJF (preemptive)		

P1	P2	P3	P2	P4	P1
0	2	4	5	7	11
					16

c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: For each process in the ready Q, Accept Arrival time

Step 4: Start the Ready Q according the shortest Burst time by sorting according to lowest to Highest burst time.

Step 5: Set the waiting time of the first process in Sorted Q as '0'.

Step 6: After every unit of time compare the remaining time of currently executing process (RT) and Burst time of newly arrived process (BT_n).

Step 7: If the burst time of newly arrived process (BT_n) is less than the currently executing process (RT) the processor will preempt the currently executing process and starts executing newly arrived process

Step 7: Calculate

(c) Average waiting time = Total waiting Time / Number of process

(d) Average Turnaround time = Total Turnaround Time / Number of process

Step 8: Stop the process

d) Program:

```
#include<iostream>

using namespace std;
int main()
{
    int a[10],b[10],x[10];
    int waiting[10],turnaround[10],completion[10];
    int i,j,smallest,count=0,time,n;
    double avg=0,tt=0,end;

    cout<<"\nEnter the number of Processes: "; //input
    cin>>n;
    for(i=0; i<n; i++)
    {
        cout<<"\nEnter arrival time of process: "; //input
        cin>>a[i];
    }
    for(i=0; i<n; i++)
    {
        cout<<"\nEnter burst time of process: "; //input
        cin>>b[i];
    }
    for(i=0; i<n; i++)
        x[i]=b[i];

    b[9]=9999;
    for(time=0; count!=n; time++)
    {
        smallest=9;
        for(i=0; i<n; i++)
        {
            if(a[i]<=time && b[i]<b[smallest] && b[i]>0 )
                smallest=i;
        }
        b[smallest]--;

        if(b[smallest]==0)
        {
```

```

        count++;
        end=time+1;
        completion[smallest] = end;
        waiting[smallest] = end - a[smallest] - x[smallest];
        turnaround[smallest] = end - a[smallest];
    }
}
cout<<"Process"<<"\t"<< "burst-time"<<"\t"<<"arrival-time" <<"\t"<<"waiting-time"
    <<"\t"<<"turnaround-time"<< "\t"<<"completion-time"<<endl;
for(i=0; i<n; i++)
{

    cout<<"p"<<i+1<<"\t\t"<<x[i]<<"\t\t"<<a[i]<<"\t\t"<<waiting[i]<<"\t\t"<<turnaround[i]<<"
    \t\t"<<completion[i]<<endl;
    avg = avg + waiting[i];
    tt = tt + turnaround[i];
}
cout<<"\n\nAverage waiting time ="<<avg/n;
cout<<" Average Turnaround time ="<<tt/n<<endl;
}

```

e) Output:

```
C:\Users\burha\OneDrive\Documents\operating system\srtf\srtf.exe

Enter the number of Processes: 4
Enter arrival time of process: 3
Enter arrival time of process: 1
Enter arrival time of process: 4
Enter arrival time of process: 0
Enter burst time of process: 1
Enter burst time of process: 3
Enter burst time of process: 2
Enter burst time of process: 5
Process burst-time    arrival-time    waiting-time    turnaround-time    completion-time
p1                    1              3              0              1              4
p2                    3              1              1              4              5
p3                    2              4              1              3              7
p4                    5              0              6              11             11

Average waiting time =2  Average Turnaround time =4.75

-----
Process exited after 28.37 seconds with return value 0
Press any key to continue . . .
```

f) Result:

Average Waiting Time

Average Turnaround Time ...

Experiment-10

PRIORITY SCHEDULING

Name of Student: Burhanuddin Sahil		Class: CSIT-1
Enrollment No: 0827CI201055		Batch: 2020-24
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c program to implement Priority scheduling.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of Priority Scheduling:

A priority is associated with each process, and the CPU is allocated to the process with the highest priority. Equal-priority processes are scheduled in FCFS order.

An SJF algorithm is simply a priority algorithm where the priority (p) is the inverse of the (predicted) next CPU burst. The larger the CPU burst, the lower the priority, and vice versa.

As an example, consider the following set of processes, assumed to have arrived at time 0, in the order

	Burst		Waiting	Turnaround
Process	Time	Priority	Time	Time
P_1	10	3	6	16
P_2	1	1	0	1
P_3	2	4	16	18
P_4	1	5	18	19

P_5	5	2	1	6
Average	-	-	8.2	12



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Sort the ready queue according to the priority number.

Step 5: Set the waiting of the first process as '0' and its burst time as its turn around time

Step 6: For each process in the Ready Q calculate

(e) $\text{Waiting time for process}(n) = \text{waiting time of process } (n-1) + \text{Burst time of process}(n-1)$

(f) $\text{Turn around time for Process}(n) = \text{waiting time of Process}(n) + \text{Burst time for process}(n)$

Step 7: Calculate

(i) $\text{Average waiting time} = \text{Total waiting Time} / \text{Number of process}$

(j) $\text{Average Turnaround time} = \text{Total Turnaround Time} / \text{Number of process}$

Step 8: Stop the process

d) Program:

```
#include<iostream>
using namespace std;
int main() {
    int bt[20],p[20],wt[20],tat[20],pr[20],i,j,n,total=0,pos,temp,avg_wt,avg_tat;
    cout<<"Enter Total Number of Process:"<<endl;
    cin>>n;
    cout<<"Enter Burst Time and Priority"<<endl;
    for (i=0;i<n;i++) {
        cin>>bt[i];
        cin>>pr[i];
        p[i]=i+1;
    }

    for (i=0;i<n;i++) {
        pos=i;
        for (j=i+1;j<n;j++) {
            if(pr[j]<pr[pos])
                pos=j;
        }
        temp=pr[i];
        pr[i]=pr[pos];
        pr[pos]=temp;
        temp=bt[i];
        bt[i]=bt[pos];
        bt[pos]=temp;
        temp=p[i];
        p[i]=p[pos];
        p[pos]=temp;
    }
    wt[0]=0;

    for (i=1;i<n;i++) {
        wt[i]=0;
        for (j=0;j<i;j++)
            wt[i]+=bt[j];
        total+=wt[i];
    }
    avg_wt=total/n;

    total=0;

    for (i=0;i<n;i++) {
        tat[i]=bt[i]+wt[i];
```



```
        total+=tat[i];  
    }  
    avg_tat=total/n;  
  
    cout<<"Average Waiting Time ="<<avg_wt<<endl;  
    cout<<"Average Turnaround Time ="<<avg_tat;  
    return 0;  
}
```

e) Output:

```
C:\Users\burha\OneDrive\Documents\operating system\priority\priority.exe
Enter Total Number of Process:
4
Enter Burst Time and Priority
3 4
7 1
6 2
1 3
Average Waiting Time =8
Average Turnaround Time =12
-----
Process exited after 18.78 seconds with return value 0
Press any key to continue . . .
```

f) Result:

Average Waiting Time

Average Turnaround Time