

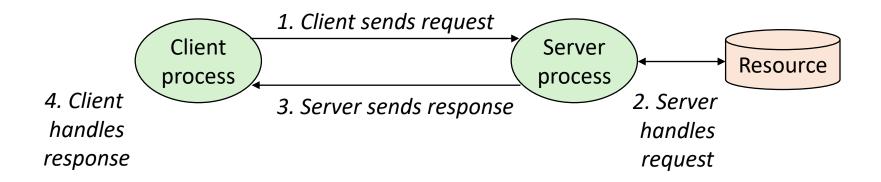
# Computer Architecture and Operating Systems Lecture 12: Basics of Networking

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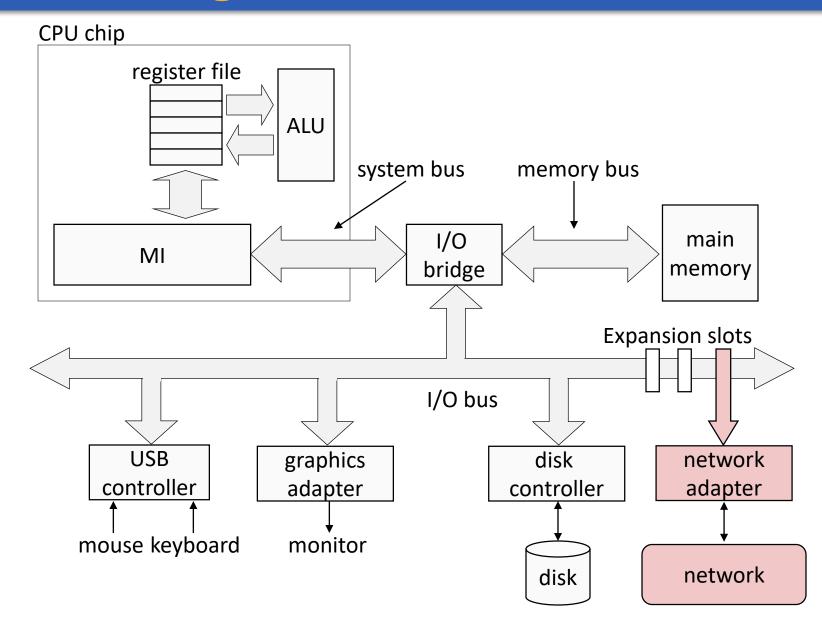
#### Client-Server Architecture

- Most network applications are based on the client-server model:
  - A *server* process and one or more *client* processes
  - Server manages some resource
  - Server provides service by manipulating resource for clients
  - Server activated by request from client (vending machine analogy)



Note: clients and servers are processes running on hosts (can be the same or different hosts)

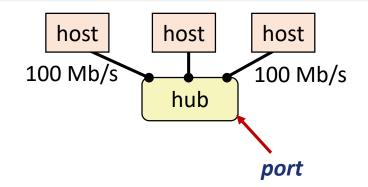
# Hardware Organization of a Network Host



## Computer Networks

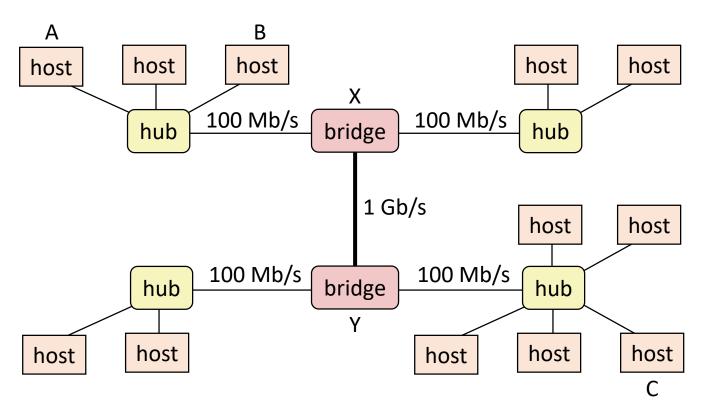
- A network is a hierarchical system of boxes and wires organized by geographical proximity
  - SAN (System Area Network) spans cluster or machine room
    - Switched Ethernet, Quadrics QSW, ...
  - LAN (Local Area Network) spans a building or campus
    - Ethernet is most prominent example
  - WAN (Wide Area Network) spans country or world
    - Typically high-speed point-to-point phone lines
- An internetwork (internet) is an interconnected set of networks
  - The Global IP Internet (uppercase "I") is the most famous example of an internet (lowercase "i")
- Let us see how an internet is built from the ground up

# Lowest Level: Ethernet Segment



- Ethernet segment consists of a collection of hosts connected by wires (twisted pairs) to a hub
- Spans room or floor in a building
- Operation
  - Each Ethernet adapter has a unique 48-bit address (MAC address)
    - E.g., 00:16:ea:e3:54:e6
  - Hosts send bits to any other host in chunks called frames
  - Hub slavishly copies each bit from each port to every other port
    - Every host sees every bit
    - Note: Hubs are on their way out. Bridges (switches, routers) became cheap enough to replace them

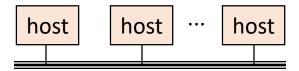
# Next Level: Bridged Ethernet Segment



- Spans building or campus
- Bridges cleverly learn which hosts are reachable from which ports and then selectively copy frames from port to port

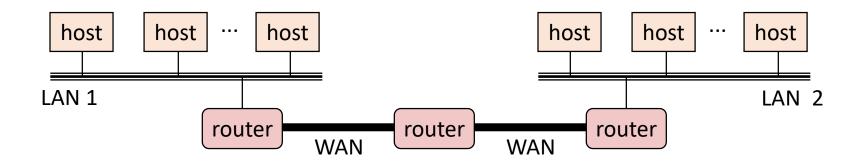
## Conceptual View of LANs

For simplicity, hubs, bridges, and wires are often shown as a collection of hosts attached to a single wire:



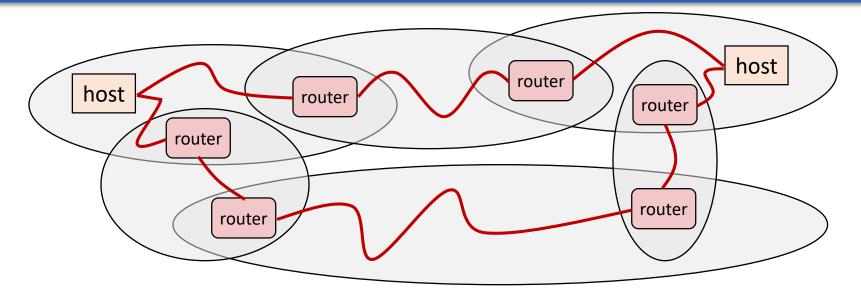
#### **Next Level: Internets**

- Multiple incompatible LANs can be physically connected by specialized computers called *routers*
- The connected networks are called an *internet* (lower case)



LAN 1 and LAN 2 might be completely different, totally incompatible (e.g., Ethernet, Fibre Channel, 802.11\*, T1-links, DSL, ...)

# Logical Structure of Internet



- Ad hoc interconnection of networks
  - No particular topology
  - Vastly different router & link capacities
- Send packets from source to destination by hopping through networks
  - Router forms bridge from one network to another
  - Different packets may take different routes

## The Notion of an Internet Protocol

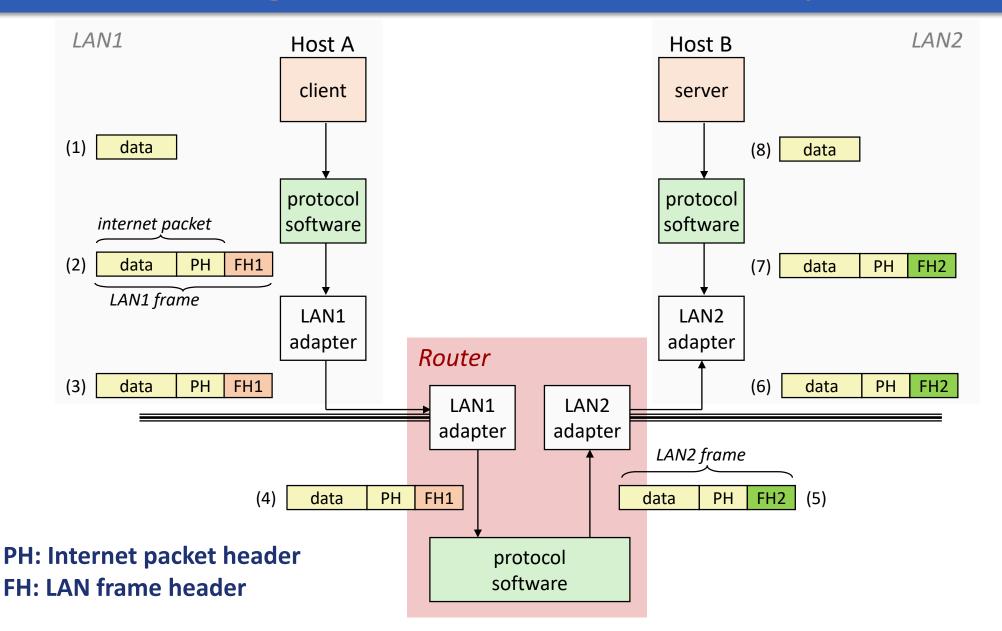
• How is it possible to send bits across incompatible LANs and WANs?

- Solution: protocol software running on each host and router
  - Protocol is a set of rules that governs how hosts and routers should cooperate when they transfer data from network to network.
  - Smooths out the differences between the different networks

## What Does an internet Protocol Do?

- Provides a naming scheme
  - An internet protocol defines a uniform format for host addresses
  - Each host (and router) is assigned at least one of these internet addresses that uniquely identifies it
- Provides a delivery mechanism
  - An internet protocol defines a standard transfer unit (packet)
  - Packet consists of *header* and *payload* 
    - Header: contains info such as packet size, source and destination addresses
    - Payload: contains data bits sent from source host

## Transferring Internet Data Via Encapsulation



### Other Issues

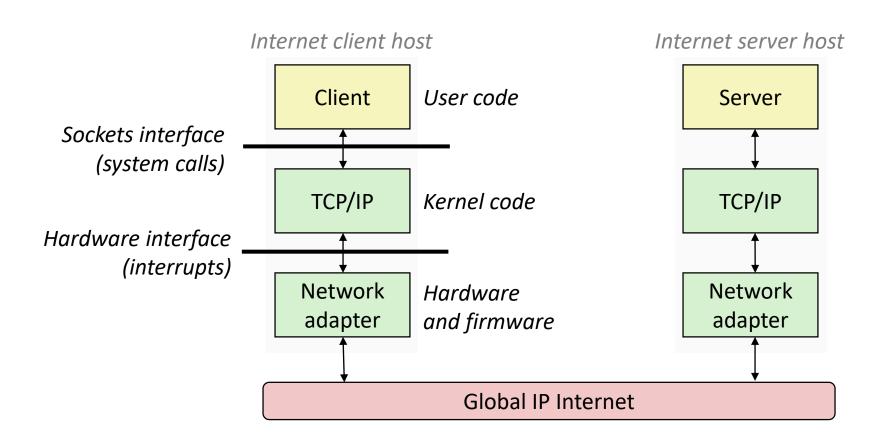
- We are glossing over a number of important questions:
  - What if different networks have different maximum frame sizes? (segmentation)
  - How do routers know where to forward frames?
  - How are routers informed when the network topology changes?
  - What if packets get lost?

These (and other) questions are addressed by the area of systems known as computer networking

# Global IP Internet (upper case)

- Most famous example of an internet
- Based on the TCP/IP protocol family
  - IP (Internet Protocol) :
    - Provides basic naming scheme and unreliable delivery capability of packets (datagrams) from host-to-host
  - UDP (Unreliable Datagram Protocol)
    - Uses IP to provide unreliable datagram delivery from process-to-process
  - TCP (Transmission Control Protocol)
    - Uses IP to provide *reliable* byte streams from *process-to-process* over connections
- Accessed via a mix of Unix file I/O and functions from the sockets interface

## Organization of an Internet Application



# A Programmer's View of the Internet

- 1. Hosts are mapped to a set of 32-bit *IP addresses* 
  - **1**28.2.203.179

- 2. The set of IP addresses is mapped to a set of identifiers called Internet *domain names* 
  - 128.2.203.179 is mapped to www.cs.cmu.edu
- 3. A process on one Internet host can communicate with a process on another Internet host over a connection

#### Aside: IPv4 and IPv6

- The original Internet Protocol, with its 32-bit addresses, is known as *Internet Protocol Version 4* (IPv4)
- 1996: Internet Engineering Task Force (IETF) introduced Internet Protocol Version 6 (IPv6) with 128-bit addresses
  - Intended as the successor to IPv4
- As of 2015, vast majority of Internet traffic still carried by IPv4
  - Only 4% of users access Google services using IPv6.
- We will focus on IPv4, but will show you how to write networking code that is protocol-independent.

#### IP Addresses

- 32-bit IP addresses are stored in an IP address struct
  - IP addresses are always stored in memory in *network byte order* (big-endian byte order)
  - True in general for any integer transferred in a packet header from one machine to another.
    - E.g., the port number used to identify an Internet connection.

```
/* Internet address structure */
struct in_addr {
   uint32_t s_addr; /* network byte order (big-endian) */
};
```

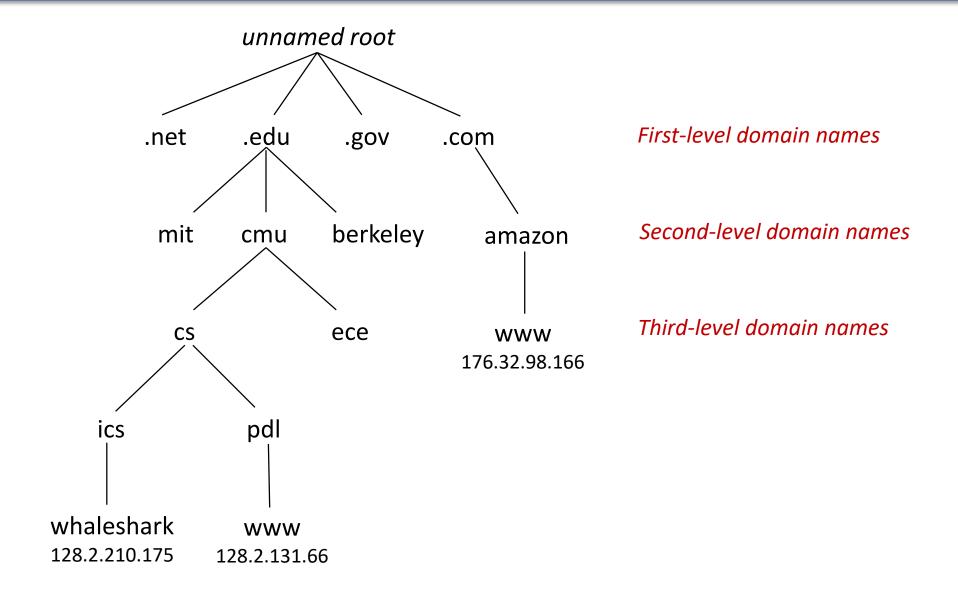
## **Dotted Decimal Notation**

 By convention, each byte in a 32-bit IP address is represented by its decimal value and separated by a period

■ IP address: 0x8002C2F2 = 128.2.194.242

•Use getaddrinfo and getnameinfo functions (described later) to convert between IP addresses and dotted decimal format.

## Internet Domain Names



# Domain Naming System (DNS)

The Internet maintains a mapping between IP addresses and domain names in a huge worldwide distributed database called DNS

- Conceptually, programmers can view the DNS database as a collection of millions of host entries.
  - Each host entry defines the mapping between a set of domain names and IP addresses.
  - In a mathematical sense, a host entry is an equivalence class of domain names and IP addresses.

## Properties of DNS Mappings

- Can explore properties of DNS mappings using nslookup
  - Output edited for brevity

```
linux> nslookup localhost
Address: 127.0.0.1
```

■ Each host has a locally defined domain name localhost which always maps to the loopback address 127.0.0.1

```
linux> hostname
whaleshark.ics.cs.cmu.edu
```

Use hostname to determine real domain name of local host:

# Properties of DNS Mappings (cont)

Simple case: one-to-one mapping between domain name and IP address:

```
linux> nslookup whaleshark.ics.cs.cmu.edu
Address: 128.2.210.175
```

• Multiple domain names mapped to the same IP address:

```
linux> nslookup cs.mit.edu
Address: 18.62.1.6
linux> nslookup eecs.mit.edu
Address: 18.62.1.6
```

# Properties of DNS Mappings (cont)

• Multiple domain names mapped to multiple IP addresses:

```
linux> nslookup www.twitter.com
Address: 199.16.156.6
Address: 199.16.156.70
Address: 199.16.156.102
Address: 199.16.156.230

linux> nslookup twitter.com
Address: 199.16.156.102
Address: 199.16.156.230
Address: 199.16.156.6
Address: 199.16.156.70
```

Some valid domain names don't map to any IP address:

```
linux> nslookup ics.cs.cmu.edu

*** Can't find ics.cs.cmu.edu: No answer
```

#### **Internet Connections**

- Clients and servers communicate by sending streams of bytes over connections. Each connection is:
  - Point-to-point: connects a pair of processes.
  - Full-duplex: data can flow in both directions at the same time,
  - Reliable: stream of bytes sent by the source is eventually received by the destination in the same order it was sent.
- A socket is an endpoint of a connection
  - Socket address is an IPaddress:port pair
- A port is a 16-bit integer that identifies a process:
  - *Ephemeral port*: Assigned automatically by client kernel when client makes a connection request.
  - Well-known port: Associated with some service provided by a server (e.g., port 80 is associated with Web servers)

## Well-known Ports and Service Names

Popular services have permanently assigned well-known ports and corresponding well-known service names:

echo server: 7/echo

ssh servers: 22/ssh

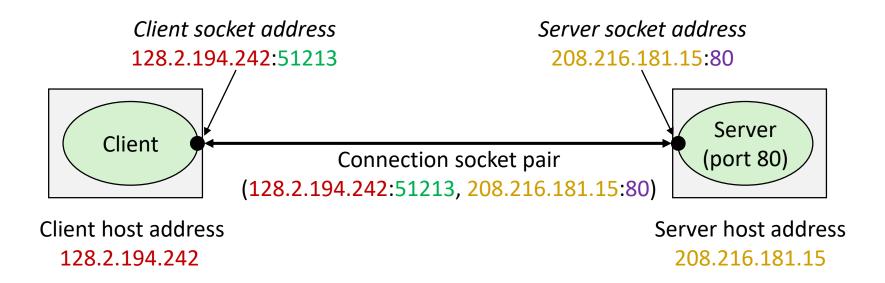
email server: 25/smtp

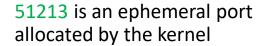
Web servers: 80/http

• Mappings between well-known ports and service names is contained in the file /etc/services on each Linux machine.

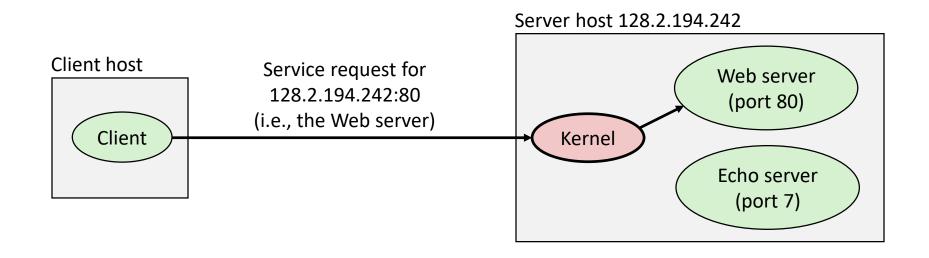
# Anatomy of a Connection

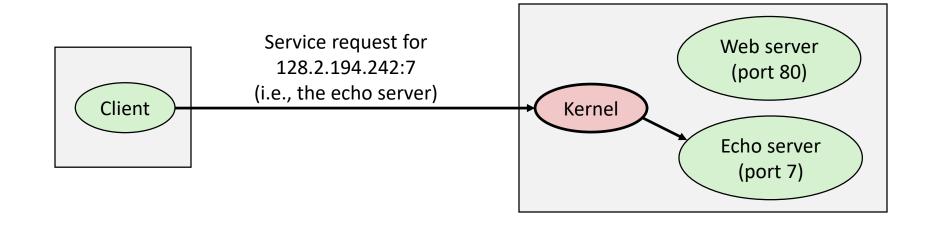
- A connection is uniquely identified by the socket addresses of its endpoints (socket pair)
  - cliaddr:cliport, servaddr:servport)





# Using Ports to Identify Services

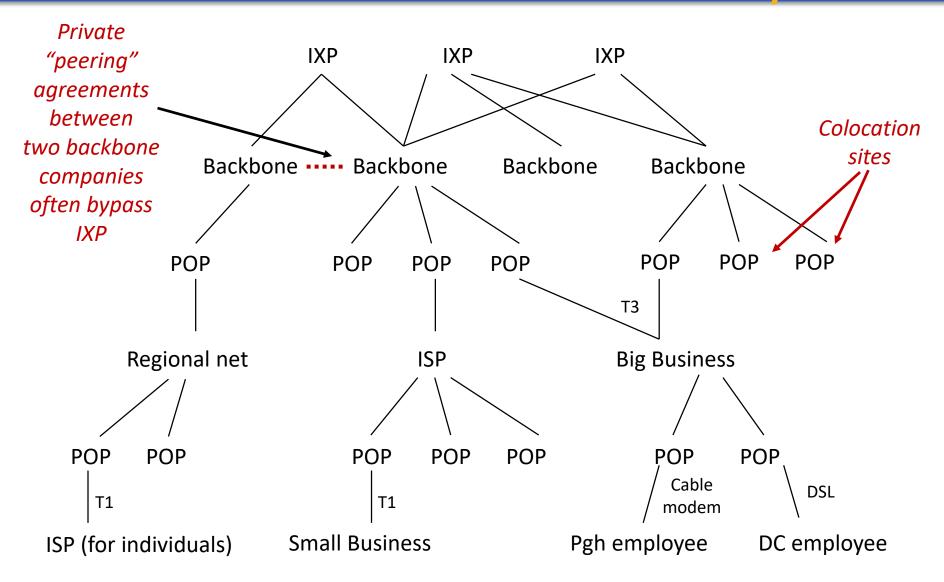




## **Basic Internet Components**

- Internet backbone:
  - collection of routers (nationwide or worldwide) connected by high-speed point-to-point networks
- Internet Exchange Points (IXP):
  - router that connects multiple backbones (often referred to as peers)
  - Also called Network Access Points (NAP)
- Regional networks:
  - smaller backbones that cover smaller geographical areas (e.g., cities or states)
- Point of presence (POP):
  - machine that is connected to the Internet
- Internet Service Providers (ISPs):
  - provide dial-up or direct access to POPs

#### Internet Connection Hierarchy



## **IP Address Structure**

■ IP (V4) Address space divided into classes:

	0	1	2 3	3	8	16	!	24		<u>31</u>
Class A	0	Net ID				Host ID				
Class B	1	0			Ne	t ID	ŀ	Host ID		
Class C	1	1	0 Net ID						Host ID	
Class D	1	1	1 0 Multicast address							
Class E	1	1	1	1 Reserved for experiments						

- Network ID Written in form w.x.y.z/n
  - n = number of bits in host address
  - E.g., CMU written as 128.2.0.0/16
    - Class B address
- Unrouted (private) IP addresses:
  10.0.0.0/8 172.16.0.0/12 192.168.0.0/16