

Report on

"Title of the project"

Submitted in partial fulfillment of the requirements for Sem VI

Compiler Design Laboratory

Bachelor of Technology in Computer Science & Engineering

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6.	 IMPLEMENTATION DETAILS (TOOL AND DATA STRUCTURES USED in order to implement the following): SYMBOL TABLE CREATION ABSTRACT SYNTAX TREE (internal representation) INTERMEDIATE CODE GENERATION CODE OPTIMIZATION ASSEMBLY CODE GENERATION ERROR HANDLING - strategies and solutions used in your Mini-Compiler implementation (in its scanner, parser, semantic analyzer, and code generator). Provide instructions on how to build and run your program. 	
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Introduction

A compiler is a computer program that transforms source code written in a programming language to another computer language. We wrote a mini-compiler which compiles C source code to machine code. The compiler itself is written in Java, using Jflex and Byaccj as tools for parsing the source code.

Sample Input: Sample Output: $oxedsymbol{ox{oxed}}}}}}$ ASSEMBLY CODE: 2 int main() MOV R0, 5 SW R0, a MOV RO, 6 SW RO, b int a=5; int b=6; .L1: while (a < 20)LW R1, a MOV R2, 20 CMP R1, R2 b=b+1;a=a+1;JGE L2 LW R1, b MOV R2, 1 ADD R0, R1, R2 SW R0, t2 return 0; LW R0, t2 SW R0, b LW R1, a MOV R2, 1 ADD R0, R1, R2 SW R0, t3 LW R0, t3 SW R0, a JMP L1

Architecture Of Language

This mini-compiler supports IF-ELSE, WHILE, DO-WHILE and FOR clauses. It has proper syntax and semantics handling, for example :

- It supports the basic data types: INT, FLOAT, VOID.
- It supports basic unary operators like: '~' '|' ' & ' ^'.
- It supports binary operators like: '+' '*' '-' '/' '%' '==' '<=' etc.
- It supports the assignment operator '='
- It has scope handling.
- It checks for undeclared variables.
- It checks for redeclaration of variables.
- Checks for '; 'at required places.
- Checks the type of the variable.
- Checks if the return type of a function matches with the datatype of variable/constant being returned.
- Checks if the function call matches its prototype.
- Checks if the index of an array is a positive integer or not.

Literature Survey

- For Grammar: https://www.lysator.liu.se/c/ANSI-C-grammar-y.html
- Compilers Principles, Techniques and Tools By Alfred V. Aho, Monica S. Lam, Ravi Sethi, Jeffrey D. Ullman
- JFlex: https://jflex.de/manual.html
- ByaccJ: http://byaccj.sourceforge.net/
- ICG: https://www.geeksforgeeks.org/intermediate-code-generation-in-c ompiler-design/
- Code Generation: https://web.cs.ucdavis.edu/~pandey/Teaching/ECS142/Lects/final.codegen.pdf

Context Free Grammar

```
"begin: external declaration",
"begin: begin external declaration",
"begin: Define begin",
"begin: error",
"primary_expression: IDENTIFIER",
"primary expression: CONSTANT",
"primary_expression: STRING_LITERAL",
"primary expression: '(' expression ')",
"Define: DEFINE",
"postfix expression: primary expression",
"postfix expression: postfix expression'['expression']",
"postfix expression: postfix expression '('')",
"postfix expression: postfix_expression '(' argument_expression_list ')",
"postfix expression: postfix expression'.' IDENTIFIER",
"postfix expression: postfix_expression PTR_OP IDENTIFIER",
"postfix expression: postfix expression INC OP",
"postfix expression: postfix expression DEC OP".
"argument_expression_list: assignment_expression",
"argument expression list: argument expression list',' assignment expression",
"unary expression: postfix expression",
"unary expression: INC OP unary expression",
"unary_expression : DEC_OP unary_expression",
"unary expression: unary operator cast expression",
"unary expression: SIZEOF unary expression",
"unary expression: SIZEOF'('type name')",
"unary operator: '&",
"unary operator: '*'",
"unary operator: '+'",
"unary_operator: '-'",
"unary_operator: '~'",
"unary_operator: '!'",
"cast expression: unary_expression",
"cast_expression: '(' type_name ')' cast_expression",
"multiplicative expression: cast expression",
"multiplicative_expression: multiplicative_expression '*' cast_expression",
"multiplicative expression: multiplicative expression', cast expression",
"multiplicative expression: multiplicative expression '%' cast expression",
"additive expression: multiplicative expression",
"additive expression: additive expression '+' multiplicative expression",
"additive expression: additive_expression'-' multiplicative_expression",
"shift expression: additive expression",
"shift_expression: shift_expression LEFT_OP additive_expression",
"shift expression: shift expression RIGHT OP additive expression",
```

```
"relational expression: shift expression",
"relational expression: relational expression'<' shift expression",
"relational expression: relational expression'>' shift expression",
"relational expression: relational expression LE OP shift expression",
"relational expression: relational expression GE OP shift expression",
"equality expression: relational expression",
"equality expression: equality expression EQ OP relational expression",
"equality expression: equality expression NE OP relational expression",
"and expression: equality expression",
"and expression: and expression'&' equality expression",
"exclusive or_expression: and_expression",
"exclusive or expression: exclusive or expression 'A' and expression",
"inclusive or expression: exclusive or expression",
"inclusive or expression: inclusive_or_expression" | exclusive_or_expression",
"logical and expression: inclusive_or_expression",
"logical_and_expression: logical_and_expression AND_OP inclusive_or_expression",
"logical or expression: logical and expression",
"logical_or_expression: logical_or_expression OR_OP logical_and_expression",
"conditional expression: logical or expression",
"conditional_expression: logical_or_expression'?' expression':' conditional_expression",
"assignment expression: conditional expression",
"assignment expression: unary expression assignment operator assignment expression",
"assignment operator: '='",
"assignment operator: MUL ASSIGN",
"assignment_operator: DIV_ASSIGN",
"assignment operator: MOD ASSIGN",
"assignment_operator : ADD_ASSIGN",
"assignment operator: SUB ASSIGN",
"assignment operator: LEFT ASSIGN".
"assignment operator: RIGHT ASSIGN",
"assignment_operator: AND_ASSIGN",
"assignment_operator: XOR_ASSIGN",
"assignment_operator: OR_ASSIGN",
"expression: assignment expression",
"expression: expression', assignment expression",
"constant expression: conditional expression",
"declaration : declaration specifiers ';"",
"declaration : declaration_specifiers init_declarator_list ';'",
"declaration specifiers: storage class specifier",
"declaration specifiers: storage_class_specifier declaration_specifiers",
"declaration specifiers: type specifier",
"declaration specifiers: type_specifier declaration_specifiers",
"init declarator list: init declarator",
"init_declarator_list : init_declarator_list ',' init_declarator",
"init declarator: declarator",
"init declarator : declarator '=' initializer",
"storage class specifier: TYPEDEF",
"storage class specifier: EXTERN",
```

```
"storage class specifier: STATIC",
"storage class specifier: AUTO",
"storage class specifier: REGISTER",
"type specifier: VOID",
"type specifier: CHAR",
"type specifier: SHORT",
"type_specifier: INT",
"type specifier: LONG",
"type_specifier: FLOAT",
"type specifier: DOUBLE",
"type specifier: SIGNED",
"type specifier: UNSIGNED",
"type specifier: struct or union specifier",
"specifier qualifier_list: type_specifier specifier_qualifier_list",
"specifier_qualifier_list: type_specifier",
"specifier_qualifier_list : CONST specifier_qualifier_list",
"specifier_qualifier_list : CONST",
"struct or union specifier: struct_or_union IDENTIFIER '{' struct_declaration_list '}' ';"",
"struct_or_union_specifier: struct_or_union '{' struct_declaration_list '}' ';"",
"struct_or_union_specifier: struct_or_union IDENTIFIER';",
"struct or union: STRUCT",
"struct_or_union: UNION",
"struct declaration list: struct declaration",
"struct declaration list: struct declaration list struct declaration",
"struct declaration: specifier qualifier list struct declarator list;",
"struct declarator_list : declarator",
"struct_declarator_list : struct_declarator_list ',' declarator",
"declarator: pointer direct_declarator",
"declarator: direct_declarator",
"direct declarator: IDENTIFIER".
"direct_declarator: '(' declarator')",
"direct_declarator: direct_declarator'[' constant_expression']",
"direct_declarator: direct_declarator'['']"",
"direct_declarator : direct_declarator '(' parameter_list ')",
"direct_declarator : direct_declarator '(' identifier_list ')",
"direct declarator: direct declarator '(' ')"",
"pointer: '*",
"pointer: '*' pointer",
"parameter list: parameter declaration",
"parameter list: parameter list', parameter declaration",
"parameter declaration: declaration_specifiers declarator",
"parameter_declaration : declaration_specifiers",
"identifier_list: IDENTIFIER",
"identifier_list: identifier_list',' IDENTIFIER",
"type name: specifier qualifier list",
"type name: specifier_qualifier_list declarator",
"initializer: assignment expression",
"initializer: '{' initializer list '}",
```

```
"initializer: '{' initializer list ',' '}",
"initializer list: initializer",
"initializer list: initializer list; initializer",
"statement: compound statement",
"statement: expression statement",
"statement: selection statement",
"statement: iteration_statement",
"statement: jump statement",
"compound_statement: '{' '}",
"compound_statement : '{' statement_list '}",
"compound_statement : '{' declaration_list '}",
"compound statement: '{' declaration list statement list '}",
"declaration list: declaration",
"declaration_list: declaration_list declaration",
"statement_list: statement",
"statement_list : statement_list statement",
"expression statement: ';'",
"expression_statement : expression ';"",
"selection statement: IF '(' expression ')' statement",
"selection_statement : IF '(' expression ')' statement ELSE statement",
"iteration statement: WHILE '(' expression ')' statement",
"iteration statement: FOR '(' expression statement expression statement'), statement",
"iteration statement: FOR '(' expression statement expression statement expression ')'
statement",
"jump_statement : CONTINUE ';",
"jump_statement : BREAK ';",
"jump_statement : RETURN ';"",
"jump statement: RETURN expression ';"",
"external declaration: function_definition",
"external declaration: declaration",
"function definition: declaration specifiers declarator declaration list
compound statement",
"function _definition : declaration_specifiers declarator compound_statement",
"function definition: declarator declaration_list compound_statement",
"function_definition: declarator compound_statement",
```

Design Strategy

Symbol Table Creation:

The symbol table is populated in the syntax phase and updated in semantic phase. It contains the following:

- Serial No The number of entries in the table.
- Identifier Name of the variables.
- Scope Scope of the variable.
- Value Mathematical value of the variable if defined.
- Type Data Type of the variable.
- Parameter type Data type of the function parameters.

Intermediate Code Generation:

The intermediate code is generated as it passes through the grammar, with help of some special sub-routines defined for the language clauses such as FOR, WHILE..etc. These routines are called as mid-rule actions.

Subroutine example for the WHILE Clause:

Code Optimization:

The input Intermediate Code is scanned and any values that are available at compile time are propagated forward. Then another pass is made to check if any arithmetic instructions can be done before moving to the nest phase.

Error Handling:

Scanner:

For any tokens which don't match the lexical rules, a -1 is returned to the parser.

Parser:

The error recovery is done using the 'error' keyword which is included in the grammar at certain places. There is a function yyerror() which handles the error, if the grammar isn't matched. It prints the line number at which the error occurs. It handles the error and continues parsing the rest of the code.

Semantic Analyzer:

The error handling here is done by writing certain mid-rule actions, which are basically java code, to check if the token passed makes semantic sense according to the language.

Target Code Generation:

There is no error handling in this phase. It is dependent on the previous phases to catch all errors.

Implementation Details

Symbol Table Creation:

The object of the following class was populated accordingly.

```
class Sym{
   public int sno;
   public int[] type = new int[100];
   public int[] paratype = new int[100];
   public int tn;
   public int pn;
   public int index;
   public int scope;
   public String token;
   public float fvalue;
   public Sym(int sno, int tn, int pn,
        int index, int scope, String token, float fvalue) {
           this.sno = sno;
            this.tn = tn;
            this.pn = pn;
            this.index = index;
           this.scope = scope;
           this.token = "";
           this.fvalue = fvalue;
```

Intermediation Code Generation:

The intermediate code while getting generated is printed out directly. But the ideal way to do it would be to store it in a Quadruples table.

Code Optimization:

First the input icg gets scanned so that all known values are stored in a dictionary which can later be propagated to other lines in the code.

Then the icg is scanned again to check for any arithmetic instructions that can be evaluated then and there. This step is done after the variable propagation is done.

Assembly Code Generation:

re

Error Handling:

As mentioned in the Design Strategy.

Results and Shortcomings

The mini-compiler does a good job in compiling real basic C-code to a hypothetical target code. It handles the error well, and has proper semantic

handling. It also does certain code optimizations.

Few shortcoming would be that it doesn't store the ICG in Quadruples table. Also, due to our lack of familiarity with Java, we weren't able to implement abstract syntax tree and hence scraped it.

Snapshots

Lexical Analysis

```
labhishek@hope -d/git/c-compiler/lexical_analysis|5 cat test case2.c
//with error - for loop syntax error
#include <stdio.h>
int main()
{
    int a=4, i;
    for(i=0;i<10)
    {
        printf("%d",i);
    }
        int x;
}
labhishek@hope -d/git/c-compiler/lexical_analysis|$</pre>

I
```

```
******** SYMBOL TABLE *******

line: 2 type: KEYWORD token: int
line: 2 type: IDENTIFIER token: main()
line: 3 type: SPECIAL SYMBOL token: {
line: 4 type: IDENTIFIER token: a
line: 4 type: OPERATOR token: =
line: 4 type: SPECIAL SYMBOL token: 4
line: 4 type: SPECIAL SYMBOL token: ,
line: 4 type: SPECIAL SYMBOL token: ,
line: 4 type: SPECIAL SYMBOL token: ;
line: 5 type: SPECIAL SYMBOL token: ;
line: 5 type: KEYWORD token: for
line: 5 type: SPECIAL SYMBOL token: (
line: 5 type: OPERATOR token: <
line: 5 type: INTGER CONSTANT token: 0
line: 5 type: OPERATOR token: <
line: 5 type: SPECIAL SYMBOL token: )
line: 7 type: SPECIAL SYMBOL token: )
line: 7 type: SPECIAL SYMBOL token: )
line: 7 type: SPECIAL SYMBOL token: )
line: 8 type: SPECIAL SYMBOL token: }
line: 9 type: IDENTIFIER token: X
```

Syntax Analysis: (wrong for loop syntax)

```
labhishek@hope ~d/git/c-compiler/syntax_analysis $ ./run test case2.c
//with error - for loop syntax error
#include <stdio.h>
int main()
{
    int a=4, i;
    for(i=0;i<10)
    {
        printf("%d",i);
    }
        int x;
}

Syntax Error at: line: 5 token: )
Syntax Error at: line: 10 token: }

Parsing Failed!
line: 2 type: FUNCTION token: main
line: 4 type: INT token: a
line: 4 type: INT token: i
line: 7 type: FUNCTION token: printf
line: 7 type: STRING token: %d
line: 9 type: INT token: x
labhishek@hope ~d/git/c-compiler/syntax_analysis $ []</pre>
```

Semantic Analysis: (scope of variable b, semicolon missing)

Intermediate Code Generation:

```
abhishek@hope ~d/git/c-compiler/icg $ ./run test_case2.c
#include <stdio.h>
int main()
{
    int a=s;
    int b=s;
    int b=s;
    int b=s;
    if(a<-7) {
        b=b-4;
    }
    else {
        b=b+3;
    }
    return 0;
}

a = 5
b = 6
t0 = a <- 7
t1 = not t0
if t1 goto L1
t2 = b - 4
b = t2
goto L2
L1
t3 = b + 3
b = t3
L2
Parsing Complete and Ok!
abhishek@hope ~d/git/c-compiler/icg $</pre>
```

Assembly Code Generation:

```
[abhishek@hope ~d/git/c-compiler/assembly_gen]$ ./run test case2.c
| abhishek@hope
| a = 5
| b = 6
| t0 = a <= 7
| t1 = not t0
| if t1 goto L1
| t2 = b - 4
| b = t2
| goto L2
| L1
| t3 = b + 3
| b = t3
| L2
 ASSEMBLY CODE:
 MOV RØ, 5
SW RØ, a
MOV RØ, 6
SW RØ, b
LW R1, a
MOV R2, 7
CMP R1, R2
 LW R1, b
 MOV R2, 4
SUB R0, R2, R1
SW R0, t2
 LW R0, t2
SW R0, b
 LW R1, b
MOV R2, 3
ADD R0, R1, R2
SW R0, t3
 LW R0, t3
SW R0, b
```

Conclusions:

The mini-compiler supports the clauses mentioned at the start of the project. It parses the source code, displays any errors with line number mentioned if found. Optimizes the code if possible and generates the target machine code. This concludes our project.

Future Enhancements:

In future, we can extend it to support a larger subset of the C-grammar. Also we could generate actual machine code for some architecture like x86 or amd64.