



Computer Architecture and Operating Systems Lecture 13: Data-level parallelism: Vector, SIMD, GPU

Andrei Tatarnikov

atatarnikov@hse.ru
@andrewt0301

Instruction and Data Streams

An alternate classification

		Data Streams	
		Single	Multiple
Instruction Streams	Single	SISD: Intel Pentium 4	SIMD: SSE instructions of x86
	Multiple	MISD: No examples today	MIMD: Intel Xeon e5345

- ■SPMD: Single Program Multiple Data
 - A parallel program on a MIMD computer
 - Conditional code for different processors

Types of Parallel Processing

- Single instruction, single data (SISD) stream: A single processor executes a single instruction stream to operate on data stored in a single memory. Uniprocessors fall into this category.
- Single instruction, multiple data (SIMD) stream: A single machine instruction controls the simultaneous execution of a number of processing elements on a lockstep basis. Each has an associated data memory, so that instructions are executed on different sets of data by different processors. Vector and array processors fall into this category.
- Multiple instruction, single data (MISD) stream: A sequence of data is transmitted to a set of processors, each of which executes a different instruction sequence. Not commercially implemented.
- Multiple instruction, multiple data (MIMD) stream: A set of processors simultaneously execute different instruction sequences on different data sets. SMPs, clusters, and NUMA systems fit into this category.

Vector Processors

- Highly pipelined function units
- Stream data from/to vector registers to units
 - Data collected from memory into registers
 - Results stored from registers to memory
- Example: Vector extension to RISC-V
 - v0 to v31: 32 × 64-element registers, (64-bit elements)
 - Vector instructions
 - fld.v, fsd.v: load/store vector
 - fadd.d.v: add vectors of double
 - fadd.d.vs: add scalar to each element of vector of double
- Significantly reduces instruction-fetch bandwidth

Example: DAXPY $(Y = a \times X + Y)$

Conventional RISC-V code:

```
f0,a(x3) # load scalar a
     fld
              x5, x19, 512
     addi
                           # end of array X
loop: fld
              f1,0(x19)
                           # load x[i]
             f1,f1,f0
                           # a * x[i]
     fmul.d
     fld
              f2,0(x20)
                           # load y[i]
     fadd.d
              f2, f2, f1
                           \# a * x[i] + y[i]
              f2,0(x20)
     fsd
                           # store y[i]
     addi
              x19, x19, 8
                           # increment index to x
              x20,x20,8
                           # increment index to y
     addi
              x19, x5, loop
                           # repeat if not done
     bltu
```

Vector RISC-V code:

```
fld f0,a(x3) # load scalar a
fld.v v0,0(x19) # load vector x
fmul.d.vs v0,v0,f0 # vector-scalar multiply
fld.v v1,0(x20) # load vector y
fadd.d.v v1,v1,v0 # vector-vector add
fsd.v v1,0(x20) # store vector y
```

Vector vs. Scalar

- Vector architectures and compilers
 - Simplify data-parallel programming
 - Explicit statement of absence of loop-carried dependences
 - Reduced checking in hardware
 - Regular access patterns benefit from interleaved and burst memory
 - Avoid control hazards by avoiding loops
- More general than ad-hoc media extensions (such as MMX, SSE)
 - Better match with compiler technology

SIMD

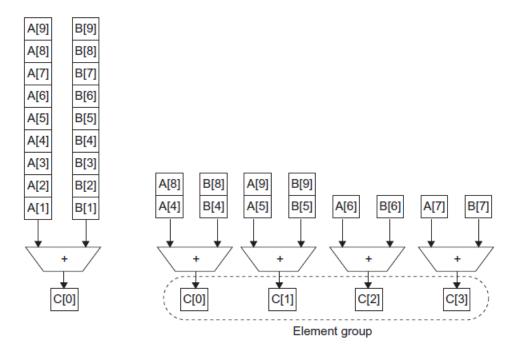
- Operate elementwise on vectors of data
 - E.g., MMX and SSE instructions in x86
 - Multiple data elements in 128-bit wide registers
- •All processors execute the same instruction at the same time
 - Each with different data address, etc.
- Simplifies synchronization
- Reduced instruction control hardware
- Works best for highly data-parallel applications

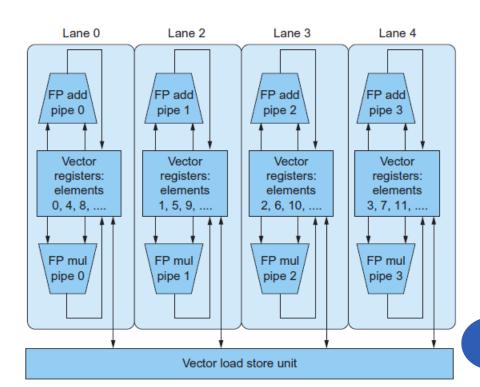
Vector vs. Multimedia Extensions

- Vector instructions have a variable vector width, multimedia extensions have a fixed width
- Vector instructions support strided access, multimedia extensions do not

Vector units can be combination of pipelined and arrayed functional

units:





GPU Architectures

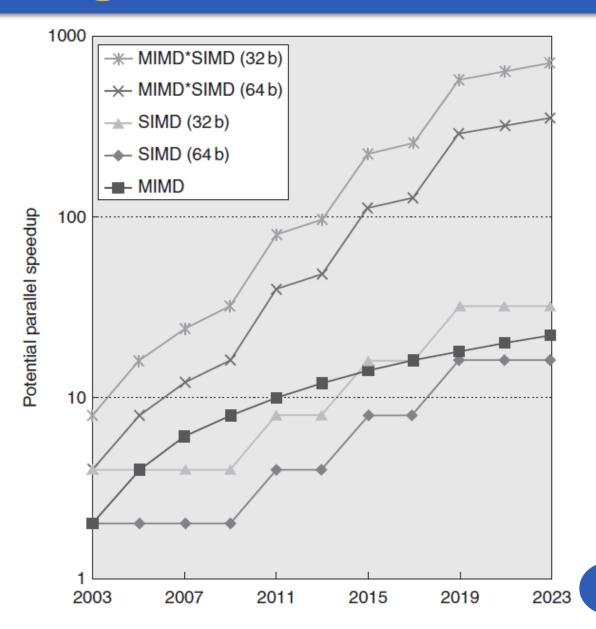
- Processing is highly data-parallel
 - GPUs are highly multithreaded
 - Use thread switching to hide memory latency
 - Less reliance on multi-level caches
 - Graphics memory is wide and high-bandwidth
- Trend toward general purpose GPUs
 - Heterogeneous CPU/GPU systems
 - CPU for sequential code, GPU for parallel code
- Programming languages/APIs
 - DirectX, OpenGL
 - C for Graphics (Cg), High Level Shader Language (HLSL)
 - Compute Unified Device Architecture (CUDA)

History of GPUs

- Early video cards
 - Frame buffer memory with address generation for video output
- 3D graphics processing
 - Originally high-end computers (e.g., SGI)
 - Moore's Law ⇒ lower cost, higher density
 - 3D graphics cards for PCs and game consoles
- Graphics Processing Units
 - Processors oriented to 3D graphics tasks
 - Vertex/pixel processing, shading, texture mapping, rasterization

Concluding Remarks

• SIMD and vector operations match multimedia applications and are easy to program



Any Questions?

```
__start: addi t1, zero, 0x18
    addi t2, zero, 0x21

cycle: beg t1, t2, done
    slt t0, t1, t2
    bne t0, zero, if_less
    nop
    sub t1, t1, t2
    j cycle
    nop

if_less: sub t2, t2, t1
    j cycle

done: add t3, t1, zero
```