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# Computer Architecture and Operating Systems

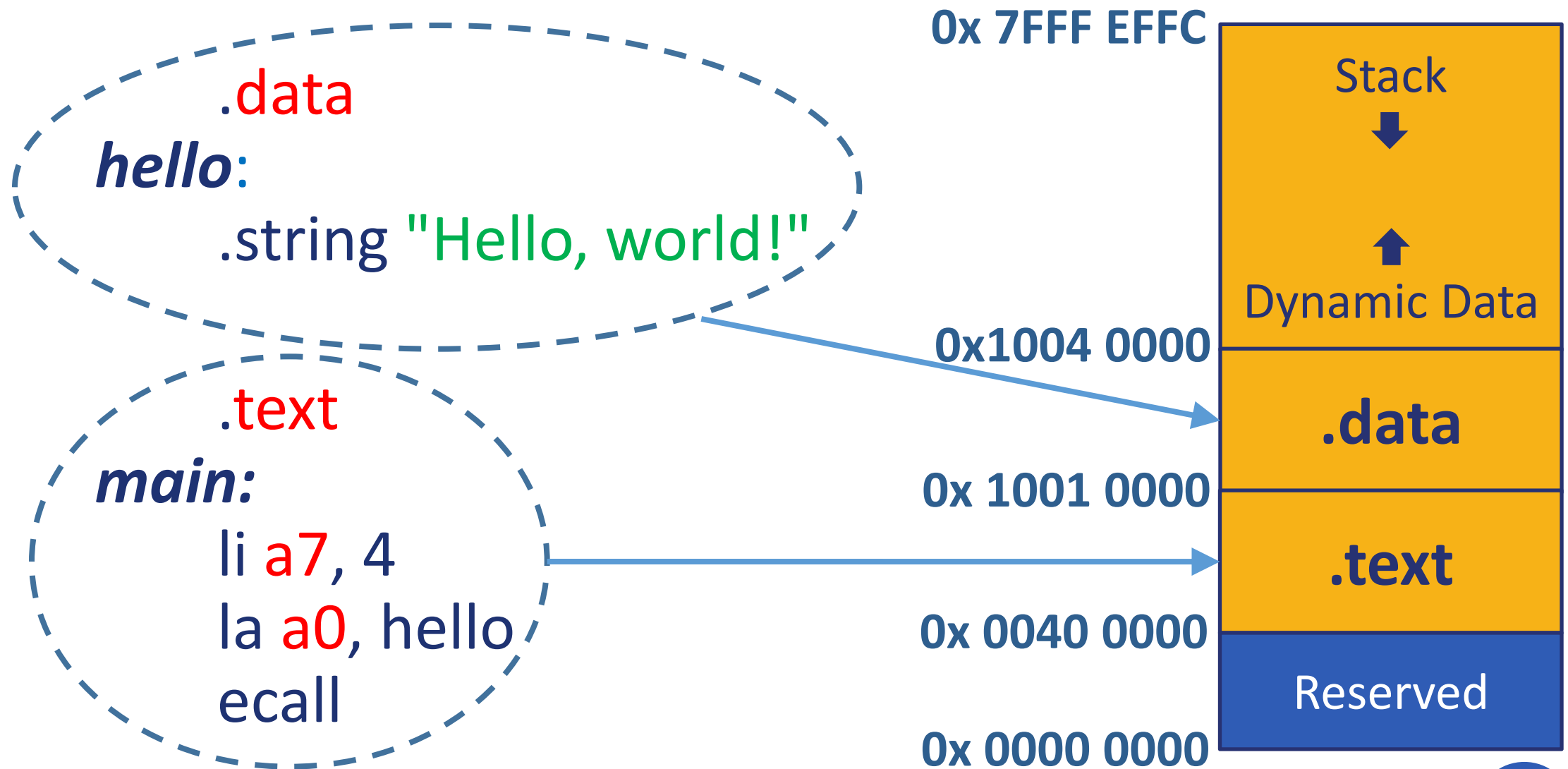
## Lecture 5: Assembly Programming – Branches and Arrays

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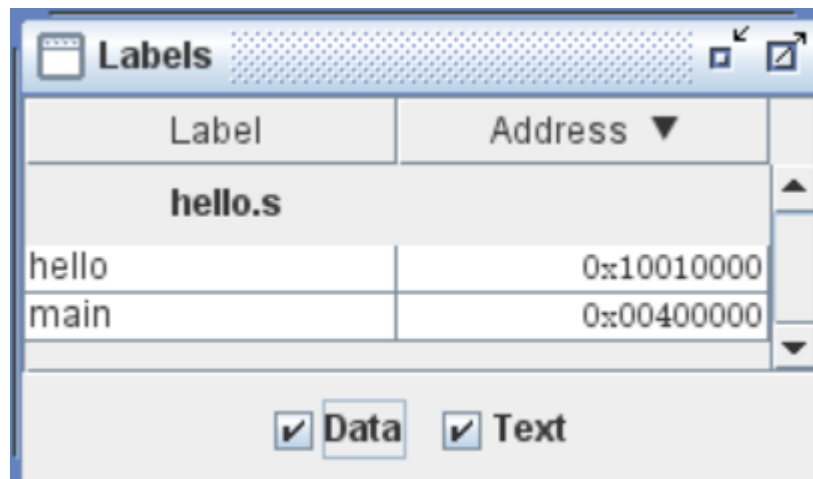
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# Program Structure and Memory Layout



# Labels

- **Labels** are symbolic names for addresses (in the .data or .text segment).
- **Labels** are used by control-flow instructions (branches and jumps).
- **Labels** are used by load and store instructions.



The screenshot shows a window titled 'Labels' with a table of labels. The table has two columns: 'Label' and 'Address'. The file 'hello.s' is selected. The table lists two labels: 'hello' at address '0x10010000' and 'main' at address '0x00400000'. At the bottom, there are checkboxes for 'Data' and 'Text', both of which are checked.

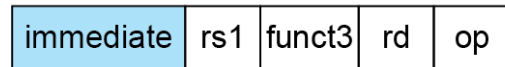
Label	Address ▼
<b>hello.s</b>	
hello	0x10010000
main	0x00400000

☒ Data ☒ Text

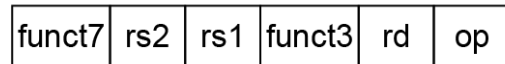
# Addressing

## Addresses can be represented in several ways

### 1. Immediate addressing



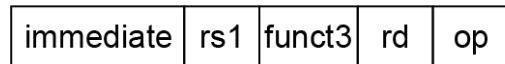
### 2. Register addressing



Registers

Register

### 3. Base addressing



Memory

Register

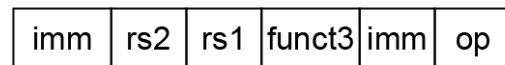


Byte Halfword

Word

Doubleword

### 4. PC-relative addressing



Memory

PC



Word

# Program Counter

- **Program Counter (PC)** is a special register that stores the address of the currently executed instruction.
- When an instruction is executed, the **PC** is incremented by the size of the instruction (4 bytes) to point to the next instruction.
- Branch and jump instructions assign to the **PC** new addresses to change the control flow.
- Branch instructions use **PC**-relative addresses (increment or decrement current value by an offset).

# Branch Instructions

## Branch Instructions

- Branch = `beq rs1, rs2, label`
- Branch  $\neq$  `bne rs1, rs2, label`
- Branch  $<$  `blt rs1, rs2, label`
- Branch  $\geq$  `bge rs1, rs2, label`
- Branch  $<$  Unsigned `bltu rs1, rs2, label`
- Branch  $\geq$  Unsigned `bgeu rs1, rs2, label`

# Branch Pseudoinstructions

## Branch Pseudoinstructions

▪ Branch unconditionally	<code>b</code>	<code>label</code>
▪ Branch = 0	<code>beqz</code>	<code>rs1, label</code>
▪ Branch $\geq$ 0	<code>bgez</code>	<code>rs1, label</code>
▪ Branch >	<code>bgt</code>	<code>rs1, rs2, label</code>
▪ Branch > Unsigned	<code>bgtu</code>	<code>rs1, rs2, label</code>
▪ Branch > 0	<code>bgtz</code>	<code>rs1, label</code>
▪ Branch $\leq$	<code>ble</code>	<code>rs1, rs2, label</code>
▪ Branch $\leq$ Unsigned	<code>bleu</code>	<code>rs1, rs2, label</code>
▪ Branch $\leq$ 0	<code>blez</code>	<code>rs1, label</code>
▪ Branch < 0	<code>bltz</code>	<code>rs1, label</code>
▪ Branch $\neq$ 0	<code>bnez</code>	<code>rs1, label</code>

# Branches and Program Counter

- Branch instructions are **PC**-relative
- They add a **12-bit** signed immediate to **PC**
- The immediate is an offset from **PC** to the target label
- The branch address range is  $\pm 2^{12}$  (4096 B = 4 KB)
- **PC** can be read with the **auipc** instruction

main:

```
auipc a0, 0    # a0 = PC + 0
li     a7, 34  # Print as hex
ecall                # Print a0
```



# Assembly Code for “If-Then-Else”

```
if (t0 == 0) {  
    t1 = 1;  
} else if (t0 < 0) {  
    t1 = 2;  
} else if (t0 >= 10) {  
    t1 = 3;  
} else {  
    t1 = 4;  
}
```

Assembly code for the above C code:

```
if_0:  
    bnez t0, if_less_0  
    li   t1, 1  
    b    end_if  
if_less_0:  
    bgtz t0, if_greater_10  
    li   t1, 2  
    b    end_if  
if_greater_10:  
    li   t3, 10  
    ble  t0, t3, else  
    li   t1, 3  
    b    end_if  
else:  
    li   t1, 4  
end_if:
```

# Assembly Code for “While”

```
while((t0 = read_int()) != 0) {  
    print_int(t0)  
    print_char('\n')  
}
```



```
while:  
    li    a7, 5  
    ecall  
    mv    t0, a0  
    beqz  a0, end_while  
    li    a7, 1  
    ecall  
    li    a7, 11  
    li    a0, '\n'  
    ecall  
    b     while  
end_while:
```

# Assembly Code for “For”

```
for (t0 = 0; t0 < t1; ++t0) {  
    print_int(t0)  
    print_char('\n')  
}
```



```
for:  
    li    a7, 5  
    ecall  
    mv    t1, a0  
    mv    t0, zero  
next:  
    beq    t0, t1, end_for  
    mv    a0, t0  
    li    a7, 1  
    ecall  
    li    a7, 11  
    li    a0, '\n'  
    ecall  
    addi   t0, t0, 1  
    b      next  
end_for:
```

# Assembly Code for Nested “For”

```
for (t0 = 0; t0 < s0; ++t0) {  
    for (t1 = 0; t0 < s1; ++t1) {  
        print_int(t0)  
        print_char(':')  
        print_int(t1)  
        print_int(' ')  
    }  
    print_char('\n')  
}
```

Assembly code for the nested for loop:

```
mv t0, zero  
next_t0:  
beq t0, s0, end_for_t0  
mv t1, zero  
next_t1:  
beq t1, s1, end_for_t1  
print_int(t0)  
print_char(':')  
print_int(t1)  
print_char(' ')  
addi t1, t1, 1  
b next_t1  
end_for_t1:  
print_char('\n')  
addi t0, t0, 1  
b next_t0  
end_for_t0:
```

Arrows indicate the mapping between the C code and the assembly code:

- The first `for` loop header `for (t0 = 0; t0 < s0; ++t0) {` maps to `mv t0, zero`, `next_t0:`, and `beq t0, s0, end_for_t0`.
- The second `for` loop header `for (t1 = 0; t0 < s1; ++t1) {` maps to `mv t1, zero`, `next_t1:`, and `beq t1, s1, end_for_t1`.
- The inner loop body `print_int(t1)` maps to `print_int(t1)`.
- The inner loop body `print_int(' ')` maps to `print_char(' ')`.
- The inner loop body `print_char('\n')` maps to `print_char('\n')`.

# Macros

**Macro** is a pattern-matching and replacement facility that provides a simple mechanism to name a frequently used sequence of instructions.

```
.macro print_int (%x)
li    a7, 1
mv    a0, %x
ecall
.end_macro
```

```
.macro read_int (%x)
li    a7, 5
ecall
mv    %x, a0
.end_macro
```

Use Macros to  
Simplify Your Code



***main:***

```
read_int(t0)
print_int(t0)
```

# Including Macro Libraries

It is possible to place macros in a **library** file and **include** it in other assembly programs.

```
.include "macrolib.s"
```

```
main:
```

```
    read_int(t0)
```

```
    print_int(t0)
```

The *read\_int* and *print\_int* macros are defined in the *macrolib.s* file.

The file must be in the same directory as the program.

# Macro Constants and Single-Line Macros

The **.eqv** directive can be used to define macro constants and single-line macros.

```
.eqv VAL 0x123
```

```
.eqv X t0
```

```
.eqv Y t1
```

```
.eqv SUM addi Y, X, VAL
```

```
main:
```

```
li    X, 0x111
```

```
SUM
```

# Data Segment

Segment **.data** stores static data (global variables and constants), which are described with the following directives:

**.data**

**.word** 0xDEADBEEF # 32-bit value

**.half** 0x1234, 0x4567 # 16-bit values

**.byte** 0x98, 0x76, 0x65, 0x43 # 8-bit values

**.space** 8 # 8 bytes of empty space

**.ascii** "Hello " # String

**.asciz** "World! " # Zero-terminated string



# Data Alignment

Data items are aligned in memory by their size for convenience of access. This means ***address is multiple of size***. Default alignment is as follows:

- **.byte**    # 1 byte
- **.half**    # 2 bytes
- **.word**    # 4 bytes

It is possible to specify a ***custom alignment by  $2^n$  bytes*** for a next data item with the .align directive.

- **.align 0**    # 1 byte
- **.align 1**    # 2 bytes
- **.align 2**    # 4 bytes
- **.align 3**    # 8 bytes
- etc.

# Data Alignment Example

**.data**  
**.space 3**

**word1:**  
**.word 0x12345678**  
**half1:**  
**.half 0x1234**  
**byte1:**  
**.byte 0x12**  
**.align 4**

**word2:**  
**.word 0x12345678**  
**.align 3**  
**half2:**  
**.half 0x1234**  
**.align 3**  
**byte2:**  
**.byte 0x12**  
**.align 0**  
**word3:**  
**.word 0x12345678**

**Default  
Alignment**

**Custom  
Alignment**

Labels	
Label	Address ▲
data.s	
word1	0x10010004
half1	0x10010008
byte1	0x1001000a
word2	0x10010010
half2	0x10010018
byte2	0x10010020
word3	0x10010021
<input checked="" type="checkbox"/> Data <input checked="" type="checkbox"/> Text	

# Load and Store Instructions

## Load instructions

lb **t1**, offset(**t2**) # t1 <- sign-extended 8-bit value from address t2 + offset  
lbu **t1**, offset(**t2**) # t1 <- zero-extended 8-bit value from address t2 + offset  
lh **t1**, offset(**t2**) # t1 <- sign-extended 16-bit value from address t2 + offset  
lhu **t1**, offset(**t2**) # t1 <- zero-extended 16-bit value from address t2 + offset  
lw **t1**, offset(**t2**) # t1 <- contents of address t2 + offset

## Store Instructions

sb **t1**, offset(**t2**) # Store low-order 8 bits (byte) of t1 to address t2 + offset  
sh **t1**, offset(**t2**) # Store low-order 16 bits (half) of t1 to address t2 + offset  
sw **t1**, offset(**t2**) # Store contents of t1 to address t2 + offset

## Load Address Pseudoinstruction

la **t2**, label # t1 <- address of label

# Load and Store Example

# x, y, and z are static variables

**int** x, y, z;

x = read\_int();

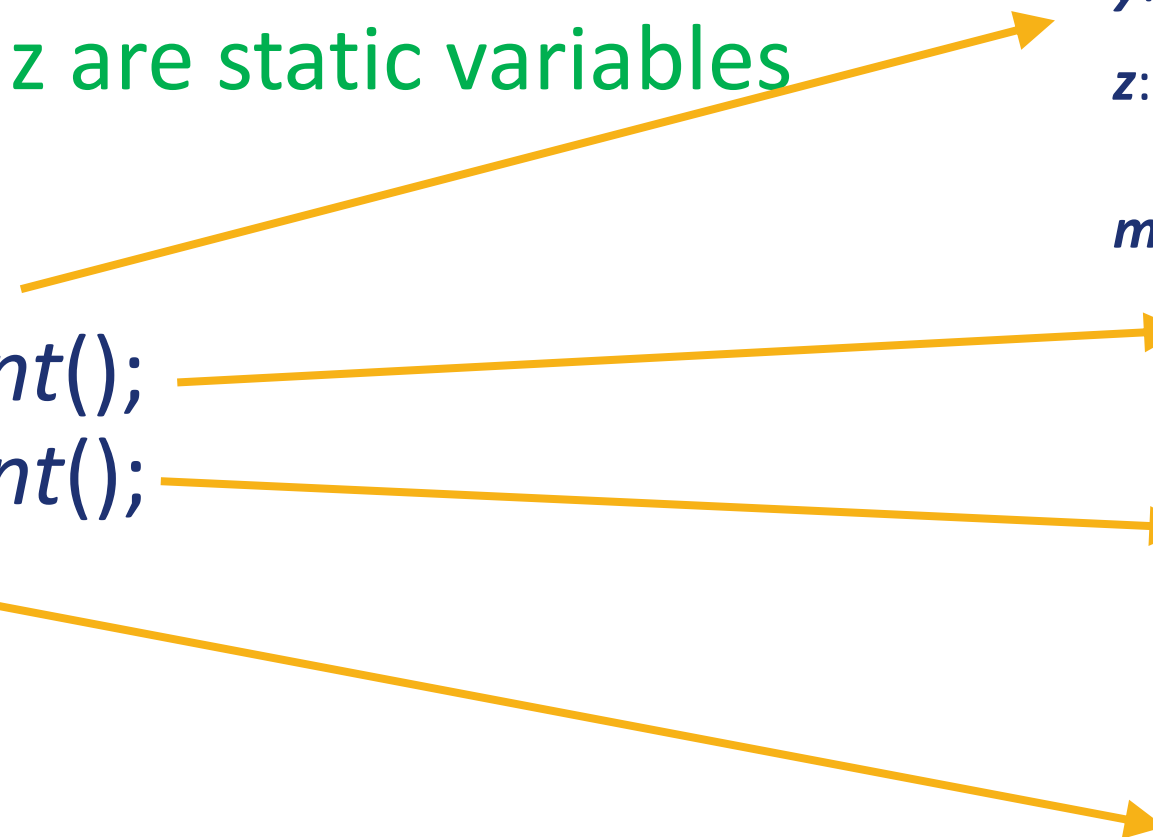
y = read\_int();

z = x + y;

```
.data
x:
.word 0
y:
.word 0
z:
.word 0
.text
main:
    read_int(t0)
    la    t2, x
    sw    t0, 0(t2)

    read_int(t0)
    la    t2, y
    sw    t0, 0(t2)

    la    t2, x
    lw    t0, 0(t2)
    la    t2, y
    lw    t1, 0(t2)
    add   t3, t0, t1
    la    t2, z
    sw    t3, 0(t2)
```



# Load and Store With Offset Example

# data[3] is a static array that stores three integer variables

**int** data[3];

*x = read\_int();*

*y = read\_int();*

*z = x + y;*

**.data**  
*data:*  
    *.word* 0, 0, 0  
    *.text*  
**main:**  
    *la* t2, data  
  
    *read\_int*(t0)  
    *sw* t0, 0(t2)  
  
    *read\_int*(t0)  
    *sw* t0, 4(t2)  
  
    *lw* t0, 0(t2)  
    *lw* t1, 4(t2)  
    *add* t3, t0, t1  
    *sw* t3, 8(t2)

# Any Questions?

```
                .text
__start:      addi t1, zero, 0x18
                addi t2, zero, 0x21
cycle:        beq t1, t2, done
                slt t0, t1, t2
                bne t0, zero, if_less
                nop
                sub t1, t1, t2
                j cycle
                nop
if_less:      sub t2, t2, t1
                j cycle
done:         add t3, t1, zero
```