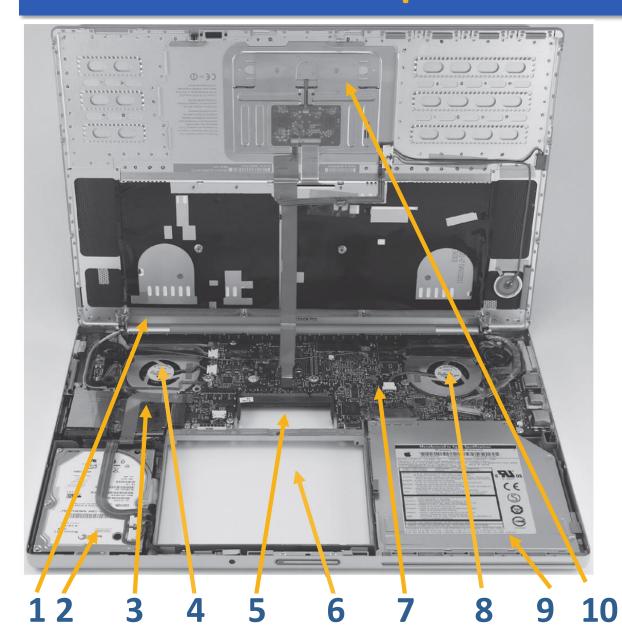


# Computer Architecture and Operating Systems Lecture 3: Computer Architecture

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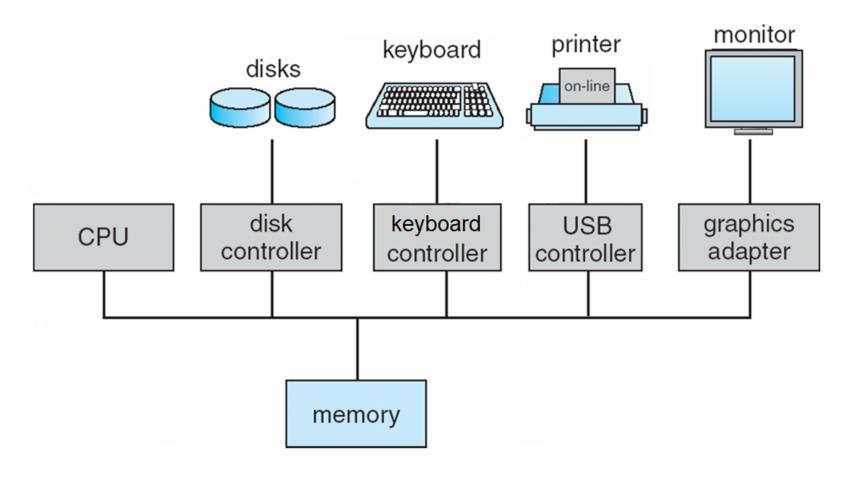
#### Computer Under Cover



- 1. Monitor
- 2. Hard drive
- 3. CPU (Processor)
- 4. Fan with cover
- 5. Spot for memory DIMMs
- 6. Spot for battery
- 7. Motherboard
- 8. Fan with cover
- 9. DVD drive
- 10.Keyboard

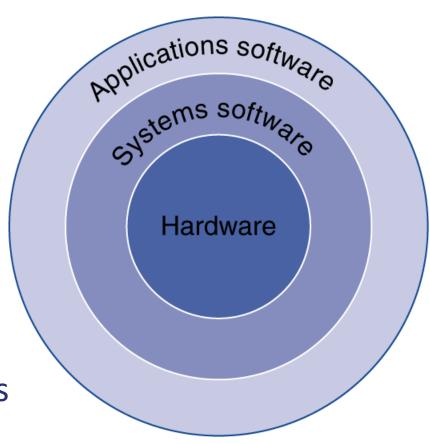
# Computer Organization

 One or more CPUs and device controllers connected through a bus providing access to shared memory



## Program Under Hood

- Application software
  - Written in high-level language
- System software
  - Compiler: translates high-level language code to machine code
  - Operating System: service code
    - Handling input/output
    - Managing memory and storage
    - Scheduling tasks & sharing resources
- Hardware
  - CPU, memory, I/O controllers



# Levels of Program Code

High-level language program (in C)

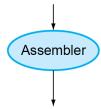
```
swap(int v[], int k)
{int temp;
    temp = v[k];
    v[k] = v[k+1];
    v[k+1] = temp;
}

Compiler
```

Assembly language program (for RISC-V)

```
swap:

slli x6, x11, 3
add x6, x10, x6
ld x5, 0(x6)
ld x7, 8(x6)
sd x7, 0(x6)
sd x5, 8(x6)
jalr x0, 0(x1)
```

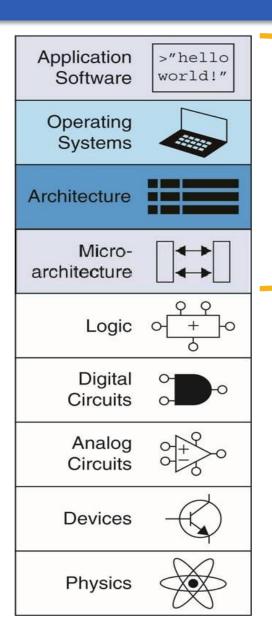


Binary machine language program (for RISC-V) 

- High-level language
  - Level of abstraction closer to problem domain
  - Provides productivity and portability
- Assembly language
  - Textual representation of instructions
- Hardware representation
  - Binary digits (bits)
  - Encoded instructions and data

#### Abstractions

- Abstraction helps us deal with complexity
  - Hide lower-level detail
- Instruction set architecture (ISA)
  - The hardware/software interface
- Application binary interface (ABI)
  - The ISA plus system software interface
- Implementation (microarchitecture)
  - The details underlying the interface



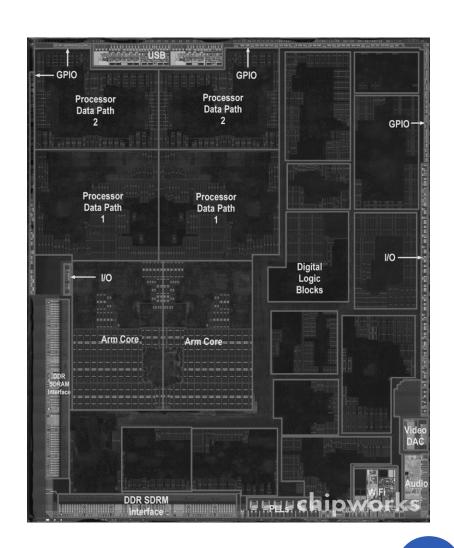
Focus of this course

## Inside the Processor (CPU)

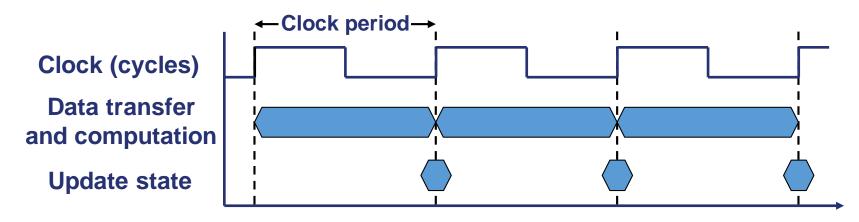
Central Processing Unit (CPU) is the heart of any computer system.

Main components:

- Register file: small fast memory for immediate access to data
- Datapath: performs operations on data
- Control unit: sequences datapath, memory, etc.



# **CPU Clocking**



- Operation of digital hardware governed by a constantrate clock
- Clock period: duration of a clock cycle
  - e.g., 250 ps = 0.25 ns =  $250 \times 10^{-12}$  s
- Clock frequency (rate): cycles per second
  - e.g.,  $4.0 \text{ GHz} = 4000 \text{ MHz} = 4.0 \times 10^9 \text{ Hz}$

#### **CPU Time**

$$CPUTime = \frac{Instructions}{Program} \times \frac{Clock \, cycles}{Instruction} \times \frac{Seconds}{Clock \, cycle}$$

- Performance depends on
  - Algorithm: affects IC, possibly CPI
  - Programming language: affects IC, CPI
  - Compiler: affects IC, CPI
  - Instruction set architecture: affects IC, CPI, T<sub>c</sub>

### Instruction Set Architecture (ISA)

Instruction set architecture (ISA) is the interface between the hardware and the lowest-level software. This is one of the most important abstractions.

#### **ISA Classification**

- Complex instruction set computer (CISC)
  - x86/x64 (Intel and AMD)
- Reduced instruction set computer (RISC)
  - ARM, PowerPC, MIPS, RISC-V
- Very long instruction word (VLIW)
  - Itanium, Elbrus

# Reduced Instruction Set Computing (RISC)

Reduced Instruction Set Computing (RISC) concept was proposed by teams of researchers at Stanford University (John Hennessy) and University of California Berkeley (David Paterson) in early 1980s as an alternative of Complex Instruction Set Computing (CISC) dominating at that time.

#### **RISC Principles**

- All instructions are executed by hardware
- Maximize the rate at which instructions are issued
- Instructions should be easy to decode
- Only loads and stores should reference memory
- Provide plenty of registers

#### RISC-V ISA

- Simple ISA by UC Berkeley (2010)
- Open and Free
- Wide-Purpose Configurable ISA (from IoT to mainframes)
- Maintained by RISC-V Foundation (moved to Switzerland)
- Supported by many IT Companies and Universities





### RISC-V Community

Wide Support of IT Companies (except Intel and ARM) and Universities



















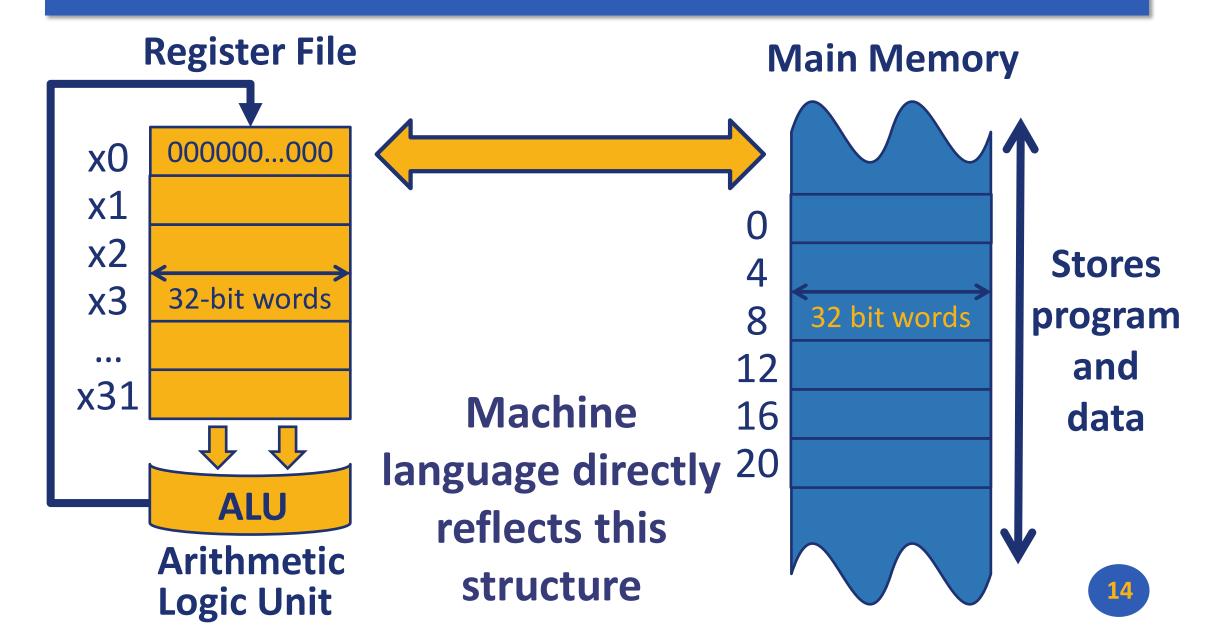






and many others...

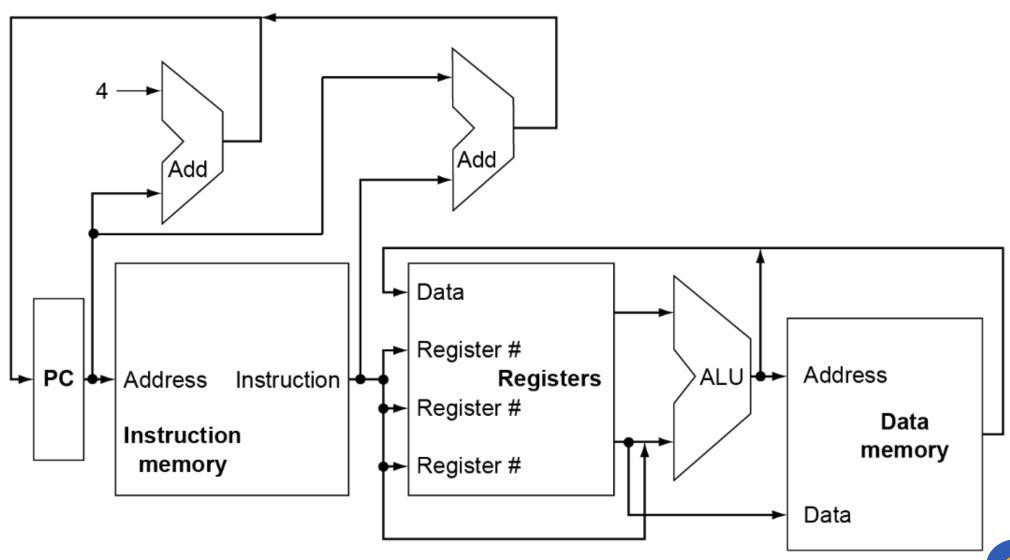
#### How CPU Works



#### Instruction Execution

- 1. Fetch next instruction from memory into instruction register
- 2. Change program counter to point to next instruction
- 3. Determine type of instruction just fetched
- 4. If instructions uses word in memory, determine where Fetch word, if needed, into CPU register
- 5. Execute the instruction
- 6. Go to step 1 to begin executing following instruction

## RISC-V CPU Scheme



# RISC-V General-Purpose Registers

Register	Name	Use	Saver
х0	zero	constant 0	n/a
<b>x1</b>	ra	return addr	caller
x2	sp	stack ptr	callee
х3	gp	gbl ptr	
x4	tp	thread ptr	
x5-x7	t0-t2	temporaries	caller
x8	s0/fp	saved/ frame ptr	callee
х9	s1	saved	callee
x10-x17	a0-a7	arguments	caller
x18-x27	s2-s11	saved	callee
x28-x31	t3-t6	temporaries	caller

**32 Registers** 

32 (or 64) Bits Wide

#### **RISC-V Instructions**

- Fixed-size 32 bit instructions
- Always three operands: d -> op(s, t)
- Instruction types
  - Computational instructions
  - Load-store instructions
  - Control-transfer instructions
  - System instructions
- •All operations done with registers

### Any Questions?

```
__start: addi t1, zero, 0x18
addi t2, zero, 0x21

cycle: beg t1, t2, done
slt t0, t1, t2

kne t0, zero, if_less

nop
sub t1, t1, t2

j cycle

nop

if_less: sub t2, t2, t1

j cycle

done: add t3, t1, zero
```