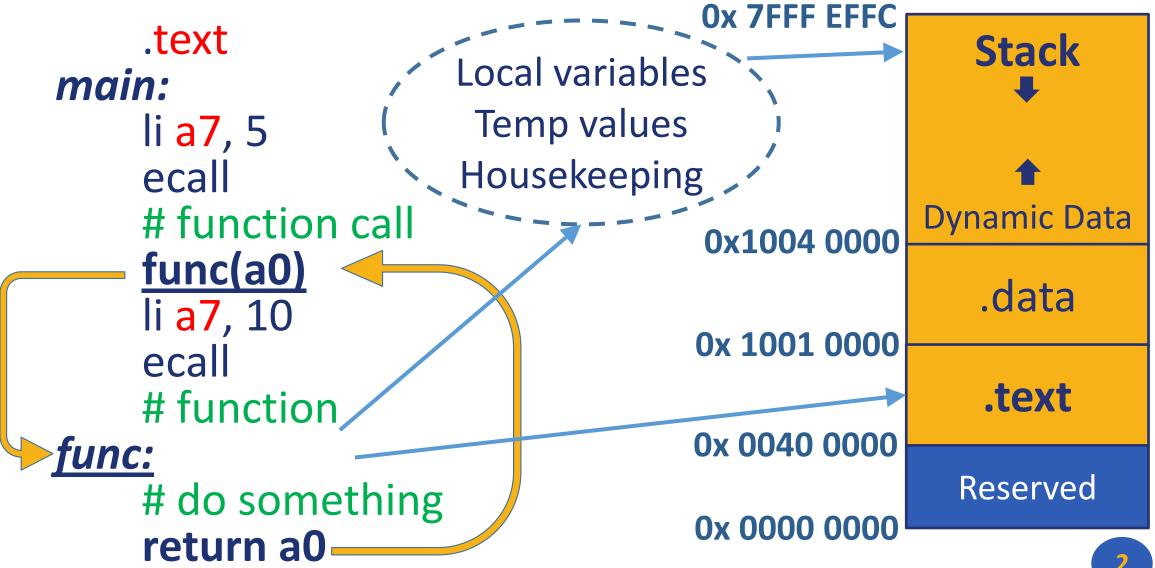


# Computer Architecture and Operating Systems Lecture 6: Assembly Programming – Stack

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# Program Structure and Memory Layout



#### **Notion of Function**

- Function (procedure) is a code that performs some task based on the arguments with which it is provided
- Caller is a code that calls a function and provides it with the necessary arguments
- Callee is a function that executes instructions based on arguments provided by the caller and then returns control to the caller
- Return address is a link that allows the callee to return control to the caller
- Jump-and-link instruction is an instruction that branches to an address and simultaneously saves the address of the next instruction in to a register

# **Function Call Steps**

- ■Place arguments in registers a0 (x10) to a7 (x17)
- Save return address in ra (x1) and jump to function
- •Allocate stack memory for the function
- Perform function's operations
- Free stack memory allocated for the function
- Place result in register a0 for caller
- Return to place of call (address in ra)

# RISC-V Register Conventions

Register	Name	Use	Saver
х0	zero	constant 0	n/a
<b>x1</b>	ra	return address	caller
<b>x2</b>	sp	stack pointer	callee
х3	gp	global pointer	
х4	tp	thread pointer	
x5-x7	t0-t2	temporaries	caller
<b>8</b> x	s0/fp	saved/ frame pointer	callee
х9	<b>s1</b>	saved	callee
x10-x17	a0-a7	arguments	caller
x18-x27	s <b>2</b> -s <b>1</b> 1	saved	callee
x28-x31	t3-t6	temporaries	caller

## Jump-and-Link Instructions

- Function call: jump and link
   jal ra, FunctionLabel (UJ-type)
  - Address of the next instruction is put in ra (x1)
  - Jumps to target address
- Function return: jump and link register jalr zero, 0(ra) (I-type)
  - Like **jal**, but jumps to 0 + address in **ra** (x1)
  - Use zero (x0) as rd (zero cannot be changed)
  - Can also be used for computed jumps
    - e.g., for case/switch statements

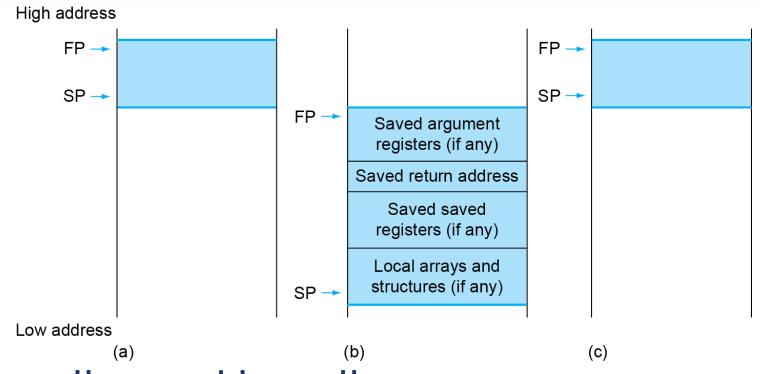
#### Jump-and-Link Pseudo Instructions

i *label* # Jump to label and do not save return address jal label # Jump to label and set return address to ra jalr t0 # Jump to address in t0 and set return address to ra jalr t0, -100 # Jump to address t0-100 and set return address to ra jr t0 # Jump Register: Jump to address in t0 jr t0, -100 # Jump Register: Jump to address t0-100

#### Stack

- Stack is a data structure for spilling registers organized as a last-in-first-out queue
- Dynamic memory for storing data (such as local variables) for function calls is organized as a task
- Stack pointer is a value denoting the most recently allocated address on the stack
- Push means to add element to stack
- Pop means to remove element from stack

#### Local Data on Stack



- Local data allocated by callee
  - Local variables, arrays, etc.
- Function frame (activation record)
  - Segment of stack containing function's saved registers and local variables

## Saving Registers

A function can overwrite values of registers.

Sometimes is undesirable. There are special rules to handle this issues. They specify who is responsible for saving the registers.

- Callee-saved register is a register saved by the routine being called
- Caller-saved register is register saved by the routine making a function call

### Function Example

```
int leaf_example (int g, int h, int i, int j) {
  int f = (g + h) - (i + j);
  return f;
}
```

#### **Requirements:**

```
arguments g, ..., j in a0(x10)...a3(x13)f in s4 (x20)
```

- temporaries t0(x5), t1(x6)
- need to save t0, t1, s4 on stack

### Function Assembly Code

```
main:
                                                leaf_example:
    read int(t0) # read g
                                                     addi sp, sp, -12
    read int(t1) # read h
                                                     sw t0, 8(sp)
    read int(t2) # read i
                                                     sw t1, 4(sp)
    read int(t3) # read j
                                                     sw s4, 0(sp)
    mv a0, t0
                                                     add t0, a0, a1
    mv a1, t1
                                                     add t1, a2, a3
    mv a2, t2
                                                     sub s4, t0, t1
    mv a3, t3
                                                     mv a0, s4
    jal ra, leaf_example
                                                     lw s4, O(sp)
    mv t4, a0
                                                     lw t1, 4(sp)
    print int(t0, t1, t2, t3, t4)
                                                     lw t0, 8(sp)
    li a7, 10
                                                     addi sp, sp, 12
    ecall
                                                     jalr x0, 0(ra)
```

# Preserving Callee-Saved Registers

#### Preserve registers:

```
addi sp, sp, -20 # make room on stack for 5 registers sw ra, 16(sp) # save ra (x1) on stack sw s1, 12(sp) # save s1 (x9) on stack sw s2, 8(sp) # save s2 (x18) on stack sw s3, 4(sp) # save s3 (x19) on stack sw s4, 0(sp) # save s4 (x20) on stack
```

#### Restore registers:

```
lw s4, 0(sp) # restore s4 (x20) from stack
lw s3, 4(sp) # restore s3 (x19) from stack
lw s2, 8(sp) # restore s2 (x18) from stack
lw s1, 12(sp) # restore s1 (x9) from stack
lw ra, 16(sp) # restore ra (x1) from stack
addi sp, sp, 20 # restore stack pointer
jalr zero, 0(ra) # return to caller
```

## Preserving Caller-Saved Registers

#### Preserve registers:

```
addi sp, sp, -16 # make room on stack for 4 registers sw t0, 12(sp) # save t0 (x5) on stack sw t1, 8(sp) # save t1 (x6) on stack sw t2, 4(sp) # save t2 (x7) on stack sw t3, 0(sp) # save t3 (x28) on stack
```

#### Restore registers:

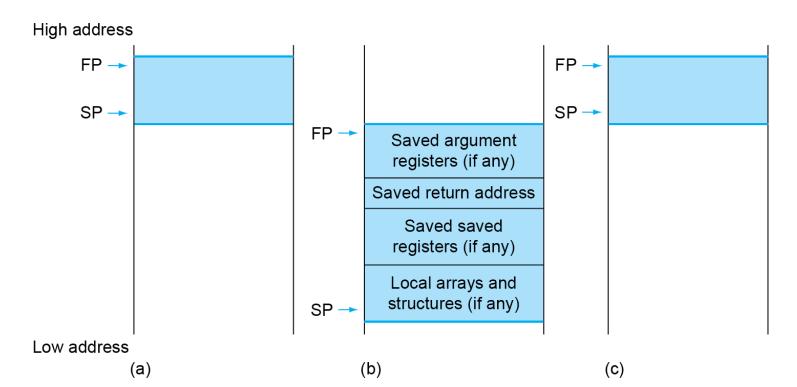
```
jal ra, callee # jump to callee
lw t3, 0(sp) # restore t3 (x28) from stack
lw t2, 4(sp) # restore t2 (x7) from stack
lw t1, 8(sp) # restore t1 (x6) from stack
lw t0, 12(sp) # restore t0 (x5) from stack
addi sp, sp, 16 # restore stack pointer
```

#### Recursive Function Example

```
fact:
                                            addi t0, a0, -1
                                            bgez t0, fact else
                                                  a0, 1
int fact (int n) {
                                            jalr zero, O(ra)
  if (n < 1) {
                                         fact else:
                                           addi sp, sp, -8 sw ra, 4(sp)
     return 1;
  } else {
                                            sw a0, 0(sp)
     return n * fact(n - 1);
                                                  a0, a0, -1
                                            addi
                                                   ra, fact
                                            ial
                                                   t1, a0
                                            mv
                                                   a0, 0(sp)
                                            lw
                                                   ra, 4(sp)
                                            lw
                                            addi
                                                   sp, sp, 8
                                            mul
                                                   a0, a0, t1
                                            jalr
                                                   zero, O(ra)
```

#### Frame Pointer

Frame pointer is a value denoting the location of the saved registers and local variables for a given procedure. Simplifies programming because when stack-pointer changes programmers have to use different offsets to access the same values.



## **Using Frame Pointer**

```
main:
  mv fp, sp
  jal ra, func
  li a7, 10
  ecall
func:
  addi sp, sp, -4
  li t0, 1
  sw t0, 0(sp)
  addi sp, sp, -4
  li t0, 2
  sw t0, 0(sp)
  addi sp, sp, -4
  li t0, 3
  sw t0, 0(sp)
  lw t0, 0(fp)
  print_int(t0)
  lw t0, -4(fp)
  print_int(t0)
  lw t0, -8(fp)
  print_int(t0)
  addi sp, sp, 12
  jalr zero, O(ra)
```

### Any Questions?

```
__start: addi t1, zero, 0x18
addi t2, zero, 0x21

cycle: beq t1, t2, done
slt t0, t1, t2
bne t0, zero, if_less
nop
sub t1, t1, t2
j cycle
nop

if_less: sub t2, t2, t1
j cycle
done: add t3, t1, zero
```