



NATIONAL RESEARCH
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Computer Architecture and Operating Systems

Lecture 2: Data Representation

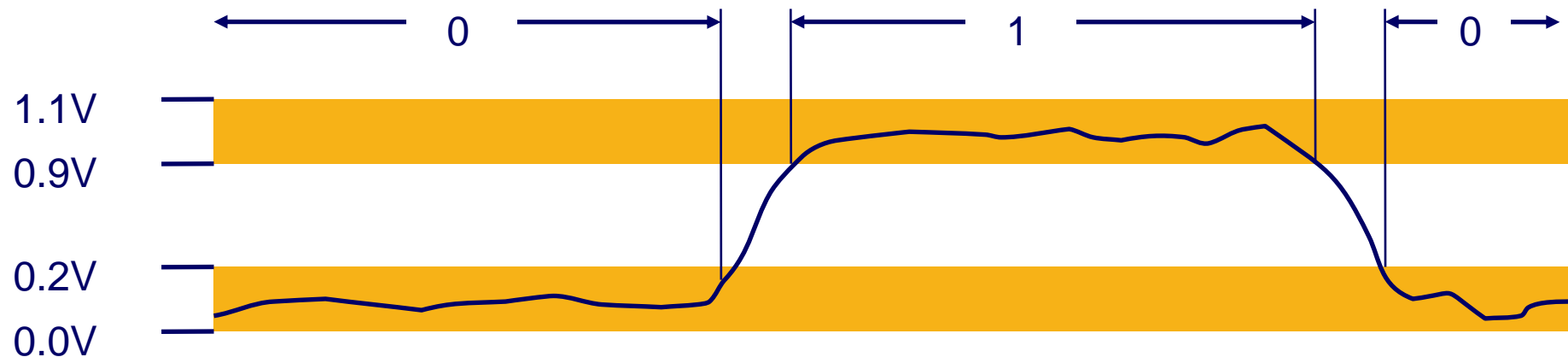
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Everything is Bits

- Each bit is 0 or 1
- By encoding/interpreting sets of bits in various ways
 - Computers determine what to do (instructions)
 - ... and represent and manipulate numbers, sets, strings, etc...
- Why bits? Electronic implementation
 - Easy to store with bistable elements
 - Reliably transmitted on noisy and inaccurate wires



Number Systems

■ Decimal numbers

1's column
10's column
100's column
1000's column

$$5374_{10} = 5 \times 10^3 + 3 \times 10^2 + 7 \times 10^1 + 4 \times 10^0$$

five thousands three hundreds seven tens four ones

■ Binary numbers

1's column
2's column
4's column
8's column

$$1101_2 = 1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 13_{10}$$

one eight one four no two one one

Powers of Two

- $2^0 = 1$
- $2^1 = 2$
- $2^2 = 4$
- $2^3 = 8$
- $2^4 = 16$
- $2^5 = 32$
- $2^6 = 64$
- $2^7 = 128$
- Handy to memorize up to 2^{10}
- $2^8 = 256$
- $2^9 = 512$
- $2^{10} = 1024$
- $2^{11} = 2048$
- $2^{12} = 4096$
- $2^{13} = 8192$
- $2^{14} = 16384$
- $2^{15} = 32768$

Number Conversion

- Decimal to binary conversion:
 - Convert 10011_2 to decimal
 - $16 \times 1 + 8 \times 0 + 4 \times 0 + 2 \times 1 + 1 \times 1 = 19_{10}$
- Decimal to binary conversion:
 - Convert 47_{10} to binary
 - $32 \times 1 + 16 \times 0 + 8 \times 1 + 4 \times 1 + 2 \times 1 + 1 \times 1 = 101111_2$

Binary Values and Range

- N -digit decimal number
 - How many values? 10^N
 - Range? $[0, 10^N - 1]$
 - Example: 3-digit decimal number:
 - $10^3 = 1000$ possible values
 - Range: $[0, 999]$
- N -bit binary number
 - How many values? 2^N
 - Range: $[0, 2^N - 1]$
 - Example: 3-digit binary number:
 - $2^3 = 8$ possible values
 - Range: $[0, 7] = [000_2 \text{ to } 111_2]$

Hexadecimal Numbers

- Base 16
- Shorthand for binary

Hex Digit	Decimal Equivalent	Binary Equivalent
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

Hexadecimal to Binary Conversion

- Hexadecimal to binary conversion:
 - Convert $4AF_{16}$ (also written 0x4AF) to binary
 - $0100\ 1010\ 1111_2$
- Hexadecimal to decimal conversion:
 - Convert $4AF_{16}$ to decimal
 - $16^2 \times 4 + 16^1 \times 10 + 16^0 \times 15 = 1199_{10}$

Bits, Bytes, Nibbles...

- Bits

10010110

most significant bit least significant bit

- Bytes & Nibbles

byte

10010110

nibble

- Bytes

CEBF9AD7

most significant byte least significant byte

Encoding Byte Values

- Byte = 8 bits
 - Binary 00000000_2 to 11111111_2
 - Decimal: 0_{10} to 255_{10}
 - Hexadecimal 00_{16} to FF_{16}
 - Base 16 number representation
 - Use characters '0' to '9' and 'A' to 'F'
 - Write $FA1D37B_{16}$ in C as
 - `0xFA1D37B`
 - `0xfa1d37b`

Example Data Representations

C Data Type	Typical 32-bit	Typical 64-bit
char	1	1
short	2	2
int	4	4
long	4	8
float	4	4
double	8	8
long double	–	–
pointer	4	8

Byte-Oriented Memory Organization

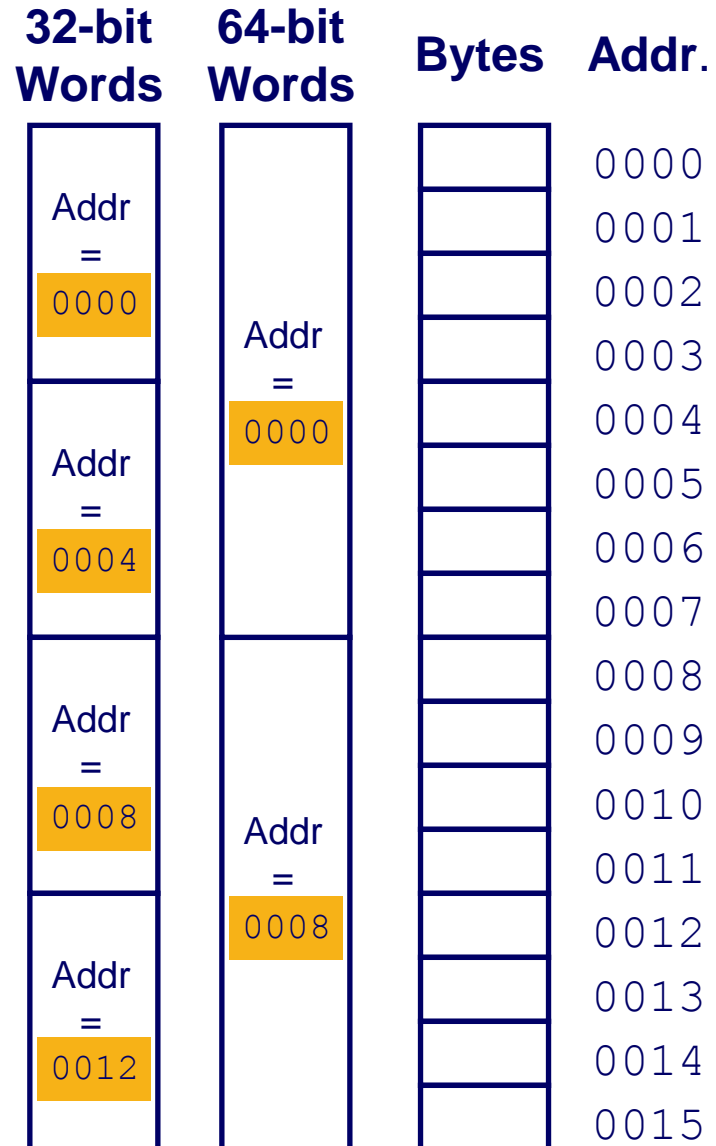
- Programs refer to data by address
 - Conceptually, envision it as a very large array of bytes
 - In reality, it's not, but can think of it that way
 - An address is like an index into that array
 - and, a pointer variable stores an address
- Note: system provides private address spaces to each “process”
 - Think of a process as a program being executed
 - So, a program can clobber its own data, but not that of others



Machine Words

- Word is a native unit of information handled by computer
- Any computer has a “Word Size”
 - Nominal size of integer-valued data
 - and of addresses
 - Until recently, most machines used 32 bits (4 bytes) as word size
 - Limits addresses to 4GB (2^{32} bytes)
 - Increasingly, machines have 64-bit word size
 - Potentially, could have 18 EB (exabytes) of addressable memory
 - That's 18.4×10^{18}
 - Machines still support multiple data formats
 - Fractions or multiples of word size
 - Always integral number of bytes

Word-Oriented Memory Organization



- Addresses Specify Byte Locations
 - Address of first byte in word
 - Addresses of successive words differ by 4 (32-bit) or 8 (64-bit)

Byte Ordering

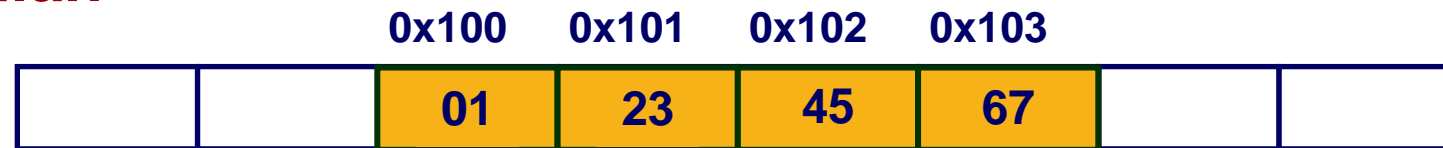
- How are the bytes within a multi-byte word ordered in memory?
- Conventions
 - **Big Endian:** Sun, PPC Mac, Internet
 - Least significant byte has highest address
 - **Little Endian:** x86, ARM processors running Android, iOS, and Windows, RISC-V
 - Least significant byte has lowest address

Byte Ordering Example

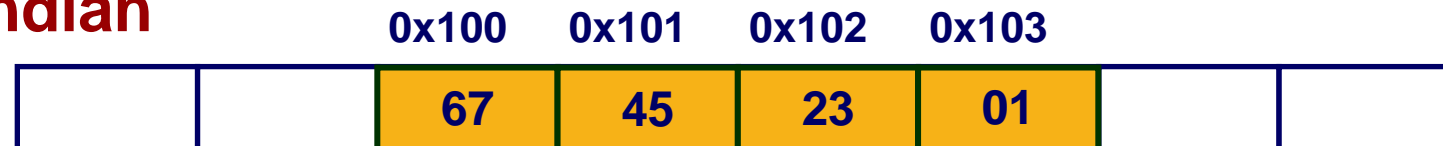
■ Example

- Variable x has 4-byte value of 0x01234567
- Address given by &x is 0x100

Big Endian



Little Endian



Encoding Integers

Unsigned


$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

C short 2 bytes long

	Decimal	Hex	Binary
x	15213	3B 6D	00111011 01101101
y	-15213	C4 93	11000100 10010011

Signed (two's complement)

$$B2T(X) = -x_{w-1} \cdot 2^{w-1} + \sum_{i=0}^{w-2} x_i \cdot 2^i$$

Sign Bit 

■ Sign Bit

- For 2's complement, most significant bit indicates sign
 - 0 for nonnegative
 - 1 for negative

Two-complement Encoding Example

$x =$ 15213: 00111011 01101101
 $y =$ -15213: 11000100 10010011

Weight	15213		-15213	
1	1	1	1	1
2	0	0	1	2
4	1	4	0	0
8	1	8	0	0
16	0	0	1	16
32	1	32	0	0
64	1	64	0	0
128	0	0	1	128
256	1	256	0	0
512	1	512	0	0
1024	0	0	1	1024
2048	1	2048	0	0
4096	1	4096	0	0
8192	1	8192	0	0
16384	0	0	1	16384
-32768	0	0	1	-32768
Sum	15213		-15213	

Boolean Algebra

- Developed by George Boole in 19th Century

- Algebraic representation of logic

- Encode “True” as 1 and “False” as 0

And

- $A \& B = 1$ when both $A=1$ and $B=1$

$\&$	0	1
0	0	0
1	0	1

Not

- $\sim A = 1$ when $A=0$

\sim	
0	1
1	0

Or

- $A | B = 1$ when either $A=1$ or $B=1$

$ $	0	1
0	0	1
1	1	1

Exclusive-Or (Xor)

- $A \wedge B = 1$ when either $A=1$ or $B=1$, but not both

\wedge	0	1
0	0	1
1	1	0

Bitwise Operations

- Operate on Bit Vectors
 - Operations applied bitwise

01101001	01101001	01101001	
& <u>01010101</u>	<u>01010101</u>	^ <u>01010101</u>	~ <u>01010101</u>
01000001	01111101	00111100	10101010

- All of the Properties of Boolean Algebra Apply

Logic Operations

- Operations `&&`, `||`, `!`
 - Different from similar bitwise operations
 - View 0 as “False”
 - Anything nonzero as “True”
 - Always return 0 or 1
 - Early termination
- Examples (8-bit data type)
 - `!0x41` \Rightarrow `0x00`
 - `!0x00` \Rightarrow `0x01`
 - `!!0x41` \Rightarrow `0x01`
 - `0x69 && 0x55` \Rightarrow `0x01`
 - `0x69 || 0x55` \Rightarrow `0x01`

Sign-Extension

- Extend number from N to M bits ($M > N$)
- Sign bit is copied to most significant bits
- Number value is same
- Example 1:
 - 4-bit representation of 3 = 0011
 - 8-bit sign-extended value: 00000011
- Example 2:
 - 4-bit representation of -5 = 1011
 - 8-bit sign-extended value: 11111011

Zero-Extension

- Extend number from N to M bits ($M > N$)
- Zeros are copied to most significant bits
- Value changes for negative numbers
- Example 1:
 - 4-bit value = $0011 = 3_{10}$
 - 8-bit zero-extended value: $00000011 = 3_{10}$
- Example 2:
 - 4-bit value = $1011 = -5_{10}$
 - 8-bit zero-extended value: $00001011 = 11_{10}$

Shift Operations

■ Left Shift: $x \ll y$

- Shift bit-vector x left y positions
 - Throw away extra bits on left
- Fill with 0's on right

■ Right Shift: $x \gg y$

- Shift bit-vector x right y positions
 - Throw away extra bits on right
- Logical shift
 - Fill with 0's on left
- Arithmetic shift
 - Replicate most significant bit on left

■ Undefined Behavior

- Shift amount < 0 or \geq word size

X	01100010
$\ll 3$	00010000
Log. $\gg 2$	00011000
Arith. $\gg 2$	00011000

X	10100010
$\ll 3$	00010000
Log. $\gg 2$	00101000
Arith. $\gg 2$	11101000

Any Questions?

```
                .text
__start:      addi t1, zero, 0x18
                addi t2, zero, 0x21
cycle:        beq t1, t2, done
                slt t0, t1, t2
                bne t0, zero, if_less
                nop
                sub t1, t1, t2
                j cycle
                nop
if_less:      sub t2, t2, t1
                j cycle
done:         add t3, t1, zero
```