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# Computer Architecture and Operating Systems

## Lecture 4: Instruction Set Architecture

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# Stored Program Concept

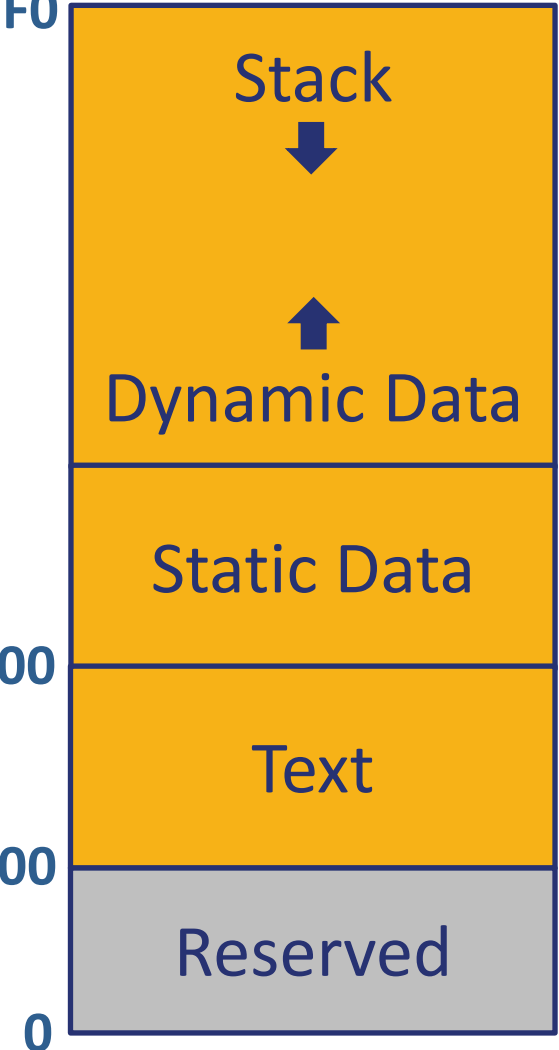
# Memory Layout

- Text: program code
- Static data: global variables
  - E.g., static variables in C, constant arrays and strings
- Dynamic data: heap
  - E.g., malloc in C, new in Java
- Stack: automatic storage

SP = 0x 0000 003F FFFF FFFF0

GP = 0x 0000 0000 1000 0000

PC = 0x 0000 0000 0040 0000



# RISC-V ISA Base and Extensions

Name	Description	Version	Status
Base			
RVWMO	Weak Memory Ordering	2.0	Ratified
RV32I	Base Integer Instruction Set, 32-bit	2.1	Ratified
RV64I	Base Integer Instruction Set, 64-bit	2.1	Ratified
RV128I	Base Integer Instruction Set, 128-bit	1.7	Open
Extensions			
M	Standard Extension for Integer Multiplication and Division	2.0	Ratified
A	Standard Extension for Atomic Instructions	2.1	Ratified
F	Standard Extension for Single-Precision Floating-Point	2.2	Ratified
D	Standard Extension for Double-Precision Floating-Point	2.2	Ratified
G	Shorthand for the base integer set (I) and above extensions (MAFD)	N/A	N/A
Q	Standard Extension for Quad-Precision Floating-Point	2.2	Ratified
C	Standard Extension for Compressed Instructions	2.0	Ratified
ZiCSR	Control and Status Register (CSR)	2.0	Ratified
Zifencei	Instruction-Fetch Fence	2.0	Ratified
And more standard and custom extensions...			

# ISA Design Principles

- **Design Principle 1:** Simplicity favors regularity
  - Regularity makes implementation simpler
  - Simplicity enables higher performance at lower cost
- **Design Principle 2:** Smaller is faster
  - 32 registers, fewer instructions
- **Design Principle 3:** Good design demands good compromises
  - Different formats complicate decoding, but allow 32-bit instructions uniformly
  - Keep formats as similar as possible

# Six Instruction Formats

- **R-format:** instructions using 3 register inputs
  - add, xor, mul - arithmetic/logical ops
- **I-format:** instructions with immediates, loads
  - addi, lw, jalr, slli
- **S-format:** store instructions
  - sw, sb
- **SB-format:** branch instructions
  - beq, bge
- **U-format:** instructions with upper immediates
  - lui, auipc - upper immediate is 20-bits
- **UJ-format:** the jump instruction
  - jal

# R-format Instructions

funct7	rs2	rs1	funct3	rd	opcode
7 bits	5 bits	5 bits	3 bits	5 bits	7 bits

**add x9, x20, x21**

0	21	20	0	9	51
0000000	10101	10100	000	01001	0110011

**0000 0001 0101 1010 0000 0100 1011 0011<sub>two</sub> = 015A04B3<sub>16</sub>**

## ■ Arithmetic Instructions

- opcode: operation code
- rd: destination register number
- funct3: 3-bit function code (additional opcode)
- rs1 and rs2: first and second source register 5-bit numbers
- funct7: 7-bit function code (additional opcode)

# I-format Instructions

immediate	rs1	funct3	rd	opcode
12 bits	5 bits	3 bits	5 bits	7 bits

**addi t0, t1, 123**

0x123	6	0	5	19
000100100011	00110	000	00101	0010011

**0001 0010 0011 0011 0000 0010 1001 0011<sub>two</sub> = 0x12330293<sub>16</sub>**

- Immediate arithmetic and load instructions
  - rs1: source or base address register number
  - immediate: constant operand, or offset added to base address
    - 2s-complement, sign extended



# S-format Instructions

imm[11:5]	rs2	rs1	funct3	imm[4:0]	opcode
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7 bits

5 bits

5 bits

3 bits

5 bits

7 bits

**sw t0, 4(t1)**

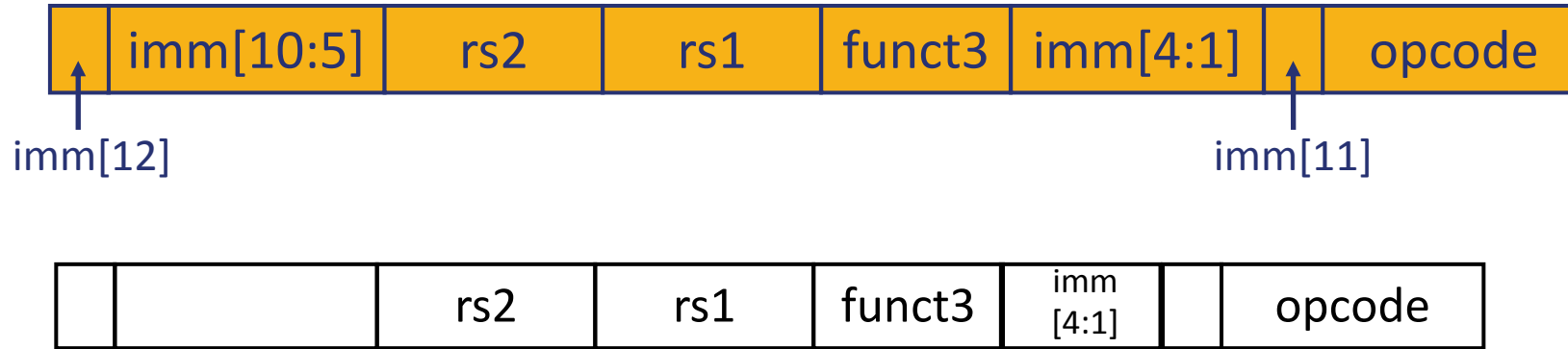
0	5	6	2	4	35
0000000	00101	00110	010	00100	100011

**0000 0000 0101 0011 0010 0010 0010 0011<sub>two</sub> = 0x00532223<sub>16</sub>**

## ■ Different immediate format for store instructions

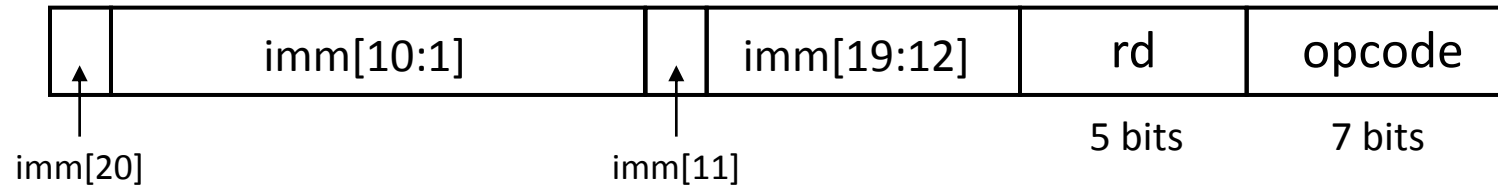
- rs1: base address register number
- rs2: source operand register number
- immediate: offset added to base address
  - Split so that rs1 and rs2 fields always in the same place

# SB-format Instructions



- Branch instructions specify
  - Opcode, two registers, target address
- Most branch targets are near branch
  - Forward or backward
- PC-relative addressing
  - Target address =  $PC + \text{immediate} \times 2$

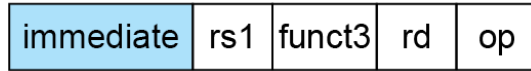
# UJ-format Instructions



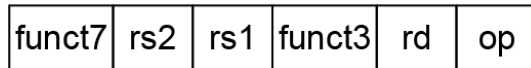
# U-format Instructions

# RISC-V Addressing Summary

## 1. Immediate addressing



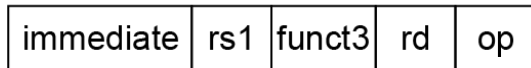
## 2. Register addressing



Registers

Register

## 3. Base addressing



Memory

Register

+

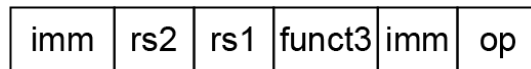
Byte

Halfword

Word

Doubleword

## 4. PC-relative addressing



Memory

PC

+

Word

# RISC-V Encoding Summary

Name (Field Size)	7 bits	5 bits	5 bits	3 bits	5 bits	7 bits	Comments
R-type	funct7	rs2	rs1	funct3	rd	opcode	Arithmetic instruction format
I-type	immediate[11:0]		rs1	funct3	rd	opcode	Loads & immediate arithmetic
S-type	immed[11:5]	rs2	rs1	funct3	immed[4:0]	opcode	Stores
SB-type	immed[12,10:5]	rs2	rs1	funct3	immed[4:1,11]	opcode	Conditional branch format
UJ-type	immediate[20,10:1,11,19:12]				rd	opcode	Unconditional jump format
U-type	immediate[31:12]				rd	opcode	Upper immediate format

# Any Questions?

```
                .text
__start:        addi t1, zero, 0x18
                addi t2, zero, 0x21
cycle:          beq t1, t2, done
                slt t0, t1, t2
                bne t0, zero, if_less
                nop
                sub t1, t1, t2
                j cycle
                nop
if_less:        sub t2, t2, t1
                j cycle
done:           add t3, t1, zero
```