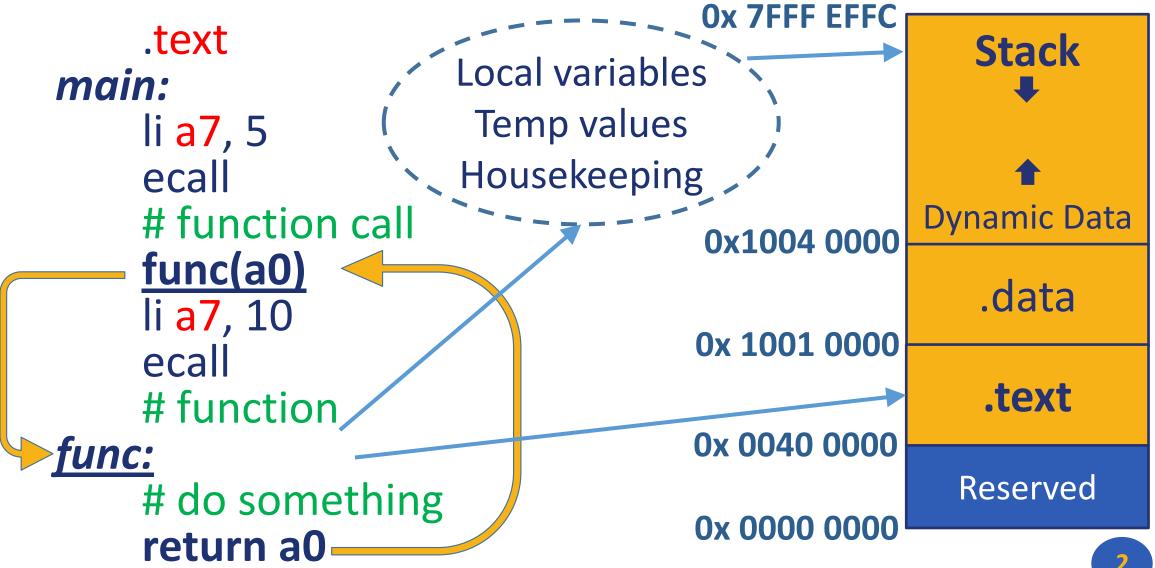


Computer Architecture and Operating Systems Lecture 6: Assembly Programming – Stack

Andrei Tatarnikov

atatarnikov@hse.ru
@andrewt0301

Program Structure and Memory Layout



Notion of Function

- Function (procedure) is a code that performs some task based on the arguments with which it is provided
- Caller is a code that calls a function and provides it with the necessary arguments
- Callee is a function that executes instructions based on arguments provided by the caller and then returns control to the caller
- Return address is a link that allows the callee to return control to the caller
- Jump-and-link instruction is an instruction that branches to an address and simultaneously saves the address of the next instruction in to a register

Function Call Steps

- ■Place arguments in registers a0 (x10) to a7 (x17)
- Save return address in ra (x1) and jump to function
- •Allocate stack memory for the function
- Perform function's operations
- Free stack memory allocated for the function
- Place result in register a0 for caller
- Return to place of call (address in ra)

RISC-V Register Conventions

Register	Name	Use	Saver
х0	zero	constant 0	n/a
x1	ra	return address	caller
x2	sp	stack pointer	callee
х3	gp	global pointer	
x4	tp	thread pointer	
x5-x7	t0-t2	temporaries	caller
8 x	s0/fp	saved/ frame pointer	callee
х9	s1	saved	callee
x10-x17	a0-a7	arguments	caller
x18-x27	s 2 -s 1 1	saved	callee
x28-x31	t3-t6	temporaries	caller

Jump-and-Link Instructions

- Function call: jump and link
 jal ra, FunctionLabel (UJ-type)
 - Address of the next instruction is put in ra (x1)
 - Jumps to target address
- Function return: jump and link register jalr zero, 0(ra) (I-type)
 - Like **jal**, but jumps to 0 + address in **ra** (x1)
 - Use zero (x0) as rd (zero cannot be changed)
 - Can also be used for computed jumps
 - e.g., for case/switch statements

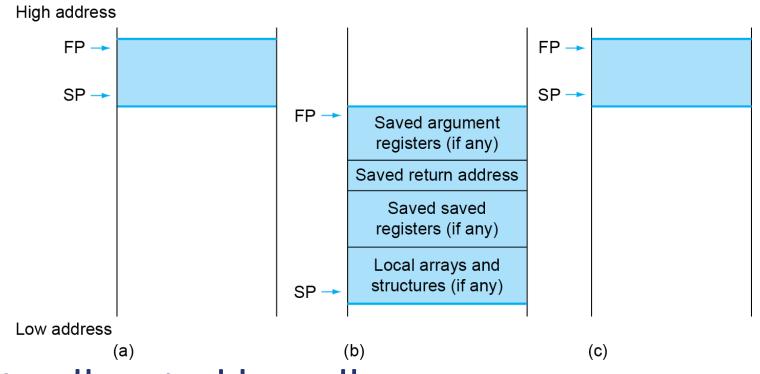
Jump-and-Link Pseudo Instructions

i *label* # Jump to label and do not save return address jal label # Jump to label and set return address to ra jalr t0 # Jump to address in t0 and set return address to ra jalr t0, -100 # Jump to address t0-100 and set return address to ra jr t0 # Jump Register: Jump to address in t0 jr t0, -100 # Jump Register: Jump to address t0-100

Stack

- Stack is a data structure for spilling registers organized as a last-in-first-out queue
- Dynamic memory for storing data (such as local variables) for function calls is organized as a task
- Stack pointer is a value denoting the most recently allocated address on the stack
- Push means to add element to stack
- Pop means to remove element from stack

Local Data on Stack



- Local data allocated by callee
 - Local variables, arrays, etc.
- Function frame (activation record)
 - Segment of stack containing function's saved registers and local variables

Saving Registers

A function can overwrite values of registers.

Sometimes is undesirable. There are special rules to handle this issues. They specify who is responsible for saving the registers.

- Callee-saved register is register saved by the routine making a function call
- Caller-saved register is a register saved by the routine being called

Function Example

```
int leaf_example (int g, int h, int i, int j) {
  int f = (g + h) - (i + j);
  return f;
}
```

Requirements:

```
arguments g, ..., j in a0(x10)...a3(x13)
f in s4 (x20)
temporaries t0(x5), t1(x6)
```

- need to save t0, t1, s4 on stack

Function Assembly Code

```
main:
leaf_example:
                                      read int(t0) # read g
     addi sp, sp, -12
                                      read int(t1) # read h
     sw t0, 8(sp)
                                      read int(t2) # read i
     sw t1, 4(sp)
                                      read int(t3) # read j
     sw s4, 0(sp)
                                      mv a0, t0
     add t0, a0, a1
                                      mv a1, t1
     add t1, a2, a3
                                      mv a2, t2
     sub s4, t0, t1
                                      mv a3, t3
     mv a0, s4
                                      jal ra, leaf example
     lw s4, 0(sp)
                                      mv t4, a0
     lw t1, 4(sp)
                                      print int(t0, t1, t2, t3, t4)
     lw t0, 8(sp)
                                      li a7, 10
     addi sp, sp, 12
     ialr x0, 0(x1)
                                      ecall
```

Preserving Callee-Saved Registers

Preserve registers:

```
addi sp, sp, -20 # make room on stack for 5 registers sw ra, 16(sp) # save ra (x1) on stack sw s1, 12(sp) # save s1 (x9) on stack sw s2, 8(sp) # save s2 (x18) on stack sw s3, 4(sp) # save s3 (x19) on stack sw s4, 0(sp) # save s4 (x20) on stack
```

Restore registers:

```
sw s4, 0(sp) # restore s4 (x20) from stack
sw s3, 4(sp) # restore s3 (x19) from stack
sw s2, 8(sp) # restore s2 (x18) from stack
sw s1, 12(sp) # restore s1 (x9) from stack
sw ra, 16(sp) # restore ra (x1) from stack
addi sp, sp, 20 # restore stack pointer
jalr zero, 0(ra) # return to caller
```

Preserving Caller-Saved Registers

Preserve registers:

```
addi sp, sp, -16 # make room on stack for 4 registers sw t0, 12(sp) # save t0 (x5) on stack sw t1, 8(sp) # save t1 (x6) on stack sw t2, 4(sp) # save t2 (x7) on stack sw t3, 0(sp) # save t3 (x28) on stack
```

Restore registers:

```
jal ra, callee # jump to callee
sw t3, 0(sp) # restore t3 (x28) from stack
sw t2, 4(sp) # restore t2 (x7) from stack
sw t1, 8(sp) # restore t1 (x6) from stack
sw t0, 12(sp) # restore t0 (x5) from stack
addi sp, sp, 16 # restore stack pointer
```

Recursive Function Example

```
fact:
                                            addi t0, a0, -1
                                            bgez t0, fact else
                                                  a0, 1
int fact (int n) {
                                            jalr zero, O(ra)
  if (n < 1) {
                                         fact else:
                                           addi sp, sp, -8 sw ra, 4(sp)
     return 1;
  } else {
                                            sw a0, 0(sp)
     return n * fact(n - 1);
                                                  a0, a0, -1
                                            addi
                                                   ra, fact
                                            ial
                                                   t1, a0
                                            mv
                                                   a0, 0(sp)
                                            lw
                                                   ra, 4(sp)
                                            lw
                                            addi
                                                   sp, sp, 8
                                            mul
                                                   a0, a0, t1
                                            jalr
                                                   zero, O(ra)
```

Any Questions?

```
__start: addi t1, zero, 0x18
    addi t2, zero, 0x21

cycle: beq t1, t2, done
    slt t0, t1, t2
    bne t0, zero, if_less
    nop
    sub t1, t1, t2
    j cycle
    nop

if_less: sub t2, t2, t1
    j cycle

done: add t3, t1, zero
```