Chap 1: Solving Nonlinear Eq.

- Interval Halving (Bisection)
- Linear Interpolation Methods
- Newton's Method
- Muller's Method
- Fixed-Point Iteration
- Order of convergence
- Multiple Roots
- Nonlinear systems

Introduction -1

- Problem: Solve f(x)=0
 - A system of nonlinear equations

$$f_1(x_1, x_2,, x_n) = 0$$

 \vdots
 $f_n(x_1, x_2,, x_n) = 0$

- Nonlinear function can be
 - Polynomials

$$f(x, y) = \sum_{i+j=n} a_{ij} x^i y^j$$

- Functions involve transcendental functions
 - sin, cos, exponentials

Introduction -2

- Closed form solutions
 - Only available for polynomials of degree less than 5.
 - Quadratic formula for degree 2
 - Complicated for degrees 3 and 4
 - For nonlinear polynomial system
 - Algebraic methods such as Grobner basis can be applied to eliminate variables and reduce the system to a triangular form

Introduction -3

- Numerical solutions
 - Via iteration procedure
 - Starting point
 - Compute iterates
 - Check for termination
 - Need to consider
 - Convergence
 - Rate of convergence
 - Stability
 - Early error magnified or not?

Bisection method 1

- An ancient but effective method for solving f(x)=0
- Starting with [x₁, x₂]: an interval that bracket a root, with f(x₁)*f(x₂) < 0
 - f changes signs at x₁ and x₂
 - If f is continuous, there must be at least one root between x₁ and x₂
- Divides the interval in half, finds in which half the root must lie, and repeat

Bisection method -2

Pseudo code

Repeat

$$x_3 = (x_1 + x_2)/2$$

if $f(x_3) f(x_1) < 0$
 $x_2 = x_3$
else $x_1 = x_3$
Until $(|x_1 - x_2|/2 < \text{tolerance})$

- The final value of x₃ approximates the root,
 with error no more than |x₁-x₂|/2
- The method may produce a false root if f(x) is discontinuous on [x₁,x₂]

Bisection Method -3

Advantages

- It is guaranteed to work if f(x) is continuous in the initial interval, and if the interval actually brackets a root
- The number of iterations to achieve a specified accuracy is known in advance, since the interval [a, b] is halved each time, so

error after n iterations
$$< \left| \frac{b-a}{2^n} \right|$$

$$\left| \frac{b-a}{2^n} \right|$$
 < Tolerance \Rightarrow a bound on n

Disadvantages

- Slow to converge
 - Why? No information about f(x) is used

Bisection Method 4

An example, a=0, b=1, tolerance=1E-4

$$f(x) = 3x + \sin(x) - e^x = 0$$

Table 1.1 The bisection method for $f(x) = 3x + \sin(x) - e^x = 0$, starting from $x_1 = 0$, $x_2 = 1$, using a tolerance value of 1E-4

Iteration	$X_{\mathbf{I}}$	X_2	X_3	$F(X_3)$	Maximum error	Actual error
1	0.00000	1.00000	0,50000	0.33070	0,50000	0.13958
	0.00000	0,50000	0.25000	-0.28662	0.25000	-0.11042
3	0.25000	0.50000	0.37500	0.03628	0.12500	0.01458
	0.25000	0.37500	0.31250	-0.12190	0.06250	-0.04792
4 5 6	0.31250	0.37500	0.34375	-0.04196	0.03125	-0.01667
6	0.34375	0.37500	0.35938	-0.00262	0.01563	-0.00105
7	0.35938	0.37500	0.36719	0.01689	0.00781	0.00677
8	0.35938	0.36719	0.36328	0.00715	0.00391	0.00286
8 9	0.35938	0.36328	0.36133	0.00227	0.00195	0.00091
10	0.35938	0.36133	0.36035	-0.00018	0.00098	-0.00007
11	0.36035	0.36133	0.36084	0.00105	0.00049	0.00042
12	0.36035	0.36084	0.36060	0.00044	0.00024	0.00017
13	0.36035	0.36060	0.36047	0.00013	0.00012	0.00005

Bisection Method .5

Some comments

- Recommended: used for finding an approximate value for the root, and the value is refined by more efficient methods
 - Most other root-finding methods require a starting value near to a root – lacking this, they may fail completely
 - Good practice:
 - Graph the function first
 - Search for interval that with sign change at ends
 - Apply Bisection to get an initial starting value
 - Apply better methods
- Not applicable to the case of multiple roots
 - Find the root by working on with f'(x), which will be zero at a multiple root.

Linear Interpolation Methods

- Approximate the function by a straight line
 - Interpolated line
 - The secant method
 - Two x-values nearest to the root
 - False position method
 - Similar to the Bisection method
 - Two points need to bracket the root
 - Tangent line
 - Newton method

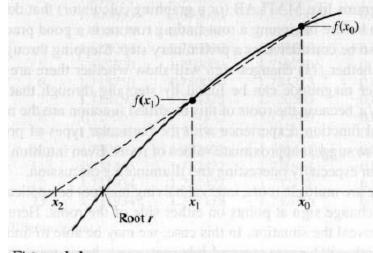


Figure 1.1

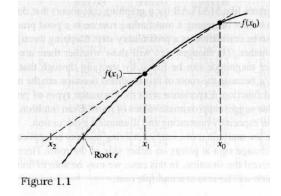
Secant method 1

Secant method

- Find two points on the curve near to the root
 - Draw a graph or apply a few iteration of bisection.
 - Two point may both be on one side of the root, or on opposite sides
- Find the line through these two points and find the point it intersects the x-axis

Repeat the process until the intersection point

is close enough to the root.



Secant method -2

From the similar triangles, we have

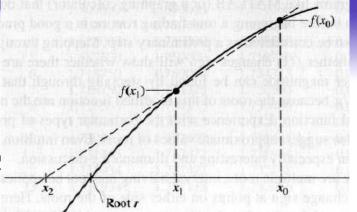
$$\frac{(x_1 - x_2)}{f(x_1)} = \frac{(x_0 - x_1)}{f(x_0) - f(x_1)}$$

Solve for x_2 :

$$x_2 = x_1 - f(x_1) \frac{(x_0 - x_1)}{f(x_0) - f(x_1)}$$

Repeat the iteration, we have

$$x_{n+1} = x_n - f(x_n) \frac{(x_{n-1} - x_n)}{f(x_{n-1}) - f(x_n)}$$



Feach newly computed value should be nearer to the root.

After second iteration:

Always using the last two computed values.

After first iteration:

There aren't two last computed values.

Swap x0 and x1 if Necessary, such that x1 is closer to the root.

The Secant Method -3

Start with x₀ and x₁ near the root

if $|f(x_0)| < |f(x_1)|$ then swap x_0 and x_1

Repeat

$$x_2 = x_1 - \frac{f(x_1)(x_0 - x_1)}{f(x_0) - f(x_1)}$$

$$x_0 = x_1$$

$$x_1 = x_2$$

Until $|f(x_2)|$ < tolerance value

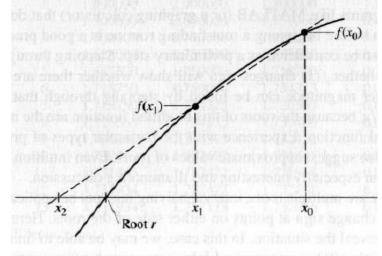


Figure 1.1

Note:

If f(x) is not continuous, the method may fail.

The Secant Method 4 An Example

Table 1.2 Secant method on $f(x) = 3x + \sin(x) - e^x$

Iteration	x_0	x_1	x_2	$f(x_2)$	
1	1	0	0.4709896	0.2651588	
2	0	0.4709896	0.3722771	2.953367E-02	
3	0.4709896	0.3722771	0.3599043	-1.294787E-03	
4	0.3722771	0.3599043	0.3604239	5.552969E-06	
5	0.3599043	0.3604239	0.3604217	3.554221E-08	

At x = .3604217, tolerance of .0000001 met!

 Fewer iterations are required compared to bisection: 5 iterations

The Secant Method ₋₅ Problems

 If the function is far from linear near the root, the successive iterates can fly off to points far from the root

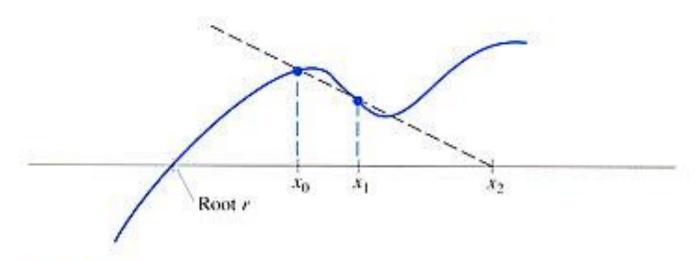


Figure 1.2
A pathological case for the secant method

False Position Method -1

- Avoid problems of secant method
 - Ensure that the root is bracketed between two starting values and remain between the successive pairs.
 - Similar to bisection method
 - x0 and x1 bracket a root

Differences:

- Next iterate is taken at the intersection of a line between the pair of x-values and the x-axis rather than the midpoint
- Gives faster convergence than does bisection, but at expense of a more complicated algorithm

False Position Method -2

 x_0 and x_1 bracket a root

Repeat

$$x_2 = x_1 - f(x_1) \frac{(x_0 - x_1)}{f(x_0) - f(x_1)}$$
if $f(x_2)$ opposite sign to $f(x_0)$

$$x_1 = x_2$$
else $x_0 = x_2$
Until $|f(x_2)| <$ tolerance value

Note:

if f(x) is not continuous, the method may fail.

A Comparison

Speed of convergence

- Secant (best), false position, then bisection

Table 1.3 Comparison of methods, $f(x) = 3x + \sin(x) - e^x = 0$, $x_0 = 0$, $x_1 = 1$

Iteration	Interval halving		False position		Secant method	
	х	f(x)	X	f(x)	x	f(x)
1	0.5	0.330704	0.470990	0.265160	0.470990	0.265160
	0.25	-0.286621	0.372277	0.029533	0.372277	0.029533
2	0.375	0.036281	0.361598	$2.94 * 10^{-3}$	0.359904	$-1.29 * 10^{-3}$
4	0.3125	-0.121899	0.360538	$2.90 \circ 10^{-4}$	0.360424	$5.55 * 10^{-6}$
5	0.34375	-0.041956	0.360433	$2.93 \cdot 10^{-5}$	0.360422	$3.55 * 10^{-7}$
Error after 5 iterations	0.01667		$-1.17 * 10^{-5}$		<-1 * 10 ⁻⁷	

- Takes a single initial x₀ (not too far from a root), and the intersection of the tangent line and x-axis as the next
- Is the most widely used method

- More rapidly convergent than bisection, secant

and false position.

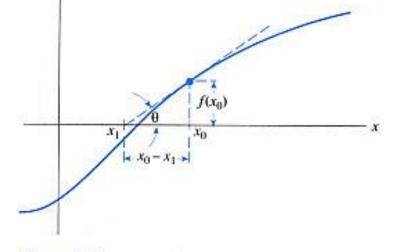


Figure 1.3

Iteration

$$\tan \theta = f'(x_0) = \frac{f(x_0)}{x_0 - x_1}$$

$$x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}$$

$$x_2 = x_1 - \frac{f(x_1)}{f'(x_1)}$$

•

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

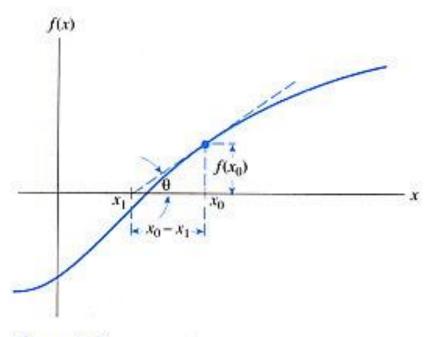


Figure 1.3

Given a x₀ reasonably close to the root.

Compute
$$f(x_0)$$
 and $f'(x_0)$
if $(f(x_0) \neq 0)$ and $(f'(x_0) \neq 0)$
Repeat

$$x_0 = x_0 - \frac{f(x_0)}{f'(x_0)}$$
Until $(|x_0 - x_1| < \text{tolerance1})$ or $(|f(x_0)| < \text{tolerance2})$

Note:

 $x_1 = x_0$

The method may converge to a root different from the expected one or diverge if the starting value is not close enough to the root.

Rate of convergence

- Quadratically convergent
 - Error of each step approaches a constant K times the square of the error of the previous step
 - The net result of this is that the number of decimal places of accuracy nearly doubles at each iteration
 - Fewer steps than previous methods
 - But, at the cost of two function evaluations at each step
 - » Previous methods need only one at each step (after the first step)
 - Faster than any of the methods discussed so far

Newton's Method -5 - Example

$$f(x) = 3x + \sin x - e^x = 0$$

Starting with $x_0 = 0.0$

$$x_1 = 0.33333$$

$$x_2 = 0.36017$$

$$x_3 = 0.3604217$$

True
$$x = 0.360421703$$

Note:

After 3 iterations, the solution is correct to 7 digits.

The error after an iteration is about 1/3 of the square of the previous error.

In some cases

- May converge to a different root, diverge, or oscillating
 - Converge
 - Wandering: jump around, and then jump to a point near to the root, and then converge rapidly
 - In endless loop, i.e., cycling
 - Starting point near to an inflection point or a turning point (local minimum/maximum)
 - Overshooting
 - Fly off to infinity when ever reach the minimum or maximum of the curve
 - X-intercept x1 is far from both x0 and the desired root, may converge to different root

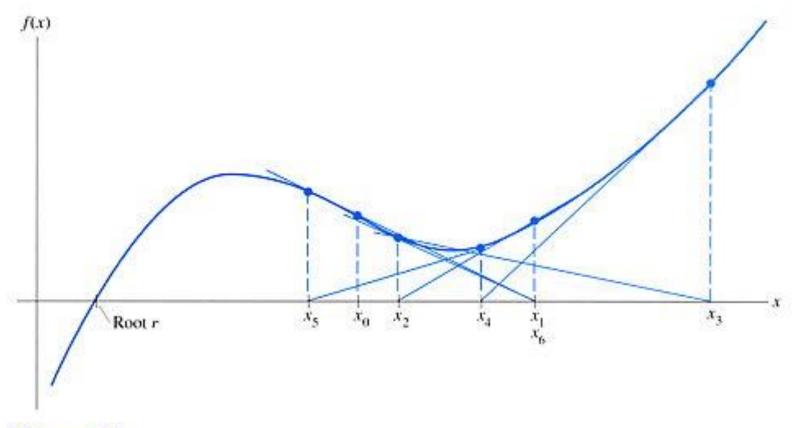


Figure 1.4

Starting point near to a turning point

Newton's Method 3 vs. interpolated methods

Interpolated methods:

$$x_{n+1} = x_n - f(x_n) \frac{(x_n - x_{n-1})}{f(x_n) - f(x_{n-1})}$$

$$= x_n - \underbrace{\frac{f(x_n)}{f(x_n) - f(x_{n-1})}}_{x_n - x_{n-1}}$$

Difference quotient approximates the derivative!! So closely resemble to Newton method!!

Newton's Method 5 Complex roots

 Newton's method works with complex roots if a complex starting value is given.

$$f(x) = x^3 + 2x^2 - x + 5$$

$$f(x) = 0$$
 has a real root at $x = -3$

It has two complex roots because x - axis is not crossed again.

Start with $x_0 = 1 + i$, we have

$$x_1 = 0.486238 + 1.04587i, x_2 = 0.448139 + 1.23665i$$

$$x_3 = 0.462720 + 1.22242i, x_4 = 0.462925 + 1.22253i$$

$$x_5 = 0.462925 + 1.22253i$$

Since agree to 6 significant digits, we have an estimate good to at least 6 significant digits.

Start with $x_0 = 1 - i$, we have the conjugate

$$0.462925 - 1.22253i$$

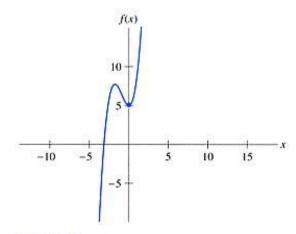


Figure 1.5 Plot of $f(x) = x^3 + 2x^2 - x + 5$

Newton's Method 10 for polynomials

- Polynomials have nice behavior:
 - They are everywhere continuous
 - They are smooth
 - Their derivatives are also continuous and smooth.
 - They are readily evaluated.
 - Number of roots can be predicted
 - Evaluation requires only +, -, *.

Newton's Method -11 for polynomials

- For previous methods, expect Newton method, there is nothing new (gain from the properties of polynomials)
- Polynomial evaluation of f(x) and f'(x) can be done by the use of synthetic division, which is based on the well known remainder theorem

Newton's Method -12 for polynomials

- Evaluating a polynomial by nested form
 - can be done by (or equivalent to)synthetic division
 - Evaluate f(x) at x=2

$$f(x) = 2x^3 + x^2 - 3x - 3$$
$$= ((2x+1)x - 3)x - 3$$

⇒ Process of synthetic division

Synthetic division for x = 2:

Note: Nested form evaluation:

$$((2x+1)x-3)x-3$$

with
$$x - 2: ((2 \times 2 + 1) \times 2 + 3) \times 2 - 3 \rightarrow 11$$

Re min der Theorem:

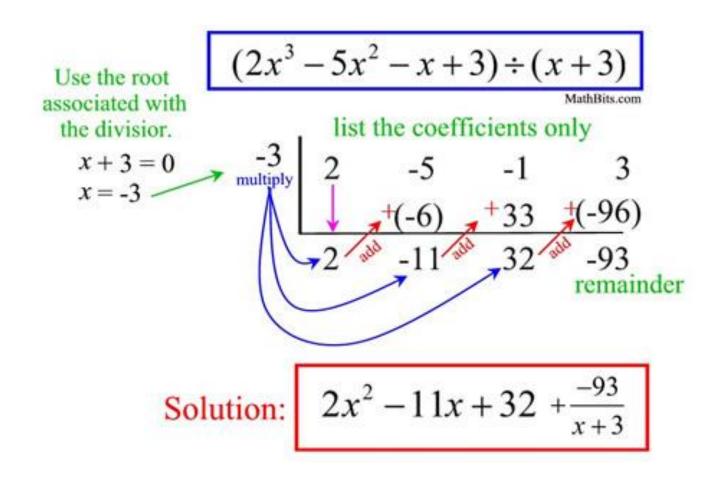
$$\frac{2x^3 + x^2 - 3x - 3}{x - 2} = 2x^2 + 5x + 7 + \frac{11}{x - 2}$$

That is,

$$2x^3 + x^2 - 3x - 3 = (x - 2)(2x^2 + 5x + 7) + 11$$

Newton's Method -13 for polynomials

Example on synthetic division



Newton's Method -14 for polynomials

 If the reduced (quotient) polynomial is divided by x-2 again, the remainder is the value of the derivative at x=2

$$f(x) = (x-2)(2x^2 + 5x + 7) + 11$$

$$x = 2$$

$$2 5 7$$

$$f'(x) = (2x^2 + 5x + 7) + (x-2)(4x + 5)$$

$$2 9 (25) \leftarrow f'(6)$$

$$f'(2)$$
 is the value of $2x^2 + 5x + 7$ at $x = 2$

$$f'(2) = 25$$

 $x_1 = 2 - \frac{11}{25} = 1.56$

Newton's Method -15 for polynomials

Synthetic division algorithm:

A way of obtaining $Q_{n-1}(x)$ and R.

$$P_n(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

$$P_n(x)$$
= $(x-a)Q_{n-1}(x) + R$
= $(x-a)(b_{n-1}x^{n-1} + \dots + b_1x + b_0) + R$

Multiplyin g out and equating codfficients of like terms in x, we have

$$\begin{aligned} & a_n = b_{n-1} \\ & P_n(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0 & a_{n-1} = b_{n-2} - a b_{n-1} \\ & a_{n-2} = b_{n-3} - a b_{n-2} \\ & \vdots \\ & P_n(x) \\ & = (x-a)Q_{n-1}(x) + R \end{aligned} \Rightarrow \begin{cases} b_{n-1} = a_n \\ b_{n-2} = a_{n-1} + a b_{n-1} \\ b_{n-3} = a_{n-2} + a b_{n-2} \\ \vdots \\ b_0 = a_1 + a b_1 \\ R = a_0 + a b_0 \end{aligned}$$

b_i and R are is the form of systhetic division

Newton's Method -16 for polynomials

 If the quotient polynomial is divided by x-a again, the remainder is the value of the derivative at x=a. WHY?

$$P_n(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

Dividing $P_n(x)$ by x-a, we have

$$\frac{P_n(x)}{x-a} = Q_{n-1}(x) + \frac{R}{x-a}$$

That is,

$$P_n(x) = (x-a)Q_{n-1}(x) + R$$

So
$$P_n(a) = R$$

Differentiate $P_n(x)$, we get

$$P'_{n}(x) = (x-a)Q'_{n-1}(x) + (1)Q_{n-1}(x)$$

Letting x = a, we have

$$P_n'(a) = Q_{n-1}(a)$$
 = remainder on dividing $Q_{n-1}(x)$ by $(x-a)$.

Comparisons Newton vs. secant

For simple root (i.e., f'(r) = 0)

- Rate of convergence
 - Newton method: quadratically, order: 2
 - Secant method: order: 1.62
 - Faster than linear convergence
 - Both converge linearly to multiple roots

Robustness

- Newton method: converge, wandering, overshooting, and cycling
- Secant method: converge, wandering, overshooting. But more robust than Newton

Muller's method -1

 A quadratic polynomial approximation is made to fit 3 points near a root.

A quadratic polynomial
$$p_2(v) = av^2 + bv + c$$

to fit $[x_1, f(x_1)], [x_0, f(x_0)], [x_2, f(x_2)]$

Using the quadratic rule to obtain the proper zero

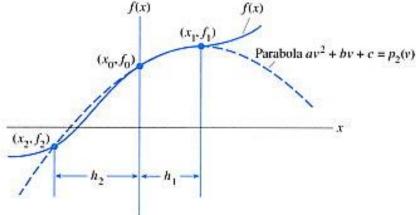


Figure 1.7

Simplify the development by transforming axes such that axes pass through the middle point.

So let $v = x - x_0$, and we try to fit the 3 points with $y = p_2(v) = av^2 + bv + c$.

Let $h_1 = x_1 - x_0$, $h_2 = x_0 - x_2$,

Evaluate the coefficients by evaluating p(v) at the three points :

$$v = 0$$
: $a(0)^2 + b(0) + c = f_0 \Rightarrow c = f_0$

$$v = h_1$$
: $a(h_1)^2 + b(h_1) + c = f_1$

$$v = -h_2$$
: $a(h_2)^2 + b(-h_2) + c = f_2$

Let $\gamma = h_2/h_1$, solving two linear equations for a and b:

$$a = \frac{\gamma f_1 - f_0(1+\gamma) + f_2}{\gamma h_1^2 (1+\gamma)}$$

$$b = \frac{f_1 - f_0 - ah_1^2}{h_1}$$

• Solve for the root of $p_2(v) = av^2 + bv + c = 0$ by the quadratic formula, choosing the root nearest to the middle point \mathbf{x}_0 by making the absolute value of denominate as large as possible in the following form:

root =
$$x_0 - v = x_0 - \frac{2c}{b \pm \sqrt{b^2 - 4ac}} \left(= x_0 + \frac{-b \pm \sqrt{b^2 - 4ac}}{2a} \right)$$

- WHY? See next slide
 - More stable compared to standard one
 - Choose the sign to give the largest absolute value of denominator: b>0: +, b<0: --, b=0: either

$$p_{2}(v) = av^{2} + bv + c$$

$$= a(x - x_{0})^{2} + b(x - x_{0}) + c = 0$$

$$v = x - x_{0} = \frac{-b \pm \sqrt{b^{2} - 4ac}}{2a}$$

$$x = x_0 + \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

with \pm chosen to minimize the modulus of the numerator (x is closest to x_0): b > 0 --> take +, b < 0 --> take -.

But minimizing the numerator may cancel significant digits when b^2 is much larger than 4ac (since $\sqrt{b^2 - 4ac} \sim b$). 39

Example:

$$p(v) = v^2 + 62.10v + 1 = 0$$

Using 4-digit rounding arithmetic

$$b > 0$$
, so take +,

$$\sqrt{b^2 - 4ac} = 62.06 \approx b = 62.10$$

and
$$x_1 = -0.01610723$$
,

with relative error: 2.4×10^{-1}

cancellation error → Large error!

Other root:

$$x_2 = -62.08390,$$

with relative error: 3.2×10^{-4}

Equivalent form:

$$\left(\frac{-b+\sqrt{b^2-4ac}}{2a}\right)\left(\frac{-b-\sqrt{b^2-4ac}}{-b-\sqrt{b^2-4ac}}\right)$$

$$=\frac{b^2-(b^2-4ac)}{2a(-b-\sqrt{b^2-4ac})} = \frac{-2c}{b+\sqrt{b^2-4ac}}$$

$$\left(\frac{-b - \sqrt{b^2 - 4ac}}{2a}\right)\left(\frac{-b + \sqrt{b^2 - 4ac}}{-b + \sqrt{b^2 - 4ac}}\right)$$

$$= \frac{-2c}{b - \sqrt{b^2 - 4ac}}$$

To guard against this, root is calculated in the equivalent form:

$$x = x_0 - \frac{2c}{b \pm \sqrt{b^2 - 4ac}}$$
with \pm chosen to maximize
the modulus of the
denominator:
$$b > 0 --> \text{take} +$$

$$b < 0 --> \text{take} -$$

There will be no case of substracting two nearly equal numbers.

Using the equivalent form:

$$x = x_0 - \frac{2c}{b \pm \sqrt{b^2 - 4ac}}$$

$$b > 0$$
, take +

$$x_1 = -0.01610,$$

relative error: 6.2×10^{-4}

No cancellation error → Smaller error!

Other root:

$$x_2 = -50.00,$$

relative error: 1.9×10^{-1}

- Start from three initial point x1, x0, x2
- Find a quadratic form passing through three point and find the roots
- Choose next 3 points that are most closely spaced
 - The root is to the right of x_0 : x_0 , x_1 , root
 - The root is to the left of $x_0 : x_0, x_2$, root

Procedure:

Given points $x_2 < x_0 < x_1$,

- Evaluate the functions f_2, f_0, f_1 .
- Find the coefficients of $p_2(v)$.
- Compute the roots:

$$x_r = x_0 - \frac{2c}{b \pm \sqrt{b^2 - 4ac}}$$

Choose the root closest to x_0

by making the denominator as largeas possible, and label it x_r

If
$$x_r > x_0$$

then rearrange to x_0, x_1, x_r

else rearrange to x_0, x_2, x_r

Until $|f(x_r)| < \text{Ftol}$

- Converge rate: similar to that for Newton method (actually order 1.85)
- No derivative evaluation and only one function evaluation per iteration (after we have obtained the next point)
- Will find a complex root if complex starting value is given
- May fail under some conditions
 - What will make the denominator of root's expression zero or nearly zero

EXAMPLE 1.2 Find a root between 0 and 1 of the same transcendental function as before: $f(x) = 3x + \sin(x) - e^x$. Let

$$x_0 = 0.5, \quad f(x_0) = 0.330704$$
 $h_1 = 0.5,$
 $x_1 = 1.0, \quad f(x_1) = 1.123189$ $h_2 = 0.5,$
 $x_2 = 0.0, \quad f(x_2) = -1$ $\gamma = 1.0.$

Then

$$a = \frac{(1.0)(1.123189) - 0.330704(2.0) + (-1)}{1.0(0.5)^{2}(2.0)} = -1.07644,$$

$$b = \frac{1.123189 - 0.330704 - (-1.07644)(0.5)^{2}}{0.5} = 2.12319,$$

$$c = 0.330704,$$

and

root =
$$0.5 - \frac{2(0.330704)}{2.12319 + \sqrt{(2.12319)^2 - 4(-1.07644)(0.330704)}}$$

= 0.354914 .

For the next iteration, we have

$$x_0 = 0.354914$$
, $f(x_0) = -0.0138066$ $h_1 = 0.145086$, $x_1 = 0.5$, $f(x_1) = 0.330704$ $h_2 = 0.354914$, $x_2 = 0$, $f(x_2) = -1$ $\gamma = 2.44623$.

Then

$$a = \frac{(2.44623)(0.330704) - (-0.0138066)(3.44623) + (-1)}{2.44623(0.145086)^2(3.44623)} = -0.808314,$$

$$b = \frac{0.330704 - (-0.0138066) - (-0.808314)(0.145086)^2}{0.145086} = 2.49180,$$

$$c = -0.0138066,$$

$$root = 0.354914 - \frac{2(-0.0138066)}{2.49180 + \sqrt{(2.49180)^2 - 4(-0.808314)(-0.0138066)}}$$

$$= 0.360465.$$

After a third iteration, we get 0.3604217 as the value for the root, which is identical to that from Newton's method after three iterations.

- A useful way for root finding
- Basis for some important theory
- **Fixed point** $f(x) = 0 \Rightarrow x = g(x)$ (in several forms) r: a fixed point of g if r = g(r)
- Fixed-point iteration $x_{n+1} = g(x_n), n = 0, 1, 2, 3,...$
- Under suitable conditions, the fixed-point iteration converges to the fixed point r, a root of f(x)=0
- Different rearranges will converge at different rate, or converge to different root, or diverge

$$f(x) = x^2 - 2x - 3 = 0$$
 roots: -1, 3

1.
$$x = g_1(x) = \sqrt{2x+3}$$

Start with $x_0 = 4$:
 $x_1 = 3.31662, x_2 = 3.10375$
 $x_3 = 3.03439, x_4 = 3.01144$
 $x_5 = 3.00381 \Rightarrow \text{converges to } x = 3$

2.
$$x = g_2(x) = \frac{3}{x-2}$$

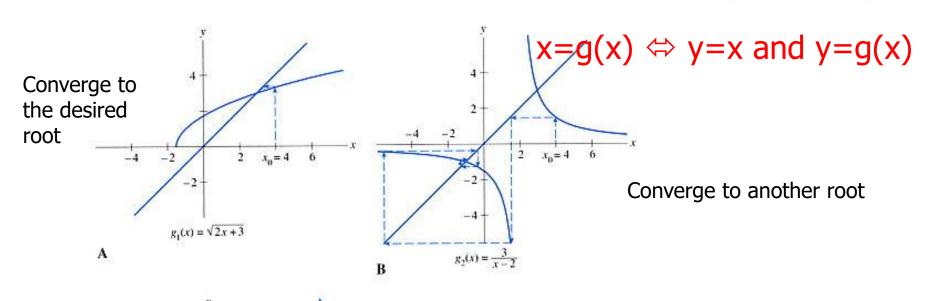
 $x_1 = 1.5, x_2 = -6, x_3 = -0.375,$
 $x_4 = -1.263158, x_5 = -0.919355$
 $x_6 = -1.02762, x_7 = -0.990876$
 $x_8 = -1.00305 \implies \text{converges to } x = -1$

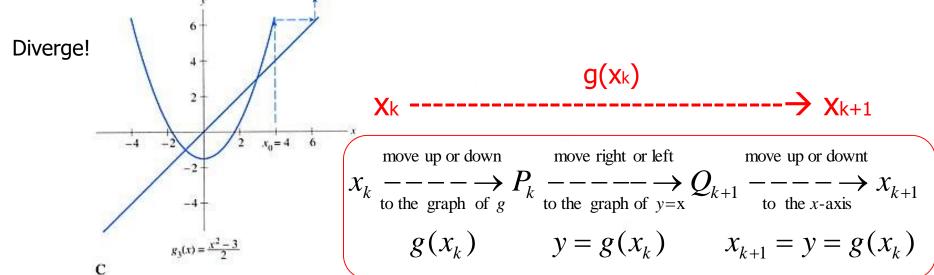
3.
$$x = g_3(x) = \frac{x^2 - 3}{2}$$

 $x_0 = 4$,
 $x_1 = 6.5$
 $x_2 = 19.625$,
 $x_3 = 191.070 \Rightarrow \text{diverges}$

Different arrangements have different convergence behavior. Why?

Look into this problem by using the graph of two intersecting plots of y = x and y = g(x).





Fixed-Point Iteration ₋₄ Rate of Convergence

Assume $\{x_n\}_{n=0}^{\infty}$ converges to r.

Let $e_n = x_n - r$.

If $|e_n|$ approaches to $c|e_{n-1}|^k$ as n becames infinite, we say that the sequence converges to r with order of convergence k, or formally:

The sequence $\{x_n\}_{n=0}^{\infty}$ is called linearly converge to r if $e_n \to 0$ in such a way that

$$\lim_{n\to\infty}\frac{e_n}{e_{n-1}}=C_L, \text{ where } 0<\left|C_L\right|<1$$

The sequence $\{x_n\}_{n=0}^{\infty}$ is called superlinearly converge to r if $e_n \to 0$ in such a way that

$$\lim_{n\to\infty}\frac{e_n}{e_{n-1}}=0$$

The sequence $\{x_n\}_{n=0}^{\infty}$ is called quadratically converge to r if $e_n \to 0$ in such a way that

$$\lim_{n\to\infty} \frac{e_n}{e_{n-1}^2} = C_Q, \text{ where } C_Q \neq 0$$

• Error

en+1 and en

$$x_{n+1} = g(x_n)$$
 and r is the root (fixed point)

$$x_{n+1} - r = g(x_n) - r = \frac{g(x_n) - g(r)}{x_n - r}(x_n - r)$$

If g(x) and g'(x) are continuous on the interval from r to x_n , the mean value theorem implies that

$$x_{n+1} - r = g'(\tau_n)(x_n - r),$$

where τ_n lies between x_n and r, and hence

$$\left|\mathbf{e}_{\mathbf{n}+1}\right| = \left|g'(\tau_n)\right| \left|e_n\right|$$

Condition for convergence

Suppose |g'(x)| < K < 1 for all x in the interval of size h around the root.

If x_0 is chosen in this interval, fixed point iteration will converge because

$$|e_{n+1}| < K|e_n| < K^2|e_{n-1}| < \cdots < K^{n+1}|e_0|.$$

Note that all succeeding iterates will lie in this interval and will converge to r.

Fixed-Point Iteration 4 Convergence Rate

Order of convergence: linear

$$\begin{aligned} |e_{n+1}| &= |x_{n+1} - r| \\ &= |g'(\tau_n)(x_n - r)| \\ &= |g'(\tau_n)||e_n|, \end{aligned}$$

where τ_n lies between x_n and r.

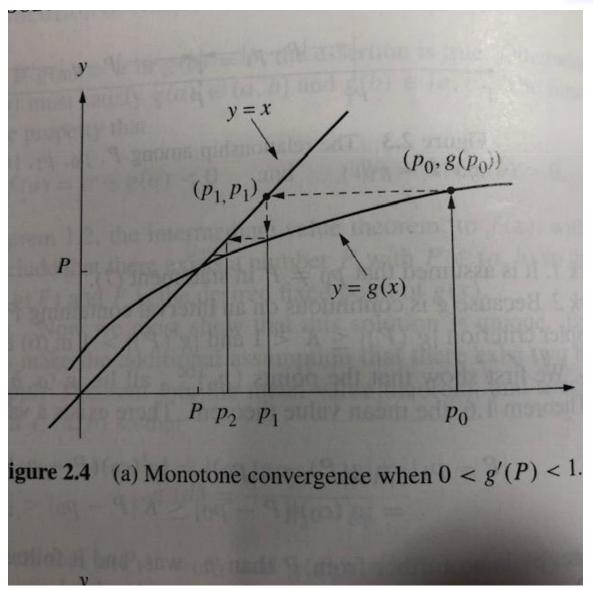
So if |g'(x)| < K < 1 for all x in the interval,

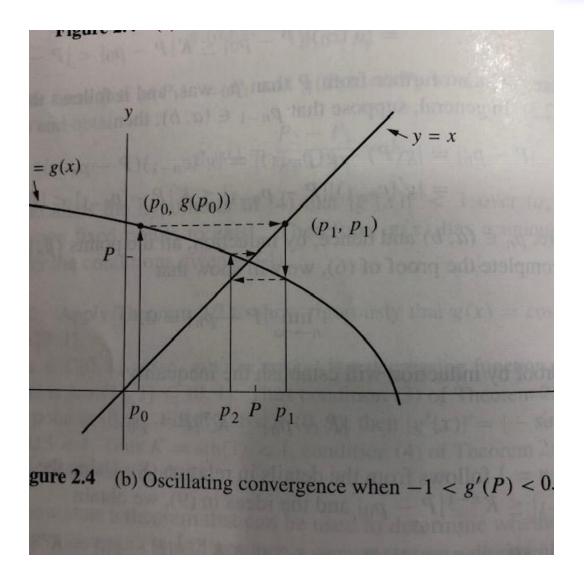
$$\lim_{n \to \infty} \frac{|e_{n+1}|}{|e_n|} = \lim_{n \to \infty} |g'(\tau_n)|, \text{ where } \tau_n \text{ lies between } x_n \text{ and } r.$$

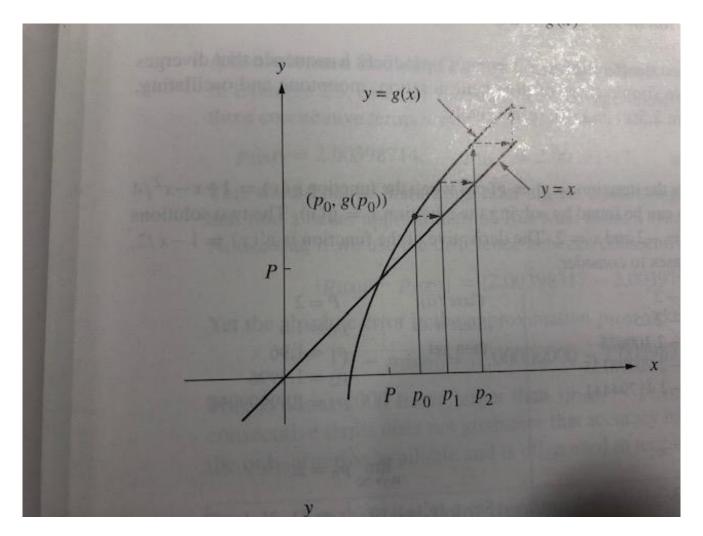
$$= |g'(r)| < 1 \text{ (since } x_n \text{ converges to } r, \tau_n \text{ converges to } r)$$

Converge linearly!

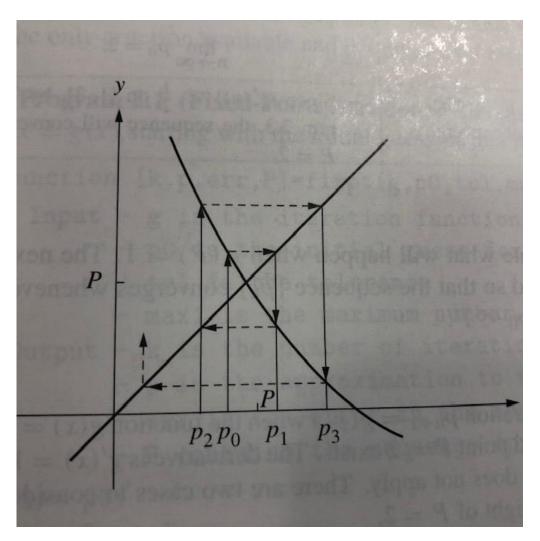
- Converge | g'(p) | < 1
 - Monotone converging
 - 0 < g'(p) < 1
 - Oscillating converging
 - -1 < g'(p) < 0
- Diverge | g'(p) | > 1
 - Monotone diverging
 - 1 < g'(p)
 - Oscillating diverging
 - g'(p) < -1







1 < g'(x): Monotone divergent



g'(x) < -1: Oscillating divergent

Convergence Condition for Newton Method

Rewrite the Newton method as a fixed-point iteration :

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} = g(x_n)$$

$$g'(x) = \frac{f(x)f''(x)}{\left[f'(x)\right]^2}$$

Based on the condition for convergence of fixed-point iteration:

Newton method converges if, on an interval about the root r,

$$\left|g'(x)\right| = \left|\frac{f(x)f''(x)}{\left[f'(x)\right]^2}\right| < 1$$

provided that x_0 is in the interval.

Convergence Order for Newton Method (for simple root)

By Taylor expansion of g(x) at r:

$$g(x) = g(r) + g'(r)(x - r) + \frac{g''(\xi)}{2}(x - r)^{2}$$

where ξ lies in the interval from x_n to r.

Since
$$f(r) = 0$$
, $g'(r) = \frac{f(r)f''(r)}{[f'(r)]^2} = 0$,

we have

$$g(x_n) = g(r) + \frac{g''(\xi_n)}{2}(x_n - r)^2$$

$$e_{n+1} = x_{n+1} - r = g(x_n) - g(r) = \frac{g''(\xi_n)}{2} e_n^2$$

As the iterates x_n approach to r when n approaches to infinity, so does ξ_n . That is,

$$\lim_{n\to\infty} \frac{\left|e_{n+1}\right|}{\left|e_{n}\right|^{2}} = \lim_{n\to\infty} \left|\frac{g''(\xi_{n})}{2}\right|$$

$$= \left| \frac{g''(r)}{2} \right| \neq 0$$

(for simple root)

So Newton method is quadratically convergent.

Convergence Condition of Bisection Method

[From Mathews and Fink, p 54]

Theorem

Assume that $f \in C[a,b]$ and that there exists a number $r \in [a,b]$ such that f(r) = 0. If f(a)f(b) < 0, the sequence $\{x_n\}_{n=0}^{\infty}$ generated by the bisection process, then

$$|x_n - r| \le \frac{b - a}{2^{n+1}}$$
 for $n = 0,1,2,...$

Therefore

$$\lim_{n\to\infty}x_n=r.$$

Convergence Condition of Bisection Method

Pf : Since the root r and midpoint x_n lie in [a, b],

$$\left|x_n - r\right| \le \frac{b_n - a_n}{2}$$
, for all n .

Observe that $b_1 - a_1 = \frac{b_0 - a_0}{2}$ $(a_0 = a, b_0 = b)$

$$b_2 - a_2 = \frac{b_1 - a_1}{2} = \frac{b_0 - a_0}{2^2}$$

By induction, we have

$$b_n - a_n = \frac{b_0 - a_0}{2^n}$$

So

$$|x_n - r| \le \frac{b_n - a_n}{2} \le \frac{b_0 - a_0}{2^{n+1}}$$

Convergence Condition of Bisection Method

Pf: Since the root r and midpoint x_n lie in [a, b],

$$|x_n - r| \le \frac{b_n - a_n}{2}$$
, for all n .

Basic step: Observe that $b_1 - a_1 = \frac{b_0 - a_0}{2}$ $(a_0 = a, b_0 = b)$

$$b_2 - a_2 = \frac{b_1 - a_1}{2} = \frac{b_0 - a_0}{2^2}$$

Induction step: Assuming it is true for n-1, we have

$$b_n - a_n = \frac{b_{n-1} - a_{n-1}}{2} = \frac{1}{2} \frac{b_0 - a_0}{2^{n-1}} = \frac{b_0 - a_0}{2^n}$$

So

$$|x_n - r| \le \frac{b_n - a_n}{2} \le \frac{b_0 - a_0}{2^{n+1}}$$

Proof by induction:

Basic step: prove it is true for n=1Induction step:

Assume it is true for n=kProve it is true for n=k+1

$$(a_0 = a, b_0 = b)$$

Convergence Order of Bisection Method

- One-half of the current interval is an upper bound to the error, which can serve as the estimate of the error
- So $|e_{n+1}|=0.5$ $|e_n|$ Linearly convergent!!

Since

$$e_{n+1} = |x_{n+1} - r| \le \frac{b_{n+1} - a_{n+1}}{2}, \quad e_{n+1} \approx \frac{b_{n+1} - a_{n+1}}{2}$$

$$e_n = |x_n - r| \le \frac{b_n - a_n}{2}, \quad e_n \approx \frac{b_n - a_n}{2}$$

$$\lim_{n \to \infty} \frac{e_{n+1}}{e_n} = \left(\frac{b_{n+1} - a_{n+1}}{2}\right) / \left(\frac{b_n - a_n}{2}\right) = \frac{1}{2}$$
(since $b_{n+1} - a_{n+1} = \frac{b_n - a_n}{2}$)

Convergence Order of Secant Method

Consider $x_{n+1} = g(x_n, x_{n-1})$.

Apply Taylor series and derive

$$\lim_{n=0}^{\infty} \frac{|e_{n+1}|}{|e_n e_{n-1}|} = \frac{1}{2} \left| \frac{f''(r)}{f'(r)} \right|$$

Further error analysis leads to

$$\lim_{n\to\infty}\frac{\left|e_{n+1}\right|}{\left|e_{n}\right|^{\alpha}}=K,$$

where

$$\alpha = \frac{1+\sqrt{5}}{2} = 1.62, \ K = \frac{1}{2} \left| \frac{f''(r)}{f'(r)} \right|^{1/\alpha}$$

So convergence order is 1.618.

Convergence Order for All Methods: summary

- Fixed-point iteration: Order: linear
- Bisection: Order: linear
 - One-half of the current interval is an upper bound to the error, which can serve as the estimate of the error
 - So $|e_{n+1}| = 0.5 |e_n|$
- False Position: Order: linear
 - $-X_{n+1}=g(x_{n}, x_{n-1})$
 - The root is always bracketed, so x_{n-1} can be thought as a constant
 - It is exactly as the fixed-point iteration, so it is linearly convergent

Convergence order for All Methods: summary

- Secant method: order: 1.62
 - $-x_{n+1}=g(x_{n}, x_{n-1})$
 - Similar analysis as the Newton method results in

$$e_{n+1} = \frac{g''(\xi_1)g''(\xi_2)}{2}e_n e_{n-1}$$

- Better than linear convergence but poorer than quadratic convergence. It has been shown the order is 1.62.
- Muller method: order=1.85
- Newton method: order= 2 (quadratic)

Starting Point Issues

- Newton method requires the initial point to be close to a root. General cases:
 - Converge, or scattering and then converge (wandering)
 - Oscillating
 - Diverge
- Newton method for polynomials (all roots)
 - Obtain an approximate for the first root
 - Proceed to obtain additional roots from the reduced polynomial, in which synthetic division can be employed to improve an initial estimate.
 - The process is repeated until the reduced

Multiple roots

- Examples
 - $f(x)=(x+1)^3$: triple root at x=-1
 - $f(x)=(x-2)^2$: double root at x=2

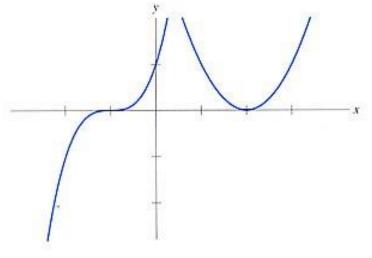


Figure 1.9

Previous methods do not work well

- Newton/secant method converges only linearly to a double/triple root
- Bisection and false position methods fail to get a double root. (because no sign changes)
- Muller's method is fastest, then Newton method, and then Secant method

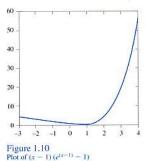


Table 1.7 Getting a double root, for $f(x) = (x-1)*(e^{(x-1)}-1)$					
posting to the post	Secant method	Newton's method	Muller's method		
Estimate after 9 iterations Start value(s)	1.00331 1.2, 1.5	1.00126 2.0	1.00058 0, 1.2, 1.5		

Double root at 1

Problems for Newton method

- Slow convergence: Converge linearly at double/triple roots with ratio (k-1)/k
 - Ratio of errors=1/2 for double roots
 - Ratio of errors=2/3 for triple roots
- Another disadvantage: Imprecision
 - Curve is flat near the root
 - f'(x) will always be 0 at a root

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

- There is a "neighborhood of uncertainty" around the root, where f(x) are very small around the root
 - » Computers will find f(x) equal to 0 throughout the neighborhood of the root
 - » Program cannot distinguish which value is really the root
 - » Using double precision helps to decrease the neighborhood of uncertainty

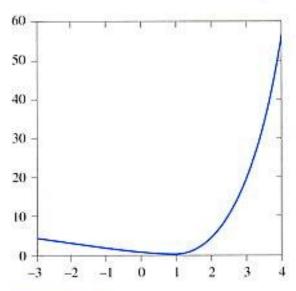


Figure 1.10 Plot of $(x - 1) (e^{(x-1)} - 1)$

Table 1.4 Errors when finding a double root

Iteration	Error	Ratio
I	0.3679	
2	0.1666	0.453
3	0.0798	0.479
4	0.0391	0.490
5	0.0193	0.494
6	0.0096	0.497
7	0.0048	0.500
8	8 0.0024	

Converge linearly to a double root with ratio=1/2

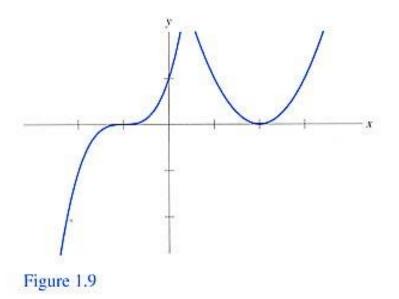


Table 1.10 Successive errors with Newton's method, for $f(x) = (x + 1)^3 = 0$

Iteration	Error	Iteration	Error
0	0.5	6	0.0439
1	0.3333	7	0.0293
2	0.2222	8	0.0195
3	0.1482	9	0.0130
4	0.0988	10	0.00867
5	0.0658		

Converge linearly to a triple root with ratio=2/3

Multiple Roots -6

- Linear convergence of Newton method to a double root and triple root
- For double roots
 - Bisection/false position methods fail for double roots
- For triple roots
 - All methods work, but even slower

Multiple Roots -7

Why linearly converge with error ratio (k-1)/k?

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} = g(x_n)$$

Taylor series of g(x) about r:

$$g(x) = g(r) + g'(r)(x-r) + (g''(\xi)/2)(x-r)^2$$

For simple roots, g'(r) = 0 leads to quadratic convergence.

For a root r of multiplicity k,

we cannot say
$$g'(r) = \frac{f(r)f''(r)}{(f'(r))^2} = 0$$
. Why? 0/0

We can factor out $(x-r)^k$ from f(x) to get

$$f(x) = (x-r)^k Q(x)$$
, where $Q(r) \neq 0$,

even though
$$f'(r) = f''(r) = \dots = f^{(k-1)}(r) = 0$$
.

Multiple Roots -8

Consider
$$g'(x) = \frac{f(x)f''(x)}{[f'(x)]^2}$$
 and let $h = x - r$

$$g'(x)$$

$$= \frac{h^k Q(x) [h^k Q''(x) + 2kh^{k-1} Q'(x) + k(k-1)h^{k-2} Q(x)]}{[h^k Q'(x) + kh^{k-1} Q(x)]^2}$$

$$=\frac{h^{2k-2}Q(x)\Big[h^2Q''(x)+2khQ'(x)+k(k-1)Q(x)\Big]}{h^{2k-2}\left[hQ'(x)+kQ(x)\right]^2}$$

When x = r, both denominator and numerator are 0, we cannot say g'(r) = 0.

Actually dividing denominator and numerator by h^{2k-2} , we have

$$\lim_{h \to 0} g'(x) = g'(r)$$

$$= \frac{Q(r) [0 \cdot Q''(r) + 0 \cdot Q'(r) + k(k-1)Q(r)]}{[0 \cdot Q'(r) + kQ(r)]^2}$$

$$= \frac{k(k-1)Q(r)^2}{k^2 Q(r)^2} = \frac{k-1}{k} \neq 0$$

$$g(x_n) = g(r) + g'(\xi_n)(x_n - r)$$

$$e_{n+1} = x_{n+1} - r = g(x_n) - g(r) = g'(\xi_n)e_n$$

$$\lim_{n \to \infty} \frac{|e_{n+1}|}{|e_n|} = \lim_{n \to \infty} g'(\xi_n) = g'(r) = \frac{k-1}{k}.$$

So the convergence is linear with

$$\lim_{n\to\infty}\frac{\left|e_{n+1}\right|}{\left|e_{n}\right|}=\frac{k-1}{k}.$$

Remedies for Multiple Roots with Newton Method 1

1. For a root of multiplicity *k*, restore quadratic convergence by modifying the formula to (if k is known)

$$x_{n+1} = x_n - k \frac{f(x_n)}{f'(x_n)} \equiv g_k(x_n)$$

Why?

At
$$f(r) = 0$$
, $g_{\nu}(r) = r$. Based on

$$f(x) = (x-r)^k Q(x)$$
, where $Q(r) \neq 0$,

we have

$$g'_k(x) =$$

$$\frac{(r-x)\{k(r-x)QQ''+Q'[2kQ-(k-1)(r-x)Q']\}}{[(r-x)Q'+kQ]^2}$$

We can easily see that $g'_k(r) = 0$, so the modified Newton method converge quadratically at multiple roots (including simple root where k = 1).

Remedies for Multiple Roots with Newton Method -2

With modified Newton method:

$$f(x) = (x-1)(e^{(x-1)}-1)$$

The third iterate is

$$x_3 = 1.00088$$
 with $f(x_3) = 0.00000$

Note: we also find that

$$e_{n+1} = 0.24e_n^2$$

confirming quadratic convergence.

Remedies for Multiple Roots with Newton Method -3

- But k is unknown in advance!!
- Guess a k by comparing a graph of f(x) with the plots of (x-r)k using an approximate r and various values of k.
 - The flatness will be similar, but this is not justified
- Divide f(x) by (x-r) and deflate the function, reducing the multiplicity by one.
 - But r is unknown
 - Dividing f by (x-s), where s is an approximate of r, does almost the same thing.

Remedies for Multiple Roots with Newton Method 4

2. If k is known

- k=2
 - Newton method can be applied to f'(x)=0 to converge quadratically to a double root of f(x) =0. (simple root of f'(x)=0)
- k=3
 - Similarly, the method can be applied to f''(x) = 0 to get quadratic convergence to a triple root. (simple root of f''(x)=0)

$$f(x) = (x-r)^{2}Q(x), \text{ where } Q(r) \neq 0,$$

$$f'(x) = 2(x-r)Q(x) + (x-r)^{2}Q'(x)$$

$$= (x-r)[2Q(x) + (x-r)Q'(x)]$$
where $2Q(r) + (r-r)Q'(x) = 2Q(r) \neq 0$

Remedies for Multiple Roots with Newton Method -5

- 3. If k is known [Acton 1970]
 A most tempting scheme:
 - When f(x) has a root r of multiplicity k, we have

$$f(x) = (x-r)^k Q(x) \text{ and } Q(r) \neq 0$$

$$f'(x) = k(x-r)^{k-1} Q(x) + (x-r)^k Q'(x)$$

$$= (x-r)^{k-1} [kQ(x) + (x-r)Q'(x)]$$

 If we divide f(x) by f'(x), we effectively deflate f(x) n-1 times and we now work with a new function that has only a single root at x=r.
 Why? See next page.

Remedies for Multiple Roots with Newton Method -6

If f(x) has a root of multiplicity k at x = r, $f(x) = (x - r)^k Q(x)$, where $Q(r) \neq 0$. Let

$$S(x) = \frac{f(x)}{f'(x)} = \frac{(x-r)^k Q(x)}{k(x-r)^{k-1} Q(x) + (x-r)^k Q'(x)}$$
$$= (x-r) \frac{Q(x)}{kQ(x) + (x-r)Q'(x)}$$

which has a simple root at x = r (since $Q(r) \neq 0$).

When S(x) used in Newton formula, we have

$$x_{n+1} = x_n - \frac{S(x_n)}{S'(x_n)}$$

$$= x_n - \frac{f(x_n)f'(x_n)}{[f'(x_n)]^2 - f(x_n)f''(x_n)} \equiv g(x_n)$$

 Used to speed up convergence of any sequence that is linearly convergent.

Definitions:

Given the sequence $\{p_n\}_{n=0}^{\infty}$, define the forward difference

$$\Delta p_n = p_{n+1} - p_n$$
, for $n \ge 0$.

Higher powers $\Delta^k p_n$ is defined recursively by

$$\Delta^k p_n = \Delta^{k-1}(\Delta p_n)$$
, for $n \ge 2$.

Theorem:

[From Friedman/Kandel, p90]
Assume that the sequence $\{p_n\}_{n=0}^{\infty}$ converge linearly to p and that $p_n - p \neq 0$ for all $n \geq 0$.

If there exists a real number

A with |A| < 1

such that

$$\lim_{n\to\infty}\frac{p_{n+1}-p}{p_n-p}=A,$$

then the sequence $\{q_n\}_{n=0}^{\infty}$ defined by

$$q_{n} = p_{n} - \frac{(\Delta p_{n})^{2}}{\Delta^{2} p_{n}}$$

$$= p_{n} - \frac{(p_{n+1} - p_{n})^{2}}{p_{n+2} - 2p_{n+1} + p_{n}}$$

converges to p faster than $\{p_n\}_{n=0}^{\infty}$ in the sense that

$$\lim_{n\to\infty}\left|\frac{q_n-p}{p_n-p}\right|=0.$$

Proof:

Since there exists a real number

A with |A| < 1 such that

$$\lim_{n\to\infty}\frac{p_{n+1}-p}{p_n-p}=A,$$

we have

$$\frac{p_{n+1} - p}{p_n - p} \approx A \approx \frac{p_{n+2} - p}{p_{n+1} - p}$$

when n is large.

This implies

$$(p_{n+1}-p)^2 \approx (p_n-p)(p_{n+2}-p)$$

Expanding the terms and cancelling p^2 yields

$$p \approx \frac{p_{n+2}p_n - p_{n+1}^2}{p_{n+2} - 2p_{n+1} + p_n}$$

$$= p_n - \frac{(p_{n+1} - p_n)^2}{p_{n+2} - 2p_{n+1} + p_n} = q_n.$$

We need to prove that $\{q_n\}_{n=0}^{\infty}$ converges to p faster than $\{p_n\}_{n=0}^{\infty}$ in the sense that

$$\lim_{n\to\infty}\left|\frac{q_n-p}{p_n-p}\right|=0.$$

Proof:(continued)

Let
$$e'_n = q_n - p$$
, then

$$e'_{n} = p_{n} - p - \frac{(p_{n+1} - p_{n})^{2}}{p_{n+2} - 2p_{n+1} + p_{n}}$$

$$= e_n - \frac{(e_{n+1} - e_n)^2}{e_{n+2} - 2e_{n+1} + e_n}$$

$$=\frac{e_n e_{n+2} - e_{n+1}^2}{e_{n+2} - 2e_{n+1} + e_n}$$

Note that

Based on
$$\lim_{n\to\infty} \frac{e_{n+1}}{e_n} = A$$
,

$$e_{n+1} = (A + \theta_n)e_n$$

$$e_{n+2} = (A + \theta_{n+1})e_{n+1}$$

= $(A + \theta_{n+1})(A + \theta_n)e_n$

where
$$\theta_n, \theta_{n+1} \to 0$$

as
$$n \to \infty$$
.

Proof:(continued)

Therefore

$$e'_n$$

$$= \frac{(A + \theta_{n+1})(A + \theta_n)e_n^2 - (A + \theta_n)^2 e_n^2}{(A + \theta_{n+1})(A + \theta_n)e_n - 2(A + \theta_n)e_n + e_n}$$

$$= \frac{(A + \theta_{n+1})(A + \theta_n) - (A + \theta_n)^2}{(A + \theta_{n+1})(A + \theta_n) - 2(A + \theta_n) + 1} e_n$$

$$= \frac{(A + \theta_n)[(A + \theta_{n+1}) - (A + \theta_n)]}{(A^2 - 2A + 1) + A\theta_{n+1} + A\theta_n + \theta_{n+1}\theta_n - 2\theta_n}$$

$$= \frac{(A+\theta_n)}{(A-1)^2 + A(\theta_{n+1}+\theta_n) + \theta_n(\theta_{n+1}-2)} (\theta_{n+1}-\theta_n) e_n$$

Hence

$$\frac{e'_{n}}{e_{n}} = \frac{(A + \theta_{n})}{(A - 1)^{2} + A(\theta_{n+1} + \theta_{n}) + \theta_{n}(\theta_{n+1} - 2)}(\theta_{n+1} - \theta_{n})$$

and

$$\lim_{n\to\infty} \left| \frac{e'_n}{e_n} \right| = \frac{A}{(A-1)^2} \lim_{n\to\infty} (\theta_{n+1} - \theta_n) = 0.$$

Advantages:

- No significant additional computing time, because the time for computing q_n is negligible, once p_n , p_{n+1} , p_{n+2} are already given.
- Aitken acceleration speeds up the convergence for any sequence that is linearly convergent.

Rate of convergence of Aitken method

– Quadratic? NO!! Still linear, but faster than that for original sequence. WHY??

Theorem:

Consider $p_{n+1} = g(p_n)$. If the sequence

 $\{p_n\}_{n=0}^{\infty}$ converge linearly to p and

$$\lim_{n\to\infty}\frac{e_{n+1}}{e_n}=A\qquad (=g'(p)),$$

for |A| < 1, then the Aitken sequence $\{q_n\}_{n=0}^{\infty}$

behaves asymptotically according to

$$\lim_{n \to \infty} \frac{e'_{n+1}}{e'_n} = A^2 < A$$

Proof: Since

$$e'_{n} = \frac{(A + \theta_{n+1})(A + \theta_{n})e_{n}^{2} - (A + \theta_{n})^{2}e_{n}^{2}}{(A + \theta_{n+1})(A + \theta_{n})e_{n} - 2(A + \theta_{n})e_{n} + e_{n}}$$

$$= \frac{(A + \theta_{n})}{(A - 1)^{2} + A(\theta_{n+1} + \theta_{n}) + \theta_{n}(\theta_{n+1} - 2)}(\theta_{n+1} - \theta_{n})e_{n}$$

Hence

$$\frac{e'_{n+1}}{e'_{n}} = \frac{(A + \theta_{n+1})[(A - 1)^{2} + A(\theta_{n} + \theta_{n+1}) + \theta_{n}(\theta_{n+1} - 2)]}{(A + \theta_{n})[(A - 1)^{2} + A(\theta_{n+1} + \theta_{n+2}) + \theta_{n+1}(\theta_{n+2} - 2)]} \frac{e_{n+1}(\theta_{n+2} - \theta_{n+1})}{e_{n}(\theta_{n+1} - \theta_{n})}$$

Since

$$\lim_{n\to\infty}\theta_n=0 \text{ and } \lim_{n\to\infty}\frac{e_{n+1}}{e_n}=A$$

We have

$$\lim_{n\to\infty}\frac{e'_{n+1}}{e'_n}=A\lim_{n\to\infty}\frac{(\theta_{n+2}-\theta_{n+1})}{(\theta_{n+1}-\theta_n)}.$$

We need to show that

$$\lim_{n\to\infty}\frac{(\theta_{n+2}-\theta_{n+1})}{(\theta_{n+1}-\theta_n)}=A.$$

Proof:

Consider

$$g(x) = g(r) + (x - r)g'(r) + \frac{(x - r)^2}{2}g''(\xi),$$

where ξ is between x and r.

$$e_{n+1} = x_{n+1} - r = g(x_n) - g(r)$$

$$= (x_n - r)g'(r) + \frac{(x_n - r)^2}{2}g''(\xi_n),$$

where ξ_n is between x_n and r, so $\lim_{n\to\infty} \xi_n = r$.

Due to the continuity of g", we can write

$$g''(\xi_n) = g''(r) + \theta'_n$$
, $\lim_{n \to \infty} \theta'_n = 0$.

Hence

$$e_{n+1} = e_n A + \frac{e_n^2}{2} (A' + \theta_n'),$$

where g'(r) = A, g''(r) = A'

Because $e_{n+1} = (A + \theta_n)e_n$, we obtain

$$e_{n+1} = e_n A + \frac{e_n^2}{2} (A' + \theta_n') = (A + \theta_n) e_n$$

and
$$\theta_n = \frac{e_n}{2} (A' + \theta_n')$$
.

Finally,

$$\lim_{n\to\infty} \frac{\theta_{n+1}}{\theta_n} = \lim_{n\to\infty} \frac{\frac{e_{n+1}}{2} \left(A' + \theta'_{n+1} \right)}{\frac{e_n}{2} \left(A' + \theta'_n \right)}$$

$$= \lim_{n \to \infty} \left(\frac{e_{n+1}}{e_n} \frac{A' + \theta'_{n+1}}{A' + \theta'_n} \right) = \lim_{n \to \infty} \frac{e_{n+1}}{e_n} = A.$$

Therefore

$$\lim_{n \to \infty} \frac{\theta_{n+2} - \theta_{n+1}}{\theta_{n+1} - \theta_n} = \lim_{n \to \infty} \frac{\frac{\theta_{n+2}}{\theta_{n+1}} - 1}{1 - \frac{\theta_n}{\theta_{n+1}}} = \frac{A - 1}{1 - \frac{1}{A}} = A$$

Theorem:

Given a tolerance $\varepsilon > 0$, let N, N' denote the number of iterations required for convergence using fixed point iteration and Aitken's iteration, respectively. Then

$$\lim_{\varepsilon \to 0} \frac{N'}{N} = \frac{1}{2}$$

N=11, N'=5

$$|e_4/e_3|=0.312 \sim |A|=g'(s)=(3/8) s^2=0.308$$

 $|e_4/e_3|=0.096 \sim A^2=0.095$

Example 3.3.4.

Let $f(x) = 1 - (x^3/8)$, $0 \le x \le 1$ (Fig. 3.3.1) and consider the equation x = f(x), given a first approximation $x_0 = 0$ and a tolerance $\epsilon = 10^{-6}$. The exact solution computed to nine significant digits is s = 0.906795303. A comparison between the performances of the SIM and ATKN is given in Table 3.3.3.

Table 3.3.3. $x = 1 - (x^3/8), \ 0 \le x \le 1, \ x_0 = 0, \ \epsilon = 10^{-6}$

n	x_n	x'_n	$ \epsilon_n/\epsilon_{n-1} $	$\epsilon'_n/\epsilon'_{n-1}$
0	0.0000000000	0.88888889		
1	1.000000000	0.906020558	0.103	0.043
2	0.875000000	0.906717286	0.341	0.101
3	0.916259766	0.906788044	0.298	0.093
4	0.903846331	0.906794608	0.312	0.096
:				
12	0.906796083		H-STATE OF	

SIM: Fixed point iteration

Example 3.3.1.

Consider $x = e^{-x}$, $1/2 \le x \le 2/3$, $x_0 = 0.5$. Table 3.3.1 contains the first eight iterations, using the SIM, compared with eight iterations based on Aitken's scheme.

Table 3.3.1. $x = e^{-x}$ (Aitken's Method)

n	X_R	x' _n	
0	0.500000	0.567624	
1	0.606531	0.567299	
2	0.545239	0.567193	
3	0.579703	0.567159	
4	0.560065	0.567148	
5	0.571172	0.567145	
6	0.564863	0.567144	
7	0.568438	0.567143	
8	0.566409		
22	0.567143		

We see that ATKN yields the first 6 digits of the exact solution after 8 iterations, compared with the 22 that are needed for the SIM (to provide the same accuracy).

Example 3.3.2

	ITERATIO	Exact solution to 8 decimal digits: 0.73908513
11- 2	$32. x = \cos x$	(Aitken's Method)
Table 5		X'n
n	χ_n	0.72801036
0	1.00000000	
The section of the section of	0.54030231	0.73366516
2	0.85755322	0.73690629
3	0.65428979	0.73805042
4	0.79348036	0.73863610
5	0.70136877	0.73887038
6	0.76395968	0.73899224
7	0.72210243	0.73904251
8	0.75041776	0.73906595
9	0.73140404	0.73907638
10	0.74423735	0.73908118
11	0.73560474	To approximate r within error 10^-5:
		Aitken method: requires 11 iterations
_26	0.73909441	Fixed-point iteration: requires 25 iterations
		The second secon

Nonlinear systems

Nonlinear system

$$f_1(x_1, x_2, x_3,..., x_n) = 0$$

$$f_2(x_1, x_2, x_3,..., x_n) = 0$$

$$\vdots$$

$$\vdots$$

$$f_n(x_1, x_2, x_3,..., x_n) = 0$$

2x2 system

$$f(x, y) = 0$$
$$g(x, y) = 0$$

Solution set: Regarded as the intersections of two algebraic curves.

Nonlinear systems

Example

$$\begin{cases} x^2 + y^2 = 4 \\ e^x + y = 0 \end{cases}$$

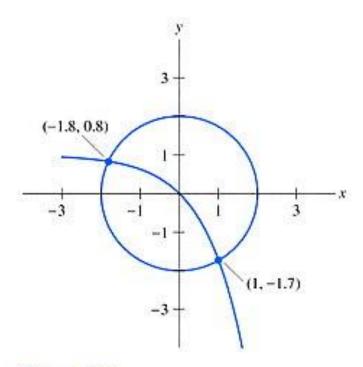


Figure 1.11

Nonlinear Systems Methods

Similar to solving f(x)=0

- Fixed point iteration
 - Different formulae behave differently
 - Converge linearly
- Newton method
 - Derivative vs. Jacobian matrix
 - Converge quadratically

Nonlinear Systems Fixed Point Iteration

Consider the 2x2 system:

$$\begin{cases} f_1(x, y) = 0 \\ f_2(x, y) = 0 \end{cases}$$

Fixed point equation:

$$\begin{cases} x = g_1(x, y) \\ y = g_2(x, y) \end{cases}$$

If the starting point is sufficiently close to the fixed point (p, q), and if

$$\left| \frac{\partial g_1}{\partial x}(p,q) \right| + \left| \frac{\partial g_1}{\partial y}(p,q) \right| < 1$$
 and

$$\left| \frac{\partial g_2}{\partial x}(p,q) \right| + \left| \frac{\partial g_2}{\partial y}(p,q) \right| < 1$$

then the iteration converges to the fixed point.

Note:

If the conditions are not met, the iteration might diverge.

Nonlinear Systems Fixed Point Iteration

$$\begin{cases} x^2 - 2x - y + 0.5 = 0 \\ x^2 + 4y^2 - 4 = 0 \end{cases}$$

Fixed point formulation:

$$\begin{cases} x = \frac{x^2 - y + 0.5}{2} \\ y = \frac{-x^2 - 4y^2 + 8y + 4}{8} \end{cases}$$

Starting with (0,1), it converges to (-0.2,1.0).

Observe that

$$\left| \frac{\partial g_1}{\partial x}(p,q) \right| + \left| \frac{\partial g_1}{\partial y}(p,q) \right|$$

$$= |x| + |-0.5| < 1, \text{ and}$$

$$\left| \frac{\partial g_2}{\partial x}(p,q) \right| + \left| \frac{\partial g_2}{\partial y}(p,q) \right|$$

$$<\frac{|-x|}{4}+|-y+1|<1$$

for
$$-0.5 < x < 0.5$$
 and $0.5 < y < 1.5$.

Starting with (2,0), it diverges away from the solution.

Nonlinear Systems Newton Method

Newton method for solving $\begin{cases} f(x, y) = 0 \\ g(x, y) = 0 \end{cases}$

Let (r, s) be a root and expand both functions as a Taylor series about (x_i, y_i) , where (x_i, y_i) is a point near the root:

$$\begin{cases} f(r,s) = 0 = f(x_i, y_i) + f_x(x_i, y_i)(r - x_i) \\ + f_y(x_i, y_i)(s - y_i) + \cdots \\ g(r,s) = 0 = g(x_i, y_i) + g_x(x_i, y_i)(r - x_i) \\ + g_y(x_i, y_i)(s - y_i) + \cdots \end{cases}$$

Truncating both series give

$$\begin{cases} 0 = f(x_i, y_i) + f_x(x_i, y_i)(r - x_i) + f_y(x_i, y_i)(s - y_i) \\ 0 = g(x_i, y_i) + g_x(x_i, y_i)(r - x_i) + g_y(x_i, y_i)(s - y_i) \end{cases}$$

Nonlinear systems Newton method

The system is rewritten as

$$\begin{cases} f_x(x_i, y_i)(r - x_i) + f_y(x_i, y_i)(s - y_i) = -f(x_i, y_i) \\ g_x(x_i, y_i)(r - x_i) + g_y(x_i, y_i)(s - y_i) = -g(x_i, y_i). \end{cases}$$

The iterate is derived by solving

$$\begin{cases} f_x(x_i, y_i) \Delta x_i + f_y(x_i, y_i) \Delta y_i = -f(x_i, y_i) \\ g_x(x_i, y_i) \Delta x_i + g_y(x_i, y_i) \Delta x_i = -g(x_i, y_i) \end{cases}$$

where $x_{i+1} = x_i + \Delta x_i$ and $y_{i+1} = y_i + \Delta y_i$ are improved estimate of the root.

Repeating this iteration until both $f(x_i)$ and $g(y_i)$ are close to zero.

Nonlinear Systems Newton Method

$$\begin{cases} f_x(x_i, y_i) \Delta x_i + f_y(x_i, y_i) \Delta y_i = -f(x_i, y_i) \\ g_x(x_i, y_i) \Delta x_i + g_y(x_i, y_i) \Delta x_i = -g(x_i, y_i) \end{cases}$$

can be rewritten as

$$J(x_i, y_i) \begin{bmatrix} \Delta x_i \\ \Delta y_i \end{bmatrix} = \begin{bmatrix} -f(x_i, y_i) \\ -g(x_i, y_i) \end{bmatrix}$$

If $J(x_i, y_i)$ is nonsingular, Δx_i and Δy_i can be found by solving the linear system.

New iterate is obtained by

$$\begin{bmatrix} x_{i+1} \\ y_{i+1} \end{bmatrix} = \begin{bmatrix} x_i \\ y_i \end{bmatrix} + \begin{bmatrix} \Delta x_i \\ \Delta y_i \end{bmatrix}.$$

Nonlinear Systems Newton Method

- The Newton method for solving nonlinear system converges quadratically, but requiring
 - n²+n function evaluations at each step
 - 6 function evaluations for a 2x2 system
 - 12 function evaluations for a 3x3 system
 - Solving an nxn linear system